

→ Recurrence Relation $T(n) = T(n-1) + 1$

void Test(int n)

{

if ($n > 0$)

{

printf ("%d", n); -1
Test (n-1); - $T(n-1)$

}

}

$f(n) = n + 1$

$O(n)$

$T(n)$ Tracing tree

Test (3)

 \

 3 Test (2)

 \

 2 Test (1)

 \

 1 Test (0)

 \

$\overline{T(n)} = T(n-1) + 1$

X

Recurrence Table

$$T(n) = \begin{cases} 1 & n=0 \\ T(n-1) + 1 & n>0 \end{cases}$$

$$T(n) = T(n-1) + 1$$

$$\therefore T(n) = T(n-1) + 1$$

$$T(n-1) = T(n-2) + 1$$

Substitute $T(n-1)$

$$T(n) = [T(n-2) + 1] + 1$$

$$T(n) = T(n-2) + 2$$

$$T(n) = [T(n-3) + 1] + 2$$

$$T(n) = T(n-3) + 3$$

continue k times

$$T(n) = T(n-k) + k$$

assume $n - k = 0$

$$\therefore n = k$$

$$T(n) = T(n-n) + n$$

$$T(n) = T(0) + n$$

$$T(n) = 1 + n$$

$\Theta(n)$

→ Recurrence Relation ($T(n) = T(n-1) + n$)

void Test(int n) $-T(n)$

{ if ($n > 0$) -1

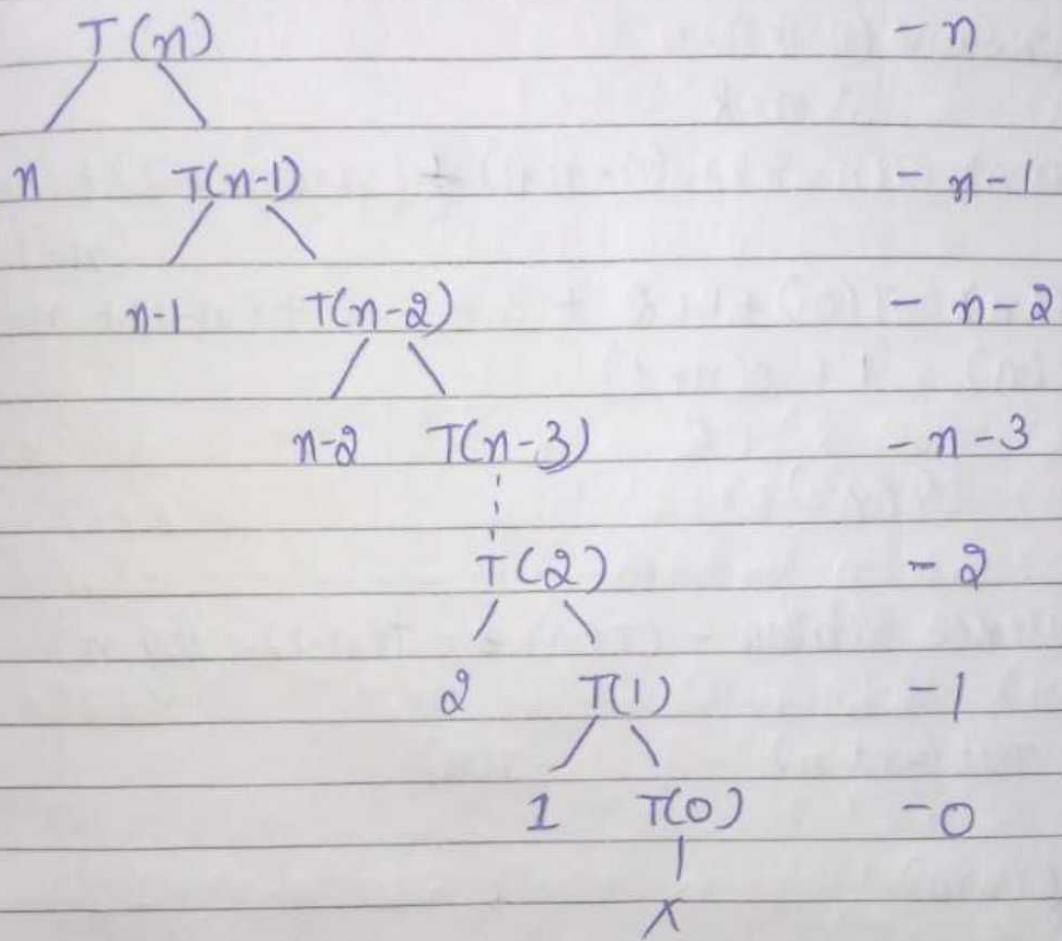
{ for ($i=0$; $i < n$; $i++$) { -n+1

printf("%d", n); } -n

Test($n-1$); $-T(n-1)$

} $\overline{T(n) = T(n-1) + 2n + 2}$

$$T(n) = \begin{cases} 1 & n=0 \\ T(n-1) + n & n>0 \end{cases}$$



$$0 + 1 + 2 + \dots + (n-1) + n = n(n+1)$$

$$T(n) = \frac{n(n+1)}{2}$$

$$\Theta = (n^2)$$

Another method

$$T(n) = T(n-1) + n$$

$$-(1) \because T(n) = T(n-1) + n$$

$$T(n) = [T(n-2) + n-1] + n \quad \therefore T(n-1) = T(n-2) + n-1$$

$$T(n) = T(n-2) + (n-1) + n \quad -(2) \quad \therefore T(n-2) = T(n-3) + n-2$$

$$T(n) = [T(n-3) + n-2] + (n-1) + n$$

$$T(n) = [T(n-3) + (n-2) + (n-1) + n]$$

$$T(n) = T(n-k) + (n-(k-1)) + (n-(k-2)) + \dots + (n-1) + n$$

Assume $(n-k) = 0$
 $\therefore n = k$

$$T(n) = T(n-n) + (n-n+1) + (n-n+2) + \dots + (n-1) + n$$

$$T(n) = T(0) + 1 + 2 + 3 + \dots + (n-1) + n$$

$$T(n) = 1 + \frac{n(n+1)}{2}$$

$$\Theta(n^2)$$

→ Recurrence Relation - $T(n) = T(n-1) + \log n$

void Test(int n) $-T(n)$

{

 if ($n > 0$)

 for ($i=1; i < n; i=i+2$)

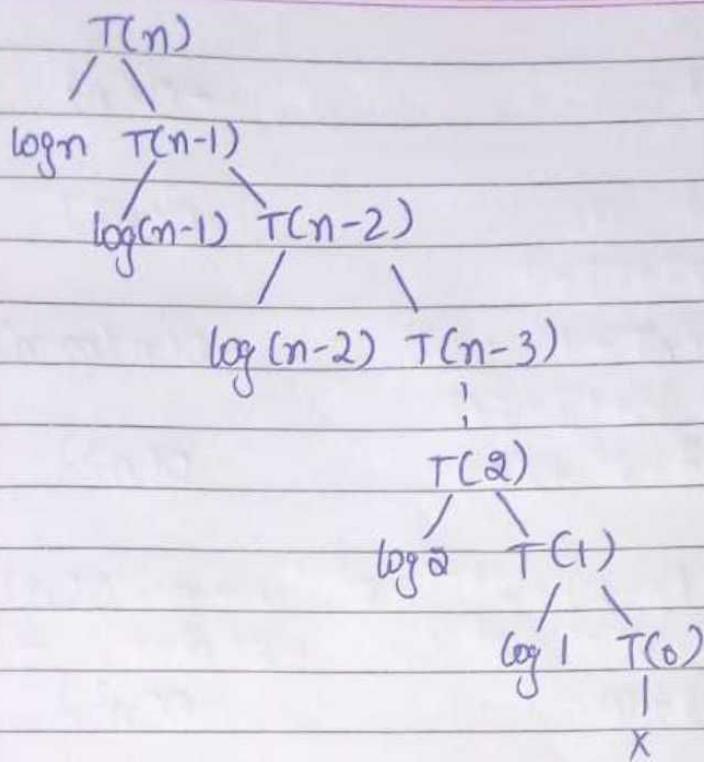
 printf("%d", i); $-\log(n)$

 Test(n-1); $-T(n-1)$

}

$$\underline{T(n) = T(n-1) + \log(n)}$$

$$T(n) = \begin{cases} 1 & n=0 \\ T(n-1) + \log(n) & n>0 \end{cases}$$



$$\log n + \log(n-1) + \log(n-2) + \dots + \log 2 + \log 1$$

$$= \log [n \times (n-1) \times \dots \times 2 \times 1]$$

$$= \log n!$$

$O(n \log n)$

Substitution Method

$$T(n) = T(n-1) + \log n \quad \text{--- (1)}$$

$$T(n) = [T(n-2) + \log(n-1)] + \log n$$

$$T(n) = T(n-2) + \log(n-1) + \log n \quad \text{--- (2)}$$

$$T(n) = [T(n-3) + \log(n-2)] + \log(n-1) + \log n$$

$$T(n) = T(n-3) + \log(n-2) + \log(n-1) + \log n \quad \text{--- (3)}$$

$$T(n) = T(n-k) + \log 1 + \log 2 + \log 3 + \dots + \log(n-1) + \log n$$

$$\therefore n-k=0$$

$$n=k$$

$$T(n) = T(0) + \log n!$$

$$T(n) = 1 + \log n!$$

~~$O(n \log n)$~~ $O(n \log n)$

$$\rightarrow T(n) = T(n-1) + 1 \quad - O(n)$$

$$T(n) = T(n-1) + n \quad - O(n^2)$$

$$T(n) = T(n-1) + \log(n) \quad - O(n \log n)$$

$$T(n) = T(n-1) + n^2 \quad - O(n^3)$$

$$T(n) = T(n-2) + 1 \quad - \frac{n}{2} O(n)$$

$$T(n) = T(n-100) + n \quad - O(n^2)$$

$$T(n) = 2 [T(n-1) + 1] \quad - O(2^n)$$

 \rightarrow Recurrence Relation $T(n) = 2 [T(n-1) + 1]$

Algorithm Test (int n) $- T(n)$

{

if ($n > 0$)
{

printf("y.o.d", n); $- 1$

Test (n-1); $- T(n-1)$

Test (n-1); $- T(n-1)$

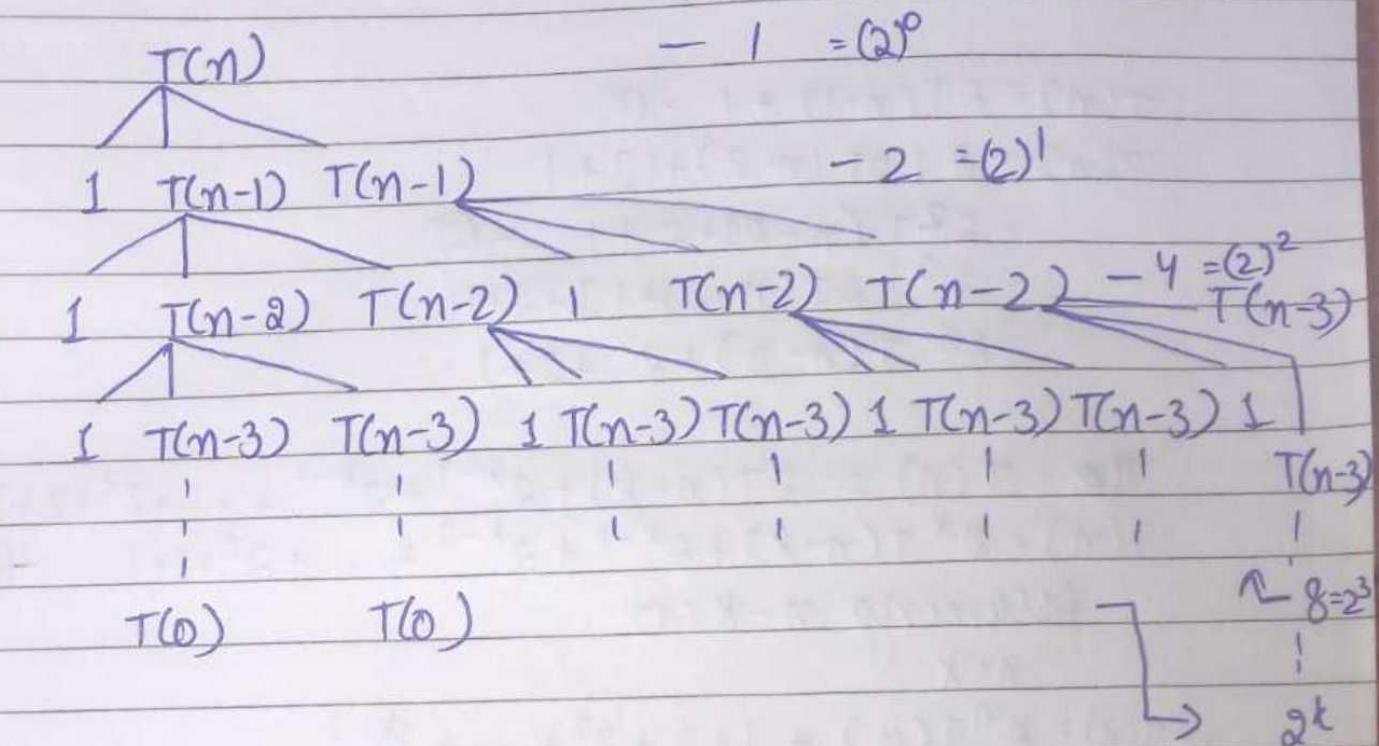
}

}

$$T(n) = 2 [T(n-1) + 1]$$

$$T(n) = \begin{cases} 1 & n=0 \\ 2 T(n-1) + 1 & n>0 \end{cases}$$

Recursion Tree Method



$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^k = 2^{k+1} - 1$$

$$a=1 \quad x=2 \quad \left[\text{using formula: } a + ax + ax^2 + ax^3 + \dots + ax^k = \frac{a(x^{k+1} - 1)}{x - 1} \right]$$

$$\therefore = \frac{1(2^{k+1} - 1)}{2 - 1}$$

$$= 2^{k+1} - 1$$

Assume $n = k = 0$

$$\begin{aligned} n &= k \\ \therefore 2^{n+1} - 1 &\Rightarrow O(2^n) \end{aligned}$$

Substitution Method :

$$T(n) = 2T(n-1) + 1 \quad -\textcircled{1}$$

$$T(n) = 2[2T(n-2) + 1] + 1$$

$$= 2^2 + (n-2) + 2 + 1 \quad -\textcircled{2}$$

$$= 2^2 [2T(n-3) + 1] + 2 + 1$$

$$= 2^3 T(n-3) + 2^2 + 2 + 1$$

!

$$\textcircled{3} \quad T(n) = 2T(n-k) + 2^{k-1} + 2^{k-2} + \dots + 2^2 + 2 + 1$$

$$T(n) = 2^k T(n-k) + 2^{k-1} + 2^{k-2} + \dots + 2^2 + 2 + 1 \quad -\textcircled{4}$$

Assuming $n-k=0$

$$n=k$$

$$\therefore T(n) = 2^n T(0) + 1 + 2 + 2^2 + \dots + 2^{k-1}$$

$$= 2^n \times 1 + 2^{k-1}$$

$$= 2^n + 2^n - 1$$

$$= 2^{n+1} - 1$$

$$\Theta(2^n)$$

