

→ Algorithms → sequence of steps on data using efficient data structures to solve a given problem.

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→ Algorithm

Program

• Design

• Implementation

• Person with Domain knowledge writes Algorithm.

• Programmer writes the program.

• Any language

• Programming language

• Hardware and Operating System Independent.

• Hardware and Operating System Dependent.

• Analysis

• Testing

→ Priori Analysis

Posteriori Testing

• Algorithm

• Program

• Independent of language

• language Dependent

• Hardware Independent

• Hardware Dependent

• Time and Space Function

• Watch Time and Bytes

## → Characteristics of algorithm

1. Input - 0 or more
  2. Output - atleast one output
  3. Definiteness
  4. Finiteness
  5. Effectiveness
- \* ----- \*

## → How to write an algorithm

- algorithm swap (a, b)

start

```
temp = a;  
a = b;  
b = temp;
```

end

----- \* ----- \*

## → How to analyze an algorithm

1. Time

2. Space

3. Network

3. Power

6. CPU Registers

----- \* ----- \*

start

Time

temp = a; - 1

a = b; - 1

b = temp; - 1

end

 $f(n) = 3$ 

Space

a - 1

b - 1

temp - 1

 $S(n) = 3 \text{ word}$ 

## → Frequency Count Method

- sum of array

A is array, where n is number

Algorithm sum(A, n)

start

S = 0; — 1  $n = 5$ 

for (i=0; i &lt; n; i++) { - n+1 space

S = S + A[i]; - n  $A = n$ }  $S = 1$ return S; - 1  $n = 1$ 

end

 $f(n) = 2n+3$   $i = 1$  $O(n)$   $S(n) = \frac{n}{1} = n+3$ 

- sum of two square matrices

 $n \times n \Rightarrow \text{Matrix}$ 

Algorithm Add(A, B, n)

start

for (i=0, i &lt; n, i++) - n+1

{

for (j=0; j < n; j++) - n(n+1)  
 $c[i, j] = A[i, j] + B[i, j];$  - n(n)

{

y

 $f(n) = 2n^2 + 2n + 1$   
~~or~~  $O(n^2)$ Space -  $A = n^2$ ,  $B = n^2$ ,  $C = n^2$ ,  $i = 1, n = 1, j = 1 = S(n) = 3(n^2+3)$   
 $O(n^2)$

• Multiplication of two square matrices

Algorithm multiply (A, B, n)

start

for ( $i=0$ ;  $i < n$ ;  $i++$ )       $-n+1$

    for ( $j=0$ ;  $j < n$ ;  $j++$ )       $-n(n+1)$

$c[i, j] = 0$ ;       $-n \times n$

        for ( $k=0$ ;  $k < n$ ;  $k++$ )       $(n+1) \times n \times n$

$c[i, j] = c[i, j] + A[i, k] * B[k, j];$

end

$$f(n) = \underline{8n^3 + 8n^2 + 2n + 1}$$

$\underline{O(n^3)}$

Space :  $A - n^2$

$B - n^2$

$C - n^2$

$n - 1$

$i - 1$

$j - 1$

$k - 1$

$$S(n) = 3n^2 + 4$$

$$O(n^2)$$

→ Time Complexity :

1. start

for ( i=0; i<n; i++ ) →  $O(n)$

start;

y

end

2.

start

for ( i=n; i>0; i-- ) →  $O(n)$

start;

y

end

3. start

for ( i=1; i<n; i=i+2 )

↳ if this is  $i = i + 20$

start;  $y = \frac{n}{2}$

y

end

$t(n) = n$   
 $\frac{n}{2}$   
still order will be same

$O(n)$