Class Starting at Tt 10 mins.

# features of a dedicated DBML.

- 1) Data Backups
- 2) Concurring
- 2) Security -> Enoughton
- 4) Efficient Storage and Retrieval of data

# Non- Relational DBMS. (Nosal DBMS)

- Don't follow relational model of storing data.

### Students Table.

B   10000	
1	] ? Rows in this table.
2	4
3	
ን	] <i>J</i>

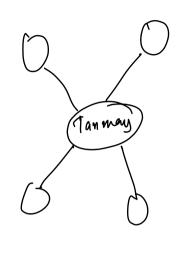
### facebook friends.

M		

1 <del>2611</del> .	
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### Tangmay: { 4, 48, 1001, 1000 5, 1006]



# Types of NOSAL DOMS.

- 1) Graph Rased (NWYT)
- 2) Colymnan (Cassandra.)
- 3) Document Type ( Manyo DB, elastic Jeanch, Amazon Doannent DB, etc.).
- 4) Key-Value jois. (Redis)

Relational Databases. -> Most widely used model.

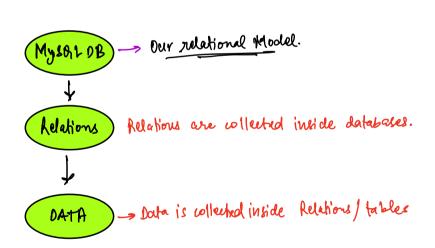
-> Adahonal Model is the model behind MysaL.

Relational Model segresents data as collection of relations

Database: - Collection of data.

Relation: Nothing but a table of values.

	Students table	<b>'</b>	_	Instructors table
bi	first name	Last-name		
l	- Klok	Gruph.	-1	<u> </u>
2	Subba	fleady.	Entity	-
	<i></i>			~



### Terms !-

- 1) Relation: Juble.
- 2) Attributes of Relation :- Columns of a table.
- 3) Tuple: Rew in a table. (set of values) (1, Tanmay, Upadhayay)
- y) Degree of a table/ Relation: No. of attributes of a relation No. of columns of a table
- 5) Cardinality of a table/febation! No of taples in a relation No of rows in a table.

# Properties of relations / tables.

1) Every ron should be unique (one column should have different value)

#### Students table:

	Stages of a large			
bi	first_name	Last-name	PsP	
1	Samit	Kumar	90	
2	Jumit	Kumar.	90	

2) Values in cells should be atomic conly single values.

Notit / JSON/ Collection.

first-Name	Last-Name	Phone Mumbers		
Alok	Singh	[8115, 85.5]	stoogy it as a sta	rivy is a hack.

#### Students table:

	Stopperel's leaf or			
	λí	first_name	Last-name	929
	1	Samit	Kumar	90
	2	Judan?	Sridhau	90
	,			
l				

#### Phone Numbers:

id	shu-jol.	Phone No	
1	1	8195910235	
2	1	8195910337	
M	2	8605	
4	2	85°5	

Stu-id is a foreign key

### 3) Order of columns should not matter.

scheet flame, psp, last nam from Students;

4) Order of now should not matter.

(.g. select \* from students where name = 'Alok'

Low 2 Now I Rowl Row 2 Row 3 low3 d rod 125

### SET THEORY

Adetion = set = Table.

Set {A,B3 = Set {B,A3

List [A,B] + List [B,A]

Every tuple is a value inside a set

Table = Set of (), (), (), (), (). ?

The order of there tuples / nows in this set table does not

Cardinality of a set = # of values in a set landinality of a table = # rows in a fable.

gel me,

## A Identify types in a Relation/now in a table.

### Students Table.

first_name	Last_name	Emeil)	Thore
Samif	Kumar	<b>\</b>	5
Sushi!	Sridhav	~	1

# Rey: Affribertes | let of attributes that can uniquely identify a now in a relation.

Affributes.	key. Kuper key	Candidate Key
f-name	*	X
l. name	X	× ×
(f_name, l-name) (f_name, l-name, e_mail)	7	×
17. name, 1-name, e-min),	V 4	×
email		
phone -no.		<b>○</b>
(email, phone)		×

# Types of Keys.

- 1) Super Key: A set of attributes that can uniquely identify a son.
- 2) Candidate key: A Key of min size that com uniquely identify a row. Candidate kay is a super key of win. size.
- 3) Primary (ley: One of the candidate key that is chosen by the Database Architect as the lay of the table.

few problems with choosing evail phone no. as PK.

- 1) It can change
- 2) It will be updated in mapping ) referring tables where storing lang etaing risn't very ideal.

Ideally, a PK should never change Create another column called student id. Super kay, candidate key, PK.

- Key with  $\geq 2$  attributes. 4) Composite Key:
- 5) Foreign Key:-

(9:22) <u>Resuming</u>.