

Secure Sockets Layer

Secure Sockets

Secure Sockets Layer

SSL/TLS provides endpoint authentication and communications privacy over the Internet using cryptography.

For web browsing, email, faxing, other data transmission.

In typical use, only the server is authenticated while the client remains unauthenticated

Mutual authentication requires PKI deployment to clients.

Protocols allow client/server applications to communicate in a way designed to prevent eavesdropping, tampering, and message forgery.

Secure Sockets Layer

Applications HTTP, FTP, ...

Presentation

SSL

Transport

TCP

Network

IP

EH

User

OS

Kerberos

TCP

IPSec

IP

...

Presentation Layer:

Provides independence from differences in data representation among applications. OS functions translate data formats (app protocol syntax) to a uniform network format (bit stream to be transmitted), and vice versa so as to eliminate network compatibility problems. Takes care of encryption and data compression.

Secure Sockets Layer

Developed by Netscape Communications Corp (1995)

Ensures data privacy: transmission of data via encryption

Supports Server and Client authentication

Supports authentication of service via certificate

Ensures data integrity

Application independent - ftp, http, telnet are layered on top of it. Mainly used in https applications.

Can negotiate encryption keys

Sits on top of TCP/IP, does not require OS changes

Can be used to create a tunnel for VPN

Encryption and compression apply only to application layer

Secure Sockets Layer

Three phases:

1. Peer negotiation for algorithm support (see below)
2. **Public key encryption** based key exchange and certificate based authentication
3. **Symmetric cipher** based traffic encryption

Secure Sockets Layer

Security features:

- Numbering all the records and using the sequence number in the Msg. Authen. Codes. (to prevent replay attacks)
- Uses HMAC for message integrity check
- Protection against several other known attacks
 - MIM attack involving a downgrade of the protocol to a less secure version via integrity check of initial handshake
 - identity fraud via digital signatures
 - known plaintext attacks via strong cryptographic algorithms
- The message that ends the handshake ("Finished") sends a hash of all the exchanged data in handshake (Integrity check to prevent Man in the Middle and Truncation attacks).

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Uses several protocols organized into layers as follows:

SSL handshake protocol	SSL cipher change protocol	SSL alert protocol	Application Protocol (eg. HTTP)
SSL Record Protocol			
TCP			
IP			

SSL Record Protocol: handles data security and integrity; encapsulates data sent by higher level protocols

Handshake, Cipher change, Alert: establish a connection; session management, crypto management, SSL message transfer

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Definitions:

Connection: logical 2-node peer-to-peer link – provides a service

Session: association between peers defining crypto algorithms, sequence numbers, etc. Created by handshake protocol. Used to avoid renegotiation of parameters from connection to connection

Session State:

session identifier: generated by the receiver

peer certificate: X.509 spec

compression method: prior to encryption

CipherSpec: encryption, integrity, and hash algorithms

MasterSecret: 48 byte shared secret

Secure Sockets Layer

Definitions:

Connection State:

Random numbers - chosen by server and client to make crypto breaking harder

Server write MAC secret - used on data from server

Client write MAC secret - used on data from client

Server write secret key - server encryption, decryption by client

Client write secret key - client encryption, decryption by server

Initialization vectors - for CBC ciphers

Sequence number - for both transmitted and received messages on both client and server sides

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SSL Record Protocol:

Fragment the data that needs to be sent: create records

Encapsulate them with appropriate headers

Create an encrypted object that can be sent over TCP

Header of each record: length of record and of data block

Contents of record after header: data, padding, MAC

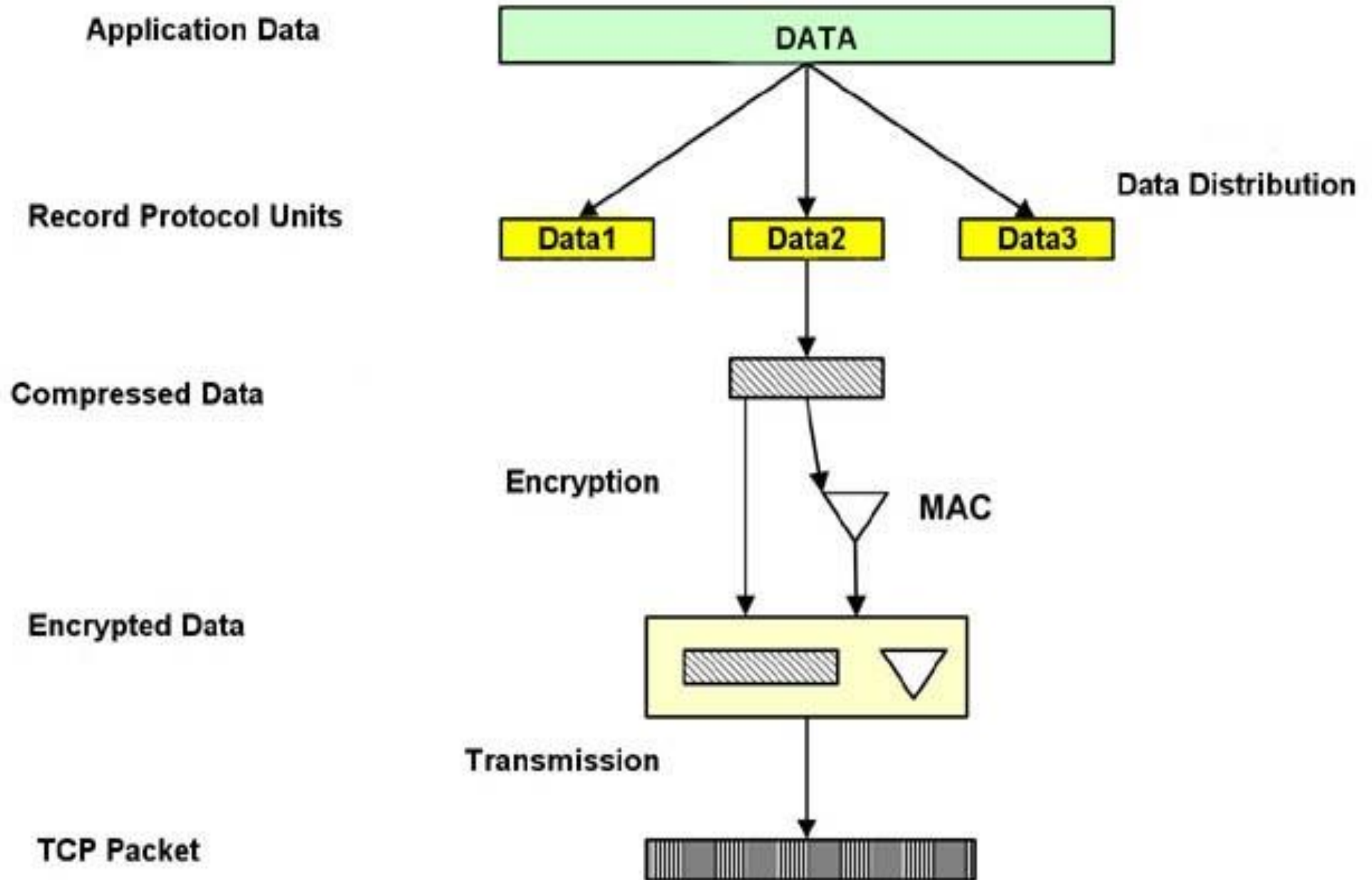
MAC = hash {secret key, data+padding, sequence number}

where hash uses specified (negotiated) algorithm like SHA-1

Encrypted Object: encrypt record plus MAC

Header of EO: content-type: which of four protocols to use to handle the data in the EO after decryption.
Protocol Major and Minor version numbers.

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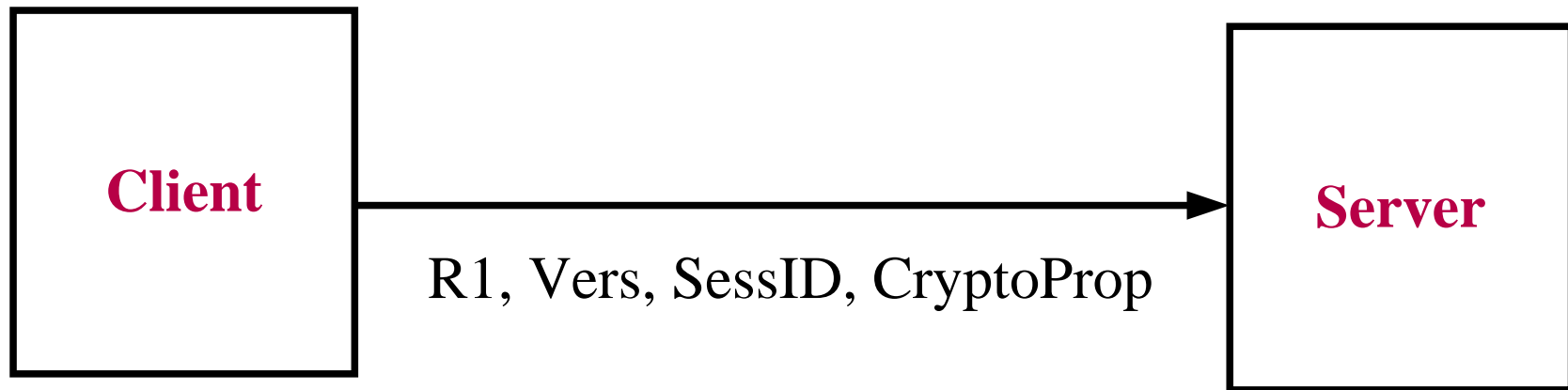
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Handshake Protocol:

To initiate a session: crypto negotiation

phase 1: initiate (client_hello), identity not revealed!

Establish logical connection, negotiate session parms



Version: Highest SSL version supported by client

R1: random number

SessionID: non-0: resume earlier session, modify parms of this session,
spawn an independent connection without handshake
0: initiate a session

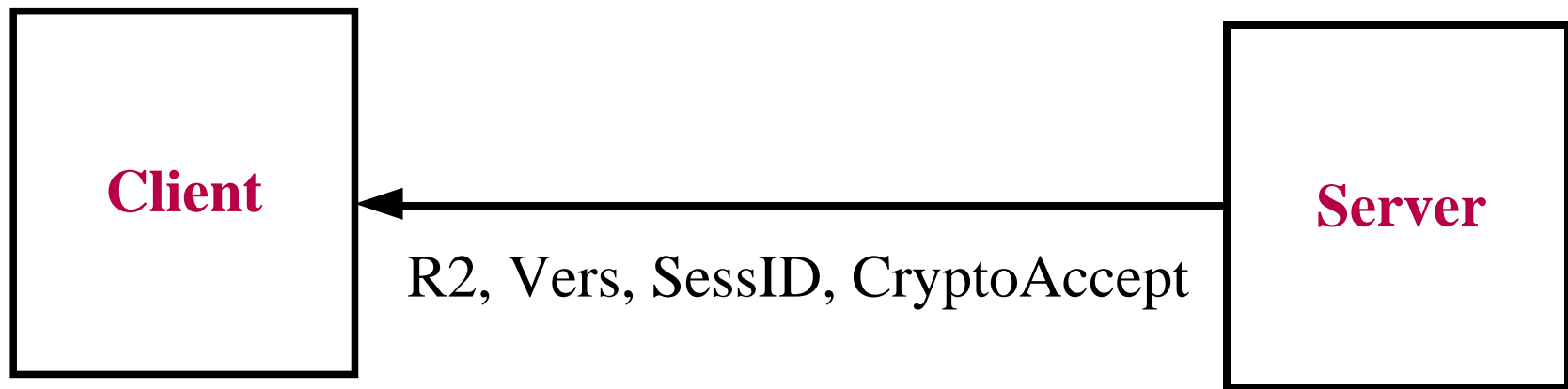
Secure Sockets Layer

Handshake Protocol:

To initiate a session: crypto negotiation

phase 1: initiate (client_hello), identity not revealed!

Establish logical connection, negotiate session parms



Version: Lowest SSL version supported by server

R2: a random number

SessionID: same as client's, if non-0, otherwise an ID decided by server

CryptoAccept: key exchange method, encrypt algorithm, hash function

Secure Sockets Layer

Handshake Protocol:

To initiate a session: crypto negotiation

phase 2: authenticate (server_hello)



Certificate(s): chain of certificates to a trusted CA to authenticate server

Diffie-Hellman value: optional

Request for certificate from client: optional

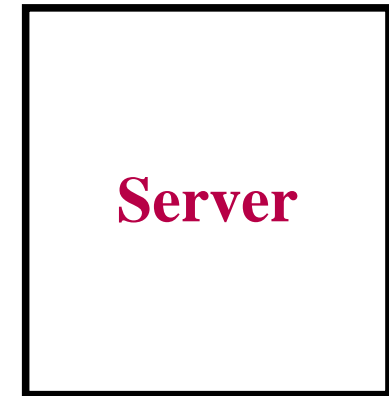
server_done: completes the message sequence

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Handshake Protocol:

To initiate a session: crypto negotiation

phase 2: authenticate (server_hello)

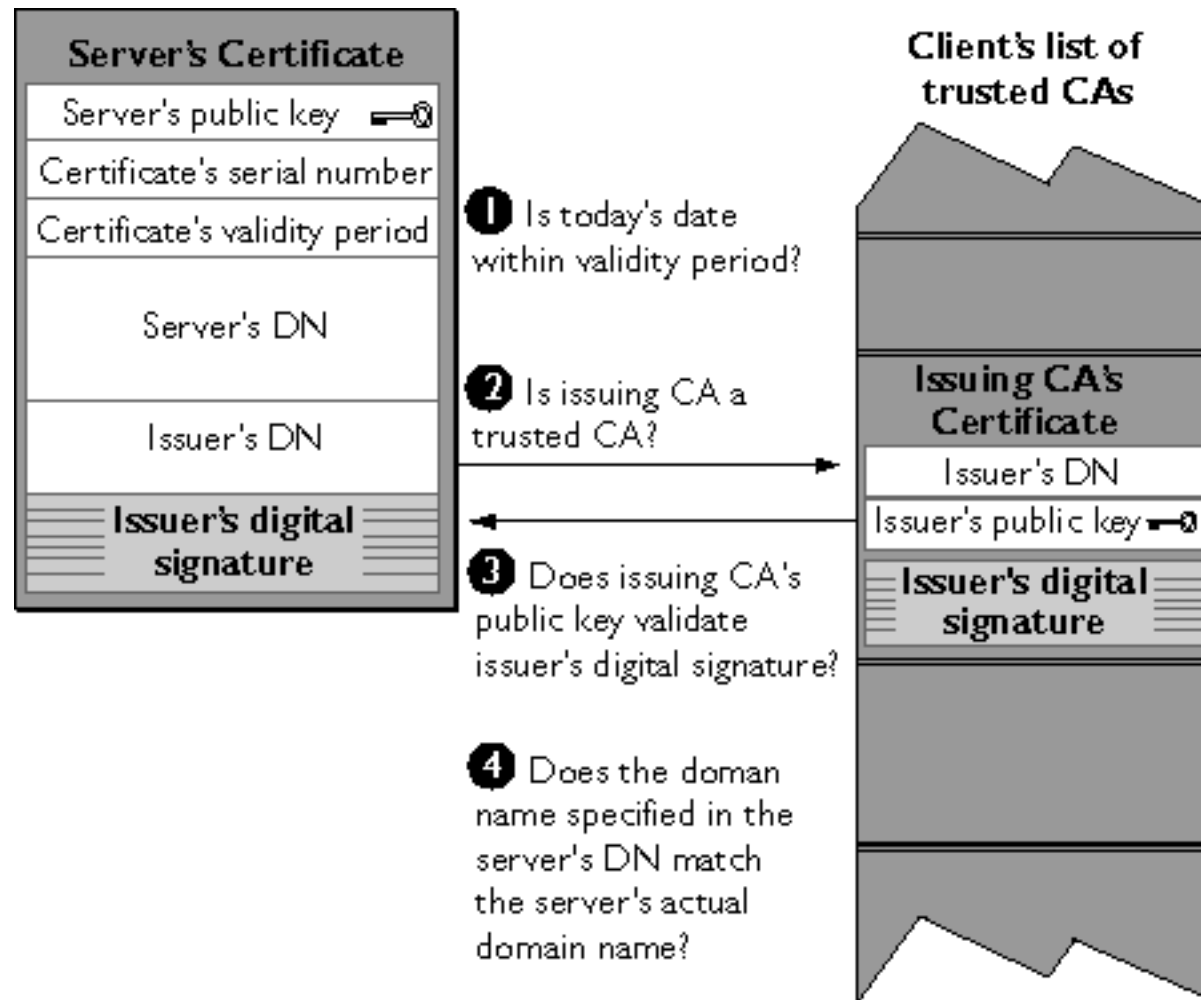


Client verifies certificate, checks date and invalidation lists

Client checks that the certifying authority is trusted

Client checks the CA's public key against that of the certificate

Client checks that the domain name in the certificate matches that of server

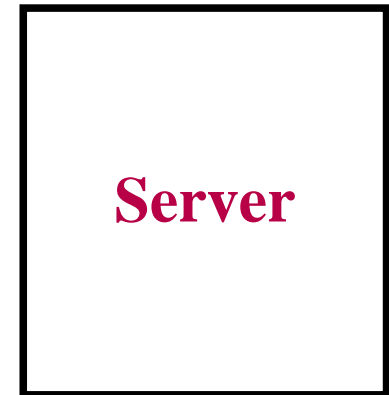


How the client authenticates the server

Secure Sockets Layer

Handshake Protocol:

To initiate a session: too many secrets



Client chooses random number S (pre-master secret) and encrypts it with server's public key ($\text{server}\{S\}$)

Client computes master key as $K=f(S,R1,R2)$, computes $\text{hash}(K+msgs)$

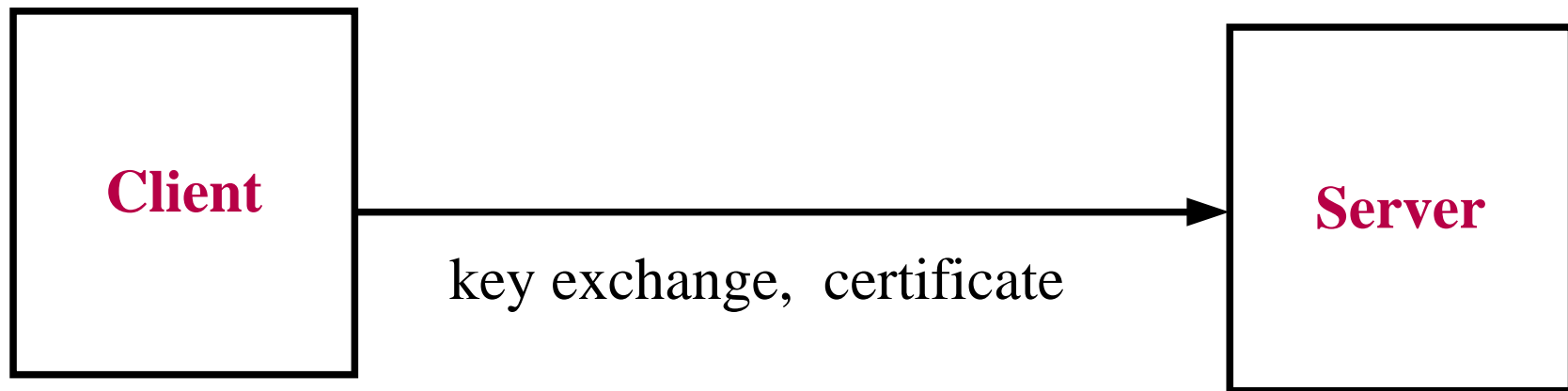
Six secret keys - 3 from client to server and 3 from server to client
integrity, encryption, initialization vector (derived from S)

Secure Sockets Layer

Handshake Protocol:

To initiate a session: crypto negotiation

phase 3: start key exchange



S sent encrypted by server's public key, server computes master secret

Key exchange msg: delivers the keys, depends on the agreed key exchange method

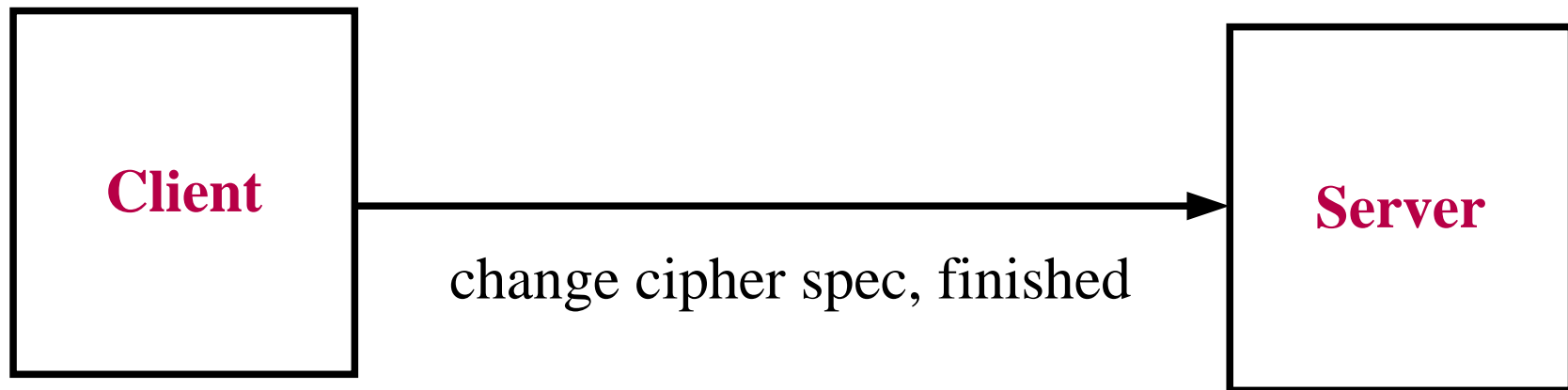
Send certificate of client: if requested by server

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Handshake Protocol:

To initiate a session: crypto negotiation

phase 4: confirmation and setup



Change cipher spec msg: crypto spec now considered agreed

Setup of algorithms: make cryptosystems ready to go

Finished msg: encrypted with agreed upon crypto algorithms and keys
so server can verify that communication is possible.
Is the hash of all the messages in the handshake.

Secure Sockets Layer

Handshake Protocol:

To initiate a session: crypto negotiation

phase 4: confirmation and setup



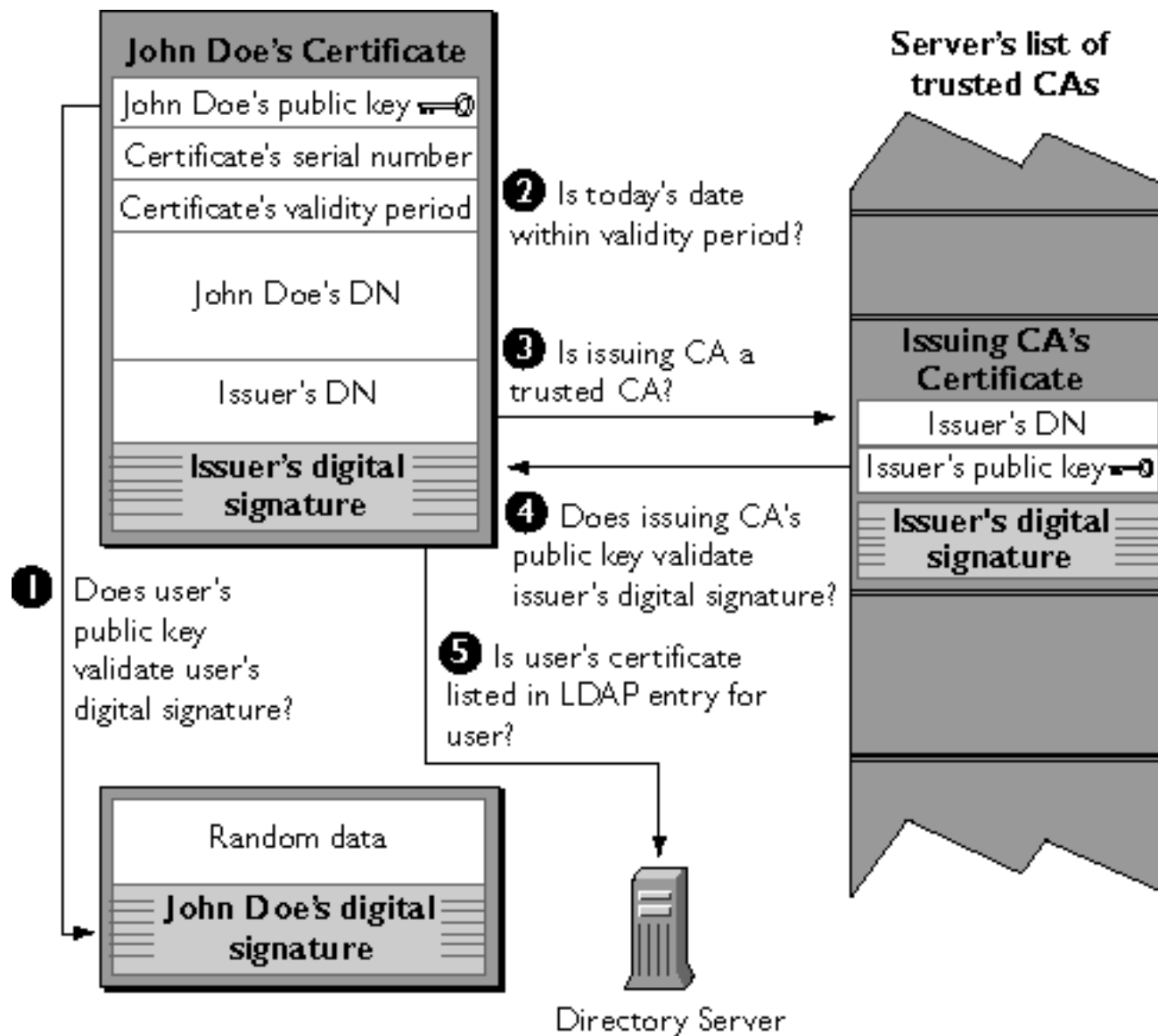
Same finished message is sent back to client

Session is terminated and the TCP connection is closed but the "state" of the session is saved to be reopened later with same parameters.

I client

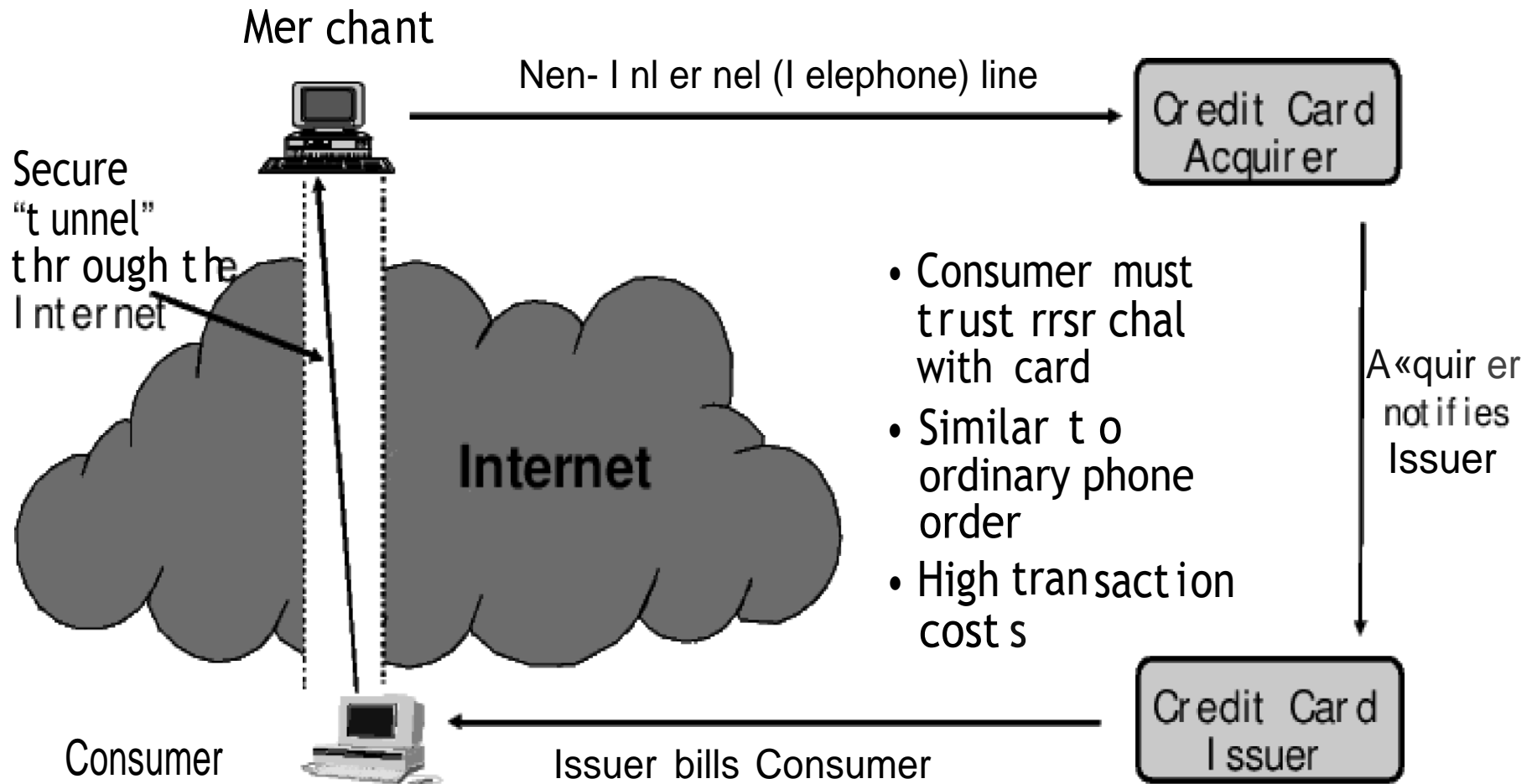
Server





How the server authenticates the client

The Main Usage of SSL



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Session Resumption:

Per-session master key is established using expensive public key cryptography

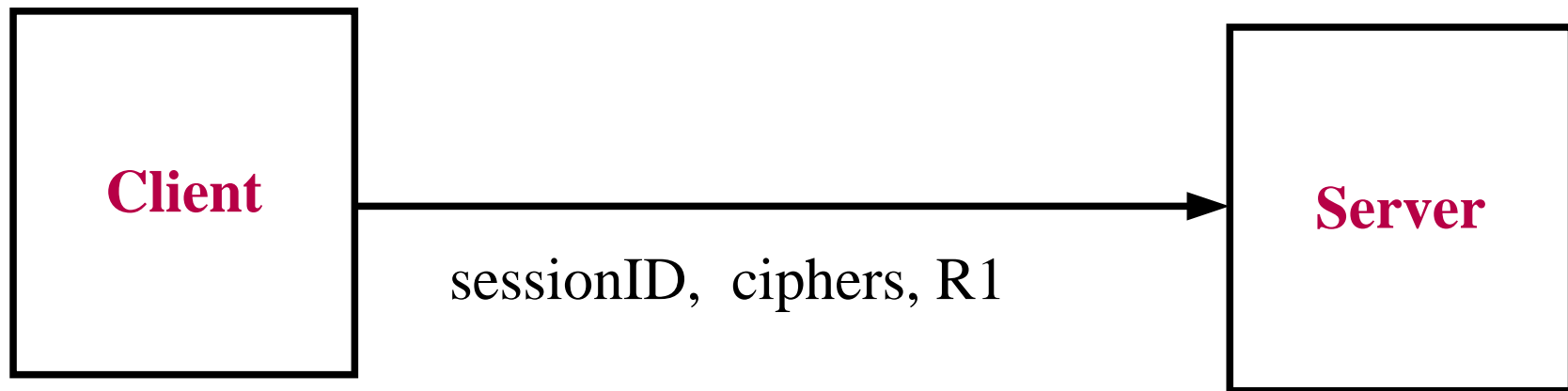
Connections cheaply derived from master key with handshake involving nonces, not public keys

SessionID and master key for the session is stored by server to support resumption

If server loses the state, it can be reestablished by the client sending S encrypted with server's public key

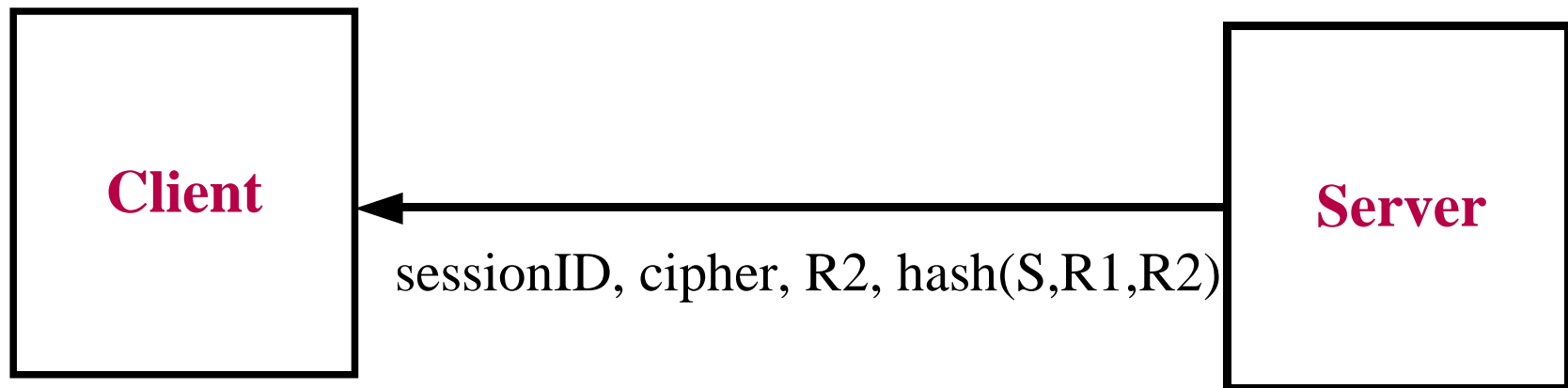
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Handshake Protocol (Resumption): Using Session-ID



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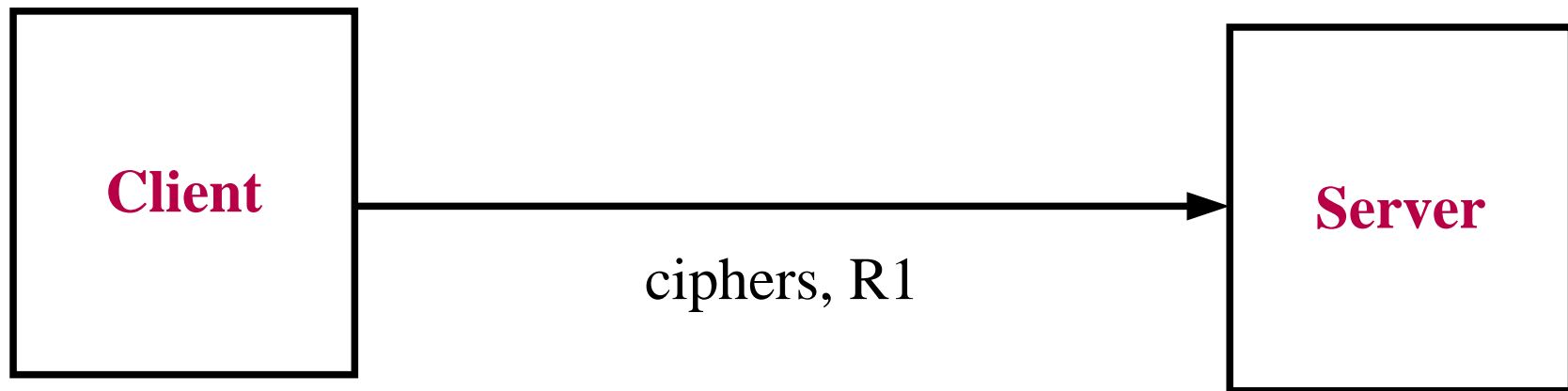
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Handshake Protocol (Resumption): Using Session-ID



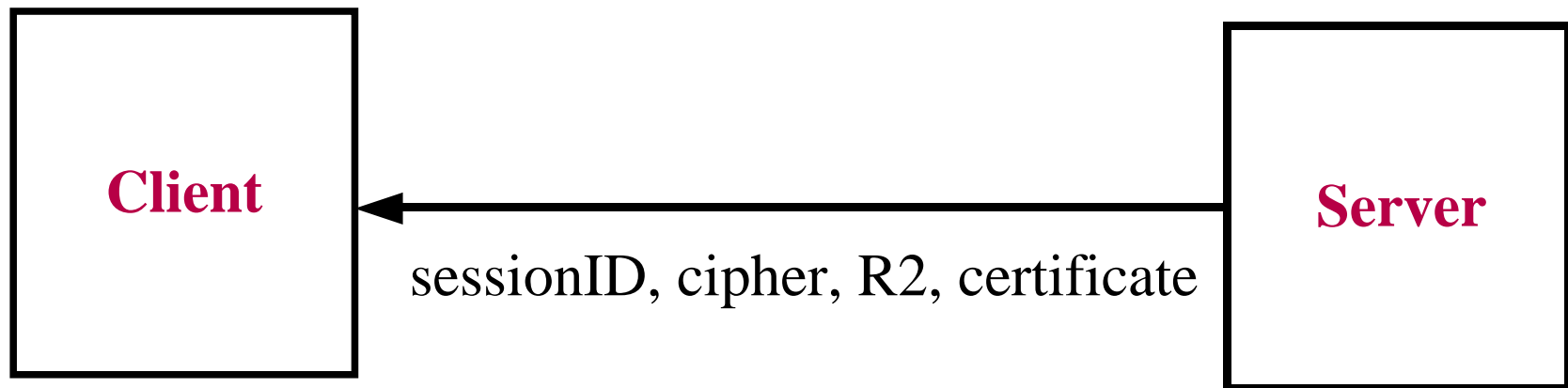
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Handshake Protocol (Resumption): Forgot Session-ID



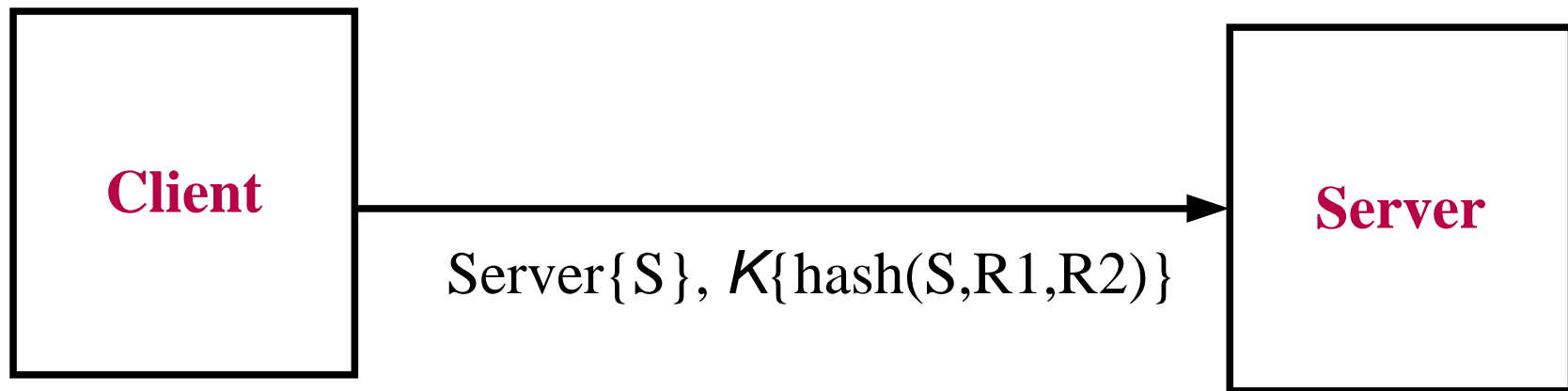
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Handshake Protocol (Resumption): Forgot Session-ID



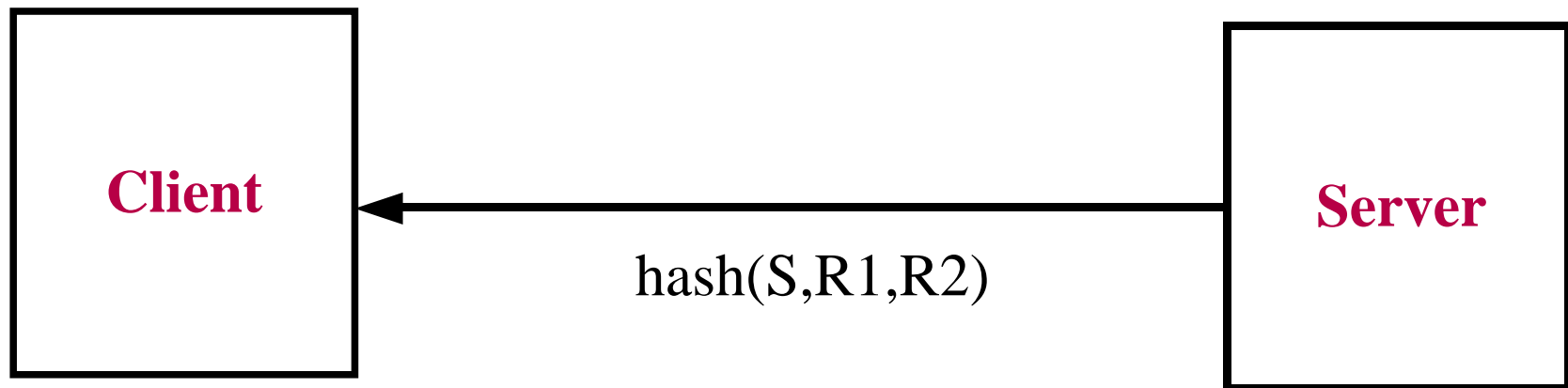
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Handshake Protocol (Resumption): Forgot Session-ID



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Handshake Protocol (Resumption): Forgot Session-ID



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Computing Keys:

If Diffie-Hellman is used to compute pre-master Secret S, fixed DH is allowed! Same session key repeated!!
Ephemeral DH, ephemeral elliptic curve DH are also allowed for perfect forward secrecy (e.g. gmail)

If kept around in memory, a fixed DH key could compromise future communications

So it is hashed with nonces to get the master secret
(Stealing master secret only affects this communication)

