

ASSIGNMENT-6

Q1. 0-1 Knapsack Problem, Total weight limit = 8

Items 1 2 3 4
Value 8 40 30 54
Weight 1 2 3 6

$$\text{Sum}[K, s] \leftarrow \max[\text{Sum}[K-1, s], \text{Sum}[K, s-w[K] + v[K]]]$$

	s=0	1	2	3	4	5	6	7	8
K=0	0	0	0	0	0	0	0	0	0
1	0	8	8	8	8	8	8	8	8
2	0	8	40	48	48	48	48	48	48
3	0	8	40	48	48	48	78	78	78
4	0	8	40	48	48	48	78	78	(94)

Hence maximum profit is 94 with the items 2 and 4 with respective weights $2+6=8$.

Q2. Trace back algorithm.

Step 1: First find max value

Step 2: Start at bottom right cell (max value)

Step 3: Check current value in cell and compare with value of the cell above. If its equal current corresponding to current value is not selected.

Step 4: If its not equal, current row is selected

Add its index to selected items

Subtract its W and value.

Move cell diagonally above and to the left

Step 5: Repeat Step 3 & 4 until you reach top of the table

Step 6: Final selected items are the required

Q2. Traceback algorithm to determine the exact subset of items

```
def knap(weights, values, W):
    n = len(weights)

    # 2D array
    table = []
    for i in range(n+1):
        t = [0] * (W+1)
        table.append(t)

    # sub-problems
    for i in range(1, n+1):
        for w in range(1, W+1):
            if weights[i-1] > w:
                table[i][w] = table[i-1][w]
            else:
                table[i][w] = max(table[i-1][w], values[i-1] + table[i-1][w-weights[i-1]])

    # Choose Max value
    max_value = table[-1][-1]

    # last cell
    return table, max_value

def backtrack(table, weights, values, W):
    n = len(weights)

    # empty list
    items = []
    i, w = n, W
    while i > 0 and w > 0:
        if table[i][w] != table[i-1][w]:
            items.append(i-1)
            w = w - weights[i-1]
        i = i - 1
    items.reverse()

    return [(weights[i], values[i]) for i in items]

# Inputs from previous Question
weights = [1, 2, 3, 6]
values = [8, 40, 30, 54]
W = 8
table, max_value = knap(weights, values, W)
items = backtrack(table, weights, values, W)

print("Max Value:", max_value)
print("Items:", items)
```

Max Value: 94
Items: [(2, 40), (6, 54)]

here the max value that does exceed weight limit of 8 is 94 and the items are 2 and 4 with weights 2 and 6 with corresponding values of 40 and 54

Q3. Dynamic programming algorithm to find longest monotonically increasing subsequence of a sequence of n integers

(11P) Step 1: Create an array X of size " n " and initialize its value to 1.

Step 2: Iterate through the input sequence.

Step 3: For each element, Find all ~~element~~ previous elements that are smaller than. ~~current~~ value and update it.

Step 4: Find max value in array X and return it

(E) Time Complexity:

Time complexity is $O(n^2)$

Since we are using loop from

1 to $n-1$

Q3.Longest monotonically increasing subsequence of a sequence of n integers.

```
def subseq(a):  
    n = len(a)  
    x = [1] * n  
    for i in range(1, n):  
        for j in range(i):  
            if a[j] < a[i]:  
                x[i] = max(x[i], x[j] + 1)  
    return max(x)
```

```
# Sequence  
a = [53, 91, 50, 1, 70, 88, 12, 19]  
print("Length of subsequence is ", subseq(a))
```

Length of subsequence is 3

(a) Recurrence Relation:

$$x[i] = \max \{ x[j] + 1 \} \quad \left. \begin{array}{l} j < i \\ a[j] < a[i] \end{array} \right\}$$

Base case is $x[i] = 1$, as subsequence ending at index i is element $a[i]$.

Q4) Dynamic Programming algorithm, to find change of n cents using Fewest coins for K denominations of coins.

(a) $D[1 \dots K]$ are Denominations
 $c[n]$ be coins required

$$c[n] = \min \{ c[n - D[i]] + 1 \} \quad \left. \begin{array}{l} \text{for } i = 1 \text{ to } K \\ D[i] \leq n. \end{array} \right\}$$

Base cases

$$c[0] = 0$$

$$c[n] = \infty \text{ for } n < 0.$$

(b) Outer loop iterates n times, for K different denominations.

\therefore Time complexity $\boxed{O(nk)}$.

(b) Step 1: Initialize array c with length $amt+1$
(from 0 to amt), $c[0] = 0$ and $c[1] = \dots = c[amt] = \infty$

Step 2: $c[0] = 0$ for the base case with 0 coins

Step 3: Compute minimum no. of coins
 $c[i] = \min(c[i], c[i-1] + 1)$

Step 4: Return $c[amt]$.

Q4. Algorithm to find the change of n cents using the fewest number of coins given any set of k different coin denominations

```
: def Coin(D,n):  
    # Initializing an array to store upto n+1  
    C=[float('inf')]*(n+1)  
    C[0] = 0  
  
    # Computing minimum number of coins  
    for i in range(1,n+1):  
        for j in range(len(D)):  
            if D[j]<=i:  
                C[i]=min(C[i],C[i-D[j]]+1)  
    return C[n]
```

```
: coins=[1, 5, 10, 25]  
amt=17  
mincoins=Coin(coins,amt)  
print(f"minimum coins required for {amt} cents are {mincoins}")
```

minimum coins required for 17 cents are 4

```
: coins=[1, 5, 10, 25]  
amt=75  
mincoins=Coin(coins,amt)  
print(f"Minimum coins required for {amt} cents are {mincoins}")
```

Minimum coins required for 75 cents are 3

```
: coins=[1, 5, 10, 25]  
amt=81  
mincoins=Coin(coins,amt)  
print(f"Minimum coins required for {amt} cents are {mincoins}")
```

Minimum coins required for 81 cents are 5