



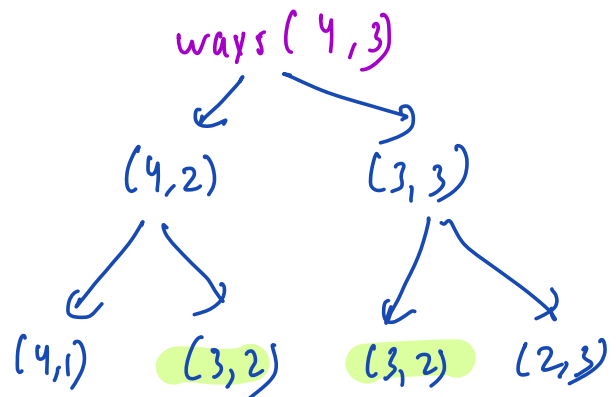
Q1  $\Rightarrow$  No of ways to go from  $(0,0)$   $\rightarrow$  (BR cell)  
 TL BR

	0	1	2
0	TL		
1			
2			BR

$(0,0), (0,1), (0,2), (1,2), (2,2)$   
 $(0,0), (0,1), (1,1), (1,2), (2,2)$   
 $(0,0), (0,1), (1,1), (2,1), (2,2)$   
 $(0,0), (1,0), (1,1), (1,2), (2,2)$   
 $(0,0), (1,0), (1,1), (2,1), (2,2)$   
 $(0,0), (1,0), (2,0), (2,1), (2,2)$

	0	1	2	3
0	1	1	1	1
1	1	2		
2	1			
3	1			
4	1			

$(3,3)$   
 $(4,2)$   
 $(4,3)$



$$dp[i][j] = dp[i-1][j] + dp[i][j-1]$$

```
int ways (int N, int M)
{
```

```
    int dp[N][M];
```

```
    for (i=0; i < N; i++)
    {
```

```
        for (j=0; j < M; j++)
        {
```

```
            if (i==0 || j==0)
                dp[i][j]=1;
```

```
            else
```

```
                dp[i][j] = dp[i-1][j] + dp[i][j-1];
```

```
        }
```

```
    return dp[N-1][M-1];
```

```
}
```

TC:  $O(N * M)$

SC:  $O(N * M)$

$N \times M$

	0	1	2	3
0	1	1	1	1
1	1	2	3	4
2	1	3	6	10
3	1			
4	1			

$(3,3)$   
 $(4,2)$   
 $(4,3)$

$M$

1	23	26	40
---	----	----	----

$O(2) \Rightarrow$  Number of ways to go from  $(0,0) \rightarrow (DR)$

$mat[i][j] = 0 \rightarrow$  it means blocked

	0	1	2	3
0	1	1	1	1
1	1	<del>2</del>	3	<del>4</del>
2	<del>1</del>	3	6	10
3	1	<del>2</del>	3	4
4	1	1	1	1

$dp$

	0	1	2	3
0	1	1	1	1
1	1	0	1	0
2	0	0	1	1
3	0	0	1	2
4	0	0	1	3

$$dp[i][j] = \begin{cases} \text{if } (mat[i][j] \neq 0) \{ dp[i][j] = 0 \} \\ \text{else } \{ dp[i-1][j] + dp[i][j-1] \} \end{cases}$$

```

int ways (int mat[][], int N, int M)
{
    // col

```

```

for (i=0 ; i < n ; i++)
{
    if ( mat[i][0] == 1 )    dp[i][0] = 1
    else
        break;
}

```

// row

```

j = 0 ; j < m ; j++
if ( mat[0][j] == 1 )    dp[0][j] = 1
else break;

```

}

```

for (i=1 ; i < N ; i++)
{

```

```

    for (j=1 ; j < n ; j++)
    {

```

```

        if ( mat[i][j] == 0 )
            dp[i][j] = 0;

```

else

```

        dp[i][j] = dp[i-1][j] +

```

```

            dp[i][j-1]

```

}

}

```

return dp[n-1][m-1];

```

}

TC:  $O(N \times M)$

```
int ways ( int mat[10][10], int N, int M)
{
```

```
    int dp[N][M] = {-1}
    if (mat[0][0] == 0) return 0;
```

```
    func (mat, dp, N-1, M-1)
```

TODO  
 (0,0) → 1

```
    }
    int func (int mat[10][10], int dp[10][10], int i, int j)
```

```
    if (i < 0 || j < 0) return 0;
    if (mat[i][j] == 0) return 0;
```

```
    if (dp[i][j] == -1)
```

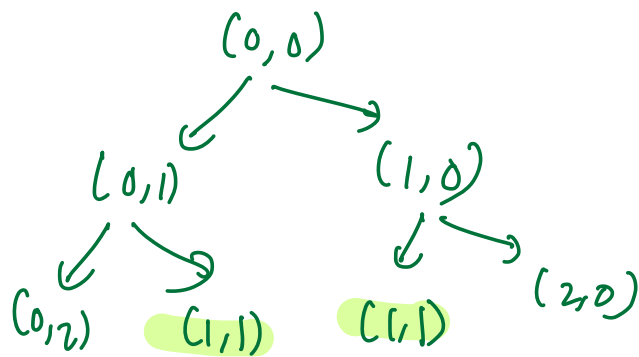
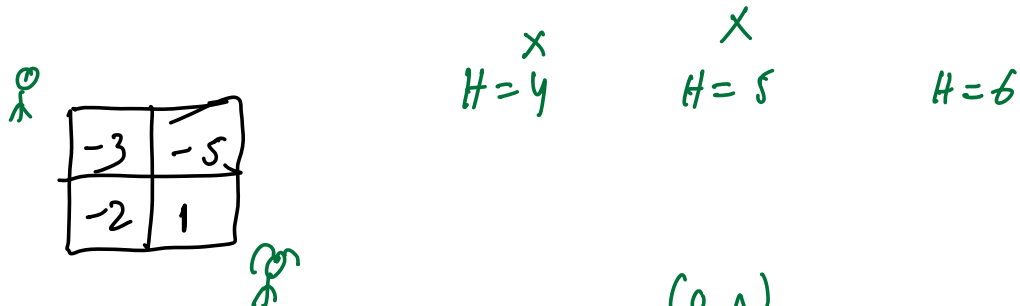
```
        dp[i][j] = {
            func (mat, dp, i-1, j)
            +
            func (mat, dp, i, j-1)
        }
```

```
    }
    return dp[i][j];
```

```
}
```

Q3  $\Rightarrow$  Min Cost to reach  $(0, 0) \rightarrow BR$

// Given mat  $(N)(M)$ , where each cell indicates health gained, min-health required at  $0,0$  such that we can reach  $(N-1)(M-1)$



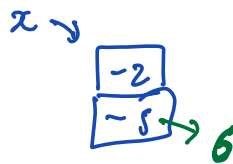
Case ~ I



$$x - 5 = 1$$

$$x = 6$$

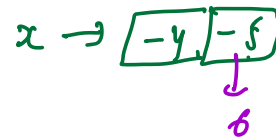
Case 2



$$x + (-2) = 6$$

$$x = 8$$

Case 3



$$x + (-4) = 6$$

$$x = 10$$

Case 4

-3	-2
-4	-5

(Arrows: from -2 to 0, from -5 to 6, from -4 to 10)

$$x + (-3) = \min(0, 10)$$

$$x = 0 + 3$$

$$= 3$$

Case 5

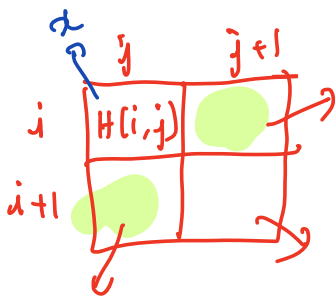
6	-2
-4	-5

(Arrows: from 6 to 0, from -5 to 6, from -4 to 10)

$$x + 6 = \min(0, 10)$$

$$x + 6 = 0$$

$$x = -6$$



$$dp[i, j+1]$$

$$dp[i+1, j]$$

$$x + H[i, j] = \min(dp[i, j+1], dp[i+1, j])$$

$$x = \min(dp[i, j+1], dp[i+1, j]) - H[i, j]$$



$dp[i][j] = \text{min health req } C[i, j] \rightarrow [N-1, M+1]$

$$dp[i][j] = \max(1, \min(dp[i, j+1], dp[i+1, j]) - H[i, j])$$

```

func (int H[][], int N, int M, int dp[][],
      int i, int j)
{

```

```

    if (dp[i][j] == -1)
    {

```

```

        int a = fun(H, N, M, dp, i, j+1)

```

```

        int b = fun(H, N, M, dp, i+1, j)

```

```

        dp[i][j] = max(1, min(a, b) - H[i][j]),

```

```

    }

```

---



}