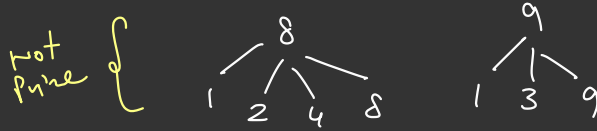


# Prime Numbers

Prime  $\rightarrow$  exactly 2 divisors  $\rightarrow 1 \& N$



Prime Check

$O(N)$

I)  $1 \xrightarrow{\text{Count}} N$   $\text{div}$

```

    cnt = 0
    for (i = 1; i <= N; i++) {
        if (N % i == 0) {
            cnt++;
        }
    }
  
```

II)

```

    if (cnt == 2) {
        Prime
    }
  
```

$(x, 2, 3, \dots, N-1, x)$

as soon as you atleast 1 div in the range  $(2, N-1)$

break

Algo-2

$O(N)$  in worst case

$N=7$

$7 \times 2$  NO  
 $7 \times 3$  NO  
 $7 \times 4$  NO  
 $7 \times 5$  NO  
 $7 \times 6$  NO  
 $7 \times 7$  YES

for  $(i=2, i \leq N-1, i++)$

if  $(N \% i == 0)$  {  
 break;  
 }

$i=3$   
 if  $(i == n)$  {  
 prime  
 }

loop ends  
 $i = n$

else {  
 not prime  
 break the loop  
 $i = n-1$

$2, 3, \dots, 8$   
 $x$

$2, 3, 4, 5, 6, 7$   
 $i++$

$2, 3, \dots, 8, 9$

$N=9$

$9 \times 2$  NO  
 $9 \times 3$  YES Stop

faster for non-prime NOs

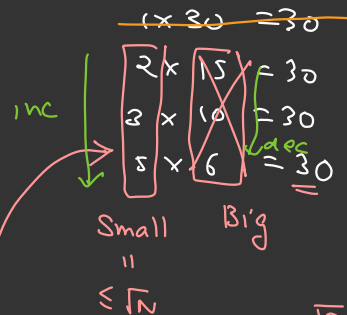
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# Algo -3      Root N optimisation



Examples -

$N=30$       ~~1~~, 2, 3, 5, 6, 10, 15      ~~30~~



Not Prime → any 1 divisor

Yes  
 $\sqrt{30} = 5.5 \times \times$

$$a = b$$

$$\Rightarrow a^2 = N$$

$$\Rightarrow a = \sqrt{N}$$

$$\boxed{a} \cdot \boxed{b} = N$$

$$> \sqrt{N} \quad > \sqrt{N}$$

$$= > N$$

$$\Rightarrow \begin{matrix} a \times b = N \\ 1 \times b = N \\ 2 \times b = N \\ \vdots \\ \sqrt{N} \end{matrix}$$

$$36 = 1, 2, 3, 4, 6, 9, 12, 18, 36$$

$\leq \sqrt{N}$



$$1 \times 36 = 36$$

$$2 \times 18 = 36$$

$$3 \times 12 = 36$$

$$4 \times 9 = 36$$

$$6 \times 6 = 36$$

Pairs

$$N = 10^4$$

$$\sqrt{N} = 100 \text{ steps}$$

$$N = 10^6$$

$$\sqrt{N} = 1000 \text{ steps}$$

isPrime(n) {

$$i * i \leq N$$

$$\Rightarrow i^2 \leq N$$

$$O(\sqrt{N})$$

isPrime = true;

for (i = 2; i \* i ≤ N, i++) {

if (N % i == 0) {

isPrime = false;

break;

}

}

if (isPrime) → "Prime"

else → not Prime

$$\underline{\underline{O(1)}}$$



Problems involving  $\Rightarrow$  Range of Prime Numbers

Q) Generate all Primes in Range 1 to N.

N = 100

2, 3, 5, 7, 11, . . . . . 97

Brute force

$O(N \sqrt{N})$

[

3

$\xrightarrow{\quad}$   
for (i = 1, i <= N; i++) {  
if (isPrime(i)) { print(i) };  
           $\uparrow$   
           $\sqrt{N}$

$\xrightarrow{\quad}$   
 $\sqrt{N}$

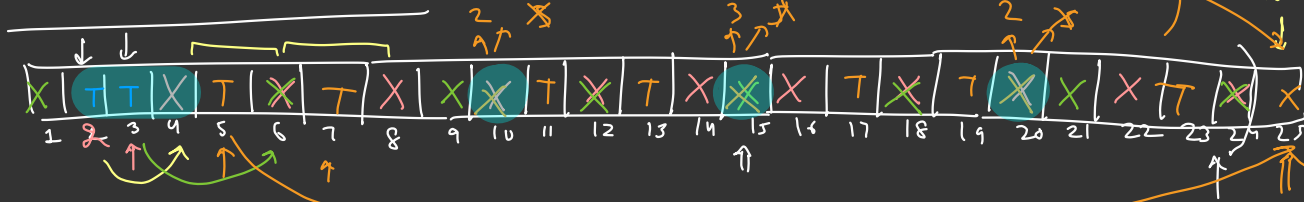
# Prime Sieve - Sieve of Eratosthenes.

(famous technique)

Init  $\Rightarrow$  "all Prime."

$\{2, 3, 7, 11, 13, 17, 19, 23\}$

boolean  
primes



2  $\rightarrow$  4

3  $\rightarrow$  6, 9, 12

4  $\rightarrow$  8, 12, 16

$i=2$

$i=3$

$i=4$

for ( $i=2, i \leq \sqrt{N}; i++$ ) {

if (primes[i] = true) {

for ( $j=2i; j \leq N, j=j+i$ ) {

primes[j] = false,

3

for  
prime  
no's

}

}

optimise it

$j = i^2$

Mark all multiples of  
2 as Not Prime

4, 6, 8, 10, --  
↑ ↑ ↑

3  $\rightarrow$  15

3  $\rightarrow$  6, 9, 12, 15, --

5  $\rightarrow$  15

5  $\rightarrow$  10, 15, 20, 25, --

Print-  
for ( $i=2 \text{ --- } N$ ) {  
if (isPrime(i)) print i  
}

2 --- N

3, 6, 9, 12, 15, 18, 21, 24, 27, 30, 33, 36, 39, 42, 45, 48, 51, 54, 57, 60, 63, 66, 69, 72, 75, 78, 81, 84, 87, 90, 93, 96, 99

$$\text{Work} = \left( \frac{N}{2} + \frac{N}{3} + \frac{N}{5} + \dots + \frac{N}{p} \right) + \sqrt{N}$$

$= O(N \log \log N)$  largest prime  $\leq \sqrt{N}$

$i = 2$   
 $i = 3 \rightarrow$   
 $i = 4$   
 $i = 5 \rightarrow$

6, 9, 12, 15, 18, 21, 24  
 10, 15, 20, 24

$\approx O(N)$   $\rightarrow$  very small  
 Practically  $\approx$

25

first multiple i will work

$j = 2$  ;  $j \leq n$  ;  $j = j + 1$   
 $i = 2$

$N$   
 $\uparrow$   
 $2 \dots \sqrt{25}$   
 $= 2 \dots 5$

$2 \dots \sqrt{4}$   
 $= 2, 3, 4$



primes[N+1] = true,

for (i = 2, i <=  $\sqrt{N}$ , i++) {

if (isPrime[i]) {

→ for (j = i<sup>2</sup>; j <= N; j = j + i) {  
isPrime[j] = false;  
}

}

for (i = 2 ————  $\sqrt{N}$ ) {

if (isPrime[i]) { print(i); }

}

Q

given Q queries, of the form  $[a, b]$  find out the  
Count of Primes in range  $[a, b]$ .  $a, b \leq 10^6$

optimised

Algo  $\rightarrow$

$Q = 5$

$a = 5$

$[3, 10]$

$\Rightarrow$

$[10, 100]$

$[5, 20]$

$[2000, 10000]$

$[1500, 3500]$

Overall

$N \log \log N + Q$

(done)

① Precompute sieve + prefix array  $N \log \log N + N$



## Prime Factorisation



10.25

↳ No = Product of Powers of Primes.

$$48 = 2^4 \cdot 3^1$$

$$51 = 3 \cdot 17$$

$$20 = 2, 2, 5$$

2	48
2	24
2	12
2	6
3	3
	1

## Algorithm - 1

~~for~~ (i = 2; i <= n; i++) {
$$\text{if } (n \% i == 0) \{$$
$$\text{cnt} = 0$$

```

    cnt = 0
    while (n % i == 0) {

```

$$h = n/i$$
$$\text{cnt} = \text{cnt} + 1;$$

3

```
print(i, cnt);
```

3

worst case

no is prime

→  $O(n)$

$$k=11$$

2x

 $3 \times$ 

9 x

5x

6x

 $\gamma_x$ 
$$f(x)$$

9x

10x

11 ✓

4 steps.

$$\frac{11}{11} = 1$$

**N**

$$i = 2$$

27

54

yes

 $i = 3$ 

1

27

cnt = 1

yes

$$r_H t = 0$$

Q.

9. 1-2

3

$r_{b+1} = 2$

2 1

1.  $1 - 2$

$$\bar{j} = 4$$

1





# Algo-2

optimised

Not True  $\sqrt{N}$  Why?

$$4 \leq \sqrt{11}$$

for ( $i=2$ ;  $i \leq n$ ;  $i++$ ) {

if ( $n \% i == 0$ ) {  
 cnt = 0  
 while ( $n \% i == 0$ ) {  
 n = n / i  
 cnt = cnt + 1;  
 }  
 print (i, cnt);

if ( $n > 1$ ) { print (n, 1); }

outside for loop

$$N \leq \sqrt{N}$$

N=44

$$2 \cdot 11 = 4 \times 11 = 44$$

i	n	cnt
2	44	0
2	22	1
2	11	2
3	11	-
4	11	-
5	11	-
6	11	-
7	11	-
8	11	-
9	11	-
10	11	-

Stop

$\sqrt{N}$

4

11

must be prime no as we couldn't find any more div  $\sqrt{N}$

↑

$$\frac{11}{1} \leq \frac{11}{1}$$

$$\text{cnt} = 0$$

$$\text{cnt} = 1$$

$$N = 10$$

⇒

	$n$	$\text{cnt}$
<u>2</u>	<u>10</u>	0
2	5	1

↓

$\sqrt{5}$   
"  
2.2

$3 < \sqrt{5} \longrightarrow \text{stop}$

```

if (n > 1) {
    print(n)
}
    
```

2 x (5)

2 — 16

8	1
4	2
2	3

$N = \boxed{1}$  4

→ (2, 4)

$$5^2 \leq 55 \Rightarrow 25 \leq 55$$

$$6^2 \leq 36 \leq 11$$

$$\boxed{6^2 \leq 11}$$

for (i=2,  $i^2 \leq 11$  ;  $i++$ ) {  
Stop

3

if (n > 1) {  
print(1, 1);

3

i=2

i=3

i=4

i=5

~~i=6~~

i=6

N = 110

110

55

55

55

55

11

→ 11

↑  
Prime NO

cnt=0

cnt=1

x

x

cnt=0

cnt=1

x

No

(2, 5, 11)

= 2 x 5 x 11

= 110

$\sqrt{11} = 3$

JW → single

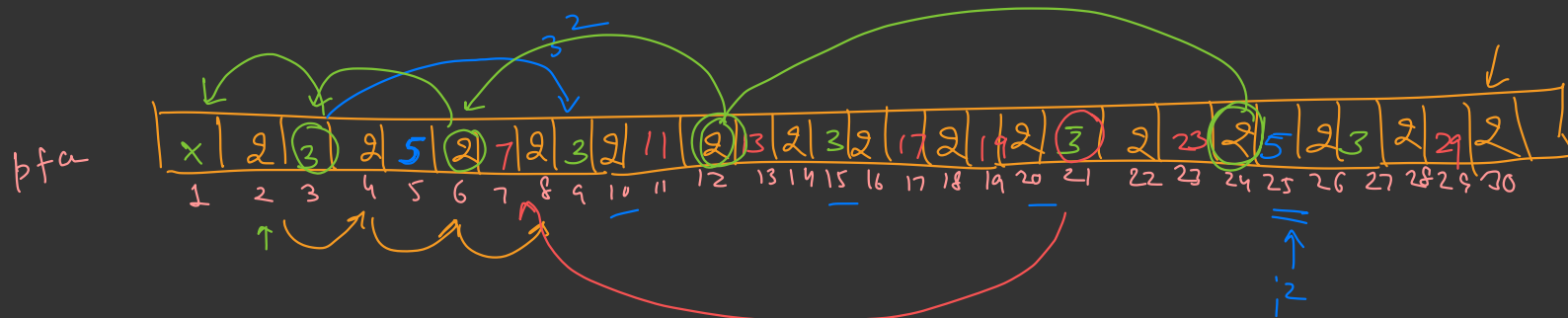
NJN → n inputs



# Prime Factorisation for a Range

for any 'no'  
store the smallest p.f

N = 30



① Build a sieve storing the smallest pf at a array?

② Prime factorisation

```

no = input()
while (no > 1) {
    print (pfa[no]);
    no = no / pfa[no];
}
    
```

3

N = 24

↓ / 2

12 / 2

↓

6 / 2 → 3 / 3 → 1 stop

21 → 3

1

$2^3 \cdot 3 = 24$





$$(3+1)(1+1) = 4 \times 2 = 8 \text{ ways}$$

$$\underline{56} = \rightarrow \underline{2^{(3)} \cdot 7^1} \quad (\text{Prime factorisation})$$



$$5 \times 3 \times 2$$

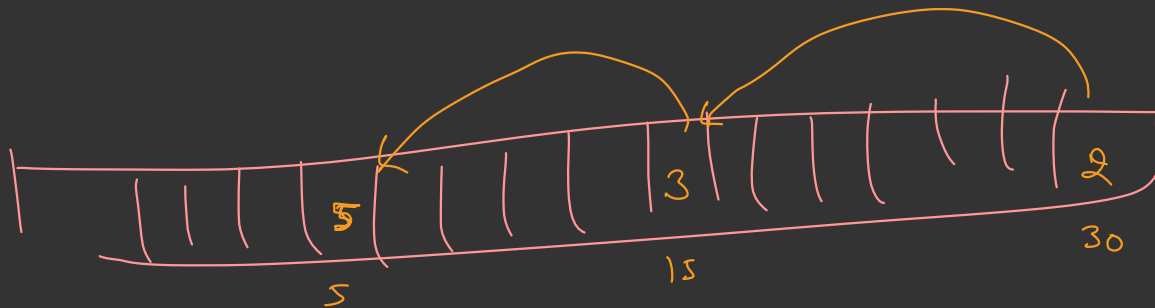
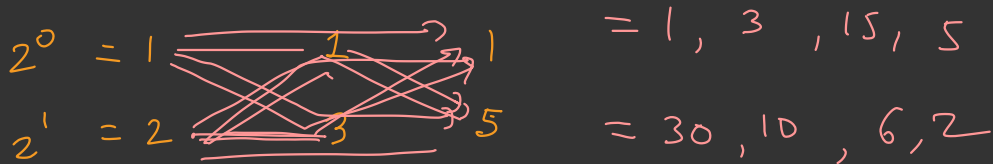
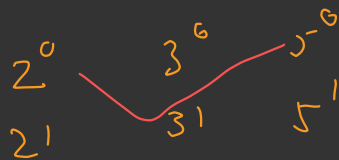
$$30 = 2^1 \cdot \underline{3^1} \cdot 5^1$$

$$= (1+\underline{1})(1+\underline{1})(1+\underline{1})$$

$$= 8$$

$$\begin{array}{l} 2^0 \\ 2^1 \\ 2^2 \\ 2^3 \end{array} \begin{array}{l} 7^0 \\ 7^1 \end{array}$$

$$\begin{array}{l} = 1 \\ = 7 \\ = 2 \\ = 14 \\ = 4 \\ = 28 \\ = 8 \\ = 56 \end{array}$$



$$30 = 2 \times 3 \times 5$$