

# Remove Node From BST

Saturday, 26 June 2021 2:34 PM

if (data is found) →

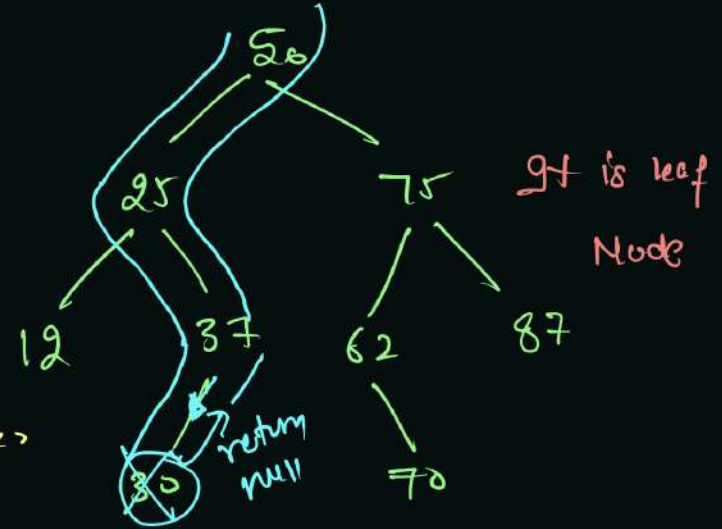
① leaf Node ] return null

② Node have only left child ] return left child of node

③ Node have only right child ] return right child of node

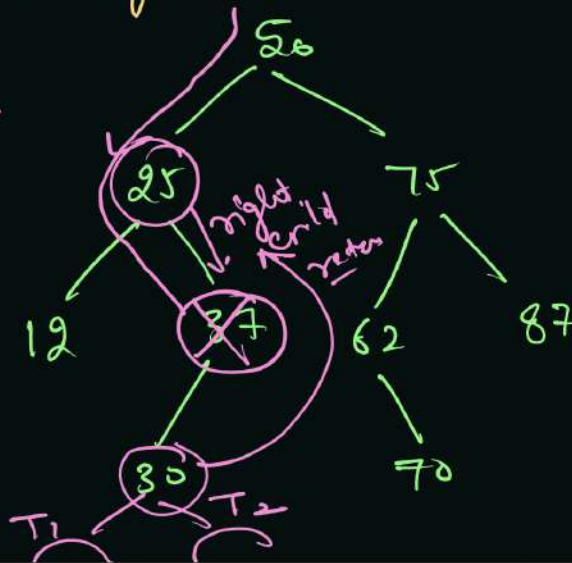
④ In the middle, have both left and right child

data = 30



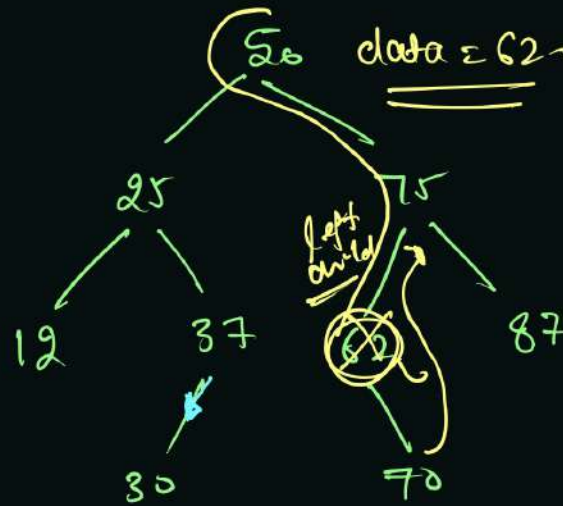
data = 37

have left child



data = 62

have right child

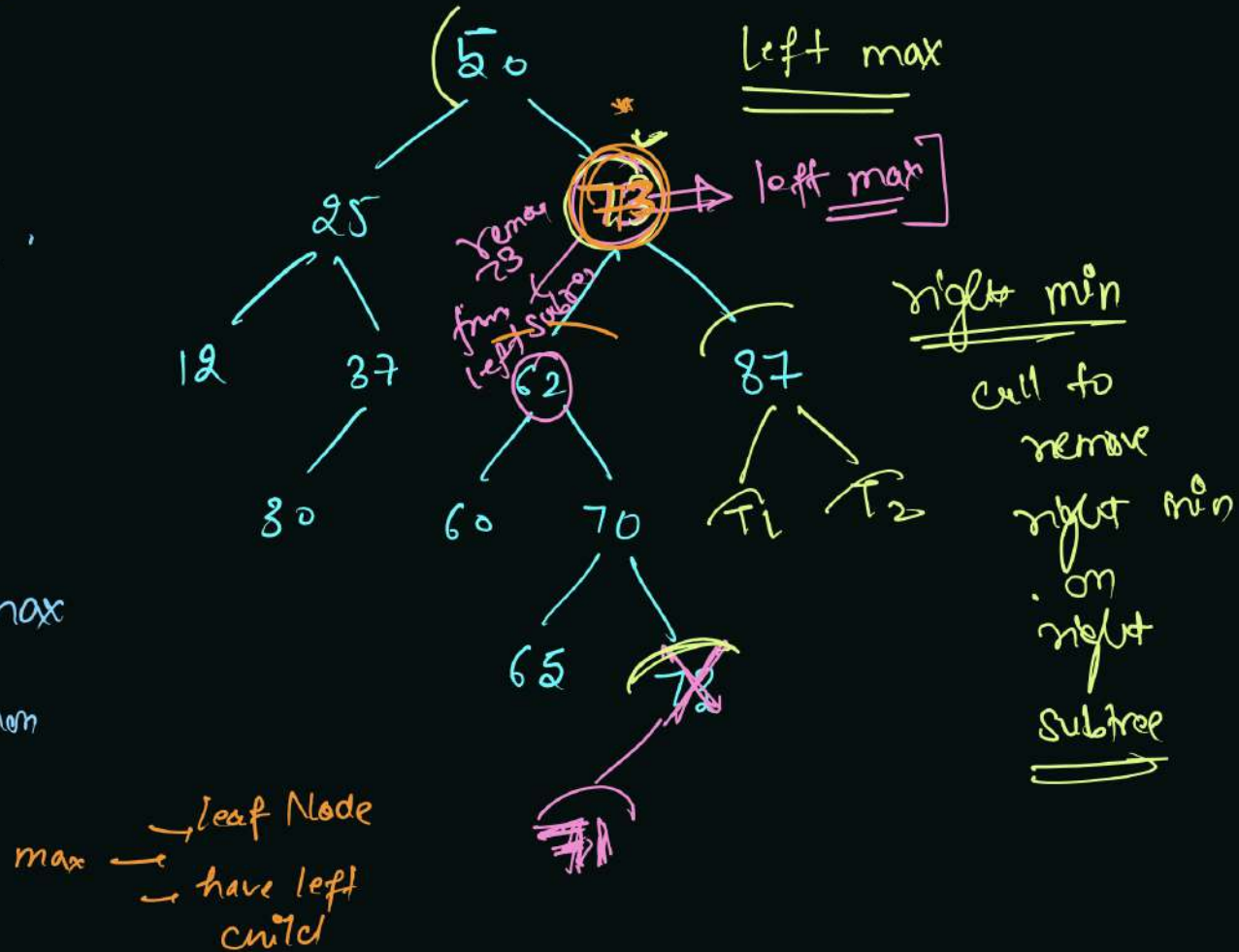


Node having both left and right child

data = 75.

BST  $\rightarrow$  property - Maintain  
Maintain BST property.

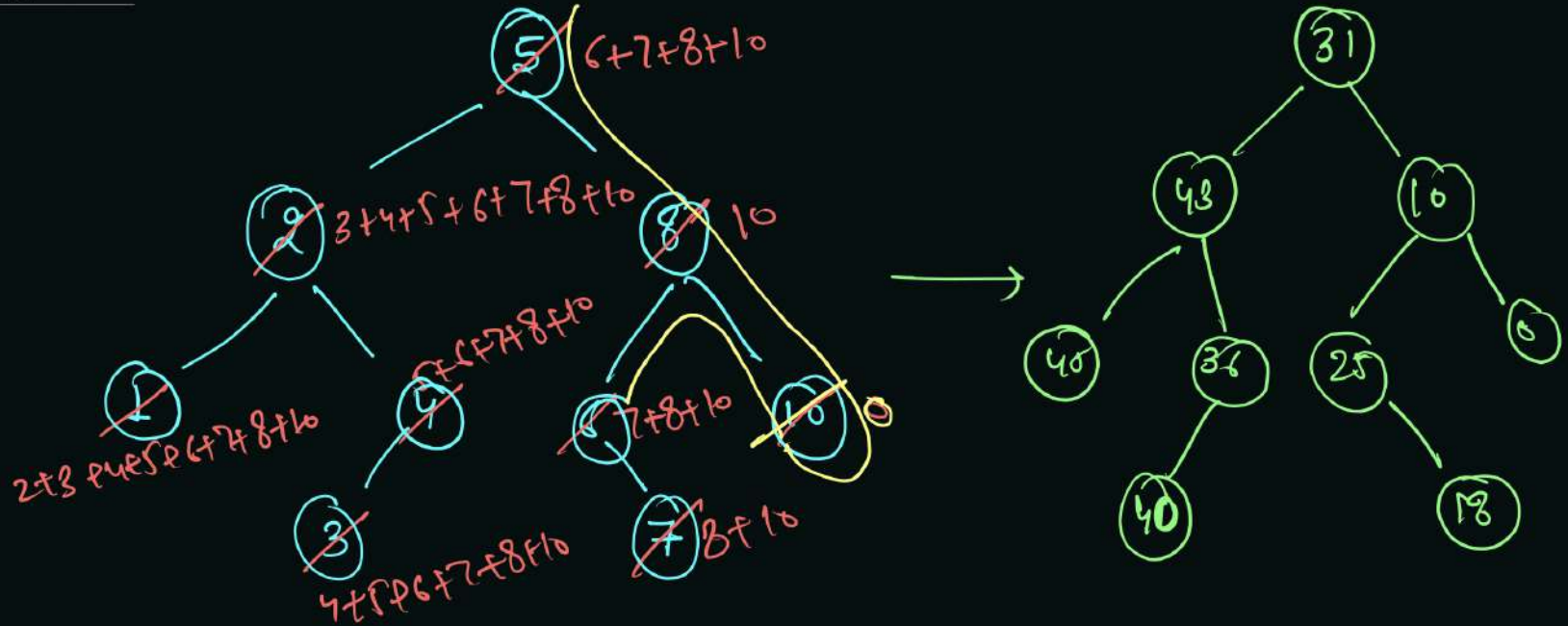
- steps
- ① max from left subtree
  - ② Set the data of node as left max
  - ③ Remove max, from left subtree



## Replace with Sum of Larger

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sum = 10



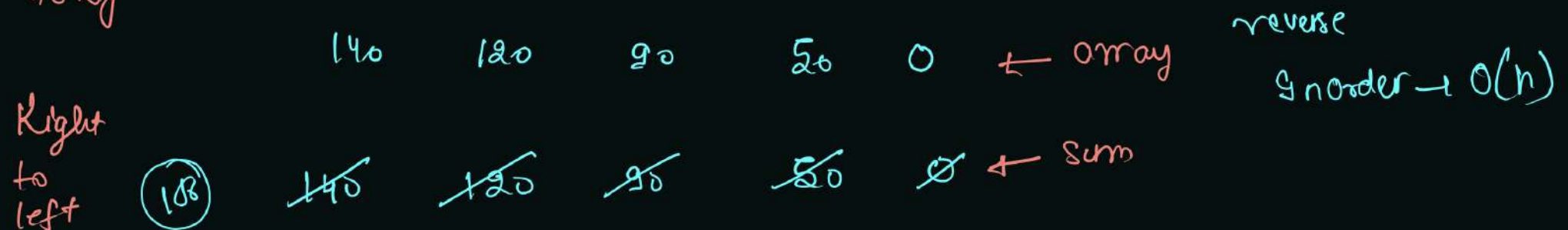
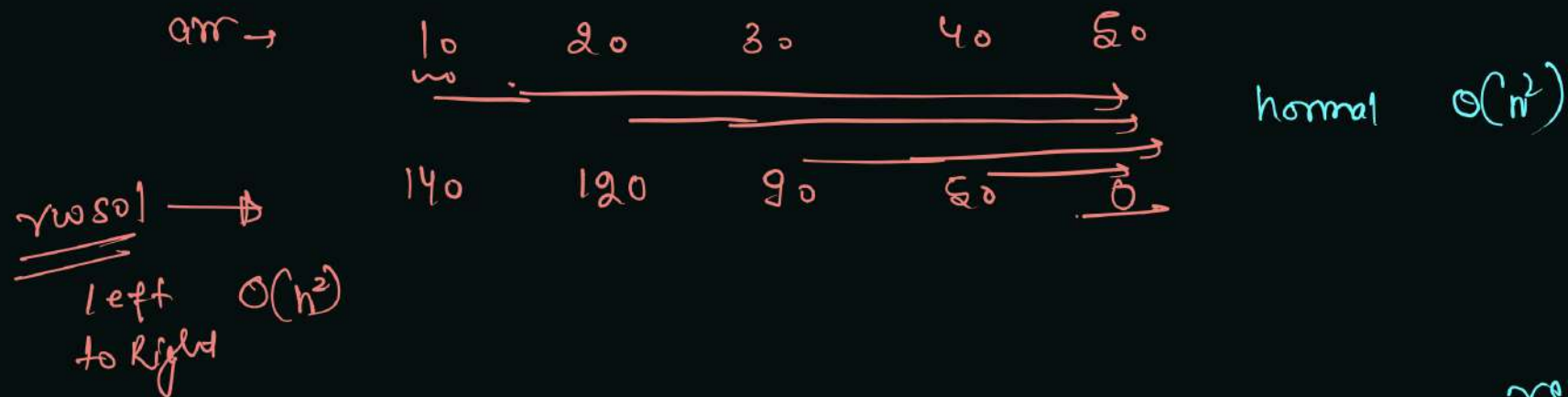
Space  $\rightarrow O(1)$  Except recursive space

Time  $\rightarrow O(n) \rightarrow$  Single traversal on Tree

In Order of BST  $\rightarrow$  Sorted, (Increasing order)

Inorder of BST  $\rightarrow$  Sorted (Increasing)

Reverse inorder of BST  $\rightarrow$  Sorted (Decreasing)





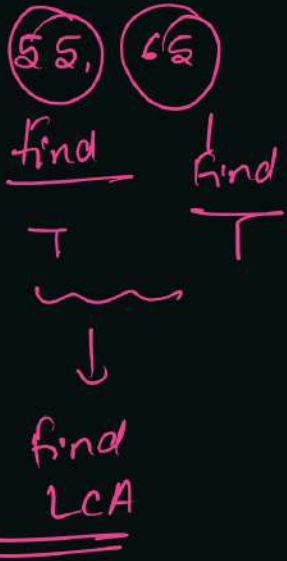
# LCA of BST

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LCA  $\rightarrow$  (55, 65)  $\rightarrow$  62.

space  $\rightarrow$  recursive space

time  $\rightarrow$   $O(\text{height})$



✓ node.data < {d1, d2}

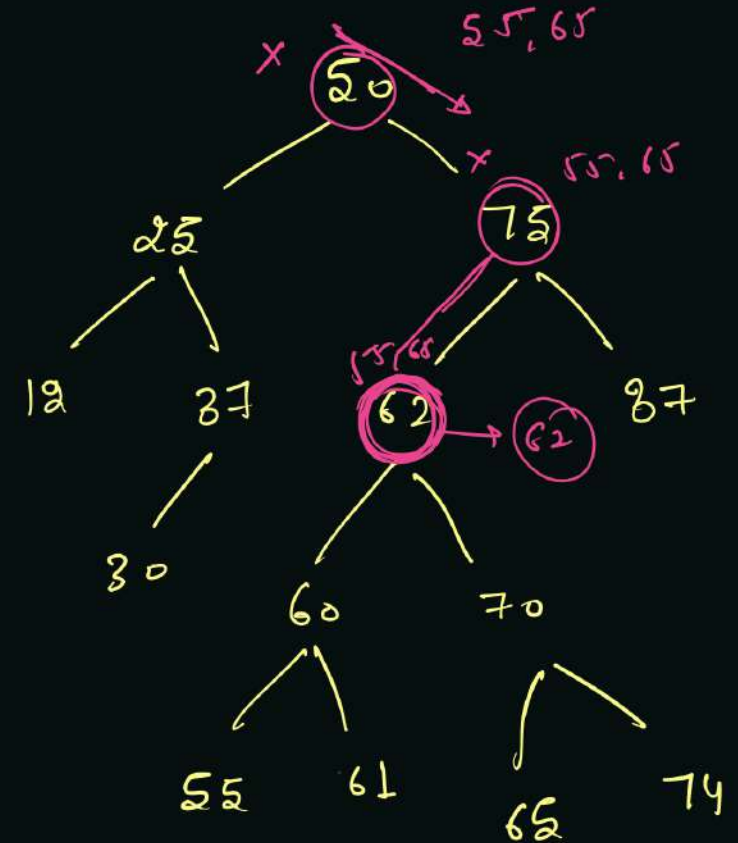
↳ Right side (LCA)

✓ {d1, d2} < node.data

↳ Left side (LCA)

✓ {d1 ≤ node.data ≤ d2 || d2 ≤ node.data ≤ d1}

else  $\rightarrow$  LCA is node.data



## Print In Range

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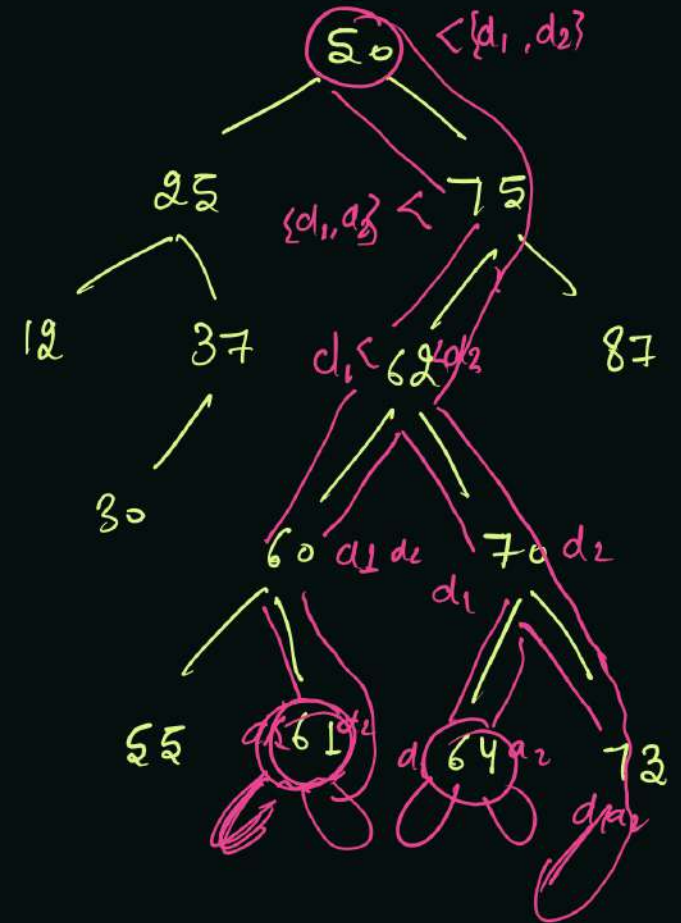
$lo = 26$   
 $upper = 72$   $\rightarrow$  print values

Time complexity  $\rightarrow$  less than  $O(n)$   
but also depend on  
Range.]

output  $\rightarrow$

[25 30 37 50 55 60 61 62 64 70]

61 62 64 70



## Target Sum Pair In BST

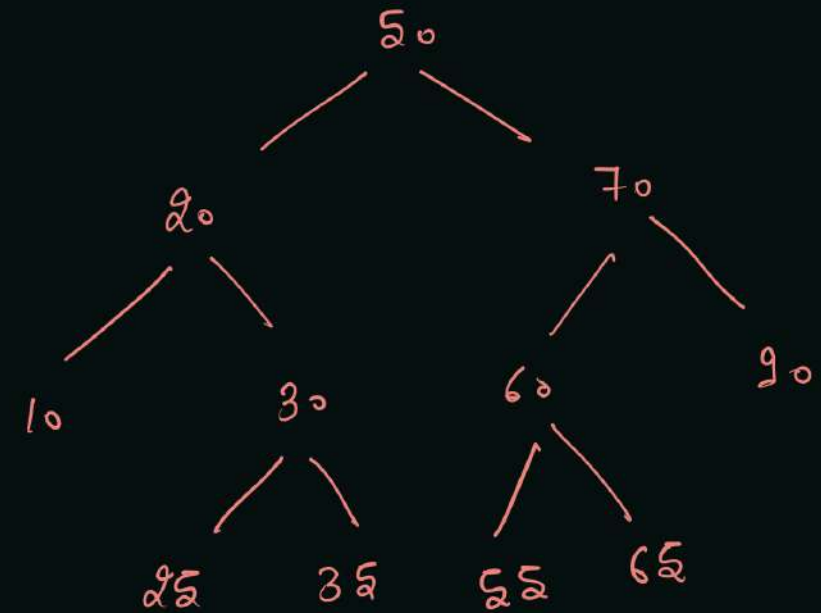
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target = 100

print all pairs from  
which we can achieve

target  $\Rightarrow$  100

Ex - 10 90  
20 80  
30 70  
⋮ ⋮  
⋮ ⋮  
⋮ ⋮  
⋮ ⋮  
⋮ ⋮



Time

space

Method 1

$O(nh)$

$O(h)$

Method 2

$O(n)$

$O(n)$

Method 3

$O(n)$

$O(h)$

Time is optime but required more space

Time is optimise as well as space optimise.



Method 1 →

Space →  $O(1)$

Time :  $O(nh)$

For find. original root

In Order.

val1 = ~~10~~ ~~20~~ ~~25~~ ~~30~~ ~~35~~ ~~50~~ ~~55~~ ~~60~~ ~~65~~ ~~70~~ ~~90~~

val2 = target - val1

if present  
in the  
tree.

Traversal →  $O(n)$

find →  $O(1)$

Total →  $O(nh)$

~~90~~ → find.

val1 val2.

✓ 10 - 90  
✓ 30 - 70  
✓ 35 - 65

val1 < val2

65 - 35  
70 - 30  
90 - 10

Repetition pair  
val2 & val1

Repeated

Repetition

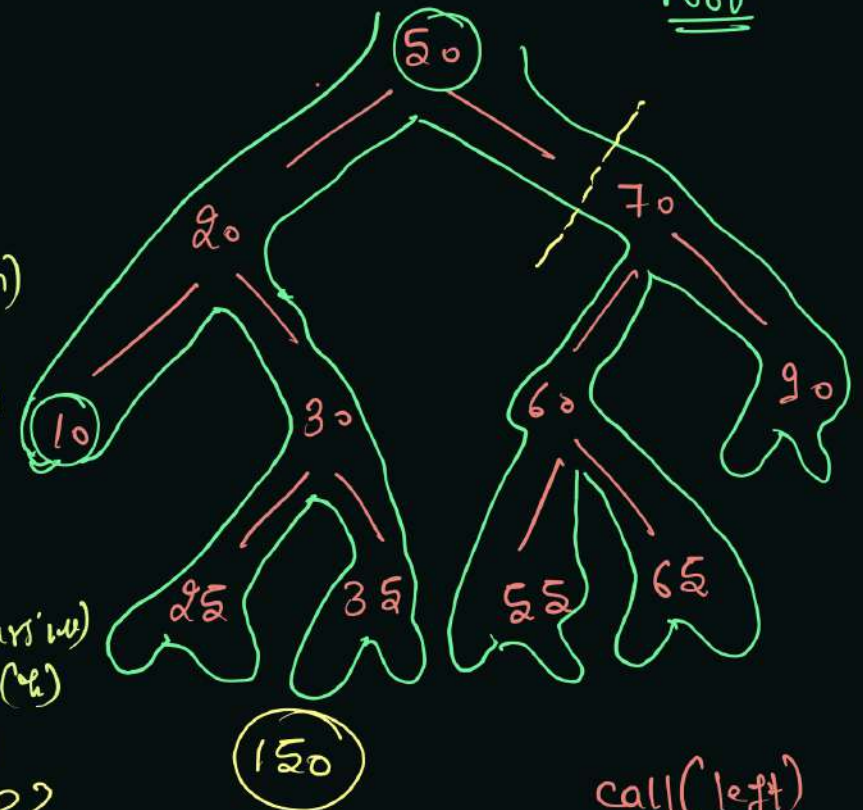
How to avoid these  
repeated pair??

make a find only  
when val2 > val1

call(left)

work.

call(right)

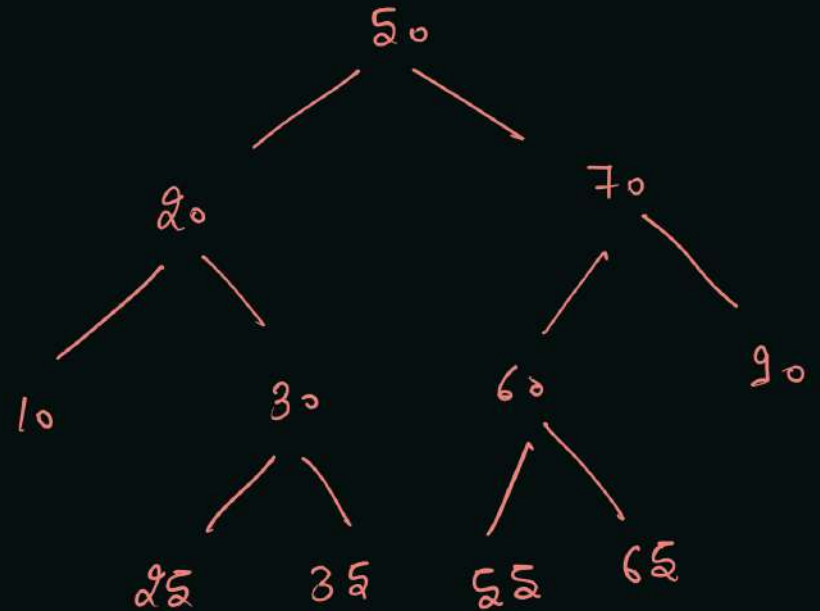




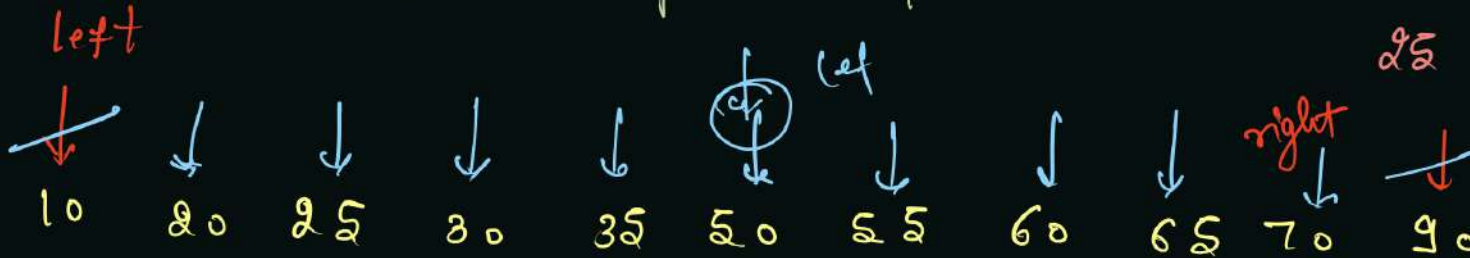
## Method 2.

Make an array and fill it in  
InOrder,  
→ target sum pair.

10 90  
30 70  
35 65



TNS → Target sum pair



$$\text{Sum} = \text{arr}[\text{left}] + \text{arr}[\text{right}]$$

if (Sum == target)  
print(arr[left], arr[right]);  
left++;  
right--;

else if (Sum > target)  
right--;

else  
left++;

condition for  
iteration → left & right

## Target Sum Pair 2

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target = 100

left = 10 20 30 40 50

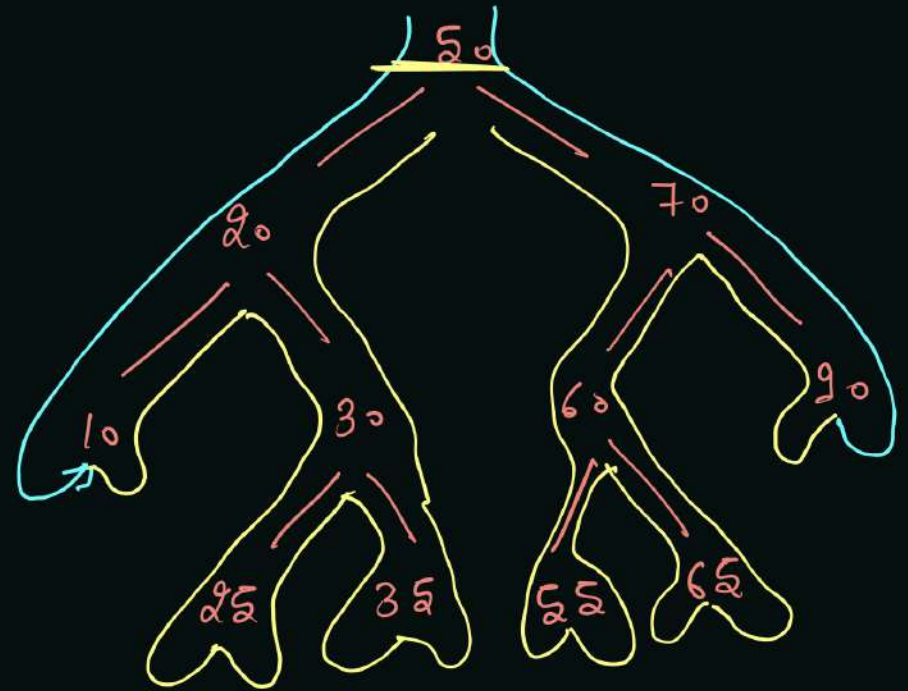
right = 90 70 65 55

sum = left + right

```
if (sum == target) {  
    print pair →  
    left++ // right--  
} else if (sum > target) {  
    right--  
}
```

```
} else {  
    left++  
}  
}
```

How we  
can handle  
left and  
right in  
recursive  
Tree ?



10 — 90  
30 — 70  
35 — 65

## Iterative Traversal :->

```

public static void targetSumPair3(Node node, int target) {
    ✓ Stack<IPair> lstack = new Stack<>();
    ✓ Stack<IPair> rstack = new Stack<>();

    ✓ lstack.push(new IPair(node, 0));
    rstack.push(new IPair(node, 0));

    ✓ Node left = itrInOrder(lstack);
    ✓ Node right = revItrInOrder(rstack);

    while(left.data < right.data) {
        int sum = left.data + right.data;
        if(sum == target) {
            System.out.println(left.data + " " + right.data);
            ✓ left = itrInOrder(lstack);
            ✓ right = revItrInOrder(rstack);
        } else if(sum > target) {
            right = revItrInOrder(rstack);
        } else {
            ✓ sum < target
            ✓ left = itrInOrder(lstack);
        }
    }
}
    
```

State

left

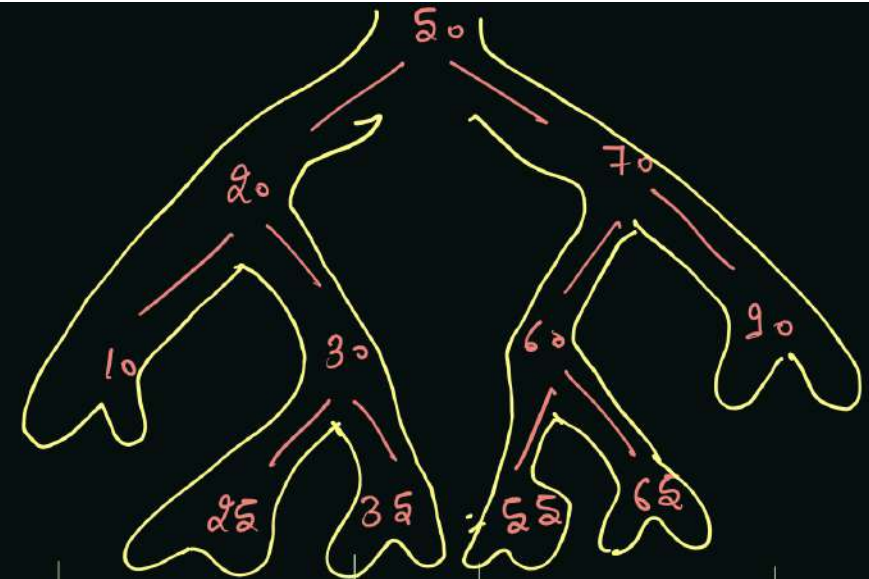
- 0 → left child
- 1 → right child
- 2 → pop

right

- 0 → right child
- 1 → left child
- 2 → pop

left = 10 20 28 30 32 (50)

right = 90 70 52 60 55 (50)



$$\begin{array}{r} 70 - 0 \\ \hline 50 - 0 \end{array}$$

Iter Inorder

10 - 90  
20 - 70  
35 - 65

$$\begin{array}{r} 70 - 0 \\ \hline 50 - 0 \end{array}$$

Rev Iter Inorder