

Today's Content:

- 1) N Sorted arrays ↗ print completed data in sorted order
↘ store complete sorted data in single array
- 2) Given 2 sorted array find k smallest pair sum

- 1) Given 2D mat[N][M], every row is sorted, merge entire data into 1D sorted list & return sorted list

$N \times M$
Ex: mat[4][6]

	0	1	2	3	4	5
0	2	7	10	17	25	34
1	-6	0	1	8	11	14
2	3	4	6	14	21	26
3	7	10	14	19	23	27

merged 1D sorted array

-6 0 1 2 3 4

Idea

- 1) Insert all elements in a list
sort & return it

// elements = $N \times M$

TC: $NM \log(NM)$ SC: $O(1)$

- 2) Apply merge $\approx N$ times

TC: $O(MN \log N)$

$r_1 \dots M \}$ $2M$
 $r_2 \dots M \}$ $3M$
 $r_3 \dots M \}$ $4M$
 $r_4 \dots M \}$ $5M$
 \vdots
 $r_{n-1} \dots M \}$ $(N-1)M$
 $r_n \dots M$
 $N \times M$

$2M + 3M + 4M + \dots + NM$

TC: $M(2 + 3 + \dots + N)$

$= M(1 + 2 + 3 + \dots + N - 1)$

$= M \left[\frac{(N)(N+1)}{2} - 1 \right]$

TC: $O(MN \log N)$

idea4: Pointer approach

0	1	2	3	4	5
2	7	10	17 ^{P1}	25	34
-6	0	1	8	9	10 ^{P2}
3	4	6	14 ^{P3}	21	26
7	10 ^{P4}	14	19	23	27

Idea:

N pointers, every iteration print pointer with min val & move pointer. continue process till all pointers reaches end

output: -6 0 1 2 3 4 6 7 7 8 9 10 10 ... ?

How to implement? → U & min heap

0	1	2	3	4	5
2	7	10	17	25	34
-6	0	1	8	11	16
3	4	6	14	21	26
7	10	14	19	23	27

1. for every number we should also store its row number.
2. for every number we should also store its col number.

<2, 0, 0>	<0, 1, 1>
<-6, 1, 0>	<1, 1, 2>
<3, 2, 0>	<8, 1, 3>
<7, 3, 0>	<7, 0, 1>

output:

-6 0 1 2 ...

$\begin{matrix} \downarrow r \\ \downarrow c \end{matrix}$	$\begin{matrix} \downarrow r \\ \downarrow c \end{matrix}$	$\begin{matrix} \downarrow r \\ \downarrow c \end{matrix}$	$\begin{matrix} \downarrow r \\ \downarrow c \end{matrix}$
<-6, 1, 0>	<0, 1, 1>	<1, 1, 2>	<2, 0, 0>

→ Each data point:

: pair<data, pair<row, col>>

↳ pair<int, pair<int, int>>

↳ // please check how to do in your language of choice?

minheap < ^{data} pair<int, ^{row, col} pair<int, int>> >> mh.

- a) Single element in minheap } a) custom class
- b) It is sorted based on data } b) override comparator

int[] Merge(int mat[N][M]) {

minheap< pair<int, pair<int, int>>> mh

i = 0; i < N; i++) {

// value we need insert: mat[i][0]

mh.insert(pair(mat[i][0], pair(i, 0)))

}

list<int> ans;

while(mh.size() > 0) {

pair<int, pair<int, int>> data = mh.getMin()
^{first} ^{row col} ^{second}

mh.deleteMin()

int val = data.first

int r = data.second.first, c = data.second.second

ans.add(val) // insert in list

if(c+1 < M) {

mh.insert(pair(mat[r][c+1], pair(r, c+1)))

}

}
 return ans;

{ 8:50 }

Obs: heap size = N TC: $NM \{ \log N \} + NM \{ \log N \} \Rightarrow O(NM \log N)$
SC: $O(N)$

28) k Smallest Pair sums:

Given 2 sorted arrays, print k smallest pair sums, a single pair only 1 time

N 0 1 2 3 4 5 6 $A[i] + B[j]$
a[s]: 2 5 8 11 13

M
b[t]: 3 7 9 12 15 16 20

K=5 output

		A[]	B[]
Pairs:	Sum	i	j
1	5	0	0
2	8	1	0
3	9	0	1
4	11	0	2
5	11	2	0

Idea 1: Store all pair sum in a list sort it & get first k elements.

pair = $N \times M$

TC: $N \times M \log N \times M$ SC: $O(N \times M)$

Idea 2: Optimization using min heap & size = k

Insert all pair in min heap & get only min k pair sums in heap

TC: $N \times M \log k$ SC: $O(k)$

Idea: 0 1 2 3 4 5 6

a[s]: 2 5 8 11 13

b[t]: 3 7 9 12 15 16 20

<u>Pairs</u> :	<u>Sum</u>	i	j
1	5	0	0
2	8	1	0
3	9	0	1

Dabba: operations

<5, 0, 0>

<8, 1, 0>

<9, 0, 1>

<11, 2, 0>

<12, 1, 1>

<12, 1, 1>

<11, 0, 2>

[Issue, duplicate insertion is not okay. check if pair is already present]

delete Insert

i	j	i+1	j	i	j+1
0	0	1	0	0	1

1 0 → 2 0 1 1

0 1 → 1 1 0 2

Dabba:

→ getmin()	} minheap
→ deletemin()	
→ insert()	
→ search()	→ <u>hashset</u>

Q Use both minheap & hashset

→ // Store i, j as a pair

→ // Convert pair into single string

→ String s = "i_j"

```
void kpairs(int a[], int b[]) {
```

```
    int n = a.length, m = b.length;
```

```
    minheap < pair<int, pair<int, int>>> mh;
```

```
    HashSet<String> hs // To store pairs
```

```
    mh.insert(pair(a[0]+b[0], pair(0, 0))); // Insert min pair:
```

```
    hs.insert("0 0"); // inserting in hashset
```

```
    for (int i = 1; i <= k; i++) {
```

```
        pair<int, pair<int, int>> d = mh.getmin();
```

```
        mh.delete(d);
```

```
        int val = d.first, r = d.second.first, c = d.second.second
```

```
        print(val)
```

```
        // for r, c possibilities → {r+1, c}, {r, c+1}
```

```
        if (c+1 < m && hs.contains(to_string(r) + " " + to_string(c+1)) == false)
```

```
            val = a[r] + b[c+1]
```

```
            mh.insert(pair(val, pair(r, c+1)))
```

```
            hs.insert(to_string(r) + " " + to_string(c+1))
```

```
        if (r+1 < n && hs.contains(to_string(r+1) + " " + to_string(c)) == false)
```

```
            val = a[r+1] + b[c]
```

```
            mh.insert(pair(val, pair(r+1, c)))
```

```
            hs.insert(to_string(r+1) + " " + to_string(c))
```

```
    }
```

```
}
```

TC: $k * (\log k + \log k + \log k) = O(k \log k)$ SC: k

2 insertions 1 deletion