Arrays : Carry Forward

Dec 1, 2023

AGENDA

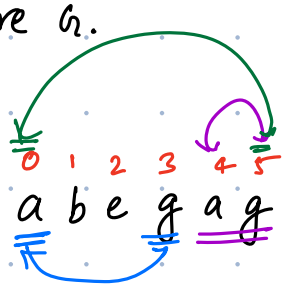
- Carry forward concept with 2 interesting questions
- Brief on subarrays

Q. **Count pairs 'AG'**

Given a string 's' of lowercase characters, return the count of (i,j) such that $i<j$ and $s[i]=='a'$ and $s[j]=='g'$.

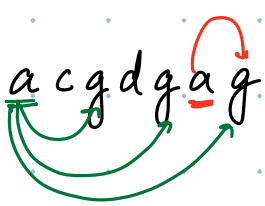
e.g. → 0 1 2 3 4 5
a b e g a g

Simplify : Count no. of AG pairs such that A comes before G.

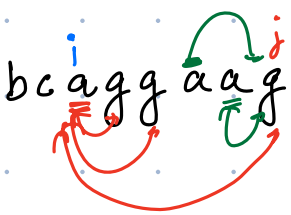


i	j		
0	3	i<j	s[i]==a s[j]==g
4	5	i<j	s[i]==a s[j]==g
0	5	i<j	s[i]==a s[j]==g

3 pairs

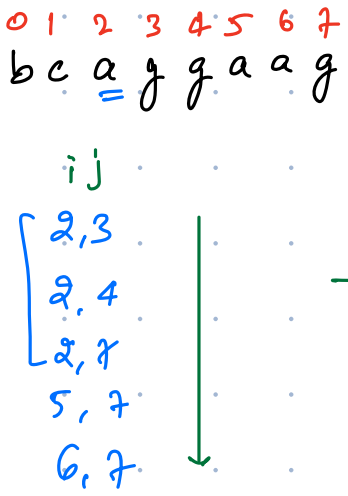


3 + 1 = 4 pairs.



5 pairs.

:



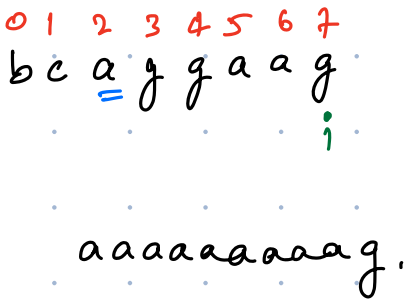
B.Fo

* Run a nested loop.
* For every 'a' encountered, count no. of 'g's to the right.

```
cnt=0                          You can also write  i<n-1.
for(int i=0; i<n;i++)
{
    if(s[i]!='a')
        continue;

    for( int j=i+1;j<n; j++)
    {
        if(s[j]=='g')
            cnt++

    }
}
return cnt
```



T.C.

↓

T.C. = $O(N^2)$

S.C. = $O(1)$

not good enough :(

Optimisation !

a d g a g a a f g a a g g

+5 +4 +3 +3 +2 +2

ans = 5 + 4 + 3 + 3 + 2 + 2

= 19

0 1 2 3 4 5 6 7

b c a g g a a g

cnt_g = ~~1~~ ~~2~~ 3

ans = ~~1~~ ~~2~~ 5

a a a a a a a a a a g

a a a a a a a a a a g

cnt = 1+1+1+1+...

Start from end.

If you see a g,

cnt_g++

If you see a a,

ans += cnt_g

If you see another character,

ignore.

Code.

```
ans=0, cnt_g=0
for (int i=n-1; i>=0; i--)
{
    if (s[i]=='g')
        cnt_g++
    else if (s[i]=='a')
        ans += cnt_g
    else
        // Ignore.
}
return ans
```

Not needed. Ignore.

T.C. = O(N)

S.C. = O(1)

"Carry forward"

H.W. :) Modify the approach so that you run loop in forward direction.

Subarrays.

↓

A contigous (continuous one) part of an array.

	0	1	2	3	4	5	6	7	8
Original arr:	4	1	2	3	-1	6	9	8	12

* 2 3 -1 6 ✓ 2 - 5
* 8 7 - 7 [A single element is also a subarray.]
* 4 1 2 3 -1 6 9 8 12 0 - 8 [Complete array is also a subarray.]
* 4 12 XX Not a subarray. (Not contigous)
* 2 3 1 XX Order is not same.
* { } XX Empty array is not a subarray.

0	1	2	3	4	5	6	7	8
4	1	2	3	-1	6	9	8	12

Representation of a subarray

Subarray can be represented by

(start point, end point)

(2,5)

Count no. of subarrays

0	1	2	3	4	5	6
4	2	10	3	12	-2	15
↑						

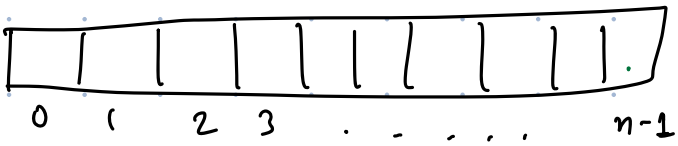
No. of subarrays starting from index 0 ?

0-0
0-1
0-2
0-3
0-4
0-5
0-6

7 subarrays.

0	1	2	3	4	5	6
4	2	10	3	12	-2	15
	↑					

No. of subarrays starting from index 1 = (1,1) (1,2) (1,3) (1,4) (1,5) (1,6) | 6 subarrays.



Consider all start points one by one.

No. of subarray starting from 0 = n

1 = n-1

2 = n-2

3 = n-3

⋮

n-1 = 1

$1+2+3 \ldots\ldots n-1+n = \frac{n(n+1)}{2}$ subarrays in array of size N.

Q. Print a subarray whose start point is i and end point is j.

```
for(int k=i ; k<=j ; k++)
{
    print(arr[k])
}
```

Q Print all subarrays of an array.

Generate all subarrays;

```
for(int i=0; i<n; i++)
{
    // i is fixed as a start point.
    for (int j=i ; j<n; j++)
    {
        // j is fixed as end point
        // i-j is the range of subarray.
        for (int k= i ; k<=j ; k++)
        {
            print (arr[k])
        }
    }
}
```

Printing 1 subarray.

T.C. = $O(N^3)$

Approach:

Break till 8:30 AM

1. Generate all possible subarrays
   - Outer loop iterates over all start points
   - Inner loop " " end points.

2. Once you get the range, iterate over the range to print it

Q. Given an array of N integers, return the length of the smallest subarray which contains both maximum and minimum of the array.

1 2 3 1 3 4 6 4 6 3

max=6
min=1

ans=4

2 2 6 4 5 1 5 2 6 4 1

min=1
max=6

contains both 1 and 6 and is smallest.

ans=3

1 3 0 4 2 7 0 8 0 0 9

ans=9

B. F.

* Find max of the array. → O(N)
* Find min of the array. → O(N)
* Generate every subarray, → O(N²) for each subarray:
  check if both max and min is present. → O(N)
* Store the minimum length of such subarrays in an ans variable

T.C → O(N³)

2 2 6 4 5 1 5 2 6 4 1

Observations

1. Min and max will always occur at the boundaries of your final ans.

   ↓

   If not at boundaries, you can always cut-short it.

   4 3 1 2 2 1 6 8 2 7 3

2. Min and max will always occur once in your final subarray.

2 2 6 4 6 5 1 5 2 6 4 1

min = 1
max = 6

ans = ~~5~~ 3

Approach:

1. Iterate over the array,
   If you encounter "min", iterate towards the right to find max.
   If you encounter "max", iterate " " to find min.
   Otherwise ignore.

T.C. $O(N^2)$

Not good enough.

0	1	2	3	4	5	6	7	8	9	10	11	12
1	6	4	6	5	1	5	2	6	4	4	2	1

Amax = 6 i

Amin = 1

max-i = -1 [Nearest Index of max that I have seen]

min-i = -1 (" " min .]

i=12

min-i = 12

max-i = -1

i=11,10,9 ignore

i=8

max-i = 8

min-i = 12

ans = 12-8+1 = 5

i=7,6 ignore

i=5

max-i = 8

min-i = 5

ans = 8-5+1 = 4

i=4 ignore

i=3

max-i = 3

min-i = 5

ans = 5-3+1 = 3

i=2 ignore

i=1

max-i = 1

min-i = 5

↓

ans won't be updated.

i=0

min-i = 0

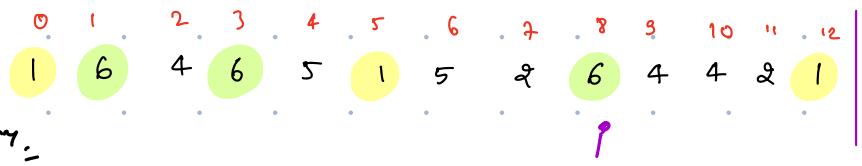
max-i = 1

ans = 1-0+1 = 2 ✓

Code.

```
// Find the min and max of the array.
// Amin and Amax.

ans = INT_MAX (+∞)        or length of the array.
min_i = -1
max_i = -1
```




```
if( Amin == Amax ) return 1.

for( int i=n-1 ; i>=0 ; i-- )
{
    if (arr[i] == Amin)
    {
        min_i = i
        if (max_i != -1)
        {
            ans = min( ans, max_i - min_i + 1 )
        }
    }
    else if (arr[i] == Amax)
    {
        max_i = i
        if (min_i != -1)
        {
```

min_i = 12
max_i = ~~-1~~ 8

Or take |max_i - min_i| absolute diff.

```
            ans = min(ans, min_i - max_i + 1)
        }
    }
}
return ans.
```

{ 8, 8, 8, 8, 8, 8, 8 }

T.C. → O(N)

S.C. → O(1)