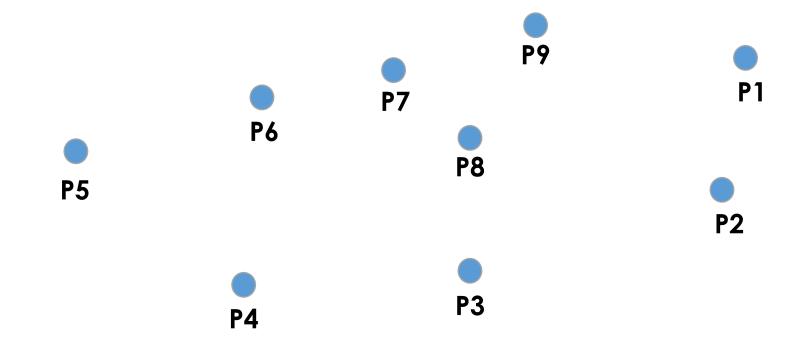
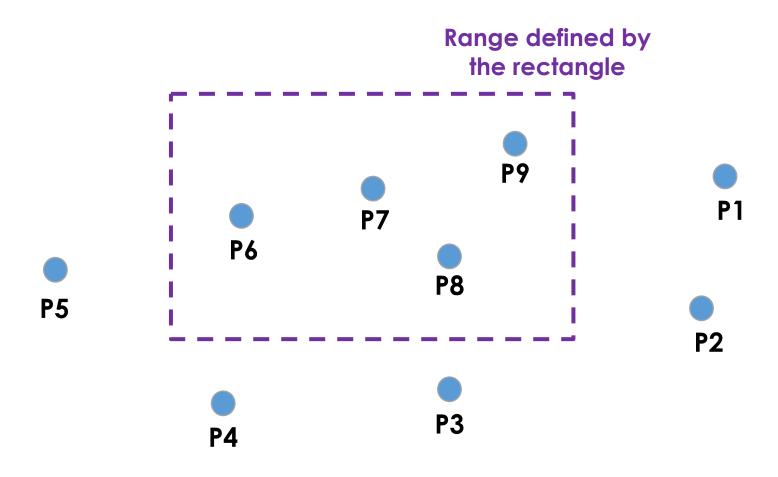
# Spatial Data structures For Point Data

# Points in 2D space



#### Range Queries

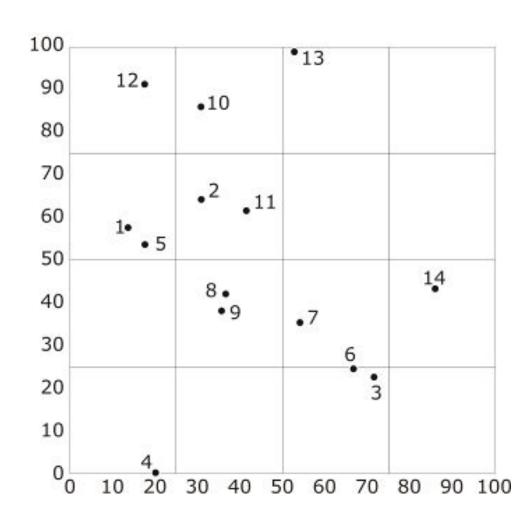


Retrieve all points inside the rectangle E.g., return all restaurants in GK-1, New Delhi Range can also be defined by a rectangle How to store 2D points in a secondary memory?

Data structure should support efficient query algorithms for range queries

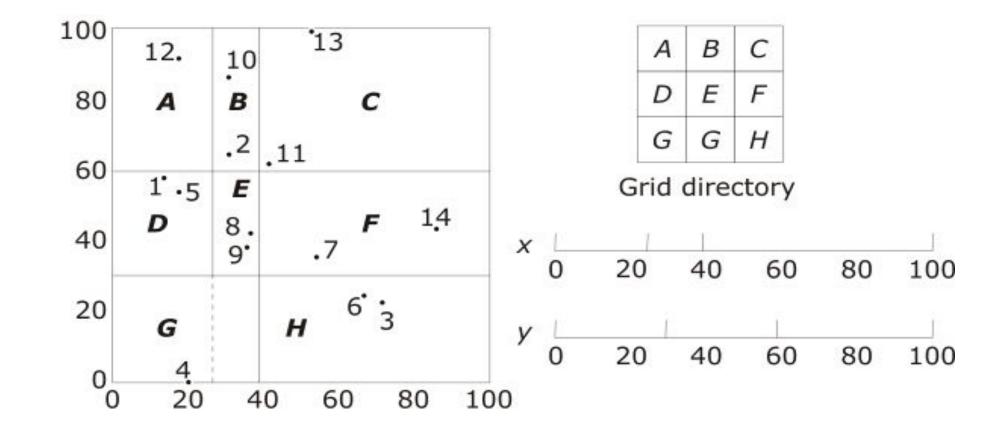
#### Uniform Grids over 2D space

- Divide space into cells by a grid
- Store data in each cell in distinct disk page
- A directory structure needed
- Efficient for find, insert, range and nearest neighbor
- But wastes a lot of disk storage space
- Non-uniform data distribution over space ??

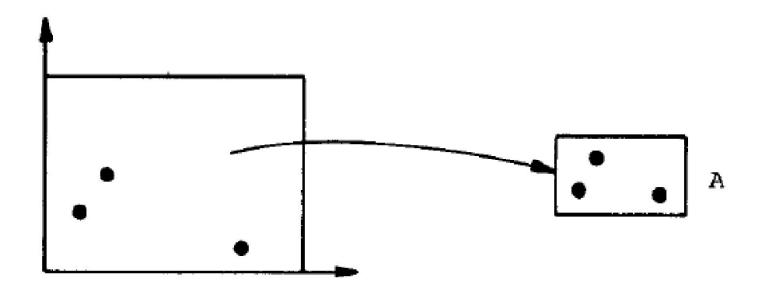


#### Refinement of basic idea into Grid Files

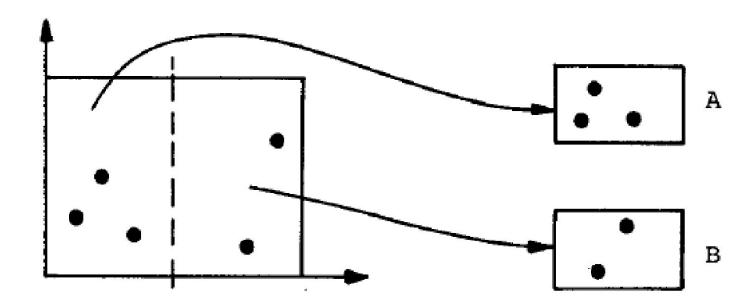
- Use non-uniform grids
- Linear scale store row and column boundaries
- Allow sharing of disk pages across grid cells



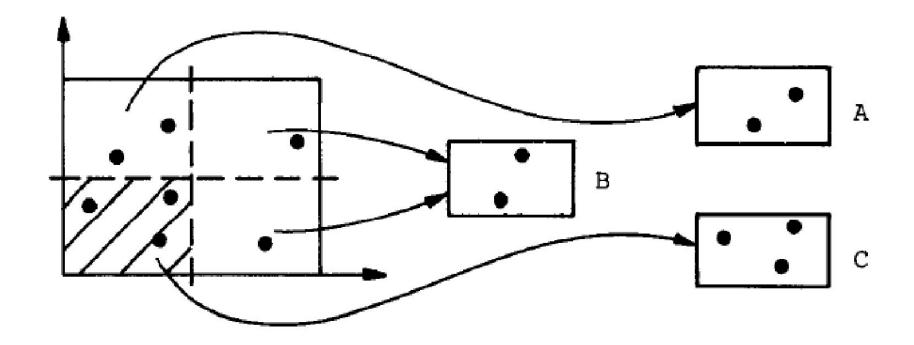
Capacity of bucket = 3

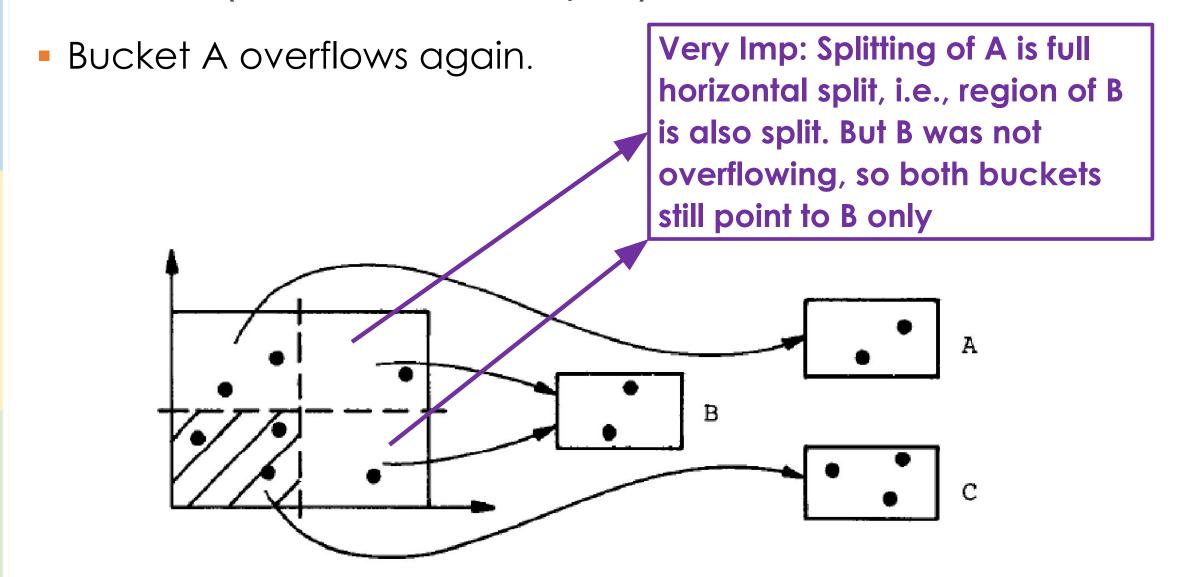


- When the bucket overflows we split it.
- A new bucket is made.
- Records that lie in one half of the space are moved to the new bucket.



Bucket A overflows again.

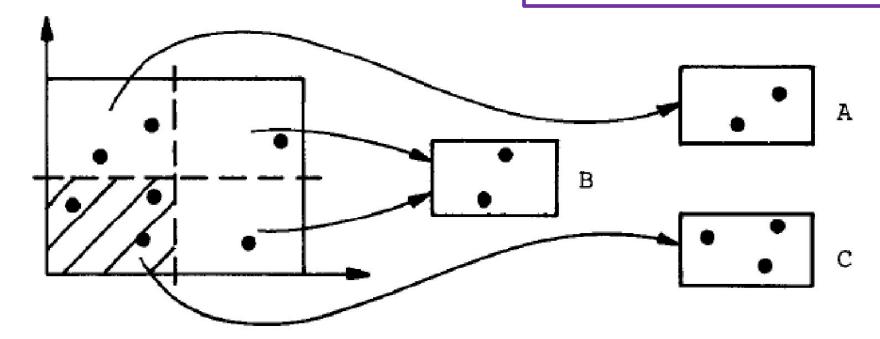




J. Nievergelt and H. Hinterberger. The Grid File: An Adaptable, Symmetric Multikey File Structure. ACM Transactions on Database Systems, Vol. 9, No. 1, March 1994

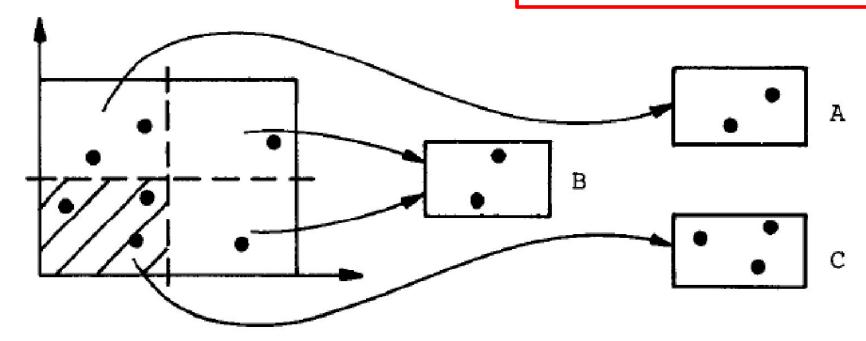
Bucket A overflows again.

In Grid files, data space which are the buckets is different from the geographic spread of the data.

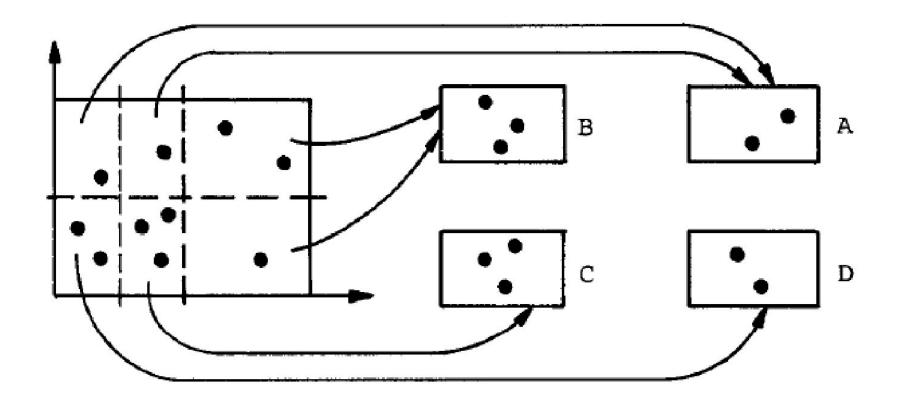


Bucket A overflows again.

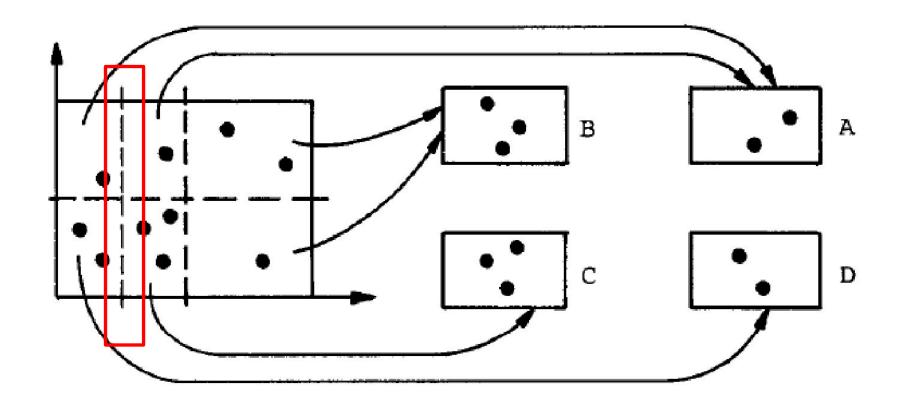
Splits in any dimension are made through and trough out. This makes the task of maintain linear scales easy

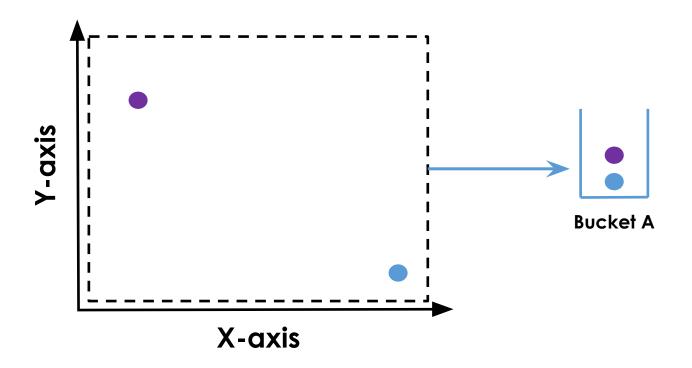


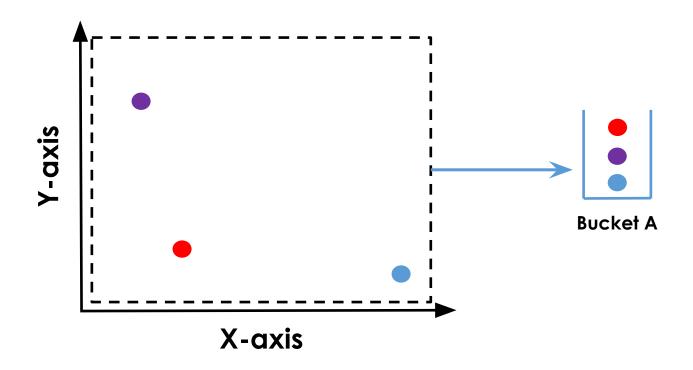
One more split.

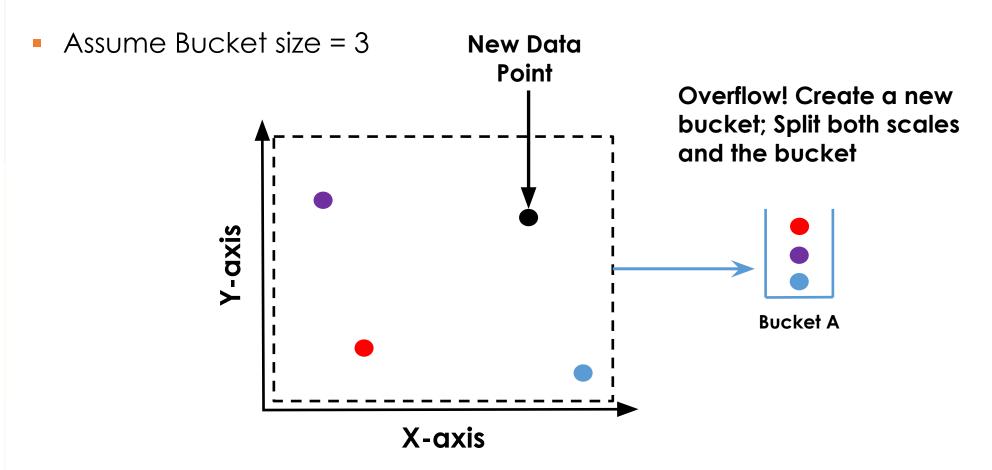


- One more split.
- Note that splits in any dimension are made through and trough.



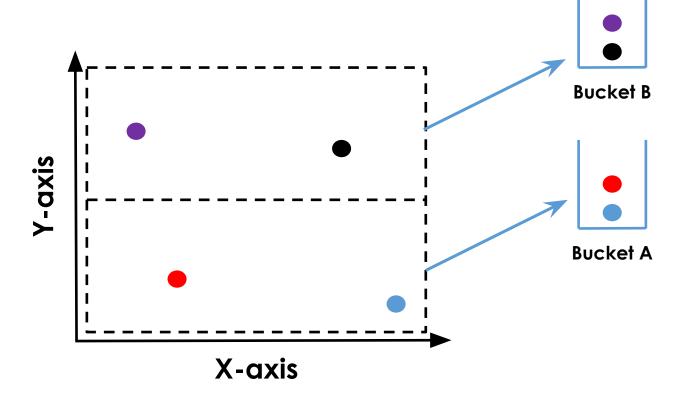




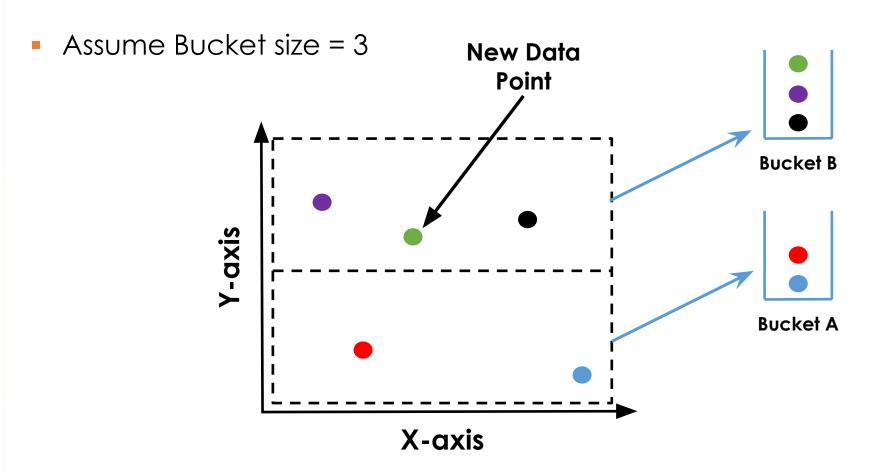


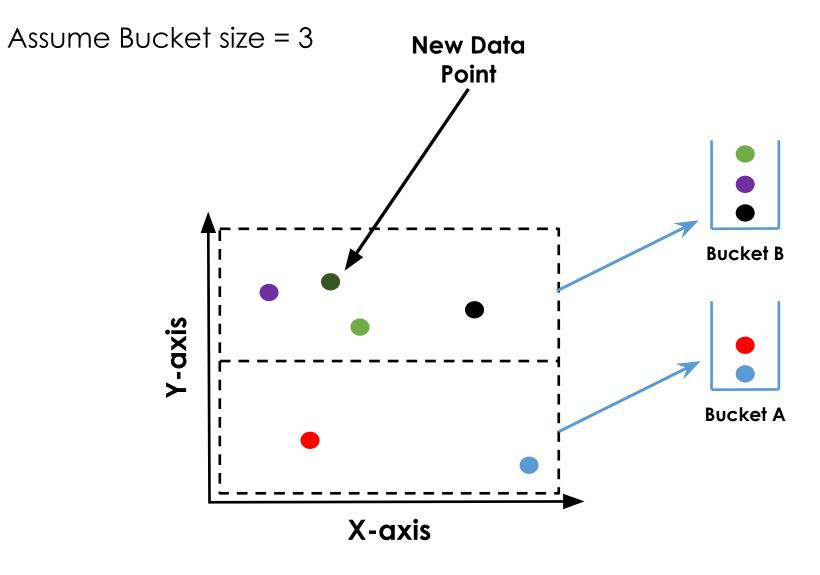
#### Resulting Grid File

Assume Bucket size = 3



Both Directory (x & y scales) and the Bucket are split



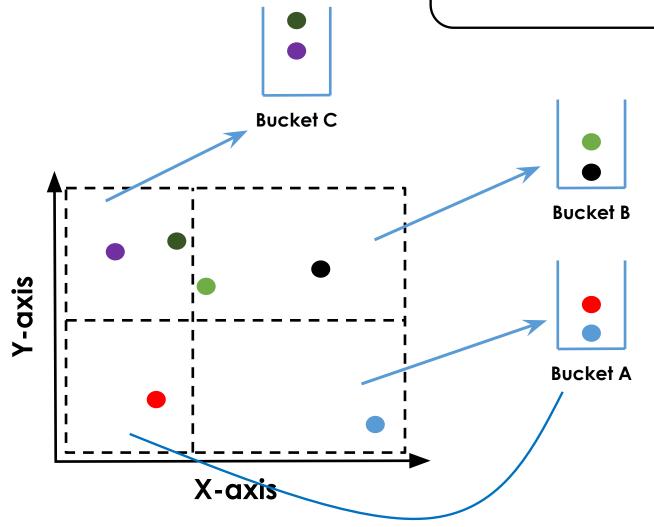


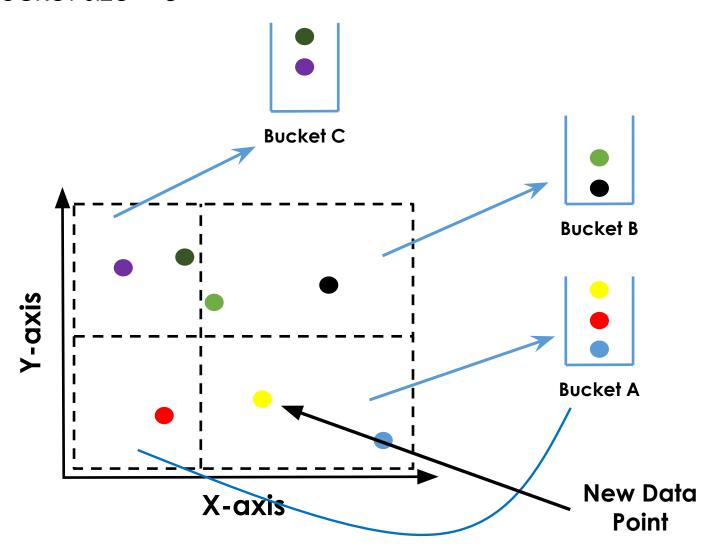
Overflow! Create a new bucket; Split both scales and the bucket.

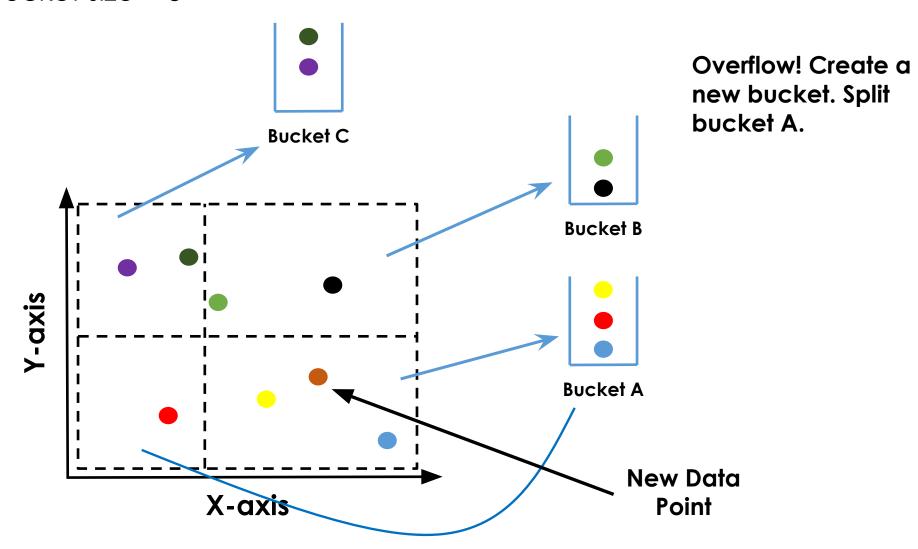
# Resulting Grid File

Assume Bucket size = 3

Both Directory (x & y scales) and the Bucket are split



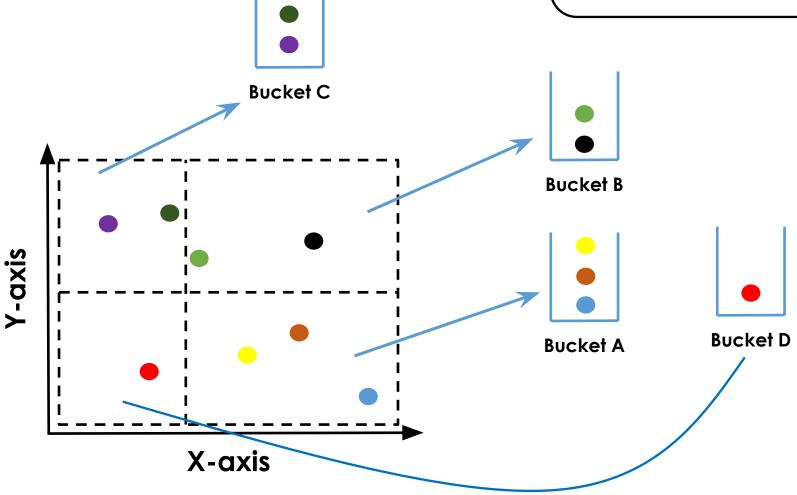


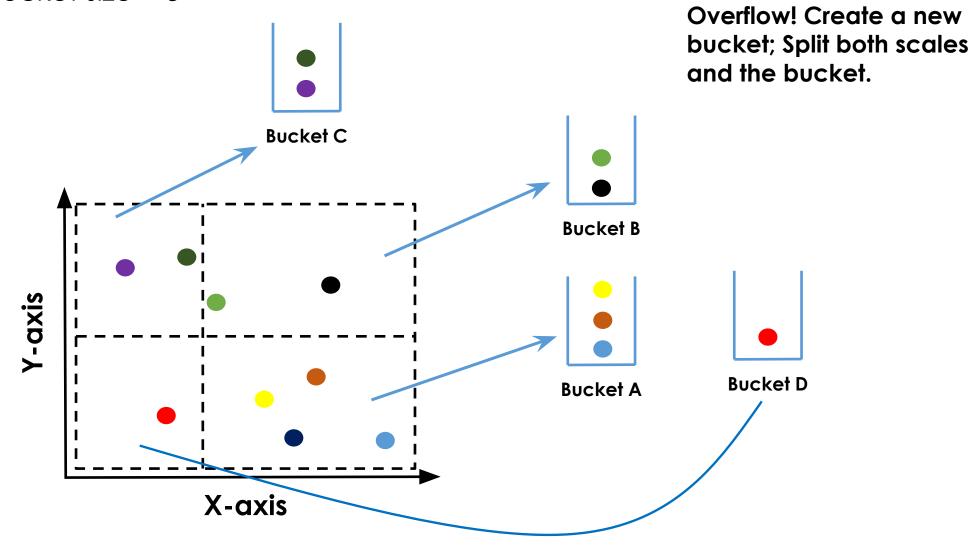


## Resulting Grid File

Assume Bucket size = 3

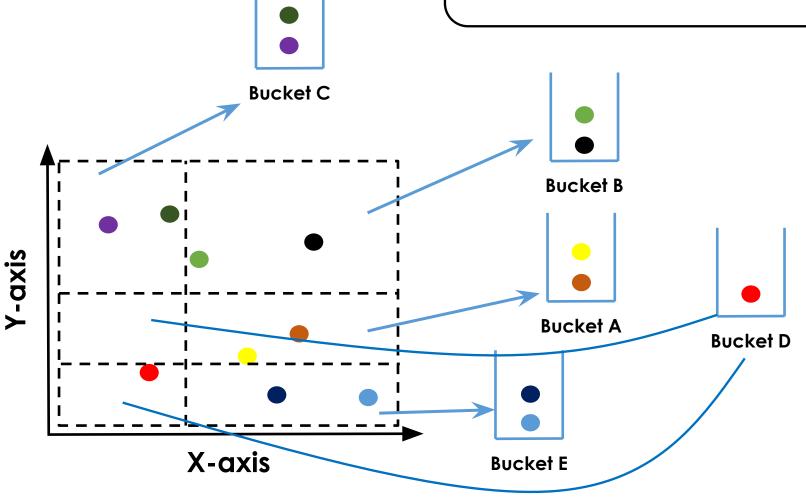
Only Bucket was split.
Directory was not split in this case.





Assume Bucket size = 3

Both Directory (x & y scales) and the Bucket are split



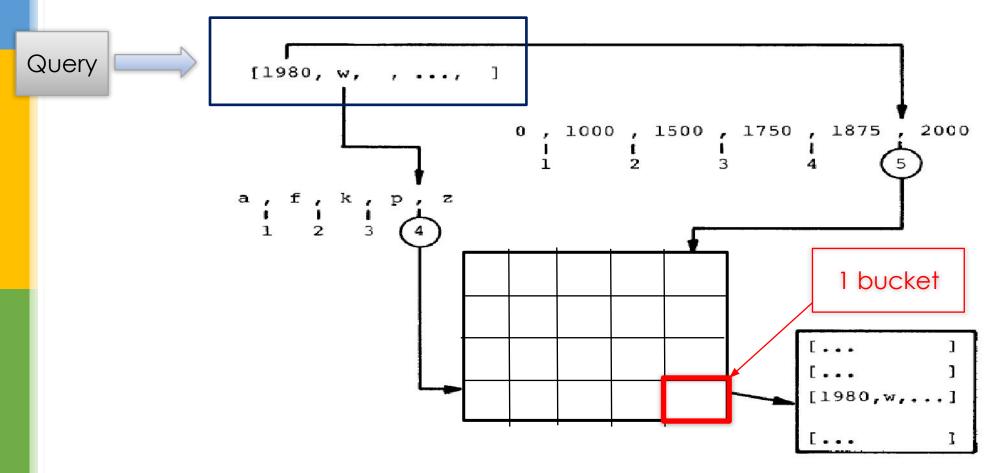
#### Grid Files (Splitting Policies)

#### Splits:

- Can happen during insertion.
- Overflow of a bucket corresponding to a grid partition leads to a split.
- Can also happen if bucket containing records from several grid partition fills up.
- Splitting dimension can be changed alternatively.
- Splitting point may not always be the middle point, other algorithms are also possible.
   J. Nievergelt and H. Hinterberger. The Grid File: An Adaptable, Symmetric Multikey File Structure. ACM Transactions on Database

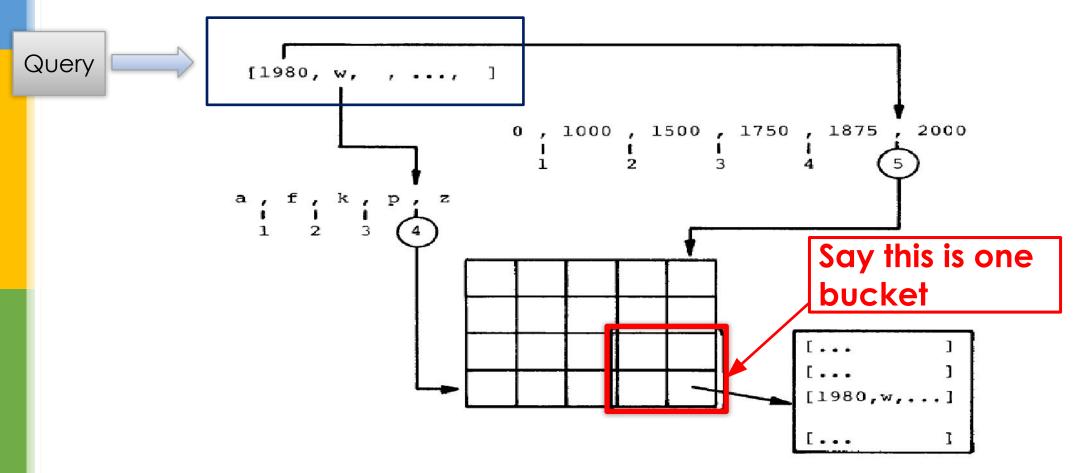
## Grid Files (Querying example)

- X-partitions (0,1000,1500,1750,1875,2000)
- Y-partitions (a, f, k, p, z).

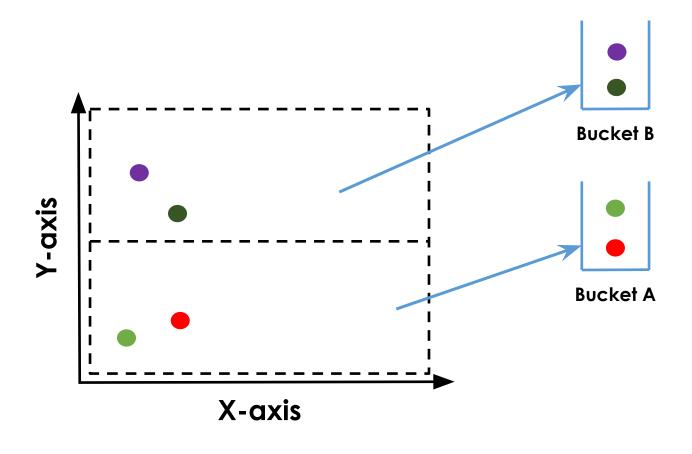


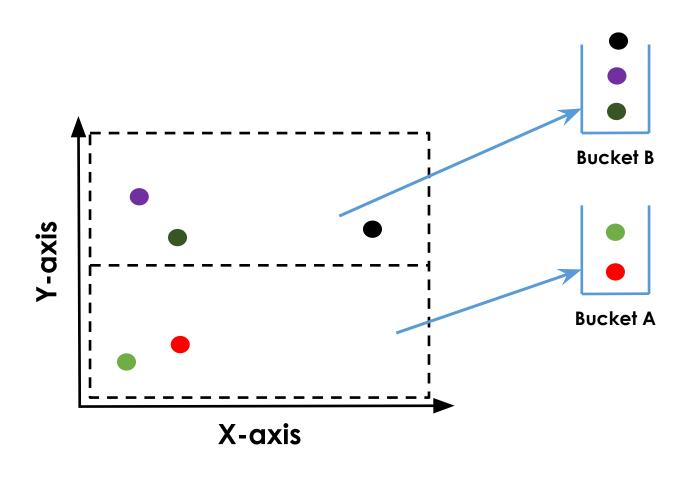
## Grid Files (Querying example)

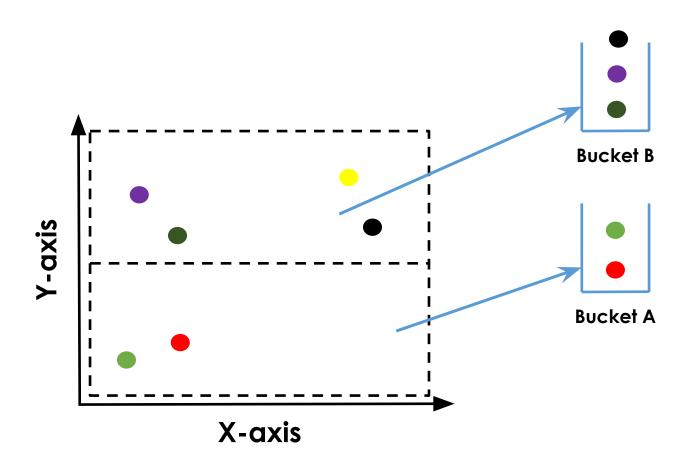
- X-partitions (0,1000,1500,1750,1875,2000)
- Y-partitions (a, f, k, p, z).

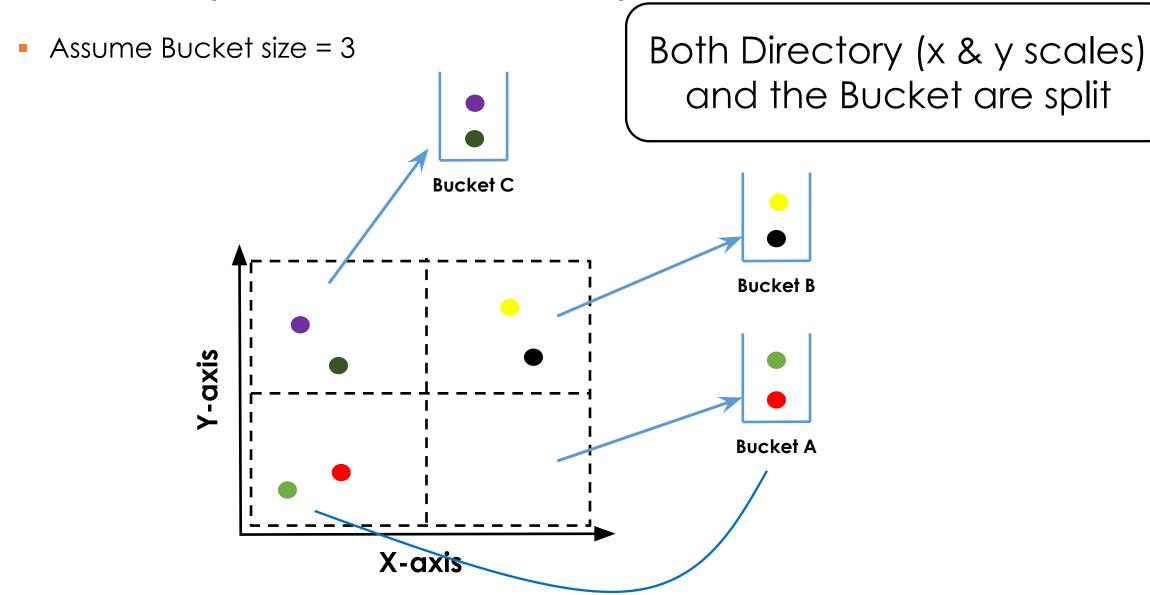


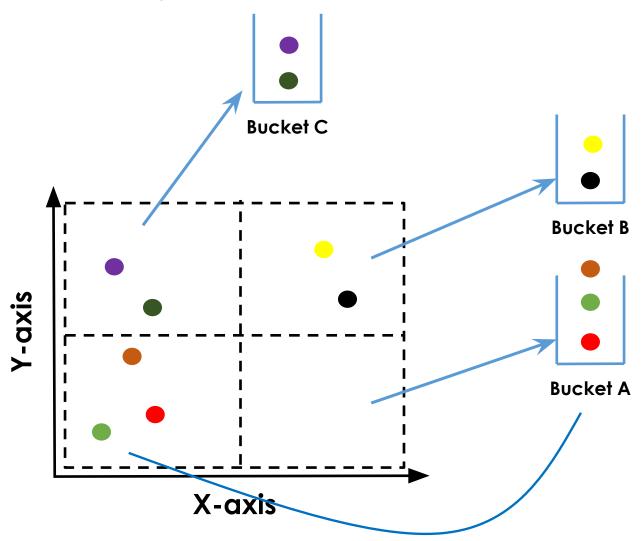


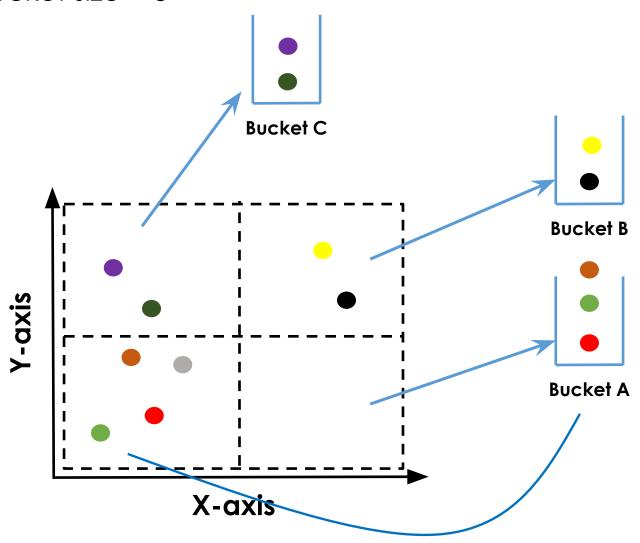


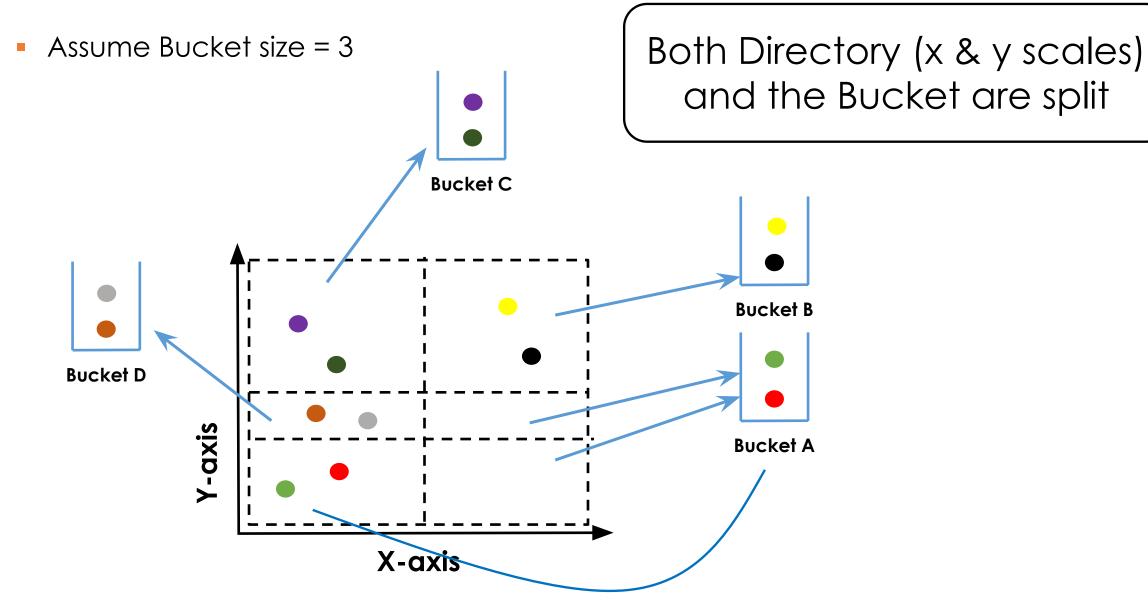












Assume Bucket size = 3

In general, there can be multiple options while splitting. For e.g., cell C1 could be pointing to bucket D as well. Also it is ok for multiple cells to point to the same bucket.

