

## Lec 8/22

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We talked about <sup>"regular"</sup>  $\forall$  induction (I)  
and complete m.l. (II)

clearly  $II \Rightarrow I$   
but  $I \Rightarrow II$  too.

$$\text{Let } Q(n) := P(1) \wedge P(2) \wedge \dots \wedge P(n)$$

$$\text{So } I \text{ on } Q \Leftrightarrow I \text{ on } P.$$

Let  $S \subseteq \mathbb{N}$ . If  $S \neq \emptyset$  then  $S$  has a least element  
(that is,  $\exists n_0 \in S$  s.t.  $\forall n \in S, n_0 \leq n$ ).

This statement is equivalent to induction (well-ordering principle).  <sup>$\nearrow W$</sup>

Proof that  $II \Rightarrow W$ : by contrapositive (prove  $S$  has no least elt  $\Rightarrow S = \emptyset$ )

(A)  $1 \notin S$  since  $1 \leq n \forall n \in \mathbb{N}$ .

(B) if  $1, 2, \dots, n \notin S$  then  $n+1 \notin S$  since o.w.  $n+1$  would be the least elt.

so  $S = \emptyset$ .

Proof that  $W \Rightarrow I$ : (by contrapositive) Let  $S = \{n \in \mathbb{N} \mid P(n) \text{ false}\}$

If  $S \neq \emptyset$  then  $\exists$  a least element  $k$ . if  $k=1$  then  $P(1)$  false, else  $P(k-1) \Rightarrow P(k)$   
since  $P(k-1)$  true but  $P(k)$  false.