

PROBLEM - 14.4 =

Multi class classifier via matrix least squares.

Show that the coefficient vectors $\beta_1, \beta_2 - \beta_k$ can be found by solving the matrix least squares problem of minimizing $\|X\beta - Y\|^2$ where β is the $n \times k$ matrix with columns $\beta_1 - \beta_k$ and Y is an $N \times k$ matrix

while training the model =

- ① = N (n -vectors) are given for training.
- ② = correct outcomes of labels for these inputs are also given.
- ③ = for these values of X and y we can say least square Problem is minimized.
- ④ = Hence differentiating the problem and putting $\hat{\beta}$ in the equation will be equal to 0.

⑤ = Hence using the given α and γ
we can calculate β .

Part-a =

Matrix } tells us the Prediction
represent } of multiclass classifier.
labels.

The i th row of matrix γ represent
the Prediction value of the
 i th datapoint among N -datapoints
for all the k -labels which the
classifier classifier.

Similarly γ_{ij} represent that whether
its i th datapoint is near to j th
label or not.

Part. b =

It is given that the rows of X are linearly independent.

It means that columns of transposed matrix X^T will be linearly independent.

We know that at minima of any function the value of differentiation is equal to 0.

It means to find minima, we have to solve (differentiation of $f(x) = 0$)

$$\text{here to minimize } f(x) = \| X^T \beta - y \|^2$$

for known X and y values. we can calculate $\hat{\beta}$ which will minimize the problem.

$\hat{\beta}$ is the minima, hence $\nabla f(\hat{x}) = 0$

$$\nabla f(\hat{\beta}) = 2(X^T)^T (X^T \hat{\beta} - y) = 0$$

$$\Rightarrow 2X(X^T \hat{\beta} - y) = 0$$

$$\Rightarrow 2X X^T \hat{\beta} - 2y = 0$$

$$= X X^T \hat{\beta} = y$$

$$\Rightarrow \hat{\beta} = (X X^T)^{-1} y$$

$$\Rightarrow \hat{\beta} = (X^T)^+ y$$

Hence $\boxed{\hat{\beta} = (X^T)^+ y}$

PROBLEM 15.4 =

Robust approximate solution of linear equation.

given that $Ax = b$

If A is invertible then $x = A^{-1}b$

for given k versions of A =

$$A^{(1)}, A^{(2)} \dots A^{(k)}$$

choose x to minimize =

$$\|A^{(1)}x - b\|^2 + \dots + \|A^{(k)}x - b\|^2$$

we can represent this equation in form of Norm of another matrix L .

$$L = \left\| \begin{bmatrix} A^{(1)}x - b \\ A^{(2)}x - b \\ \vdots \\ A^{(k)}x - b \end{bmatrix} \right\|$$

we can define -

$$A' = \begin{bmatrix} A^{(1)} \\ A^{(2)} \\ A^{(3)} \\ \vdots \\ A^{(k)} \end{bmatrix}_{K \times 1}, \quad b' = \begin{bmatrix} b \\ b \\ b \\ \vdots \\ b \end{bmatrix}_{K \times 1}$$

for a general equation $Cy - d = 0$

we have $y = (C^T C)^{-1} C^T d$

Hence in our case we can write

$$x^{rob} = (A'^T \quad A')^{-1} \quad A'^T \quad b'$$

which reduces to -

$$x^{rob} = \left(\sum_{k=1}^K (A^k)^T A^k \right)^{-1} \left(\sum_{k=1}^K (A^k)^T b \right)$$

When $|k|=1$ then

$$x^{obs} = \left(\left(A^{(1)} \right)^T A^{(1)} \right)^{-1} \left(\left(A^{(1)} \right)^T b \right)$$

$$x^{obs} = \left[\left(A^{(1)} \right)^T \left(A^{(1)} \right)^T \right]^{-1} \left[A^{(1)} \right]^T b$$

$$x^{obs} = \left(A^{(1)} \right)^{-1} \left[\left(A^{(1)} \right)^T \left(A^{(1)} \right)^T \right]^{-1} b$$

$$x^{obs} = \left(A^{(1)} \right)^{-1} b$$

proved

PROBLEM 15.1) = General Pseudo Inverse

Part a = what is the Pseudo Inverse
of $n \times n$ zero matrix.

given that A is zero matrix of $n \times n$

$$A = [0]_{n \times n}$$

also pseudo Inverse By kernel
trick identity is given as -

$$A^+ = \lim_{\lambda \rightarrow 0} (A^T A + \lambda I)^{-1} A^T$$

$$A^+ = \lim_{\lambda \rightarrow 0} \left[[0]_{n \times n}^T [0]_{n \times n} + \lambda I \right]^{-1} [0]_{n \times n}^T$$

$$A^+ \Rightarrow \lim_{\lambda \rightarrow 0} [0_{n \times n} + \lambda I]^{-1} [0]_{n \times n}$$

$$A^+ \Rightarrow 0_{n \times n}$$

$$\text{Since } A^T = O_{n \times m} = [0]_{m \times n}^T$$

we can say that for a zero matrix
the pseudo Inverse could be defined
as its transpose.

Part-b = given that A has linearly independent columns.

$$A^T = \lim_{\gamma \rightarrow 0} (A^T A + \gamma I)^{-1} A^T$$

as γ reaches 0, objective function

$$\|Ax - b\|^2 + \gamma \|x - x^{des}\|^2$$

reaches to $\|Ax - b\|^2$

So if $[A]_{m \times n}$ matrix has linearly independent columns.

Then it means that,

$$\hat{x} = A^+ b$$

and $A^+ = (A^T A)^{-1} A^T$

Matrix inverse is continuous function

Hence $\lim_{\lambda \rightarrow 0} (A^T A + \lambda I)^{-1}$

$$= \left(\lim_{\lambda \rightarrow 0} A^T A + \lim_{\lambda \rightarrow 0} \lambda I \right)^{-1}$$

$$= \left(\lim_{\lambda \rightarrow 0} A^T A + 0 \right)^{-1}$$

$$= (A^T A)^{-1} \longrightarrow \text{eqtn ①}$$

NOTE - given that =

$$\lim (f(x))^{-1} = (\lim f(x))^{-1}$$

if $f(x)$ = continuous function

and $\lambda \rightarrow 0$, also given that A has linearly independent columns.

$$A^+ = \lim_{\lambda \rightarrow 0} (A^T A + \lambda I)^{-1} A^T$$

$$A^+ = (A^T A)^{-1} A^T$$

[Putting values from equation ①]

Part-C - A has linearly independent rows.

As we solved in part b for linearly independent columns. Using the same method we can solve this part.

Here $A A^T$ is invertible, $\lambda \rightarrow 0$.

$$\text{also } \|Ax - b\|^2 + \lambda \|x - x_{\text{des}}\|^2 \approx \|Ax - b\|^2$$

we'll be having =

$$\hat{x} = A^T b = A^T (A A^T)^{-1} b$$

we can use the same idea in part b
that since matrix inverse is continuous
hence limit of the inverse of matrix
will be equal to inverse of limit.

$$\text{so } \lim_{\lambda \rightarrow 0} (A A^T + \lambda I)^{-1}$$

$$= \left(\lim_{\lambda \rightarrow 0} A A^T + \lim_{\lambda \rightarrow 0} \lambda I \right)^{-1}$$

$$= (A A^T)^{-1} \longrightarrow \text{eqtn ①.}$$

Since, $A^T = \lim_{\lambda \rightarrow 0} A^T (A A^T + \lambda I)^{-1}$

using eqtn ① =

$$A^T = A^T (A A^T)^{-1}$$

PROBLEM 1b.1) = Least distance Problem.

minimizing $\|x-a\|^2$
is subject to $Cx=d$.

We have to determine n-vector x .

Given = n-vector a .

2) = $P \times n$ matrix C

3) = P vector d .

We have to prove that =

Solution of this Problem is -

$$\hat{x} = a - C^T (C a - d)$$

Let's assume $y = x - a$

if, $(\gamma = x - a)$ Then $(x = \gamma + a)$

and we have to minimize $= \|\gamma\|^2$

Then according to the Problem
it will be subjected to,

$$\Rightarrow Cx = d$$

$$\Rightarrow C(\gamma + a) = d$$

$$\Rightarrow C\gamma + Ca = d$$

$$\Rightarrow C\gamma = d - Ca$$

After using least norm solution
we get =

$$\hat{\gamma} = C^\dagger (d - Ca)$$

$$\Rightarrow \hat{x} - a = C^\dagger (d - Ca)$$

$$\Rightarrow \hat{x} = a + C^\dagger (d - Ca)$$

$$\Rightarrow \hat{x} = a - C^\dagger (Ca - d)$$

Proved

PROBLEM. 5 =

Solving Least Square \Rightarrow minimize $\|Ax - b\|^2$
Problem

Given that columns of A are linearly independent.

Solution for the Problem is given by =

$$\hat{x} = A^T b = (A^T A)^{-1} A^T b$$

A^T = Pseudo Inverse of A.

① = QR factorization = The least square problem could be solved using QR factorization. In order to get \hat{x} we need to figure out Pseudo inverse of matrix A first which is (A^T) .

Using QR factorization we can determine the A^T . and then $\hat{\lambda}$ values could be calculated.

we know that A could be written as

$$A = QR.$$

so A^T could be found as follow =

$$A^T = (A^T A)^{-1} A^T$$

$$A^T = \left((QR)^T (QR) \right)^{-1} (QR)^T$$

$$A^T = (R^T Q^T Q R)^{-1} (QR)^T$$

$$A^T = (R^T R)^{-1} (QR)^T$$

$$A^T = (R^{-1} R^{-T} R^T Q^T)$$

$$A^T = R^{-1} Q^T$$

$$\text{Hence } \hat{x} = A^T b$$

$$\hat{x} = (R^{-1} Q^T) b$$

$$x = R^{-1} Q^T b$$

So QR-factorization algorithm will have following steps =

①= find the QR factorization of matrix A. ($A = QR$)
(requires $2n^3$ flops)

②= Algorithm will calculate $Q^T b$ first. (requires $2n^2$ flops)

③= Using Back Substitution method
solve the equation $Rx = Q^T b$
(requires n^2 flops)

This will give us \hat{x} value.

Complexity of QR-factorization algo =

$$= (2n^3 + 2n^2 + n^2) \text{ flops}$$

$$= (2n^3 + 3n^2) \text{ flops}$$

complexity = $O(n^3)$. flops.

② = cholesky factorization algo =

we can see that A is positive definite. Hence cholesky factorizat'n

can be used to solve the least squares problem.

It breaks given matrix A as follows

$$A = L L^T$$

where L is lower triangular matrix.

In order to calculate $\hat{x} = A^T b$
the algorithm follows below

Steps =

1) = Do cholesky factorization.

$$A = LL^T$$

(requires $(\frac{1}{3})n^3$ flops)

2) = Do forward substitution
and solve

$$Lr = b$$

(requires n^2 flops)

3) = Do backward Substitution
and solve

$$L^T \sigma L = V$$

(requires n^2 flops)

complexity of Cholesky Algorithm =

$$= \left(\frac{1}{3}n^3 + n^2 + n^2 \right) \text{ flops}$$

$$= \left(\frac{1}{3}n^3 + 2n^2 \right) \text{ flops}$$

$$= O(n^3) \text{ flops.}$$

PROBLEM-6 = Solving Regularized Least Squares.

The least square Problem

$$\underset{x}{\text{minimize}} \quad \|Ax - b\|^2 + \gamma \|x\|^2$$

could be solved in $\mathcal{O}(mn^2 + n^3)$.

if A is wide ($m < n$) we have to
reduce it to $\mathcal{O}(m^2n + m^3)$.

we know that,

$$\hat{x} = (A^T A + \gamma I)^{-1} A^T b$$

also matrix inversion lemma tells
us that =

$$(A^T A + \gamma I)^{-1} A^T = A^T (A A^T + \gamma I)^{-1}$$

Hence we can rewrite \hat{x} as -

$$\hat{x} = A^T(AA^T + \gamma I)^{-1}b$$

This method uses following =

① = It takes $2(m+n)m^2$ flops
for QR factorization.

② = A^T is a matrix
 $(AA^T + \gamma I)^{-1}b \rightarrow$ a vector.

Hence to calculate \hat{x} it takes
 mn flops for matrix vector
multiplication.

Total complexity =

$$\begin{aligned} & 2(m+n)m^2 + mn \\ &= 2m^3 + 2nm^2 + mn \end{aligned}$$

Hence complexity = $O(m^2n + m^3)$

PROBLEM 7 = Companion Matrix.

Part a = what is characteristic polynomial of C_2 .

Given $C = \begin{bmatrix} 0 & 0 & \dots & 0 & -c_0 \\ 1 & 0 & \dots & 0 & -c_1 \\ 0 & 1 & 0 & 0 & -c_2 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & 0 & \dots & 1 & -c_{n-1} \end{bmatrix}$

as given in the Slides the characteristic polynomial for any matrix A

is given as $P(x) = \det(xI - C)$

Hence for companion matrix we can write =

$$P(\gamma) = \det(\gamma I - C)$$

where $I = \text{Identity matrix}$.

Matrix $(\gamma I - C) =$

$$\begin{bmatrix} \gamma & 0 & 0 & \dots & 0 \\ 0 & \gamma & 0 & \dots & 0 \\ \vdots & & & & \\ 0 & 0 & \dots & \dots & \gamma \end{bmatrix} - \begin{bmatrix} 0 & 0 & \dots & 0 & -c_0 \\ 0 & 0 & \dots & 0 & -c_1 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \dots & 1 & -c_{n-1} \end{bmatrix}$$

$$= \begin{bmatrix} \gamma & 0 & \dots & -1 & 0 & c_0 \\ -1 & \gamma & \dots & -1 & 0 & c_1 \\ 0 & -1 & \dots & -1 & 0 & c_2 \\ \vdots & \vdots & & & & \\ 0 & 0 & \dots & -1 & \gamma + c_{n-2} & \end{bmatrix}$$

Then

$$\det(\gamma I - C) = \begin{vmatrix} \gamma & 0 & \dots & -1 & 0 & c_0 \\ -1 & \gamma & \dots & -1 & 0 & c_1 \\ 0 & -1 & \dots & -1 & 0 & c_2 \\ \vdots & \vdots & & & & \\ 0 & 0 & \dots & -1 & \gamma + c_{n-2} & \end{vmatrix}$$

using cofactor expansion correspondingly

to find now we obtain that =

$$\det(\gamma I - C) =$$

$$\Rightarrow \begin{vmatrix} \gamma & 0 & \cdots & 0 & q \\ -1 & \gamma & \cdots & 0 & c_2 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & \cdots & \gamma & \gamma + c_{n-1} \end{vmatrix}$$

$$+ (-1)^{n+1} c_0 \begin{vmatrix} -1 & \gamma & \cdots & 0 \\ 0 & -1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & -1 \end{vmatrix}$$

By Induction we know the first determinant value is =

$$\gamma^{n-1} + c_{n-1}\gamma^{n-2} + \cdots + c_2\gamma + q$$

the second determinant is $(-1)^{n-1}$
 Since it is an $(n-1) \times (n-1)$ triangular matrix determinant will be
 product of diagonal entries.

$$\text{2nd determinant} = (-1)^{n+1} c_0 (-1)^{n-1}$$

Therefore

$$\det(\gamma I - C) =$$

$$= \gamma (\gamma^{n-1} + c_{n-1} \gamma^{n-2} + \dots + c_2 \gamma + c_1) \\ + (-1)^{n+1} c_0 (-1)^{n-1}$$

$$\Rightarrow \gamma^n + c_{n-1} \gamma^{n-1} + \dots + c_1 \gamma + c_0$$

$$P(x) = \gamma^n + c_{n-1} \gamma^{n-1} + \dots + c_1 \gamma + c_0$$

Part b =

given that C has n -distinct eigenvalues $\lambda_1, \lambda_2, \lambda_3, \dots, \lambda_n$.

Show that V is the matrix of left eigen vectors of C .

i.e. $VC = \Lambda V$

let's calculate $VC =$

$$VC = \begin{bmatrix} 1 & \lambda_1^2 & \cdots & \lambda_1^{n-1} \\ 1 & \lambda_2^2 & \cdots & \lambda_2^{n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & \lambda_n^2 & \cdots & \lambda_n^{n-1} \end{bmatrix} \begin{bmatrix} 0 & 0 & \cdots & 0 & -c_0 \\ 1 & 0 & \cdots & 0 & -c_1 \\ 0 & 1 & \cdots & 0 & -c_2 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & \cdots & 1 & -c_{n-1} \end{bmatrix}$$

$$VC = \begin{bmatrix} \lambda_1 & \lambda_1^2 & \cdots & \lambda_1^{n-1} & c_0 - c_1\lambda_1 - c_2\lambda_1^2 - \cdots - c_{n-1}\lambda_1^{n-1} \\ \lambda_2 & \lambda_2^2 & \cdots & \lambda_2^{n-1} & c_0 - c_1\lambda_2 - \cdots - c_{n-1}\lambda_2^{n-1} \\ \vdots & \vdots & & & \vdots \\ \lambda_n & \lambda_n^2 & \cdots & \lambda_n^{n-1} & c_0 - c_1\lambda_n - c_2\lambda_n^2 - \cdots - c_{n-1}\lambda_n^{n-1} \end{bmatrix}$$

Here d_i are roots of characteristic equation hence characteristic equation will be 0 at d_i .

So V_C reduces to -

$$V_C = \begin{pmatrix} d_1 & d_1^2 & \cdots & d_1^n \\ d_2 & d_2^2 & \cdots & d_2^n \\ \vdots & \vdots & \ddots & \vdots \\ d_n & d_n^2 & \cdots & d_n^n \end{pmatrix}$$

also

$$\Lambda V = \begin{bmatrix} d_1 & 0 & 0 & \cdots & 0 \\ 0 & d_2 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & d_n \end{bmatrix} \quad \begin{bmatrix} 1 & d_1 & \cdots & d_1^{n-1} \\ 1 & d_2 & \cdots & d_2^{n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & d_n & \cdots & d_n^{n-1} \end{bmatrix}$$

$$= \begin{pmatrix} d_1 & d_1^2 & \cdots & d_1^n \\ d_2 & d_2^2 & \cdots & d_2^n \\ \vdots & \vdots & \ddots & \vdots \\ d_n & d_n^2 & \cdots & d_n^n \end{pmatrix} = V_C.$$

Hence $V_C = \Lambda V$

Problem.b = Circulant Matrix

given are circulant matrix and DFT matrix. C and F.

Part a =

from previous problem.7 we proved that.

$$VC = \Lambda V$$

also $\pi = \text{Diag}(\lambda_1, \dots, \lambda_n)$

Using the same analogy we can say that =

$$\Rightarrow FC = \Lambda F$$

where Λ should be $\text{Diag}(FC)$.

$$\Rightarrow C = F^{-1} \Lambda F$$

We are given that $F^{-1} = F^*$

Hence

$$C = F^* \Delta F$$

proved

Part-b = Solve linear equation -
 $Cx = b$ where C is circulant matrix.

$$Cx = b$$

$$= (F^{-1} \Delta F) x = b \quad \left[\begin{array}{l} \text{from} \\ \text{part.a} \end{array} \right]$$

$$\Rightarrow x = (F^{-1} \Delta F)^{-1} b \quad C = F^{-1} \Delta F$$

$$\Rightarrow x = (F)^{-1} (F^{-1} \Delta F)^{-1} b$$

$$x = F^{-1} \Delta^{-1} (F^{-1})^{-1} b$$

$$x = F^{-1} \Delta^{-1} F b$$

b/c $(AB)^{-1} = B^{-1} A^{-1}$

To calculate \Rightarrow

$$\overbrace{F^{-1} \Lambda^{-1} F b}^Q \xrightarrow{\rho} R$$

① $P = F b$ (multiplication)

$n(\log n)$ flops are required b/c FFT is used in calculation.

② = Then $Q = \Lambda^{-1} P$

This multiplication requires n flops.

③ = Then $R = F^{-1} B$

b/c inverse FFT is used so that's why we will need $n \log n$ flops.

Total complexity = for $\boxed{R = F^{-1} \Lambda^{-1} F b}$

$$n \log(n) + n + n \log n$$

$$= 2n \log n + n$$

$$= O(n \log n)$$

PROBLEM-9 = General Least Squares.

Part-a - for thin SVD of $A = U \Sigma V^T$

we can give the solution as -

$$\hat{x} = A^T b = V \Sigma^{-1} U^T b \rightarrow ①$$

Transforming equation ① by multiplying
By $A^T A$. we get -

$$A^T A \hat{x} = A^T A (A^T b) = A^T A (V \Sigma^{-1} U^T b)$$

now we know that =

$$A^T A \hat{x} = A^T A (V \Sigma^{-1} U^T b)$$

Substituting the value of $A = U \Sigma V^T$
in the right hand side we get,

$$A^T A \hat{x} = A^T \cdot (U \Sigma V^T) \cdot (V \Sigma^{-1} U^T b)$$

$$A^T A \hat{x} = A^T [U \Sigma V^T] V \Sigma^{-1} U^T b$$

$$A^T A \hat{x} = A^T [U \Sigma \Sigma^{-1} U^T b]$$

$$A^T A \hat{x} = A^T [U U^T b]$$

$$A^T A \hat{x} = A^T b \quad (\text{Proved})$$

Part.b - verify that it minimizes
the objective function.

lets solve $\|Ax - b\|^2$

$$\begin{aligned} \|Ax - b\|^2 &= \|Ax - A\hat{x} + A\hat{x} - b\|^2 \\ &= \|Ax - A\hat{x}\|^2 + \|A\hat{x} - b\|^2 \\ &\quad + 2 \cdot (Ax - A\hat{x})^T (A\hat{x} - b) \end{aligned}$$

↳ equation ①.

Simplify, $2(Ax - A\hat{x})^T (A\hat{x} - b)$

$$= 2(Ax - A\hat{x})^T (A\hat{x} - b)$$

$$= 2(x - \hat{x})^T A^T \cdot (A\hat{x} - b)$$

$$= 2(x - \hat{x}) (A^T A \hat{x} - A^T b)$$

$$= 2(x - \hat{x}) (A^T b - A^T b)$$

$$= 2(x - \hat{x}) (0)$$

$$= 0$$

$$\begin{aligned} & b(c) \\ & A^T A \hat{x} \\ & = A^T b \end{aligned}$$

Hence putting this value in equation ① =

$$\|Ax - b\|^2 = \|Ax - A\hat{x}\|^2 + \|A\hat{x} - b\|^2$$

This implies that =

$$\|A\hat{x} - b\|^2 \leq \|Ax - b\|^2$$

It means \hat{x} has minimized the objective function. Hence \hat{x} minimizes the objective function.

Problem.10 = QR SVD.

Given $m \times n$ matrix A . ($m \geq n$)

The QR SVD algorithm is given.

Part-b =

Show that $Q_1(R, R_0)$ is the QR decomposition of $R_0^T R_0$ where Q_1 is the orthonormal factor and R, R_0 is upper triangular factor.

From the given algorithm =

QR factorization is given as =

$$R_0^T = Q_1 R_1 \longrightarrow \textcircled{1}$$

$$R_1^T = Q_2 R_2$$

$$R_2^T = Q_3 R_3 \quad \text{and so on.}$$

So lets calculate $R_0^T R_0$.

$$R_0^T R_0 = (Q_1 R_1) R_0 \quad (\text{from eqtn})$$

$$R_0^T R_0 = Q_1 R_1 R_0 \quad \underline{\text{proved}}$$