

UNIT-02

Relational Algebra.

Two formal languages associated with relational data models are ① Relational Algebra

② Relational calculus.

Operators

① Selection (σ): - To select the rows or records or tuples of a table

$C \rightarrow$ condition

Syntax: $\sigma_C (T/E)$. $T \rightarrow$ table name / Expression

Ex- Find the details of all sailors with rating above 7

$\sigma_{\text{rating} > 7}(\text{sailors})$

② projection (π) \rightarrow To select columns of a table

Syntax:-

$\pi_{A_1, A_2, A_3, \dots, A_n} (T/E)$

$A_1, A_2, \dots, A_n \rightarrow$ Column

$T \rightarrow$ table name

Ex- Find the sid, name of sailors

$\pi_{\text{sid, name}}(\text{sailors})$

③ Set operators: - $\cup, \cap, -$

$\cup \rightarrow$ union

$\cap \rightarrow$ intersection

$- \rightarrow$ set difference

Ex-

Division - It is used to find records in one relation that are associated with all records in another relation. - It is commonly used when we want to identify entities that satisfy certain conditions across multiple related data sets.

~~Ex~~ R_1, S are two relations. The division operator, the division operators, $R \div S$ can be applied if

- 1) the attributes of S are proper subset of the attributes of R

result will include all attributes of R except those that are in S . It returns the tuples from R that are associated with every tuple in S .

Ex:- Find the sailors who resumed all boats

$$\pi_{\text{sid, bid}} (\text{Resumes}) \div \pi_{\text{bid}} (\text{Boats})$$

Find the sailors who resume all red boats

$$\pi_{\text{sid, bid}} (\text{Resumes}) \div \pi_{\text{bid}} (\sigma_{\text{color} = 'red'} (\text{Boats}))$$

Generalised projection

It extends the projection operations by allowing operations such as arithmetic and string functions to be used in the projection list.

Syntax

$$\pi_{f_1(A), f_2(A), \dots, f_n(A)} (Table)$$

Ex:- Find the independent incremented ratings of all sailors.

$$\pi_{\text{rating}+1} (\text{Sailors})$$

Compute the annual salaries of all employees

$$\pi_{\text{salary} * 12} (\text{employee})$$

Aggregation operators

↳ per means the use of aggregate function

Syntax:-

$$G_{op(Att)} (T(E))$$

$$G (T(E))$$

$$C_1, C_2, C_3, \dots, op(Att)$$

This is to apply aggregate functions on different groups

Ex:-

① Find the no. of sailors:

\int Sailors.
count(s.id)

② Find the avg age of sailors for each rating level

\int Sailors.
rating avg(age)

TRC [Tuple Relational calculus]

Syntax

$\{T \mid P(T)\}$

T - A tuple variable

P(T) - Formula that describes a tuple variable

Ex:-

① Find all sailors with rating above 7

$\{S \mid S \in \text{sailors} \wedge S.\text{rating} > 7\}$

② Find the all details of sailors with age < 40

$\{S \mid S \in \text{sailors} \wedge S.\text{age} < 40\}$

③ Find s.id, s.name of the sailors with rating > 7

$\{P \mid \exists S \in \text{sailors} \wedge (S.\text{rating} > 7 \wedge P.\text{s.id} = S.\text{s.id} \wedge P.\text{s.name} = S.\text{s.name})\}$

Ex-① Find the sailor name, boat id, and reservation day with each reservation.

$\{P \mid \exists S \in \text{sailors} \exists R \in \text{Reservations} . (S.sid = R.sid \wedge P.name = R.name \wedge P.bid = R.bid \wedge P.day = R.day)\}$

② Find the names of sailors who have reserved boat id

$\{P \mid \exists S \in \text{sailors} . \exists R \in \text{Reservations} (S.sid = R.sid \wedge P.name = S.name \wedge R.bid = 103)\}$

open to k

Syn

Ex-2

Comp

Agg

lp

Syn