CS 164 Discussion Notes

Earley's Algorithm

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Algorithm

The state set at input position k is called S(k). The parser is seeded with S(0) consisting of only the top-level rule. The parser then repeatedly executes three operations:

- **Prediction**: For every state in S(k) of the form $(X \to \alpha \cdot Y \beta, j)$ (where j is the origin position as above), add $(Y \to \cdot \gamma, k)$ to S(k) for every production in the grammar with Y on the left-hand side $(Y \to \gamma)$.
- **Scanning**: If a is the next symbol in the input stream, for every state in S(k) of the form $(X \to \alpha \cdot a \beta, j)$, add $(X \to \alpha a \cdot \beta, j)$ to S(k+1).
- **Completion**: For every state in S(k) of the form $(X \to \gamma \cdot, j)$, find states in S(j) of the form $(Y \to \alpha \cdot X \beta, i)$ and add $(Y \to \alpha X \cdot \beta, i)$ to S(k).

Note: duplicate states are not added to the state set, only new ones. These three operations are repeated until no new states can be added to the set.

https://en.wikipedia.org/wiki/Earley_parser



Example

Consider the following grammar:

```
P \rightarrow E \rightarrow
E \rightarrow E + E
E \rightarrow E \times E
E \rightarrow ID
```



Example (con't)

Consider the following input string:

$$ID + ID * ID$$

And show whether or not:

- it is part of the language defined by our grammar;
- our grammar is ambiguous.

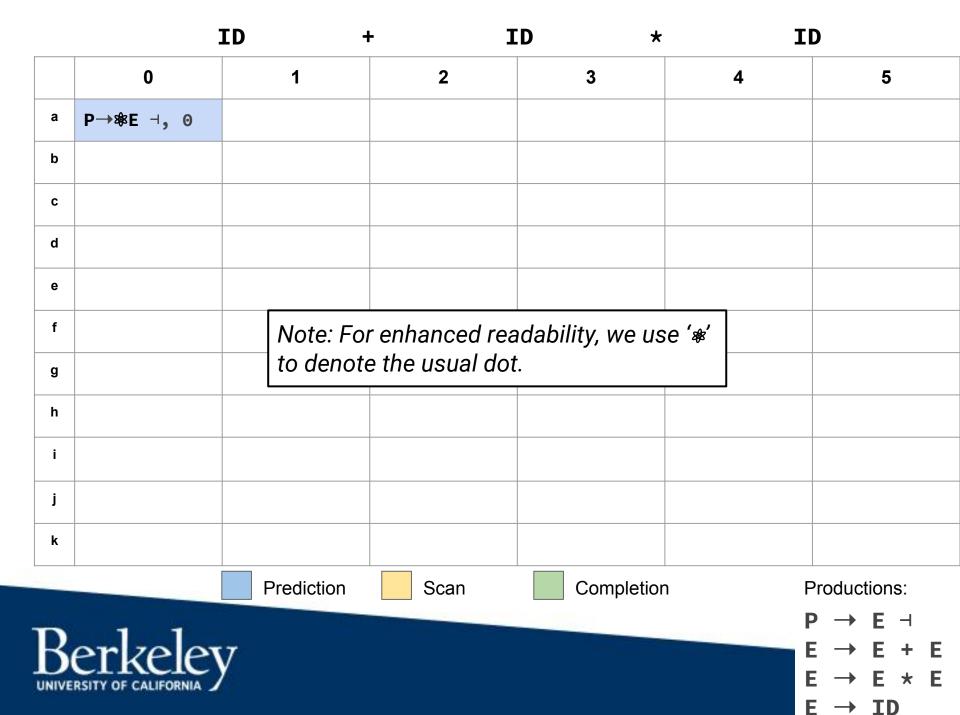


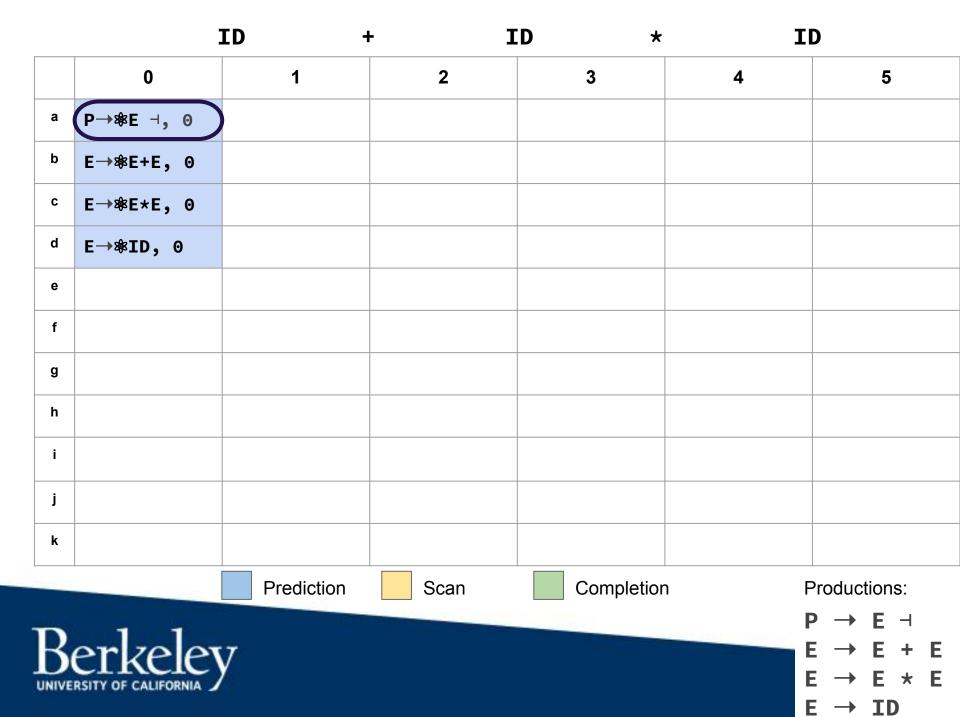
Representation

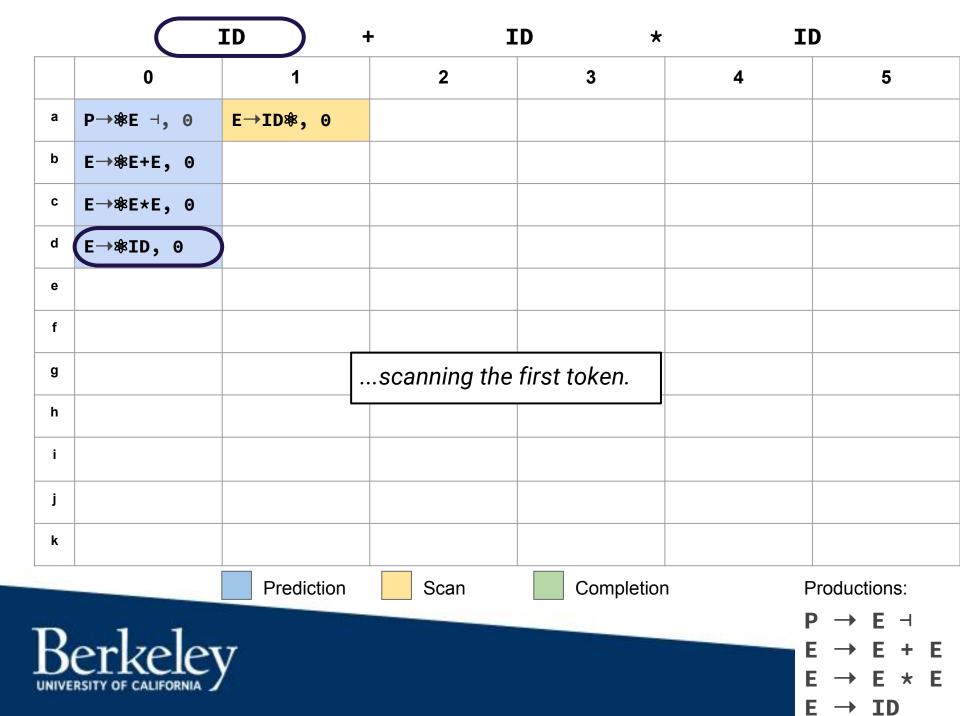
We represent our state sets in a table, where each column denotes a state set, and we use the row indexes (a-k) to trace the origin of newly produced completions. Our first item in S(0) is based on the start symbol of our grammar:

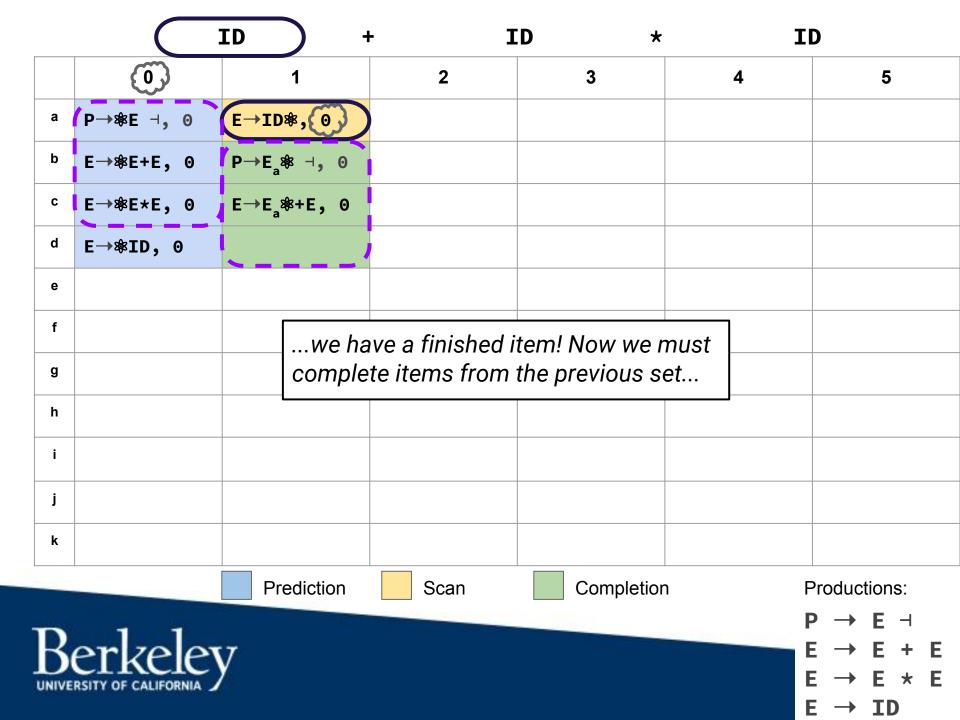
$$P \rightarrow E \dashv$$

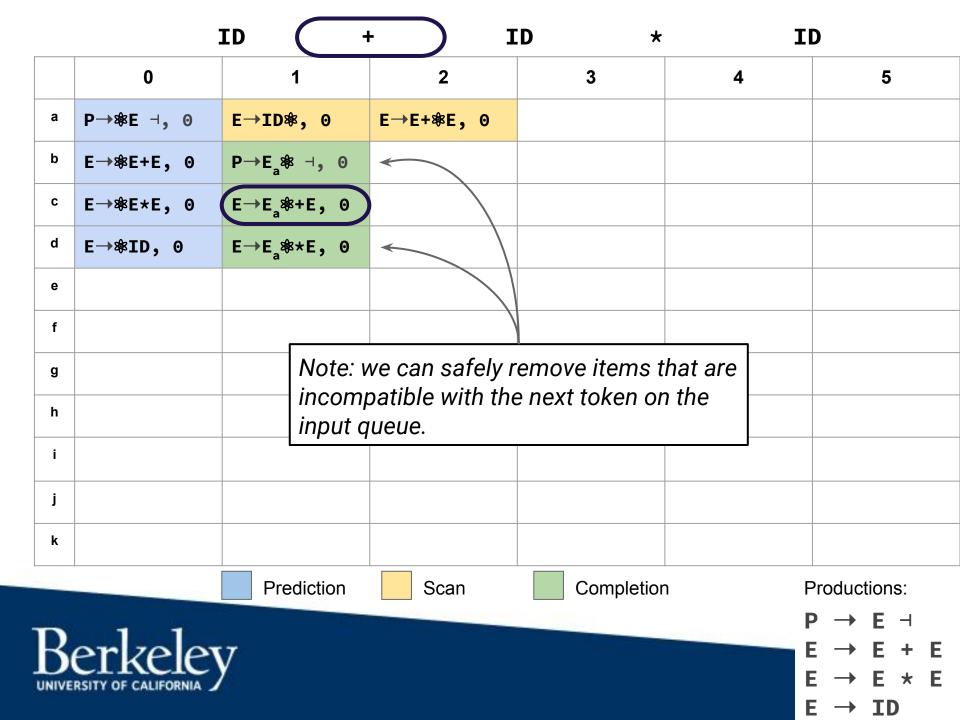


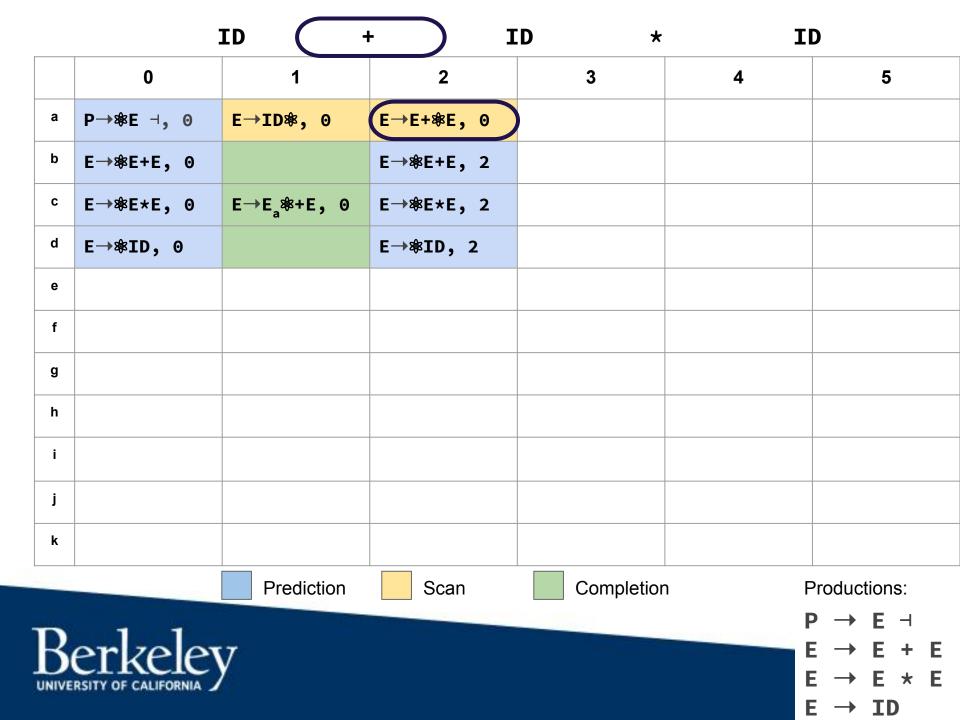


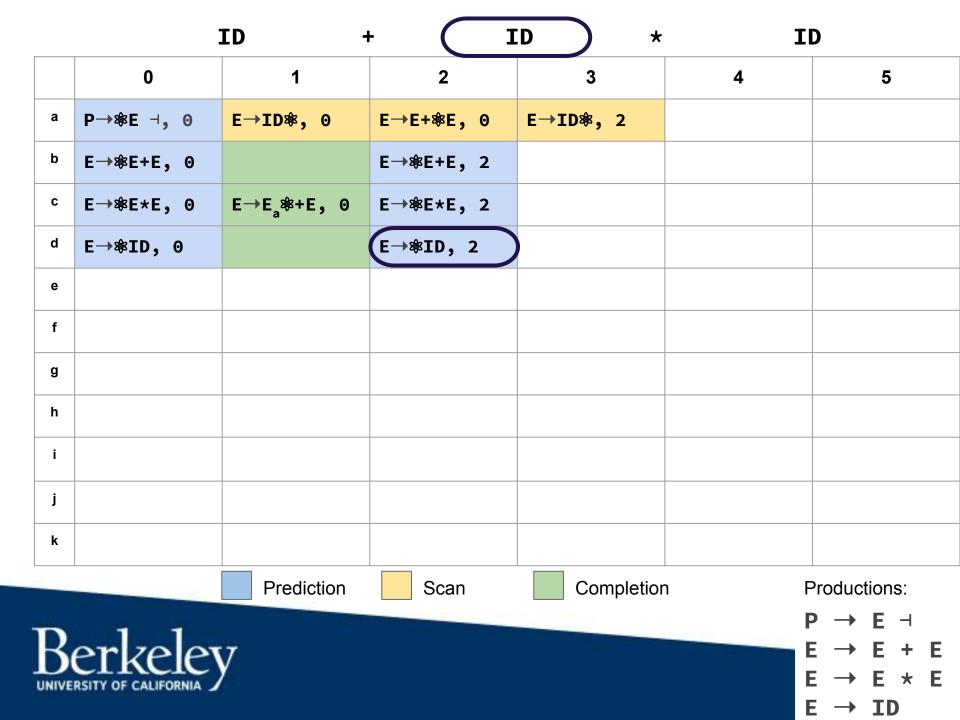


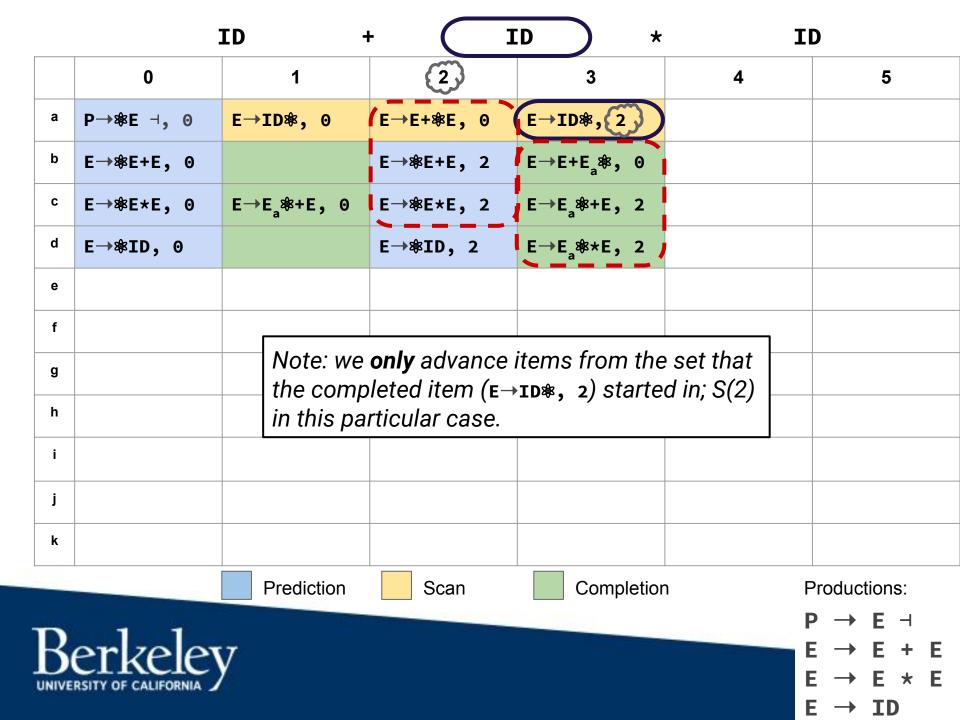


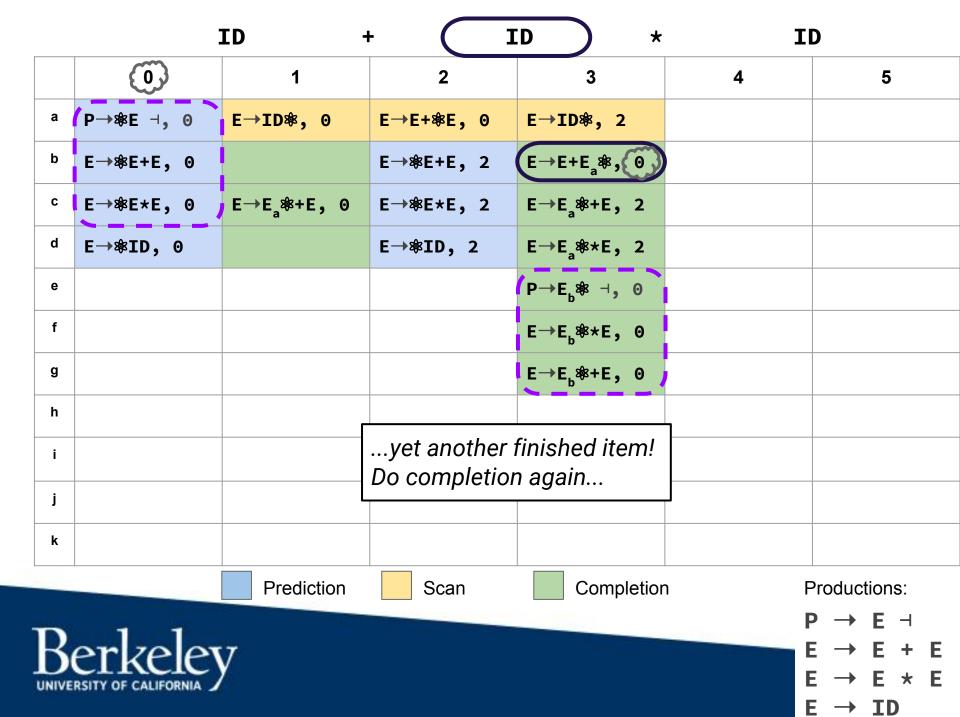


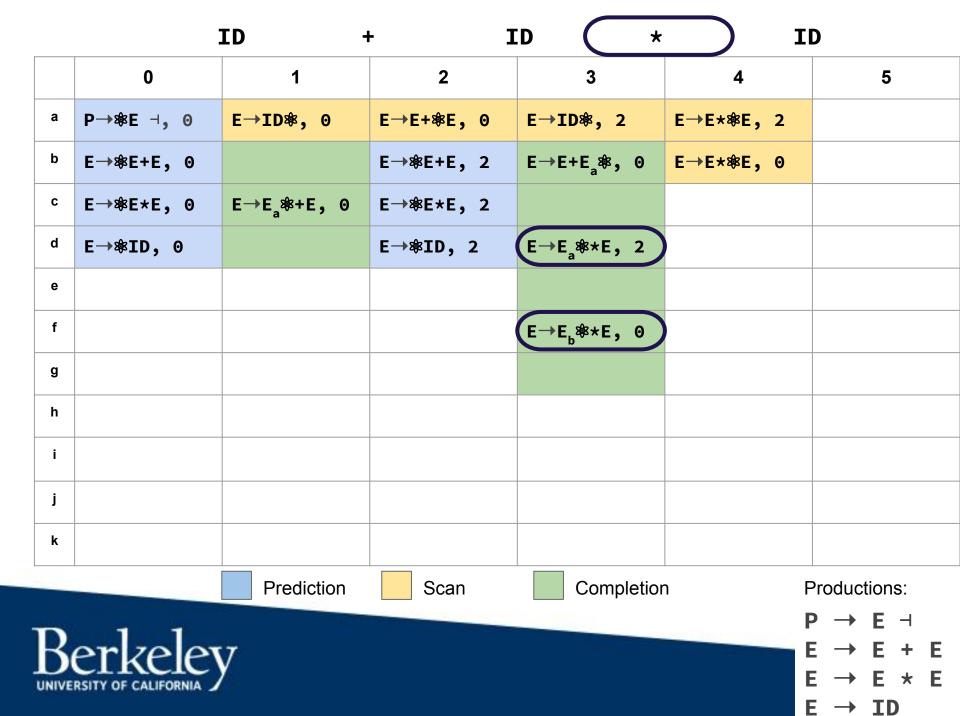


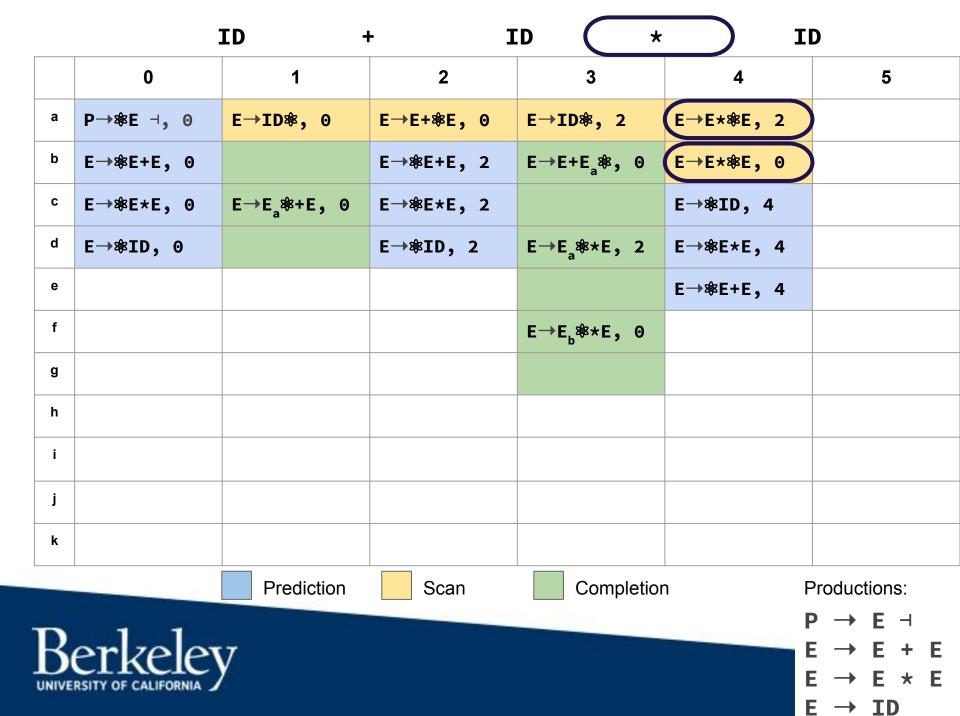












		ID +	· I	:D *		
	0	1	2	3	4	5
а	P→ 寒E ⊣, 0	E→ID*, 0	E→E+*E, 0	E→ID*, 2	E→E* *E, 2	E→ID*, 4
b	E→ \$E+E, 0		E→ \$E+E, 2	E→E+E _a ≉, 0	E→E**E, 0	
С	E→ ≇E*E, 0	E→E _a *+E, 0	E→ 8E*E, 2		E→*ID, 4	
d	E→ \$ID, 0		E→ \$ID, 2	E→E _a **E, 2	E→ \$ E*E , 4	
е					E→ * E+E, 4	
f				E→E _b ***E, 0		
g						
h						
i						
j						
k						
		Prediction	Scan	Completion	n F	Productions:
D	11					$P \rightarrow E +$

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 $E \rightarrow E + E$ $E \rightarrow E * E$ $E \rightarrow ID$

	:	ID -	:D *		D		
	0	1	2	3	4)	5	
а	P→ 8E ⊣, 0	E→ID≉, 0	E→E+ \$E, 0	E→ ID ≉, 2	E→E* *E, 2	E→ID*,4	
b	E→ * E+E, 0		E→ * E+E, 2	E→E+E _a ≉, 0	E→E*®E, 0	E→E*E _a ≉, 2	
С	E→ * E*E, 0	E→E _a *+E, 0	E→ ≉E*E, 2		E→ %ID , 4	E→E*E _a ≉, 0	
d	E→≉ID, 0		E→ *ID , 2	E→E _a **E, 2	E→ \$E*E, 4	E→E _a **E, 4	
е					E→ \$E+E, 4	E→E _a *+E, 4	
f				E→E _b ***E, 0			
g							
h							
i							
j							
k							
		Prediction	Scan	Completion	ำ	Productions:	
B	$\begin{array}{c} P \to E \dashv \\ E \to E + E \\ E \to E \star E \end{array}$						

 $\mathsf{E} \to \mathsf{ID}$

ID + ID * ID							
	0	1	(2)	3	4	5	
а	P→ 寒E ⊣, 0	E→ID≉, 0	E→E+ \$E, 0	E→ID≉, 2	E→E* \$E, 2	E→ID ≉, 4	
b	E→ \$E+E, 0		E→ 8E+E, 2	E→E+E _a ≉, 0	E→E* \$E, 0	E→E*E _a *, 2	
С	E→ \$E*E, 0	E→E _a *+E, 0	E→ ®E *E, 2		E→ *ID , 4	E→E*E _a ≉, 0	
d	E→ *ID , 0		E→ *ID , 2	E→E _a **×E, 2	E→ \$E*E, 4	E→E _a **E, 4	
е					E→ 8E+E, 4	E→E _a *+E, 4	
f				E→E _b **×E, 0		E→E+E _b ≉, 0	
g						E→E _b *+E, 2	
h						E→E _b ®×E, 2	
i							
j							
k							
	Prediction Scan Completion Productions:						
BUNIVE	$\begin{array}{c} \text{Berkeley} \\ \text{university of california} \end{array}$						

 $E \rightarrow ID$

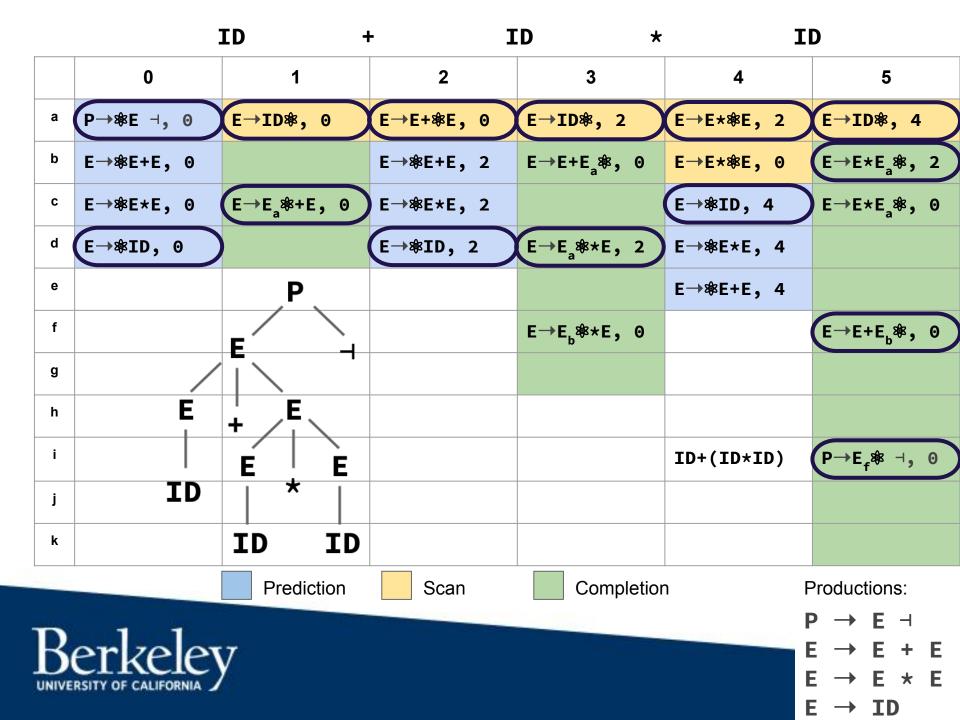
		ID +	· I	:D *	I	D	
	(0)	1	2	3	4	5	
а	/P→ ≉E ⊣, 0	E→ID≉, 0	E→E+ *E, 0	E→ID*, 2	E→E* \$E, 2	E→ ID ≉, 4	
b	E→ ≉E+E, 0		E→ \$E+E, 2	E→E+E _a ≉, 0	E→E* \$E, 0	E→E*E _a ≉, 2	
С	E→®E*E, 0	E→E _a ≉+E, 0	E→ \$E*E, 2		E→≋ID, 4	E→E*E _a ≉, 0	
d	E→≉ID, 0		E→ \$ID, 2	E→E _a ***E, 2	E→ 8E*E, 4	E→E _a ***E, 4	
е					E→ ®E+E , 4	E→E _a ≉+E, 4	
f				E→E _b **×E, 0		E→E+E _b *, 0	
g						E→E _b *+E, 2	
h						E→E _b ***E, 2	
i						P→E _f ≉ ⊣, 0	
j						E→E _f ≉+E, 0	
k						E→E _f ≉*E, 0	
		Prediction	Scan	Completion	ו [Productions:	
D	$P \rightarrow E \dashv$						
В	erkelev	V				$E \to E + E$	
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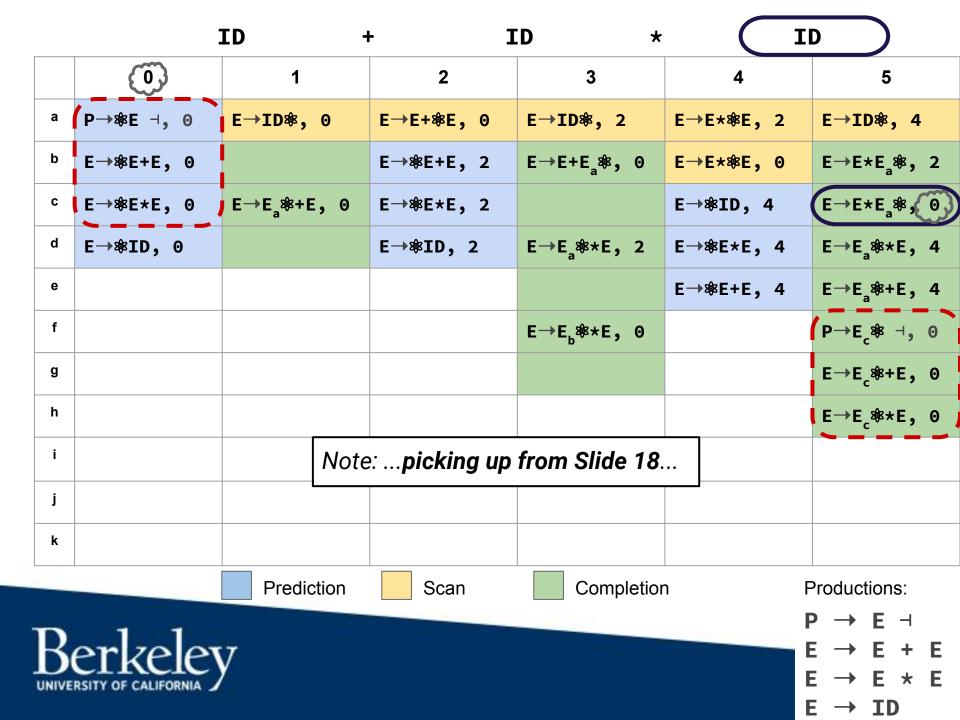
 $E \rightarrow ID$

		ID +	· I	:D *		
	0)	1	2	3	4	5
а	P→ 8E ⊣, 0	E→ID≉, 0	E→E+ % E, 0	E→ID*, 2	E→E* *E, 2	E→ID*, 4
b	E→ \$E+E, 0		E→ \$E+E, 2	E→E+E _a ≉, 0	E→E**E, 0	E→E*E _a *, 2
С	E→ \$E*E, 0	E→E _a %+E, 0	E→ \$E*E, 2		E→ %ID , 4	E→E*E _a ≉, 0
d	E→ \$ID, 0		E→ \$ID, 2	E→E _a ***E, 2	E→ 8E*E, 4	
е					E→ 8E+E, 4	
f				E→E _b ≉*E, 0		E→E+E _b ≉, 0
g						
h						
i			we recogn	nized	ID+(ID*ID)	$P \rightarrow E_f \gg \neg, 0$
j			the input str	ring!		
k						
		Prediction	Scan	Completion	า F	Productions:
n	1 1					$P \rightarrow E \dashv$



 $P \rightarrow E + E$ $E \rightarrow E * E$ $E \rightarrow ID$



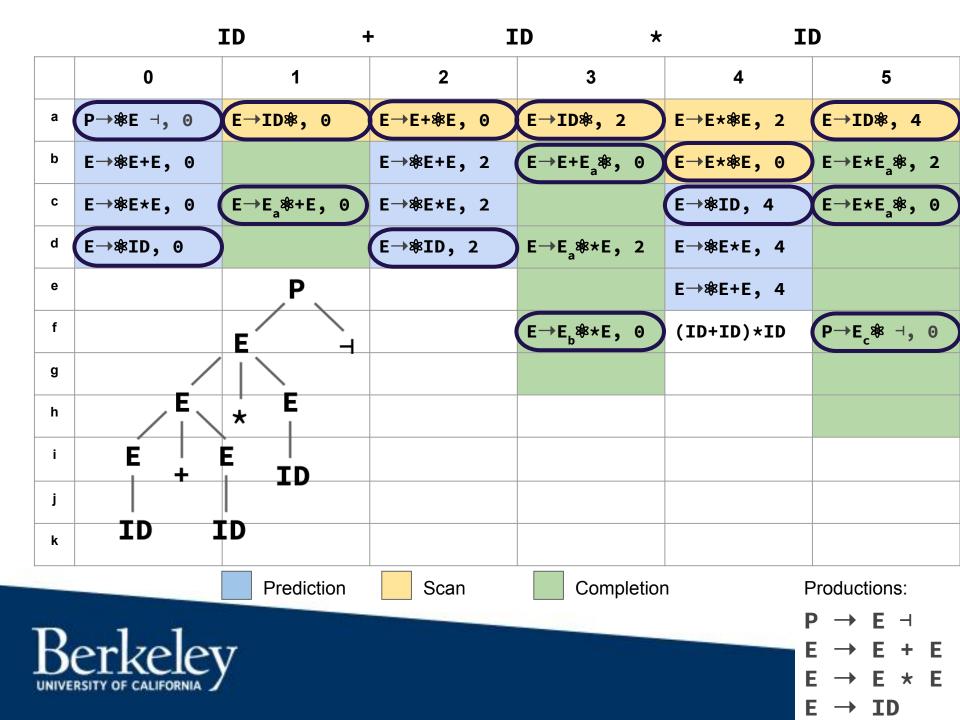


		ID -	+ 1	:D *	I	D
	(0)	1	2	3	4	5
а	P→8 E ⊣, 0	E→ID*, 0	E→E+ \$E, 0	E→ID*, 2	E→E* *E, 2	E→ID*, 4
b	E→ \$E+E, 0		E→ ≉E+E, 2	E→E+E _a ≉, 0	E→E* *E, 0	E→E*E _a ≉, 2
С	E→ \$E*E, 0	E→E _a ≉+E, 0	E→ ≉E*E, 2		E→≉ID, 4	E→E*E _a ≉, 0
d	E→≋ID, 0		E→≉ID, 2	E→E _a ***E, 2	E→ ≉E*E, 4	
е					E→ 8E+E, 4	
f				E→E _b ***E, 0	(ID+ID)*ID	$P \rightarrow E_c * \dashv , \emptyset$
g						
h						
i		we	found another	valid parse tre	ee!	
j						
k						
		Prediction	Scan	Completion	n F	Productions:
B	erkelev	J				$P \rightarrow E + E$

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 $E \rightarrow E \star E$

 $E \rightarrow ID$



Acceptance

We found two *complete* Earley items that end with '¬' in our last state set. Because these items span the entire input string (they starts at *zero*), the input string is accepted and hence is part of the language.



Parse Tree Construction

We constructed a parse tree for each accepting state by tracing back the items that led to that accepting state. Since we found more than one tree, our grammar is *ambiguous*.

