

1

1.1

| Buchstabe | Kodierung | " k_i/r " |
|-----------|-----------|----------------|
| L | 000 | $\frac{3}{13}$ |
| E | 001 | $\frac{1}{13}$ |
| O | 010 | $\frac{2}{13}$ |
| | 011 | $\frac{1}{13}$ |
| Ä | 100 | $\frac{3}{13}$ |
| S | 101 | $\frac{1}{13}$ |
| T | 110 | $\frac{1}{13}$ |

1.2

$$\begin{aligned} I &= - \sum_{i=1}^n p_i \log p_i \\ &= - \sum_{i=1}^7 p_i \log p_i \\ &= - \left(\frac{1}{7} \cdot \log \frac{1}{7} \right) \cdot 7 = - \log \frac{1}{7} \approx 2.81 \end{aligned}$$

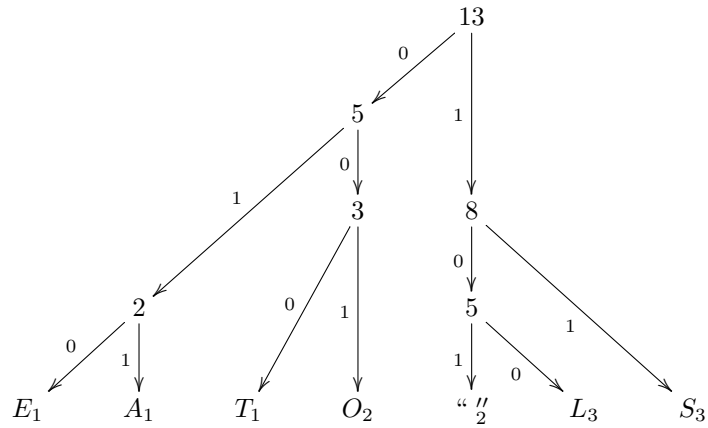
1.3

$$\begin{aligned} I &= - \sum_{i=1}^n p_i \log p_i \\ &= - \left(\frac{3}{13} \cdot \log \frac{3}{13} + \frac{1}{13} \cdot \log \frac{1}{13} + \frac{2}{13} \cdot \log \frac{2}{13} + \frac{2}{13} \cdot \log \frac{2}{13} \right. \\ &\quad \left. + \frac{1}{13} \cdot \log \frac{1}{13} + \frac{3}{13} \cdot \log \frac{3}{13} + \frac{1}{13} \cdot \log \frac{1}{13} \right) \approx 2.06 \end{aligned}$$

1.4

Der Informationsgehalt ist um etwa 0.75 verringert.

1.5



1.6

| Buchstabe | Kodierung |
|-----------|-----------|
| S | 11 |
| L | 100 |
| | 101 |
| O | 001 |
| T | 000 |
| A | 011 |
| E | 010 |

1.7

$$N = n_S + n_L + n + n_O + n_T + n_A + n_E = 3 \cdot 2 + 3 \cdot 3 + 2 \cdot 3 + 3 + 3 + 3 = 30$$

1.8

$$b = r \cdot I \approx 13 \cdot 2.06 = 26.78$$

1.9

Der Huffman-Kode weicht um etwa 12% ab.

1.10

Der Huffman-Kode braucht $\frac{30}{13} \approx 2.31$ Zeichen, dies weicht um etwa 12% vom Informationsgehalt $I \approx 2.06$ ab.

2

2.1

| a | b | c | $\neg a \vee b$ | $(\neg a \vee b) \wedge \neg c$ | $(a \vee c)$ | $b \wedge (a \vee c)$ | $f(a, b, c)$ |
|---|---|---|-----------------|---------------------------------|--------------|-----------------------|--------------|
| 0 | 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 0 | 1 | 1 | 1 | 1 | 1 |
| 1 | 1 | 1 | 1 | 0 | 1 | 1 | 1 |

2.2

$$(\neg a \wedge \neg b \wedge \neg c) \vee (\neg a \wedge \neg b \wedge c) \vee (\neg a \wedge b \wedge c) \vee (a \wedge b \wedge \neg c) \vee (a \wedge b \wedge c) = m_0 \vee m_2 \vee m_3 \vee m_6 \vee m_7 = \sum m(0, 2, 3, 6, 7)$$

2.3

$$\neg y = \sum m(1, 4, 5)$$

2.4

$$\neg y = \prod M(0, 2, 3, 6, 7)$$

2.5

$$y = \prod M(1, 4, 5)$$

2.6

?

3

3.1

$$f_{x_1=0} = f(0, x_2, x_3) = \underbrace{[(x_2 \vee x_3) \wedge 0]}_0 \vee \underbrace{[(x_2 \wedge (x_2 \vee x_3) \wedge x_3) \wedge 0]}_0 = 0$$

$$f_{x_1=1} = f(0, x_2, x_3) = \underbrace{[(x_2 \vee x_3) \wedge 1]}_{x_2 \vee x_3} \vee \underbrace{[(x_2 \wedge (x_2 \vee x_3) \wedge x_3) \wedge 1]}_{x_2 \wedge (x_2 \vee x_3) \wedge x_3} = (x_2 \vee x_3) \vee [x_2 \wedge (x_2 \vee x_3) \wedge x_3]$$

3.2

$$f_{x_1=1, x_2=0} = f(0, 0, x_3) = \underbrace{(0 \vee x_3)}_{x_3} \vee \underbrace{[0 \wedge (0 \vee x_3) \wedge x_3]}_0 = x_3$$

$$f_{x_1=1, x_2=1} = f(0, 1, x_3) = \underbrace{(1 \vee x_3)}_1 \vee \underbrace{[1 \wedge (1 \vee x_3) \wedge x_3]}_{x_3} = x_3$$

3.3

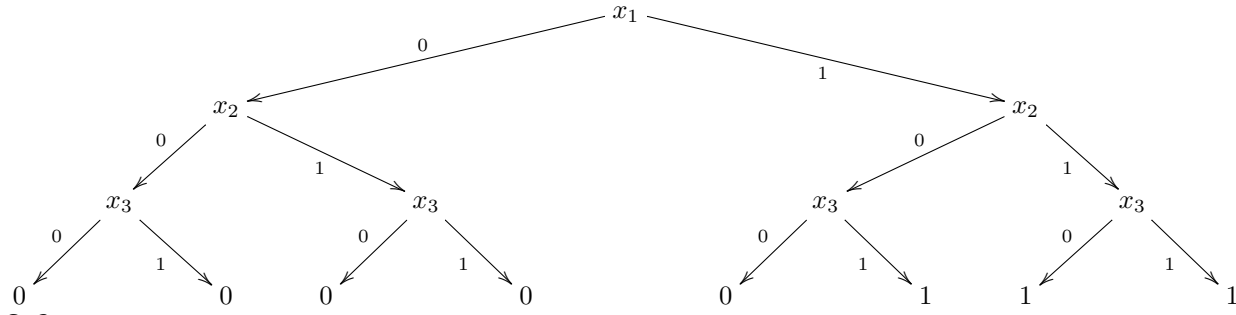
$$f_{x_1=1, x_2=1, x_3=0} = f(0, 1, 0) = 0$$

$$f_{x_1=1, x_2=1, x_3=1} = f(0, 1, 1) = 1$$

3.4

$$y = \underbrace{(x_1 \wedge \neg x_2 \wedge x_3)}_{101} \vee \underbrace{(x_1 \wedge x_2 \wedge x_3)}_{111} \vee \underbrace{(x_1 \wedge x_2 \wedge \neg x_3)}_{110} = \sum m(5, 6, 7)$$

3.5



3.6

