

Prime Decomposition

$$K = K_1 \# \dots \# K_m$$

$K_i = \text{prime} \neq \text{unknot}$

$$\left(\begin{array}{l} \text{Recall: } K = A \# B \\ \Rightarrow g(K) = g(A) + g(B) \end{array} \right)$$

Theorem Suppose K is a (nontrivial, tame) knot in S^3 .

Then K admits a prime decomposition.

Moreover, if it has two decompositions

$$\begin{aligned} K &\approx K_1 \# \dots \# K_m \\ &\approx K'_1 \# \dots \# K'_n \end{aligned}$$

Then $m=n$ and $\exists \sigma \in S_n$ s.t. $K_i = K'_{\sigma(i)}$.

Theorem Suppose $K = P \# Q$ w/ P a prime knot, $Q \neq \text{unknot}$.
and $K = K_1 \# K_2$. Then P is a factor of K_1 or K_2 .

so either $K_1 \approx P \# R_1$, $Q \approx R_1 \# K_2$

or $K_2 \approx P \# R_2$, $Q \approx R_2 \# K_1$

(This implies uniqueness by induction).

Proof Let $B = \text{ball containing } P$,



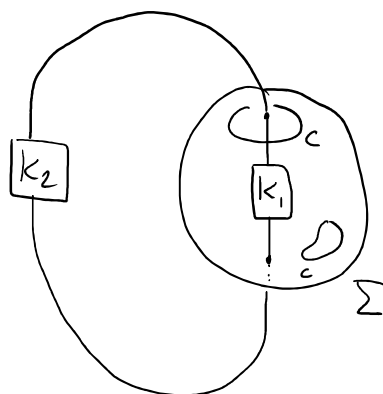
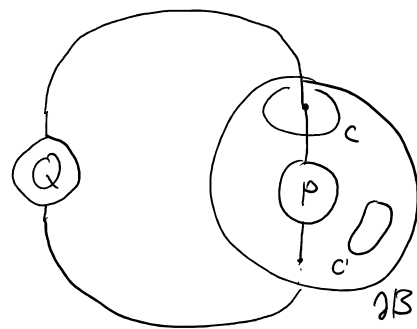
Proof Let B = ball containing P ,
and $\partial B \cap K = 2$ points.

Similarly, let Σ be separating
Sphere for $K_1 \# K_2$ - decomp.

Put ∂B and Σ in gen-pos

$\partial B \nparallel \Sigma$. That is,

$$\partial B \cap \Sigma = \bigcup_{\text{either } \partial B \text{ or } \Sigma} \text{circles on}$$



To each circle $C \subset \partial B \cap \Sigma$,

assign $l_c = |Lk(c, K)|$ (either 0 or 1, two types).

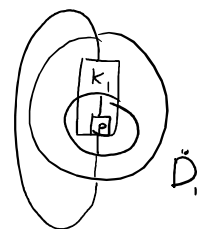
$l_c = 0 \rightsquigarrow C$ bounds disc in ∂B and in Σ

$l_c = 1 \rightsquigarrow C$ separates points in $\partial B \cap K$ and in $\Sigma \cap K$.

Suppose $\partial B \cap \Sigma = \emptyset$. $\partial B \subset S^3 - \Sigma$
 $= \dot{D}_1 \cup \dot{D}_2$

so B entirely in one of \dot{D}_1 or \dot{D}_2

if $B \subset \dot{D}_i$, P is a factor of K_i . We are done.



Strategy: get rid of circles (via surgery).

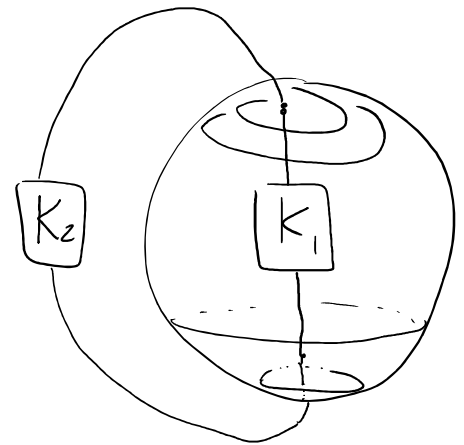
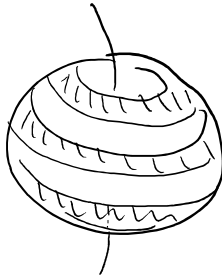
if $l_c = 0$, just shrink or grow B to $\text{depending on where the knot is}$

enclose or outclose Σ along c
(start w/ innermost).

Keep going until $l_c = 1 \forall c$.

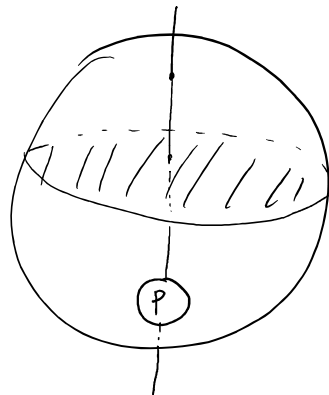
get "striped shirt":

$$\Sigma \cap B =$$



$\Sigma \cap B$ consists of at most 2 discs
each intersecting K in 1 pt, as well
as several annuli.

Suppose $\Delta \subset \Sigma \cap B$ is a disc component

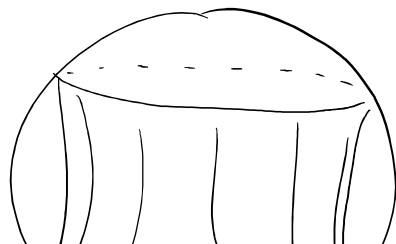


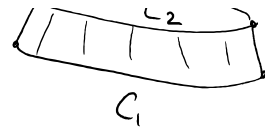
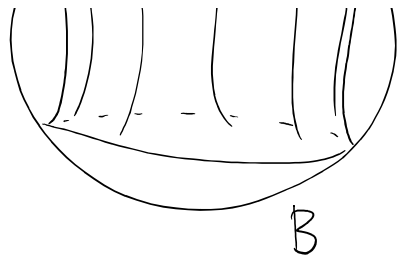
B

P is prime, so it's
either above or below Δ .

So chop off the top
or bottom to
eliminate Δ .

If A is an annular component,





Annulus splits B into $\text{Core} \cong \text{Ball}$ and full torus.

Chop out full torus to eliminate C_1 & C_2 .

(Prime knot is entirely inside the core).

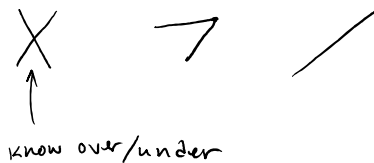
This removes all intersections so $\sum \cap B = \emptyset$.

□

Theorem $K = K_1 \# K_2$, then K is fibered
iff both K_1 & K_2 are fibered.

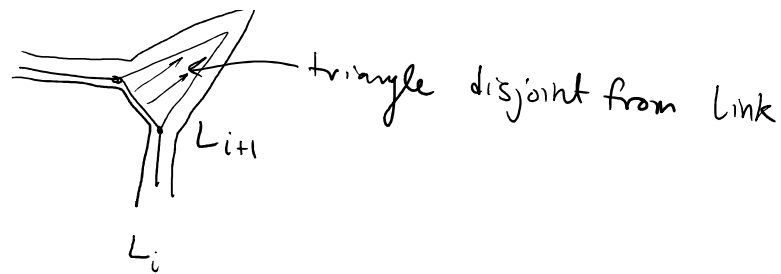
Presentations of Knots & Links

Generic proj for PL-knot



Combinatorial equivalence for PL-links





(triangle move. Can go either way).

$$L \sim L' \Leftrightarrow L = L_0 \rightarrow L_1 \rightarrow L_2 \rightarrow \dots \rightarrow L_N = L'$$

all Δ -moves

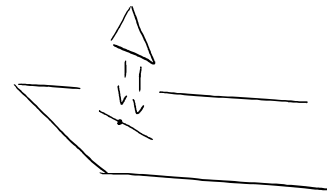
find projection for which all L_j are in general position at once.

$$L_j \xrightarrow{T} L_{j+1} \quad (\text{PL-links})$$

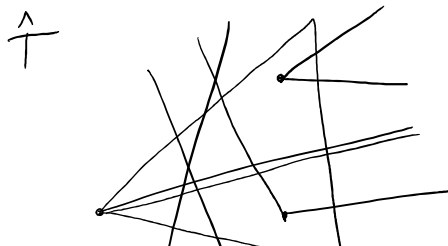
$$\hat{L}_j \xrightarrow{\hat{T}} \hat{L}_{j+1} \quad (\text{Diagrams})$$

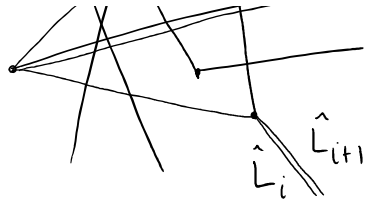
Degenerate case: \hat{T} is a line segment

$$\hat{L}_j = \hat{L}_{j+1} \quad (\text{modulo one subdivision pt}).$$



Generic case:

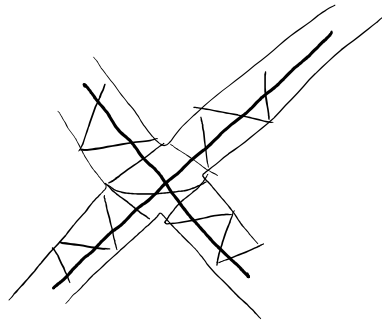
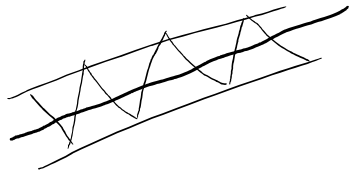




Subdivide \hat{T} by triangulation
with nice local pictures.

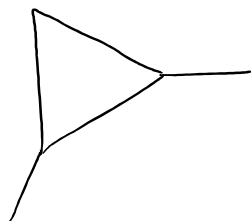
then \hat{T} -move is the same as doing all the little moves.

Strategy for a good triangulation.

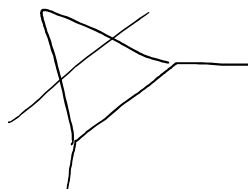


only 4 pictures:

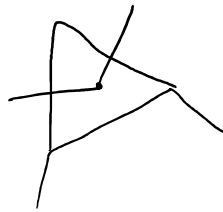
(0)



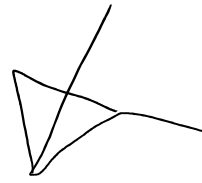
(1)



(2)



(3)



(4)

