



EE290C - Fall 2018

Advanced Topics in Circuit Design

VLSI Signal Processing

Lecture 3: Finite Wordlength
Effects

CORDIC



Announcements

- Assignment 1 – due today
- Assignment 2 – due on Thursday 9/6
 - Chisel bootcamp 2.3-2.5, 3.1-3.2



Projects

- OFDM software radio (LTE, WiFi)
 - Bluetooth modem
 - GPS receiver
 - Smart NIC
 - Wellness monitor
-
- #students on project = #unique blocks



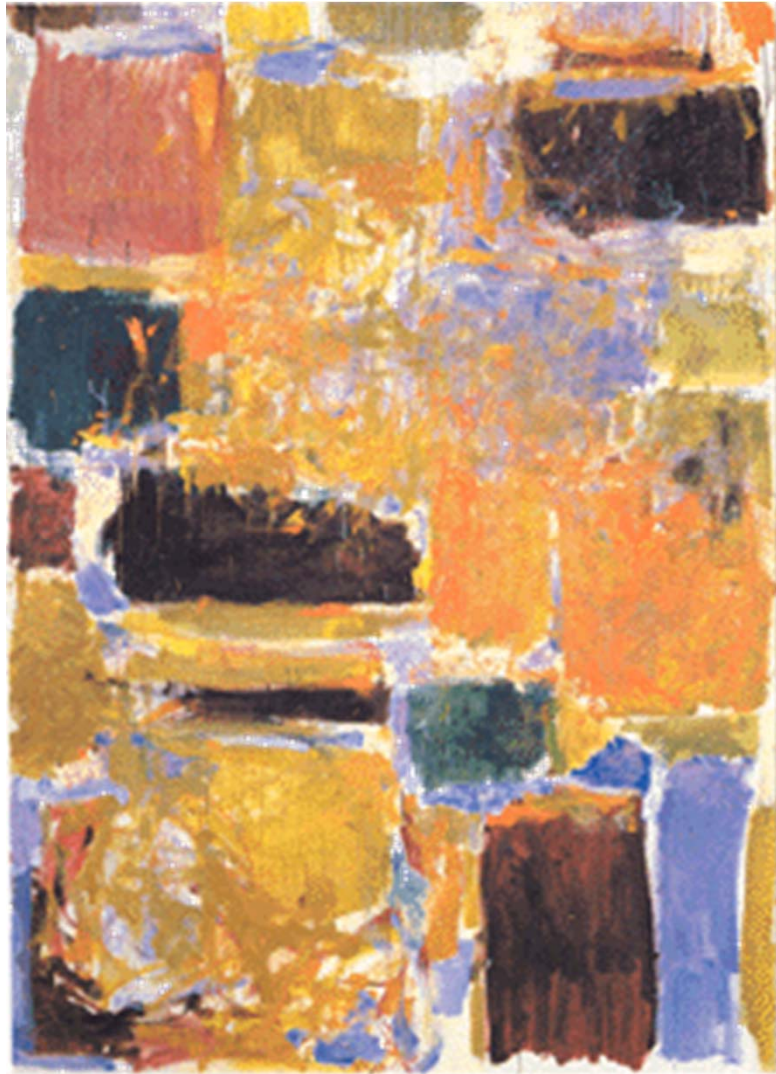
Reading

- **Numbers**

- Markovic, Brodersen, DSP Architecture Design Essentials, 2012, (Ch. 5)

- **CORDIC**

- Markovic, Brodersen, DSP Architecture Design Essentials, 2012, (Ch. 6)
 - Any other book above

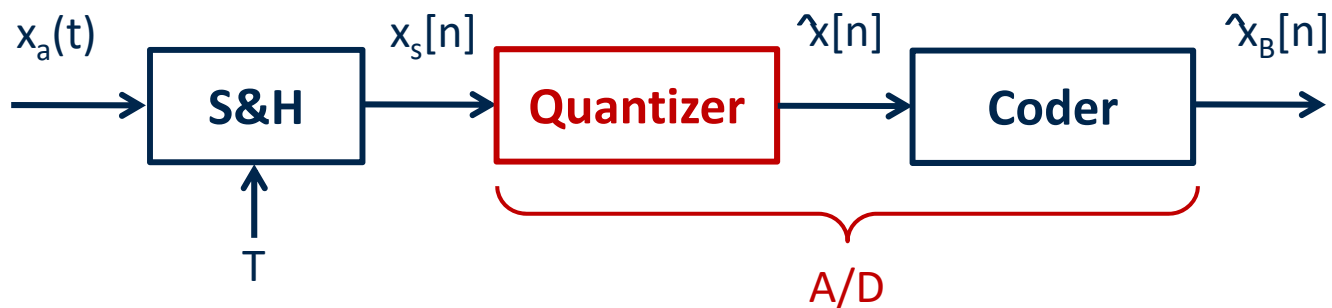
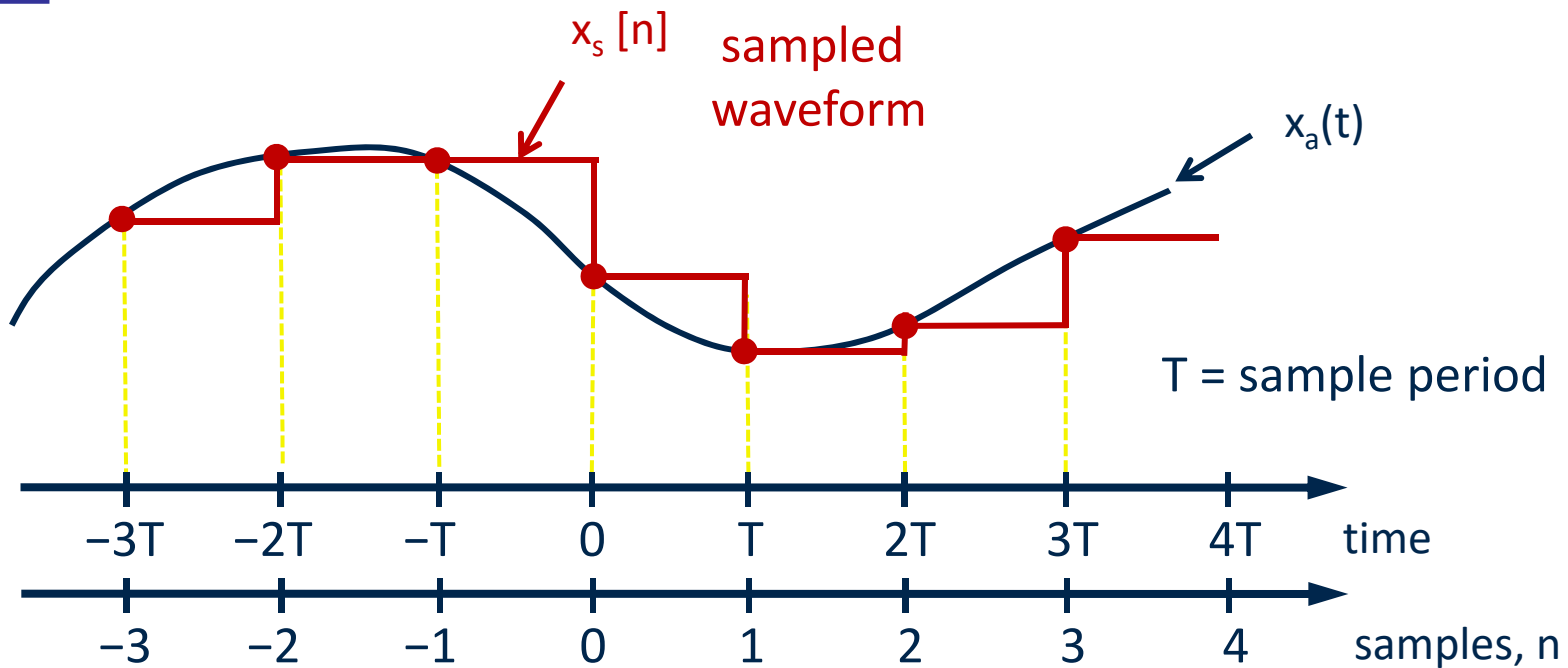


Advanced Topics in Circuit Design

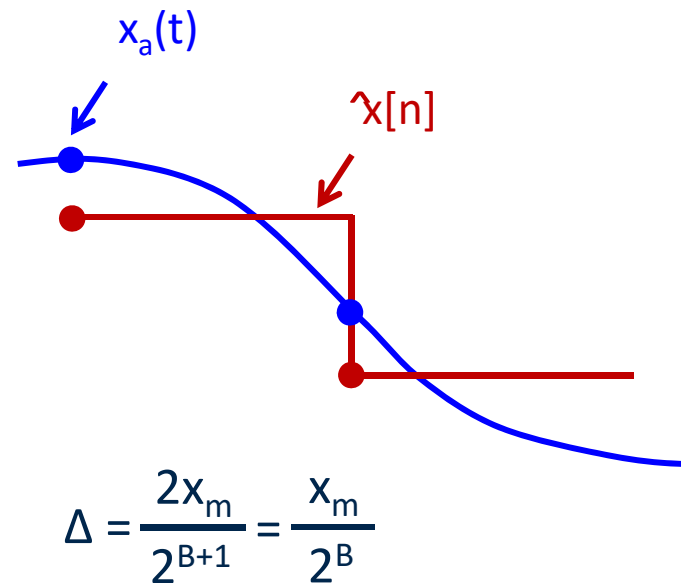
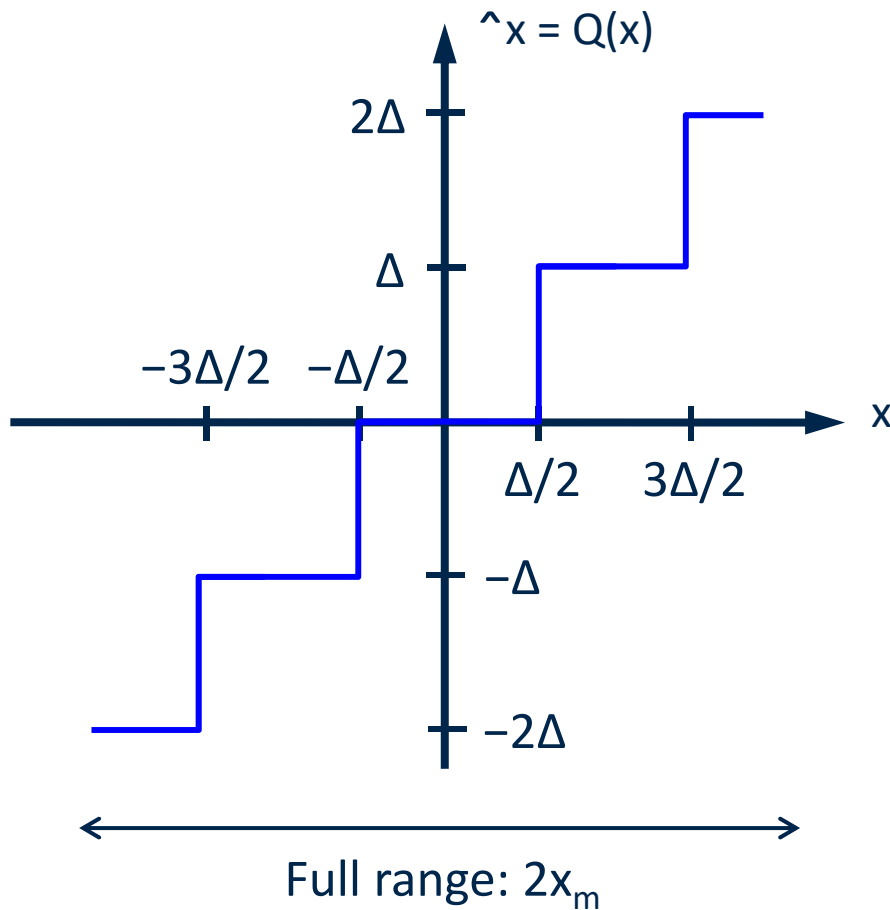
VLSI Signal Processing

Finite Wordlength Effects

Quantization Effects



Quantization



$B = \#$ bits of quantization

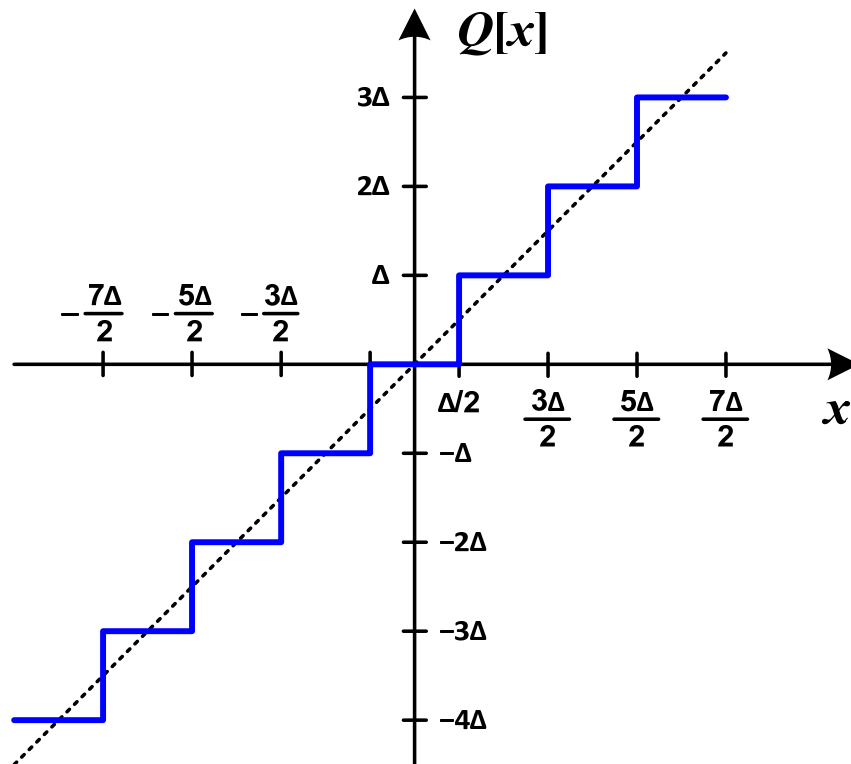
$$e[n] = \hat{x}[n] - x[n]$$
$$-\Delta/2 < e[n] \leq \Delta/2$$

(AWN process)

2's complement representation

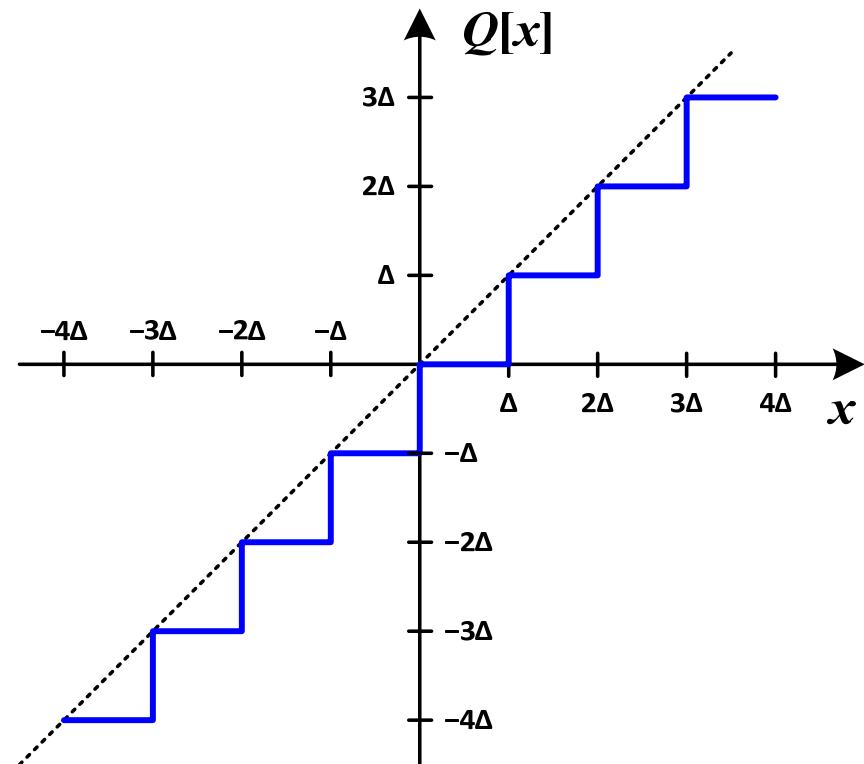
Quantization: Rounding, Truncation

Rounding

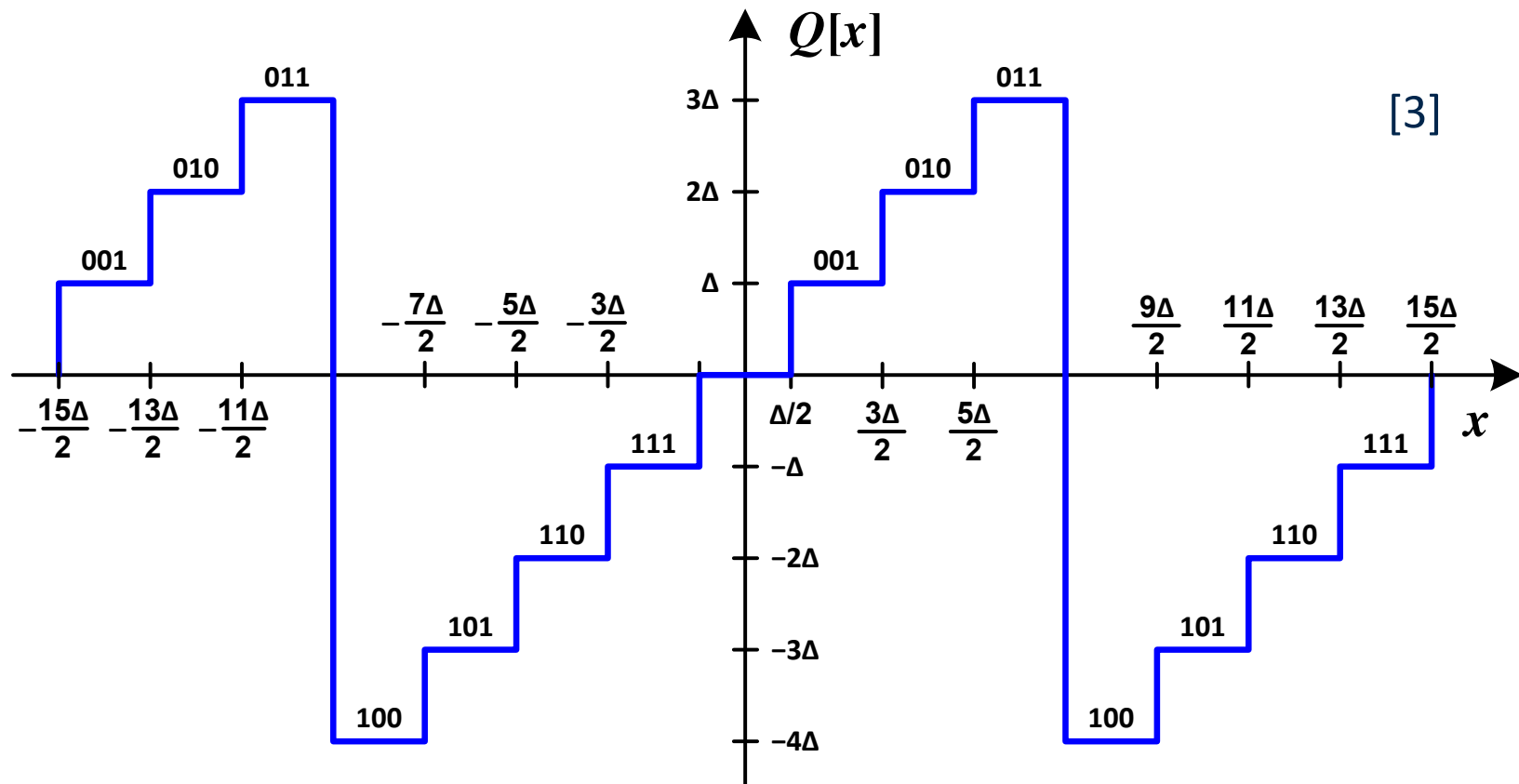


Feedback systems use
rounding

Truncation

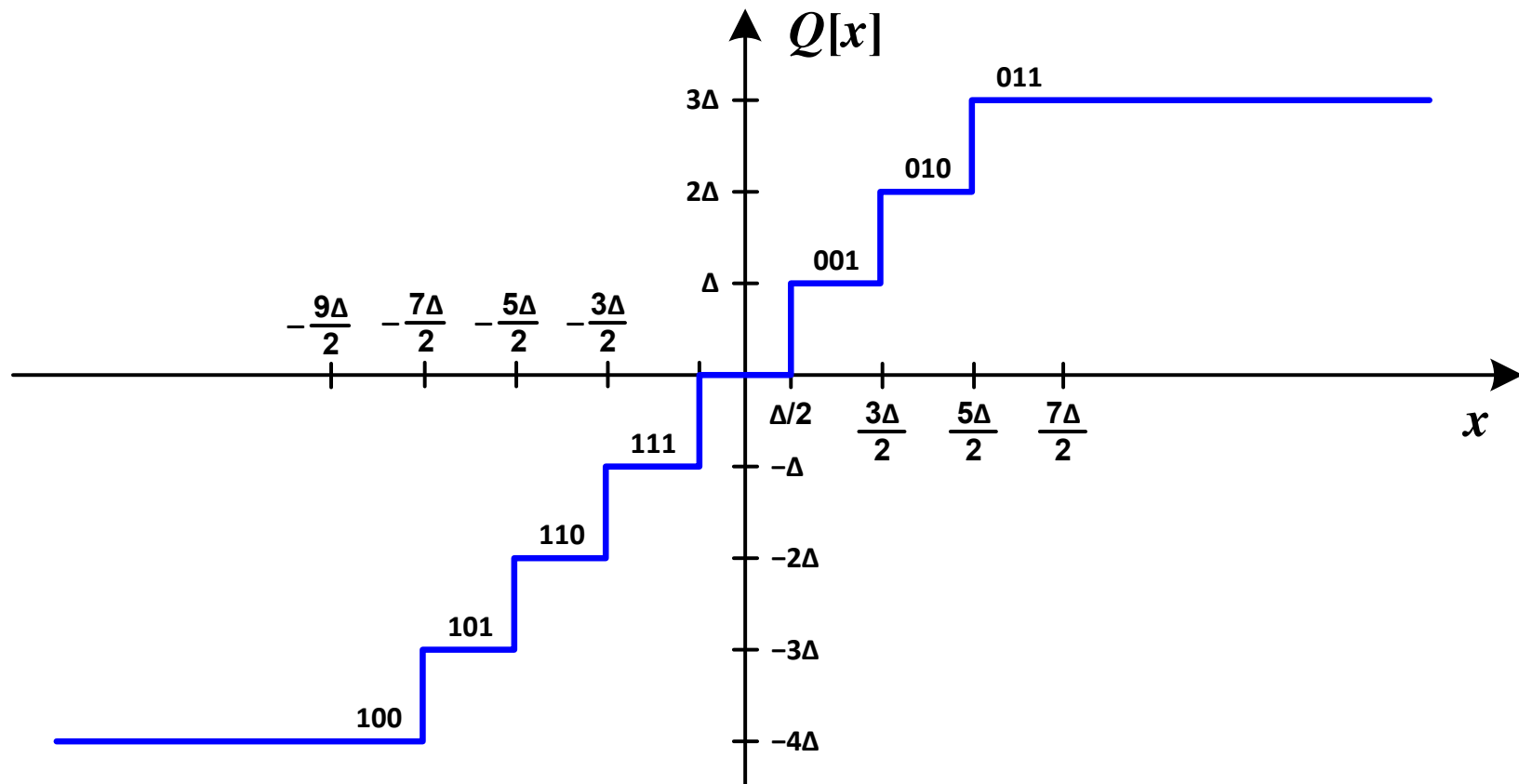


Overflow Modes: Wrap Around



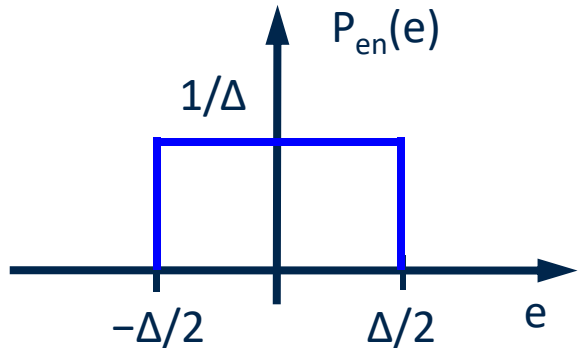
[3] A.V. Oppenheim, R.W. Schafer, with J.R. Buck, Discrete-Time Signal Processing, (2nd Ed), Prentice Hall, 1998.

Overflow Modes: Saturation



Feedback systems use saturation

Quantization Noise



$$\sigma_e^2 = \int_{-\Delta/2}^{\Delta/2} e^2 \frac{1}{\Delta} de = \frac{\Delta^2}{12}$$

x_m : full-scale signal
 $B + 1$ quantizer

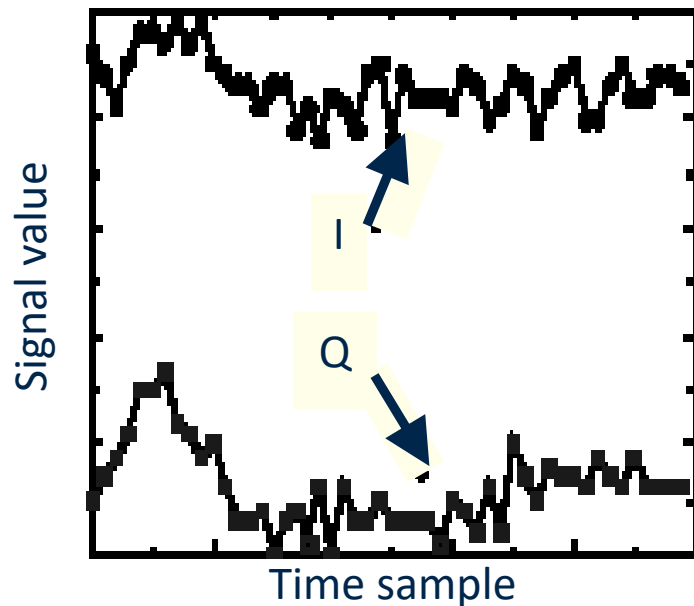
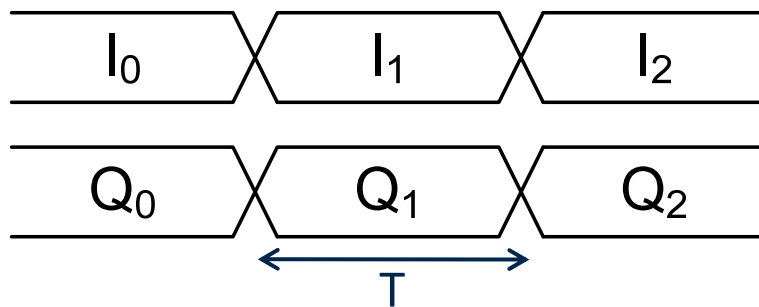
$$\sigma_e^2 = \frac{2^{-2B} x_m^2}{12}$$

$$SQNR = 10 \log_{10} \left(\frac{\sigma_x^2}{\sigma_e^2} \right) = 10 \log_{10} \left(\frac{12 \cdot 2^{2B} \cdot \sigma_x^2}{x_m^2} \right)$$

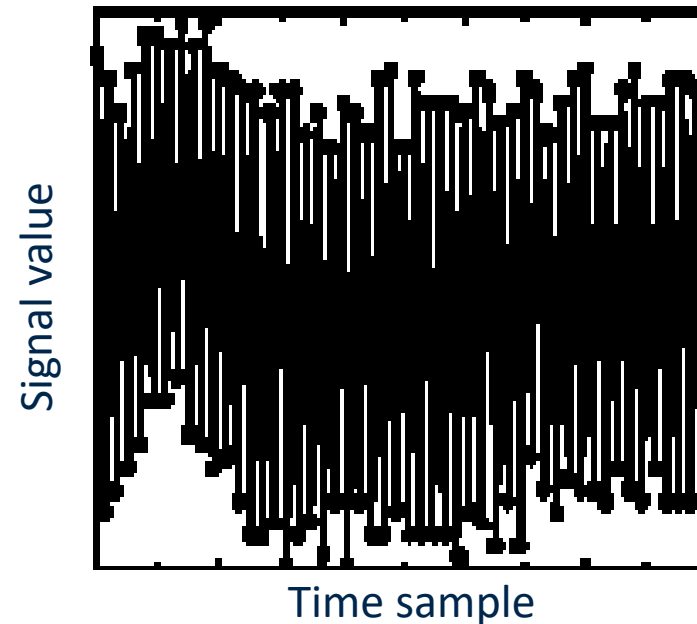
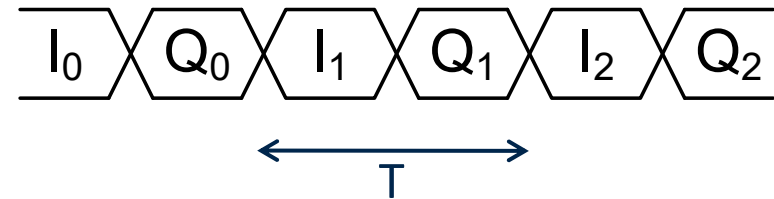
$$= 6.02 \cdot B + 10.8 - 20 \log_{10} \left(\frac{x_m}{\sigma_x} \right)$$

Time-Multiplexed Architectures

Parallel bus for I,Q



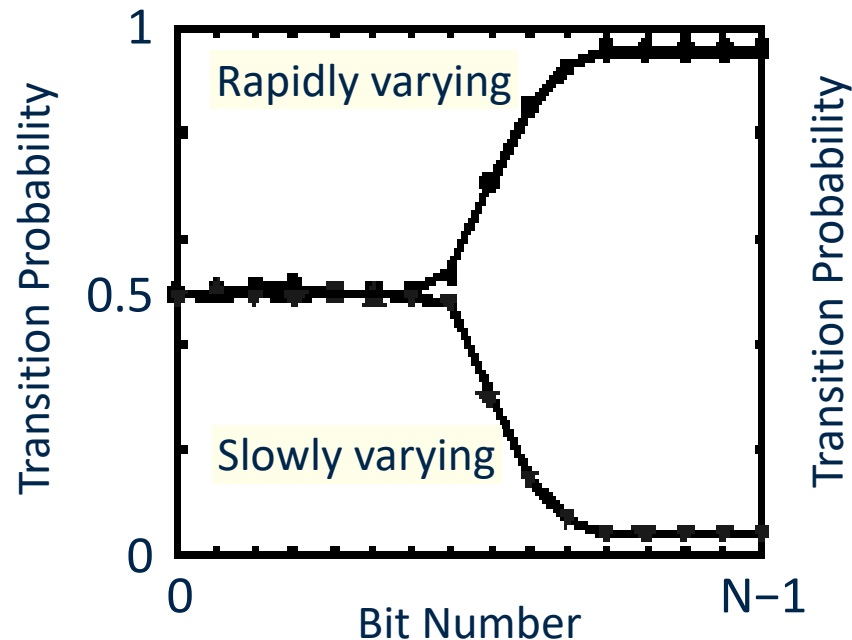
Time-shared bus for I,Q



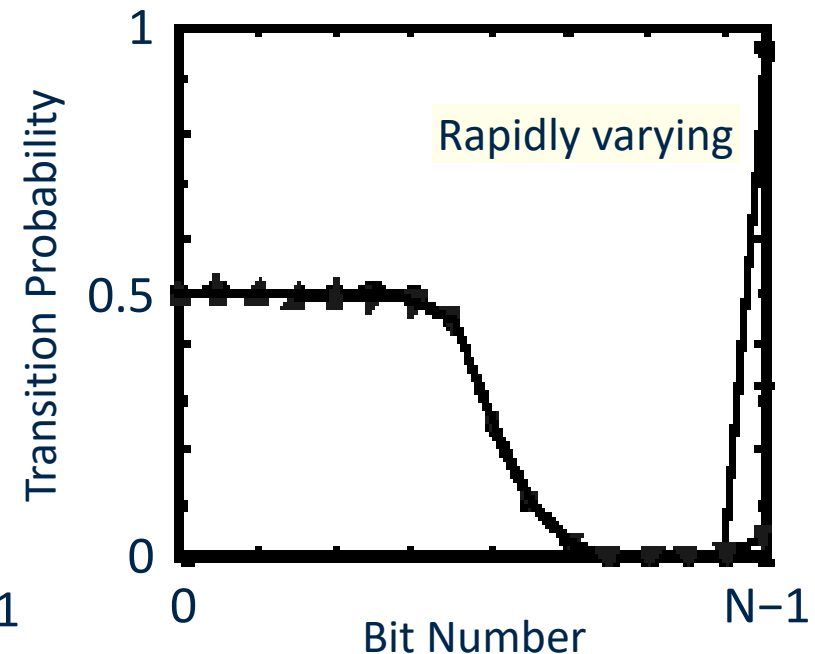
- ◆ Time-shared bus destroys signal correlations and increases switching activity

Number Representation

2's complement



Sign Magnitude

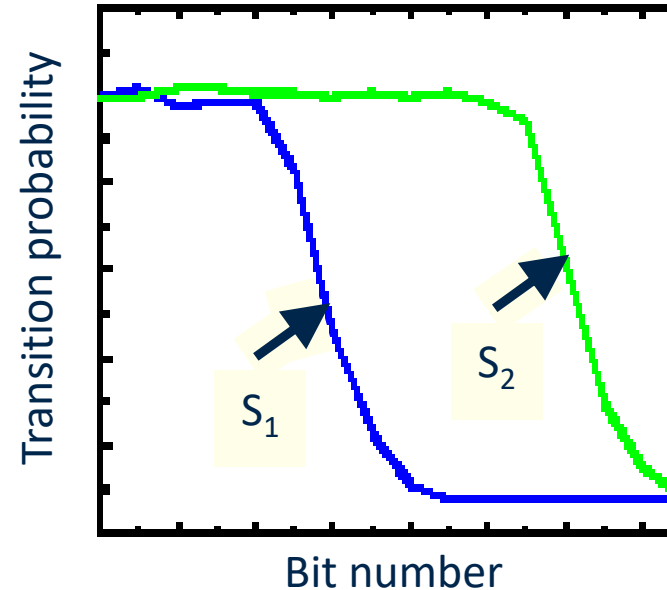
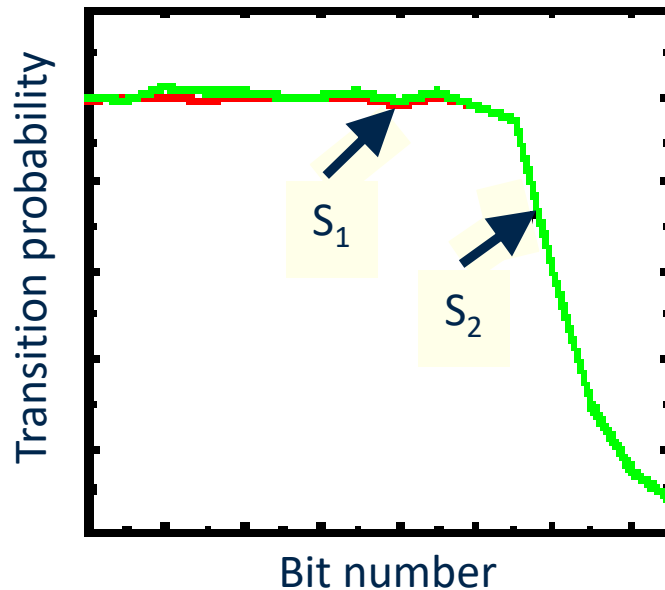
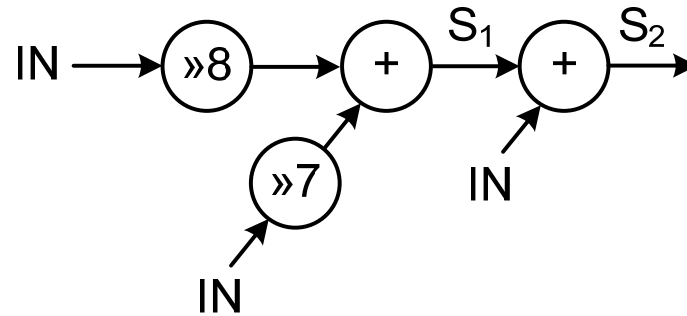
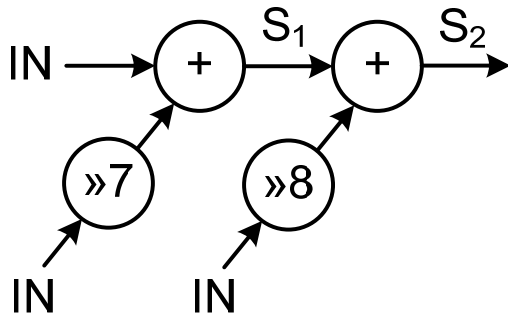


- ◆ **Sign-extension activity is significantly reduced using sign-magnitude representation**

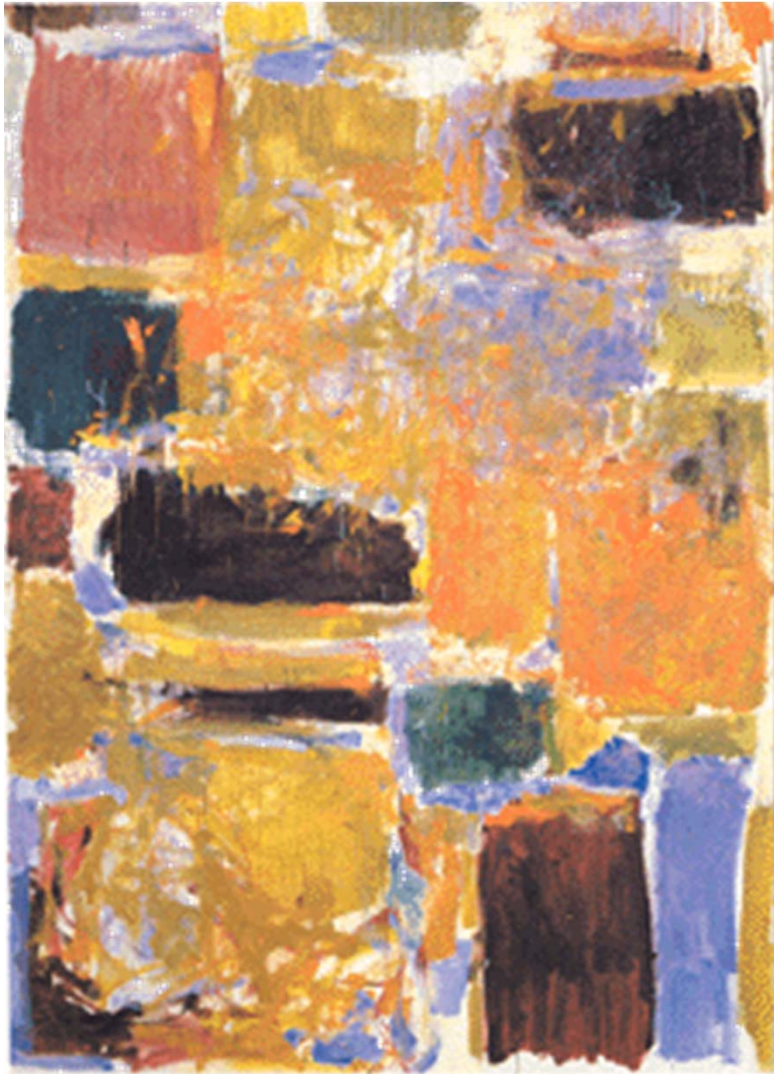
[5]

[5] A. Chandrakasan, Low Power Digital CMOS Design, Ph.D. Thesis, University of California, Berkeley, 1994.

Reducing Activity by Reordering Inputs



30% reduction in switching energy

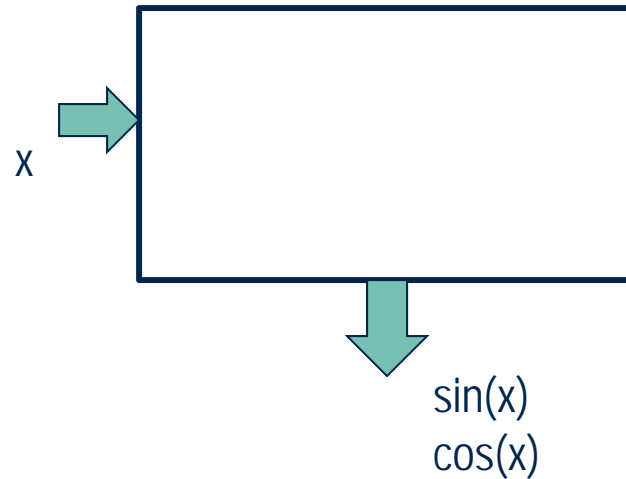


Advanced Topics in Circuit Design VLSI Signal Processing

CORDIC Algorithm and Implementation

Trigonometric Functions

- How to implement sine, cosine, etc?
- Lookup table:



- Or CORDIC

CORDIC

To perform the following transformation

$$\mathbf{y}(t) = y_R + j \cdot y_I \rightarrow |\mathbf{y}| \cdot e^{j\phi}$$

and the inverse, we use the CORDIC algorithm

CORDIC:

COordinate ROtation DIgital Computer

CORDIC: Idea

Use **rotations** to implement a variety of functions

Examples:

$$x + j \cdot y \Leftrightarrow |\sqrt{x^2 + y^2}| e^{j \cdot \tan^{-1}(y/x)}$$

$$z = \sqrt{x^2 + y^2}$$

$$z = \cos(y / x)$$

$$z = \tan(y / x)$$

$$z = x / y$$

$$z = \sin(y / x)$$

$$z = \sinh(y / x)$$

$$z = \tan^{-1}(y / x)$$

$$z = \cos^{-1}(y)$$

CORDIC: How to Do It?

- Start with **general rotation by ϕ**

$$x' = x \cdot \cos(\phi) - y \cdot \sin(\phi)$$

$$y' = y \cdot \cos(\phi) + x \cdot \sin(\phi)$$

$$x' = \cos(\phi) \cdot [x - y \cdot \tan(\phi)]$$

$$y' = \cos(\phi) \cdot [y + x \cdot \tan(\phi)]$$

- The key is to only do rotations by values of **$\tan(\phi)$** which are **powers of 2**

CORDIC: **An Iterative Process**

- To rotate to any arbitrary angle, we do **a sequence of rotations** to get to that value

Rotation Number
↓

| φ | $\tan(\varphi)$ | k | i |
|-----------|-----------------|-----|-----|
| 45° | 1 | 1 | 0 |
| 26.565° | 2 ⁻¹ | 2 | 1 |
| 14.036° | 2 ⁻² | 3 | 2 |
| 7.125° | 2 ⁻³ | 4 | 3 |
| 3.576° | 2 ⁻⁴ | 5 | 4 |
| 1.790° | 2 ⁻⁵ | 6 | 5 |
| 0.895° | 2 ⁻⁶ | 7 | 6 |

Basic CORDIC Iteration

Gain Direction Angle

$$\begin{aligned}x_{i+1} &= (K_i) \cdot [x_i - y_i \cdot d_i \cdot 2^{-i}] \\y_{i+1} &= (K_i) \cdot [y_i + x_i \cdot d_i \cdot 2^{-i}]\end{aligned}$$

$$K_i = \cos(\tan^{-1}(2^{-i})) = 1/(1 + 2^{-2i})^{0.5}$$

$$d_i = \pm 1 \text{ (rotate by } \pm\phi)$$

- If we don't multiply by K_i we get a gain error which is independent of the direction of the rotation
 - The error converges to 0.61
 - May not need to compensate for it

Basic CORDIC Iteration

Gain Direction Angle

$$\begin{aligned}x_{i+1} &= (K_i) \cdot [x_i - y_i \cdot d_i \cdot 2^{-i}] \\y_{i+1} &= (K_i) \cdot [y_i + x_i \cdot d_i \cdot 2^{-i}]\end{aligned}$$

$$K_i = \cos(\tan^{-1}(2^{-i})) = 1/(1 + 2^{-2i})^{0.5}$$

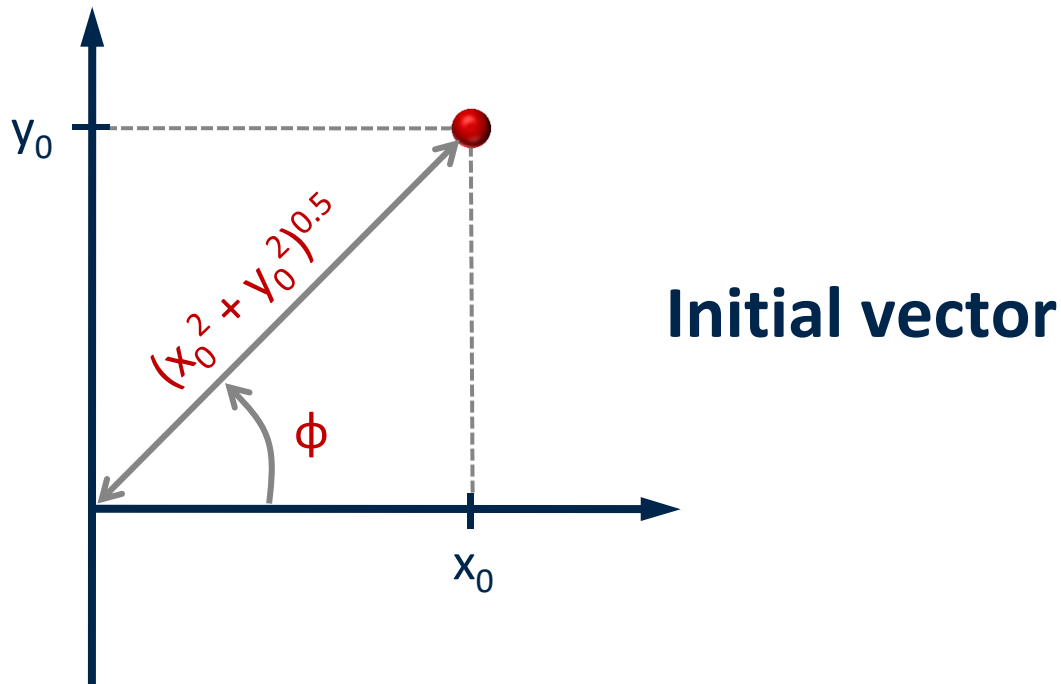
$$d_i = \pm 1 \text{ (rotate by } \pm\phi\text{)}$$

- We can also **accumulate the rotation angle**:

$$z_{i+1} = z_i - d_i \cdot \tan^{-1}(2^{-i})$$

Cartesian to Polar Conversion

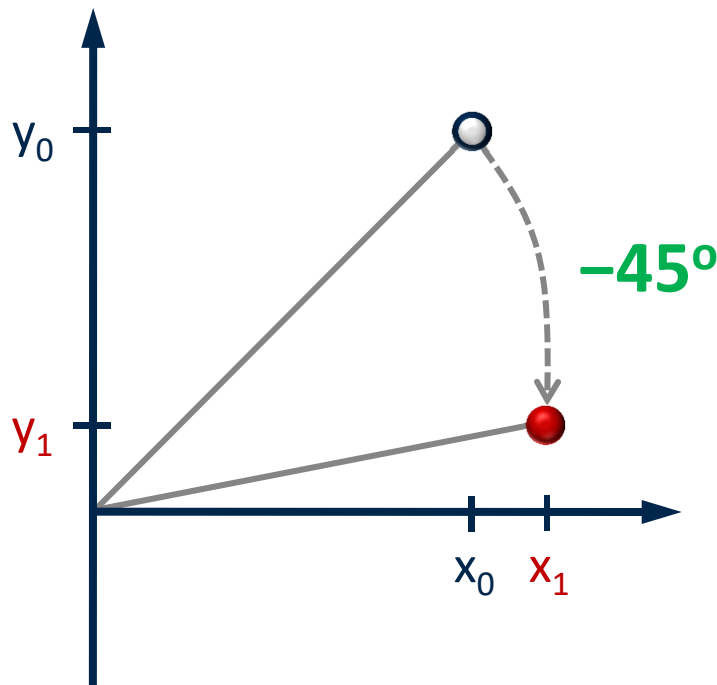
Initial vector is described by x_0 and y_0 coordinates



➡ Find ϕ and $(x_0^2 + y_0^2)^{0.5}$

Step 1: Check the Angle / Sign of y_0

- If positive, rotate by -45°
- If negative, rotate by $+45^\circ$



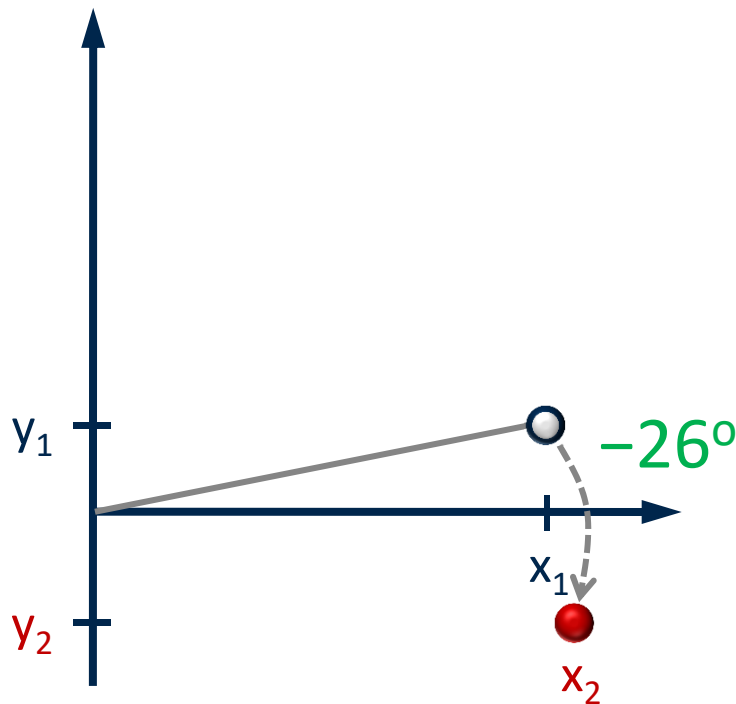
$$d_1 = -1 \quad (y_0 > 0)$$

$$x_1 = x_0 + y_0$$

$$y_1 = y_0 - x_0$$

Step 2: Check the Sign of y_1

- If positive, rotate by -26.57°
- If negative, rotate by $+26.57^\circ$



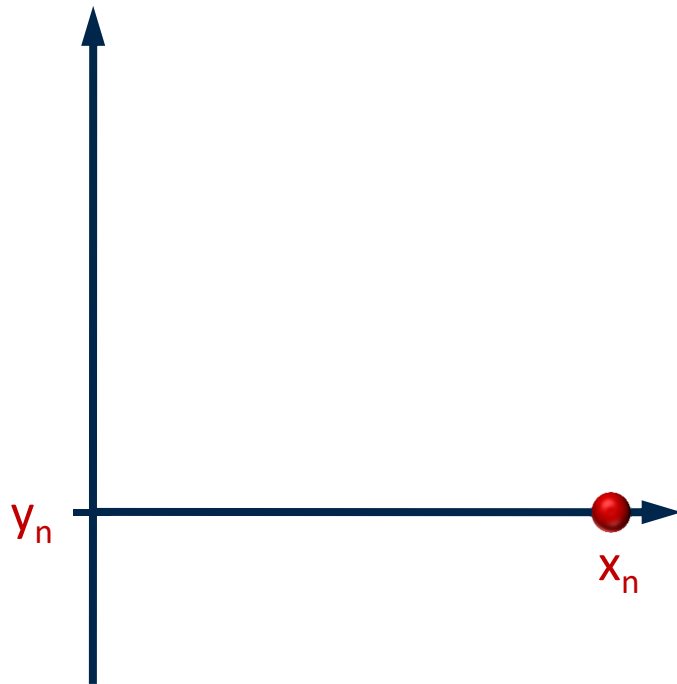
$$d_2 = -1 \quad (y_1 > 0)$$

$$x_2 = x_1 + y_1 / 2$$

$$y_2 = y_1 - x_1 / 2$$

Repeat Step 2 for Each Rotation k

- Until $y_n = 0$



$$y_n = 0$$

$$x_n = A_n \cdot (x_0^2 + y_0^2)^{0.5}$$

accumulated gain

The Gain Factor

- Gain accumulation (when you don't multiply by K_i)

$$G_0 = 1$$

$$G_0 G_1 = 1.414$$

$$G_0 G_1 G_2 = 1.581$$

$$G_0 G_1 G_2 G_3 = 1.630$$

$$G_0 G_1 G_2 G_3 G_4 = 1.642$$

- So, start with x_0, y_0 ; end up with:

$$z_4 = 71^\circ$$

$$(x_0^2 + y_0^2)^{0.5} = 1.642 \dots$$

Shift & adds of x_0, y_0



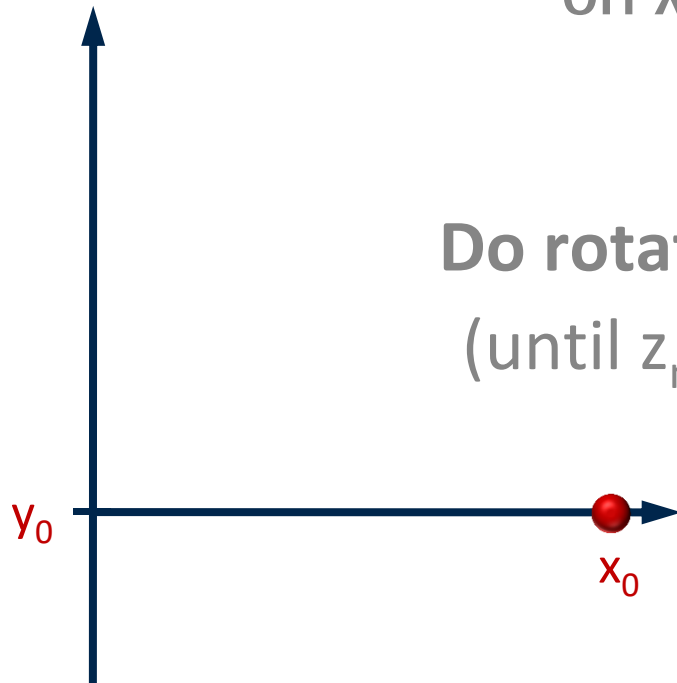
- We did the **Cartesian-to-polar** coordinate conversion

Polar-to-Rectangular Conversion

Start with vector
on x-axis



Do rotations
(until $z_n = 0$)



$$\mathbf{A} = |\mathbf{A}| \cdot e^{j\phi}$$

$$x_0 = |\mathbf{A}|$$

$$y_0 = 0, z_0 = \phi$$

$$z_i < 0, d_i = -1$$

$$z_i > 0, d_i = +1$$

$$z_{i+1} = z_i - d_i \cdot \tan^{-1}(2^{-i})$$

CORDIC Algorithm

$$\begin{aligned}x_{i+1} &= x_i - y_i \cdot d_i \cdot 2^{-i} \\y_{i+1} &= y_i + x_i \cdot d_i \cdot 2^{-i} \\z_{i+1} &= z_i - d_i \cdot \tan^{-1}(2^{-i})\end{aligned}$$

$$d_i = \begin{cases} -1, & z_i < 0 \\ +1, & z_i > 0 \end{cases}$$

Rotation mode

(rotate by specified angle)

Minimize residual angle

$$\begin{aligned}x_n &= A_n \cdot [x_0 \cdot \cos(z_0) - y_0 \cdot \sin(z_0)] \\y_n &= A_n \cdot [y_0 \cdot \cos(z_0) + x_0 \cdot \sin(z_0)] \\z_n &= 0\end{aligned}$$

$$d_i = \begin{cases} -1, & y_i > 0 \\ +1, & y_i < 0 \end{cases}$$

Vectoring mode

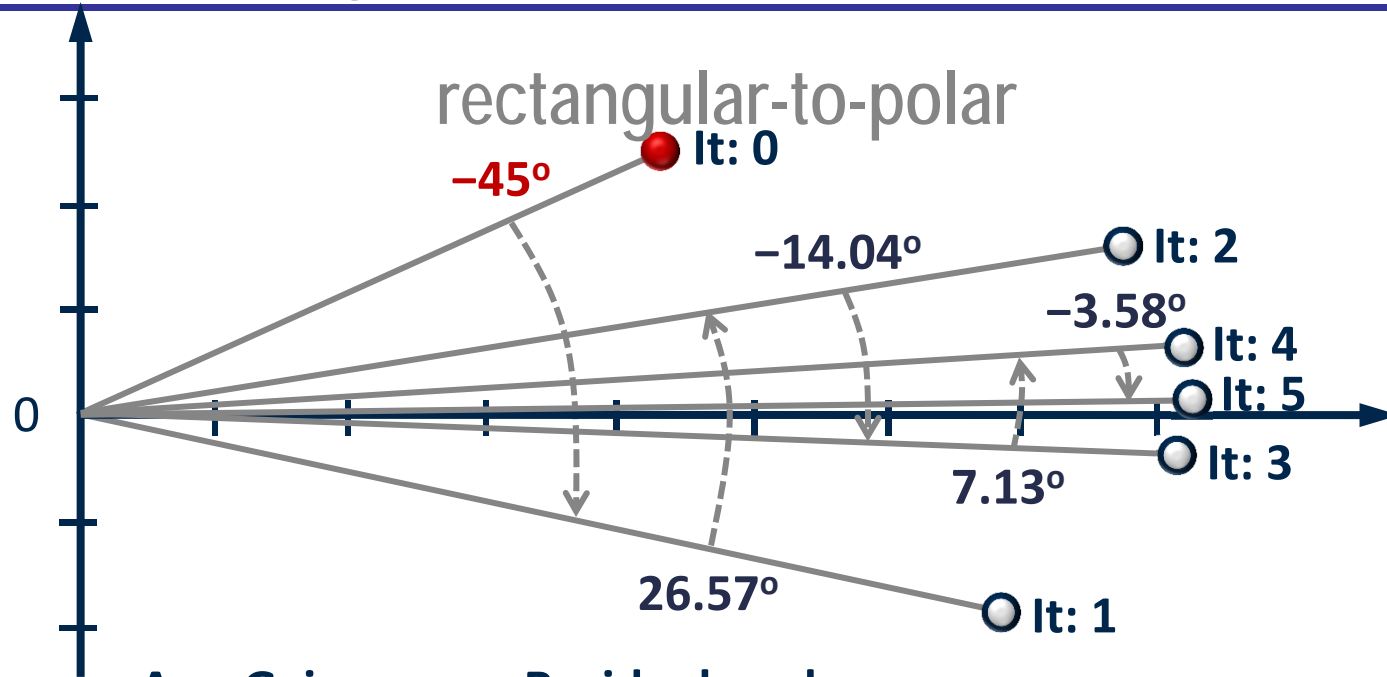
(align with the x-axis)

Minimize y component

$$\begin{aligned}x_n &= A_n \cdot (x_0^2 + y_0^2)^{0.5} \\y_n &= 0 \\z_n &= z_0 + \tan^{-1}(y_0/x_0)\end{aligned}$$

$$A_n = \prod_{i=0}^n (1 + 2^{-2i}) \cdot 0.5 \rightarrow 1.647$$

Vectoring Example



Acc. Gain

$K_0 = 1$

$K_1 = 1.414$

$K_2 = 1.581$

$K_3 = 1.630$

$K_4 = 1.642$

$K_5 = 1.646$

Etc.

Residual angle

$\phi = 30^\circ$

$\phi = -15^\circ$

$\phi = 11.57^\circ$

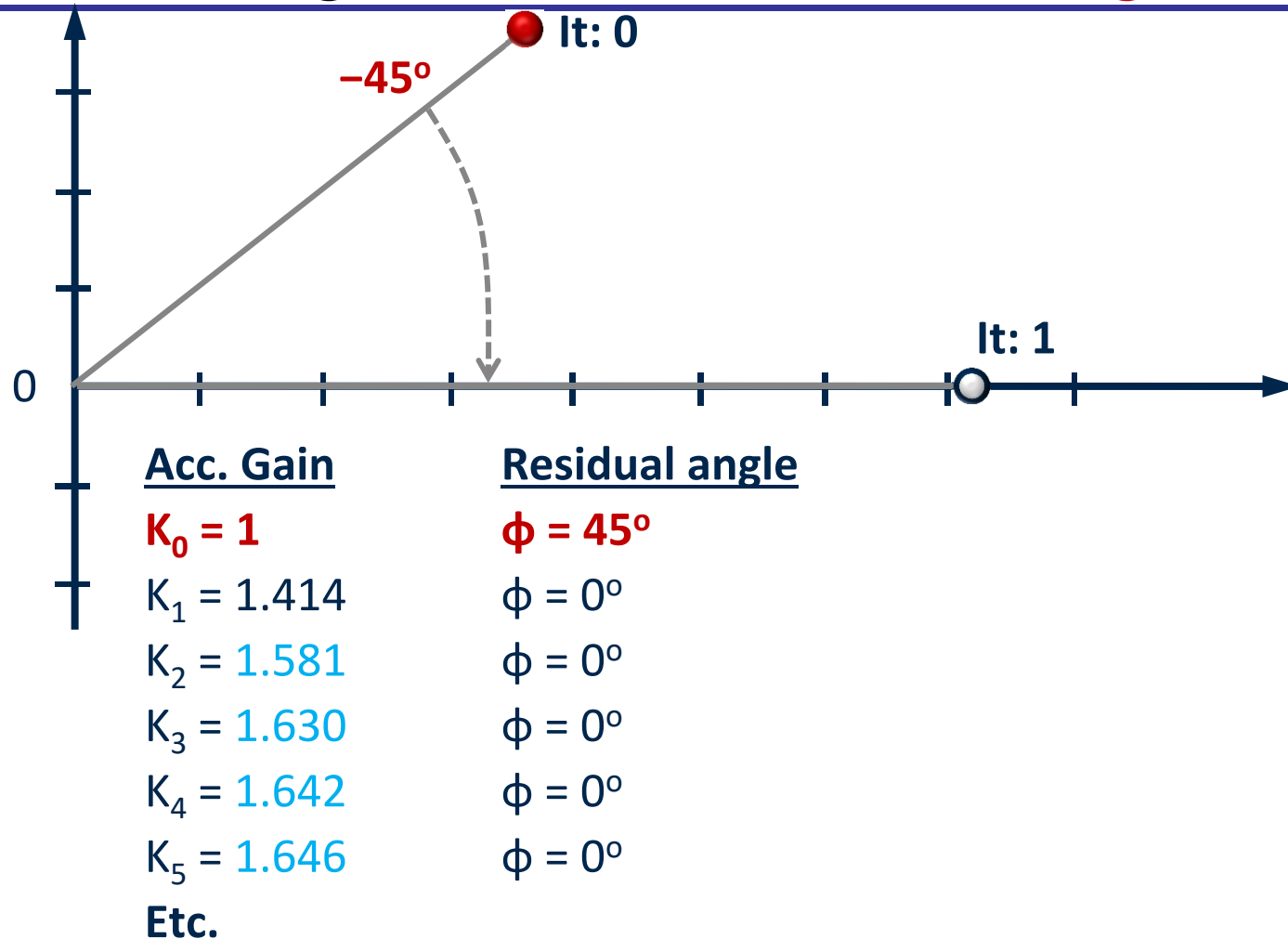
$\phi = -2.47^\circ$

$\phi = 4.65^\circ$

$\phi = 1.08^\circ$

Etc.

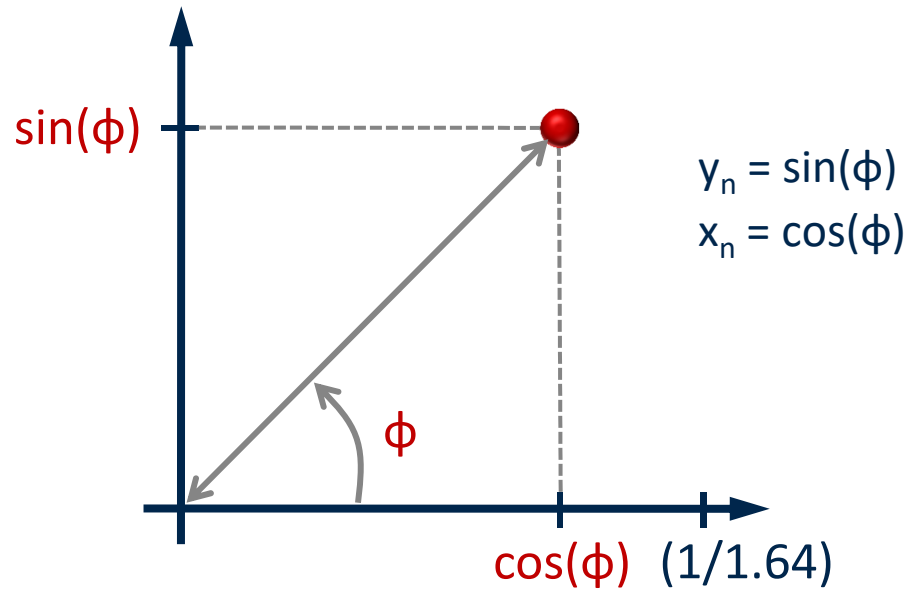
Vectoring: Best-Case Convergence



In the best case ($\phi = 45^\circ$), we can converge in 1 iteration

Calculating Sine and Cosine

- Start with $x_0 = 1/1.64$, $y_0 = 0$
- Rotate by ϕ



Functions

Rotation mode

sin/cos

$$z_0 = \text{angle}$$

$$y_0 = 0, x_0 = 1/A_n$$

$$x_n = A_n \cdot x_0 \cdot \cos(z_0)$$

$$y_n = A_n \cdot x_0 \cdot \sin(z_0)$$

(=1)

Polar → Rectangular

$$x_n = r \cdot \cos(\phi)$$

$$y_n = r \cdot \sin(\phi)$$

$$x_0 = r$$

$$z_0 = \phi$$

$$y_0 = 0$$

Vectoring mode

\tan^{-1}

$$z_0 = 0$$

$$z_n = z_0 + \tan^{-1}(y_0/x_0)$$

Vector/Magnitude

$$x_n = A_n \cdot (x_0^2 + y_0^2)^{0.5}$$

Rectangular → Polar

$$r = (x_0^2 + y_0^2)^{0.5}$$

$$\phi = \tan^{-1}(y_0/x_0)$$

CORDIC Divider

- To do a divide, change CORDIC rotations to a **linear function** calculator

$$x_{i+1} = x_i - 0 \cdot y_i \cdot d_i \cdot 2^{-i} = x_i$$

$$y_{i+1} = y_i + x_i \cdot d_i \cdot 2^{-i}$$

$$z_{i+1} = z_i - d_i \cdot (2^{-i})$$

Generalized CORDIC

$$x_{i+1} = x_i - m \cdot y_i \cdot d_i \cdot 2^{-i}$$

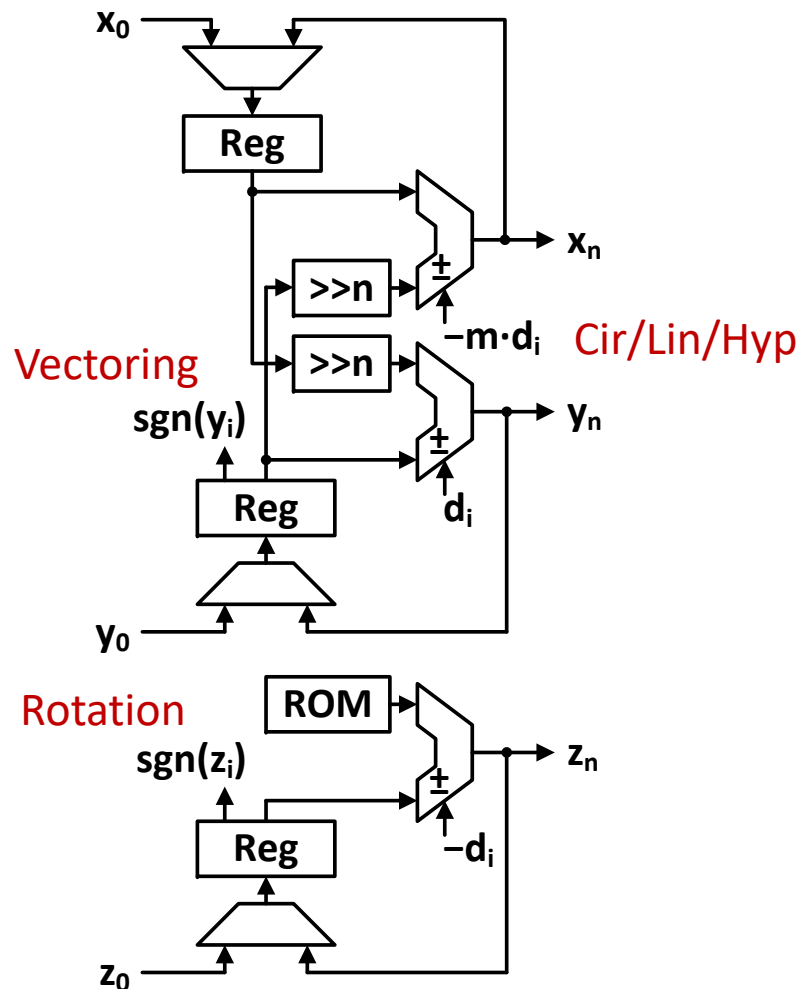
$$y_{i+1} = y_i + x_i \cdot d_i \cdot 2^{-i}$$

$$z_{i+1} = z_i - d_i \cdot e_i$$

| d_i | |
|--|--|
| Rotation | Vectoring |
| $d_i = -1, z_i < 0$ $d_i = +1, z_i > 0$ | $d_i = -1, y_i > 0$ $d_i = +1, y_i < 0$ |
| $\text{sign}(z_i)$ | $-\text{sign}(y_i)$ |

| Mode | m | e_i |
|------------|-----|----------------------|
| Circular | +1 | $\tan^{-1}(2^{-i})$ |
| Linear | 0 | 2^{-i} |
| Hyperbolic | -1 | $\tanh^{-1}(2^{-i})$ |

An FPGA Implementation



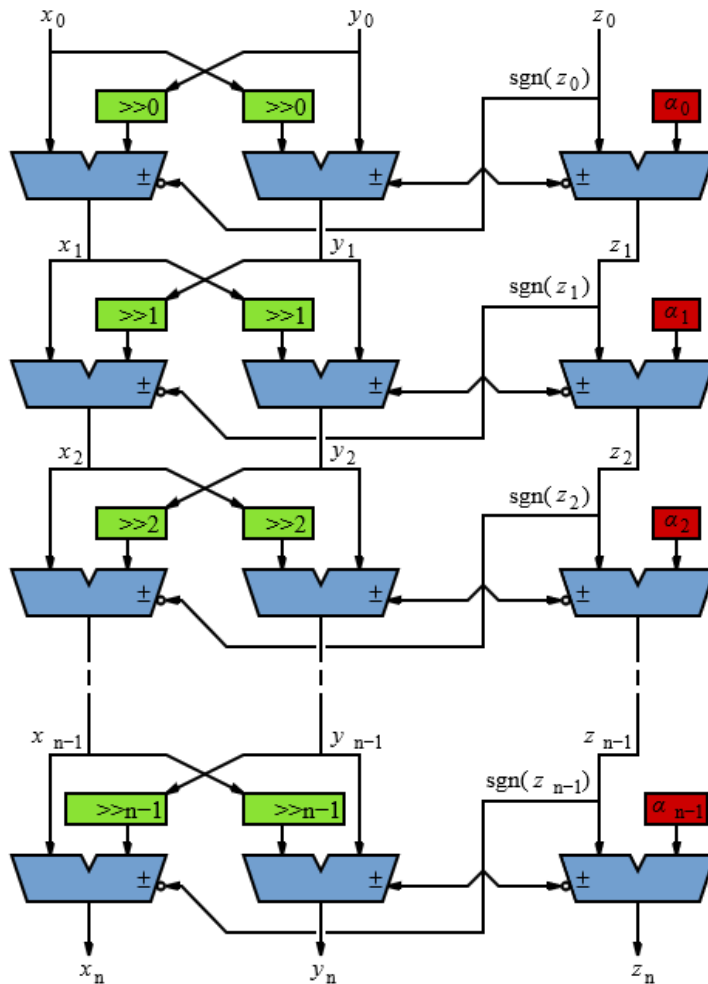
- Three difference equations directly mapped to hardware
- The decision d_i is driven by the sign of the y or z register
 - Vectoring: $d_i = -sign(y_i)$
 - Rotation: $d_i = sign(z_i)$
- The initial values loaded via muxes
- On each clock cycle
 - Register values are passed through shifters and add/sub and the values placed in registers
 - The shifters are modified on each iteration to cause the desired shift (state machine)
 - Elementary angle stored in ROM

R. Andraka, "A survey of CORDIC algorithms for FPGA based computers," in Proc. Int. Symp. FPGAs, Feb. 1998, pp. 191-200.

Last iteration: results read from reg — 36

Unrolled CORDIC

Single-cycle loop - unroll





Next Lecture

- Processor interfacing