



Data Structure & Algorithms

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Insertion Sort

6 5 3 8 2 4

$$i + r = 1 + 2 + 3 + \dots + (n-1)$$

$$T \propto \frac{n(n-1)}{2}$$

$$T \propto n^2 - n$$

$$T \propto n^2$$

$$O(n^2)$$

pass 1 (1)

pass 2 (2)

pass 3 (3)

pass 4 (4)

pass 5 (5)

total: 15

0	1	2	3	4	5
6	5	3	8	2	4

5	6				
3	5	6			
3	5	6	8		
2	3	5	6	8	
2	3	4	5	6	8

```
for (i=1; i<n; i++) {
```

```
    temp = a[i];
```

```
    for (j=i-1; j>=0 && a[j]>temp; j--)
```

```
        a[j+1] = a[j];
```

```
    a[j+1] = temp;
```

```
}
```

① Calculate time complexity - general case



Insertion Sort

```
for (i = 1; i < n; i++) {  
    temp = a[i];  
    * for (j = i - 1; j >= 0 && a[j] > temp; j--)  
        a[j + 1] = a[j];  
    a[j + 1] = temp;  
}  
iter = n - 1  
T ∝ n - 1  
T ∝ n  
O(n) ← best case
```

② calculate time complexity - best case
1, 2, 3, 4, 5, 6

0	1	2	3	4	5
1	2	3	4	5	6

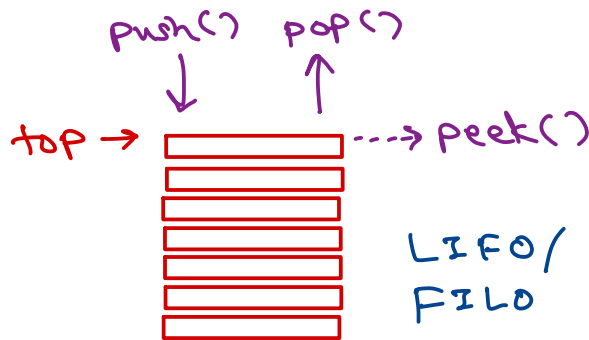
1 2
1 2 3
1 2 3 4
1 2 3 4 5
1 2 3 4 5 6



Stack and Queue

- Stack & Queue are utility data structures. - *data processing (not storage)*
- Can be implemented using array or linked lists.
- Usually time complexity of stack & queue operations is $O(1)$. *$T=k$*
- Stack is Last-In-First-Out structure.
- Stack operations

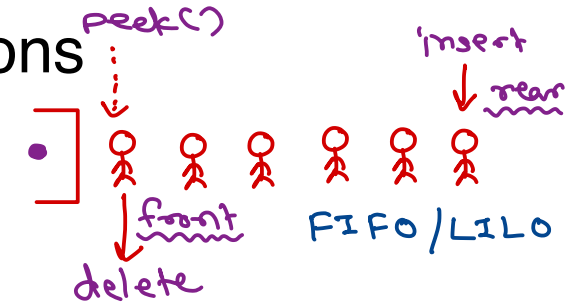
- push()
- pop()
- peek()
- isEmpty() ✓
- isFull()* ✓



- Simple queue is First-In-First-Out structure.

- Queue operations

- push()
- pop()
- peek()
- isEmpty() ✓
- isFull()* ✓



- Queue types

- Linear queue ✓
- Circular queue ✓
- Deque - *Double ended queue*
- Priority queue → *ele with highest priority comes out first. (internally eles are maintained in order) time complexity $O(1)$ X*
efficiently impl using heap DS.



Linear Queue - using array

init:

$$f = -1$$

$$r = -1$$

push:

$$r++;$$

$$arr[r] = val;$$

full:

$$r == \text{max} - 1$$

empty:

$$f == r$$

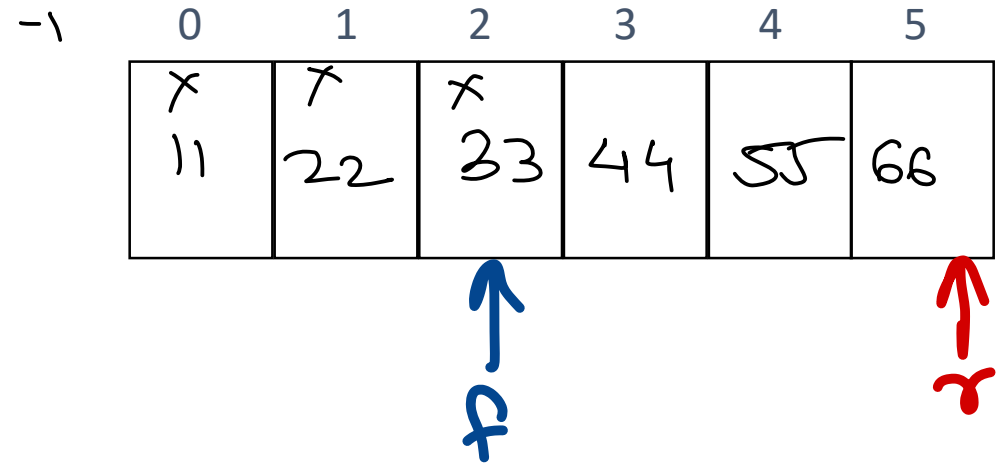
pop:

$$f++;$$

$$val = arr[f];$$

peek:

$$val = arr[f+1];$$



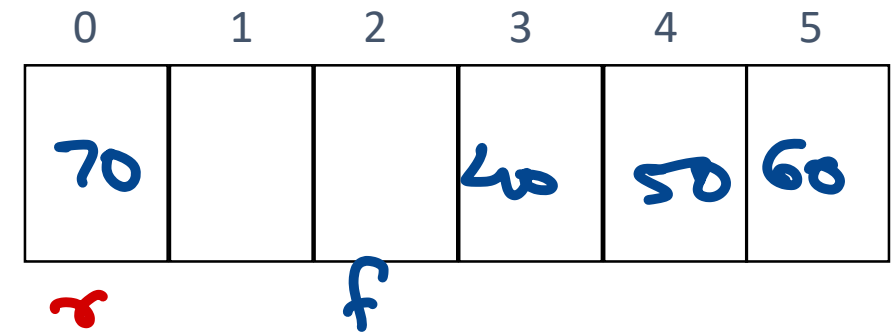
In linear queue, when we reach last ele ($\text{max} - 1$), then further eles cannot be added i.e. queue full condition.

In this case there may be few spaces empty at start of array. Thus lin que not doing mem. utilization properly \rightarrow limitation.



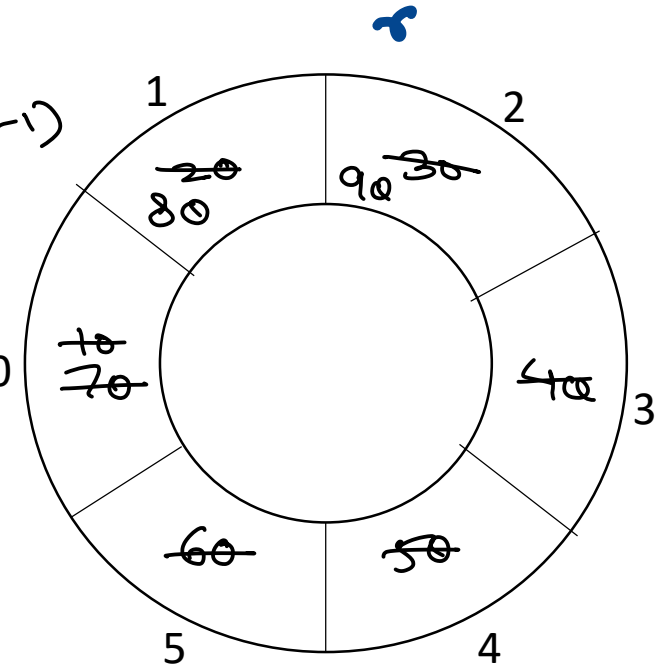
Circular Queue

- In linear queue (using array) when rear reaches last index, further elements cannot be added, even if space is available due to deletion of elements from front. Thus space utilization is poor.
- Circular queue allows adding elements at the start of array if rear reaches last index and space is free at the start of the array.
- Thus rear and front can be incremented in circular fashion i.e. 0, 1, 2, 3, ..., n-1. So they are said to be circular queue. $\rightarrow 0, 1, 2, 3, \dots$
- However queue full and empty conditions become tricky.



Circular incr
 $r++;$
if ($r == \text{max} - 1$)
 $r = 0;$

Circular incr
 $r = (r + 1) \% \text{max}$
 $\begin{cases} (-1 + 1) \% 6 = 0 \\ (0 + 1) \% 6 = 1 \\ (1 + 1) \% 6 = 2 \\ (2 + 1) \% 6 = 3 \\ (3 + 1) \% 6 = 4 \\ (4 + 1) \% 6 = 5 \\ (5 + 1) \% 6 = 0 \end{cases}$



Circular Queue

effective queue: $f+1$ to r

init:

$$r = -1$$

$$f = -1$$

push:

$$r = (r + 1) \% \text{max}$$

$$\text{arr}[r] = \text{val};$$

pop:

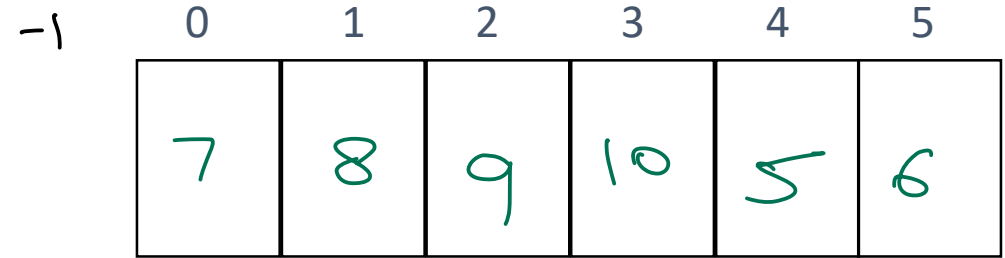
$$f = (f + 1) \% \text{max}$$

$$\text{val} = \text{arr}[f];$$

peek:

$$i = (f + 1) \% \text{max}$$

$$\text{val} = \text{arr}[i];$$



f r

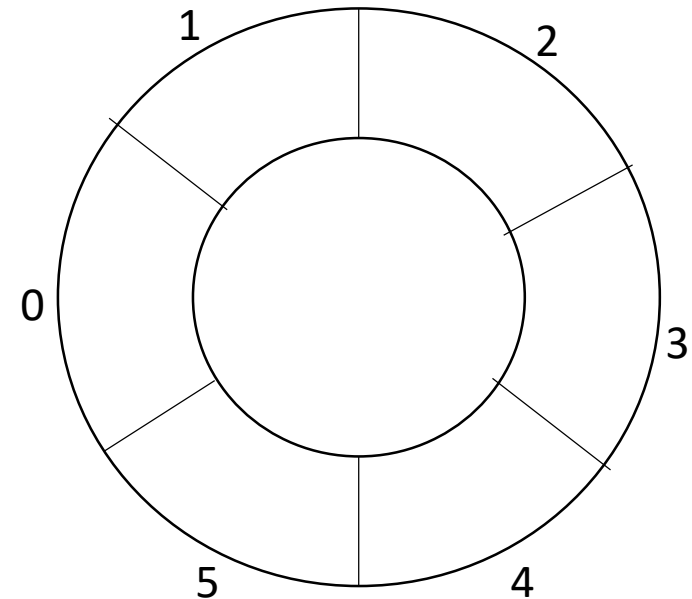
empty:

$$f == r$$

full:

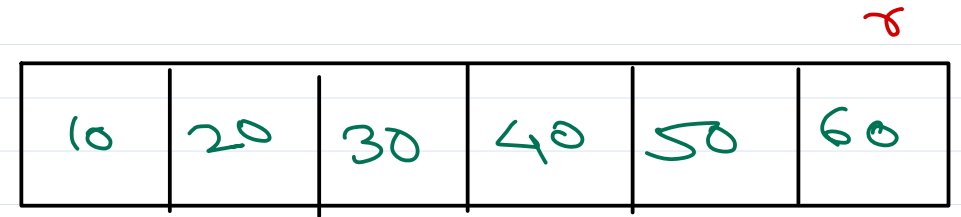
$$f == r$$

Same
Cond'n



que full

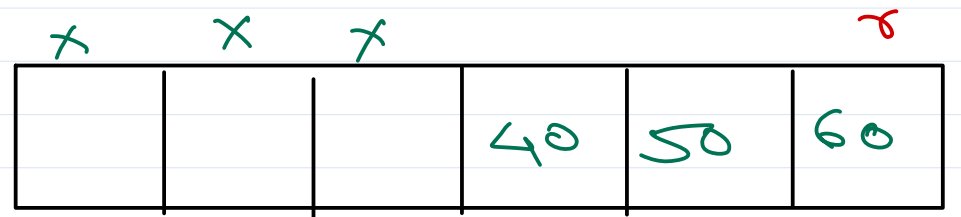
$(f == -1 \ \&\& \ r == \text{max} - 1)$ ←



4

push:

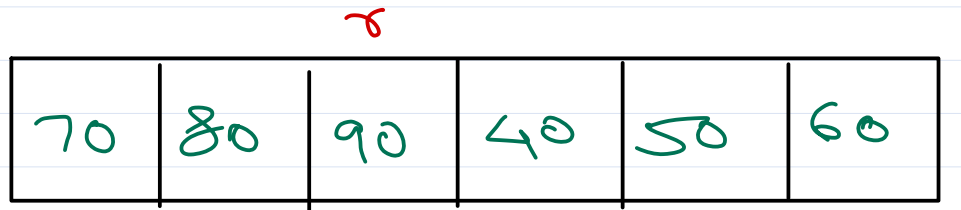
$r = (r + 1) \% \text{max}$
 $a[r] = \text{val};$



4

OR

← $(r == f \ \&\& \ r != -1)$



4



que empty

$(r == f \ \&\& \ r == -1)$ init end

pop:

$f = (f + 1) \% \text{max}$

$\text{val} = \text{arr}[f];$

if $(f == r)$ {

$f = -1;$

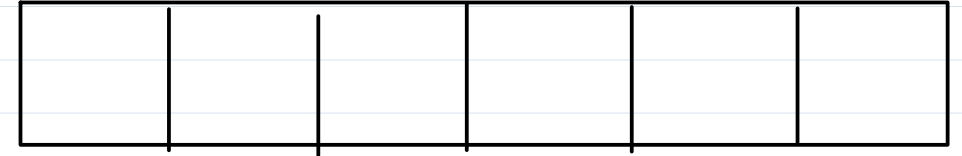
$r = -1;$

3

r

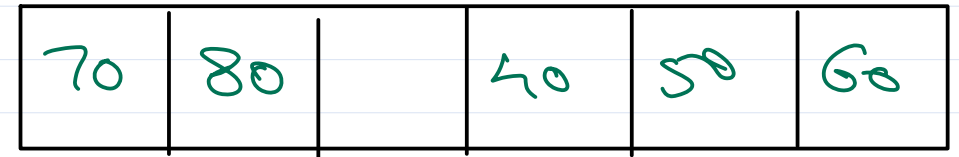
-1

f



r

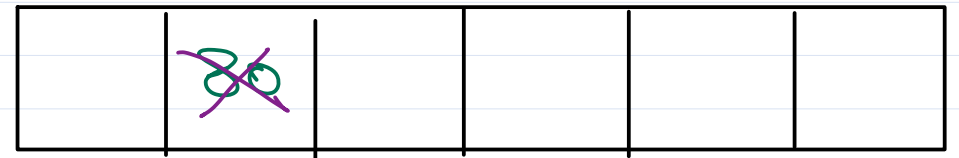
-1



f



-1



f

Circular Queue

effective queue: $f+1$ to s

①

init:

$$r = -1$$

$$f = -1$$

$$\text{cnt} = 0$$

push:

$$r = (r + 1) \% \text{max}$$

$$\text{arr}[r] = \text{val};$$

$$\text{cnt}++;$$

pop:

$$f = (f + 1) \% \text{max}$$

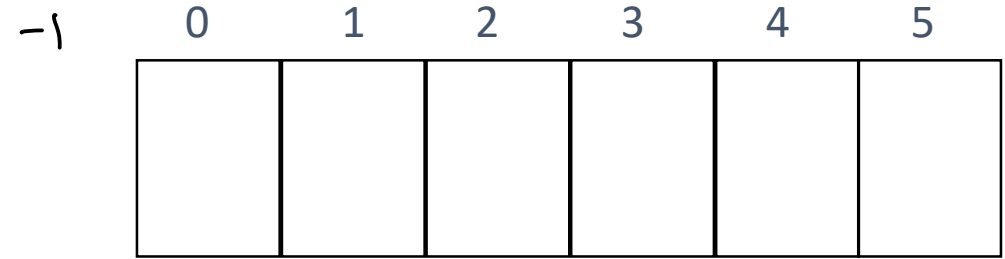
$$\text{val} = \text{arr}[f];$$

$$\text{cnt}--;$$

peek:

$$i = (f + 1) \% \text{max}$$

$$\text{val} = \text{arr}[i];$$



fr

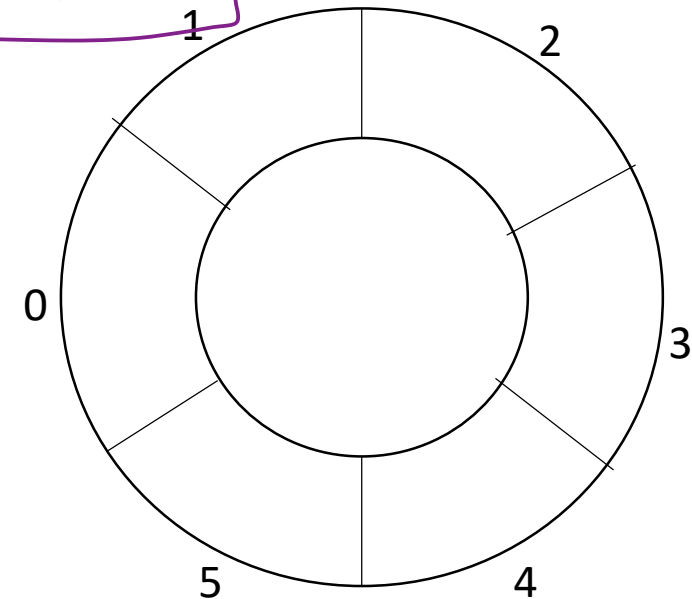
Homework

empty:

$$\text{cnt} == 0$$

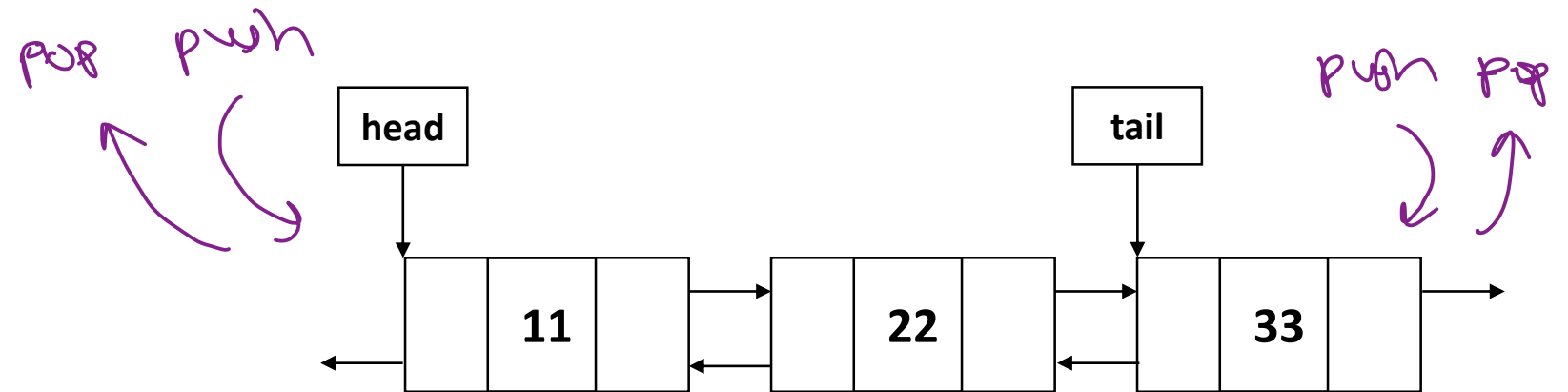
full:

$$\text{cnt} == \text{max}$$



DeQueue

- In double ended queue, values can be added or deleted from front end or rear end.



Priority queue

- In priority queue, element with highest priority is removed first.

efficiently impl
using heap DS.

→ ele with highest priority
comes out first. (internally
eles are maintained in order)
time complexity $O(1)$ X



Stack effective stack ele: 0 to top.

init:

$top = -1$

is full:

$top == max - 1$

push:

$top++;$

$arr[top] = val;$

is empty:

$top == -1$

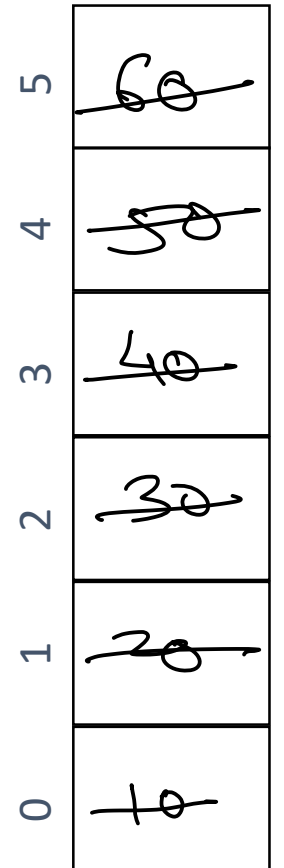
pop:

$val = arr[top]$

$top--;$

peek:

$val = arr[top]$



top → 1





Thank you!

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