



University of Central Punjab

Faculty of Information Technology

Compiler Construction

Assignment 02

Total Marks: 10

Due Date: 1 January 2026

Instructions:

1. **This is your individual handwritten assignment.**
2. **Your assignment should be plagiarism free.** In case of plagiarism, straight Zero would be rewarded.
3. **Your assignment is projection of your project phase 02.**
4. **No late submissions would be entertained in any case.**

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Assignment 2

Task 1 : Purpose:

Avian is a minimalist educational programming language designed to demonstrate compiler construction principle while maintaining C++ familiarity. The word 'avian' comes from 'birds' or 'flight' aiming it to be flight of compiler construction.

Style of Syntax:

Avian adopts a hybrid syntax style. Combines C++ structure mostly. It used underscores before and after of usual C++ style of syntax with avn- i.e., avn-if, avn-for, avn-int and so on. Operators were given underscores as well for simplicity and readability. i.e., -op- eg: -+-, --, -*---- and so on.

Reasons for Keywords:

The syntax and name conventions were chosen for following reasons:

→ Namespace Isolation, prevents collisions between keywords and identifiers.

- enhanced readability.
- Simplified Parsing
- Educational Demonstration.
- Avian self identity ("Avian" → avn_)

Language Elements from Phase 1:

Keyword	Purpose	Example
avn_int	integer variable	avn_int n = 10;
avn_if	conditional branch	avn_if (x > 0) { ... }
avn_main	program entry	avn_int avn_main()
avn_cin	input stream	avn_cin >> n;
avn_break	loop break	avn_break;
avn_for	for loop	avn_for (i=0; i<5; i++) { }
Operators	Description	Example
-+	addition	sum = a + b;
=	equal (assignment)	x = 10;
>>	exression output	avn_cont << x;

Punctuations	Purpose	Example
- , -	Comma	avr_int add(a,-,b)
;	Statement Terminator	x = 5 ;
()	Parathesces	avr_if(x=5){...}

Task 2 : CFG

// Start

P

P → GL FL

GL → D GL | ε

FL → F FL | ε

// function definition

F → avn_func T id (para) B,

F → avn_func T avn_main () B

// DTypr

T → avn_int | avn_float | avn_double | avn_char

T → avn_void

// Parameters

Para → PL | ε

PL → PR PL'

PL → , PR PL' | ε

PR → T id

// Bracket

B → { SL }

SL → S SL | ε

// Statement Expression

$S \rightarrow D \mid A \mid IF \mid WH \mid FOR \mid OUT \mid IN$

$S \rightarrow RET \mid ES \mid \text{avn_end};$

// Variable Declaration

$D \rightarrow T \ id \ DI;$

$DI \rightarrow = E \mid \epsilon$

// Assignment Operator

$A \rightarrow id = E;$

// Conditional Statement

$IF \rightarrow \text{avn_if}(C) \ B \ ELSE$

$ELSE \rightarrow \text{avn_else} \ B \mid \epsilon$

// Loops

$WH \rightarrow \text{avn_while}(C) \ B$

$FOR \rightarrow \text{avn_for}(A_{\text{for}}, C_{\text{opt}}, V_{\text{opt}}) \ B$

$A_{\text{for}} \rightarrow id = E \mid \epsilon$

$C_{\text{opt}} \rightarrow C \mid \epsilon$

$V_{\text{opt}} \rightarrow V \mid \epsilon$

// Operators (Unary)

$U \rightarrow id_++ \mid id__-- \mid --id \mid -++id$

// Input / output

$OUT \rightarrow \text{avn_cout} \ll OL;$

$OL \rightarrow OI \ OL'$

$OL' \rightarrow , \ OL \ OL' \mid \epsilon$

$OI \rightarrow \underline{\text{str}} \mid id \mid \underline{\text{num}}$

// Input

IN → avn-avn ⇒ IO:

IO → id IO'

IO' → id IO' | ε

// Return

RET → avn-return Eopt;

Eopt → E | ε

// Expression

ES → CALL | V ;

CALL → id (ALopt) .

ALopt → AL | ε

AL → E AL'

AL' → , E AL' | ε

// Conditions + Operators

C → C&R | R

R → E RO E | E

RO → < | > | = | !=

// Arithmetic Operators

E → E AO T | T

AO → - + | - -

T → T MO F | F

MO → * | / | -

F → (E) | id | num | ch | str | CALL

Task 3:

Production Rules

$P \rightarrow F$
 $F \rightarrow \text{avn-func} \text{ avn-int } id() B$
 $B \rightarrow \{ SL \}$
 $SL \rightarrow S \text{ } SL | \epsilon$
 $S \rightarrow D | A | IF | WH | OUT$
 $D \rightarrow \text{avn-int } id ;$
 $A \rightarrow id = num ;$
 $IF \rightarrow \text{avn-if (num) } B$
 $WH \rightarrow \text{avn-while (num) } B$
 $OUT \rightarrow \text{avn-cout} << id ;$

Task 4 : FIRST & FOLLOWS

Non Terminal	First	Follow
$P(\text{start})$	$\{ \text{avn-func} \}$	$\{ \$ \}$
SL	$\{ \text{avn-int}, id, \text{avn-if}, \text{avn-while}, \text{avn-cout}, \epsilon \}$	$\{ \text{avn-int}, id, \text{avn-if}, \text{avn-while}, \text{avn-cout} \}$
SL	$\{ \text{avn-int}, id, \text{avn-if}, \text{avn-while}, \text{avn-cout}, \epsilon \}$	$\{ \text{avn-int}, id, \text{avn-if}, \text{avn-while}, \text{avn-cout} \}$
IF	$\{ \text{avn-if} \}$	$\{ \text{avn-int}, \text{avn-if}, \text{avn-while}, \text{avn-cout} \}$

Task 5

Ambiguity Check

The ambiguity occurs in the if-else statement.

IF → $\text{avn_if}(\text{num}) \text{B} \mid \text{avn_if}(\text{num}) \text{B} \text{avn_else} \text{B}$

This creates a problem classically "dangling else problem" where it is unclear which `avn_if` and `avn_else` belongs to when nested if statements are used.

example :

```
avn_if (num)
    avn_if (num) B
    avn_else B
```

Here, 'avn_else' can be associated with either inner if and outer if.

Solution :

In implementation is resolved by attaching `avn_else` to nearest unmatched `avn_if`.

Parser generator like YACC/BISON handle this using procedural rules or by separating matched and unmatched if statements.

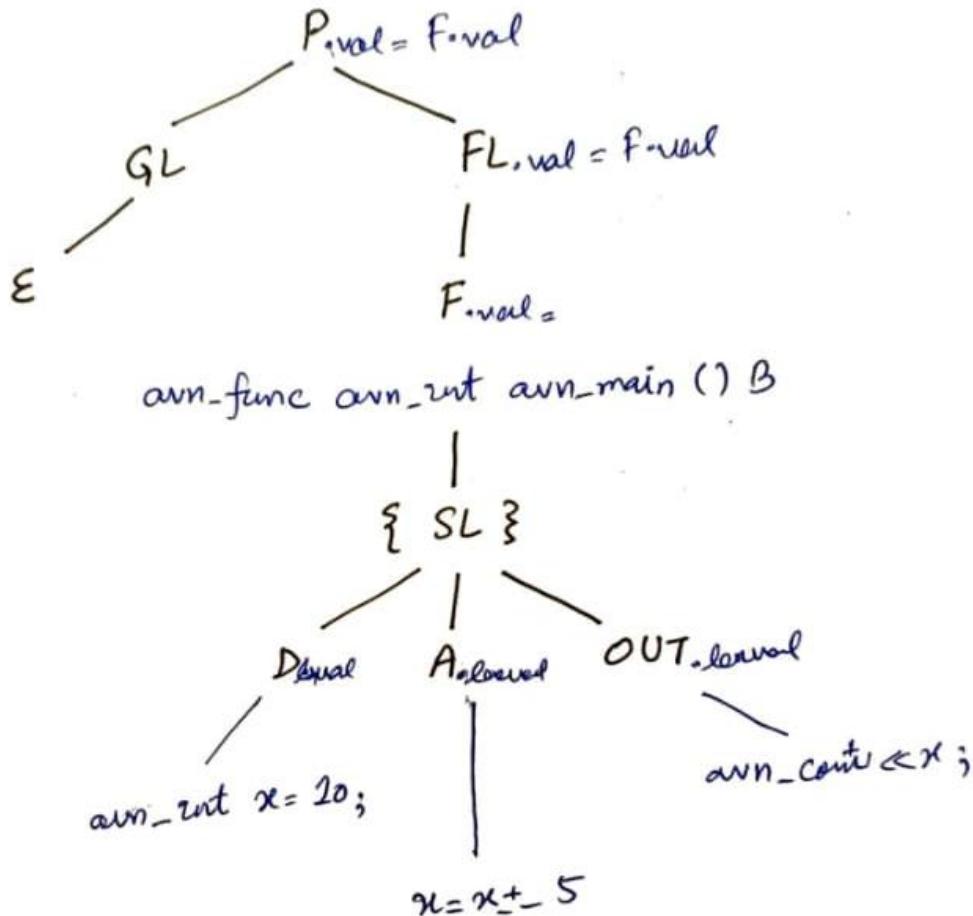
Task 6:

Parse Tree

avn-vid $x = 10;$

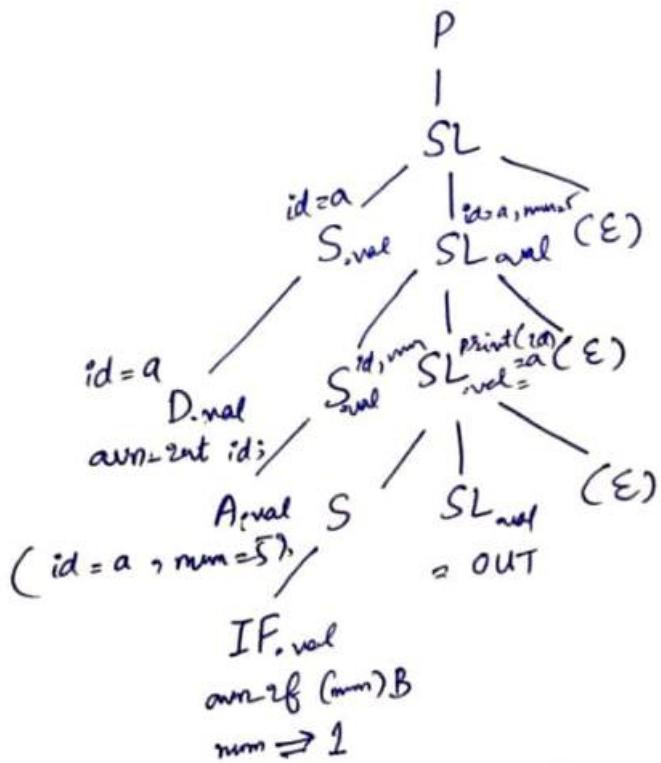
$x = x + 10;$

avn-cont $\ll x;$



Task 6: Parse Tree

```
avn-int a;  
a = 5;  
avn-if (1) { avn-print a; }
```



Task 7: Error Scenarios

E1: missing Semi-Colon

avn-int a

Rule violated: $D \rightarrow \text{avn_int} \ id;$

Expected token: ;

E2: missing closing Paranthesis

avn-if ({avn-cont < a; }

Rule Violated: avn-if (num) B

Expected Token:)