

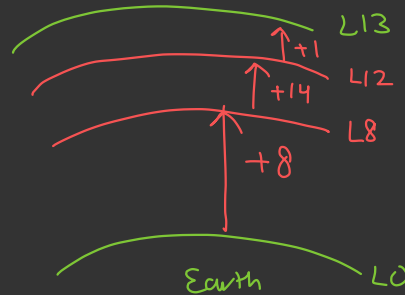
Bitmasking

warm up

There is an enemy at N levels above the surface of the earth. You are a superhero, standing on earth (level 0) and you can take jumps in powers of 2. Min jumps that would be needed to reach the enemy.

1, 2, 4, 8, 16, 32, ...

$N = 13$

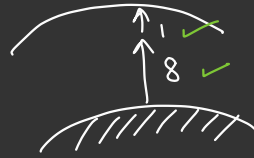
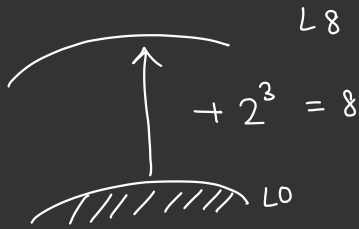


$$\begin{aligned} 13 &= 1 \cdot 2^3 + 1 \cdot 2^2 + 0 \cdot 2^1 + 1 \cdot 2^0 \\ &= 8 + 4 + 1 \\ &= \boxed{13} \rightarrow \text{Always be the minimum} \end{aligned}$$

Logic



Count the
no of set
bits in \boxed{N} (Binary Rep)



$$9 = \boxed{1001}$$

$8 + 1 = 9$

Set bits in bin rep in N

$$N = 13$$

$$\begin{array}{r} 110\cancel{1} \\ 2 \quad \quad \quad 1 \\ \hline \quad \quad \quad 1 \end{array}$$

$$\begin{array}{r} 11\cancel{0} \\ 2 \quad 1 \\ \hline + \quad 0 \end{array}$$

$$\begin{array}{r} 1\cancel{1} \quad 1 \\ 2 \quad 1 \quad 2 \quad 1 \\ \hline + \quad 1 \quad + \quad 1 \end{array}$$

$$= \textcircled{3} \text{ jumps}$$

$N \& 1 \rightarrow$ gives me the value of last bit in N

$$N = \boxed{6}$$

$$6 \& \underline{1} \rightarrow 0$$

$$\begin{array}{r} 0000000011\textcircled{0} \\ \& 0000000001 \\ \hline 00000000\boxed{0} \end{array}$$

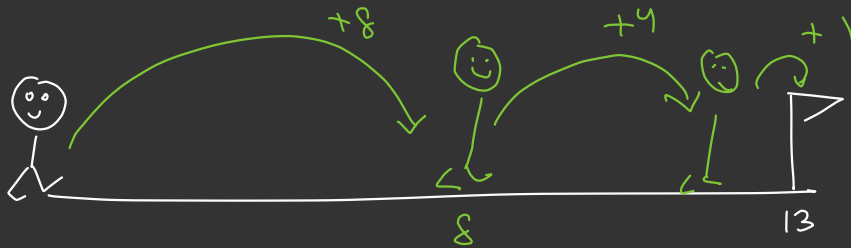
$\textcircled{0}$

$$N = 9$$

$$N \& 1$$

$$\begin{array}{r} 0000100\textcircled{1} \\ \& 000000001 \\ \hline 000000\boxed{1} \end{array}$$

$\textcircled{1}$



Common Operations involving Bitwise Operators

```
int getIth bit (no, i) {
    return (no >> i) & 1;
}
```

```
int getIth bit (no, i) {
    mask = 1 << i
    if (no & mask > 0) {
        return 1;
    }
    else {
        return 0;
    }
}
```

$N = 13$

$i = 2$ ^{ith bit}

				2	2	1	0
	0	0	0	1	0	0	1
mask →	0	0	0	0	1	0	0
				0000 1 00		> 0	

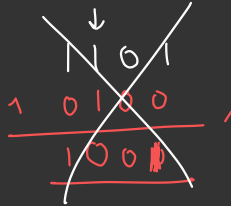
i=3	i=2	i=1	i=0
←			
1	1	0	0
xor 1			
1			

Set I^{th} Bit (N, i) {

$\text{mask} = 1 \ll i$

return $N \mid \text{mask}$

}



Toggle Bit

$1 \rightarrow 0$
 $0 \rightarrow 1$

} XOR with mask

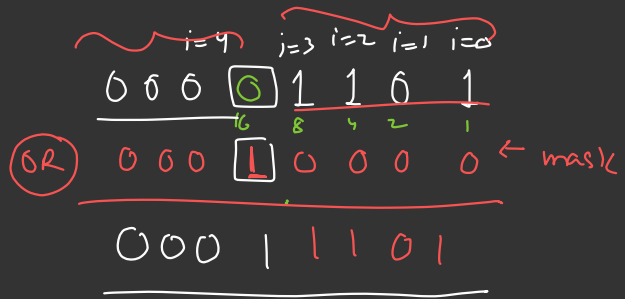
$0 \wedge 1 = 1$
 $1 \wedge 1 = 0$

Toggle I^{th} Bit (N, i) {

$\text{mask} = 1 \ll i$

return $N \wedge \text{mask}$

}



0 \rightarrow clearing bit & 0

1 \rightarrow setting bit 1

$13 + 16 \Rightarrow 29$

$N + (1 \ll i)$

ToDo:

Update i^{th} Bit (N, i, v) {

step1 [$\text{mask} = \sim(1 \ll i);$
 $N = N \& \text{mask};$

step2 $N = N | (v \ll i)$] $v=0 \rightarrow \text{No change}$

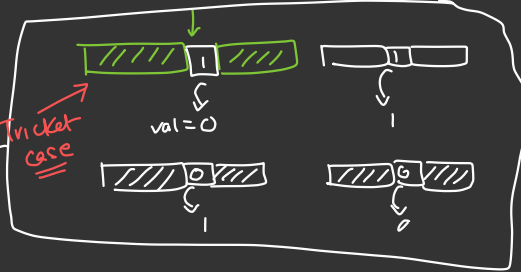
return N;

3

update (13, 2, 0)

update (13, 1, 1)

$1 \ll i$
 $10000 \quad 00000$
 $v \rightarrow 0 \text{ or } 1$



i 3 2 1 0
 1 1 0 1 $\rightarrow 1001 = 9$
 1 1 0 1 $\rightarrow 1111 = 15$

Step-1 make i^{th} bit as 0 (Clear i^{th} Bit)

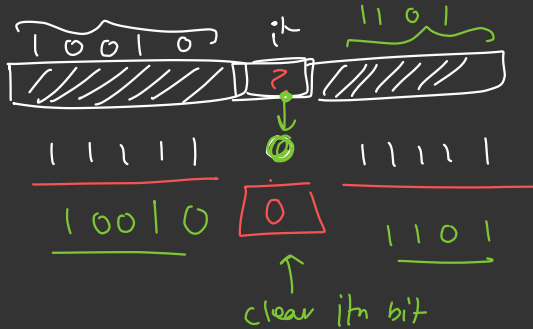
clear Bit i (N, i) {

N $\sim(1 \ll i)$
 10000
Step-2 Set the val at loc i .

N $v \ll i$
 11110111
 0001000
 11111111
 00000000
 11111111

Step-1

(8)
SR



clear i-th bit

$$\text{mask} = \sim(1 \ll i)$$

$$N = N \& \text{mask};$$

Challenge

Given a No, clear last i bits?

$$1101 \xrightarrow{i \text{ bits}} 1100$$

$$N = 13$$

$$i = 2$$

$$\begin{array}{r} 3 \ 2 \ 1 \ 0 \\ 1101 \\ \downarrow \\ 1100 \\ (12) \end{array}$$

great

$$(N \gg i) \ll i$$

Retain

$$110110110$$

$$\begin{array}{r} \& \\ 1111110000 \\ \hline 1101100000 \end{array}$$

$$\text{mask} = (\sim 1 \ll i)$$

$$N = N \& \text{mask}$$

Another way

2's Complement

$-1 = \boxed{11111111}$
 $-1 \ll 1 = \cancel{11111111}0000$

Flip all bits + 1

$$\begin{array}{r} \Rightarrow \quad 00000000 \\ + \quad 1 \\ \hline 00000001 \end{array}$$

Clear Bits In Range (N, i, j)

$a = -1 \ll (j+1)$

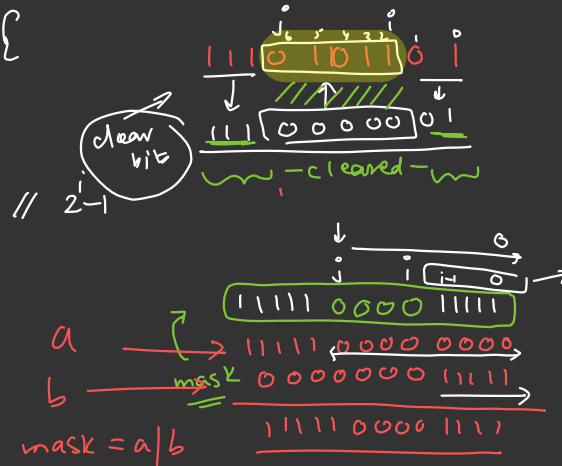
$b = (1 \ll i) - 1$

$mask = a | b$

$n = n \& mask$

return n

}

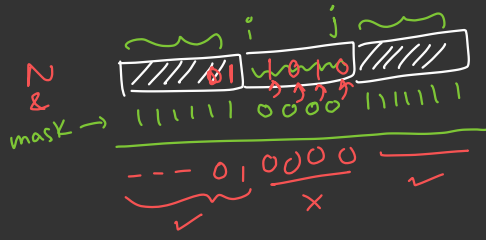


$1 \ll i \Rightarrow 2^i$

- 1 → 01
- 2 → 10
- 3 → 11
- 4 → 100
- 5 → 101
- 6 → 110
- 7 → 111
- 8 → 1000

$$2^2 - 1 = 3 = \boxed{11}$$

$$2^3 - 1 = \boxed{111}$$



Δ1 → Retain Bit
 &0 → Clear Bit

- 15 → 1111
- 16 → 10000



$$2^4 - 1 = \boxed{1111}$$

$$2^4 - 1 \Rightarrow 16 - 1 \Rightarrow \boxed{15}$$

$$\begin{aligned} k &< 4 \rightarrow 1 \\ &= 10000 - 1 \\ &= 01111 \end{aligned}$$

$$31 = \boxed{\begin{matrix} 16+8+4+2+1 \\ 11111 \end{matrix}}$$

$$2^5 - 1$$



$$2^k - 1 \Rightarrow \boxed{\begin{matrix} k \\ 00001111 \end{matrix}}$$

Q

Given $2N + 2$ numbers where every no is
coming twice except 2 unique no
find out the unique no's Bitwise

15 Min
Break



✓ 2, ✓ 5, ✓ 7, ✓ 7, (9), 5, (3), 2 ✓

once

9, 3

Welcome
back

$$XOR = 2 \cancel{15} \cancel{17} \cancel{17} \wedge 9 \cancel{15} \cancel{13} \wedge 2$$

↓
res = 9, 3 // (16)

↓
will atleast
one set
bit

→ find out the
index of
the any
set bit

a 1 1 0 1
b ^ 1 0 1 1

0 1 1 0

Any set bit → must be present only in one of the no's

Tricky

$$\begin{array}{r} 1001 \\ 0011 \\ \hline 1010 \end{array} \leftarrow$$

= (16)

i=3 i=2 i=1 i=0

$$\begin{array}{|c|c|c|c|} \hline 1 & 0 & 1 & 0 \\ \hline \end{array}$$

idx

→ sorting

$\overline{10}$ $\overline{101}$ $\overline{111}$ $\overline{111}$ $\overline{1001}$ $\overline{101}$ $\overline{11}$ $\overline{10}$
 2, 5, 7, 7, 9, 5, 3, 2

$2, 2, \textcircled{3}, 5, 5, 7, 7, \textcircled{9}$
 $\xrightarrow{\quad} \xrightarrow{\quad} \xrightarrow{\quad} \xrightarrow{\quad}$
 $\underline{\underline{N \log N}}$

have set bit at idx
 don't have a set bit
 XOR
 A → ~~2~~, ~~7~~, ~~7~~, $\textcircled{3}$, ~~2~~
 B → ~~5~~, ~~9~~, $\textcircled{9}$, ~~5~~

$2N+1$ $\textcircled{3}$
 $2N+1$ $\textcircled{9}$

$A = 0$
 $A \wedge \cancel{2} \wedge \cancel{7} \wedge \cancel{7} \wedge \cancel{3} \wedge \cancel{2}$
 $0 \wedge 3 = \textcircled{3}$

ARR[N] = { 2, 5, 7, 7, 9, 5, 3, 2 }

Step 1
 ↓
 $O(N)$

res = 0
 for (i=0; i < N; i++) {

$0 \wedge \underline{2} \wedge 5 \wedge 7 \wedge 9 \wedge 5 \wedge 3$
 $\wedge 2$

res = res \wedge arr[i];

$O(N)$

for (i=0, i < n; i++) {

if (getIthBit(arr[i], pos) == 1) { }

→ $A = A \wedge arr[i];$

}

}

$B = res \wedge A$

$res = \cancel{UN1} \wedge \boxed{UN2}$
 $\wedge \cancel{UN1}$

print(A, B)

overall time → $O(N)$
space → $O(1)$ } Best
soln

na
2

N = 12

pos = 0

N & 1 → 1 stop

while ((N & 1) == 0) {
pos = pos + 1
N = N >> 1

→

print(pos), // (2)

pos = 2
pos = 1 pos = 1
1 1 0 0
0 0 0 0
0 0 0 0

pos = 2 ✓
1 1 0 0

Day - 1 Day - 2 Day - 3 Day - 4 →

1 unit (2) (4) 8

✓
Break & make
+ Practice

A = 14 units →

(3 Days)

Power of (2)

8 + 4 + 2

1	1	1	0
8	4	2	1
✓	✓	✓	

Count
set
Bits

Reverse bits of 32 bit ^{32 bit} unsigned int A

↑
C++/C#

→ Java

DIY

0 1 1 0 1 0 1 0 1 0 1 0 1 0 1 0

1 2 3 4

```
while (n > 0) {
    bit = n & 1
    ans = ans << 1 + bit
    n = n >> 1
}
```

Ans = 0 << 1 + 1
 = 0 1 << 1 + 1
 = 0 1 0 + 1
 = 0 1 1
 = 0 1 1 0

