

# Vector Balancing in Lebesgue Spaces

Victor Reis

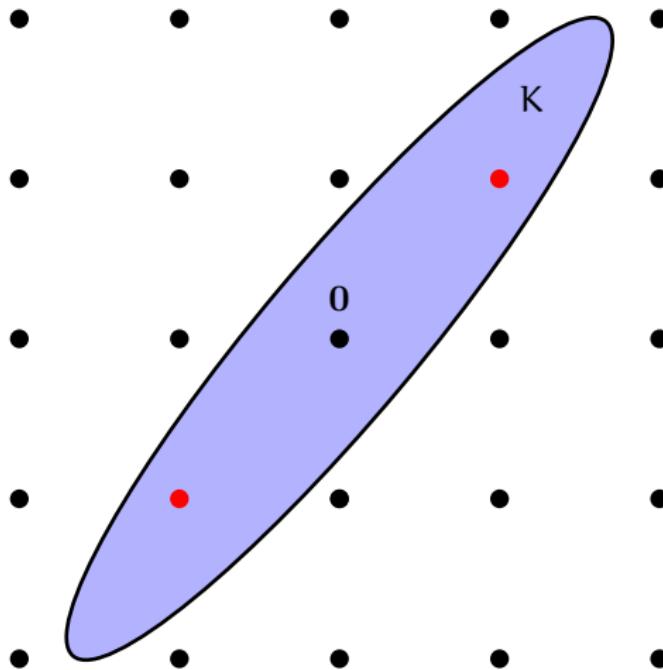
BIRS Combinatorial and Geometric Discrepancy Workshop  
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Joint work with Thomas Rothvoss



## Warmup: Minkowski's Theorem (1889)

- ▶ Any symmetric convex  $K \subset \mathbb{R}^n$  with volume  $> 2^n$  intersects  $\mathbb{Z}^n \setminus \{0\}$



# Non-constructive Partial Coloring

## Theorem (Gluskin, 1989)

Given symmetric convex  $K \subseteq \mathbb{R}^n$  with  $\text{vol}(K \cap [-1, 1]^n) \geq C^n$  for  $C > 1$  there exists  $y \in 2K \cap \{-1, 0, 1\}^n$  with  $\Omega(n)$  coordinates in  $\{-1, 1\}$ .

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- ▶ But  $\text{vol}(K \cap [-1, 1]^n) \geq 2^n \gamma_n(K)$  where

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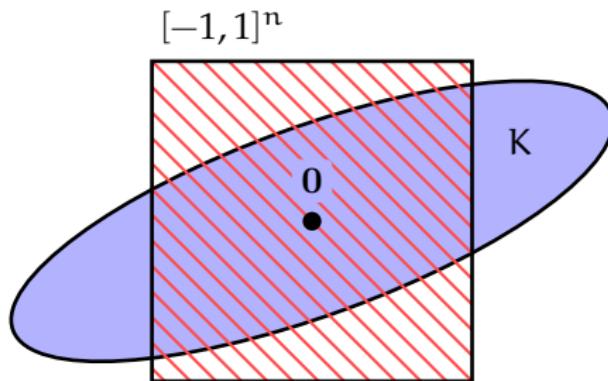
## Corollary

Given symmetric convex  $K \subseteq \mathbb{R}^n$  with  $\gamma_n(K) \geq 0.51^n$  there exists  $y \in 2K \cap \{-1, 0, 1\}^n$  with  $\Omega(n)$  coordinates in  $\{-1, 1\}$ .

# Constructive Partial Coloring

Theorem (Rothvoss, FOCS 2014)

Given symmetric convex  $K \subseteq \mathbb{R}^n$  with  $\gamma_n(K) \geq 0.999^n$  we can find  $y \in K \cap [-1, 1]^n$  with  $\Omega(n)$  coordinates in  $\{-1, 1\}$  in **polynomial time**.



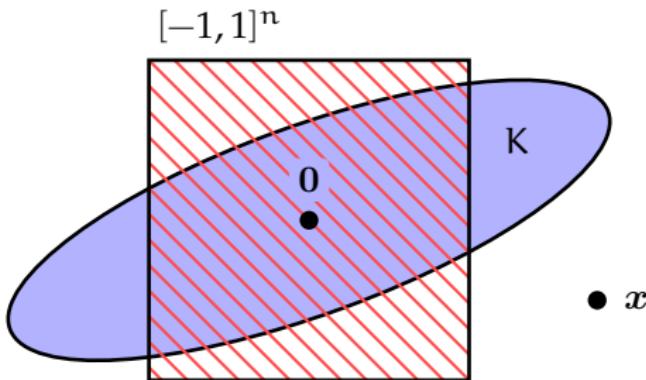
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- (1) Sample  $x \sim N(\mathbf{0}, I_n)$



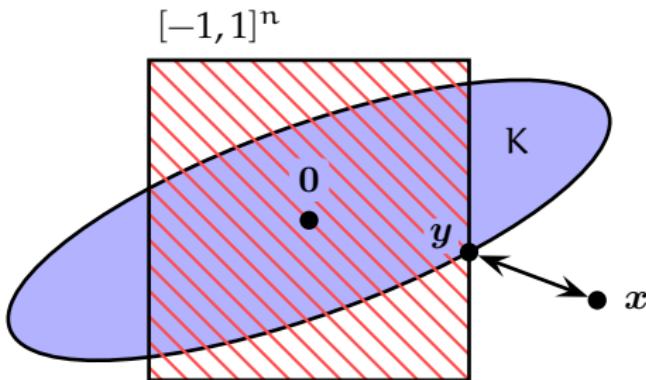
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**Algorithm:**

- (1) Sample  $x \sim N(\mathbf{0}, I_n)$
- (2) Output the projection  $y$  of  $x$  onto  $K \cap [-1, 1]^n$ .



# Constructive Partial Coloring

## Theorem (R., Rothvoss '20)

For every  $c \in (0, 1)$ , there exist  $\varepsilon, \delta \geq \text{poly}(c)$  and a poly-time alg. that given symmetric convex  $K \subseteq \mathbb{R}^n$  with  $\gamma_n(K) \geq c^n$  returns  $y \in (\frac{1}{\varepsilon}K) \cap [-1, 1]^n$  with  $\delta n$  coordinates in  $\{-1, 1\}$ .

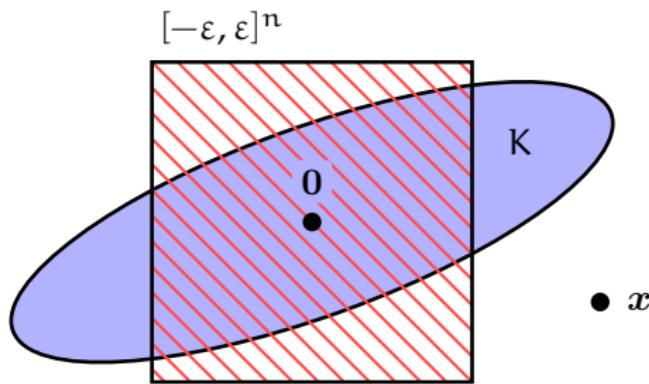
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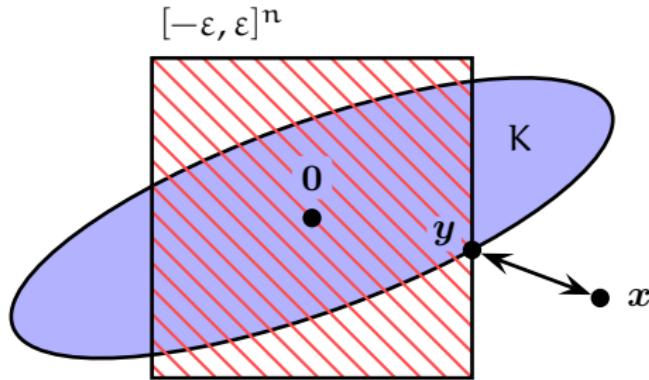
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## Algorithm:

- (1) Sample  $x \sim N(\mathbf{0}, I_n)$
- (2) Output  $1/\varepsilon$  times the projection  $y$  of  $x$  onto  $K \cap [-\varepsilon, \varepsilon]^n$ .



## Key technical lemma

### Lemma (R., Rothvoss '20)

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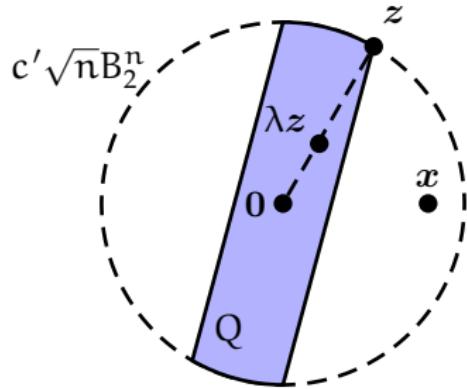
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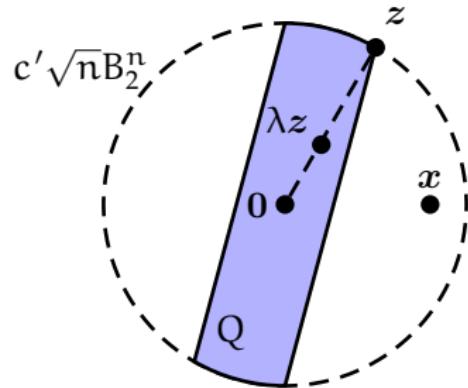


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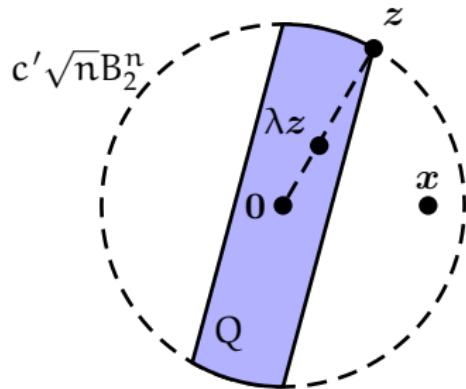
- ▶ Let  $Q := K \cap c' \sqrt{n} B_2^n$  and  $z := \operatorname{argmax}_Q \langle x, z \rangle$

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- ▶ Let  $Q := K \cap c' \sqrt{n} B_2^n$  and  $z := \operatorname{argmax}_Q \langle x, z \rangle$
- ▶ Upper bound  $d(x, K) \leq \|x - \lambda z\|_2$  for  $\lambda = \text{poly}(c)$ .

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- ▶ In fact, this already implies  $\gamma_n(K) \geq \text{poly}(c)^n$
- ▶ Proof uses polarity, Kubota integral formula and M-ellipsoids

## Application: vector balancing in $\ell_p$ norms

- Recall  $\|a\|_p := (|a_1|^p + \dots + |a_n|^p)^{1/p}$  (think  $p \in [2, \underbrace{\log n}_{\approx \infty}]$ )

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Given  $a_1, \dots, a_n \in \mathbb{R}^m$  with  $\|a_i\|_p \leq 1$ , find the minimum  $C_p(m, n)$  for which there always exist signs  $x \in \{-1, 1\}^n$  with

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## Theorem (Spencer, 1985)

$$C_\infty(m, n) \lesssim \sqrt{n \log(2m/n)}.$$

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- ▶ Combine Khintchine with an  $\ell_{p_0} \rightarrow \ell_\infty$  measure bound via

$$n^{1/p_0} B_{p_0}^m \cap B_\infty^m \subset n^{1/p} B_p^m$$

# Measure bound for $\ell_p \rightarrow \ell_\infty$ balancing

## Key Lemma (Measure bound for $\ell_p \rightarrow \ell_\infty$ )

Given  $A \in \mathbb{R}^{m \times n}$  with columns  $a_1, \dots, a_n \in B_p^m$  and rows

$A_1, \dots, A_m \in \mathbb{R}^n$ , the set  $K := \{x \in \mathbb{R}^n : \| \sum_{i=1}^n x_i a_i \|_\infty \leq \sqrt{p} \cdot n^{\frac{1}{2} - \frac{1}{p}} \}$  has

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$$\gamma_n(\{x \in \mathbb{R}^n : |\langle x, v \rangle| \leq 1\}) \geq 1 - \exp(-\|v\|_2^{-2}/2)$$

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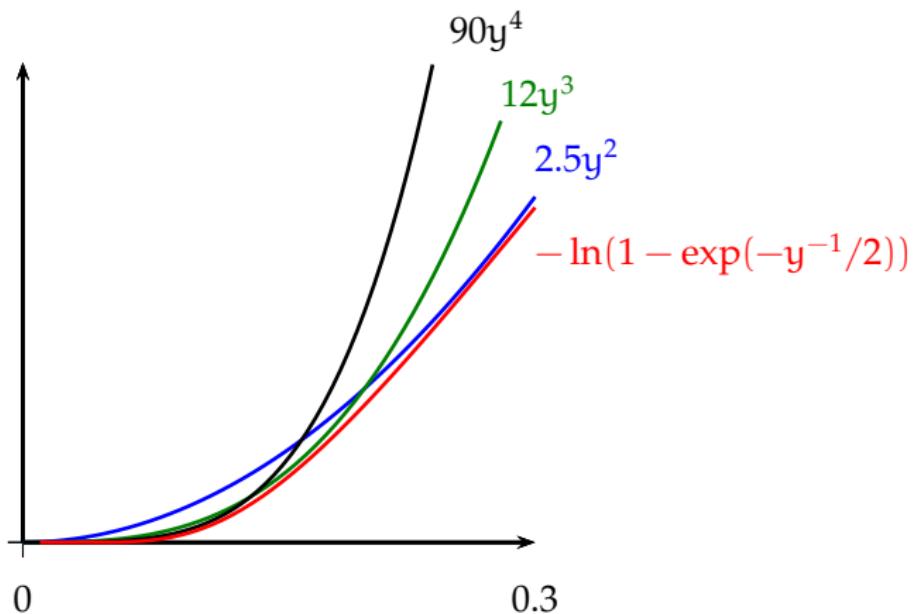
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$$\begin{aligned}\gamma_n(\{x \in \mathbb{R}^n : |\langle x, v \rangle| \leq 1\}) &\geq 1 - \exp(-\|v\|_2^{-2}/2) \\ &\geq \exp(-c' p^{p/2} \|v\|_2^p)\end{aligned}$$

and multiply!

# Key inequality

$$-\ln(1 - \exp(-y^{-1}/2)) \lesssim p^{p/2} y^{p/2}$$



## Last step: bootstrapping

- ▶ Combining the previous Lemma with Khintchine we can show

$$\gamma_n \left( \left\{ x \in \mathbb{R}^n : \left\| \sum_{i=1}^n x_i a_i \right\|_p \leq \sqrt{p_0} \cdot n^{1/2 - 1/p_0 + 1/p} \cdot m^{1/p_0 - 1/p} \right\} \right) \geq c^n$$

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- ▶ For larger  $p$ , we can take  $p_0 := \max(2, 2 \log(m/n))$
- ▶ In any case we get our improved measure lower bound:

$$\gamma_n \left( \left\{ x \in \mathbb{R}^n : \left\| \sum_{i=1}^n x_i a_i \right\|_p \leq \sqrt{n \cdot \min(p, \log(\frac{2m}{n}))} \right\} \right) \geq c^n$$

# Conclusion

## Main theorem (R., Rothvoss' 20)

Let  $n \leq m$  and  $1 \leq p \leq q \leq \infty$ . Given  $\mathbf{a}_1, \dots, \mathbf{a}_n \in \mathbb{R}^m$  with  $\|\mathbf{a}_i\|_p \leq 1$  we have in poly-time  $\mathbf{x} \in [-1, 1]^n$  with  $n/2$  coordinates in  $\pm 1$  so that

$$\left\| \sum_{i=1}^n x_i \mathbf{a}_i \right\|_q \lesssim \sqrt{\min\left(p, \log\left(\frac{2m}{n}\right)\right)} \cdot n^{\max(0, 1/2 - 1/p) + 1/q}.$$

## Theorem: Full coloring

Let  $n \leq m$  and  $1 \leq p \leq q \leq \infty$  with  $\max(0, 1/2 - 1/p) + 1/q > 0$ . Given  $\mathbf{a}_1, \dots, \mathbf{a}_n \in \mathbb{R}^m$  with  $\|\mathbf{a}_i\|_p \leq 1$  we have in poly-time  $\mathbf{x} \in \{-1, 1\}^n$  so that

$$\left\| \sum_{i=1}^n x_i \mathbf{a}_i \right\|_q \lesssim \frac{\sqrt{\min\left(p, \log\left(\frac{2m}{n}\right)\right)}}{\max(0, 1/2 - 1/p) + 1/q} \cdot n^{\max(0, 1/2 - 1/p) + 1/q}.$$

## Corollary: Beck-Fiala with many ones

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- Take  $p = 4$  and  $q = \infty$  in previous theorem:  $\|Ax\|_\infty \lesssim n^{1/4} t^{1/4} \leq \sqrt{t}$ .

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- ▶ Take  $p = 4$  and  $q = \infty$  in previous theorem:  $\|Ax\|_\infty \lesssim n^{1/4} t^{1/4} \leq \sqrt{t}$ .
- ▶ For  $t \leq n$  we also get an upper bound  $\sqrt{t} \cdot \log(2n/t)$ .

# Open problems

- ▶ Generalized Komlós Conjecture:

Let  $n \leq m$  and  $2 \leq p \leq q \leq \infty$ . Given  $a_1, \dots, a_n \in \mathbb{R}^m$  with  $\|a_i\|_p \leq 1$  can we find  $x \in \{-1, 1\}^n$  so that

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- ▶ Generalized Matrix Spencer Conjecture:

Let  $n \leq m^2$  and  $2 \leq p \leq q \leq \infty$ . Given  $A_1, \dots, A_n \in \mathbb{R}^{m \times m}$  with  $\|A_i\|_p \leq 1$  can we find  $x \in \{-1, 1\}^n$  so that

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Thanks for your attention!