## proof by contradiction.

## Greedy Algorithm for scheduling.

Lay 2 jobs i, i mith pronty/meights wi, wi and time taken 1,° & 4,° respectively. So based on Greedy algorithm, me define a ratio w/1 for every fob. Higher me ratio, higher/earlier tre job mill be enemed.

Day Willig Willig so job i mouded be encented

first.

 $\frac{\omega_i^{\circ}}{l_i^{\circ}} > \frac{\omega_j^{\circ}}{l_j^{\circ}} \rightarrow 0$  (me call it  $\sigma^{\bullet}$ - Hgo soln).

now lets arbitarily say a solution with j being \*
enewted first and i later is what we now call epsimal or 5.

If we prove of is not appirmal, then we can say Algo soln- o was better, and no better solution is available.

so by surfering It!.

the benefit is for job i cout is for job i

the benefit is for job i.

Benefit = Wi x 4°

w, '4' > w, '4' D=> \( \frac{\omega\_i^0}{40} > \frac{\omega\_i^0}{40} \) cost > Benefit

so me soln ot has more cost than Benefit 10 eur contradict 5\*, and conclude mat o is

better colution.

QED.