

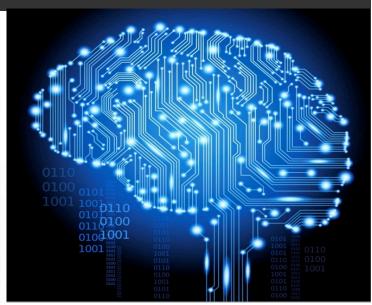
Information Technology

FIT1008/2085 MIPS – Intro to assembly programs

Prepared by: Maria Garcia de la Banda Revised by A. Aleti, D. Albrecht, G. Farr, J. Garcia and P. Abramson







Where are we at

- We now know why MIPS is useful
- We know the basics of MIPS R2000 architecture
 - CPU: ALU, GPRs, Control, and the special purpose registers
 - Main memory, its segments and their purpose
- We know how programs are executed in this architecture
 - The fetch-decode-execute cycle



Learning objectives for this lesson block

- To become familiar with MIPS assembly programs
 - Main components (labels, comments, instructions, directives, etc)
- To understand the role of MIPS assembly directives and to know how to use them in a program



In the previous lesson we said...

The machine code to compute the factorial of a number for a "MIPS" R2000 CPU consists of the following bit patterns:

Each line of 32 bits (each word) is a different machine code instruction



```
# Computes factorial of number (n) in $t0
# and returns result ("Result") in $v0
      .text
fact:
       ori $v0, $0, 1
                       \# Result = 1
       addi $s0, $0, 1
                       # s0 = 1
       slt $t1, $s0, $t0 # if n <= 1
loop:
            $t1, $s0, end # goto end
       bne
            $v0, $t0
       mult
       mflo
            $v0
                               \# Result = Result * n
       addi
            $t0, $t0, -1
                               \# n = n - 1
             loop
                                # goto loop
end:
       jr
             $ra
                                # return
```

```
# Computes factorial of number (n) in $t0
   # and returns result ("Result") in $v0
            .text
   fact:
            ori
                    $v0, $0, 1
                                              # Result = 1
            addi
                     $s0, $0, 1
                                              # s0 = 1
                                              # if n <= 1
   loop:
            slt
                     $t1, $s0, $t0
                The instructions themselves
• One per line
                                              # Result = Result * n
• Human-readable instruction code, e.g. addi
• This one says "add the immediate value 1 to the contents of
register zero, and store the result in register so"
• Translated by the assembler into machine language bit patterns
```

```
Computes factorial of number (n) in $t0
 # and returns result ("Result") in $v0
         .text
 fact:
         ori
                                            let Result =
                                         # let s0 = 1
                                         # if n <= 1
Comments: start with a hash sign
                                         # goto end
(#) and go to end of line. $50, er
They are ignored by the assembler
                                         # Result = Result * n
                 $t0, $t0, -1
          addi
                                         \# n = n - 1
                 loop
                                         # goto loop
 end:
          jr
                 $ra
                                           return
```

Let's look at some features of assembly language that help you read and write it:

```
# Computes factorial of number (n) in $t0
# and returns result ("Result") in $v0
       .text
```

Defines a label

```
fact:
```

loop:

```
Labels identify lines of code
so that you can refer to
them by name. $50, $t0
```

They are translated by the assembler into addresses

\$t0, \$t0, -1 addi

end:

jr

loop

\$ra

```
# Result = 1
```

s0 = 1

if n <= 1

goto end

Result = Result * n

n = n - 1

goto loop

return

Uses a label



```
# Computes factorial of number (n) in $t0
# and returns result ("Result") in $v0
        .text
fact:
                       $0, 1
                                          # Result = 1
         ori
         addi
                                          \# c0 = 1
loop:
         slt
                   GPRs: as we saw last lecture,
         bne
              begin with $
         mul<sup>-</sup>
              please use name (e.g. $v0)
         mf1
                                            Result = Result * n
              can (but don't) use number (e.g. $2)
         add
              except for $0
                                                  loop
                 $ra
         jr
                                          # return
end:
```



```
# Computes factorial of number (n) in $t0
# and returns result ("Result") in $v0
         .text
fact:
                 $v0, $0, 1
                                           # Result = 1
         ori
         addi
              Lines beginning with a dot are assembler
         slt
loop:
              directives.
         bne
              They do not become instructions.
         mul
              Instead, they just tell the assembler to
         mf1
                                                    lt = Result * n
              do something.
         add
              This one tells it to put the code in the
                                                    lloop
              text segment.
                  $ra
         jr
                                           # return
end:
```

```
# Computes factorial of number (n) in $t0
# and returns result ("Result") in $v0
```

```
Numbers by themselves are immediate
                                                     \# Result = 1
values; this one is -1.
                                                     # s0 = 1
This instruction decrements $t0 by
                                                    # if n <= 1
                                       $t0
adding -1 to it.
                                       end
                                                     # goto end
Immediate values may appear in
decimal, hexadecimal, or octal.
                                                     # Result = Result * n
                           $t0, $t0, -1
                   addi
                                                     \# n = n - 1
                           loop
                                                     # goto loop
                   jr
                           $ra
                                                     # return
         end:
```

Assembler Directives

- Always start with . (dot)
- Assembler directives <u>don't</u> assemble to machine language instructions
 - Instead, they are interpreted by the assembler
 - Result in the assembler doing something
- Do what? Depending on the directive:
 - To allocate space/data
 - To switch modes, tell in which memory segment is working



Assembler Directives – Switch Mode

.data

- Tells the assembler it is working in the part of the program that will create things (variables) in the (static part of the) data segment
 - From then on, it will find assembler directives that allocate space

.text

- Tells the assembler it is working in the part of the program that will become machine code instructions and reside in the text segment
 - From then on, it will find assembler instructions

```
# Computes factorial of number (n) in $t0
# and returns result ("Result") in $v0
        .text
fact:
                 $v0, $0, 1
                                         # Result = 1
         ori
         addi
                 $ca $a
                                         \#s0 = 1
                                           lif n <= 1
loop:
         slt
             Lines beginning with a dot are
                                         # goto end
         bnel
             assembler directives.
         mu11
             They tell the assembler to do
                                           |Result = Result * n
         mf1¢
             something.
                                           ln = n - 1
         add:
             This one tells it to put the code in
                                           goto loop
             the text segment.
         jr
                 $ra
                                          # return
end:
```

Assembler Directives – Allocate Space

- Allocates memory in the (static part of the) data segment
 - Thus, they appear under the .data directive
- space NNot very used in our unit
 - Allocates N bytes, stores nothing
- .word w1 [, w2, w3, ...]
 - Allocates 4-byte word(s) and stores the wi value(s) in it(them)
- .asciiz "string"
 - Allocates the string as a sequence of ASCII values (1 byte each), terminated by a zero byte (null character indicating end of string)
 - So, in total: 1 byte per character in the string + 1 zero byte



Assembler Directives – Example

```
# start of data segment
        .data
                             # word at var address=3
        .word 3,11
var:
                             # at var+4 address = 11
prompt: .asciiz "Hi!"
                             # starting at address
                             # prompt, store 4 bytes
                             # for chars H I ! and Zero byte
                             # start of text segment
        .text
                             # first instruction at
main:
     lw $t0, var
```

address main

Summary

- Became familiar with MIPS assembly programs
 - Main components (labels, comments, instructions, directives, etc)
- Understand the role of MIPS assembly directives:
 - Switch mode
 - Allocate space
- Know how to use these in a program:
 - data
 - text
 - asciiz
 - word
 - space N



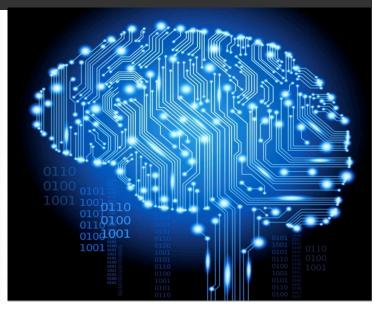


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FIT1008/2085 MIPS – Assembly Instruction Set -- I

Prepared by:
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Where are we at

- Know the basics of MIPS R2000 architecture
 - CPU: ALU, GPRs, Control, and the special purpose registers
 - Main memory, its segments and their purpose
- Understand the fetch-decode-execute cycle
- Know the main components of a MIPS program
 - Labels, comments, instructions, directives, etc
- Understand the role of assembly directives (switch mode, allocate space)
- Know how to use .data .text .asciiz .word .space N in a program



Learning objectives for this lesson block

- To become familiar with MIPS instruction set
- To learn about the MIPS instructions we will use
- To be able to use some of them in a program



MIPS Instructions (basic kinds)

- Arithmetic: add, addi, sub, mult, div
- Data movement: mfhi, mflo
- Logical: and, or, xor, nor, andi, ori, xori
- Shift: sll, sllv, sra, srav, srl, srlv
- Load/store: lw, sw
- Comparison: slt, slti
- Control transfer: beq, bne (conditional), j, jr, jal (unconditional)
- Sytem calls: syscall
- We will not see all in this lesson



Arithmetic Instructions (integer)

- <u>add</u>ition (+)
 - add \$t0, \$t1, \$t2 # \$t0 = \$t1 + \$t2
- addi immediate addition (+)
 - addi \$t0, \$t1, 5 # \$t0 = \$t1 + 5
- subtraction (-)
 - sub \$t0, \$t1, \$t2 # \$t0 = \$t1 \$t2
- multiplication (*)
 - mult \$t1, \$t2 # LO=\$t1*\$t2, HI=overflow
- **■** <u>div</u>ision (//)
 - div \$t1, \$t2 # LO=\$t1//\$t2, HI=remainder

Data Movement Instructions

- move from HI
 - mfhi \$t0 # \$t0 = HI
- move from LO
 - $\ \text{mflo} \ \$t0 \ \# \ \$t0 = LO$

Input/Output

- Programs often need to communicate with users (I/O)
- The operating system manages all peripherals including the console
- MIPS programs do I/O by asking the OS using a special command called syscall
- To make a system call in MIPS you must:
 - 1. Work out which service you want
 - 2. Put service's call code in register \$v0
 - 3. Put argument (if any) in registers \$a0, \$a1
 - 4. Perform the syscall instruction
 - 5. Result (if any) will be returned in register \$v0



System Services (cont'd)

Service	Call code	Argument	Result
Print integer	1	\$a0 (int to be printed)	n/a
Print string	4	\$a0 (addr of first char of string)	n/a
Read integer	5	n/a	\$v0 (integer)
Read string	8	\$a0 (addr to put string) \$a1 (number of bytes to read)	n/a
Allocate memory	9	\$a0 (number of bytes requested)	\$v0 (addr of allocated memory)
Exit program	10	n/a	n/a

No need to memorise! It is in the MIPS sheet (Moodle Week 1)



MIPS – I/O Example

Instructions in the program
.text

main:

```
#read integer
addi $v0, $0, 5
syscall
```

```
# multiply it by 250
addi $t1, $0, 254  # put mms pe
mult $v0, $t1  # multiply,
mflo $t0  # put result
addi $t2, $0, 10  # put 10 in
div $t0, $t2  # divide $t0
mflo $t3  # put quotie
```

```
#print result & exit
add $a0, $0, $t3
addi $v0, $0, 1
syscall
addi $v0, $0, 10
syscall
```

```
Code
     Service
                               Arg
                                          Res
Print integer
                      1
                            $a0
                                        n/a
                            $a0
Print string
                                        n/a
Read integer
                            n/a
                                        $v0
                      5
                            $a0 $a1
                                        n/a
Read string
                                        $v0
Allocate memory
                            $a0
Exit program
                     10
                            n/a
                                        n/a
```

```
# put a 5 in $v0
# result is in $v0
```

```
and // by 10
# put mms per inch in $t1 ($t1=254)
# multiply, result left in LO
# put result in $t0 ($t0=n*254)
# put 10 in $t2
# divide $t0 by $t2 (int division: n*254//10)
# put quotient in $t3 ($t3= n*254//10)
```

```
# print quotient
# system call 1 (print integer)
# print
# system call 10 (exit)
# exit
```

Load/Store Instructions

- Load (read) word from memory to GPR
 - lw \$t0, address # \$t0 = content(address)
 - Loads the 4 bytes beginning at address into \$t0
- store (write) word from GPR to memory
 - sw \$t0, address # content(address) = \$t0
 - Stores the content of \$t0 into the 4 bytes beginning at address
- Question is, how do we specify an address?
 - The opcode (lw or sw) takes up 6 bits of the IR
 - The destination register (\$t0) takes up another 5
 - This leaves us with 21 bits to indicate the address
 - Addresses are 32 bits...



Five ways to specify an address

Directly (or using a label), e.g.

```
lw $t1, N # loads from label N
```

Label plus offset, e.g.

```
lw $t1, N+4  # loads from (label N + 4)
```

Using a GPR to store the address, e.g.

```
lw $t1, ($s0)  # loads from address stored in $s0
```

• GPR + offset, e.g.

```
lw $t1, 4($s0) # loads from (address stored in $s0)+4
```

Label, offset, and GPR, e.g.

```
lw $t1, N+4($s0)# loads from (label N+4)+contents of $s0
```



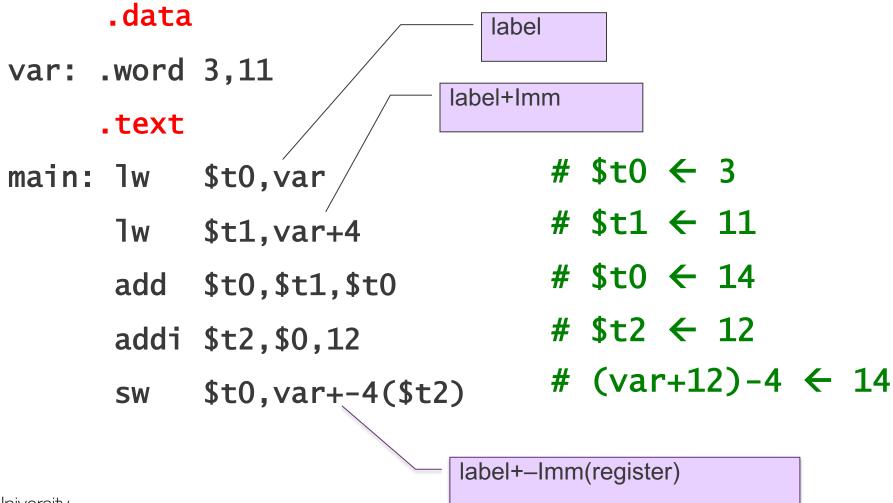
They are called Addressing Formats

Summary of addressing formats in table format:

Format	Calculation	
lmm	immediate, i.e., actual value	
label	address of label	
label+[-]Imm	address of label+[-] Imm	
(register)	contents of register	
[–]Imm(register)	contents of register+[-] Imm	
label+[-]lmm(register)	address of label+ contents of register +[-] Imm	



Load/Store Instructions – Example



Summary

- Are now familiar with MIPS instruction set
- Have learned about the MIPS instruction set we will use
- Are able to use some of them in a program:
 - Arithmetic
 - Data movement
 - Input/output
 - Load/store
- We know how to use the different addressing formats

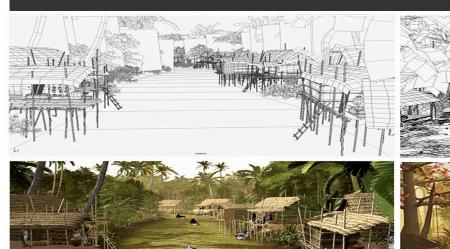




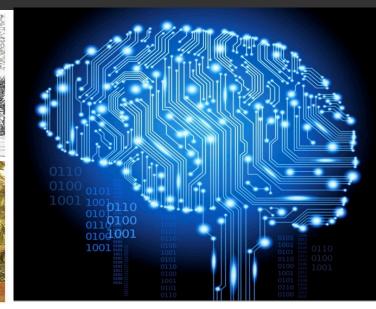
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FIT1008/2085 MIPS – Assembly Instruction Set II

Prepared by: Maria Garcia de la Banda Revised by A. Aleti, D. Albrecht, G. Farr, J. Garcia and P. Abramson







Where are we at

- Know the basics of MIPS R2000 architecture
- Understand the fetch-decode-execute cycle
- Know the main components of a MIPS program
- Can use assembly directives in a program
- Know the MIPS instruction set and can use some of them in a program:
 - Arithmetic
 - Data movement
 - Input/output
 - Load/store
- Know how to use the different addressing formats



Learning objectives for this lesson block

Continue to learn about other MIPS instructions (bitwise and shift)



Bitwise Logical Instructions

Bitwise AND (&)

```
- and $t0, $t1, $t2  # $t0 = $t1 & $t2 - andi $t0, $t1, 0xa0b1 # $t0 = $t1 & 0xa0b1

• Bitwise OR (|)
```

■ Bitwise exclusive OR (XOR) (^)

```
- xor $t0, $t1, $t2  # $t0 = $t1 \land $t2
- xori $t0, $t1, $t2  # $t0 = $t1 \land 0x0005
```

Bitwise not-OR (NOR)

```
- nor $t0, $t1, $t2  # $t0 = \sim($t1 | $t2)
```

A	В	A AND B	A OR B	A XOR B
0	0	0	0	0
0	1	0	1	1
1	0	0	1	1
1	1	1	1	0

Bitwise Logic Example

.text

Using bits to represent other things

Let's think of sequences of bits as sets:

	A	В	C	D	Е	F		
Set 1	1	0	1	1	0	0	$\{A,C,D\}$	
Set 2	1	1	0	1	0	1	$\{A,B,D,F\}$	
AND	1	0	0	1	0	0	$\{A,D\}$	Intersection
OR	1	1	1	1	0	1	$\{A,B,C,D,F\}$	Union
XOR	0	1	1	0	0	1	$\{B,C,F\}$	Difference

Very fast set operations!



Sneaky assembler tricks

- Suppose the address of label N is 0x7FFFFFF8
- That does not fit in 32 bits with the opcode and register
 - How does the assembler manage to translate loads/stores from addresses larger than 0xFFFF? (i.e., those that do not fit in 16 bits)
- Using pseudoinstructions: it translates la \$t0, N into two lines of machine code! (16 + 16 = 32 bits)
 - One sets the top 16 bits: lui \$t0, 0x7FFF
 - "load upper immediate": loads 16-bit value into top 16 bits of register
 - The other sets the bottom 16 bits: ori \$t0, \$t0, 0xFFF8

You will not be asked to use **lui**, so don't worry about the details.

More pseudoinstructions

- li *\$r1*, *n* # load immediate using 32 bits
 - puts immediate value into \$r1
 - translates that line into a lui and an ori (or addi)
- la *\$r1*, *label* # load address (again, 32 bits)
 - puts address of Label into \$r1
 - assembler knows this address, so translates it like li

There are many pseudoinstructions in MIPS but you cannot use them for FIT1008/2085. The point is for you to learn to use the basic blocks.

The only one you can use is **1a**, which is very useful for loading the address of any string we want to print.

Have a look at the MIPS reference sheet to see the instructions you are allowed to use (the one with the blue title!)

Shift Instructions

Only the lower 5 bits of the \$t2 register are used

Same as multiplying by 2⁵

shift left (logical) (<<)</p>

```
    fill with zero bits
```

- sll \$t0, \$t1, 5 # \$t0=\$t1 << 5
- sllv \$t0, \$t1, \$t2 # \$t0=\$t1 << content(\$t2)</pre>

Immediate value ≤ 31

shift right (logical) (>>> in JS, not provided by Python)

```
    fill with zero bits
```

- srl \$t0, \$t1, 5 # \$t0=\$t1 >> 5

- srlv \$t0, \$t1, \$t2 # \$t0=\$t1 >> content(\$t2)

shift right (arithmetic) (>>)

Same as dividing by 2⁵

- fill with copies of MSB (most significant bit – indicates the sign) - sra \$t0, \$t1, 5 # \$t0=\$t1 >> 5 Sa - srav \$t0, \$t1, \$t2 # \$t0=\$t1 >> content(\$t2)

```
s]] $t2, $t0, 5 #1111111111111111111111110100000
srl $t2, $t0, 5 #0000010111
sra $t2, $t0, 5
```

Bitwise Shift and Logic Example

.text

```
0001 = 1
main: ori $t0, $0, 0xa0b1
                                 # $t0 = 0000a0b1
                                                    0101 = 5
                                 # $t1 = 00000005
      ori $t1, $0, 5
                                 # $t2 = 0000a0b5
      or $t2, $t1, $t0
                                 \# $t3 = 00000001
      and $t3, $t1, $t0
                                 # $t4 = 00000050
      sll $t4, $t1, 4
                                                    1111 = f
                                 # $t5 = fffffffa
      addi $t5, $0, -6
                                                    0111 = 7
                                 # $t6 = 07ffffff
      srlv $t6, $t5, $t1
                                 # $t7 = fffffffff
      sra $t7, $t5, 4
```

Summary

- Are able to use mode MIPS instructions in a program:
 - Bitwise: and or xor nor
 - Shift: sll slv srl srlv sra srv
- We have learned about pseudoinstructions and know we will only use one:
 - la for printing strings

The following 3 slides are not examinable but might be interesting for some



Labels and symbols

- We want to use names (labels) for memory locations (addresses)
- They might appear under the .data and under .text directives
- They need to be translated into addresses before execution
- To do this, the assembler performs two passes on every program:
- The first builds a symbol table, i.e., when it sees a label being defined:
 - It puts the label name and the current address in the table
- The second pass uses the table, i.e. when it sees a label being used:
 - It looks the name up in the table to find what address it refers to



MIPS Program – Directives Example

```
# Sets current addr to 0x10000000
                 .data
                                      # Updates N in table and allocates 4+4=8 bytes
N:
                 .word 100, 72
                                      # Set contents to 100 and 72
                                      # Updates aString and allocates 7 bytes
                 .asciiz "Hello!"
aString:
                                      # Set contents to "Hello!" + null
                                      # Sets current addr to 0x00400000
                 .text
                              # Look up N in table, translate instruction
                 lw $t0, N
                 addi $t0, $t0, 10 # Update loop in table, translate instruction
loop:
                                       # Look up loop in table, translate instruction
                j loop
```

```
Symbol table (at the end of the first pass)
N: 0x10000000
aString: 0x10000008
loop: 0x00400004
```

Dark green happens in the first pass of the assembler, light green in the second pass



How a (simplified) assembler works

- First pass: check file for assembler directives and labels
 - Handle those if found by building a symbol table (see example in next slide)
- Second: go back to start of file; for each line of assembly language do:
 - Look up its operation code:
 - If valid, set first six bits of instruction to opcode, else output error
 - For each register in the line,
 - Look its number up in table and set the appropriate five bits in the instruction
 - If there is a reference to a label:
 - Look its value up in the symbol table and treat it like an immediate
 - If there is an immediate value on the line
 - Copy it into the last sixteen bits of the instruction

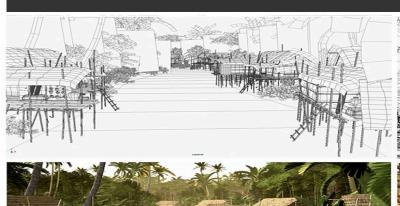




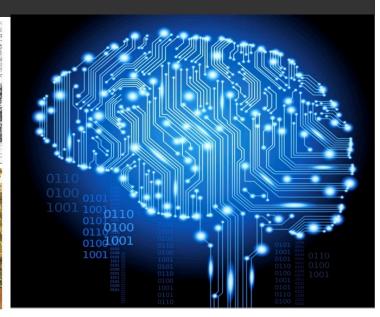
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FIT1008/2085 MIPS – Instruction Set III and Formats

Prepared by: Maria Garcia de la Banda Revised by A. Aleti, D. Albrecht, G. Farr, J. Garcia and P. Abramson







Where are we at

- Know the basics of MIPS R2000 architecture
- Understand the fetch-decode-execute cycle
- Know the main components of a MIPS program
- Can use assembly directives in a program
- Know the MIPS instruction set and can use some in a program:
 - Arithmetic
 - Data movement
 - Input/output
 - Load/store
 - Bitwise
 - Shift
- Know how to use addressing formats and the pseudoinstruction la



Learning objectives for this lesson block

- To learn how to use unconditional control transfer instructions
- To learn the three different instruction formats



Blast from the past: the goto statement

- Remember: a label is an identifier for a memory position (data/code)
- The goto statement performs an unconditional jump to its label argument
- It promotes code whose control flow is extremely difficult to understand
- That is why it is not supported by most languages, including Python



If Python had a goto statement ...

```
# Code could be this ugly!
def main():
                                 A very unclear way to print 1 2 3 4!!
      print(1)
      goto apple
  orange:
      print(3)
      goto pomegranate
  apple:
      print(2)
      goto orange
  pomegranate: 
      print(4)
```





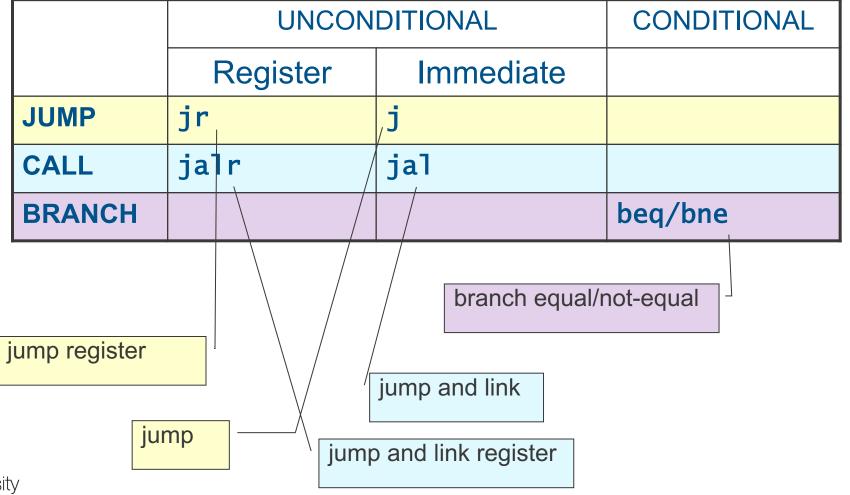


MIPS control transfer instructions

- MIPS does have several kinds of goto instructions:
 - They are called control transfer instructions
- They are needed to change the control flow
- They replace the value in the Program Counter (PC) with the address value of a specified "destination"
 - PC = destination address
- Control is transferred to this new destination point at the next "fetch instruction" phase
 - The next instruction loaded into IR is the one at the destination
- They can be conditional and unconditional
- And use a register for the address or an immediate (label)



Control Transfer Instructions



Jump and call instructions

jump (go) to label

```
j foo  # set PC = foo
# so, go to foo
```

jump to address contained in register

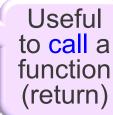
```
jr $t0  # set PC=$t0, so go to the
# address contained in $t0
```

jump to label and link (remember origin)

```
jal foo  # $ra = PC+4; PC = foo, same as j foo
# but setting a return address
```

jump to register and link (remember origin)

```
jalr $t0  # $ra = PC+4; PC = $t0, same as
# jr $t0 but setting return address
```



Jump instructions – Example

```
.text
                                                   1020
                                    # $t0 ← 5
1020
      main: addi $t0,$0,5
                                                   1024
                                                   1028
            addi $t1,$0,9
1024
                                    # $t1 ← 9
                                                   102C
1028
            addi $t2,$0,-3
                                    # $t2 ← -3
                                      $ra ←1030
                                                   ,1038
102C
            jal subr
                                      $t3 ← 34
                                                   1034
            add $t3,$t3,$t3
1030
                                                   1044
1034
                 end
                                                   103C
      subr: add $t3,$t1,$t0
1038
                                                   1040
            sub $t3,$t3,$t2
103C
                                                   1030
            jr $ra
1040
1044
       end:
```

PC

MIPS Instruction Format

- Remember: every MIPS instruction occupies 4 bytes of memory (32 bits)
- Remember: each instruction contains
 - opcode
 - operation code: specifies type of instruction
 - operands
 - values or location to perform operation on
 - registers
 - immediate (constant) numbers
 - labels (addresses of other lines of program)



MIPS Instruction Format

Three general assembler formats

```
sub $t0, $t1, $t2
```

R (for "register") format instruction: three registers

<u>sub</u>tract the contents of register <u>\$t2</u> from the contents of register <u>\$t1</u>; put the result in register <u>\$t0</u>

addi \$v0, \$a2, 742

I (for "immediate") format instruction: two registers and one immediate operand

add the immediate number 742 with the contents of register \$a2; put the result in register \$v0

j foo

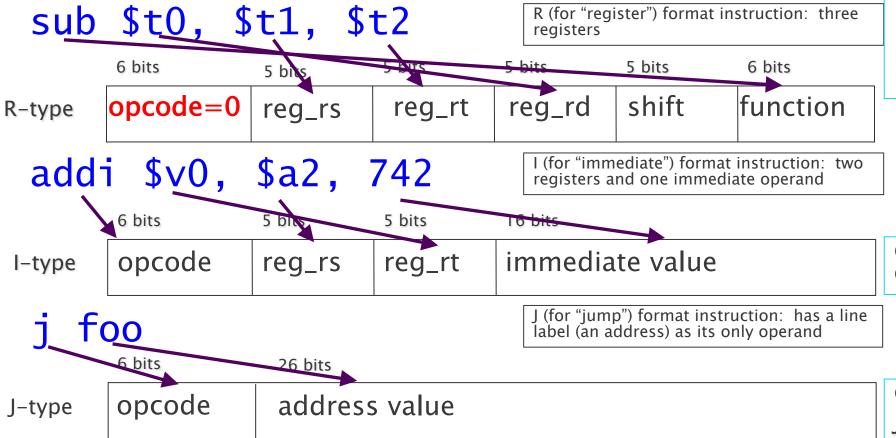
J (for "jump") format instruction: has a line label (an address) as its only operand

jump (go) to the line with the label <u>foo</u> and continue running from there



MIPS Instruction Format (cont'd)

Three general assembler formats



encodes ALU ops: +, -, &, |, ^, ~| <<, >> rs==rt?, *, / plus instr: jr, mfhi/lo, syscall,...

encodes everything else

encodes ALU ops: J and jal

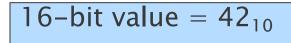
Dealing with immediate values

- Consider the instruction:
 - addi \$t0, \$t1, 123
- Immediate value 123 is represented using 16 bits (I-type instruction)
- But it must be put in 32-bit format to go into the ALU
- So how do we make a 16 bit value into a 32-bit value?
 - It depends on the meaning of the immediate
 - In turn, this depends on the actual instruction



Zero Extension

- Used by andi, ori, xori
- Zeroes are added in front



000000000101010

fill with zeroes







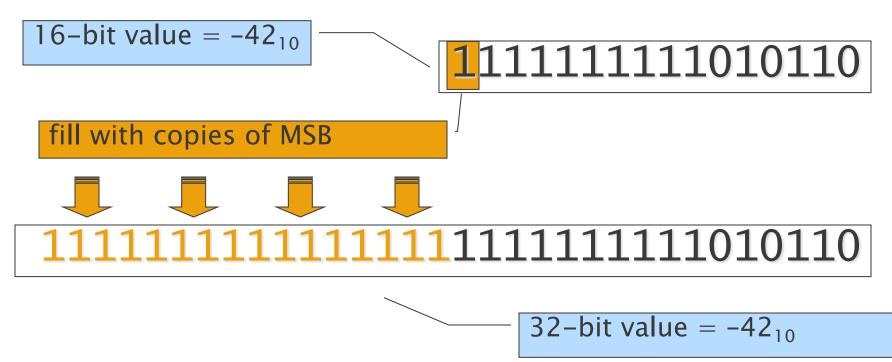


32-bit value = 42_{10}

Sign Extension

This is one of the main differences between addi (sign extended) and ori (zero extended)

- Used for addi,sll,srl,sra,slti
- Sign is extended by duplicating MSB (sign bit)



Summary

Know how to use unconditional control transfer instructions

Know the three different instruction formats

The following 4 slides are not examinable but might be interesting for some



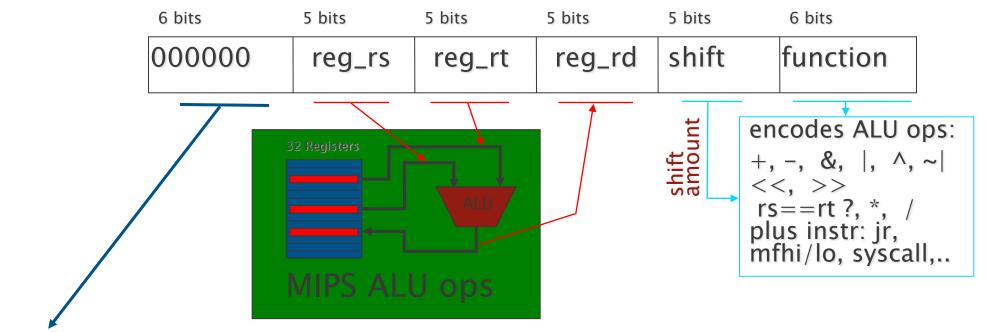
"J-type" instructions

6 bits	26 bits
opcode	address operand

- Only two MIPS instructions are j-type:
 - j label and jal label2 (opcodes 2 & 3, respectively)
 - all label addresses in text segment are multiples of 4
 - assembler encodes 26 bit operand = (label addr >> 2)
 - this encoding improves maximum range of a jump
 - during instruction execution the operand is shifted << 2 and used to replace lowest 28 bits of the PC
 - jal is used exclusively for function calls
 - before PC is modified, PC+4 (=address of instruction after jal) is saved in register \$ra (
 return address)



"R-type" Instructions



- All R-type instructions have opcode = 0
- They use register operands exclusively
- 6-bit function field specifies exact action
 - 28 different instructions encoded by this field

"I-type" Instructions

6 bits	5 bits	5 bits	16 bits
opcode	reg_rs	reg_rt	immediate operand

- All remaining instructions are I-type
- Immediate bit field used in three ways:
 - In (real) load/store instructions
 - Actual address referenced is sum of reg rs and imm field
 - In conditional branching instructions
 - imm is the "branch distance" measured in memory words from the address of the following instruction
 - In immediate ALU arithmetic or logic ops
 - Saves time where one operand is a small imm constant

I-type Instruction – Example

Instruction's components encoded in binary

