Outline

- 1 Recap
- 2 Dictionaries
- Wildcard queries
- 4 Edit distance
- Spelling correction
- 6 Soundex

Spelling correction

- Two principal uses
 - Correcting documents being indexed
 - Correcting user queries
- Two different methods for spelling correction
- Isolated word spelling correction
 - Check each word on its own for misspelling
 - Will not catch typos resulting in correctly spelled words, e.g., an asteroid that fell form the sky
- Context-sensitive spelling correction
 - Look at surrounding words
 - Can correct form/from error above

Correcting queries

- First: isolated word spelling correction
- Premise 1: There is a list of "correct words" from which the correct spellings come.
- Premise 2: We have a way of computing the distance between a misspelled word and a correct word.
- Simple spelling correction algorithm: return the "correct" word that has the smallest distance to the misspelled word.
- Example: information → information
- For the list of correct words, we can use the vocabulary of all words that occur in our collection.

Alternatives to using the term vocabulary

- A standard dictionary (Webster's, OED etc.)
- An industry-specific dictionary (for specialized IR systems)
- The term vocabulary of the collection, appropriately weighted

Distance between misspelled word and "correct" word

- We will study several alternatives.
- Edit distance and Levenshtein distance
- Weighted edit distance
- k-gram overlap

Edit distance

- The edit distance between string s_1 and string s_2 is the minimum number of basic operations that convert s_1 to s_2 .
- Levenshtein distance: The admissible basic operations are insert, delete, and replace
- Levenshtein distance dog-do: 1
- Levenshtein distance cat-cart: 1
- Levenshtein distance cat-cut: 1
- Levenshtein distance cat-act: 2
- Damerau-Levenshtein distance cat-act: 1
- Damerau-Levenshtein includes transposition as a fourth possible operation.

Levenshtein distance: Computation

		f	а	S	t
	0	1	2	3	4
С	1	1	2	3	4
a	2	2	1	2	3
t	3	3	2	2	2
S	4	4	3	2	3

```
Levenshtein Distance (s_1, s_2)
  1 for i \leftarrow 0 to |s_1|
  2 do m[i, 0] = i
  3 for j \leftarrow 0 to |s_2|
    do m[0,j] = j
  5 for i \leftarrow 1 to |s_1|
     do for j \leftarrow 1 to |s_2|
          do if s_1[i] = s_2[j]
                then m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]\}
  8
                else m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]+1\}
  9
      return m[|s_1|, |s_2|]
 10
Operations: insert (cost 1), delete (cost 1), replace (cost 1), copy
(cost 0)
```

```
LEVENSHTEINDISTANCE(s_1, s_2)
  1 for i \leftarrow 0 to |s_1|
  2 do m[i, 0] = i
    for j \leftarrow 0 to |s_2|
    do m[0,j] = j
  5 for i \leftarrow 1 to |s_1|
     do for j \leftarrow 1 to |s_2|
          do if s_1[i] = s_2[j]
                then m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]\}
  8
                else m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]+1\}
  9
      return m[|s_1|, |s_2|]
Operations: insert (cost 1), delete (cost 1), replace (cost 1), copy
(cost 0)
```

```
Levenshtein Distance (s_1, s_2)
    for i \leftarrow 0 to |s_1|
    do m[i, 0] = i
    for j \leftarrow 0 to |s_2|
    do m[0,j] = j
  5 for i \leftarrow 1 to |s_1|
     do for j \leftarrow 1 to |s_2|
          do if s_1[i] = s_2[j]
  8
                then m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]\}
                else m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]+1\}
  9
      return m[|s_1|, |s_2|]
 10
Operations: insert (cost 1), delete (cost 1), replace (cost 1), copy
(cost 0)
```

```
Levenshtein Distance (s_1, s_2)
  1 for i \leftarrow 0 to |s_1|
     do m[i, 0] = i
     for j \leftarrow 0 to |s_2|
     do m[0,j] = j
    for i \leftarrow 1 to |s_1|
     do for j \leftarrow 1 to |s_2|
          do if s_1[i] = s_2[j]
                then m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]\}
                else m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]+1\}
  9
      return m[|s_1|, |s_2|]
Operations: insert (cost 1), delete (cost 1), replace (cost 1), copy
(cost 0)
```

```
Levenshtein Distance (s_1, s_2)
  1 for i \leftarrow 0 to |s_1|
     do m[i, 0] = i
     for j \leftarrow 0 to |s_2|
     do m[0,j] = j
     for i \leftarrow 1 to |s_1|
     do for j \leftarrow 1 to |s_2|
          do if s_1[i] = s_2[j]
                then m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]\}
                else m[i,j] = \min\{m[i-1,j]+1, m[i,j-1]+1, m[i-1,j-1]+1\}
      return m[|s_1|, |s_2|]
Operations: insert (cost 1), delete (cost 1), replace (cost 1), copy
(cost 0)
```

Levenshtein distance: Example

		f	a	S	t
	0	1 1	2 2	3 3	4 4
	1	1 2	2 3	3 4	4 5
ر	1	2 1	2 2	3 3	4 4
	2	2 2	1 3	3 4	4 5
а	2	3 2	3 1	2 2	3 3
	3	3 3	3 2	2 3	2 4
L	3	4 3	4 2	3 2	3 2
	4	4 4	4 3	2 3	3 3
S	4	5 4	5 3	4 2	3 3

Each cell of Levenshtein matrix

cost of getting here from my upper left neighbor (copy or replace)	cost of getting here from my upper neighbor (delete)
cost of getting here from my left neighbor (insert)	the minimum of the three possible "movements"; the cheapest way of getting here

Levenshtein distance: Example

			f		â	à	9	5	t	
	0	- -	1 1	_		2	3	3	4	4
	1		1 2		2	3	3	4	4	5
C	1		2 1		2	2	3	3	4	4
	2	T :	2 2		1	3	3	4	4	5
а	2		3 2		3	1	2	2	3	3
	3		3 3		3	2	2	3	2	4
١	3		4 3		4	2	3	2	3	2
	4	4	4 4		4	3	2	3	3	3
S	4	_ _ !	5 4		5	3	4	2	3	3

Dynamic programming (Cormen et al.)

- Optimal substructure: The optimal solution to the problem contains within it subsolutions, i.e., optimal solutions to subproblems.
- Overlapping subsolutions: The subsolutions overlap. These subsolutions are computed over and over again when computing the global optimal solution in a brute-force algorithm.
- Subproblem in the case of edit distance: what is the edit distance of two prefixes
- Overlapping subsolutions: We need most distances of prefixes 3 times – this corresponds to moving right, diagonally, down.

Weighted edit distance

- As above, but weight of an operation depends on the characters involved.
- Meant to capture keyboard errors, e.g., m more likely to be mistyped as n than as q.
- Therefore, replacing m by n is a smaller edit distance than by q.
- We now require a weight matrix as input.
- Modify dynamic programming to handle weights

Using edit distance for spelling correction

- Given query, first enumerate all character sequences within a preset (possibly weighted) edit distance
- Intersect this set with our list of "correct" words
- Then suggest terms in the intersection to the user.
- \rightarrow exercise in a few slides

Exercise

- Compute Levenshtein distance matrix for OSLO SNOW
- 2 What are the Levenshtein editing operations that transform cat into catcat?

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	$\frac{1}{1}$				
s	2 2				
I	3 3				
o	4 4				

		S	n	0	w
	0	1 1	2 2	3 3	4 4
o	$\frac{1}{1}$	1 2 ?			
s	2 2				
I	3 3				
o	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	1 1	1 2 2 1			
s	2 2				
ı	3 3				
o	4 4				



		S	n	0	W
	0	1 1	2 2	3 3	4 4
0	$-\frac{1}{1}$	1 2 2 1	2 3 ?		
s	2 2				
I	3 3				
0	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
0	$\frac{1}{1}$	1 2 2 1	2 3 2 2		
s	2 2				
I	3 3				
o	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
0	$-\frac{1}{1}$	1 2 2 1	2 3 2 2	2 4 3 ?	
s	2 2				
I	3 3				
0	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	$\frac{1}{1}$	1 2 2 1	2 3 2	3 2	
S	2 2				
I	3 3				
o	4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
0	$-\frac{1}{1}$	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 ?
s	2 2				
I	3 3				
o	4 4				

			S		n		0		W	
	0	_	1	1	2	2	3	3	4	4
0	$-\frac{1}{1}$	_ I	1 2	2 1	2 2	3 2	3	2	⁴ / ₃	5 3
s	2 2	— I								
Ι	3	— I								
0	4	— I								

			S		n		0		W	
		0	1	1	2	2	3	3	4	4
0		1	1 2	2 1	2 2	3 2	2 3	2	⁴ / ₃	5 3
s		2 2	3	?						
Ι		3								
0	_	4								

		S	n	0	W
	0	1 1	2 2	3 3	4 4
0	$-\frac{1}{1}$	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1			
ı	3 3				
o	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	$-\frac{1}{1}$	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 ?		
I	3 3				
o	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
0	$\frac{1}{1}$	$\begin{array}{c c} 1 & 2 \\ \hline 2 & 1 \end{array}$	2 3 2 2	2 4 3 2	4 5 3 3
s	2 2	$\begin{array}{c c} 1 & 2 \\ \hline 3 & 1 \end{array}$	2 3 2 2		
Ι	3 3				
0	4 4				

			S		n		0		W	
	0	1	1		2	3	3	4	4	
o	1 1	1 2	2 1	2 2	3 2	3	2	⁴ / ₃	5 3	
s	2 2	3	2 1	2 2	3 2	3 3	3 ?			
I	3 3									
o	4 4									

		S	n	0	W	
	0	1 1	2 2	3 3	4 4	
0	$-\frac{1}{1}$	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 3	
s	2 2	1 2 3 1	2 3 2	3 3 3		
I	3 3					
o	4 4					

		S	n	0	W	
	0	1 1	2 2	3 3	4 4	
o	$-\frac{1}{1}$	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3	
s	2 2	1 2 3 1	2 3 2	3 3 3 3	3 4 4 ?	
I	3 3					
o	4 4					

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	1 1	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
S	2 2	1 2 3 1	2 3 2	3 3 3 3	3 4 4 3
I	3 3				
o	4 4				

		S	n	0	w
	0	1 1	2 2	3 3	4 4
o	$-\frac{1}{1}$	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2 2	3 3 3 3	3 4 4 3
I	3 3	3 2 4 ?			
0	4 4				

		S	n	0	w
	0	1 1	2 2	3 3	4 4
0	$-\frac{1}{1}$	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3 3	3 4 4 3
I	3 3	3 2 4 2			
0	4 4				

		S	n	0	w
	0	1 1	2 2	3 3	4 4
0	1 1	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3	3 4 4 3
ı	3 3	3 2 4 2	2 3 ?		
o	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	$-\frac{1}{1}$	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2 2	3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2		
o	4 4				

		S	n	0	w
	0	1 1	2 2	3 3	4 4
o	$-\frac{1}{1}$	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
S	2 2	1 2 3 1	2 3 2	3 3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2	3 4 3 ?	
0	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	$\frac{1}{1}$	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
S	2 2	1 2 3 1	2 3 2	3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2	3 4 3 3	
0	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	$-\frac{1}{1}$	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2	3 4 3	4 4 7
0	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	$-\frac{1}{1}$	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2 2	3 3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2	3 4 3	4 4 4 4
o	4 4				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
0	$-\frac{1}{1}$	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3 3	3 4 4 3
Ι	3 3	3 2 4 2	2 3 3 2	3 4 3 3	4 4 4 4
o	4 4	4 3 7 7			

		S	n	0	w
	0	1 1	2 2	3 3	4 4
0	$\frac{1}{1}$	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3	3 4 4 3
ı	3 3	3 2 4 2	2 3 3 2	3 4 3 3	4 4 4 4
0	4 4	4 3 5 3			

		S	n	0	W
	0	1 1	2 2	3 3	4 4
0	1 1	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 3
S	2 2	1 2 3 1	2 3 2 2	3 3 3 3	3 4 4 3
Ι	3 3	3 2 4 2	2 3 3 2	3 4 3	4 4 4 4
0	4 4	4 3 5 3	3 3 4 ?		

		9	5	r	ı	()	V	V
	0	1	1	2	2	3	3	4	4
0	1	_1	2	_2	3	_2	4	4	5
Ů	1	2	1	2	2	3	2	3	3
s	2	1	2	2	3	3	3	3	4
5	2	3	1	2	2	3	3	4	3
	3	3	2	2	3	3	4	4	4
'	3	4	2	3	2	3	3	4	4
	4	4	3	3	3				
0	4	5	3	4	3				

		S	n	0	W
	0	1 1	2 2	3 3	4 4
	1	1 2	2 3	2 4	4 5
0	1	2 1	2 2	3 2	3 3
	2	1 2	2 3	3 3	3 4
S	2	3 1	2 2	3 3	4 3
	3	3 2	2 3	3 4	4 4
ĽIJ	3	4 2	3 2	3 3	4 4
	4	4 3	3 3	2 4	
0	4	5 3	4 3	4 ?	

		S	n	0	W
	0	1 1	2 2	3 3	4 4
0	1 1	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2	3 4 3 3	4 4 4 4
О	4 4	4 3 5 3	3 3 4 3	2 4 4 2	

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	1	1 2	2 3	2 4	4 5
	1	2 1	2 2	3 2	3 3
s	2	1 2	2 3	3 3	3 4
٦	2	3 1	2 2	3 3	4 3
	3	3 2	2 3	3 4	4 4
_ '	3	4 2	3 2	3 3	4 4
	4	4 3	3 3	2 4	4 5
0	4	5 3	4 3	4 2	3 ?

		S	n	0	w
	0	1 1	2 2	3 3	4 4
o	1 1	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3 3	3 4 4 3
ı	3 3	3 2 4 2	2 3 3 2	3 4 3	4 4 4
o	4 4	4 3 5 3	3 3 4 3	2 4 4 2	4 5 3 3

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	$-\frac{1}{1}$	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2	3 4 3	4 4 4
o	4 4	4 3 5 3	3 3 4 3	2 4 4 2	⁴ ⁵ 3 3

			9	5	r	ı	()	V	V
	0	_	1	1		2	3	3	4	4
o	1		1	2	2	3	2	4	4	5
	1		2	1	2	2	3	2	3	3
s	2	_	1	2	2	3	3	3	3	4
Щ	2		3	1	2	2	3	3	4	3
	3	— 1	_3	2	2	3	3	4	4	4
	3		4	2	3	2	3	3	4	4
0	4		4	3	_3	3	_2	4	4	5
Ŭ	4		5	3	4	3	4	2	3	3

What are the editing operations that transform OSLO into SNOW?

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	1 1	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2	3 4 3 3	4 4 4
o	4 4	4 3 5 3	3 3 4 3	2 4 4 2	4 5 3 3

cost	operation	input	output
1	insert	*	W

		S	n	0	W
	0	1 1	2 2	3 3	4 4
	_ 1	1 2	2 3	2 4	4 5
0	1	2 1	2 2	3 2	3 3
	2	1 2	2 3	3 3	3 4
S	2	3 1	2 2	3 3	4 3
	3	3 2	2 3	3 4	4 4
'	3	4 2	3 2	3 3	4 4
	4	4 3	3 3	2 4	4 5
0	4	5 3	4 3	4 2	3 3

cost	operation	input	output
0	(copy)	0	0
1	insert	*	W

		S	n	0	w
	0	1 1	2 2	3 3	4 4
0	1 1	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
S	2 2	1 2 3 1	2 3 2	3 3 3	3 4 4 3
ı	3 3	3 2 4 2	2 3 3 2	3 4 3	4 4 4
0	4 4	4 3 5 3	3 3 4 3	2 4 4 2	4 5 3 3

cost	operation	input	output
1	replace		n
0	(copy)	0	0
1	insert	•	W

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	$-\frac{1}{1}$	1 2 2 1	2 3 2 2	2 4 3 2	4 5 3 3
S	2 2	1 2 3 1	2 3 2 2	3 3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2	3 4 3 3	4 4 4
0	4 4	4 3 5 3	3 3 4 3	2 4 4 2	4 5 3 3

cost	operation	input	output
0	(copy)	S	S
1	replace		n
0	(copy)	0	0
1	insert	*	W

		S	n	0	W
	0	1 1	2 2	3 3	4 4
o	1 1	1 2 2 1	2 3 2	2 4 3 2	4 5 3 3
s	2 2	1 2 3 1	2 3 2	3 3 3	3 4 4 3
I	3 3	3 2 4 2	2 3 3 2	3 4 3	4 4 4
0	4 4	4 3 5 3	3 3 4 3	2 4 4 2	4 5 3 3

СО	st operation	n input	output
1	delete	0	*
0	(copy)	S	S
1	replace	I	n
0	(copy)	0	0
1	insert	*	W

				С	á	a	1	t	(á	a	t	
	_	0	1	1		2	3	3	4	4	5	5	6	6
		1	0	2	2	3	3	4	3	5	5	6	6	7
		1	2	0	1	1	2	2	3	3	4	4	5	5
		2	2	1	0	2	2	3	3	4	3	5	5	6
а		2	3	1	2	0	1	1	2	2	3	3	4	4
+		3	3	2	2	1	0	2	2	3	3	4	3	5
١		3	4	2	3	1	2	0	1	1	2	2	3	3

		С	a	t	С	a	t
	0	1 1	2 2	3 3	4 4	5 5	6 6
С	1 1	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3	5 6 4 4	6 7 5 5
а	2 2	2 1 3 1	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3	5 6 4 4
t	3 3	3 2 4 2	2 1 3 1	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3

cost	operation	input	output
1	insert	*	С
1	insert	*	a
1	insert	*	t
0	(copy)	С	С
0	(copy)	a	a
0	(copy)	t	t

		С	а	t	С	a	t
	0	1 1	2 2	3 3	4 4	5 5	6 6
С	1 1	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3	5 6 4 4	6 7 5 5
а	2 2	2 1 3 1	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3	5 6 4 4
t	3 3	3 2 4 2	2 1 3 1	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3

cost	operation	input	output
0	(copy)	С	С
1	insert	*	a
1	insert	*	t
1	insert	*	С
0	(copy)	a	a
0	(copy)	t	t

		С	a	t	С	a	t
	0	1 1	2 2	3 3	4 4	5 5	6 6
С	1 1	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3	5 6 4 4	6 7 5 5
а	2 2	2 1 3 1	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3	5 6 4 4
t	3 3	3 2 4 2	2 1 3 1	$\begin{array}{c c} 0 & 2 \\ \hline 2 & 0 \end{array}$	2 3 1 1	3 4 2 2	3 5 3 3

cost	operation	input	output
0	(copy)	С	С
0	(copy)	a	a
1	insert	*	t
1	insert	*	С
1	insert	*	a
0	(copy)	t	t

			С	a	t	С	а	t
	0	-	1 1	2 2	3 3	4 4	5 5	6 6
С	$-\frac{1}{1}$	-	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3	5 6 4 4	6 7 5 5
а	2 2	-	2 <u>1</u> 3 <u>1</u>	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3	5 6 4 4
t	3 3	-	3 2 4 2	2 1 3 1	0 2 2 0	2 3 1 1	3 4 2 2	3 5 3 3

cost	operation	input	output
0	(copy)	С	С
0	(copy)	a	a
0	(copy)	t	t
1	insert	*	С
1	insert	*	a
1	insert	*	t

Outline

- 1 Recap
- 2 Dictionaries
- Wildcard queries
- 4 Edit distance
- Spelling correction
- 6 Soundex

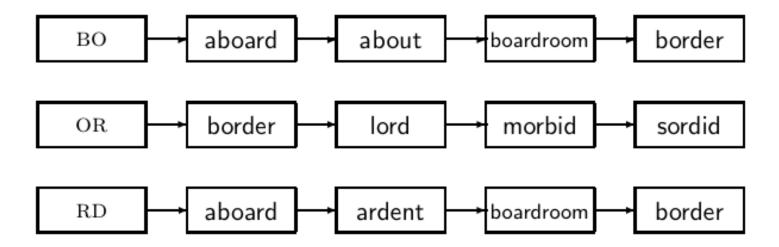
Spelling correction

- Now that we can compute edit distance: how to use it for isolated word spelling correction
- k-gram indexes for isolated word spelling correction.
- Context-sensitive spelling correction
- General issues

k-gram indexes for spelling correction

- Enumerate all k-grams in the query term
- Example: bigram index, misspelled word bordroom
- Bigrams: bo, or, rd, dr, ro, oo, om
- Use the k-gram index to retrieve "correct" words that match query term k-grams
- Threshold by number of matching k-grams
- E.g., only vocabulary terms that differ by at most 3 k-grams

k-gram indexes for spelling correction: bordroom



Context-sensitive spelling correction

- Our example was: an asteroid that fell form the sky
- How can we correct form here?
- One idea: hit-based spelling correction
 - Retrieve "correct" terms close to each query term
 - for flew form munich: flea for flew, from for form, munch for munich
 - Now try all possible resulting phrases as queries with one word "fixed" at a time
 - Try query "flea form munich"
 - Try query "flew from munich"
 - Try query "flew form munch"
 - The correct query "flew from munich" has the most hits.