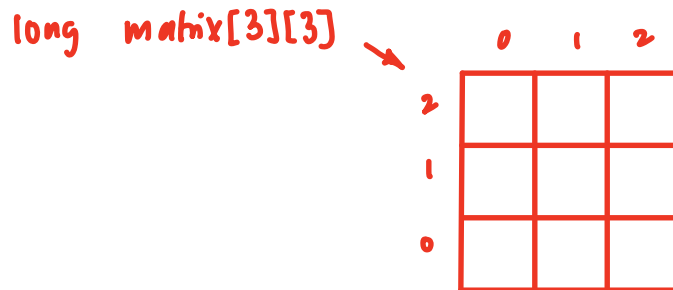


## nD Arrays (Multidimensional Arrays)

Declaration: `type name[#rows][#columns]`

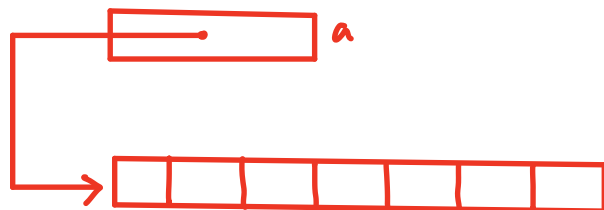


Array decay in multidimensional array:

Given an nD array `matrix[i][j]`, `matrix[i]` decays to `&matrix[i][0]`  
`&matrix[i][0]` is a pointer to long

`matrix` decays to `&matrix[0]`  
`&matrix[0]` is a pointer to an array of long values.

notation for pointer to an array:  
`long (*)[len_array]`



Passing an nD array into a function (2D-array)

`void foo(size_t rows, size_t cols, long matrix[10][20])`

`void foo(size_t rows, size_t cols, long matrix[rows][20])`

```
void foo(size_t rows, size_t cols, long matrix[][20])
void foo(size_t rows, size_t cols, long (*matrix)[20])
```

exact subarray length must be passed in to the function

`long (*a)[20]` → points to an array of 20 elements

`long *a[20]` → an array of pointers to long

`long **a[20]` → an array of pointers to long

Indexing in nD arrays:

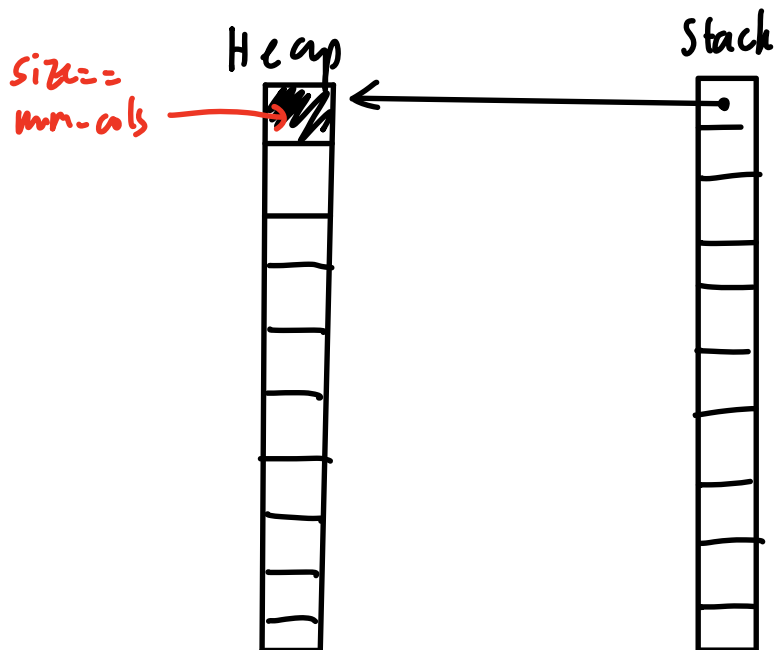
`bucket[i][j] == * (bucket[i] + j)`

### Heap allocation of memory in nD arrays

```
long *buckets[10]
for (long i = 0; i < 10; i += 1) {
    buckets[i] = malloc(k * size_of(long));
    if (buckets[i] == NULL) {
        cs1010_println_string("unable to allocate");
        for (long j = 0; j < i, j += 1) {
            free(buckets[j]); ← to free all previously allocated memory
        }
        return 1;
    }
}
```

Easier way to allocate memory

```
bucket[0] = malloc (10 * num-of-cols * sizeof(double));  
if (bucket[0] == null) {  
    // IO-error - string ("unable to allocate");  
    return 1;  
}  
for (size_t i = 1; i < 10; i += 1) {  
    bucket[i] = bucket[i-1] + num-of-cols;  
}  
free(bucket[0]);
```



If # rows & # columns are unknown

```
double ** canvas;
```

```
size_t num_rows;
```

```
size_t num_cols;
```

```
canvas = malloc (num_rows * size of (double *))
```

```
for (size_t i = 0; i < num_rows; i += 1) {
```

```
    canvas[i] = malloc (num_cols * size of (double));
```

```
    if (canvas[i] == NULL) {
```

```
        :
```

```
    }
```

```
}
```

```
for (size_t i = 0; i < num_rows; i += 1) {
```

```
    free (canvas[i]);
```

```
}
```

```
free (canvas)
```

deallocate the memory that each pointer in canvas points to, to prevent people from crashing your program

## C-Preprocessor

C Program → pre-process → C Program → compile (etc) → Machine Code → Run

clang -E → will return C program after pre-processing

`#include` } pre-processing directives  
`#define` }

`#define` pre-processing directive

Used to define constants

MACROS:     `#define fn <fn-body>`

Role of macro - purely text substitution (does not check type)

Properly defining a macro:

`#define SQUARE(x) (x)*(x)`  
                                  ↑  
                                  include parentheses

ASSERT MACRO     `#include <assert.h>`

Macros that help user find bugs in a program

We can assert that an input/parameter must meet certain condition in the function body. It will raise an error if the assertion fails.

## Program Efficiency

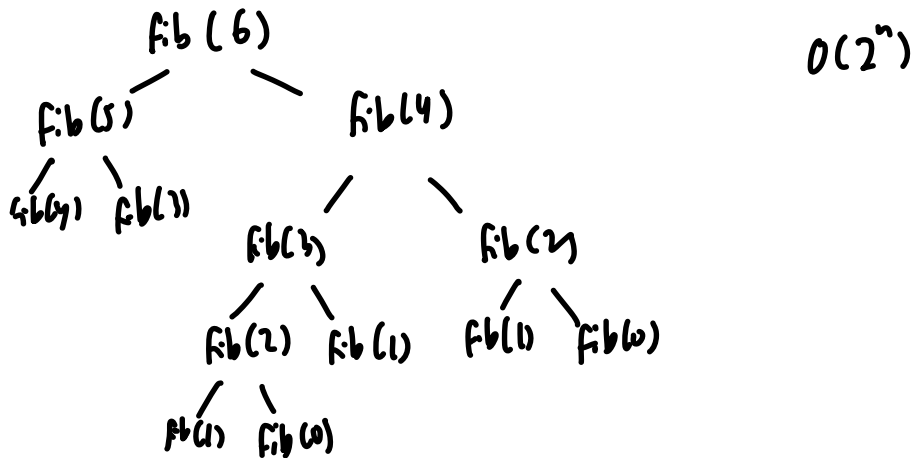
### 1) Avoid redundant work

Check program's efficiency in "WORST CASE SCENARIO". Example very large prime number in is\_prime function.

### 2) No duplication

For example, in finding range of a list, we can combine two loops (one for max one for min) into just one loop (that finds both max & min)

## Fibonacci Program (Recursive Implementation)



## Fibonacci Program (Iterative implementation)

Number of steps in terms of  $n$ :

2	+	3	+	$3(n-2)$	+	1
↑		↑		↑		↑
boolean check		initialising $a, b \& c$		looping until result		return

$$= 5 + 3n - 6 + 1$$

$$= 3n$$

we can ignore the constant

The program takes  $O(n)$  time.

## Big-O notation

To find big O, (1) find number of steps taken by a program in terms of  $n$ ,  
(2) drop the terms with lower rate of growth & (3) drop multiplicative constants

Example:  $O(\frac{n^4}{10} + n^3 + n^2 + 2n + 1) \sim O(n^4)$

### Comparing rate of growth

If  $f(n)$  grows faster than  $g(n)$ , we can find constants  $n_0$  and  $c$  such that  $f(n) > c g(n)$  for all  $n > n_0$

Example:

$$f(n) = n^n, \quad g(n) = 2^n$$

$$n_0 = 2$$

$$n_0 = 3$$

$$f(n_0) = g(n_0)$$

$$f(n_0) > g(n_0)$$

Example selection sort:  $O(n^2)$

we loop through the elements in the list  $n-2$  times to sort

We loop through the elements in the list a further  $n-2$  times to find max within the outer loop

$$(n-2) + (n-2) = n^2 - 4n + 4 \sim O(n^2)$$