Del: A Partial function f is a subset of  $\omega \times \omega$  so that if  $(a_1b)$ ,  $(a_1c) \in f$  then b=c. If  $(n_1a) \in f$  we say f(n)=a.

Thus a total function is a special case of a partial function.

If f and g are partial Anctions then f=g iff  $\forall n \left[f(n)=g(n) \Leftrightarrow \text{either both are}\right]$ 

Det A partial function is called partial recursive if it is computed by sometM.

That is, JM s.t. Yn, if f(n) is defined, then Minit with value

f(n) as its output. otherwise, Minit (may or many not halt).

Note: a computable for is partial recursive, former. repred these is applied.

We denote partial function computed by TM Mi by Pi.

 $S_{mn}$  theorem Let f be a prof. then  $\exists$  a total new residefunction of St.  $\forall i, i, f_{\sigma(i)}(i) = f(\langle i, i \rangle)$ 

Proof given i, let M be a turny machine that, given in put i, this is composable, intuitively.

Encodes Li, is on the tape of then simulates the Machine

which computes of on this in put. Let or (i) be the index of M. I

Recursion theorem (fixed-point form): Let  $\sigma$  be a total recursive function. Then  $\exists i$  s.t.  $l_i = l_{\sigma_{ii}}$ . (there is a fixed point of  $\lambda_x$ .  $l_{\sigma_{ii}}$ ).

Proof. Consider the p.r.f.  $f(c_{i,i}) = \int_{\sigma(q_{i}(i))}^{(j)} f(i)$ . This is p.r. since were action necessary to compute it is partial recursive.

(if  $q_{i}(i)$  is undefined run  $f(c_{i,i})$ ) is too). By  $S_{mn}$  theorem,

Special case of (\*),  $I_{g(n)} = I_{\sigma(\gamma_n(n))}$ . But  $Y_n = g$ Since Mn computes g. So  $I_{g(n)} = I_{\sigma(\gamma_n(n))}$ . But  $Y_n = g$ theorem is proved  $(g(n), \frac{\log \sigma(\gamma_n(n))}{\sin \sigma(\gamma_n(n))} = \exp(\frac{1}{2} \cos \frac{\sigma(\gamma_n(n))}{\sin \sigma(\gamma_n(n))})$ .