# Encryption

SECURE COMMUNICATIONS & CRYPTOGRAPHY

#### Encryption

- Cryptography can be defined as the process of concealing the contents of a message from all except those who know the key.
- Encryption is the process used within cryptography to convert plaintext into cipher text
- Symmetric and asymmetric are the two primary types of encryptions
- Symmetric encryption uses a single key, whereas asymmetric uses two keys.

#### Key Terms

- Algorithm: the set of rules or mathematical formula used to encrypt and decrypt data.
- Plaintext: Cleartext that is readable

- Ciphertext:Data is scrambled and unreadable
- Encryption: the transformation of data into an unreadable format

#### Key Terms

- Cryptographic key: A key is a piece of information that controls how the cryptographic algorithm functions. It can be used to control the transformation of Plaintext to ciphertext or ciphertext to plaintext
- Symmetric Encryption: Uses the same key to encode and decode data

Asymmetric Encryption: Uses different keys for encryption and decryption. Each participant is assigned a pair of keys, what one does the other undoes.

# Symmetric and Asymmetric Differences

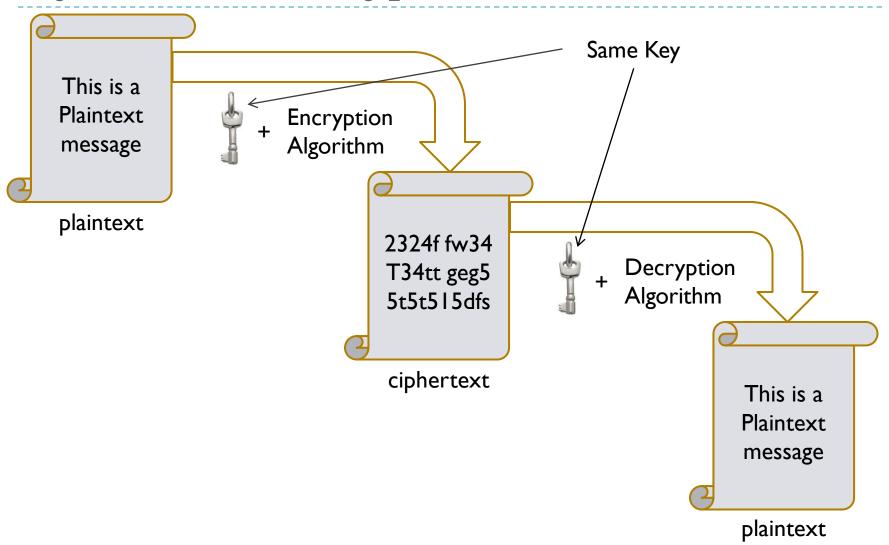
#### **Symmetric**

- Faster than Asymmetric
- Difficult key distribution
- Only provides confidentiality

#### **Asymmetric**

- Slower than Symmetric (typically hundreds to thousands times slower)
- Easy key exchange
- Can provide confidentiality and authentication

- Symmetric encryption is the older of the two forms of encryption.
- It uses a single shared secret key for encryption and decryption
- Symmetric-key algorithms can be divided into stream ciphers and block ciphers.
- Stream ciphers encrypt the bytes of the message one at a time, and block ciphers take a number of bytes and encrypt them as a single unit.



- Symmetric encryption is fast and is considered strong if large enough keys are used
- It does however have three big disadvantages
  - Key distribution:
    - We need a secure method to exchange keys
  - Key management:
    - We need a shared key for each pair of users
  - Authentication:
    - It doesn't offer us authentication

# Symmetric Algorithms

- DES Data encryption standard (still most widely used)
- Blowfish Intended as a DES replacement
- Rijndael The current AES (Advanced encryption standard)
- RC4 Rivest Cipher 4 (stream based)
- RC5 Rivest Cipher 5
- SAFER Secure and fast encryption routine

# Rivest Cipher (RC)

 RC is a general term for a family of ciphers designed by Ron Rivest (RC2, RC4, RC5, RC6)

#### RC2:

- Earliest algorithm in the series
- 64-bit block cipher that can be used with DES
- Variable key size

#### ▶ RC4:

- A stream cipher, which is faster than block mode ciphers
- ▶ The 40-bit version was originally available in WEP
- Most commonly found as 128 bit version

# Rivest Cipher (RC)

#### • RC5:

- Block based cipher
- Has a number of rounds (0-255)
- ▶ Key size range (0 2040 bits)

#### RC6:

- Variable key sizes, and rounds
- Two extra feature over RC5
  - Integer multiplication
  - ▶ Four 4-bit working registers

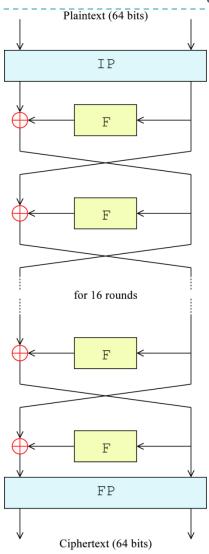
- Developed by National Bureau of Standards (NBS)
  - Now known as NIST
- Originally based on IBM algorithm called Lucifer, it was dopted as the national standard in 1976
- DES had to be recertified every five years
  - In 1993 NIST states that DES was beginning to outlive its usefulness and started looking for a replacement (AES)
- In 1998 the Electronic Frontier Foundation (EFF) managed to crack DES in 23 hours.
  - ▶ They did however use 100,000 machines

- DES is a Block Cipher and processes plaintext blocks of 64-bits into ciphertext blocks of 64-bits
- lt has a 64 bit key (56-bits are really only used)
- Because it's a symmetric algorithm it uses the same key to encrypt and decrypt
- DES performs 16 rounds
  - Each round takes the 64-bits and then uses a substitution cipher before performing a permutation on the input.

To extend the usefulness of DES, Triple DES or TDES was invented.

It is much more secure as it has a key length of 168-bits (3 \* 56-bit keys) but is three times as slow to implement

Why no double DES?

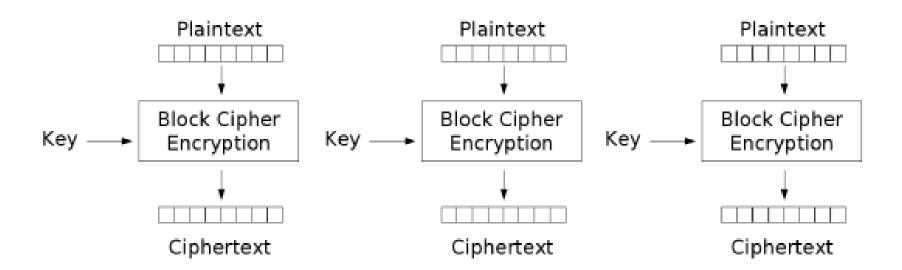


One round of the DES Cipher shown

DES repeats this 16 times

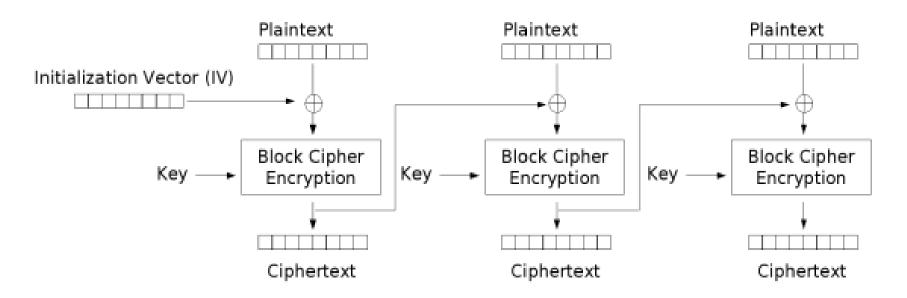
- There are four common modes or types
  - Electronic Code Book (ECB) mode
  - Cipher Block chaining (CBC) mode
  - Cipher Feedback (CFB) mode
  - Output Feedback (OFB) mode

Electronic Code Book (ECB) mode



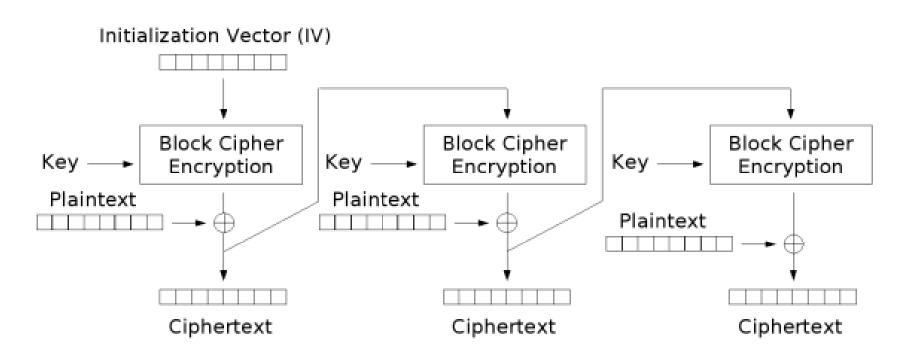
Electronic Codebook (ECB) mode encryption

Cipher Block chaining (CBC) mode



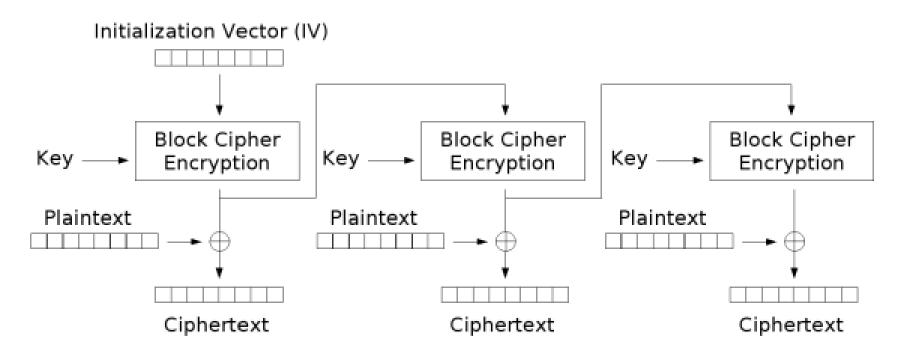
Cipher Block Chaining (CBC) mode encryption

Cipher Feedback (CFB) mode



Cipher Feedback (CFB) mode encryption

#### Output Feedback (OFB) mode



Output Feedback (OFB) mode encryption

# Advanced Encryption Standard (AES)

- The Rijndael cipher was developed by two Belgian cryptographers, Joan Daemen and Vincent Rijmen
- Chosen in 2001 By NIST as the replacement for DES
- Supports a block size of 128 and variable key lengths of 128,192 or 256 bits
- It is considered a fast, simple and robust encryption mechanism

# Advanced Encryption Standard (AES)

- lt uses four-step, parallel series of rounds
  - Byte sub:
    - ▶ Each byte is replaced by an S-box substation
  - Shift row:
    - Bytes are arranged in a rectangle and shifted
  - Mix column:
    - Matrix multiplication is performed based on the arranged rectangle
  - Add round key:
    - ▶ This rounds sub key is coded in

# Asymmetric Encryption (Aka. Public Key Cryptography)

- Unlike symmetric encryption, which uses a single shared key, asymmetric encryption uses two keys.
- They do not need a secure initial exchange as one of the keys can be made public, without any treats to the security.
- Asymmetric encryption techniques can also be used for digital signatures

- What one key does the second key undoes
- The keys are referred to as public and private keys
- The public key can be published and given to anyone, the user keeps the private key a secret
- Asymmetric encryption tends to use one way functions, but uses a trapdoor function within them to quickly reverse the operation

- An asymmetric-key cryptosystem was published in 1976 by Whitfield Diffie and Martin Hellman, disclosed a method of public-key agreement.
- This method of key exchange, which uses exponentiation in a finite field, came to be known as Diffie-Hellman key exchange.

This was the first published practical method for establishing a shared secret-key over an authenticated (but not private) communications channel without using a prior shared secret.

- Common Asymmetric functions:
  - Diffie-Hellman
  - RSA
  - **ECC**
  - ▶ EL Gamal

#### Diffie-Hellman

- Was originally developed for use as a key exchange protocol
- It is used in SSL and IPSec
- It is vulnerable to man-in-the middle attacks
- This vulnerability can be overcome if you use digital signatures
- Based on the discrete logarithm problem

#### Diffie-Hellman

#### How it works

- Alice and Bob must agree two numbers, these numbers can be sent to each other in public
  - p: some random prime number (the bigger the better)
  - g: a generator (A small prime number, usually 2,5, or 11)
- Alice and Bob now each pick some secret number
  - a: the secret number picked by Alice
  - **b**: the secret number picked by Bob
- Alice and Bob now work out  $X = g^x \mod p$  and send it to each other
  - ▶ Alice sends Bob  $A = g^a \mod p$
  - ▶ Bob sends Alice  $B = g^b \mod p$
- Alice and Bob now compute their shared secret key
  - ightharpoonup Alice computes  $s = B^a \mod p$
  - ▶ Bob computes  $s = A^b \mod p$

#### Diffie-Hellman

#### Diffie-Hellman Key Exchange



©2007 Mat-D.cor

Alice

Bob and Alice know and have the following: p = 23 (a prime number) g = 11 (a generator)

Alice chooses a secret random number a = 6

Alice receives B = 5 from Bob

Bob chooses a secret random number b = 5

Bob computes : 
$$B = g \stackrel{b}{mod} p$$
  
 $B = 11^5 \mod 23 = 5$ 

Bob receives A = 9 from Alice

$$K = 9^{5} \mod 23 = 8$$

The common secret key is: 8

N.B. We could also have written :  $K = g^{ab} mod$ 

- Invented by Rivest, Shamir and Adleman in 1977
- Based on large number factorisation problem
- RSA key sizes can be very large (RFC 2537: does limited the size to 4096 bits)

- How it works
  - Alice and Bob must agree two prime numbers, p and q
    - Alice and Bob now calculate n = pq
  - They now pick another prime number e
    - ▶ e must be less than (p-1)(q-1)
  - Alice now send her public key (n,e) to Bob
  - Bob wanting to send a message M, to Alice, sends
    - $\rightarrow$  c =  $M^e \mod n$

- Alice recovers the message by calculating
  - $M = c^d \mod n$
- We work out d as
  - b de = I mod (p-I)(q-I)
- This works becase
  - $ightharpoonup c^d = M^{ed} \mod n$

- Let's take as an example
  - p = 7 and q = 3
- Now we are going to find **n**.
  - $n = p \times q = 7 \times 3 = 21$
- Alice also needs to compute her private (or so called secret) key and the public key.
- In order to compute both keys we need to find e and d so that:
  - $e \times d \mod phi(n) = 1.$
- Let's take e = 7
- Now we need to find out **d**.
  - > 7 d mod 12.