# **Distributed Systems Homework-3**

## Team 11

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# 1. Distributed K-Nearest Neighbors

## 1.1. Problem Overview

- **Objective**: Given a set of data points and query points, find the K-nearest neighbors for each query point using Euclidean distance.
- Parallelization Strategy: Leverage MPI (Message Passing Interface) to distribute the workload of query processing across multiple processes for enhanced performance.

## 1.2. Program Components

#### 1. Data Structures:

• Point Struct: Represents a 2D point with x and y coordinates.

## 2. Custom MPI Data Type:

• **createPointType Function**: Defines a custom MPI datatype for the **Point** struct, enabling the transmission of **Point** objects between MPI processes.

### 1.3. Parallel Execution Workflow

#### 1. Initialization:

• Initialize MPI, set up the environment, and determine the total number of processes and the rank of each process.

#### 2. Input Handling:

• Use the <u>inputHandler</u> function. The root process (rank 0) reads the input data, and the data is distributed to all processes using MPI\_Bcast.

## 3. Query Processing:

• Distribute the query workload among processes. Each process calculates distances, sorts them, and identifies the K-nearest neighbors using the processqueries function.

## 4. Result Gathering:

• Collect the results from all processes with MPI\_Gatherv and aggregate them in the root process.

#### 5. Output Handling:

• Use the **outputHandler** function to print the final K-nearest neighbors from the root process.

#### 6. Performance Measurement:

Measure and print the execution time using MPI\_Wtime.

#### 7. Finalization:

• Clean up the MPI environment with MPI Finalize.

## 1.4. Complexity Analysis

- Computational Complexity: O(M×(NlogN+K)), where M is the number of queries, N is the number of data points, and K is the number of nearest neighbors.
- Space complexity:
- Message complexity:
- **Communication Complexity**: Involves broadcasting input data and gathering results, dependent on the size of the data and network performance.

#### 1.5. Code Structure

#### 1. Struct Definition:

Point struct representing a 2D point.

#### 2. Euclidean Distance Function:

• calculateDistance function computes the Euclidean distance between two points.

#### 3. MPI Datatype Creation:

• createPointType defines the MPI datatype for Point.

#### 4. Input Handling:

inputHandler reads and broadcasts input data.

#### 5. Query Processing:

processqueries function processes assigned queries to find K-nearest neighbors.

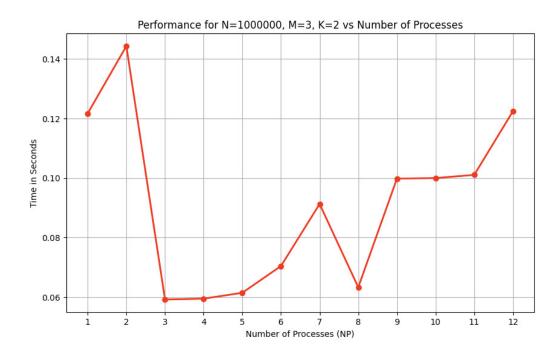
## 6. Output Handling:

• outputHandler prints the results.

#### 7. Main Function:

Coordinates initialization, processing, and finalization.

## 1.6 Performance Scaling



## 2. Julia Set

## 2.1. Problem Overview

• **Objective**: Compute the Julia set for a given complex constant c across a grid of points in the complex plane.

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• **Parallelization Strategy**: Divide the grid of points among multiple processes using MPI to compute the Julia set in parallel.

## 2.2. Program Workflow

## 1. Input Handling:

 Read parameters (grid size, max iterations, constant c) from the input file and broadcast them to all processes using MPI\_Bcast.

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## 2. Computation:

• Each process computes the Julia set for its assigned portion of the grid using the computeJuliaSet function.

## 3. Result Gathering:

• Collect results from all processes with MPI\_Gatherv and aggregate them in the root process.

## 4. Output Handling:

• Print the final results on the root process using the outputHandler function.

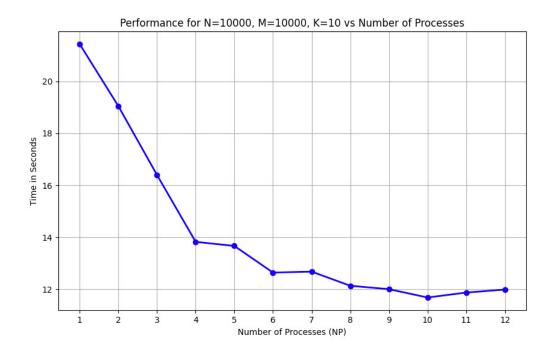
## 2.3. Complexity Analysis

- **Time Complexity**: O(N×M×K), where N×M is the grid size and K is the maximum number of iterations.
- Space Complexity: O(N×M), as the root process stores the entire grid.
- Message Complexity: O(MxP), where M is number of queries and P is number of processes

#### 2.4. Code Structure

- 1. juliaSetIterations Function:
  - Computes the number of iterations for a point to escape or confirms it belongs to the Julia set.
- 2. inputHandler Function:
  - · Reads input parameters and broadcasts them.
- 3. computeJuliaSet Function:
  - · Computes the Julia set for a subset of grid points.
- 4. outputHandler Function:
  - Gathers and prints the results.
- 5. Main Function:
  - Coordinates the execution.

## 2.5 Performance Scaling



## 3. Prefix Sum

## 3.1. Problem Overview

• **Objective**: Compute the prefix sum (cumulative sum) of an array using MPI for parallel computation.

# 3.2. Program Workflow

## 1. Input Handling:

• Read the array size and values from a file on the root process (rank 0), then broadcast the size to all processes.

## 2. Data Distribution:

 Scatter the array data among processes, with each process receiving a portion of the data.

#### 3. Local Computation:

• Each process computes the prefix sum for its local chunk.

## 4. Prefix Sum Sharing:

 Processes share cumulative sums with adjacent processes to ensure global consistency.

#### 5. Result Gathering:

Gather the local prefix sums back to the root process.

### 6. Output Handling:

• Print the final prefix sum array from the root process.

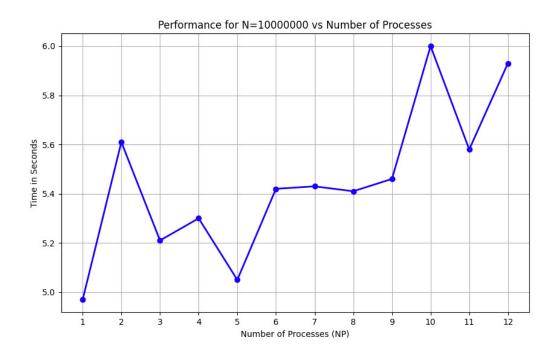
# 3.3. Complexity Analysis

- Time Complexity: O(N)
- Space Complexity: O(N)
- Message Complexity: O(1)

## 3.4. Code Structure

- 1. inputHandler Function:
  - Reads and broadcasts the input data.
- 2. scatterDataHandler Function:
  - Distributes chunks of the array to processes.
- 3. computeLocalPrefixSum Function:
  - Computes the local prefix sum.
- 4. sumsShareBetweenProcesses Function:
  - Shares cumulative sums between processes.
- 5. gatherResults Function:
  - · Gathers all prefix sums back to the root process.
- 6. outputHandler Function:
  - Prints the final prefix sum array.
- 7. Main Function:
  - Coordinates the execution.

## 3.5 Performance Scaling



# 4. Parallel Matrix Inversion using Row Reduction

## 1. Overview

This program performs matrix inversion in parallel using MPI and row reduction. It distributes the matrix across processes, leveraging parallelism for efficient computation, especially with large matrices.

# 2. Key Components and Complexities

Component	Description	Time Complexity	Space Complexity	Message Complexity
Matrix Structure ( mmatrix )	Stores matrix data and its size.	N/A	O(n²)	N/A
initialize_matrix	Initializes matrix, broadcasts size and row distribution among processes.	O(n)	O(p + n²)	O(log p) for broadcasts
input_matrix	Loads the matrix from file, sets up identity matrix, and scatters data across processes.	O(n² / p)	O(n² / p)	O(n² / p) for scatter

parallel_inverse_matrix	Performs parallel row reduction and row operations for inversion.	O(n³ / p)	O(n² / p)	O(n² / p) for gathers and broadcasts
deallocate_matrix	Frees dynamically allocated memory.	O(1)	O(n² / p)	N/A
print_result	Gathers and prints the final inverted matrix.	O(n²)	O(n²)	O(n² / p) for gather

#### 3. MPI Communication

- MPI\_Bcast: Used for broadcasting matrix size and row information.
- MPI\_Scatterv: Distributes matrix rows across processes.
- MPI\_Gather/MPI\_Gatherv: Collects parts of the inverted matrix back at the root.
- MPI\_Send/MPI\_Recv: Used during pivoting to share matrix rows between processes.

## 4. Algorithm

- **Pivoting**: The process selects the pivot row with the maximum absolute value for numerical stability.
- **Rescaling**: Rescales the pivot row so that the pivot element becomes 1.
- Row Elimination: Updates other rows by eliminating elements in the pivot column.
- **Matrix Augmentation**: The input matrix is augmented with an identity matrix to simultaneously compute the inverse.

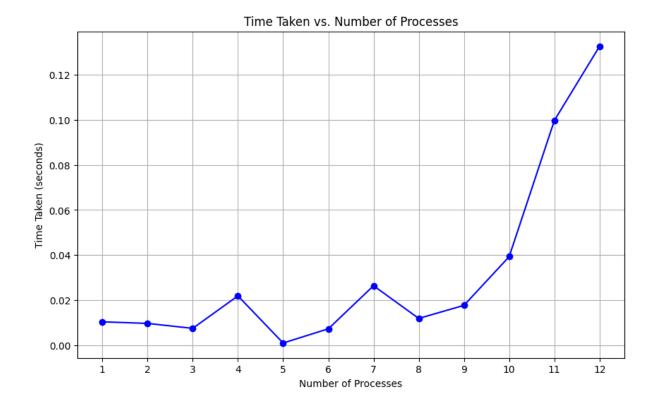
## 5. Input and Output

- **Input**: Matrix is read from command line. The file should include the matrix size followed by elements.
- Output: The inverted matrix is printed to the console by the root process.

## 6. Error Handling

• Ensures successful file opening and handles the case where no processes are available.

## 7. Performance Scaling



Increased time with increase in number of processes could potentially be due to broadcasting of matric to different processes, which causes a more severe effect as number of processes increases.

# 5. Matrix Chain Multiplication

## 1. Overview

This program solves the Matrix Chain Multiplication problem using MPI for parallel computation. It uses dynamic programming (DP) to minimize the number of scalar multiplications needed to multiply a sequence of matrices.

# 2. Key Components and Complexities

Component	Description	Time Complexity	Space Complexity	Message Complexity
Matrix Structure ( matrix_sizes )	Stores matrix chain dimensions.	N/A	O(n)	O(log p) for broadcasts
DP Table ( dp_table )	Stores the DP table for minimum multiplication costs.	N/A	O(n²)	O(log p) for broadcasts

take_input	Reads matrix dimensions from file and broadcasts to all processes.	O(n)	O(n)	O(log p) for broadcasts
<pre>init_dp_table</pre>	Initializes the DP table and broadcasts to all processes.	O(n²)	O(n²)	O(n² / p) for scatter
matrix_multiplication_mpi	Computes minimum matrix multiplication costs using parallel dynamic programming.	O(n³ / p)	O(n² / p)	O(n² / p) for reduce and broadcasts

## 3. MPI Communication

- MPI\_Bcast: Used for broadcasting matrix dimensions and DP table data.
- **MPI\_Reduce**: Gathers local minimum multiplication costs and computes the global minimum.
- MPI\_Wtime: Measures execution time.

## 4. Algorithm

- **DP Table Initialization**: Initializes the diagonal elements of the DP table to 0 (base case).
- Parallel DP Calculation: Each process computes partial costs for chain lengths, then uses MPI\_Reduce to determine the global minimum for each subproblem.
- **Result**: The minimum cost is stored in <code>dp[1][N-1]</code>, representing the optimal number of scalar multiplications.

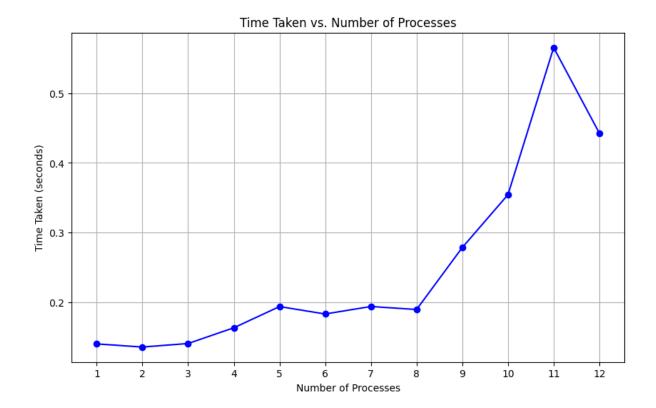
## 5. Input and Output

- Input: Matrix chain dimensions from a file input through command line.
- **Output**: The minimum number of multiplications and execution time printed to the console.

## 6. Error Handling

 Handles file opening errors and MPI initialization errors with proper termination (MPI\_Abort).

## 7. Performance Scaling



Increased time with number of processes is potentially due to the overhead of transferring matrix chain data to multiple processes.