

1.) Using empname as a clustered index is possible only when every employee will have a unique name. If this is ensured, the tuples will be organized according to empname alphabetically.

Using empid as a clustered index is definitely possible considering everyone already has a unique id assigned to them. The tuples will be organized according to empid.

Using both empname & empid as clustered indexes may not be possible but it is possible to have one clustered index and one non-clustered index.

Vikash Singh (19bcs113)

2.) DDL is important in representing information in DBMS because it is used to describe external and logical schemas.

•) DML is used to update and access data, it is not important for representing data.

Vikash Singh (196CS113)

3rd) True, A DBMS is typically shared among many users. Transactions from these users can be interleaved to improve the execution time of user's queries. By interleaving queries, users do not have to wait for other user's transactions to complete fully before their own transaction begins. Without interleaving, if user A begins a transaction that will take 10 seconds to complete, and user B wants to begin a transaction, user B would have to wait an additional 10 seconds for user A's transaction to complete before the database would begin processing user B's request.

4th a.) A user must guarantee that his or her transaction does not corrupt data or insert nonsense in the database. For example, in a banking database, a user must guarantee that a cash withdrawal transaction accurately models the amount a person removes from his or her account. A database application would be worthless if a person removed 20 dollars from an ATM but the transaction set their balance to zero!

b.) A DBMS must guarantee that transactions are executed fully and independently of other transactions. An essential property of a DBMS is that a transaction should execute atomically, or as if it is the only transaction running. Also, transactions will either complete fully, or will be aborted and the database returned to its initial state. This ensures that the database remains consistent.

Virash Singh 19BCS113

Ans)

Yes, we can determine the key of relation with the help of instance. eg. In a one to many relation we can consider the column / attribute with unique values as a primary key.

Vikash Singh (196cs113)

7th

$P(R_1, \text{catalog})$

$P(R_2, \text{catalog})$

$$\pi_{R_1 \cdot \text{pid} \wedge R_1 \cdot \text{pid} = R_2 \cdot \text{pid} \wedge R_1 \cdot \text{sid} \neq R_2 \cdot \text{sid}}(R_1 \times R_2)$$

Using the following:-

SID	PID	COST
1	1	\$10.00
2	1	\$9.00
2	3	\$34.00
3	1	\$11.00

$R_1 \times R_2$ gives

SID	PID	COST	SID	PID	COST
1	1	\$10.00	1	1	10
1	1	\$10.00	2	1	9
1	1	\$10.00	2	3	34
1	1	\$10.00	3	1	11
2	1	\$9.00	1	1	10
2	1	\$9.00	1	1	9
2	1	\$9.00	2	1	34
2	1	\$9.00	2	3	11
2	1	\$9.00	3	1	10
2	3	\$34.00	1	1	9
2	3	\$34.00	2	1	34
2	3	\$34.00	2	3	11
3	1	\$11.00	3	1	11
3	1	\$11.00	1	1	10
3	1	\$11.00	2	1	9
3	1	\$11.00	2	3	34
3	1	\$11.00	3	1	11

$\sigma R_1 \cdot \text{pid} = R_2 \cdot \text{pid}$ gives us:

SID	PID	Cost	SID	PID	Cost
1	1	100	1	1	10
1	1	10	2	1	9
1	1	10	3	1	11
2	1	9	1	1	10
2	1	9	2	1	9
2	1	9	3	1	11
2	8	34	2	3	34
3	1	11	1	1	10
3	1	11	2	1	9
3	1	11	3	1	11

$\sigma R_1 \cdot \text{pid} = R_2 \cdot \text{pid} \wedge R_1 \cdot \text{sid} \neq R_2 \cdot \text{sid}$ give us:-

SID	PID	Cost	SID	PID	Cost
1	1	10	2	1	9
1	1	10	3	1	11
2	1	9	1	1	10
2	1	9	3	1	11
3	1	11	1	1	10
3	1	11	2	1	9

Projecting on PID gives us a single part number -
eliminating the duplicates.

Vinodh 1965113

7th

SOL:-

SELECT C.sid
FROM catalog C

WHERE EXISTS (SELECT C1.sid
FROM catalog C1
WHERE C1.pid = C.pid AND
C1.sid \neq C.sid)

Vikash Singh 196CS113

(8) Invalid query

Explanation :- This relational algebra statement does not return anything because of the sequence of projection operators. Once the sid is projected, it is the only field in the set. Therefore, projection on some will not return anything.