Proper Let  $f: [0,1] \longrightarrow \mathbb{C}^{\times}$  be cts. Let  $w_0$  be a logarithm of f(0). Then there is a unique branch of of log f such that g(0) = w.

Pf By the mot result from monday, there is a bonnon of log f, say h. Then  $e^{h(0)} = f(0) = e^{h(0)}$  so  $\exists k \in \mathbb{Z}$  sit.  $h(0) = W_0 + 2\pi i k$ . Let  $g = h - 2\pi i k$ . Then g is chosen  $\forall t \in [0,i]$   $e^{h(t)} = e^{h(t) - 2\pi i k} = e^{h(t)} e^{-2\pi i k} = f(t) \cdot 1 = f(t)$ . Thus g is a branch of  $\log f$ . g is unique since if g is another one then  $e^{-1} = e^{-1}$  so  $g = g^{-1} + 2\pi i n$ . but  $g(0) = w_0 - g(0)$ .

## Winding numbers.

Let  $\gamma: S \to C^*$  be continuous. Define  $f: [0,i] \to C^*$  by  $f(t) = \chi(e^{2\pi i t})$  let  $\gamma: h$  be branches of leg f. Then since [0,i] is connected  $\beta: h$  an integer k sit.  $\forall t \in [0,i]$   $h(t) = g(t) + 2\pi i k$ . hence h(i) - h(o) = g(i) - g(o).

Let  $n = \frac{g(i) - g(o)}{2\pi i}$ .  $n \in \mathbb{Z}$  since  $e^{g(i)} = f(1) - \chi(e^{2\pi i}) = \chi(e^o) = f(o) = e^{f(o)}$ .

So  $g(i) - g(o) \in 2\pi i \mathbb{Z}$ . Furthermore,  $\frac{h(i) - h(o)}{2\pi i} = n$  as well, so this n is well defined as a fin of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  with respect to the origin  $\gamma: n$  denoted  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  with respect to the origin  $\gamma: n$  denoted  $\gamma: n$  of  $\gamma: n$  is called the winding  $\gamma: n$  of  $\gamma: n$  with respect to the origin  $\gamma: n$  denoted  $\gamma: n$  is  $\gamma: n$ .

If Let  $r: S' \xrightarrow{cto} (o, \infty)$ . Let  $n \in \mathbb{Z}$ . Define  $Y: S' \to C'$  by Y(z) = Z''Y(z).

Thus  $\operatorname{ind}(Y) = Y$ .

Pf Define 
$$f: [0, \Pi \to C^*]$$
 by  $f(t) = Y(e^{2\pi i t})$ , Thun  $\forall t \in [0, \Pi]$ , 
$$f(t) = e^{2\pi i t} r(e^{2\pi i t}) = e^{3(t)} \quad \text{where} \quad g(t) = \ln(r(e^{2\pi i t})) + 2n\pi i t$$

$$g \text{ 1S cts So } g \text{ is a brown of } \log f, \text{ So } \text{ ind } (\pi) = \frac{g(n - g(0))}{2\pi i} = \frac{h(r(0))^4}{2\pi i} = n.$$

Fact let  $\beta, \delta: S \xrightarrow{crs} C^{\times}$ . Then  $ind(\beta) = ind(\delta)$  iff  $\beta \simeq \gamma$  in  $C^{\times}$ .

Pf of  $(\Rightarrow)$   $\mathcal{L}_{ppose}$  ind $(\gamma) = ind(\beta)$ . Define  $\beta, \gamma: To_{1}T \longrightarrow C^{\times}$  by  $\beta(\beta) = p(e^{2\pi i\beta})$   $\gamma(\beta) = \dots$ .

let u and  $\gamma$  be branches of log  $\beta$  and log  $\delta$  respectively.  $\gamma$ -u is a branch of  $\log \frac{\gamma}{\beta}$ .

Define  $H: [0,1]^{2} \longrightarrow C^{\times}$  by  $H(s,t) = e^{(1-t)(\gamma(s)-u(s))}$ . Then H: s cts,  $H(s,0) = e^{\gamma(s)-u(s)} = \frac{\gamma(s)}{\beta(s)}$  and  $H(s,1) = e^{0} = 1$ .

Now define  $H: S^{\times}[0,1] \longrightarrow C^{\times}$  by  $H(e^{\pi i \cdot s},t) = H(s,t)$ . Since  $ind(\beta) = ind(\gamma)$ ,  $\gamma(s) - \gamma(s) = u(s) - u(s)$  so  $\gamma(s) - u(s) = \gamma(s) - u(s)$ . Thus  $\gamma(s) = u(s) = u(s)$  and  $\gamma(s) = u(s) = u(s)$ .  $\gamma(s) = \frac{\gamma(s)}{\beta(s)}$  and  $\gamma(s) = u(s) = u(s)$ . So  $\gamma(s) = u(s) = u(s)$ . Thus  $\gamma(s) = u(s) = u(s)$  and so  $\gamma(s) = \frac{\gamma(s)}{\beta(s)}$  and  $\gamma(s) = u(s) = u(s)$ .

Law we simplify this ???

Proph For all loops  $\beta$  and  $\chi$  in  $\Gamma$ , the product  $\beta\chi$  is a loop in  $\Gamma^{\chi}$  and ind  $(\beta\chi) = \operatorname{Ind}(\beta\chi) + \operatorname{Ind}(\chi)$ .

If let  $\tilde{\beta}(t) = \beta(e^{2\pi i t})$ ,  $\tilde{\gamma}(t) = \chi(e^{2\pi i t})$  as before. Let u, v be branches of  $lg\tilde{\beta}$  and  $lg\tilde{\chi}$  resp.

Let u = u + v.  $e^{iu} = e^{u}e^{v} = \tilde{\beta}\tilde{\chi}$ . So u is a bound of  $log(\tilde{\beta})$ .

So  $ind(\beta\chi) = u(1) - u(0) = (u(1) - u(0)) + (v(1) - v(0)) = ind(\beta) + ind(\chi)$ .

(on long:  $ind (\frac{1}{k}) = -ind (1)$ ,  $ind (\frac{p}{\delta}) = ind (p) - ind (p)$ .  $ind (p) = ind (d) \iff ind (\frac{p}{\delta}) = ind (\frac{\delta}{p}) = 0$ .

Proper Let Y be a loop in  $C^{\times}$ . Then ind (Y) = 0 iff J a branch of log(x)  $Y^f \Rightarrow \mathcal{D}$  eline  $\hat{Y} : L_0 : D \to C^{\times}$  by  $\hat{Y}(t) = Y(e^{2\pi i \cdot t})$ . Let  $\hat{g}$  be a branch of  $log \hat{g}$ .

Thus  $\tilde{g}(i) - \hat{g}(o) = 0$  so  $\tilde{g}(i) = \hat{g}(o)$ , define  $g: \hat{S} \to C^{\times}$  by  $g(e^{2\pi i \cdot t}) = \tilde{g}(t)$   $g: \hat{S} \to C^{\times}$  by  $g(e^{2\pi i \cdot t}) = \tilde{g}(t)$   $g: \hat{S} \to C^{\times}$  by  $g(e^{2\pi i \cdot t}) = \tilde{g}(t)$   $g: \hat{S} \to C^{\times}$  by  $g(e^{2\pi i \cdot t}) = \tilde{g}(t)$   $g: \hat{S} \to C^{\times}$  by  $g(e^{2\pi i \cdot t}) = \tilde{g}(t)$   $g: \hat{S} \to C^{\times}$  branch of log Y, say  $g: loc \hat{Y}(t) = Y(e^{2\pi i \cdot t})$ ,  $g: \hat{g}(t) = g(e^{2\pi i \cdot t})$ .

Thus  $\tilde{g}: \hat{S} \to C^{\times}$  branch of log Y, say  $g: loc \hat{Y}(t) = Y(e^{2\pi i \cdot t})$ ,  $g: \hat{g}(t) = g(e^{2\pi i \cdot t})$ .

Lorolley Let  $\beta$ , Y be loops in  $C^{\times}$ . Suppose in  $J(\beta) = m^{2}(Y)$ . Thun  $\beta \simeq Y$  in  $C^{\times}$ .

Pf in  $J(\frac{Y}{\beta}) \circ S \circ J$  a branch of log  $\frac{Y}{\beta}$  in  $C^{\times}$ , say W. Hence  $\frac{Y}{\beta} \simeq 1$  in  $C^{\times}$ .

Let  $H: \frac{Y}{\beta} \simeq 1$ . Then  $gH: Y \simeq \beta$ .