



Data Structure & Algorithms

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Course Introduction

- Data Structure and Algorithms

- Data Structures: Linked list, Stack, Queue, Binary search tree, Heap, Graph.
- Algorithms: Sorting, Searching, Stack, Queue & Linked list applications, Graph algorithms.

- Course Goals

- Implement each DS & Algorithms from scratch.
- Understand complexity of algorithms.

- Course Schedules 10th

- 12th Aug 2022 to ~~12th~~ Sep 2022
- Mon-Fri: Lecture – 5:00 PM to 8:00 PM

- Resource sharing

- <https://gitlab.com/sunbeam-modular/dsa-05> ✓
Recorded videos will be available for 7 days. ✓
<http://students.sunbeamapps.org> ✓

- Course Format

- Participants are encouraged to code alongside (copy code from code-sharing utility in student portal).
- Post your queries in chat box (on logical end of each topic).
- Practice assignments will be shared.
They are optional. If any doubts, share on WA group (possibly with screenshot).
Faculty members or peers can help.

- Programming language

- DS & Algorithms are language independent.
- Classroom coding will be in Java (use IDE of your choice).
- Will share C++/Python codes at the end of session.
- Language pre-requisites ?



Course Pre-requisites

Java

→ Sunbeam
youtube channel

- ✓ • Language Funda
- ✓ • Methods
- ✓ • Class & Object
- ✓ • static members
- ✓ • Arrays
- Collections

↳ ArrayList &
Hash Map &
Stack

Python

- Language Funda
- Functions
- Class & Object
- Collections

C++

- Language Funda
- Functions
- Class & Object
- Friend class
- Arrays
- Pointers

C

- Language Funda
- Functions
- Structures
- Arrays
- Pointers



Data Structure

- Data Structure
 - Organizing data in memory
 - Processing the data
- Common data structures
 - Array
 - Linked List
 - Stack
 - Queue
 - Hash Table
- Advanced data structures
 - Tree
 - Heap
 - Graph



Data Structure

- Data Structure efficient
 - Organizing data in memory
 - Processing the data
- Common data structures
 - Array
 - Linked List
 - Stack
 - Queue
- Advanced data structures
 - Tree
 - Heap
 - Graph

- Asymptotic analysis
 - It is not exact analysis
 - Big' O notation
↳ order of.

→ depend on machine, language.
→ CPU arch

Exact Analysis

- ① how many bytes?
- ② how much time (sec)?

○ Space complexity

- Unit space to store the data (Input space) and additional space to process the data (Auxiliary space).
- $O(1)$, $O(n)$, $O(n^2)$

double arr [5];
 $S \propto n$
Input Space → $O(n)$

| | | | | |
|-----|-----|-----|-----|-----|
| 1.1 | 2.2 | 3.3 | 4.4 | 5.5 |
|-----|-----|-----|-----|-----|

Sum
i
 $S = k$
Aux Space → $O(1)$

double sum = 0.0;
for(int i = 0; i < 5; i++)
sum = sum + arr[i];

○ Time complexity

- Unit time required to complete any algorithm.
- (proportionate) • Approximate measure of time required to complete any algorithm.
- Depends on loops in the algorithm.
- $O(n^3)$, $O(n^2)$, $O(n \log n)$, $O(n)$, $O(\log n)$, $O(1)$

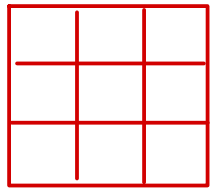


Time complexity

- Write a program to calculate factorial of given number.

$f(n)$ $x = 1;$
 $\text{for}(i=1; i \leq n; i++)$ $\text{itr} = n$
 $x = x * i$ $T \propto n$
 $\boxed{O(n)}$

- Print 2-D matrix of $n \times n$.



$\text{for}(i=0; i < n; i++) \{$
 $\text{for}(j=0; j < n; j++) \Rightarrow \text{itr} = n * n$
 $\text{print}(\text{arr}[i][j]);$ $T \propto n^2$
 $\}$ $\boxed{O(n^2)}$

- Print given number into binary.

$\text{while}(n > 0) \{$
 $\text{print}(n \% 2);$
 $n = n / 2;$
 $\}$

10 \rightarrow 5 (0)
 \rightarrow 2 (1)
 \rightarrow 1 (0)
 \rightarrow 0 (1)

31 \rightarrow 15 (1)
 7 (1)
 3 (1)
 1 (1)
 0 (1)

$\text{itr} = n$
 $\text{itr} \log 2 = \log n$
 $\text{itr} = \frac{\log n}{\log 2}$

- Print table of given number.

$\text{for}(i=1; i \leq 10; i++)$ $T = k$
 $\text{print}(n * i);$ $\boxed{O(1)}$

$T \propto \frac{\log n}{\log 2}$
 $T \propto \log n$
 $\boxed{O(\log n)}$



Linear Search

key ← ele to find

- Find a number in a list of given numbers (random order).

```
for (i = 0; i < n; i++) {  
    if (a[i] == key)  
        return i;  
}  
return -1;
```

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|----------------|----|----|----|----|----|----|----|----|
| 88 <u>m</u> | 33 | 66 | 99 | 11 | 77 | 22 | 55 | 11 |

↑

- Time complexity

- Worst case** → max num of iterations
→ finding last ele or ele not exist in array. $\text{itr} = n$
 $T \propto n$ } $O(n)$
- Best case** → min num of iterations.
→ finding the first element → $\text{itr} = 1$
 $T = k$ } $O(1)$
- Average case** → avg num of iterations
→ avg iterations = $n/2$
 $T \propto \frac{n}{2}$
 $T \propto n$ } $O(n)$



Binary Search

key = 50

```
l = 0;
r = n - 1;
while (l <= r) {
    m = (l + r) / 2;
    if (key == a[m])
        return m;
    if (key < a[m])
        r = m - 1;
    else // key > a[m]
        l = m + 1;
}
return -1;
```

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|----|----|----|----|----|----|----|----|----|
| 11 | 22 | 33 | 44 | 55 | 66 | 77 | 88 | 99 |

r
 m
 l

$$2^{\text{itr}} = n$$

$$\text{itr} \log 2 = \log n$$

$$\text{itr} = \frac{\log n}{\log 2}$$

$$T \propto \frac{\log n}{\log 2}$$

$$T \propto \log n$$

✓

$$O(\log n)$$



Recursion

- Function calling itself is called as recursive function.
- To write recursive function consider
 - Explain process/formula in terms of itself
 - Decide the end/terminating condition

- Examples:

- $n! = n * (n-1)!$

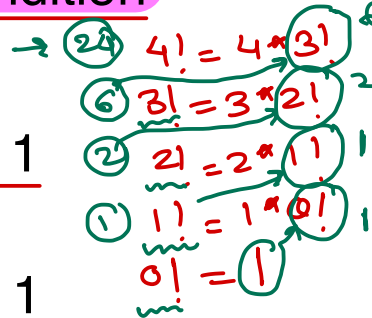
- $x^y = x * x^{y-1}$

$$\begin{aligned} 2^3 &= 2 * 2^2 \\ 2^2 &= 2 * 2^1 \\ 2^1 &= 2 * 2^0 \\ 2^0 &= 1 \end{aligned}$$

$$\underline{0! = 1}$$

$$\underline{x^0 = 1}$$

$$\underline{T_1 = T_2 = 1}$$



- On each function call, function activation record or stack frame will be created on stack.

```
int fact(int n) {
```

```
    int r;
```

```
    if(n==0)
```

```
        return 1;
```

```
    r = n * fact(n-1);
```

```
    return r;
```

```
}
```

```
res=fact(5);
```

- $T_n = T_{n-1} + T_{n-2}$
 - factors(n) = 1st prime factor of n * factors(n)



Recursion

```
int fact(int n) {  
    int r;  
    if(n == 0) ✗  
        return 1;  
    r = n * fact(n-1);  
    return r;  
}
```

③ → ② → ①

```
int fact(int n) {  
    int r;  
    if(n == 0) ✗  
        return 1;  
    r = n * fact(n-1);  
    return r;  
}
```

② → ①

```
int fact(int n) {  
    int r;  
    if(n == 0) ✗  
        return 1;  
    r = n * fact(n-1);  
    return r;  
}
```

①

```
int fact(int n) {  
    int r;  
    if(n == 0) ✓  
        return 1;  
    r = n * fact(n-1);  
    return r;  
}
```

③ → ② → ①

```
int fact(int n) {  
    int r;  
    if(n == 0) ✗  
        return 1;  
    r = n * fact(n-1);  
    return r;  
}
```

④ → ③ → ② → ①

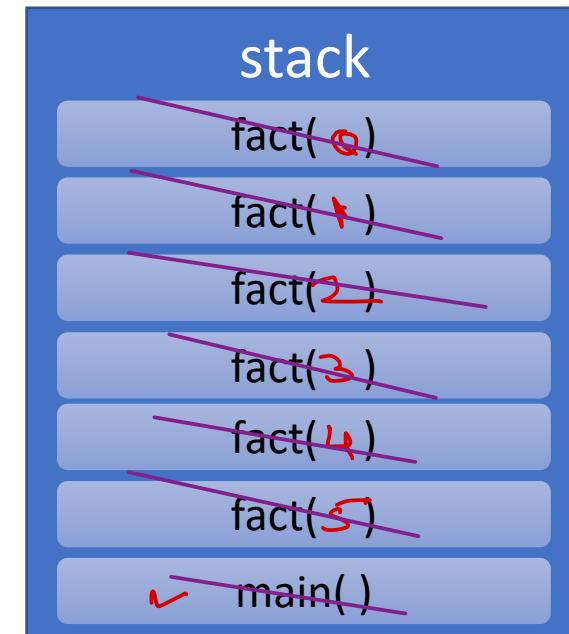
```
int fact(int n) {  
    int r;  
    if(n == 0) ✗  
        return 1;  
    r = n * fact(n-1);  
    return r;  
}
```

⑤ → ④ → ③ → ② → ①

```
int main() {  
    int res;  
    res = fact(5);  
    printf("%d", res);  
    return 0;  
}
```

⑤ → ④ → ③ → ② → ①

120





Thank you!

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