

---

# **Linux Hid Documentation**

**The kernel development community**

**Jan 15, 2023**



## CONTENTS

<b>1</b>	<b>Care and feeding of your Human Interface Devices</b>	<b>1</b>
<b>2</b>	<b>HIDRAW - Raw Access to USB and Bluetooth Human Interface Devices</b>	<b>7</b>
<b>3</b>	<b>HID Sensors Framework</b>	<b>11</b>
<b>4</b>	<b>HID I/O Transport Drivers</b>	<b>17</b>
<b>5</b>	<b>UHID - User-space I/O driver support for HID subsystem</b>	<b>23</b>
<b>6</b>	<b>ALPS HID Touchpad Protocol</b>	<b>27</b>
<b>7</b>	<b>Intel Integrated Sensor Hub (ISH)</b>	<b>31</b>
<b>8</b>	<b>AMD Sensor Fusion Hub</b>	<b>43</b>



## **CARE AND FEEDING OF YOUR HUMAN INTERFACE DEVICES**

### **1.1 Introduction**

In addition to the normal input type HID devices, USB also uses the human interface device protocols for things that are not really human interfaces, but have similar sorts of communication needs. The two big examples for this are power devices (especially uninterruptable power supplies) and monitor control on higher end monitors.

To support these disparate requirements, the Linux USB system provides HID events to two separate interfaces: \* the input subsystem, which converts HID events into normal input device interfaces (such as keyboard, mouse and joystick) and a normalised event interface - see Documentation/input/input.rst \* the hiddev interface, which provides fairly raw HID events

The data flow for a HID event produced by a device is something like the following:

```
usb.c ---> hid-core.c ----> hid-input.c ----> [keyboard/mouse/joystick/event]
                                     |
                                     --> hiddev.c ----> POWER / MONITOR CONTROL
```

In addition, other subsystems (apart from USB) can potentially feed events into the input subsystem, but these have no effect on the HID device interface.

### **1.2 Using the HID Device Interface**

The hiddev interface is a char interface using the normal USB major, with the minor numbers starting at 96 and finishing at 111. Therefore, you need the following commands:

```
mknod /dev/usb/hiddev0 c 180 96
mknod /dev/usb/hiddev1 c 180 97
mknod /dev/usb/hiddev2 c 180 98
mknod /dev/usb/hiddev3 c 180 99
mknod /dev/usb/hiddev4 c 180 100
mknod /dev/usb/hiddev5 c 180 101
mknod /dev/usb/hiddev6 c 180 102
mknod /dev/usb/hiddev7 c 180 103
mknod /dev/usb/hiddev8 c 180 104
mknod /dev/usb/hiddev9 c 180 105
mknod /dev/usb/hiddev10 c 180 106
```

```
mknod /dev/usb/hiddev11 c 180 107
mknod /dev/usb/hiddev12 c 180 108
mknod /dev/usb/hiddev13 c 180 109
mknod /dev/usb/hiddev14 c 180 110
mknod /dev/usb/hiddev15 c 180 111
```

So you point your hiddev compliant user-space program at the correct interface for your device, and it all just works.

Assuming that you have a hiddev compliant user-space program, of course. If you need to write one, read on.

### 1.3 The HIDDEV API

This description should be read in conjunction with the HID specification, freely available from <https://www.usb.org>, and conveniently linked of <http://www.linux-usb.org>.

The hiddev API uses a read() interface, and a set of ioctl() calls.

HID devices exchange data with the host computer using data bundles called “reports”. Each report is divided into “fields”, each of which can have one or more “usages”. In the hid-core, each one of these usages has a single signed 32-bit value.

#### 1.3.1 read():

This is the event interface. When the HID device’s state changes, it performs an interrupt transfer containing a report which contains the changed value. The hid-core.c module parses the report, and returns to hiddev.c the individual usages that have changed within the report. In its basic mode, the hiddev will make these individual usage changes available to the reader using a struct hiddev\_event:

```
struct hiddev_event {
    unsigned hid;
    signed int value;
};
```

containing the HID usage identifier for the status that changed, and the value that it was changed to. Note that the structure is defined within <linux/hiddev.h>, along with some other useful #defines and structures. The HID usage identifier is a composite of the HID usage page shifted to the 16 high order bits ORed with the usage code. The behavior of the read() function can be modified using the HIDIOCSFLAG ioctl() described below.

### 1.3.2 ioctl():

This is the control interface. There are a number of controls:

#### **HIDIOCGVERSION**

- int (read)

Gets the version code out of the hiddev driver.

#### **HIDIOCAPPLICATION**

- (none)

This ioctl call returns the HID application usage associated with the HID device. The third argument to ioctl() specifies which application index to get. This is useful when the device has more than one application collection. If the index is invalid (greater or equal to the number of application collections this device has) the ioctl returns -1. You can find out beforehand how many application collections the device has from the num\_applications field from the hiddev\_devinfo structure.

#### **HIDIOCGCOLLECTIONINFO**

- struct hiddev\_collection\_info (read/write)

This returns a superset of the information above, providing not only application collections, but all the collections the device has. It also returns the level the collection lives in the hierarchy. The user passes in a hiddev\_collection\_info struct with the index field set to the index that should be returned. The ioctl fills in the other fields. If the index is larger than the last collection index, the ioctl returns -1 and sets errno to -EINVAL.

#### **HIDIOCGDEVINFO**

- struct hiddev\_devinfo (read)

Gets a hiddev\_devinfo structure which describes the device.

#### **HIDIOCGSTRING**

- struct hiddev\_string\_descriptor (read/write)

Gets a string descriptor from the device. The caller must fill in the “index” field to indicate which descriptor should be returned.

#### **HIDIOCINITREPORT**

- (none)

Instructs the kernel to retrieve all input and feature report values from the device. At this point, all the usage structures will contain current values for the device, and will maintain it as the device changes. Note that the use of this ioctl is unnecessary in general, since later kernels automatically initialize the reports from the device at attach time.

#### **HIDIOCGNAME**

- string (variable length)

Gets the device name

#### **HIDIOCGREPORT**

- struct hiddev\_report\_info (write)

Instructs the kernel to get a feature or input report from the device, in order to selectively update the usage structures (in contrast to INITREPORT).

### HIDIOCSREPORT

- struct hiddev\_report\_info (write)

Instructs the kernel to send a report to the device. This report can be filled in by the user through HIDIOCSUSAGE calls (below) to fill in individual usage values in the report before sending the report in full to the device.

### HIDIOCGREPORTINFO

- struct hiddev\_report\_info (read/write)

Fills in a hiddev\_report\_info structure for the user. The report is looked up by type (input, output or feature) and id, so these fields must be filled in by the user. The ID can be absolute – the actual report id as reported by the device – or relative – HID\_REPORT\_ID\_FIRST for the first report, and (HID\_REPORT\_ID\_NEXT | report\_id) for the next report after report\_id. Without a priori information about report ids, the right way to use this ioctl is to use the relative IDs above to enumerate the valid IDs. The ioctl returns non-zero when there is no more next ID. The real report ID is filled into the returned hiddev\_report\_info structure.

### HIDIOCGFIELDINFO

- struct hiddev\_field\_info (read/write)

Returns the field information associated with a report in a hiddev\_field\_info structure. The user must fill in report\_id and report\_type in this structure, as above. The field\_index should also be filled in, which should be a number from 0 and maxfield-1, as returned from a previous HIDIOCGREPORTINFO call.

### HIDIOCGUCODE

- struct hiddev\_usage\_ref (read/write)

Returns the usage\_code in a hiddev\_usage\_ref structure, given that its report type, report id, field index, and index within the field have already been filled into the structure.

### HIDIOCGUSAGE

- struct hiddev\_usage\_ref (read/write)

Returns the value of a usage in a hiddev\_usage\_ref structure. The usage to be retrieved can be specified as above, or the user can choose to fill in the report\_type field and specify the report\_id as HID\_REPORT\_ID\_UNKNOWN. In this case, the hiddev\_usage\_ref will be filled in with the report and field information associated with this usage if it is found.

### HIDIOCSUSAGE

- struct hiddev\_usage\_ref (write)

Sets the value of a usage in an output report. The user fills in the hiddev\_usage\_ref structure as above, but additionally fills in the value field.

### HIDIOGCOLLECTIONINDEX

- struct hiddev\_usage\_ref (write)

Returns the collection index associated with this usage. This indicates where in the collection hierarchy this usage sits.



**HIDIOCGFLAG**

- int (read)

**HIDIOCSFLAG**

- int (write)

These operations respectively inspect and replace the mode flags that influence the read() call above. The flags are as follows:

**HIDDEV\_FLAG\_UREF**

- read() calls will now return struct hiddev\_usage\_ref instead of struct hiddev\_event. This is a larger structure, but in situations where the device has more than one usage in its reports with the same usage code, this mode serves to resolve such ambiguity.

**HIDDEV\_FLAG\_REPORT**

- This flag can only be used in conjunction with HIDDEV\_FLAG\_UREF. With this flag set, when the device sends a report, a struct hiddev\_usage\_ref will be returned to read() filled in with the report\_type and report\_id, but with field\_index set to FIELD\_INDEX\_NONE. This serves as additional notification when the device has sent a report.



## **HIDRAW - RAW ACCESS TO USB AND BLUETOOTH HUMAN INTERFACE DEVICES**

The hidraw driver provides a raw interface to USB and Bluetooth Human Interface Devices (HIDs). It differs from hiddev in that reports sent and received are not parsed by the HID parser, but are sent to and received from the device unmodified.

Hidraw should be used if the userspace application knows exactly how to communicate with the hardware device, and is able to construct the HID reports manually. This is often the case when making userspace drivers for custom HID devices.

Hidraw is also useful for communicating with non-conformant HID devices which send and receive data in a way that is inconsistent with their report descriptors. Because hiddev parses reports which are sent and received through it, checking them against the device's report descriptor, such communication with these non-conformant devices is impossible using hiddev. Hidraw is the only alternative, short of writing a custom kernel driver, for these non-conformant devices.

A benefit of hidraw is that its use by userspace applications is independent of the underlying hardware type. Currently, hidraw is implemented for USB and Bluetooth. In the future, as new hardware bus types are developed which use the HID specification, hidraw will be expanded to add support for these new bus types.

Hidraw uses a dynamic major number, meaning that udev should be relied on to create hidraw device nodes. Udev will typically create the device nodes directly under /dev (eg: /dev/hidraw0). As this location is distribution- and udev rule-dependent, applications should use libudev to locate hidraw devices attached to the system. There is a tutorial on libudev with a working example at:

<http://www.signal11.us/oss/udev/>  
[https://web.archive.org/web/2019\\*/www.signal11.us](https://web.archive.org/web/2019*/www.signal11.us)

### **2.1 The HIDRAW API**

#### **2.2 read()**

read() will read a queued report received from the HID device. On USB devices, the reports read using read() are the reports sent from the device on the INTERRUPT IN endpoint. By default, read() will block until there is a report available to be read. read() can be made non-blocking, by passing the O\_NONBLOCK flag to open(), or by setting the O\_NONBLOCK flag using fcntl().

On a device which uses numbered reports, the first byte of the returned data will be the report number; the report data follows, beginning in the second byte. For devices which do not use numbered reports, the report data will begin at the first byte.

### 2.3 write()

The write() function will write a report to the device. For USB devices, if the device has an INTERRUPT OUT endpoint, the report will be sent on that endpoint. If it does not, the report will be sent over the control endpoint, using a SET\_REPORT transfer.

The first byte of the buffer passed to write() should be set to the report number. If the device does not use numbered reports, the first byte should be set to 0. The report data itself should begin at the second byte.

### 2.4 ioctl()

Hidraw supports the following ioctls:

**HIDIOCGRDESCSIZE:** Get Report Descriptor Size

This ioctl will get the size of the device's report descriptor.

**HIDIOCGRDESC:** Get Report Descriptor

This ioctl returns the device's report descriptor using a hidraw\_report\_descriptor struct. Make sure to set the size field of the hidraw\_report\_descriptor struct to the size returned from HIDIOCGRDESCSIZE.

**HIDIOCGRAWINFO:** Get Raw Info

This ioctl will return a hidraw\_devinfo struct containing the bus type, the vendor ID (VID), and product ID (PID) of the device. The bus type can be one of:

- BUS\_USB
- BUS\_HIL
- BUS\_BLUETOOTH
- BUS\_VIRTUAL

which are defined in uapi/linux/input.h.

**HIDIOCGRAWNAME(len):** Get Raw Name

This ioctl returns a string containing the vendor and product strings of the device. The returned string is Unicode, UTF-8 encoded.

**HIDIOCGRAWPHYS(len):** Get Physical Address

This ioctl returns a string representing the physical address of the device. For USB devices, the string contains the physical path to the device (the USB controller, hubs, ports, etc). For Bluetooth devices, the string contains the hardware (MAC) address of the device.

**HIDIOCSFEATURE(len):** Send a Feature Report

This ioctl will send a feature report to the device. Per the HID specification, feature reports are always sent using the control endpoint. Set the first byte of the supplied buffer to the report

number. For devices which do not use numbered reports, set the first byte to 0. The report data begins in the second byte. Make sure to set len accordingly, to one more than the length of the report (to account for the report number).

**HIDIOCGFEATURE(len):** Get a Feature Report

This ioctl will request a feature report from the device using the control endpoint. The first byte of the supplied buffer should be set to the report number of the requested report. For devices which do not use numbered reports, set the first byte to 0. The returned report buffer will contain the report number in the first byte, followed by the report data read from the device. For devices which do not use numbered reports, the report data will begin at the first byte of the returned buffer.

**HIDIOCSINPUT(len):** Send an Input Report

This ioctl will send an input report to the device, using the control endpoint. In most cases, setting an input HID report on a device is meaningless and has no effect, but some devices may choose to use this to set or reset an initial state of a report. The format of the buffer issued with this report is identical to that of HIDIOCSFEATURE.

**HIDIOCGINPUT(len):** Get an Input Report

This ioctl will request an input report from the device using the control endpoint. This is slower on most devices where a dedicated In endpoint exists for regular input reports, but allows the host to request the value of a specific report number. Typically, this is used to request the initial states of an input report of a device, before an application listens for normal reports via the regular device read() interface. The format of the buffer issued with this report is identical to that of HIDIOCGFEATURE.

**HIDIOCSOUTPUT(len):** Send an Output Report

This ioctl will send an output report to the device, using the control endpoint. This is slower on most devices where a dedicated Out endpoint exists for regular output reports, but is added for completeness. Typically, this is used to set the initial states of an output report of a device, before an application sends updates via the regular device write() interface. The format of the buffer issued with this report is identical to that of HIDIOCSFEATURE.

**HIDIOCGOUTPUT(len):** Get an Output Report

This ioctl will request an output report from the device using the control endpoint. Typically, this is used to retrieve the initial state of an output report of a device, before an application updates it as necessary either via a HIDIOCSOUTPUT request, or the regular device write() interface. The format of the buffer issued with this report is identical to that of HIDIOCGFEATURE.

## 2.5 Example

In samples/, find hid-example.c, which shows examples of read(), write(), and all the ioctls for hidraw. The code may be used by anyone for any purpose, and can serve as a starting point for developing applications using hidraw.

Document by:

Alan Ott <[alan@signal11.us](mailto:alan@signal11.us)>, Signal 11 Software



## **HID SENSORS FRAMEWORK**

HID sensor framework provides necessary interfaces to implement sensor drivers, which are connected to a sensor hub. The sensor hub is a HID device and it provides a report descriptor conforming to HID 1.12 sensor usage tables.

Description from the HID 1.12 “HID Sensor Usages” specification: “Standardization of HID usages for sensors would allow (but not require) sensor hardware vendors to provide a consistent Plug And Play interface at the USB boundary, thereby enabling some operating systems to incorporate common device drivers that could be reused between vendors, alleviating any need for the vendors to provide the drivers themselves.”

This specification describes many usage IDs, which describe the type of sensor and also the individual data fields. Each sensor can have variable number of data fields. The length and order is specified in the report descriptor. For example a part of report descriptor can look like:

```
INPUT(1) [INPUT]
..
  Field(2)
    Physical(0020.0073)
    Usage(1)
      0020.045f
    Logical Minimum(-32767)
    Logical Maximum(32767)
    Report Size(8)
    Report Count(1)
    Report Offset(16)
    Flags(Variable Absolute)
..
..
```

The report is indicating “sensor page (0x20)” contains an accelerometer-3D (0x73). This accelerometer-3D has some fields. Here for example field 2 is motion intensity (0x045f) with a logical minimum value of -32767 and logical maximum of 32767. The order of fields and length of each field is important as the input event raw data will use this format.

## 3.1 Implementation

This specification defines many different types of sensors with different sets of data fields. It is difficult to have a common input event to user space applications, for different sensors. For example an accelerometer can send X,Y and Z data, whereas an ambient light sensor can send illumination data. So the implementation has two parts:

- Core HID driver
- Individual sensor processing part (sensor drivers)

### 3.1.1 Core driver

The core driver (hid-sensor-hub) registers as a HID driver. It parses report descriptors and identifies all the sensors present. It adds an MFD device with name HID-SENSOR-xxxx (where xxxx is usage id from the specification).

For example:

HID-SENSOR-200073 is registered for an Accelerometer 3D driver.

So if any driver with this name is inserted, then the probe routine for that function will be called. So an accelerometer processing driver can register with this name and will be probed if there is an accelerometer-3D detected.

The core driver provides a set of APIs which can be used by the processing drivers to register and get events for that usage id. Also it provides parsing functions, which get and set each input/feature/output report.

### 3.1.2 Individual sensor processing part (sensor drivers)

The processing driver will use an interface provided by the core driver to parse the report and get the indexes of the fields and also can get events. This driver can use IIO interface to use the standard ABI defined for a type of sensor.

## 3.2 Core driver Interface

Callback structure:

```
Each processing driver can use this structure to set some callbacks.  
    int (*suspend)(..): Callback when HID suspend is received  
    int (*resume)(..): Callback when HID resume is received  
    int (*capture_sample)(..): Capture a sample for one of its data fields  
    int (*send_event)(..): One complete event is received which can have  
                           multiple data fields.
```

Registration functions:

```
int sensor_hub_register_callback(struct hid_sensor_hub_device *hsdev,  
                               u32 usage_id,  
                               struct hid_sensor_hub_callbacks *usage_callback):
```



Registers callbacks for a usage id. The callback functions are not allowed to sleep:

```
int sensor_hub_remove_callback(struct hid_sensor_hub_device *hsdev,  
                             u32 usage_id):
```

Removes callbacks for a usage id.

Parsing function:

```
int sensor_hub_input_get_attribute_info(struct hid_sensor_hub_device *hsdev,  
                                       u8 type,  
                                       u32 usage_id, u32 attr_usage_id,  
                                       struct hid_sensor_hub_attribute_info *info);
```

A processing driver can look for some field of interest and check if it exists in a report descriptor. If it exists it will store necessary information so that fields can be set or get individually. These indexes avoid searching every time and getting field index to get or set.

Set Feature report:

```
int sensor_hub_set_feature(struct hid_sensor_hub_device *hsdev, u32 report_id,  
                          u32 field_index, s32 value);
```

This interface is used to set a value for a field in feature report. For example if there is a field `report_interval`, which is parsed by a call to `sensor_hub_input_get_attribute_info` before, then it can directly set that individual field:

```
int sensor_hub_get_feature(struct hid_sensor_hub_device *hsdev, u32 report_id,  
                          u32 field_index, s32 *value);
```

This interface is used to get a value for a field in input report. For example if there is a field `report_interval`, which is parsed by a call to `sensor_hub_input_get_attribute_info` before, then it can directly get that individual field value:

```
int sensor_hub_input_attr_get_raw_value(struct hid_sensor_hub_device *hsdev,  
                                       u32 usage_id,  
                                       u32 attr_usage_id, u32 report_id);
```

This is used to get a particular field value through input reports. For example accelerometer wants to poll X axis value, then it can call this function with the usage id of X axis. HID sensors can provide events, so this is not necessary to poll for any field. If there is some new sample, the core driver will call registered callback function to process the sample.

---

### 3.2.1 HID Custom and generic Sensors

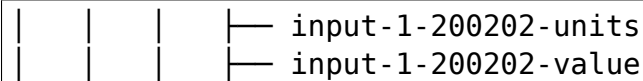
HID Sensor specification defines two special sensor usage types. Since they don't represent a standard sensor, it is not possible to define using Linux IIO type interfaces. The purpose of these sensors is to extend the functionality or provide a way to obfuscate the data being communicated by a sensor. Without knowing the mapping between the data and its encapsulated form, it is difficult for an application/driver to determine what data is being communicated by the sensor. This allows some differentiating use cases, where vendor can provide applications. Some common use cases are debug other sensors or to provide some events like keyboard attached/detached or lid open/close.

To allow application to utilize these sensors, here they are exported using sysfs attribute groups, attributes and misc device interface.

An example of this representation on sysfs:

```
/sys/devices/pci0000:00/INT33C2:00/i2c-0/i2c-INT33D1:00/0018:8086:09FA.0001/
→HID-SENSOR-2000e1.6.auto$ tree -R
```

```
.
├── enable_sensor
│   ├── feature-0-200316
│   │   ├── feature-0-200316-maximum
│   │   ├── feature-0-200316-minimum
│   │   ├── feature-0-200316-name
│   │   ├── feature-0-200316-size
│   │   ├── feature-0-200316-unit-expo
│   │   ├── feature-0-200316-units
│   │   └── feature-0-200316-value
│   ├── feature-1-200201
│   │   ├── feature-1-200201-maximum
│   │   ├── feature-1-200201-minimum
│   │   ├── feature-1-200201-name
│   │   ├── feature-1-200201-size
│   │   ├── feature-1-200201-unit-expo
│   │   ├── feature-1-200201-units
│   │   └── feature-1-200201-value
│   ├── input-0-200201
│   │   ├── input-0-200201-maximum
│   │   ├── input-0-200201-minimum
│   │   ├── input-0-200201-name
│   │   ├── input-0-200201-size
│   │   ├── input-0-200201-unit-expo
│   │   ├── input-0-200201-units
│   │   └── input-0-200201-value
│   └── input-1-200202
│       ├── input-1-200202-maximum
│       ├── input-1-200202-minimum
│       ├── input-1-200202-name
│       ├── input-1-200202-size
│       └── input-1-200202-unit-expo
```



Here there is a custom sensor with four fields: two feature and two inputs. Each field is represented by a set of attributes. All fields except the “value” are read only. The value field is a read-write field.

Example:

```
/sys/bus/platform/devices/HID-SENSOR-2000e1.6.auto/feature-0-200316$ grep -r .  
↪ *  
feature-0-200316-maximum:6  
feature-0-200316-minimum:0  
feature-0-200316-name:property-reporting-state  
feature-0-200316-size:1  
feature-0-200316-unit-expo:0  
feature-0-200316-units:25  
feature-0-200316-value:1
```

## How to enable such sensor?

By default sensor can be power gated. To enable sysfs attribute “enable” can be used:

```
$ echo 1 > enable_sensor
```

Once enabled and powered on, sensor can report value using HID reports. These reports are pushed using misc device interface in a FIFO order:

```
/dev$ tree | grep HID-SENSOR-2000e1.6.auto
| | | | └─ 10:53 -> ../HID-SENSOR-2000e1.6.auto
| └─ HID-SENSOR-2000e1.6.auto
```

Each report can be of variable length preceded by a header. This header consists of a 32-bit usage id, 64-bit time stamp and 32-bit length field of raw data.

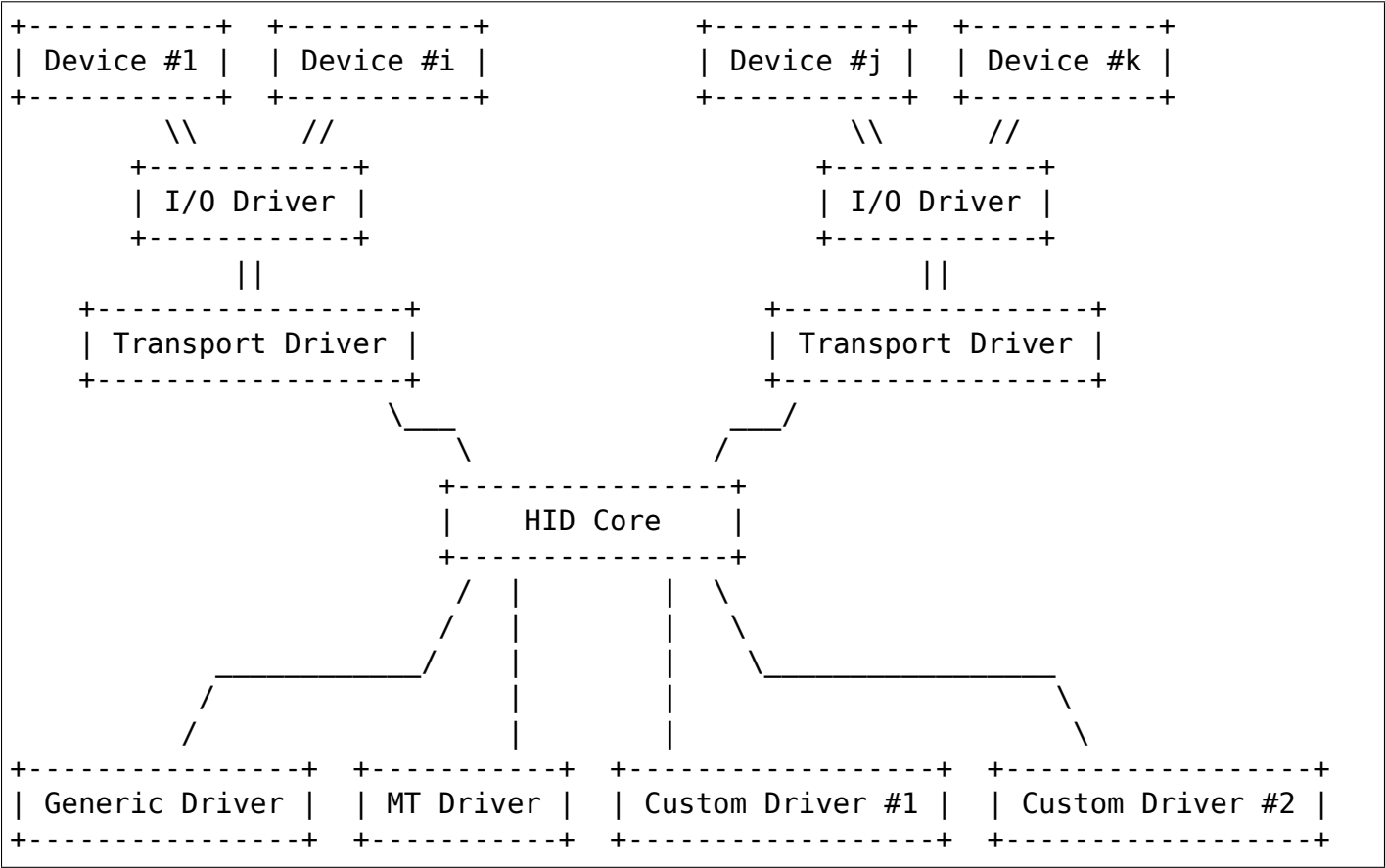


HID I/O TRANSPORT DRIVERS

The HID subsystem is independent of the underlying transport driver. Initially, only USB was supported, but other specifications adopted the HID design and provided new transport drivers. The kernel includes at least support for USB, Bluetooth, I2C and user-space I/O drivers.

4.1 1) HID Bus

The HID subsystem is designed as a bus. Any I/O subsystem may provide HID devices and register them with the HID bus. HID core then loads generic device drivers on top of it. The transport drivers are responsible for raw data transport and device setup/management. HID core is responsible for report-parsing, report interpretation and the user-space API. Device specifics and quirks are handled by all layers depending on the quirk.



Example Drivers:

- I/O: USB, I2C, Bluetooth-I2cap
- Transport: USB-HID, I2C-HID, BT-HIDP

Everything below “HID Core” is simplified in this graph as it is only of interest to HID device drivers. Transport drivers do not need to know the specifics.

### 4.1.1 1.1) Device Setup

I/O drivers normally provide hotplug detection or device enumeration APIs to the transport drivers. Transport drivers use this to find any suitable HID device. They allocate HID device objects and register them with HID core. Transport drivers are not required to register themselves with HID core. HID core is never aware of which transport drivers are available and is not interested in it. It is only interested in devices.

Transport drivers attach a constant “struct hid\_ll\_driver” object with each device. Once a device is registered with HID core, the callbacks provided via this struct are used by HID core to communicate with the device.

Transport drivers are responsible for detecting device failures and unplugging. HID core will operate a device as long as it is registered regardless of any device failures. Once transport drivers detect unplug or failure events, they must unregister the device from HID core and HID core will stop using the provided callbacks.

### 4.1.2 1.2) Transport Driver Requirements

The terms “asynchronous” and “synchronous” in this document describe the transmission behavior regarding acknowledgements. An asynchronous channel must not perform any synchronous operations like waiting for acknowledgements or verifications. Generally, HID calls operating on asynchronous channels must be running in atomic-context just fine. On the other hand, synchronous channels can be implemented by the transport driver in whatever way they like. They might just be the same as asynchronous channels, but they can also provide acknowledgement reports, automatic retransmission on failure, etc. in a blocking manner. If such functionality is required on asynchronous channels, a transport-driver must implement that via its own worker threads.

HID core requires transport drivers to follow a given design. A Transport driver must provide two bi-directional I/O channels to each HID device. These channels must not necessarily be bi-directional in the hardware itself. A transport driver might just provide 4 uni-directional channels. Or it might multiplex all four on a single physical channel. However, in this document we will describe them as two bi-directional channels as they have several properties in common.

- Interrupt Channel (intr): The intr channel is used for asynchronous data reports. No management commands or data acknowledgements are sent on this channel. Any unrequested incoming or outgoing data report must be sent on this channel and is never acknowledged by the remote side. Devices usually send their input events on this channel. Outgoing events are normally not sent via intr, except if high throughput is required.
- Control Channel (ctrl): The ctrl channel is used for synchronous requests and device management. Unrequested data input events must not be sent on this channel and are normally ignored. Instead, devices only send management events or answers to host requests on this channel. The control-channel is used for direct blocking queries to the device inde-

pendent of any events on the intr-channel. Outgoing reports are usually sent on the ctrl channel via synchronous SET\_REPORT requests.

Communication between devices and HID core is mostly done via HID reports. A report can be of one of three types:

- **INPUT Report:** Input reports provide data from device to host. This data may include button events, axis events, battery status or more. This data is generated by the device and sent to the host with or without requiring explicit requests. Devices can choose to send data continuously or only on change.
- **OUTPUT Report:** Output reports change device states. They are sent from host to device and may include LED requests, rumble requests or more. Output reports are never sent from device to host, but a host can retrieve their current state. Hosts may choose to send output reports either continuously or only on change.
- **FEATURE Report:** Feature reports are used for specific static device features and never reported spontaneously. A host can read and/or write them to access data like battery-state or device-settings. Feature reports are never sent without requests. A host must explicitly set or retrieve a feature report. This also means, feature reports are never sent on the intr channel as this channel is asynchronous.

INPUT and OUTPUT reports can be sent as pure data reports on the intr channel. For INPUT reports this is the usual operational mode. But for OUTPUT reports, this is rarely done as OUTPUT reports are normally quite scarce. But devices are free to make excessive use of asynchronous OUTPUT reports (for instance, custom HID audio speakers make great use of it).

Plain reports must not be sent on the ctrl channel, though. Instead, the ctrl channel provides synchronous GET/SET\_REPORT requests. Plain reports are only allowed on the intr channel and are the only means of data there.

- **GET\_REPORT:** A GET\_REPORT request has a report ID as payload and is sent from host to device. The device must answer with a data report for the requested report ID on the ctrl channel as a synchronous acknowledgement. Only one GET\_REPORT request can be pending for each device. This restriction is enforced by HID core as several transport drivers don't allow multiple simultaneous GET\_REPORT requests. Note that data reports which are sent as answer to a GET\_REPORT request are not handled as generic device events. That is, if a device does not operate in continuous data reporting mode, an answer to GET\_REPORT does not replace the raw data report on the intr channel on state change. GET\_REPORT is only used by custom HID device drivers to query device state. Normally, HID core caches any device state so this request is not necessary on devices that follow the HID specs except during device initialization to retrieve the current state. GET\_REPORT requests can be sent for any of the 3 report types and shall return the current report state of the device. However, OUTPUT reports as payload may be blocked by the underlying transport driver if the specification does not allow them.
- **SET\_REPORT:** A SET\_REPORT request has a report ID plus data as payload. It is sent from host to device and a device must update its current report state according to the given data. Any of the 3 report types can be used. However, INPUT reports as payload might be blocked by the underlying transport driver if the specification does not allow them. A device must answer with a synchronous acknowledgement. However, HID core does not require transport drivers to forward this acknowledgement to HID core. Same as for GET\_REPORT, only one SET\_REPORT can be pending at a time. This restriction is enforced by HID core as some transport drivers do not support multiple synchronous SET\_REPORT requests.

Other ctrl-channel requests are supported by USB-HID but are not available (or deprecated) in most other transport level specifications:

- GET/SET\_IDLE: Only used by USB-HID and I2C-HID.
- GET/SET\_PROTOCOL: Not used by HID core.
- RESET: Used by I2C-HID, not hooked up in HID core.
- SET\_POWER: Used by I2C-HID, not hooked up in HID core.

## 4.2 2) HID API

### 4.2.1 2.1) Initialization

Transport drivers normally use the following procedure to register a new device with HID core:

```
struct hid_device *hid;
int ret;

hid = hid_allocate_device();
if (IS_ERR(hid)) {
    ret = PTR_ERR(hid);
    goto err_<...>;
}

strncpy(hid->name, <device-name-src>, sizeof(hid->name));
strncpy(hid->phys, <device-phys-src>, sizeof(hid->phys));
strncpy(hid->uniq, <device-uniq-src>, sizeof(hid->uniq));

hid->ll_driver = &custom_ll_driver;
hid->bus = <device-bus>;
hid->vendor = <device-vendor>;
hid->product = <device-product>;
hid->version = <device-version>;
hid->country = <device-country>;
hid->dev.parent = <pointer-to-parent-device>;
hid->driver_data = <transport-driver-data-field>;

ret = hid_add_device(hid);
if (ret)
    goto err_<...>;
```

Once `hid_add_device()` is entered, HID core might use the callbacks provided in “`custom_ll_driver`”. Note that fields like “country” can be ignored by underlying transport-drivers if not supported.

To unregister a device, use:

```
hid_destroy_device(hid);
```

Once `hid_destroy_device()` returns, HID core will no longer make use of any driver callbacks.



## 4.2.2 2.2) hid\_ll\_driver operations

The available HID callbacks are:

```
int (*start) (struct hid_device *hdev)
```

Called from HID device drivers once they want to use the device. Transport drivers can choose to setup their device in this callback. However, normally devices are already set up before transport drivers register them to HID core so this is mostly only used by USB-HID.

```
void (*stop) (struct hid_device *hdev)
```

Called from HID device drivers once they are done with a device. Transport drivers can free any buffers and deinitialize the device. But note that `->start()` might be called again if another HID device driver is loaded on the device.

Transport drivers are free to ignore it and deinitialize devices after they destroyed them via `hid_destroy_device()`.

```
int (*open) (struct hid_device *hdev)
```

Called from HID device drivers once they are interested in data reports. Usually, while user-space didn't open any input API/etc., device drivers are not interested in device data and transport drivers can put devices asleep. However, once `->open()` is called, transport drivers must be ready for I/O. `->open()` calls are nested for each client that opens the HID device.

```
void (*close) (struct hid_device *hdev)
```

Called from HID device drivers after `->open()` was called but they are no longer interested in device reports. (Usually if user-space closed any input devices of the driver).

Transport drivers can put devices asleep and terminate any I/O of all `->open()` calls have been followed by a `->close()` call. However, `->start()` may be called again if the device driver is interested in input reports again.

```
int (*parse) (struct hid_device *hdev)
```

Called once during device setup after `->start()` has been called. Transport drivers must read the HID report-descriptor from the device and tell HID core about it via `hid_parse_report()`.

```
int (*power) (struct hid_device *hdev, int level)
```

Called by HID core to give PM hints to transport drivers. Usually this is analogical to the `->open()` and `->close()` hints and redundant.

```
void (*request) (struct hid_device *hdev, struct hid_report *report,
                int reqtype)
```

Send a HID request on the ctrl channel. "report" contains the report that should be sent and "reqtype" the request type. Request-type can be `HID_REQ_SET_REPORT` or `HID_REQ_GET_REPORT`.

This callback is optional. If not provided, HID core will assemble a raw report following the HID specs and send it via the `->raw_request()` callback. The transport driver is free to implement this asynchronously.

```
int (*wait) (struct hid_device *hdev)
```

Used by HID core before calling `->request()` again. A transport driver can use it to wait for any pending requests to complete if only one request is allowed at a time.

```
int (*raw_request) (struct hid_device *hdev, unsigned char reportnum,  
                   __u8 *buf, size_t count, unsigned char rtype,  
                   int reqtype)
```

Same as `->request()` but provides the report as raw buffer. This request shall be synchronous. A transport driver must not use `->wait()` to complete such requests. This request is mandatory and hid core will reject the device if it is missing.

```
int (*output_report) (struct hid_device *hdev, __u8 *buf, size_t len)
```

Send raw output report via intr channel. Used by some HID device drivers which require high throughput for outgoing requests on the intr channel. This must not cause SET\_REPORT calls! This must be implemented as asynchronous output report on the intr channel!

```
int (*idle) (struct hid_device *hdev, int report, int idle, int u  
            ↪ reqtype)
```

Perform SET/GET\_IDLE request. Only used by USB-HID, do not implement!

### 4.2.3 2.3) Data Path

Transport drivers are responsible of reading data from I/O devices. They must handle any I/O-related state-tracking themselves. HID core does not implement protocol handshakes or other management commands which can be required by the given HID transport specification.

Every raw data packet read from a device must be fed into HID core via `hid_input_report()`. You must specify the channel-type (intr or ctrl) and report type (input/output/feature). Under normal conditions, only input reports are provided via this API.

Responses to GET\_REPORT requests via `->request()` must also be provided via this API. Responses to `->raw_request()` are synchronous and must be intercepted by the transport driver and not passed to `hid_input_report()`. Acknowledgements to SET\_REPORT requests are not of interest to HID core.

---

Written 2013, David Herrmann <[dh.herrmann@gmail.com](mailto:dh.herrmann@gmail.com)>

## **UHID - USER-SPACE I/O DRIVER SUPPORT FOR HID SUBSYSTEM**

UHID allows user-space to implement HID transport drivers. Please see [HID I/O Transport Drivers](#) for an introduction into HID transport drivers. This document relies heavily on the definitions declared there.

With UHID, a user-space transport driver can create kernel hid-devices for each device connected to the user-space controlled bus. The UHID API defines the I/O events provided from the kernel to user-space and vice versa.

There is an example user-space application in `./samples/uhid/uhid-example.c`

### **5.1 The UHID API**

UHID is accessed through a character misc-device. The minor number is allocated dynamically so you need to rely on `udev` (or similar) to create the device node. This is `/dev/uhid` by default.

If a new device is detected by your HID I/O Driver and you want to register this device with the HID subsystem, then you need to open `/dev/uhid` once for each device you want to register. All further communication is done by `read()`'ing or `write()`'ing "struct `uhid_event`" objects. Non-blocking operations are supported by setting `O_NONBLOCK`:

```
struct uhid_event {
    __u32 type;
    union {
        struct uhid_create2_req create2;
        struct uhid_output_req output;
        struct uhid_input2_req input2;
        ...
    } u;
};
```

The "type" field contains the ID of the event. Depending on the ID different payloads are sent. You must not split a single event across multiple `read()`'s or multiple `write()`'s. A single event must always be sent as a whole. Furthermore, only a single event can be sent per `read()` or `write()`. Pending data is ignored. If you want to handle multiple events in a single syscall, then use vectored I/O with `readv()/writev()`. The "type" field defines the payload. For each type, there is a payload-structure available in the union "u" (except for empty payloads). This payload contains management and/or device data.

The first thing you should do is send a `UHID_CREATE2` event. This will register the device. UHID will respond with a `UHID_START` event. You can now start sending data to and reading

data from UHID. However, unless UHID sends the UHID\_OPEN event, the internally attached HID Device Driver has no user attached. That is, you might put your device asleep unless you receive the UHID\_OPEN event. If you receive the UHID\_OPEN event, you should start I/O. If the last user closes the HID device, you will receive a UHID\_CLOSE event. This may be followed by a UHID\_OPEN event again and so on. There is no need to perform reference-counting in user-space. That is, you will never receive multiple UHID\_OPEN events without a UHID\_CLOSE event. The HID subsystem performs ref-counting for you. You may decide to ignore UHID\_OPEN/UHID\_CLOSE, though. I/O is allowed even though the device may have no users.

If you want to send data on the interrupt channel to the HID subsystem, you send a HID\_INPUT2 event with your raw data payload. If the kernel wants to send data on the interrupt channel to the device, you will read a UHID\_OUTPUT event. Data requests on the control channel are currently limited to GET\_REPORT and SET\_REPORT (no other data reports on the control channel are defined so far). Those requests are always synchronous. That means, the kernel sends UHID\_GET\_REPORT and UHID\_SET\_REPORT events and requires you to forward them to the device on the control channel. Once the device responds, you must forward the response via UHID\_GET\_REPORT\_REPLY and UHID\_SET\_REPORT\_REPLY to the kernel. The kernel blocks internal driver-execution during such round-trips (times out after a hard-coded period).

If your device disconnects, you should send a UHID\_DESTROY event. This will unregister the device. You can now send UHID\_CREATE2 again to register a new device. If you close() the fd, the device is automatically unregistered and destroyed internally.

## 5.2 write()

write() allows you to modify the state of the device and feed input data into the kernel. The kernel will parse the event immediately and if the event ID is not supported, it will return -EOPNOTSUPP. If the payload is invalid, then -EINVAL is returned, otherwise, the amount of data that was read is returned and the request was handled successfully. O\_NONBLOCK does not affect write() as writes are always handled immediately in a non-blocking fashion. Future requests might make use of O\_NONBLOCK, though.

**UHID\_CREATE2:** This creates the internal HID device. No I/O is possible until you send this event to the kernel. The payload is of type struct uhid\_create2\_req and contains information about your device. You can start I/O now.

**UHID\_DESTROY:** This destroys the internal HID device. No further I/O will be accepted. There may still be pending messages that you can receive with read() but no further UHID\_INPUT events can be sent to the kernel. You can create a new device by sending UHID\_CREATE2 again. There is no need to reopen the character device.

**UHID\_INPUT2:** You must send UHID\_CREATE2 before sending input to the kernel! This event contains a data-payload. This is the raw data that you read from your device on the interrupt channel. The kernel will parse the HID reports.

**UHID\_GET\_REPORT\_REPLY:** If you receive a UHID\_GET\_REPORT request you must answer with this request. You must copy the “id” field from the request into the answer. Set the “err” field to 0 if no error occurred or to EIO if an I/O error occurred. If “err” is 0 then you should fill the buffer of the answer with the results of the GET\_REPORT request and set “size” correspondingly.

**UHID\_SET\_REPORT\_REPLY:** This is the SET\_REPORT equivalent of

UHID\_GET\_REPORT\_REPLY. Unlike GET\_REPORT, SET\_REPORT never returns a data buffer, therefore, it's sufficient to set the "id" and "err" fields correctly.

## 5.3 read()

read() will return a queued output report. No reaction is required to any of them but you should handle them according to your needs.

**UHID\_START:** This is sent when the HID device is started. Consider this as an answer to UHID\_CREATE2. This is always the first event that is sent. Note that this event might not be available immediately after write(UHID\_CREATE2) returns. Device drivers might require delayed setups. This event contains a payload of type `uhid_start_req`. The "dev\_flags" field describes special behaviors of a device. The following flags are defined:

- UHID\_DEV\_NUMBERED\_FEATURE\_REPORTS
- UHID\_DEV\_NUMBERED\_OUTPUT\_REPORTS
- UHID\_DEV\_NUMBERED\_INPUT\_REPORTS

Each of these flags defines whether a given report-type uses numbered reports. If numbered reports are used for a type, all messages from the kernel already have the report-number as prefix. Otherwise, no prefix is added by the kernel. For messages sent by user-space to the kernel, you must adjust the prefixes according to these flags.

**UHID\_STOP:** This is sent when the HID device is stopped. Consider this as an answer to UHID\_DESTROY.

If you didn't destroy your device via UHID\_DESTROY, but the kernel sends an UHID\_STOP event, this should usually be ignored. It means that the kernel reloaded/changed the device driver loaded on your HID device (or some other maintenance actions happened).

You can usually ignore any UHID\_STOP events safely.

**UHID\_OPEN:** This is sent when the HID device is opened. That is, the data that the HID device provides is read by some other process. You may ignore this event but it is useful for power-management. As long as you haven't received this event there is actually no other process that reads your data so there is no need to send UHID\_INPUT2 events to the kernel.

**UHID\_CLOSE:** This is sent when there are no more processes which read the HID data. It is the counterpart of UHID\_OPEN and you may as well ignore this event.

**UHID\_OUTPUT:** This is sent if the HID device driver wants to send raw data to the I/O device on the interrupt channel. You should read the payload and forward it to the device. The payload is of type "struct `uhid_output_req`". This may be received even though you haven't received UHID\_OPEN yet.

**UHID\_GET\_REPORT:** This event is sent if the kernel driver wants to perform a GET\_REPORT request on the control channel as described in the HID specs. The report-type and report-number are available in the payload. The kernel serializes GET\_REPORT requests so there will never be two in parallel. However, if you fail to respond with a UHID\_GET\_REPORT\_REPLY, the request might silently time out. Once you read a GET\_REPORT request, you shall forward it to the HID device and remember the "id" field in the payload. Once your HID device responds to the GET\_REPORT (or if it fails), you

must send a `UHID_GET_REPORT_REPLY` to the kernel with the exact same “id” as in the request. If the request already timed out, the kernel will ignore the response silently. The “id” field is never re-used, so conflicts cannot happen.

**UHID\_SET\_REPORT:** This is the `SET_REPORT` equivalent of `UHID_GET_REPORT`. On receipt, you shall send a `SET_REPORT` request to your HID device. Once it replies, you must tell the kernel about it via `UHID_SET_REPORT_REPLY`. The same restrictions as for `UHID_GET_REPORT` apply.

---

Written 2012, David Herrmann <[dh.herrmann@gmail.com](mailto:dh.herrmann@gmail.com)>

## **ALPS HID TOUCHPAD PROTOCOL**

### **6.1 Introduction**

Currently ALPS HID driver supports U1 Touchpad device.

U1 device basic information.

Vender ID	0x044E
Product ID	0x120B
Version ID	0x0121

### **6.2 HID Descriptor**

Byte	Field	Value	Notes
0	wHIDDescLength	001E	Length of HID Descriptor : 30 bytes
2	bcdVersion	0100	Compliant with Version 1.00
4	wReportDescLength	00B2	Report Descriptor is 178 Bytes (0x00B2)
6	wReportDescRegister	0002	Identifier to read Report Descriptor
8	wInputRegister	0003	Identifier to read Input Report
10	wMaxInputLength	0053	Input Report is 80 Bytes + 2
12	wOutputRegister	0000	Identifier to read Output Report
14	wMaxOutputLength	0000	No Output Reports
16	wCommandRegister	0005	Identifier for Command Register
18	wDataRegister	0006	Identifier for Data Register
20	wVendorID	044E	Vendor ID 0x044E
22	wProductID	120B	Product ID 0x120B
24	wVersionID	0121	Version 01.21
26	RESERVED	0000	RESERVED

## 6.3 Report ID

ReportID-1	(Input Reports)	(HIDUsage-Mouse) for TP&SP
ReportID-2	(Input Reports)	(HIDUsage-keyboard) for TP
ReportID-3	(Input Reports)	(Vendor Usage: Max 10 finger data) for TP
ReportID-4	(Input Reports)	(Vendor Usage: ON bit data) for GP
ReportID-5	(Feature Reports)	Feature Reports
ReportID-6	(Input Reports)	(Vendor Usage: StickPointer data) for SP
ReportID-7	(Feature Reports)	Flash update (Bootloader)

## 6.4 Data pattern

Case1	ReportID_1	TP/SP	Relative/Relative
Case2	ReportID_3 ReportID_6	TP SP	Absolute Absolute

## 6.5 Command Read/Write

To read/write to RAM, need to send a command to the device.

The command format is as below.

DataByte(SET\_REPORT)

Byte1	Command Byte
Byte2	Address - Byte 0 (LSB)
Byte3	Address - Byte 1
Byte4	Address - Byte 2
Byte5	Address - Byte 3 (MSB)
Byte6	Value Byte
Byte7	Checksum

Command Byte is read=0xD1/write=0xD2.

Address is read/write RAM address.

Value Byte is writing data when you send the write commands.

When you read RAM, there is no meaning.

DataByte(GET\_REPORT)

Byte1	Response Byte
Byte2	Address - Byte 0 (LSB)
Byte3	Address - Byte 1
Byte4	Address - Byte 2
Byte5	Address - Byte 3 (MSB)
Byte6	Value Byte
Byte7	Checksum



Read value is stored in Value Byte.

Packet Format Touchpad data byte —————

.	b7	b6	b5	b4	b3	b2	b1	b0
1	0	0	SW6	SW5	SW4	SW3	SW2	SW1
2	0	0	0	Fcv	Fn3	Fn2	Fn1	Fn0
3	Xa0_7	Xa0_6	Xa0_5	Xa0_4	Xa0_3	Xa0_2	Xa0_1	Xa0_0
4	Xa0_15	Xa0_14	Xa0_13	Xa0_12	Xa0_11	Xa0_10	Xa0_9	Xa0_8
5	Ya0_7	Ya0_6	Ya0_5	Ya0_4	Ya0_3	Ya0_2	Ya0_1	Ya0_0
6	Ya0_15	Ya0_14	Ya0_13	Ya0_12	Ya0_11	Ya0_10	Ya0_9	Ya0_8
7	LFB0	Zs0_6	Zs0_5	Zs0_4	Zs0_3	Zs0_2	Zs0_1	Zs0_0
8	Xa1_7	Xa1_6	Xa1_5	Xa1_4	Xa1_3	Xa1_2	Xa1_1	Xa1_0
9	Xa1_15	Xa1_14	Xa1_13	Xa1_12	Xa1_11	Xa1_10	Xa1_9	Xa1_8
10	Ya1_7	Ya1_6	Ya1_5	Ya1_4	Ya1_3	Ya1_2	Ya1_1	Ya1_0
11	Ya1_15	Ya1_14	Ya1_13	Ya1_12	Ya1_11	Ya1_10	Ya1_9	Ya1_8
12	LFB1	Zs1_6	Zs1_5	Zs1_4	Zs1_3	Zs1_2	Zs1_1	Zs1_0
13	Xa2_7	Xa2_6	Xa2_5	Xa2_4	Xa2_3	Xa2_2	Xa2_1	Xa2_0
14	Xa2_15	Xa2_14	Xa2_13	Xa2_12	Xa2_11	Xa2_10	Xa2_9	Xa2_8
15	Ya2_7	Ya2_6	Ya2_5	Ya2_4	Ya2_3	Ya2_2	Ya2_1	Ya2_0
16	Ya2_15	Ya2_14	Ya2_13	Ya2_12	Ya2_11	Ya2_10	Ya2_9	Ya2_8
17	LFB2	Zs2_6	Zs2_5	Zs2_4	Zs2_3	Zs2_2	Zs2_1	Zs2_0
18	Xa3_7	Xa3_6	Xa3_5	Xa3_4	Xa3_3	Xa3_2	Xa3_1	Xa3_0
19	Xa3_15	Xa3_14	Xa3_13	Xa3_12	Xa3_11	Xa3_10	Xa3_9	Xa3_8
20	Ya3_7	Ya3_6	Ya3_5	Ya3_4	Ya3_3	Ya3_2	Ya3_1	Ya3_0
21	Ya3_15	Ya3_14	Ya3_13	Ya3_12	Ya3_11	Ya3_10	Ya3_9	Ya3_8
22	LFB3	Zs3_6	Zs3_5	Zs3_4	Zs3_3	Zs3_2	Zs3_1	Zs3_0
23	Xa4_7	Xa4_6	Xa4_5	Xa4_4	Xa4_3	Xa4_2	Xa4_1	Xa4_0
24	Xa4_15	Xa4_14	Xa4_13	Xa4_12	Xa4_11	Xa4_10	Xa4_9	Xa4_8
25	Ya4_7	Ya4_6	Ya4_5	Ya4_4	Ya4_3	Ya4_2	Ya4_1	Ya4_0
26	Ya4_15	Ya4_14	Ya4_13	Ya4_12	Ya4_11	Ya4_10	Ya4_9	Ya4_8
27	LFB4	Zs4_6	Zs4_5	Zs4_4	Zs4_3	Zs4_2	Zs4_1	Zs4_0

**SW1-SW6:** SW ON/OFF status

**Xan\_15-0(16bit):** X Absolute data of the “n”th finger

**Yan\_15-0(16bit):** Y Absolute data of the “n”th finger

**Zsn\_6-0(7bit):** Operation area of the “n”th finger

## 6.6 StickPointer data byte

.	b7	b6	b5	b4	b3	b2	b1	b0
Byte1	1	1	1	0	1	SW3	SW2	SW1
Byte2	X7	X6	X5	X4	X3	X2	X1	X0
Byte3	X15	X14	X13	X12	X11	X10	X9	X8
Byte4	Y7	Y6	Y5	Y4	Y3	Y2	Y1	Y0
Byte5	Y15	Y14	Y13	Y12	Y11	Y10	Y9	Y8
Byte6	Z7	Z6	Z5	Z4	Z3	Z2	Z1	Z0
Byte7	T&P	Z14	Z13	Z12	Z11	Z10	Z9	Z8

**SW1-SW3:** SW ON/OFF status

**Xn\_15-0(16bit):** X Absolute data

**Yn\_15-0(16bit):** Y Absolute data

**Zn\_14-0(15bit):** Z

## INTEL INTEGRATED SENSOR HUB (ISH)

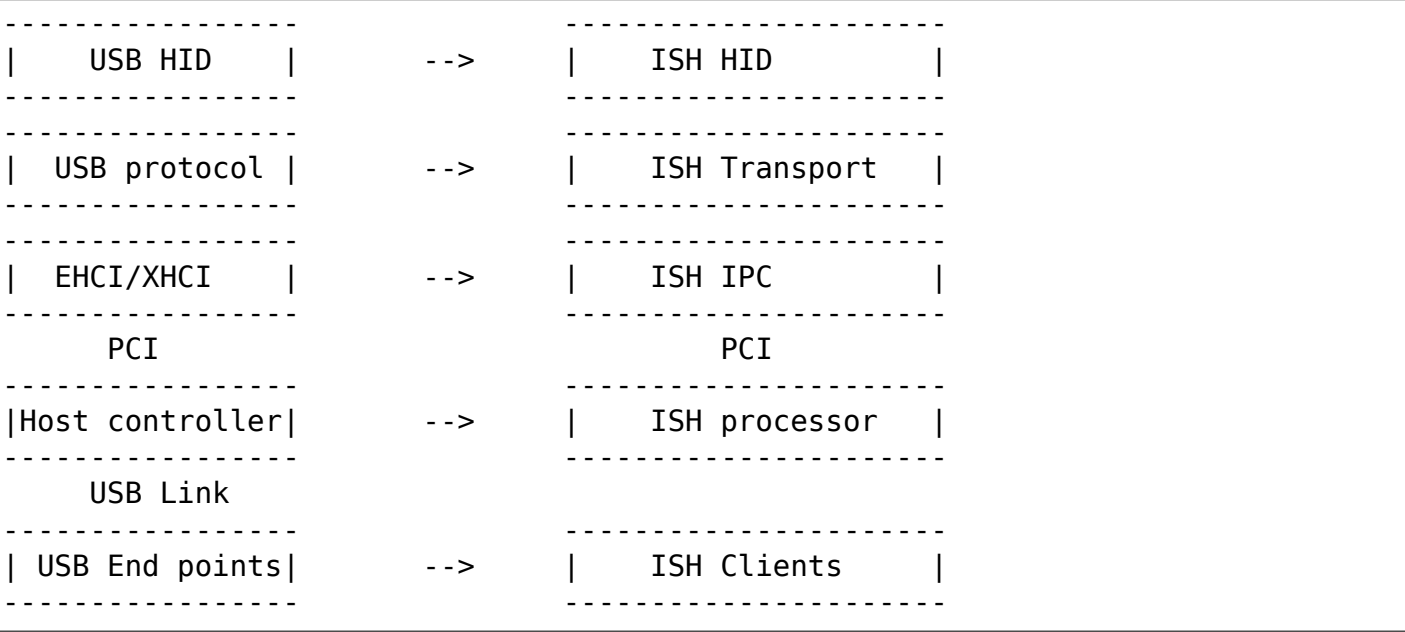
A sensor hub enables the ability to offload sensor polling and algorithm processing to a dedicated low power co-processor. This allows the core processor to go into low power modes more often, resulting in increased battery life.

There are many vendors providing external sensor hubs conforming to HID Sensor usage tables. These may be found in tablets, 2-in-1 convertible laptops and embedded products. Linux has had this support since Linux 3.9.

Intel® introduced integrated sensor hubs as a part of the SoC starting from Cherry Trail and now supported on multiple generations of CPU packages. There are many commercial devices already shipped with Integrated Sensor Hubs (ISH). These ISH also comply to HID sensor specification, but the difference is the transport protocol used for communication. The current external sensor hubs mainly use HID over I2C or USB. But ISH doesn't use either I2C or USB.

### 7.1 1. Overview

Using a analogy with a usbhid implementation, the ISH follows a similar model for a very high speed communication:



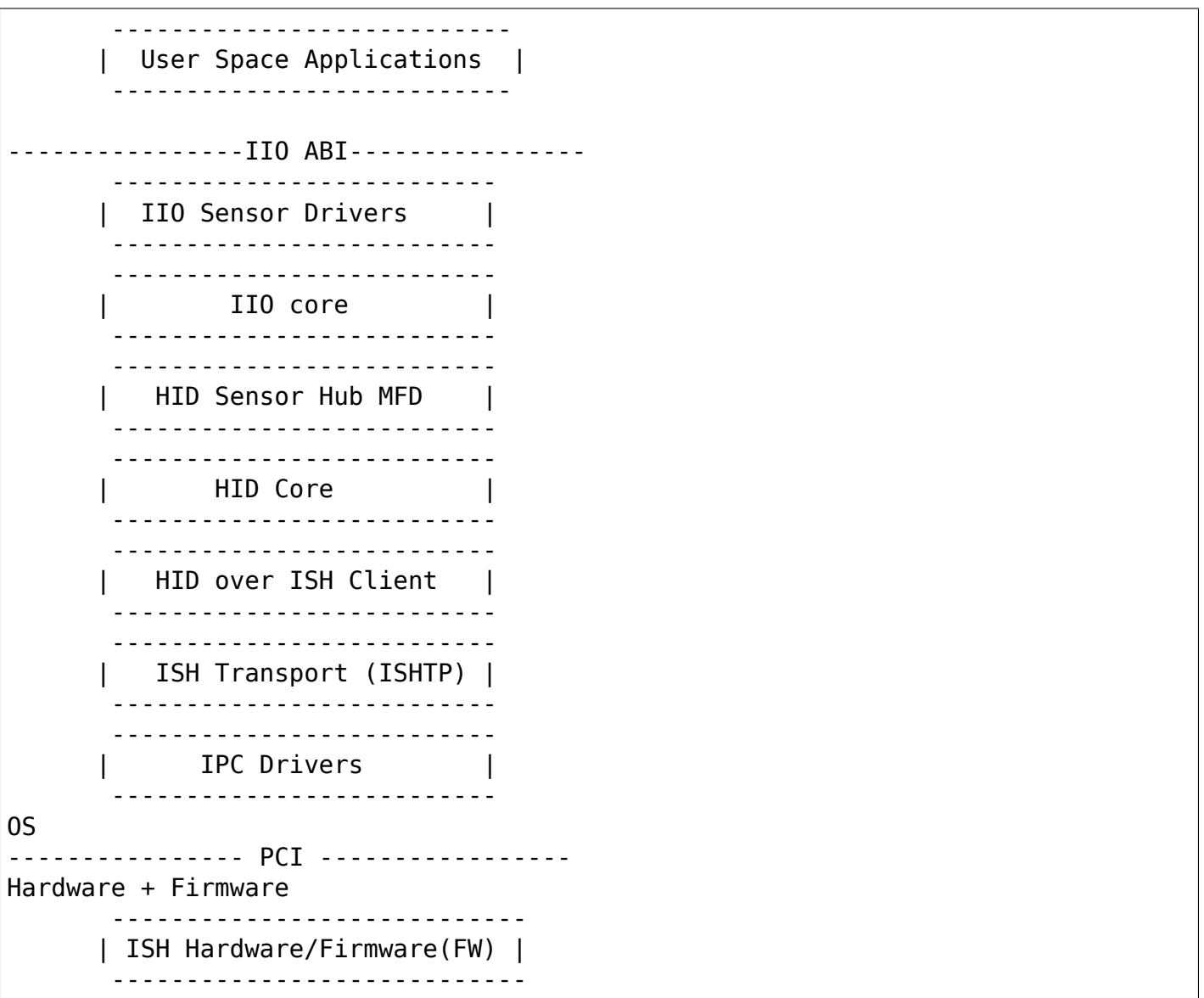
Like USB protocol provides a method for device enumeration, link management and user data encapsulation, the ISH also provides similar services. But it is very light weight tailored to

manage and communicate with ISH client applications implemented in the firmware.

The ISH allows multiple sensor management applications executing in the firmware. Like USB endpoints the messaging can be to/from a client. As part of enumeration process, these clients are identified. These clients can be simple HID sensor applications, sensor calibration applications or sensor firmware update applications.

The implementation model is similar, like USB bus, ISH transport is also implemented as a bus. Each client application executing in the ISH processor is registered as a device on this bus. The driver, which binds each device (ISH HID driver) identifies the device type and registers with the HID core.

## 7.2 2. ISH Implementation: Block Diagram



## 7.3 3. High level processing in above blocks

### 7.3.1 3.1 Hardware Interface

The ISH is exposed as “Non-VGA unclassified PCI device” to the host. The PCI product and vendor IDs are changed from different generations of processors. So the source code which enumerates drivers needs to update from generation to generation.

### 7.3.2 3.2 Inter Processor Communication (IPC) driver

Location: drivers/hid/intel-ish-hid/ipc

The IPC message uses memory mapped I/O. The registers are defined in hw-ish-regs.h.

#### 3.2.1 IPC/FW message types

There are two types of messages, one for management of link and another for messages to and from transport layers.

#### TX and RX of Transport messages

A set of memory mapped register offers support of multi-byte messages TX and RX (e.g. IPC\_REG\_ISH2HOST\_MSG, IPC\_REG\_HOST2ISH\_MSG). The IPC layer maintains internal queues to sequence messages and send them in order to the firmware. Optionally the caller can register handler to get notification of completion. A doorbell mechanism is used in messaging to trigger processing in host and client firmware side. When ISH interrupt handler is called, the ISH2HOST doorbell register is used by host drivers to determine that the interrupt is for ISH.

Each side has 32 32-bit message registers and a 32-bit doorbell. Doorbell register has the following format:

Bits 0..6: fragment length (7 bits are used)
Bits 10..13: encapsulated protocol
Bits 16..19: management command (for IPC management protocol)
Bit 31: doorbell trigger (signal H/W interrupt to the other side)
Other bits are reserved, should be 0.

#### 3.2.2 Transport layer interface

To abstract HW level IPC communication, a set of callbacks is registered. The transport layer uses them to send and receive messages. Refer to struct ishtp\_hw\_ops for callbacks.

### 7.3.3 3.3 ISH Transport layer

Location: drivers/hid/intel-ish-hid/ishtp/

#### 3.3.1 A Generic Transport Layer

The transport layer is a bi-directional protocol, which defines: - Set of commands to start, stop, connect, disconnect and flow control (see ishtp/hbm.h for details) - A flow control mechanism to avoid buffer overflows

This protocol resembles bus messages described in the following document: <http://www.intel.com/content/dam/www/public/us/en/documents/technical-specifications/dcmi-hi-1-0-spec.pdf> “Chapter 7: Bus Message Layer”

#### 3.3.2 Connection and Flow Control Mechanism

Each FW client and a protocol is identified by a UUID. In order to communicate to a FW client, a connection must be established using connect request and response bus messages. If successful, a pair (host\_client\_id and fw\_client\_id) will identify the connection.

Once connection is established, peers send each other flow control bus messages independently. Every peer may send a message only if it has received a flow-control credit before. Once it has sent a message, it may not send another one before receiving the next flow control credit. Either side can send disconnect request bus message to end communication. Also the link will be dropped if major FW reset occurs.

#### 3.3.3 Peer to Peer data transfer

Peer to Peer data transfer can happen with or without using DMA. Depending on the sensor bandwidth requirement DMA can be enabled by using module parameter ishtp\_use\_dma under intel\_ishtp.

Each side (host and FW) manages its DMA transfer memory independently. When an ISHTP client from either host or FW side wants to send something, it decides whether to send over IPC or over DMA; for each transfer the decision is independent. The sending side sends DMA\_XFER message when the message is in the respective host buffer (TX when host client sends, RX when FW client sends). The recipient of DMA message responds with DMA\_XFER\_ACK, indicating the sender that the memory region for that message may be reused.

DMA initialization is started with host sending DMA\_ALLOC\_NOTIFY bus message (that includes RX buffer) and FW responds with DMA\_ALLOC\_NOTIFY\_ACK. Additionally to DMA address communication, this sequence checks capabilities: if thw host doesn't support DMA, then it won't send DMA allocation, so FW can't send DMA; if FW doesn't support DMA then it won't respond with DMA\_ALLOC\_NOTIFY\_ACK, in which case host will not use DMA transfers. Here ISH acts as busmaster DMA controller. Hence when host sends DMA\_XFER, it's request to do host->ISH DMA transfer; when FW sends DMA\_XFER, it means that it already did DMA and the message resides at host. Thus, DMA\_XFER and DMA\_XFER\_ACK act as ownership indicators.

At initial state all outgoing memory belongs to the sender (TX to host, RX to FW), DMA\_XFER transfers ownership on the region that contains ISHTP message to the receiving side, DMA\_XFER\_ACK returns ownership to the sender. A sender need not wait for previous DMA\_XFER to be ack'ed, and may send another message as long as remaining continuous

memory in its ownership is enough. In principle, multiple DMA\_XFER and DMA\_XFER\_ACK messages may be sent at once (up to IPC MTU), thus allowing for interrupt throttling. Currently, ISH FW decides to send over DMA if ISHTP message is more than 3 IPC fragments and via IPC otherwise.

### 3.3.4 Ring Buffers

When a client initiates a connection, a ring of RX and TX buffers is allocated. The size of ring can be specified by the client. HID client sets 16 and 32 for TX and RX buffers respectively. On send request from client, the data to be sent is copied to one of the send ring buffer and scheduled to be sent using bus message protocol. These buffers are required because the FW may have not have processed the last message and may not have enough flow control credits to send. Same thing holds true on receive side and flow control is required.

### 3.3.5 Host Enumeration

The host enumeration bus command allows discovery of clients present in the FW. There can be multiple sensor clients and clients for calibration function.

To ease implementation and allow independent drivers to handle each client, this transport layer takes advantage of Linux Bus driver model. Each client is registered as device on the transport bus (ishtp bus).

Enumeration sequence of messages:

- Host sends HOST\_START\_REQ\_CMD, indicating that host ISHTP layer is up.
- FW responds with HOST\_START\_RES\_CMD
- Host sends HOST\_ENUM\_REQ\_CMD (enumerate FW clients)
- FW responds with HOST\_ENUM\_RES\_CMD that includes bitmap of available FW client IDs
- For each FW ID found in that bitmap host sends HOST\_CLIENT\_PROPERTIES\_REQ\_CMD
- FW responds with HOST\_CLIENT\_PROPERTIES\_RES\_CMD. Properties include UUID, max ISHTP message size, etc.
- Once host received properties for that last discovered client, it considers ISHTP device fully functional (and allocates DMA buffers)

## 7.3.4 3.4 HID over ISH Client

Location: drivers/hid/intel-ish-hid

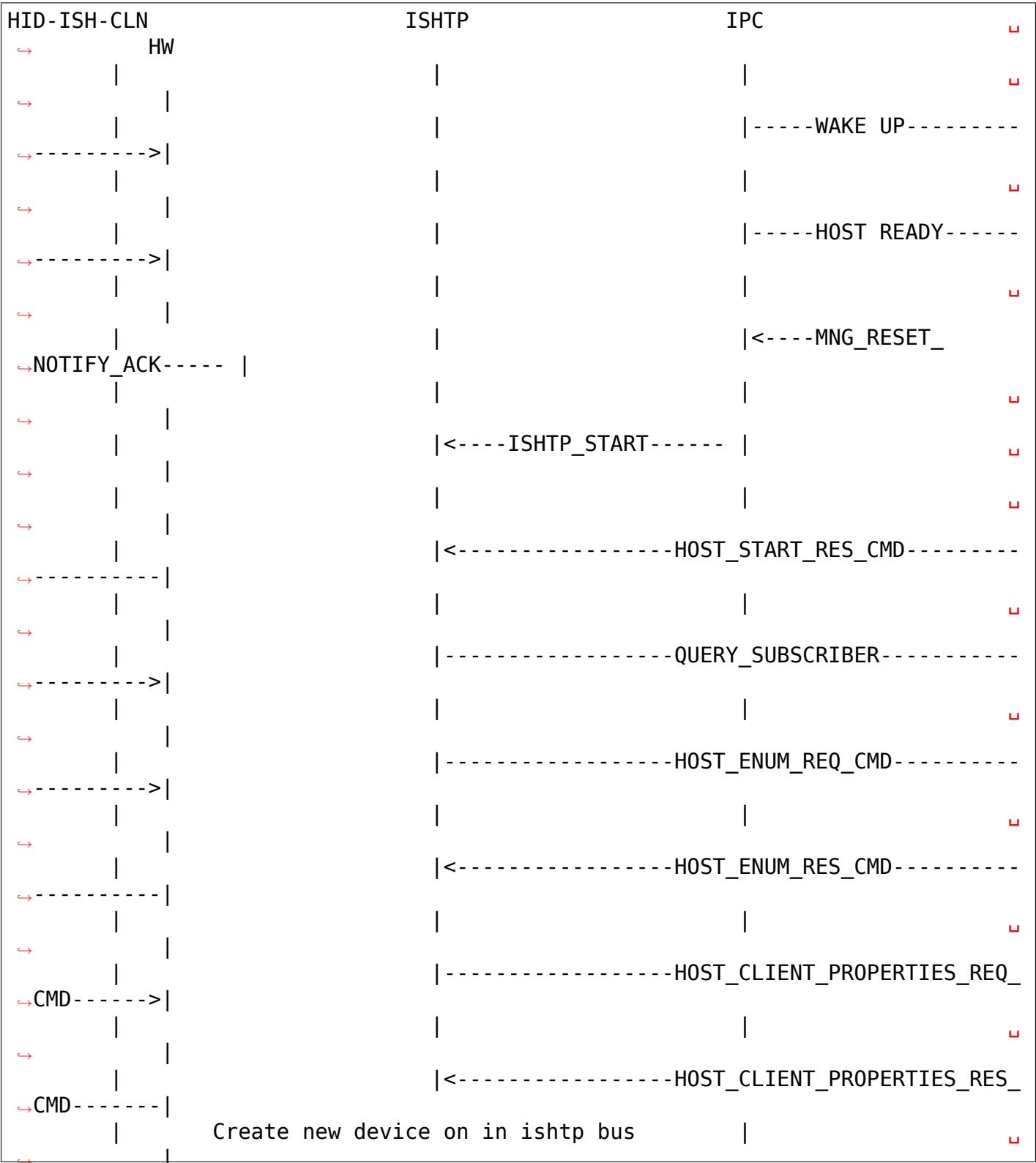
The ISHTP client driver is responsible for:

- enumerate HID devices under FW ISH client
- Get Report descriptor
- Register with HID core as a LL driver
- Process Get/Set feature request
- Get input reports

### 7.3.5 3.5 HID Sensor Hub MFD and IIO sensor drivers

The functionality in these drivers is the same as an external sensor hub. Refer to [HID Sensors Framework](#) for HID sensor Documentation/ABI/testing/sysfs-bus-iio for IIO ABIs to user space.

### 7.3.6 3.6 End to End HID transport Sequence Diagram







```

for each enumerated device
    |ishtp_cl_send(
    HOSTIF_GET_REPORT_DESCRIPTOR|-----fill ishtp_msg_hdr struct_
→write to HW-->|
    |
→    |
    |
→    |
hid_allocate_device
    |
→    |
hid_add_device
→    |
    |
→    |

```

### 7.3.7 3.7 ISH Debugging

To debug ISH, event tracing mechanism is used. To enable debug logs:

```

echo 1 > /sys/kernel/debug/tracing/events/intel_ish/enable
cat /sys/kernel/debug/tracing/trace

```

### 7.3.8 3.8 ISH IIO sysfs Example on Lenovo thinkpad Yoga 260

```

root@otcpl-ThinkPad-Yoga-260:~# tree -l /sys/bus/iio/devices/
/sys/bus/iio/devices/
├── iio:device0 -> ../../../../devices/0044:8086:22D8.0001/HID-SENSOR-200073.9.
→auto/iio:device0
│   ├── buffer
│   │   ├── enable
│   │   ├── length
│   │   └── watermark
│   ...
│   ├── in_accel_hysteresis
│   ├── in_accel_offset
│   ├── in_accel_sampling_frequency
│   ├── in_accel_scale
│   ├── in_accel_x_raw
│   ├── in_accel_y_raw
│   ├── in_accel_z_raw
│   ├── name
│   ├── scan_elements
│   │   ├── in_accel_x_en
│   │   ├── in_accel_x_index
│   │   ├── in_accel_x_type
│   │   └── in_accel_y_en

```

				in_accel_y_index
				in_accel_y_type
				in_accel_z_en
				in_accel_z_index
				in_accel_z_type
...				
			devices	
				buffer
				enable
				length
				watermark
				dev
				in_intensity_both_raw
				in_intensity_hysteresis
				in_intensity_offset
				in_intensity_sampling_frequency
				in_intensity_scale
				name
				scan_elements
				in_intensity_both_en
				in_intensity_both_index
				in_intensity_both_type
				trigger
				current_trigger
...				
				buffer
				enable
				length
				watermark
				dev
				in_magn_hysteresis
				in_magn_offset
				in_magn_sampling_frequency
				in_magn_scale
				in_magn_x_raw
				in_magn_y_raw
				in_magn_z_raw
				in_rot_from_north_magnetic_tilt_comp_raw
				in_rot_hysteresis
				in_rot_offset
				in_rot_sampling_frequency
				in_rot_scale
				name
...				
				scan_elements
				in_magn_x_en
				in_magn_x_index
				in_magn_x_type

				<ul style="list-style-type: none"> <li>└─ in_magn_y_en</li> <li>└─ in_magn_y_index</li> <li>└─ in_magn_y_type</li> <li>└─ in_magn_z_en</li> <li>└─ in_magn_z_index</li> <li>└─ in_magn_z_type</li> <li>└─ in_rot_from_north_magnetic_tilt_comp_en</li> <li>└─ in_rot_from_north_magnetic_tilt_comp_index</li> <li>└─ in_rot_from_north_magnetic_tilt_comp_type</li> </ul>
				<ul style="list-style-type: none"> <li>└─ trigger <ul style="list-style-type: none"> <li>└─ current_trigger</li> </ul> </li> </ul>
...				<ul style="list-style-type: none"> <li>└─ buffer <ul style="list-style-type: none"> <li>└─ enable</li> <li>└─ length</li> <li>└─ watermark</li> </ul> </li> <li>└─ dev</li> <li>└─ in_anglvel_hysteresis</li> <li>└─ in_anglvel_offset</li> <li>└─ in_anglvel_sampling_frequency</li> <li>└─ in_anglvel_scale</li> <li>└─ in_anglvel_x_raw</li> <li>└─ in_anglvel_y_raw</li> <li>└─ in_anglvel_z_raw</li> <li>└─ name</li> <li>└─ scan_elements <ul style="list-style-type: none"> <li>└─ in_anglvel_x_en</li> <li>└─ in_anglvel_x_index</li> <li>└─ in_anglvel_x_type</li> <li>└─ in_anglvel_y_en</li> <li>└─ in_anglvel_y_index</li> <li>└─ in_anglvel_y_type</li> <li>└─ in_anglvel_z_en</li> <li>└─ in_anglvel_z_index</li> <li>└─ in_anglvel_z_type</li> </ul> </li> <li>└─ trigger <ul style="list-style-type: none"> <li>└─ current_trigger</li> </ul> </li> </ul>
...				<ul style="list-style-type: none"> <li>└─ buffer <ul style="list-style-type: none"> <li>└─ enable</li> <li>└─ length</li> <li>└─ watermark</li> </ul> </li> <li>└─ dev</li> <li>└─ in_anglvel_hysteresis</li> <li>└─ in_anglvel_offset</li> <li>└─ in_anglvel_sampling_frequency</li> <li>└─ in_anglvel_scale</li> <li>└─ in_anglvel_x_raw</li> </ul>

```

|
|
|
|
|— in_anglvel_y_raw
|— in_anglvel_z_raw
|— name
|— scan_elements
|   |— in_anglvel_x_en
|   |— in_anglvel_x_index
|   |— in_anglvel_x_type
|   |— in_anglvel_y_en
|   |— in_anglvel_y_index
|   |— in_anglvel_y_type
|   |— in_anglvel_z_en
|   |— in_anglvel_z_index
|   |— in_anglvel_z_type
|— trigger
|   |— current_trigger
...

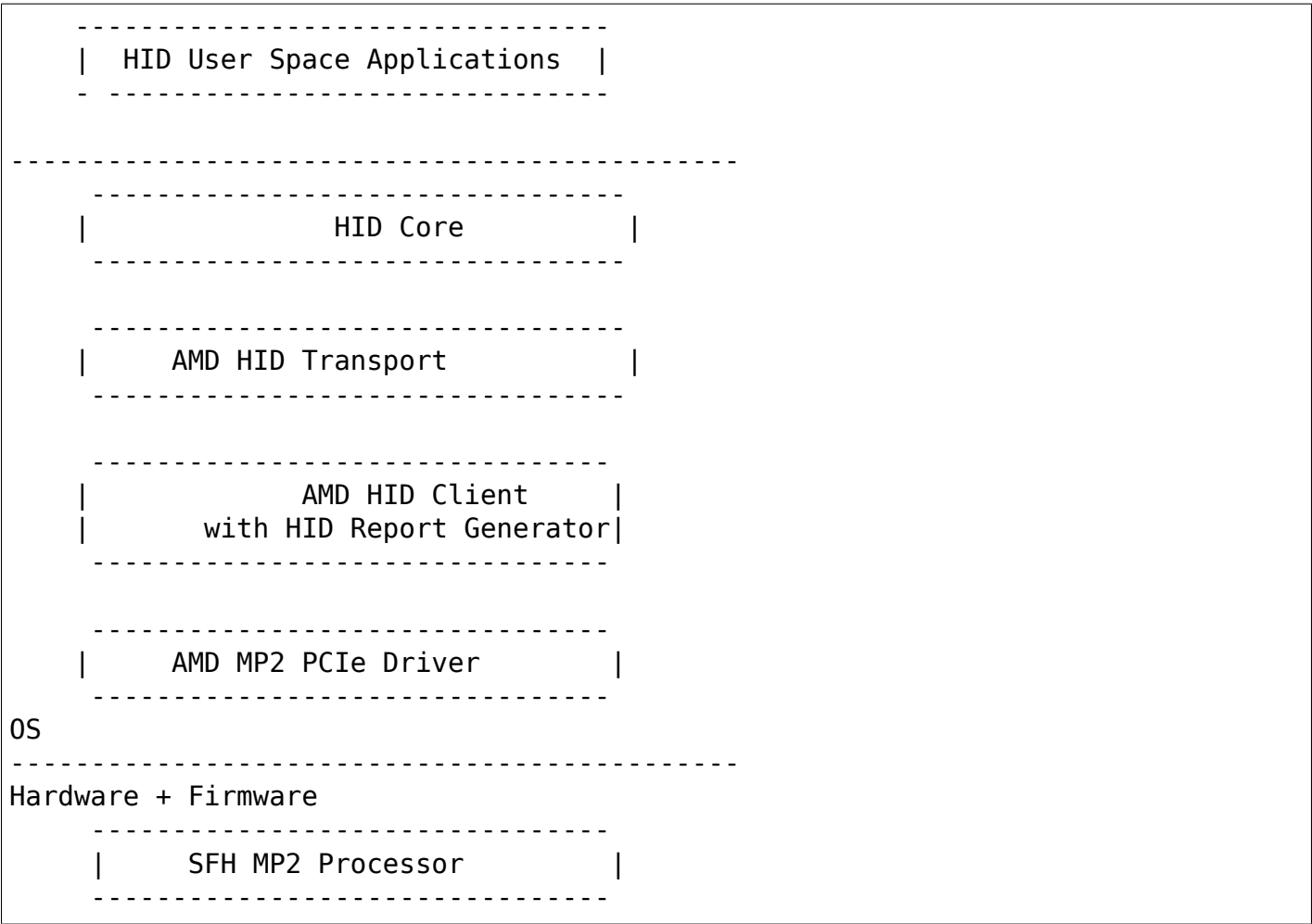
```



AMD SENSOR FUSION HUB

AMD Sensor Fusion Hub (SFH) is part of an SOC starting from Ryzen-based platforms. The solution is working well on several OEM products. AMD SFH uses HID over PCIe bus. In terms of architecture it resembles ISH, however the major difference is all the HID reports are generated as part of the kernel driver.

8.1 Block Diagram



## 8.2 AMD HID Transport Layer

AMD SFH transport is also implemented as a bus. Each client application executing in the AMD MP2 is registered as a device on this bus. Here, MP2 is an ARM core connected to x86 for processing sensor data. The layer, which binds each device (AMD SFH HID driver) identifies the device type and registers with the HID core. Transport layer attaches a constant “struct hid\_ll\_driver” object with each device. Once a device is registered with HID core, the callbacks provided via this struct are used by HID core to communicate with the device. AMD HID Transport layer implements the synchronous calls.

## 8.3 AMD HID Client Layer

This layer is responsible to implement HID requests and descriptors. As firmware is OS agnostic, HID client layer fills the HID request structure and descriptors. HID client layer is complex as it is interface between MP2 PCIe layer and HID. HID client layer initializes the MP2 PCIe layer and holds the instance of MP2 layer. It identifies the number of sensors connected using MP2-PCIe layer. Based on that allocates the DRAM address for each and every sensor and passes it to MP2-PCIe driver. On enumeration of each sensor, client layer fills the HID Descriptor structure and HID input report structure. HID Feature report structure is optional. The report descriptor structure varies from sensor to sensor.

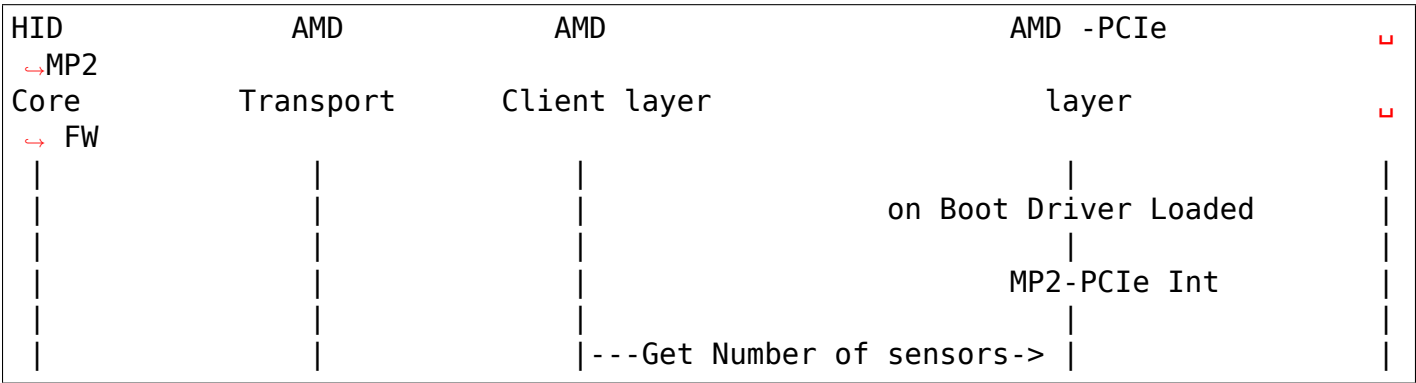
## 8.4 AMD MP2 PCIe layer

MP2 PCIe Layer is responsible for making all transactions with the firmware over PCIe. The connection establishment between firmware and PCIe happens here.

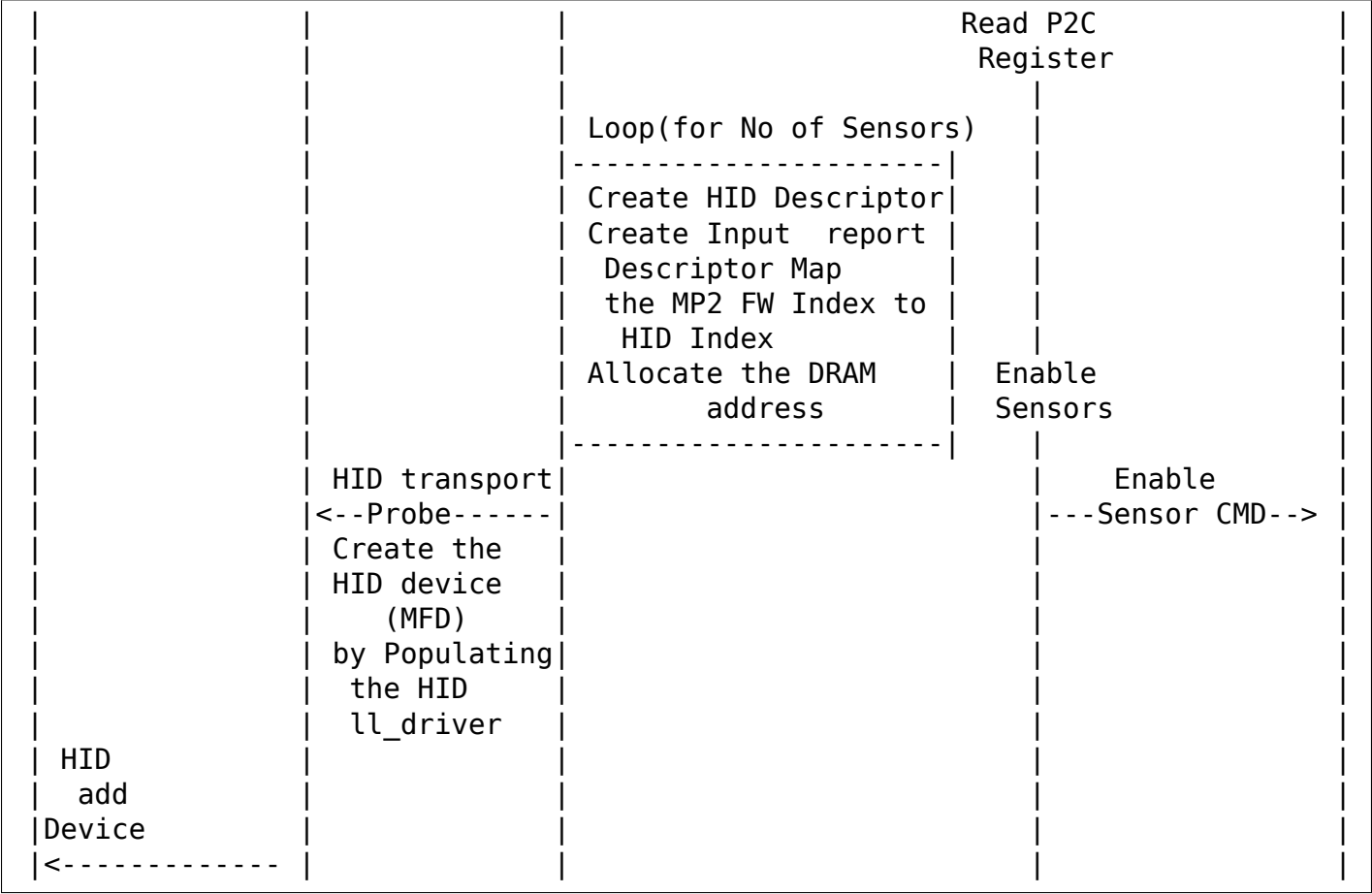
The communication between X86 and MP2 is split into three parts. 1. Command transfer via the C2P mailbox registers. 2. Data transfer via DRAM. 3. Supported sensor info via P2C registers.

Commands are sent to MP2 using C2P Mailbox registers. Writing into C2P Message registers generates interrupt to MP2. The client layer allocates the physical memory and the same is sent to MP2 via the PCI layer. MP2 firmware writes the command output to the access DRAM memory which the client layer has allocated. Firmware always writes minimum of 32 bytes into DRAM. So as a protocol driver shall allocate minimum of 32 bytes DRAM space.

## 8.5 Enumeration and Probing flow







8.6 Data Flow from Application to the AMD SFH Driver

