

3) mathematically derive the average runtime complexity of the non-random pivot version of quicksort

A: Recurrence relation is

$$T(n) = T(k) + T(n-k-1) + O(n)$$

here

$T(n)$ = Time complexity of quicksort for an array of size n

k = number of element in the left subarray

$n-k-1$ = no of element in right array

$O(n)$ = represented time array

considering Average case

$$T(n) = 2T(n/2) + O(n)$$

here $2T(n/2)$ represented avg time for recursively

$$\text{So } T(n) = O(n) + 2 \cdot T(n/2)$$

$$T(n) = O(n) + 2 \left(O\left(\frac{n}{2}\right) + 2 \cdot T\left(\frac{n}{4}\right) \right)$$

$$T(n) = O(n) + 2 \left(O\left(\frac{n}{2}\right) + 4 \cdot T\left(\frac{n}{4}\right) \right)$$

$$\therefore T(n) = k \cdot O\left(\frac{n}{2^k}\right) + 2^k \cdot T\left(\frac{n}{2^k}\right)$$

Therefore, Runtime complexity of quicksort for Average case is $O(n \log n)$

$$T(n) = O(n \log n)$$