

Review:

$P(n)$  = "if in a set of  $n$  dogs, one is a chihuahua then all are chihuahuas"

$P(1)$  is true:  $\{D_1\}$   $D_1$  is a chihuahua

induction does not hold however:  $P(n) \Rightarrow P(n+1) \forall n \geq 2$  but not for  $n=1$ .

Given  $P(n)$ :  $\{D_1, D_2, \dots, D_n\}$  are all chihuahuas if  $D_1$  is a chihuahua

then  $\{D_1, D_2, \dots, D_n, D_{n+1}\} \setminus \{D_n\}$  are all chihuahuas

$\{D_1, D_2, \dots, D_n, D_{n+1}\} \setminus \{D_{n+1}\}$  are all chihuahuas

so  $\{D_1, D_2, \dots, D_n, D_{n+1}\}$  are all chihuahuas.

$$2^n > n \quad \forall n \in \mathbb{N} = \{0, 1, 2, \dots\}$$

the induction step:

$$2 \cdot 2^n > 2 \cdot n \geq n+1 \quad \text{only works for } n \geq 1$$

a different induction step:

assume  $2^n > n$

then  $2^n \geq n+1$

$$2^{n+1} = (1+1)2^n = 2^n + 2^n > 2n+1 \geq n+1$$

$$\text{so } 2^{n+1} > n+1$$

Simple induction:

if  $P(0)$  holds

and  $P(n) \Rightarrow P(n+1) \quad \forall n$

then  $P(n)$  holds  $\forall n$

Complete induction:

if  $P(0)$  holds

and  $P(0) \wedge P(1) \wedge \dots \wedge P(n) \Rightarrow P(n+1) \quad \forall n$

then  $P(n)$  holds  $\forall n$

\* Clear that complete induction  $\Leftrightarrow$  simple induction

Conversely, simple induction for the statement  $Q(n) = P(0) \wedge P(1) \wedge \dots \wedge P(n)$  is the same as complete induction for  $P(n)$

$$Q(n) \text{ holds } \forall n \Leftrightarrow P(n) \text{ holds } \forall n$$

ex: Proof by complete induction

Generalized associative law for addition

$P(n) = a_1 + a_2 + \dots + a_n$  is independent of parenthesization

base case: Vacuously true for  $n=1, 2$   
true for  $n=3$  by PI

induction: So assume  $n \geq 3$  and we know  $P(n)$  holds

$$S = a_1 + a_2 + \dots + a_n + a_{n+1}$$

with different choices of parenthesization

$$T = a_1 + a_2 + \dots + a_n + a_{n+1}$$

want to show  $S = T$

Without loss of generality, we may assume that

$$S = (((a_1 + a_2) + a_3) + \dots) + a_n + a_{n+1}$$

left to right parenthesization

$$T = T_1 + T_2 \quad (\text{last addition performed in computation of } T \text{ w/ arbitrary parenthesization})$$

Case 1:  $T_2$  consists of a single term:  $T_2 = a_{n+1}$

then  $T_1$  is some parenthesization of  $a_1 + a_2 + \dots + a_n$

by induction,  $T_1$  is independent of parenthesization

$$T_1 =$$

$$T = T_1 + T_2 = \underbrace{(((a_1 + a_2) + a_3) + \dots) + a_n}_{T_1} + \underbrace{a_{n+1}}_{T_2} = S$$

Case 2:  $T_2$  contains more than the single term  $a_{n+1}$ .

**Case 2:**  $T_2$  contains more than the single term  $a_{n+1}$

$T_2$  has  $\leq n$  terms

by Complete induction hypothesis,

$T_2 = T_2'$  where terms are parenthesized left to right

$$= T_2'' + a_{n+1}$$

$T_1' + T_2' = T$ ,  $T_2'$  is a single term

$$\text{now } T = T_1 + T_2 = T_1 + (T_2'' + a_{n+1}) = \underbrace{(T_1 + T_2'')}_{\text{This is now in case 1 which we have already covered.}} + a_{n+1} \quad \text{by P1}$$

This completes the induction

**Illustration of induction step  $5 \rightarrow 6$  in case 2**

$$T = \underbrace{(a_1 + (a_2 + a_3))}_{T_1} + \underbrace{(a_4 + (a_5 + a_6))}_{T_2}$$

$$T_2 = \underbrace{(a_4 + a_5)}_{T_2''} + a_6 \quad \text{by case } n=3 \text{ or P1}$$

$$T = (a_1 + (a_2 + a_3)) + ((a_4 + a_5) + a_6)$$

$$(((a_1 + (a_2 + a_3)) + (a_4 + a_5)) + a_6) \quad \text{by P1}$$

$$= (((((a_1 + a_2) + a_3) + a_4) + a_5) + a_6) \quad \text{by case 1}$$

One further variation of induction:

## Well ordering principle

If  $S \subseteq \mathbb{N}$  and  $S \neq \emptyset$  then  $S$  has a least element

Prove by induction:

by contradiction: Assume  $S$  has no least element. let  $T = \{x \in \mathbb{N} : \{0, 1, \dots, x\} \subseteq \mathbb{N} \setminus S\}$

then  $0 \notin S$  because otherwise  $0$  would be the least element of  $S$ .

base case: so  $0 \in T$

induction: if  $n \in T$  then

$$\{0\} \subseteq \mathbb{N} \setminus S$$

$$\{0, 1\} \subseteq \mathbb{N} \setminus S$$

$\vdots$

$$\{0, 1, \dots, n\} \subseteq \mathbb{N} \setminus S$$

$$\text{so } \{0, 1, \dots, n\} \subseteq T$$

if  $n+1 \notin T$  then  $n+1 \in S$

but then  $n+1$  would be the least element of  $S$   
because  $0, 1, \dots, n \notin S$

$$\text{so } n+1 \in T$$

$$\text{so } S = \emptyset \text{ because } T = \mathbb{N} \setminus S = \mathbb{N}$$

**Theorem**  $\mathbb{Q}$  is dense in  $\mathbb{R}$  i.e. any open interval in  $\mathbb{R}$  contains a rational.

Proof sketch: let  $(a, b) \in \mathbb{R}$

may assume  $0 < a < b$  (why?)

choose  $q$  so  $\frac{1}{q} < b - a$

(since  $\mathbb{R}$  has no infinitesimals)

$$\text{let } S = \{n \in \mathbb{N} \mid \frac{n}{q} \geq b\}$$

$S$  is non empty

(why?)

let  $p$  be the least elem. of  $S$

$$\text{then } \frac{p-1}{q} \in (a, b)$$