(1)

* The main objectives of generics are to provide type safty of to resolve type coesting Problems.

Case 1:- Tyle Safty: -

Arrays are Type safe that is we can give the gownenty for the type of elements presents inside arrays.

for example if ourse programing suguriments is only holds the string type of objects we can shoope string array by mistake if we are trying to add any other type of object we will get compile type Ersor.

Storeg () S = onew Storing [Joso];

S[0] = "durge";

S[1] = "Ravi";

S[2] = new Integer (10); (F. Im lompay ble Tytze

(found: Java, lary. Integer

S[2] = "Shrva";

(found: Java, lary. Integer

S[2] = "Shrva";

hence soring array can contein only string type of objects. due to this we earn give the gowranty for the type of arrays elements presents maide array hence arrays are safe to eye with respects to type that in sorens are type safe.

* But Collections are not Type safe that in we can't give the government for example of object and we choose arraylist, by mistake if we are toying to add any other type of object, we will not get any compile time error. but program may fail at yun time.

Bonylist il = new Possaylist();

digd("durge");

l. 9dd("Rowi");

d. add(new integer (20));

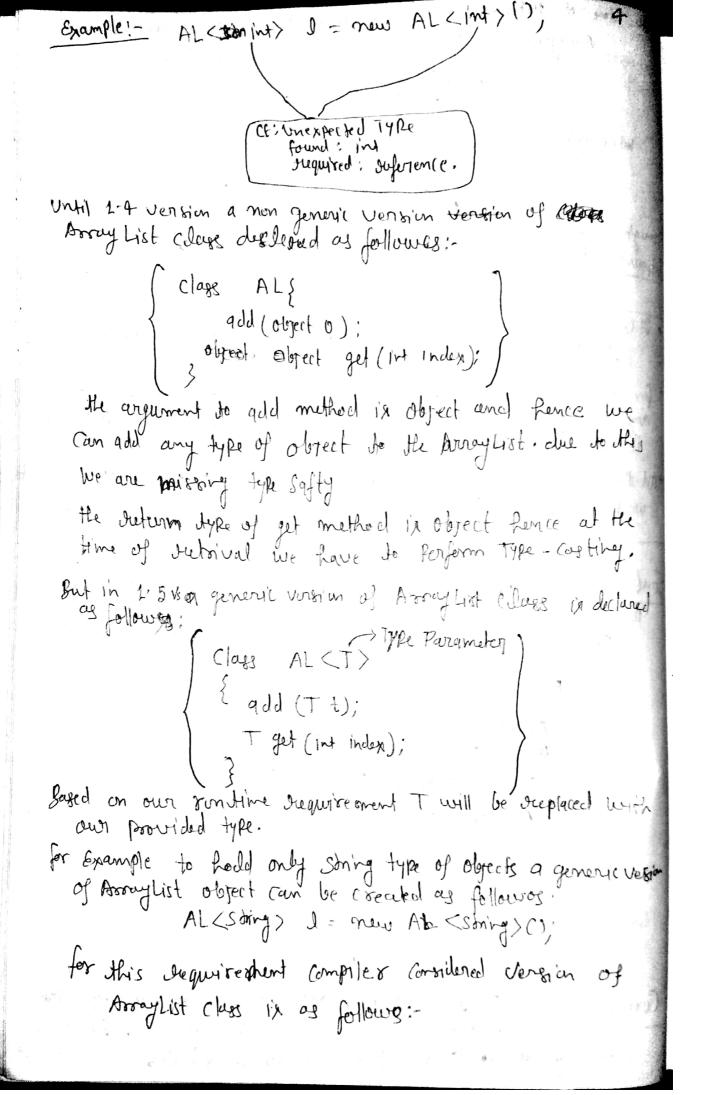
Bu fetriving object from bouray List.

Story roune! = (Story) light(0); Story name? = (Story) light(2); X RE:- (loss(ast Story name? = (Story) light(2); X RE:- (loss(ast Exception hance we can not generally for the type of Elements Asserts inside Pollection dul to this collections are not safe to use with respect to type. Heat ix collections care not Type saide.

at the time of detrival we are not sequired to Perform type Costing: 3 Boraylist (Soning> 1 = onen Braylist (Soning> () Ladd ("durger"); while substral: - Storney name 1 = , claget (0); Type-costing is not septimin Juguired Through Generic we can solve type Costing Problem.

Difference Between normal AL version & Generic AL version. ALKSoning> 1 = new ALKSoning>(); AL l= new AL(); • It is a mon generic version | • It is a generic version of AL object of AL object · fer this AL we can add · for this A.L. We can add only String type of objects & hence of in Type. any type of of object & hunce it ix not Type-safe Sife · ad the time of subsivel Compulsoy. But the stime of subsival we are not we have to perform Type-cashing orequired to perform Type-Costing Condidution 1:- Polymer phism Concept Applicable only for the Base Type But not for Parameter Type. usage of Parrents duference to hold child object ix the Concept of Polymorphism > Paramerter Type ByeTyle AL <String> il = new AL <String>(); & List (Strig) d= new AL (Strig)(); Collection (String) I = new AL (String) (); AL Colorects I = onew AL (Sonly > (); X) CE: incompatible type found; AL < String > required: AL (object) Conclusion -2:- for the type Parameter we can provide any alors or Insterface name but not Primitive. If we are trying to provide primitive then we will get CE.

Scanned by CamScanner



(close AL < son'ng>
{ add (Sonly s)
}
Sonly get(int index)
}

The argument to add method ix storing type hence we can add only storing Type of objects. by mistake if we are trying to add any other type we will get compile time cross.

l.add ("durger"); L'add ("new Integer (10));

(CE: Cannot find Symbol
Symbol; method add (J. J. Integer)
Jo(ation: class AL (String)

Lunie through generic we are getting Type Safty.
The Julian type of get method in stone and fence at
the time of Julianual we are not required to perform
type custing.

Story namel = l.get(0);

[Type Costing ix not required]

* In generic we are associating a Type parameter to the class Surk type of parameterized classes are nothing but Greneric classes or Template classes.

* Jujed on our orequirement we can define our own Generic Clarkes also:

Class Account < T >
{

Account < Chold > 91 = new Account < (nold > ();

Account < Platinum > 91 = new Account < Platinum > ().

```
Example: - Class Chen < T>
                        Tobe;
                     Cum (Tob)
                    this.ob = ob;
                    Public Void show ()
                         System.out.pointln (The tye of Util+
                                                         ob. get (does (). get Name ());
                     Public T get 06 ()
                    3 return ub.
       Class henderno
            P& V main (Shring[) angs)
{
(Grennes Shring) g1 = new Gren (Shring) ("dunga");

J1.8how(); // The type of ob; Java Leary Shring
                          soph (91.90061); // durger.
              (Gen < Integer) 32 = new hen < Integer) (10);
32. Show(); // The type of ob; Java. Loung. Integer
Soph (92. get 06 ()); // 10
             (Gren (Double) g3 = new Gren (Dobbble) (10.5);

33-Show(); // The life of ob; Java. Day. Double

Soph(33.940b()); 10.5
```

6

Bounded Types !we can bound the Type percameter for a perticular range by vertige extends Keywerd such types are Called Sounded Types. (loss TestaT> as the type parameter we can puts any Type and there are no sustriction and of is unbounded type. VTast < Integer> t1 = new Test < Integer>(); V Test < Shing> to = new Test < Shing> (); Syntex for Bounded Types !-Class Text < Textender X>] * x can be either class or anderforce. If x ix 9 class them as the type parameter we can poss either X type on its Child closses. it x ix a solarface then as the type parameter we can page Either X type or its simplementation (losses. Closs Rest (T extends Number) 1. Class Test < T externals Runnable> Test (Integer) It = new Test (Intger)(); It Test (Rumable) It = new Test (Rumable)(); Took (Thread) to = new Took (Thread) (). ()< frintes) the Test (strintes) the 1x XTest < Integer) to = new Test(Integer>(); CE: Type Parameter Java Lang-Soring) (CE: Type Parameter Java. Land Biteger 12 not within its bound in not with in it bound * we can define bounded type Even in combination also class Test (T extends Numbers of Rumable) as the type Parameter we can take anything which should te child class of Number and Should Emplements Runnable Interface

@ ml (AL (String) ! -

We can call this method by Russing Assemplist of only string type. but noith in the method we can add only string type of object to the List by mistake If we are trying to add any other type other we get compile time error.

m1 (AL <Soning> d) l. qdd ("A"); ligdd ("nuil"); 1. add (10); X

(2) ml (AL <?> d):-

We can call this method by Parshing Assaylist of every in Known type. But with in the method we can't gold anything to the List Except need null, Because we don't Know the type exectly. nuil is allowed because It is valid value for any type.

m1 (AL < ?> 1) 1 Ladd (10.5); X 1.add ("A"); X 1 add (20); X } d. add (nuil);

This type of methods are best suitable for read only Operation.

3 m1 (AL<? extends X) l):-

x can be either class or Interface. If X ix a class other we can call this method by Patring Array List of Either X type of sts child classes. If X IX an Interface other we can call this method by Passing ArrayList of Either X type or Its Implementation Classes. But with in the method we can't add anything to the List Except null. Because we don't know Type excatly. This method also best Suitable for Read only orination.

(4) ml(ALC? Super X > I) ;- x can be Either Class or contentine. If x ix a chara than we can Call this method By Besting Amonghist of wither X type OR 3ts Supper Classes of X ix can onterface other we can call this method by fassing somewhat of sither X type or Super Class of Implementation class of X.

Object Rummuble (1)

But with in the method we can add x type of object and 'null' to the list.

Knample !
AL < storing > 1 = new AL < Storing > ();

AL < ? > 1 = new AL < Integer > ();

AL < ? ? xtends Number > 1 = new AL < Integer > ();

X AL < ? extends Number > 1 = new AL < Integer > ();

X AL < ? extends Number > 1 = new AL < storing > ();

X AL < ? extends Number > 1 = new AL < storing > ();

X AL < ? sufer Storing > 1 = new AL < object > ();

X AL < ? > 1 = new AL < ? > ();

CE: unex Pecked type

found: ?

X AL < ? > 1 = new AL < ? > ();

CE: unex Pecked type

found: ?

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

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X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? > 1 = new AL < ? extends Number > ();

X AL < ? extends Number > ();

CE; unexpected type
Found; ? extends Numbers
drequired; class or Interface without bounds

we can declare Type Parameter Enther at class level of 11 method level. · Declare Type Parameter at class lovel. class Test <T> we can use'T' with in this class Based on our sugaisement. · Declaring Type Parameter at Method level. we have to declare Room Type Parameter Just before Jutorn Type. Class Test Public <T> Void m1 (T ob) We can there I'm where with in this Method bused on our Juquirement. we can define Bounded Type Even at Method level also -~ Public . < T> void m1() V Public < T etends Number > m1() : ; Public <T extends Runnable> m1(); Public <T extends Number 4 Runnable> m1 (); Public <T extends companable 4 Runnable> mL(); Noblic <T extends Numbers & Comparentes & Runnerble > m2() * public <Textends Runnable & Number> m1 (); [first we have to take class and]
then interface X Public <7 extends Number 4 thread > m1(); [we can't Extend more than one class]

If we send General Object to non General arrea the Its stant behaving like non Generu's object. Similarly of we send non Cremeric object to Gray Area Hen gts start behaving like Generic Object that is the Jocation on which object present Based his on that Belowiour will be defined.

Class Test Public State word main (Storty [] angs) Array List (String) . I = new Array List (Shigh) I add ("durger"); I add ("Rowi"): ladd (10) // CF m1(1); soph(I); / (durger, Rovi, 10, 20, 5, true) } l.add (20.5), -> CE Public S VML (Array List . I) liadd (10.5); } Non generic Area }.d.add (tom);

Conclusion; - The main purpose of Generics 13 To Provide Type Safty and To resolve Type Cashing Problem. Type Safty 4 Type Casting both are applied

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mal

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at compile time hance cremeric concept also applicable 13
  only at compile Fine but not at Rendine.
  At the time of compilation and as lost step Genic
  Combex will be sumove and hence for the JUM
 Crereric Syntex will not be quailable.
 hence the following declaration are equals:
AL J = new AL (String) ();

AL J = new AL (Integor) ();

AL J = new AL (Double) ();

AL J = new AL();
        AL I= new AL (Sting) ();
            1. add(10);
             J. 999 (10.2);
              1.add (tous):
               sopin (d); [10, 10.5, toue.
     The following Declarations are equals
   equals { AL<Satisfy ) d=new AL<Saring>();

AL<Saring> d=new AL();
      for they promylist abject we can only string type
       of object.
   · closs Test
        Public void one (ALCSmirg) &) => me (AL d) {
        Mic void m1 (AL < Sirtyen > 2) => month f. m1 (AL I);
```

Same Erog Sure.

De Compile Code Normally by Constiderity guranic

Come et Syndex.

2 Remove Cherenic Syntoxx;

3) compile once again regultant Code.