

The Subspace Flatness Conjecture and Faster Integer Programming

Victor Reis

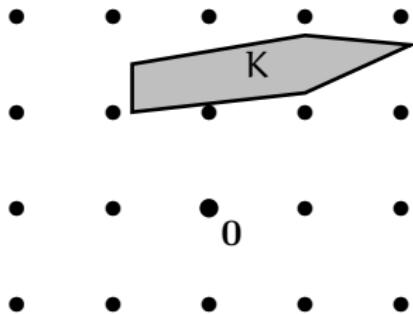
Joint work with Thomas Rothvoss



Brief History of Integer Programming

Problem

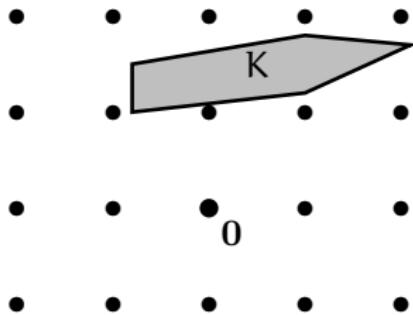
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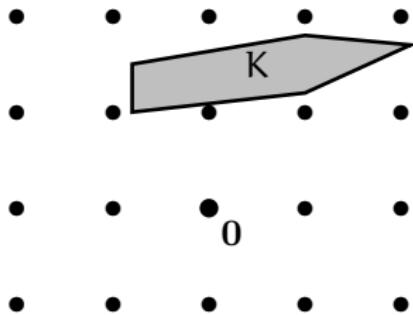


- ▶ NP-hard [Karp '72]

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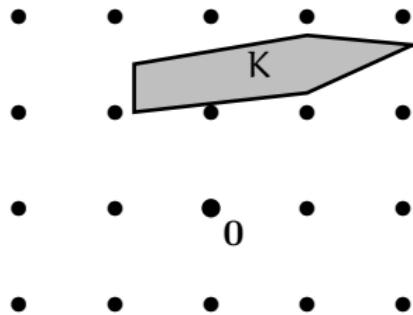


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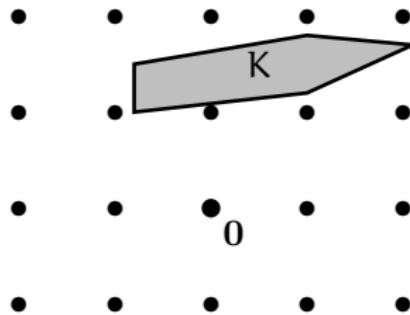


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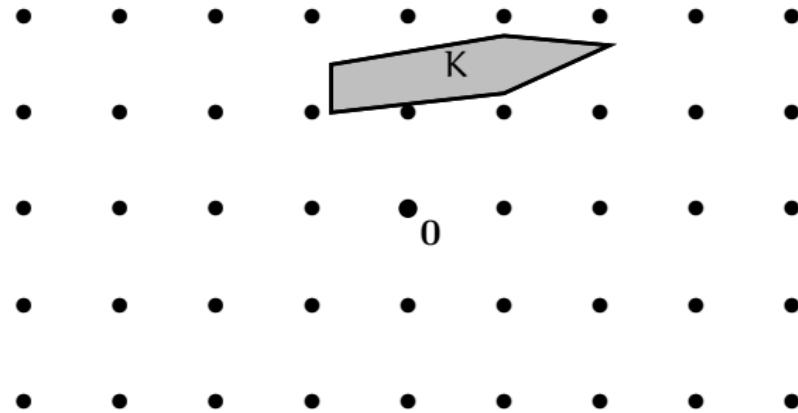
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- ▶ $(O(n))^n$ [Dadush '12, Dadush, Eisenbrand, Rothvoss '22]

Dadush's algorithm



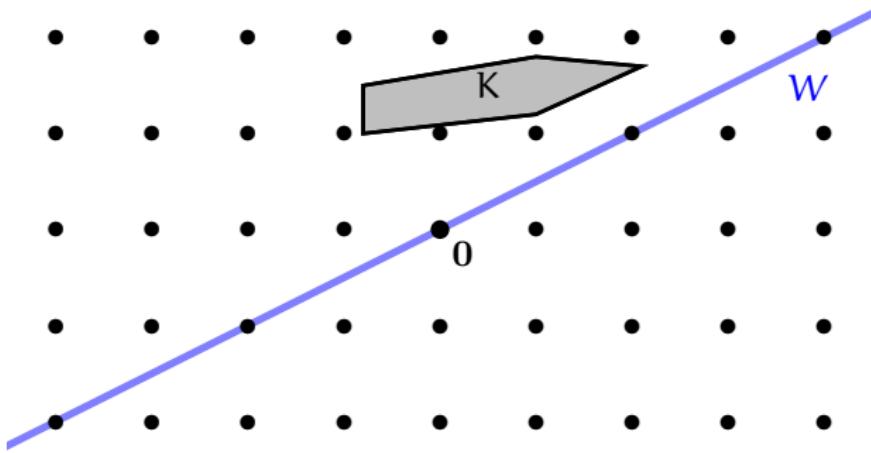
► A **lattice** is the integer span $\mathcal{L} = B\mathbb{Z}^k$ of the columns of $B \in \mathbb{R}^{n \times k}$.

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Output: Point in $K \cap \mathcal{L}$

- (1) Shrink K if 'too large' compared to \mathcal{L}
- (2) Find subspace W for which $X := \Pi_W(K) \cap \Pi_W(\mathcal{L})$ is small
- (3) Enumerate X
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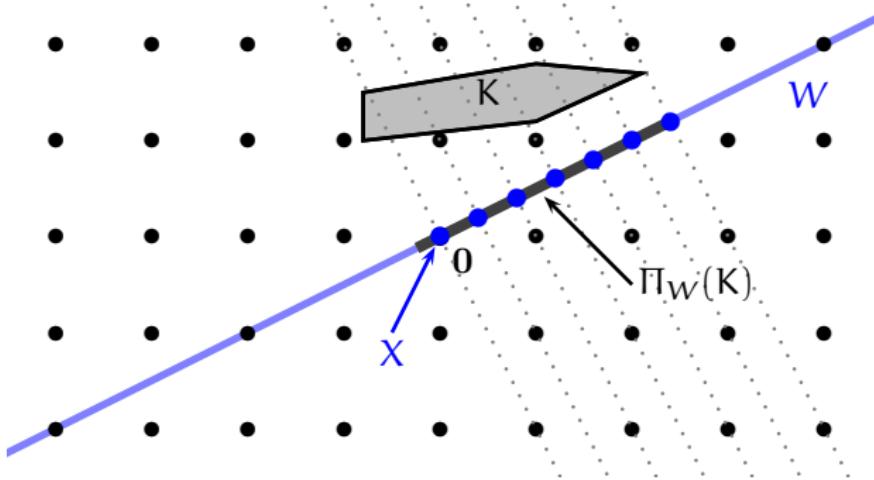
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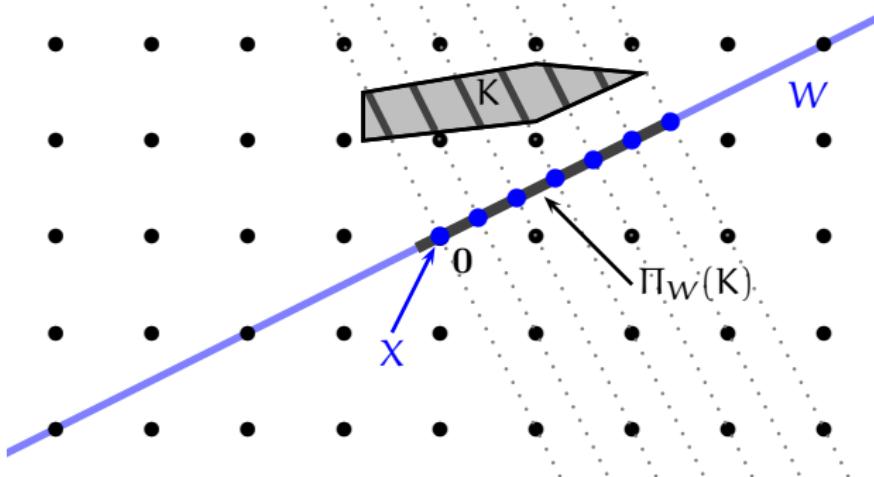
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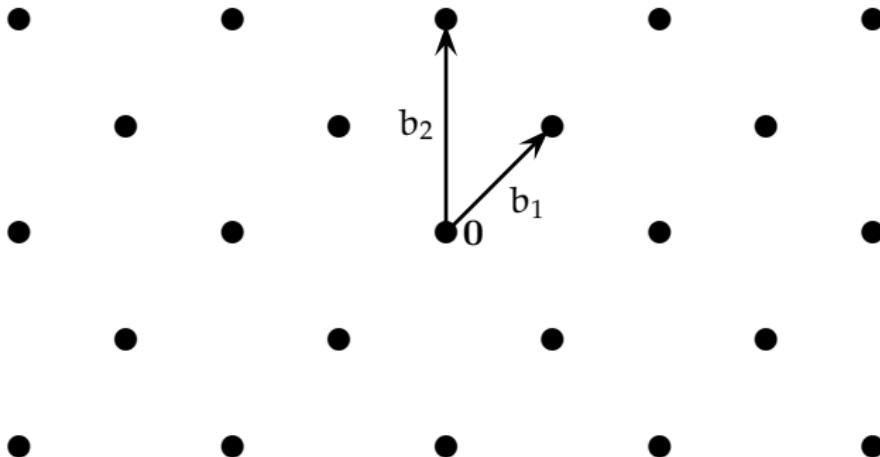
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Covering radius

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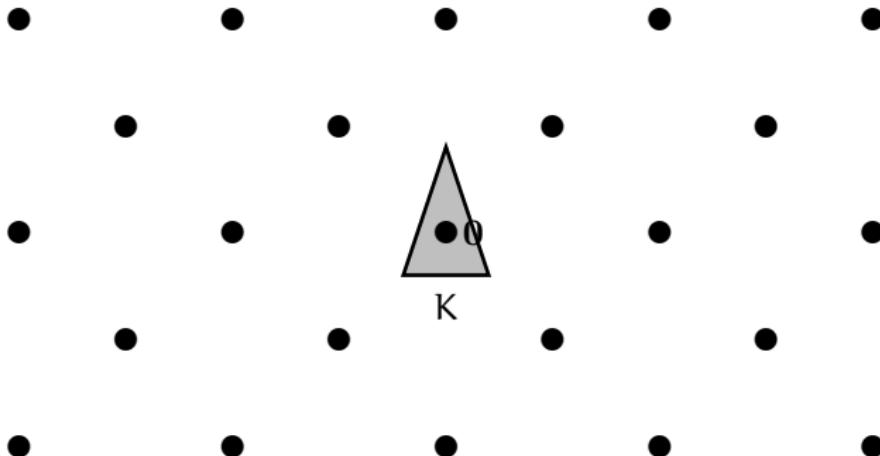
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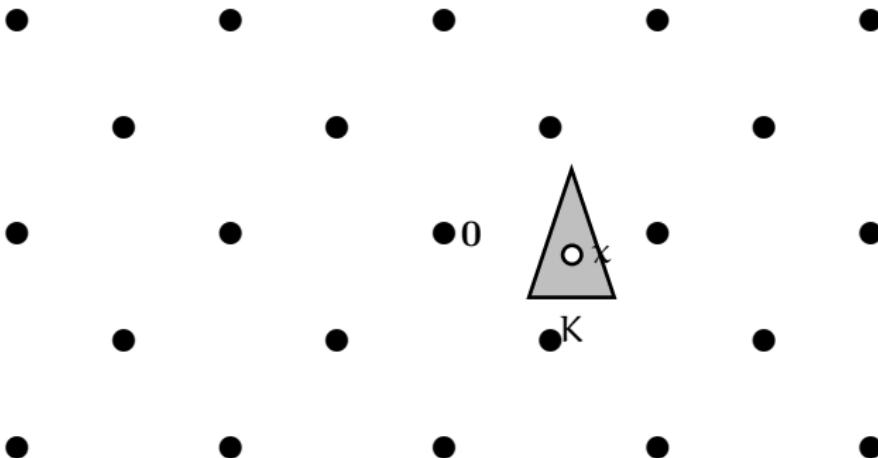
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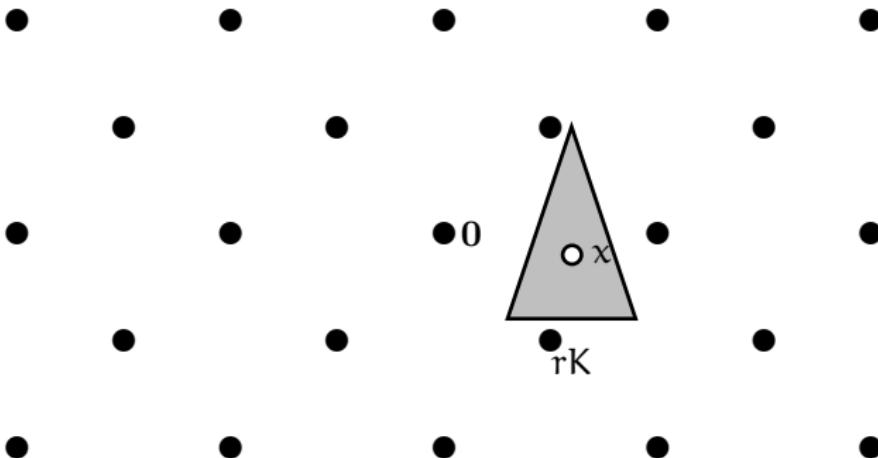
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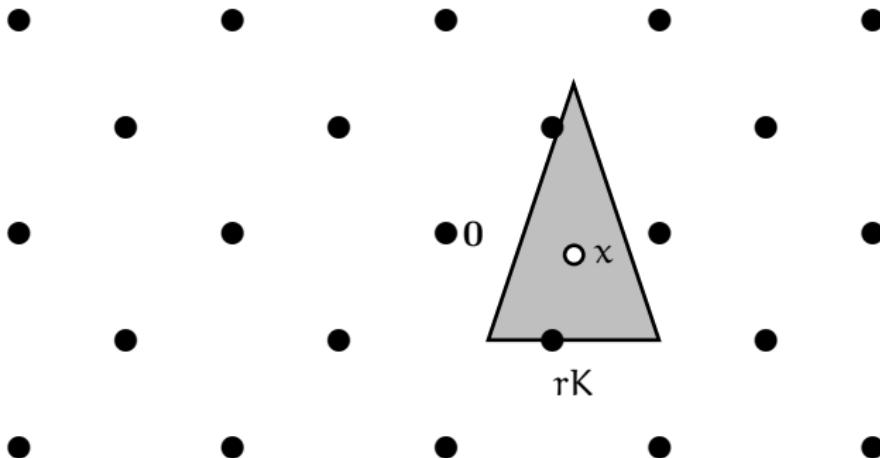
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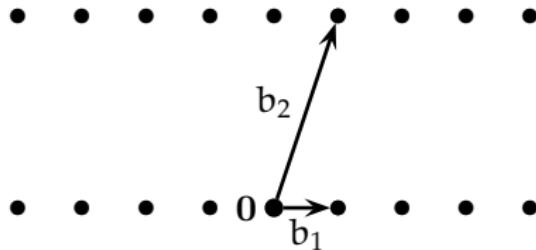
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Lower bounds on the covering radius

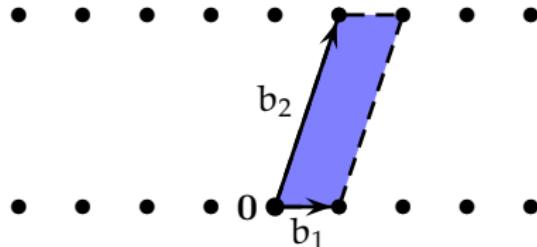


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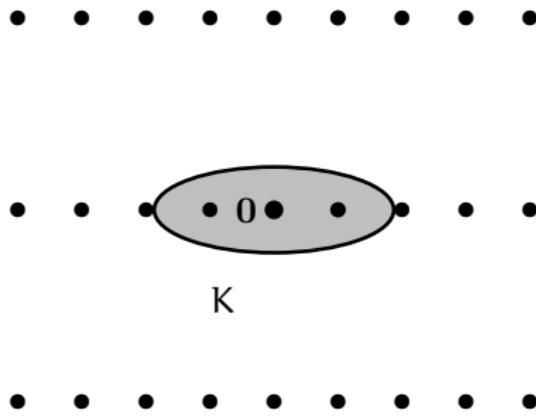
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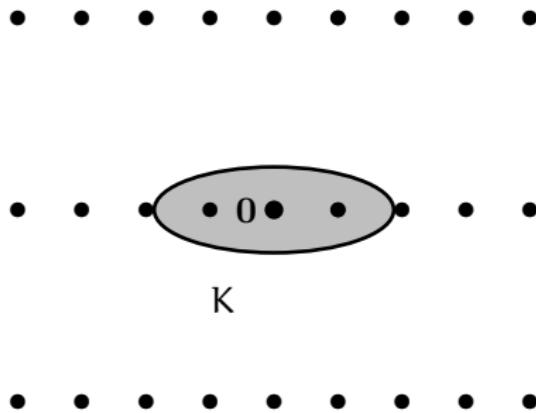
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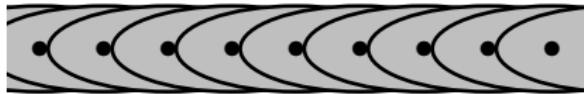
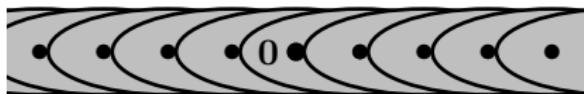
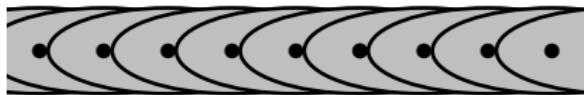
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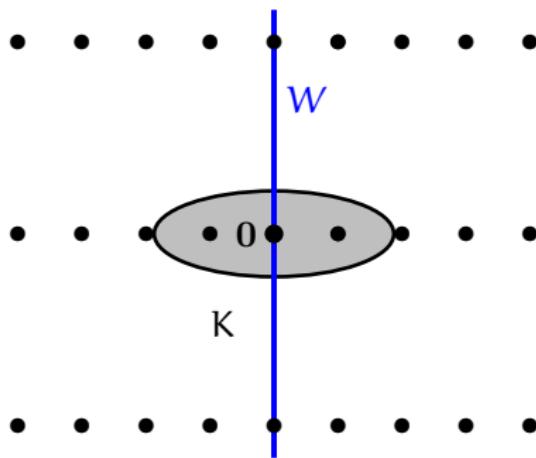
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- ▶ For any subspace $\mu(\mathcal{L}, K) \geq \mu(\Pi_W(\mathcal{L}), \Pi_W(K)) \geq (\frac{\det(\Pi_W(\mathcal{L}))}{\text{vol}(\Pi_W(K))})^{1/d}$

Main results

Theorem (R., Rothvoss '23)

Given a full rank lattice $\mathcal{L} \subseteq \mathbb{R}^n$ and convex $K \subseteq \mathbb{R}^n$, recall

$$\mu(\mathcal{L}, K) = \min\{r > 0 \mid \mathcal{L} + rK = \mathbb{R}^n\}$$

There exists a subspace $W \subseteq \mathbb{R}^n$ so that

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Theorem (R., Rothvoss '23, following Dadush '12)

For convex $K \subseteq \mathbb{R}^n$, we can find a point in $K \cap \mathbb{Z}^n$ in time $(O(\log n))^{4n}$.

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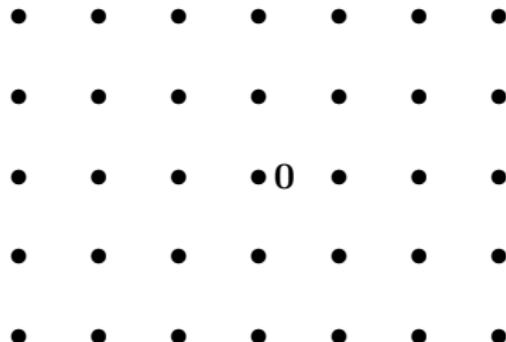
- ▶ Find 'nice' chain $\{\mathbf{0}\} = \mathcal{L}_0 \subset \cdots \subset \mathcal{L}_k = \mathcal{L}$ and take $U := \text{span}(\mathcal{L}_{i^*})$

Quotient lattices

Definition

Consider a lattice $\mathcal{L} \subseteq \mathbb{R}^n$ with a sublattice $\mathcal{L}' := \mathcal{L} \cap U$.

The **quotient lattice** is $\mathcal{L}/\mathcal{L}' := \Pi_{U^\perp}(\mathcal{L})$.

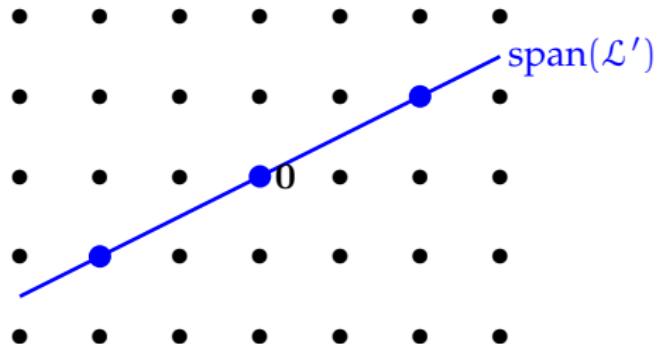


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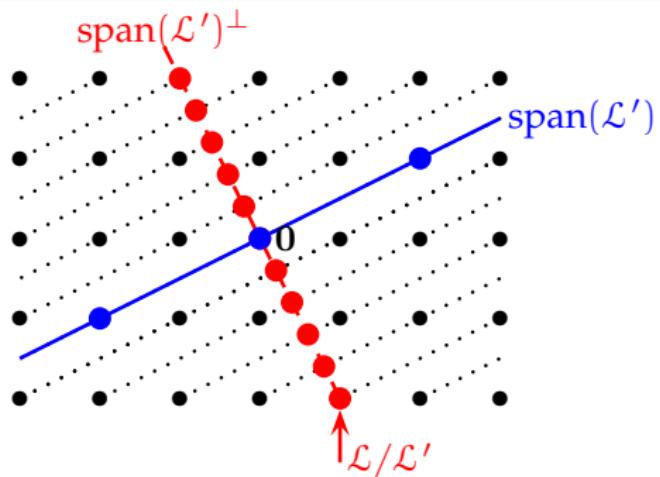


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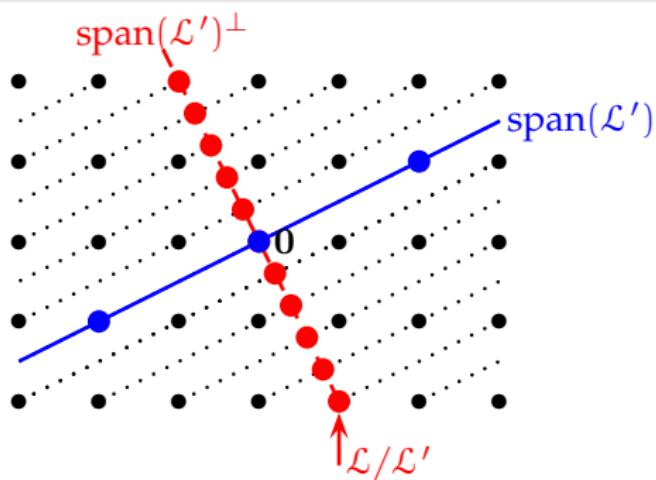


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- ▶ **Intuition:** We can factor \mathcal{L} into \mathcal{L}' and \mathcal{L}/\mathcal{L}'
- ▶ For example $\det(\mathcal{L}) = \det(\mathcal{L}') \cdot \det(\mathcal{L}/\mathcal{L}')$.

The canonical filtration

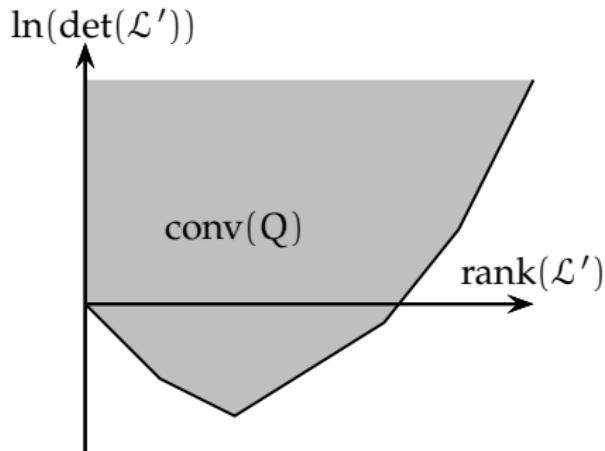
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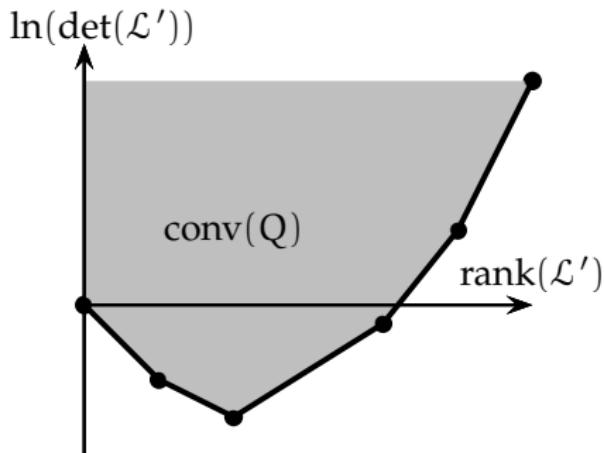


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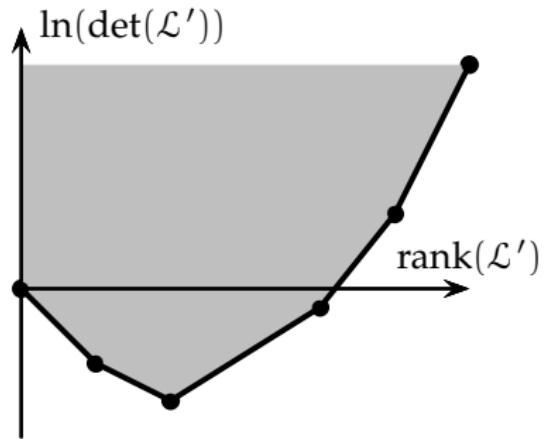
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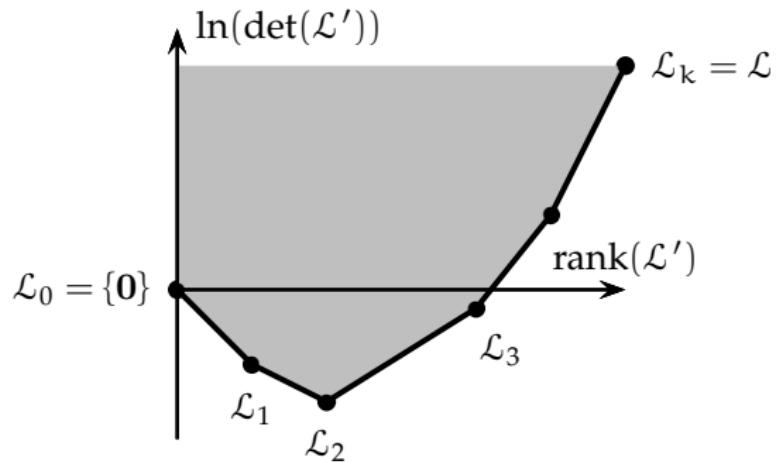
- ▶ Lower envelope of $\text{conv}(Q)$ is called **canonical polygon**



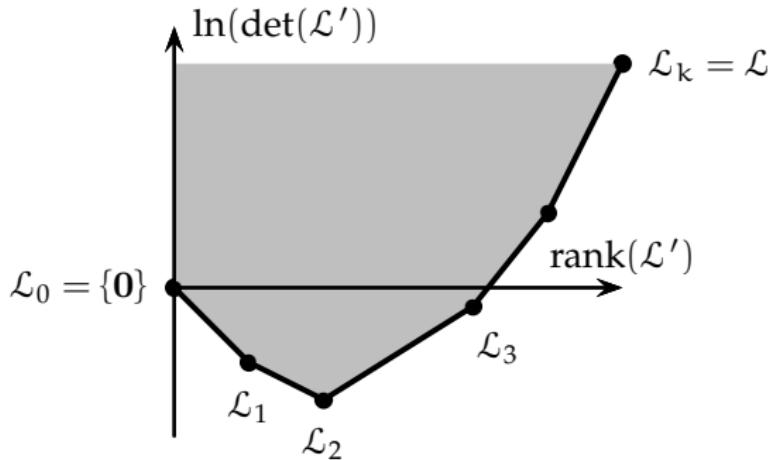
The canonical filtration (2)



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Theorem (Canonical filtration)

- (a) The vertices of the canonical plot form a chain

$$\{\mathbf{0}\} = \mathcal{L}_0 \subset \mathcal{L}_1 \subset \dots \subset \mathcal{L}_k = \mathcal{L}.$$

- (b) $r_i := \det(\mathcal{L}_i / \mathcal{L}_{i-1})^{1/rank(\mathcal{L}_i / \mathcal{L}_{i-1})}$ satisfy $r_1 < \dots < r_k$

- (c) Each $\frac{1}{r_i}(\mathcal{L}_i / \mathcal{L}_{i-1})$ has determinant = 1 and its sublattices have $\det \geq 1$.

The Reverse Minkowski Theorem

Reverse Minkowski (Regev, Stephens-Davidowitz '17)

Let $\mathcal{L} \subseteq \mathbb{R}^n$ with $\det(\mathcal{L}') \geq 1$ for all $\mathcal{L}' \subseteq \mathcal{L}$. Then for $s := C \log n$,

$$\rho_{1/s}(\mathcal{L}) = \sum_{x \in \mathcal{L}} \exp(-\pi s^2 \|x\|_2^2) \leq \frac{3}{2}$$

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- ▶ implies **Subspace Flatness** for **ellipsoids** [Dadush, Regev '16]

The Reverse Minkowski Theorem

Reverse Minkowski (Regev, Stephens-Davidowitz '17)

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Definition

A lattice \mathcal{L} is called stable if $\det(\mathcal{L}) = 1$ and $\det(\mathcal{L}') \geq 1$ for all $\mathcal{L}' \subseteq \mathcal{L}$.

Main proof

Take $U := \text{span}(\mathcal{L}_{i^*})$ for some $i^* \in [k]$ and bound

$$\mu(\mathcal{L} \cap U, K \cap U) \leq \sum_{i=1}^{i^*} \mu(\mathcal{L}_i / \mathcal{L}_{i-1}, \Pi_{\text{span}(\mathcal{L}_{i-1})^\perp}(K \cap \text{span}(\mathcal{L}_i)))$$

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Lemma

For symmetric convex K and stable lattice \mathcal{L} one has

$$\mu(\mathcal{L}, K) \lesssim \log(n) \cdot \ell_K,$$

where $\ell_K := \mathbb{E}_{x \sim N(0, I_n)} [\|x\|_K^2]^{1/2}$.

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- ▶ Therefore $\mu(\mathcal{L} \cap U, K \cap U) \lesssim \log(n) \cdot \ell_K \cdot \sum_{i=1}^{i^*} r_i$.

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- ▶ $\mu(\mathcal{L} \cap U, K \cap U) \lesssim \log(n) \cdot \ell_K \cdot \sum_{i=1}^{i^*} r_i.$
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$$r_{i^*} \leq \det(\Pi_W(\mathcal{L}))^{1/d}$$

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- ▶ Use $r_{i^*} < \dots < r_k$
- ▶ **Urysohn:** Euclidean ball minimizes width (fixing volume) where $w(K) = \mathbb{E}_{\theta \sim S^{d-1}} [\max\{\langle \theta, x - y \rangle : x, y \in K\}]$

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- ▶ Use previous lemma $\ell_{K^\circ \cap W} \leq \ell_{K^\circ}$

Main proof (3)

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Theorem (Figiel, Tomczak-Jaegerman, Pisier)

For any symmetric convex body $K \subseteq \mathbb{R}^n$, there is an invertible linear map $T: \mathbb{R}^n \rightarrow \mathbb{R}^n$ so that $\ell_{T(K)} \cdot \ell_{(T(K))^\circ} \lesssim n \log(n).$

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- ▶ Suffices to pick minimal i^* so that $\dim(U) \geq \frac{n}{2} \implies d \geq \frac{n}{2}$

Putting everything together

Theorem (R., Rothvoss '23)

Given a full rank lattice $\mathcal{L} \subseteq \mathbb{R}^n$ and symmetric convex $K \subseteq \mathbb{R}^n$, recall

$$\mu(\mathcal{L}, K) = \min\{r > 0 \mid \mathcal{L} + rK = \mathbb{R}^n\}$$

There exists a subspace $W \subseteq \mathbb{R}^n$ so that

$$\left(\frac{\det(\Pi_W(\mathcal{L}))}{\text{vol}(\Pi_W(K))} \right)^{1/\dim(W)} \leq \mu(\mathcal{L}, K) \lesssim \log^3(n) \cdot \left(\frac{\det(\Pi_W(\mathcal{L}))}{\text{vol}(\Pi_W(K))} \right)^{1/\dim(W)}.$$

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- ▶ Recall 'triangle inequality':

$$\mu(\mathcal{L}, K) \leq \underbrace{\mu(\mathcal{L} \cap U, K \cap U)}_{\text{pay a } \log^2(n)} + \underbrace{\mu(\Pi_{U^\perp}(\mathcal{L}), \Pi_{U^\perp}(K))}_{\text{induction: } \log^3(n/2)}$$

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- ▶ Remains to note $\log^2(n) + \log^3(n/2) \leq \log^3(n)$.

Generalization to non-symmetric K

- ▶ Translate K so that the barycenter of K° is at $\mathbf{0}$
- ▶ We need:

$$\text{vol}(\Pi_W(K))^{1/d} \lesssim \left(\frac{n}{d}\right)^3 \cdot \text{vol}(\Pi_W(K \cap (-K)))^{1/d}$$

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- ▶ This is a polar version of Rudelson's inequality:

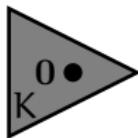
$$\text{vol}((K - K) \cap W)^{1/d} \lesssim \frac{n}{d} \cdot \max_{x \in \mathbb{R}^n} \left\{ \text{vol}(K \cap (x + W))^{1/d} \right\}$$

Other implications

Theorem (R., Rothvoss '23)

For full rank lattice $\mathcal{L} \subseteq \mathbb{R}^n$ and convex body $K \subseteq \mathbb{R}^n$ one has

$$\mu(\mathcal{L}, K) \leq O(\log^3(n)) \cdot \mu(\mathcal{L}, K - K)$$

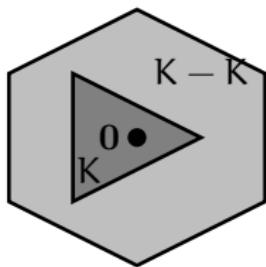


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For convex body $K \subseteq \mathbb{R}^n$ with $K \cap \mathbb{Z}^n = \emptyset$, there is a direction $c \in \mathbb{Z}^n$ with

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- ▶ Previously best known: $\lesssim n^{4/3} \log^{O(1)} n$ [Rudelson '98]

Open problem

IP in exponential time

Can we solve every n variable integer program in time $2^{O(n)}$?

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Thanks for your attention!