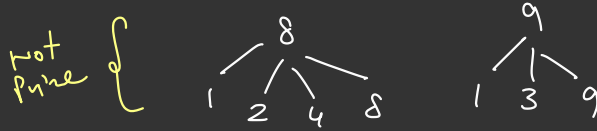


Prime Numbers

Prime \rightarrow exactly 2 divisors $\rightarrow 1 \& N$



Prime Check

$O(N)$

I) $1 \xrightarrow{\text{Count}} N$ div

```

    cnt = 0
    for (i = 1; i <= N; i++) {
        if (N % i == 0) {
            cnt++;
        }
    }
  
```

II)

```

    if (cnt == 2) {
        Prime
    }
  
```

$(x, 2, 3, \dots, N-1, x)$

as soon as you atleast 1 div in the range $(2, N-1)$

break

Algo-2

$O(N)$ in worst case

$N=7$

7×2 NO
 7×3 NO
 7×4 NO
 7×5 NO
 7×6 NO
 7×7 YES

for $(i=2, i \leq N-1, i++)$

if $(N \% i == 0)$ {
 break;

3

loop ends
 $i = N$

if $(i == N)$ {
 prime

else {
 3

not prime

break the loop
 $i = N-1$

3

$2, 3, \dots, 8$
 x

$2, 3, 4, 5, 6, 7$
 $i++$

$2, 3, \dots, 8$
 9

$N=9$

9×2 NO
 9×3 YES Stop

faster for non-prime N's

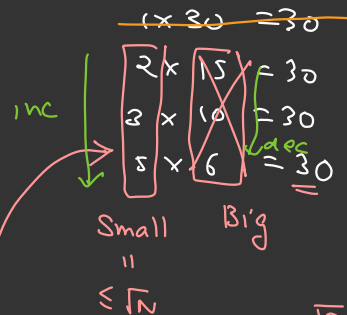
8

Algo -3 Root N optimisation



Examples -

$N=30$ ~~1~~, 2, 3, 5, 6, 10, 15 ~~30~~



Not Prime \Rightarrow any 1 divisor

$$a = b$$

$$\Rightarrow a^2 = N$$

$$a = \sqrt{N}$$

$$\boxed{a} \cdot \boxed{b} = N$$

$$> \sqrt{N} \quad > \sqrt{N}$$

$$= > N$$

$$\Rightarrow \begin{matrix} \boxed{a} \times b = N \\ 1 \times b = N \\ 2 \times b = N \\ \vdots \\ \sqrt{N} \end{matrix}$$

$$36 = 1, 2, 3, 4, 6, 9, 12, 18, 36$$

$\leq \sqrt{N}$



$$1 \times 36 = 36$$

$$2 \times 18 = 36$$

$$3 \times 12 = 36$$

$$4 \times 9 = 36$$

$$6 \times 6 = 36$$

Pairs

$$N = 10^4$$

$$\sqrt{N} = 100 \text{ steps}$$

$$N = 10^6$$

$$\sqrt{N} = 1000 \text{ steps}$$

is Prime (n) {

$$i * i \leq N$$

$$\Rightarrow i^2 \leq N$$

$$O(\sqrt{N})$$

is Prime = true;

for (i = 2; i * i ≤ N, i++) {

if (N % i == 0) {

is Prime = false;

break;

}

}

if (is Prime) → "Prime"

else → not Prime

$$\underline{\underline{O(1)}}$$

Problems involving \Rightarrow Range of Prime Numbers

Q) Generate all Primes in Range 1 to N.

N = 100

2, 3, 5, 7, 11, 97

Brute force

$O(N \sqrt{N})$

[

3

$\xrightarrow{\quad}$
for (i = 1, i <= N; i++) {
if (isPrime(i)) { print(i) };
 \uparrow
 \sqrt{N}

$\xrightarrow{\quad}$
 \sqrt{N}

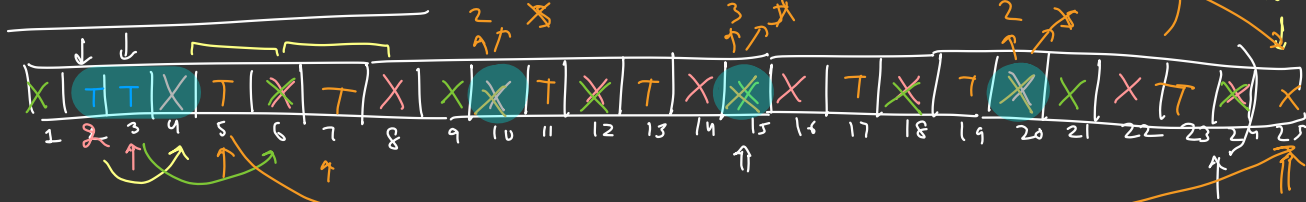
Prime Sieve - Sieve of Eratosthenes.

(famous technique)

Init \Rightarrow "all prime."

$\{2, 3, 7, 11, 13, 17, 19, 23\}$

boolean
primes



Prime (Yes)

2 \rightarrow 4

3 \rightarrow 6, 9, 12

4 \rightarrow 8, 12, 16

$i=2$

$i=3$

$i=4$

for ($i=2, i \leq \sqrt{N}; i++$) {

if (primes[i] = true) {

for ($j=2i; j \leq N, j=j+i$) {

primes[j] = false,

3

for
prime
no's

}

}

Mark all multiples of
2 as Not Prime

4, 6, 8, 10, --
↑ ↑ ↑

3 \rightarrow 15

3 \rightarrow 6, 9, 12, 15, --

5 \rightarrow 15

5 \rightarrow 10, 15, 20, 25, --

optimise it

$j = i^2$

work =

$$\frac{N}{2} + \frac{N}{3} + \frac{N}{5} + \dots + \frac{N}{p} + \sqrt{N}$$

2 --- N

3, 6, 9, 12, 15, 18, 21, 24

$= O(N \log \log N)$

largest prime $\leq \sqrt{N}$

Print

for ($i=2$ --- N) {

if (isPrime(i)) print i

}

$i = 2$
 $i = 3 \rightarrow$
 $i = 4$
 $i = 5 \rightarrow$

6, 9, 12, 15, 18, 21, 24
 10, 15, 20, 24

$\approx O(N)$ \rightarrow very small
 Practically \approx

25

first multiple i will work

$j = 2$; $j \leq n$; $j = j + 1$
 $i = 2$

N

$2 \dots \sqrt{25}$
 $= 2 \dots 5$

$2 \dots \sqrt{24}$
 $= 2, 3, 4$

primes[N+1] = true,

for (i = 2, i <= \sqrt{N} , i++) {

if (isPrime[i]) {

→ for (j = i²; j <= N; j = j + i) {
isPrime[j] = false;
}

}

for (i = 2 ———— \sqrt{N}) {

if (isPrime[i]) { print(i); }

}

(Q)

given Q queries, of the form $[a, b]$ find out the
Count of Primes in range $[a, b]$. $a, b \leq 10^6$

optimised

Algo \rightarrow

$Q = 5$

$a = 5$

$[3, 10]$



$[10, 100]$

$[5, 20]$

$[2000, 10000]$

$[1500, 3500]$

Overall

$N \log \log N + Q$

(done)

① Precompute sieve + prefix array $N \log \log N + N$



② iterate over ~~3 to 10~~ and count (b-a) ✓

for
each
query.

prefix[i] = count of primes till i $\rightarrow O(1)$ per query
 $Q(a, b) = \text{prefix}[b] - \text{prefix}[a-1]$

Prime Factorisation



10.25

↳ No = Product of Powers of Primes.

$$48 = 2^4 \cdot 3^1$$

$$51 = 3 \cdot 17$$

$$20 = 2, 2, 5$$

2	48
2	24
2	12
2	6
3	3
	1

Algorithm - 1

```
for (i = 2; i <= n; i++) {
```

```
    if (n % i == 0) {
```

```
        cnt = 0
```

```
        while (n % i == 0) {
```

```
            n = n / i
```

```
            cnt = cnt + 1;
```

```
        }
        print(i, cnt);
```

```
    }
```

$$2 \cdot 3^3$$

2	54
3	27
3	9
3	3
	1

N

(2, 1) (3, 3)

N

i = 2 27 54 yes

i = 3 1 27 cnt = 1

yes cnt = 0

9 cnt = 1

3 cnt = 2

1 cnt = 3

i = 4 1



Worst case

no is prime

→ O(N)

4 Steps.

N = 11

2x
3x
4x
5x
6x
7x

8x
9x
10x
11x

$\frac{11}{11} = 1$

Algo-2

optimised

Not True \sqrt{N} WHY?

$$4 \leq \sqrt{11}$$

for ($i=2$; $i \leq n$; $i++$) {

if ($n \% i == 0$) {
 cnt = 0
 while ($n \% i == 0$) {
 n = n / i
 cnt = cnt + 1;
 }
 print (i, cnt);

if ($n > 1$) { print (n, 1); }

outside
for
loop

$$N \leq \sqrt{N}$$

N=44

$$2 \cdot 11 = 4 \times 11 = 44$$

i	n	cnt
2	44	0
2	22	1
2	11	2
3	11	-
4	11	-
5	11	-
6	11	-
7	11	-
8	11	-
9	11	-
10	11	-

must be prime no as we couldn't find any more div \sqrt{N}

↑

$$\frac{11}{1} \leq \frac{11}{1}$$

$$\text{cnt} = 0$$

$$\text{cnt} = 1$$

N = 10

⇒

	n	cnt
2	10	0
2	5	1

↓

√5
"
2.2

3 < √5 → stop

```

if (n > 1) {
    print(n)
}
    
```

2 × (5)

2	16
8	1
4	2
2	3
N =	1
	4

→ (2, 4)

$$5^2 \leq 55 \Rightarrow 25 \leq 55$$

$$6^2 \leq 36 \leq 11$$

$$\boxed{6^2 \leq 11}$$

for (i=2, $i^2 \leq 11$; $i++$) {
Stop

3

if (n > 1) {
print(1, 1);

3

i=2

i=3

i=4

i=5

~~i=6~~

i=6

N = 110

110

55

55

55

55

11

→ 11

↑
Prime NO

cnt=0

cnt=1

x

x

cnt=0

cnt=1

x

No

(2, 5, 11)

= 2 x 5 x 11

= 110

$\sqrt{11} = 3$

JN → single

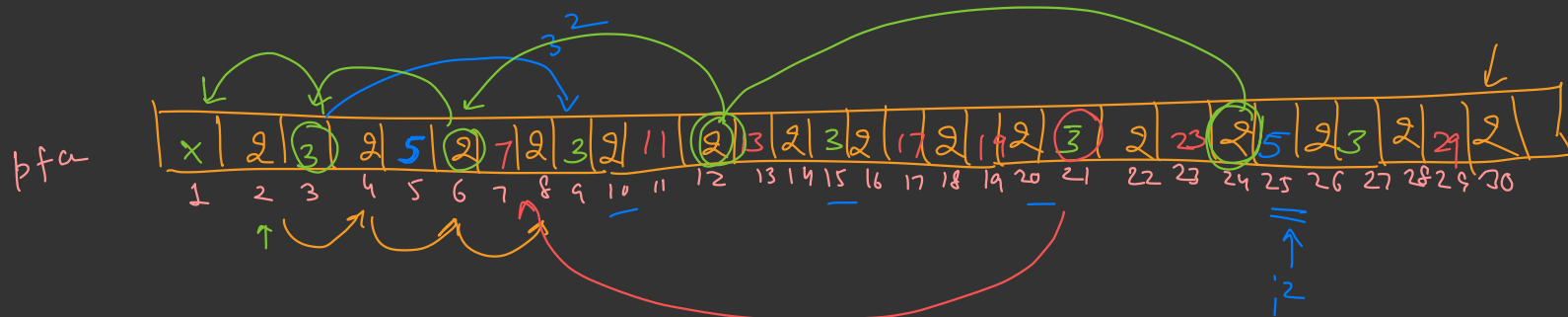
NJN → n inputs

==

Prime Factorisation for a Range

for any 'no'
store the smallest p.f

N = 30



① Build a sieve storing the smallest pf at a array?

② Prime factorisation

```

no = input()
while (no > 1) {
    print (pfa[no]),
    no = no / pfa[no];
}
    
```

3

N = 24

↓ / 2

12 / 2

↓

6 / 2 → 3 / 3 → 1 stop

21 → 3

1

$2^3 \cdot 3 = 24$

$$(3+1)(1+1) = 4 \times 2 = 8 \text{ ways}$$

$$\underline{56} = \rightarrow \underline{2^{(3)} \cdot 7^1} \quad (\text{Prime factorisation})$$



$$5 \times 3 \times 2$$

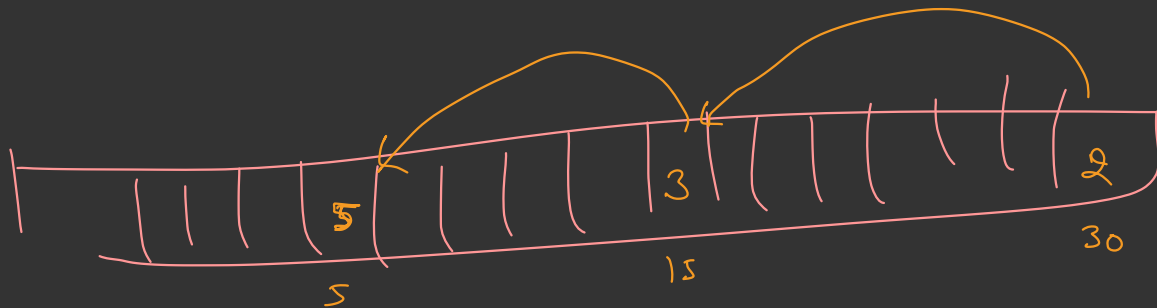
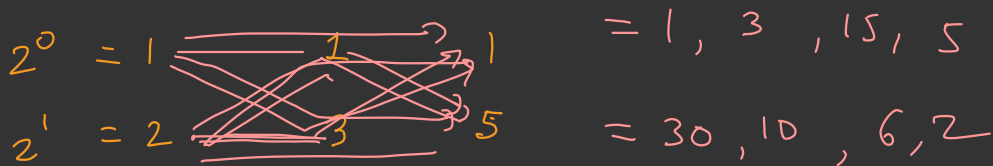
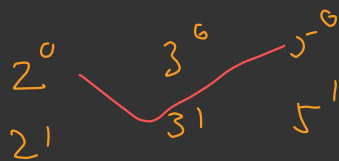
$$30 = 2^1 \cdot \underline{3^1} \cdot 5^1$$

$$= (1+\underline{1})(1+\underline{1})(1+\underline{1})$$

$$= 8$$

$$\begin{array}{l} 2^0 \\ 2^1 \\ 2^2 \\ 2^3 \end{array} \quad \begin{array}{l} 7^0 \\ 7^1 \end{array}$$

$$\begin{array}{l} = 1 \\ = 7 \\ = 2 \\ = 14 \\ = 4 \\ = 28 \\ = 8 \\ = 56 \end{array} \quad \left. \vphantom{\begin{array}{l} = 1 \\ = 7 \\ = 2 \\ = 14 \\ = 4 \\ = 28 \\ = 8 \\ = 56 \end{array}} \right\}$$



$$30 = 2 \times 3 \times 5$$