

Parallel Merging pb

①

A_1, \dots, A_n B_1, \dots, B_n two sorted arrays

$$A_1 \leq \dots \leq A_n \quad B_1 \leq \dots \leq B_n$$

merge A, B into ^{array} C C_1, \dots, C_{2n} s.t.
 $C_1 \leq \dots \leq C_{2n}$

let $n = 2^k$, $\frac{n}{\log n} = r$, k, r integers
 $\log n = k$

sequential merge algorithm $O(n)$
time

1. $i = 1$ $j = 1$ $k = 1$, $A_{n+1} = +\infty$, $B_{n+1} = +\infty$
2. while $k \leq 2n$ do

$\left[\begin{array}{l} \text{if } A_i \leq B_j \left[\begin{array}{l} \text{then } C_k = A_i \quad i = i + 1 \\ \text{else } C_k = B_j \quad j = j + 1 \end{array} \right. \\ K = k + 1 \end{array} \right.$

(2)

parallel merge algorithm

1) partition the sorted array A in r groups each having $k = \lceil \log n \rceil$ elements as follows:

Group 1 : $A_1 \dots A_k$

Group 2 : $A_{k+1} \dots A_{2k}$

\vdots

Group i : $A_{(i-1)k+1} \dots A_{ik}$

\vdots

Group r : $A_{(r-1)k+1} \dots A_{rk}$

2) find r integers $j(1) \dots j(r)$ s.t.

$j(1)$ is the greatest index s.t. $A_k \geq B_{j(1)}$

$j(2)$ is the ^{gr} ~~smallest~~ index s.t. $A_{2k} \geq B_{j(2)}$

$j(i)$ is the ^{gr} ~~smallest~~ index s.t. $A_{ik} \geq B_{j(i)}$

$j(r)$ is the ^{gr} ~~smallest~~ index s.t. $A_{rk} \geq B_{j(r)}$

③

(3) this induces a partition of array B , into r groups:

Group 1 : $B_1, B_2, \dots, B_{j_{c1}}$

Group 2 : $B_{j_{c1}+1}, \dots, B_{j_{c2}}$

\vdots

Group i : $B_{j_{c(i-1)}+1}, \dots, B_{j_{ci}}$

\vdots

Group r : $B_{j_{c(r-1)}+1}, \dots, B_{j_{cr}}$

4) Observation : every entry in group 1 of A

is \leq to every entry of group ~~1, 2, 3, ...~~ 2, 3, ... of B

assign processor i to
merge group i of A
& group i of B

Merging group 1 of A & group 1
guarantees that the elem^s of B
have reached their final
position in $C(1:n)$

④

parallel merge algorithm