

# **Authentication versus Encryption**

<https://meetup.com/Internet-Security-and-Privacy>  
Vancouver

by Boris Reitman

# Who am I?

- Internet Security & Privacy Meetup  
(Vancouver)
- Secure webmail solution
- Advocate of Capitalism
- Speaker at B|Sides



# Agenda

- WhatsApp
- Byzantine General's Problem
- No auth – Shortened URLs
- OTP
- Biometric Auth + KDF
- Digital Signatures: RSA and ECC
- PGP
- Email phishing, DomainKeys
- DNSSEC, Tack
- OPAQUE

# Encryption vs. Authentication

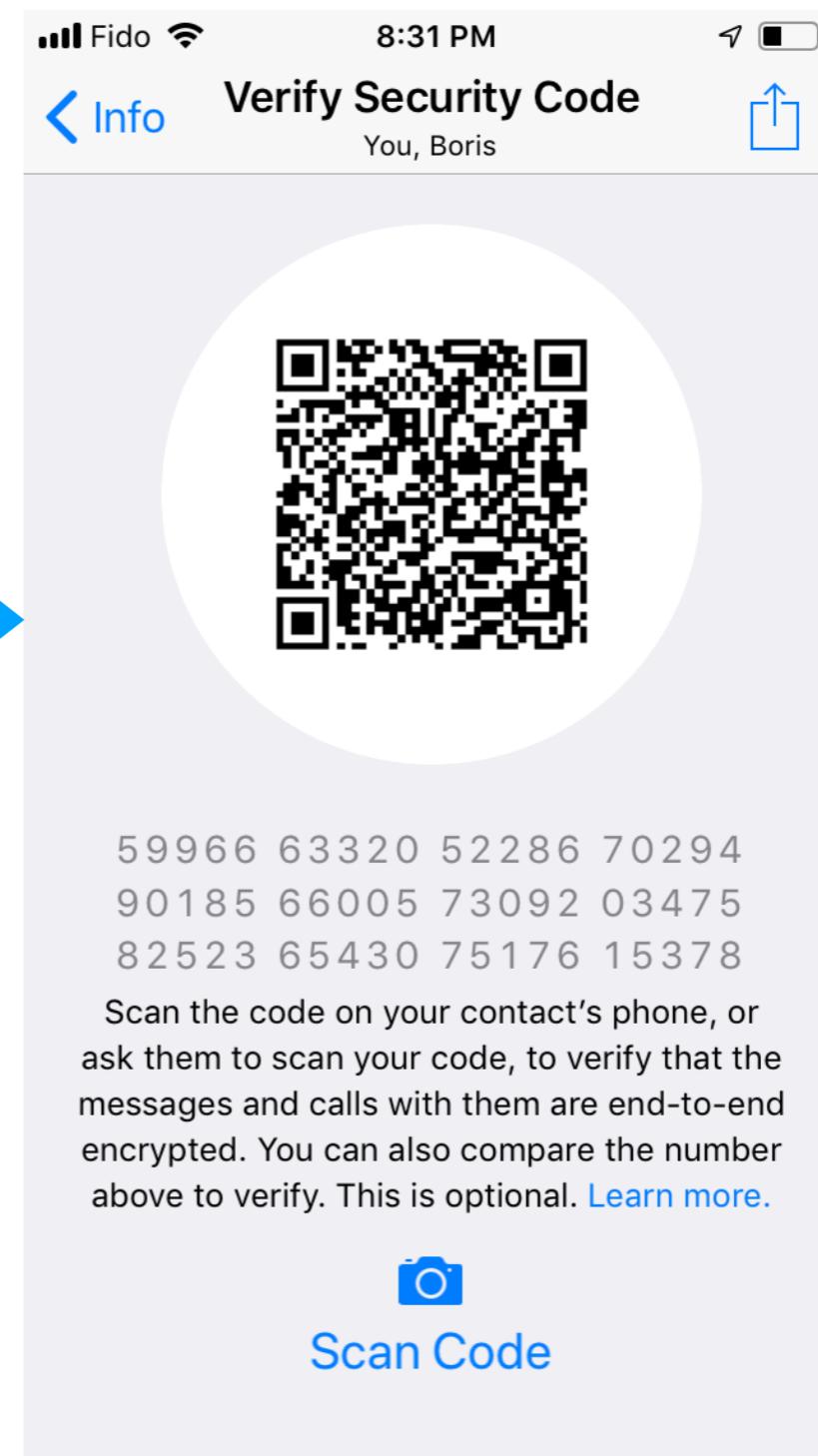
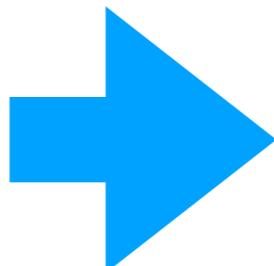
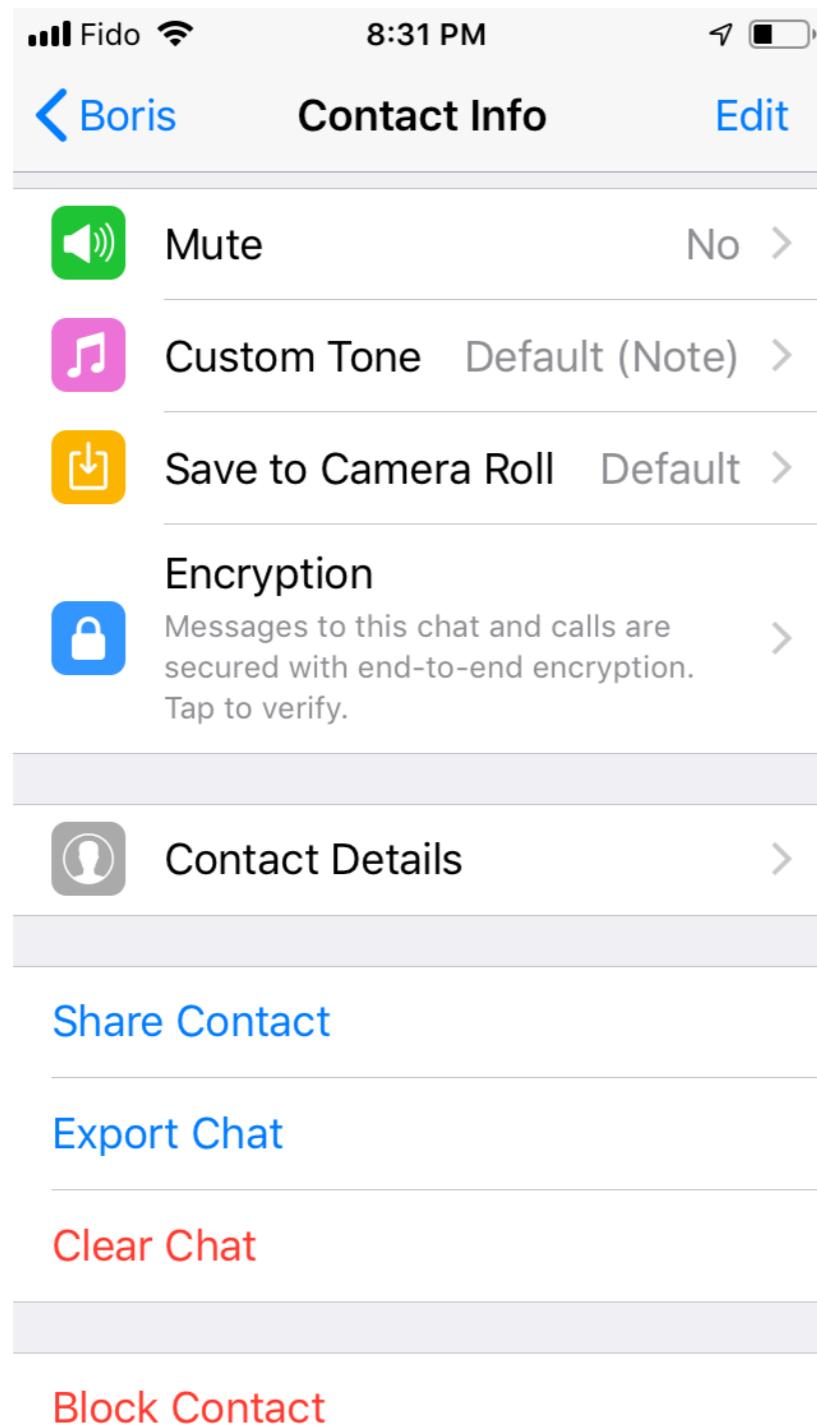
- Encryption: Guarantees “what” is said
- Authentication: Guarantees “who” said it

# **Example: WhatsApp**

# WhatsApp Security

- Information is guaranteed to be end-to-end encrypted.
- Are the public keys matching ?
- Who are the people behind the phone numbers?
- What are the phone numbers behind the names?

# WhatsApp verify keys



# No phone numbers under contacts

The screenshot shows a messaging application interface. At the top, there's a search bar with a magnifying glass icon and the placeholder text "Search or start new chat". To the right of the search bar are three icons: a circular arrow, a speech bubble with a plus sign, and a vertical ellipsis. Below the search bar, there are three message entries:

- Elizabeth** [REDACTED] Yesterday  
Photo No phone numbers shown Fishing Possible
- Ka** [REDACTED] Yesterday  
Окей
- Ana** [REDACTED] Yesterday  
I'll let you know by the weekend if that's ok

In each message entry, the contact name is followed by a redacted phone number (represented by a black rectangle). A blue status bar contains the text "No phone numbers shown". A red status bar contains the text "Fishing Possible". The timestamp "Yesterday" is at the end of each message entry. Photo placeholders are shown next to the names.

attacker

# Shoulder Surfing

victim



# Phone Numbers

- People take over phone numbers.
- The new owner of the phone number can impersonate the person.

# “Auth”

- Authentication (identification, ID)
- Authorization (ACL rules)
- OAuth2, Google Sign-In

# User Identification

- Identify user by cookies
- By localStorage and sessionStorage
- By window.name
- By auto completion of (hidden) form fields.
- By browser “fingerprint”

# Canvas Fingerprint

- Different graphics cards implement Canvas interface differently.
  - Javascript draws text with the font and size of its choice
  - Adds background colours.
  - ToDataURL method to get the canvas pixel data
  - Hash the result

# Canvas Fingerprint

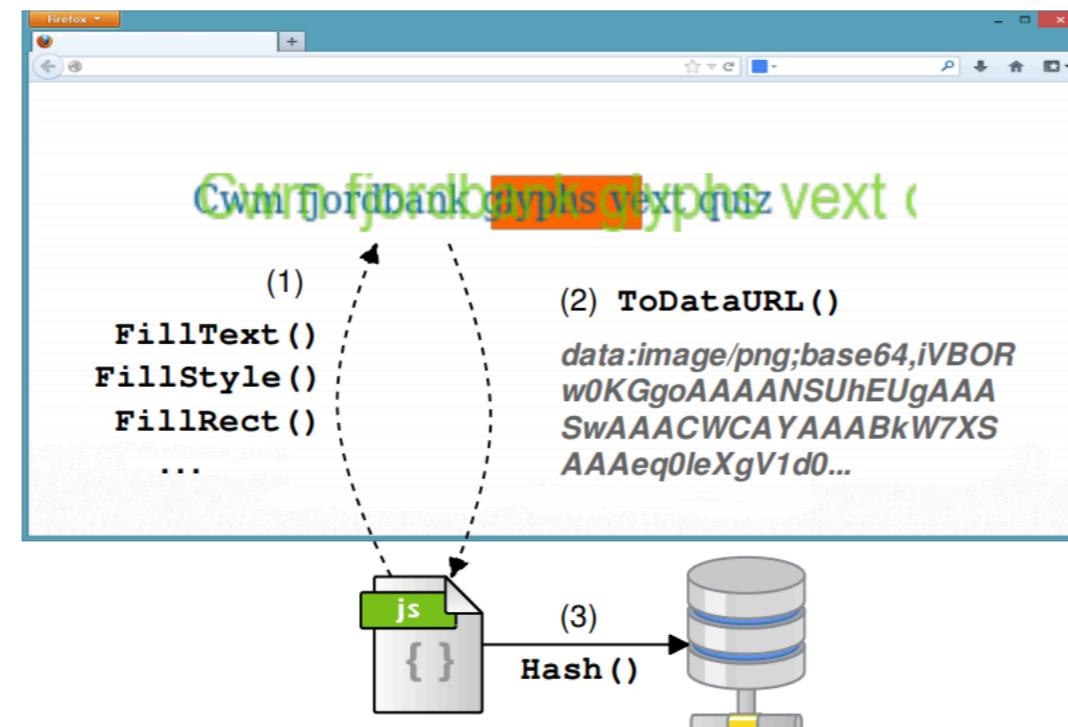


Figure 1: Canvas fingerprinting basic flow of operations

**Paper: “The web never forgets”**

# Scripts in the wild

## addthis.com

Fingerprinting script	Number of including sites	Text drawn into the canvas
ct1.addthis.com/static/r07/core130.js	5282	Cwm fjordbank glyphs vext quiz, ☺
i.ligatus.com/script/fingerprint.min.js	115	http://valve.github.io
src.kitcode.net/fp2.js	68	http://valve.github.io
admicro1.vcmmedia.vn/fingerprint/figp.js	31	http://admicro.vn/
amazonaws.com/af-bdaz/bquery.js	26	Centillion
*.shorte.st/js/packed/smeadvert-intermediate-ad.js	14	http://valve.github.io
stat.ringier.cz/js/fingerprint.min.js	4	http://valve.github.io
cya2.net/js/STAT/89946.js	3	ABCDEFGHIJKLMNOPQRSTUVWXYZ abcdefghijklmnopqrstuvwxyz0123456789+/-
images.revtrax.com/RevTrax/js/fp/fp.min.jsp	3	http://valve.github.io
pof.com	2	http://www.plentyoffish.com
*.rackcdn.com/mongoose.fp.js	2	http://api.gonorthleads.com
9 others*	9	(Various)
TOTAL	5559 (5542 unique <sup>1</sup> )	-

# More info

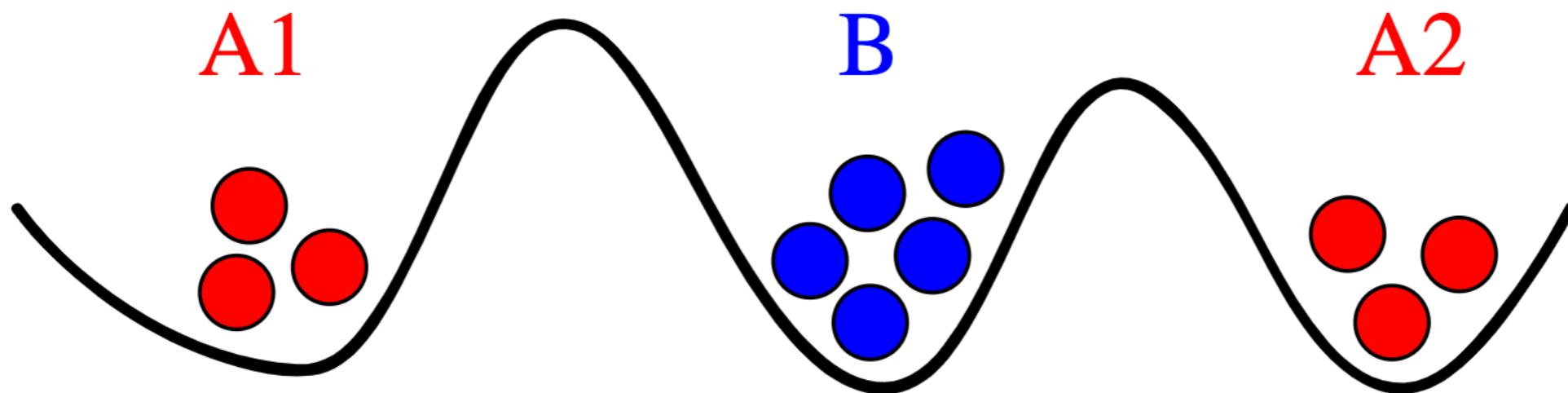
## **The Web Never Forgets: Persistent Tracking Mechanisms in the Wild**

Gunes Acar<sup>1</sup>, Christian Eubank<sup>2</sup>, Steven Englehardt<sup>2</sup>, Marc Juarez<sup>1</sup>  
Arvind Narayanan<sup>2</sup>, Claudia Diaz<sup>1</sup>

# Theory

- The Two Generals Problem
- The Byzantine Generals Problem
- Proved unsolvable

# Two Generals



Army “B” intercepts coordination messages between A1 and A2 armies.

# Shortened URLs

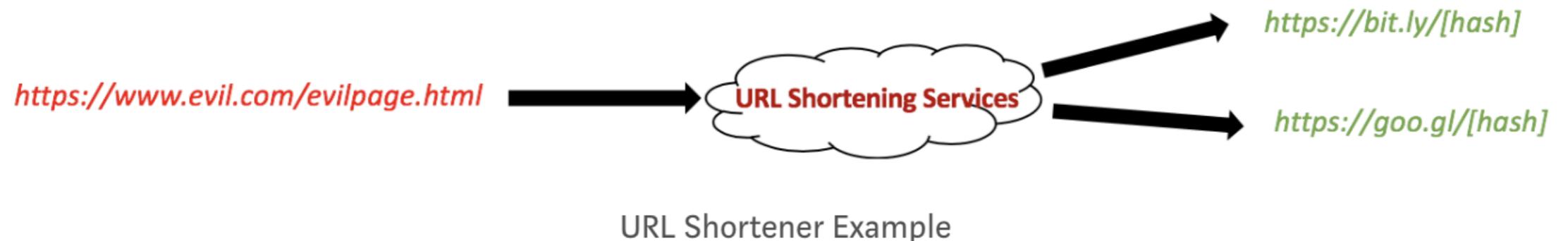
- Security by obscurity
- URLs to Google Drive and Dropbox documents
- If you have the URL, you are authenticated.

# URL Shorteners

- Convenient:
  - Can be pasted where there's a length limit.
  - Can be incorporated into a text message, in the middle of a sentence.
  - Can detect vulnerable redirects, and act like an “anti-virus”.
- Hide bad urls in good urls.
- Guess URLs (brute force search)
- Collect Analytics (Tracking)
- Dangerous if the issuing service is hacked.
- Deliberate malicious url shortener services.
- `http://` supporting shorteners + man-in-the-middle bypass SSL.

# Shortening as Obfuscation

- Shortener hides malicious URLs and allows to bypass checks.
- Can be used for SPAM emails.
- Mitigate: [checkshorturl.com](http://checkshorturl.com)



# End-to-end encryption fails

- Many web services rely on embedding an authorization key in the URL.
- The key is long and hard to guess.
- But, the shortened URL of it is easy to guess.

# Trust to Shortener Service

- The shortener service knows the same location info about you that the target site will know.
- IP & geo location, HTTP Referrer, your browser fingerprint.
- Can add its own cookie, to link all your sessions into one profile. (Knows all websites that you visit.)
- Mitigation?
  - Block cookies for known shortener services.
  - Still possible through browser fingerprinting.

# Shortener Service Hacked

- The service cli.gs got hacked (2009)
- 2.2 million URLs got redirected to a different destination.
- (according to their own blog report, later.)
- Affects all existing links, that previously were considered clean by users.

# Malicious Shortener Services

- Operate a legitimate looking URL Shortener Service (USS)
- Redirect only some clients selective to malicious pages.
  - Clients using old browsers with known vulnerabilities.
    - Advanced users would not notice, because they use the latest browsers.
    - Clients specifically targeted.
    - Sell portion of traffic to botnets.

# SSL Bypass

- If the target website is https://
- But, the shortener is http://
- Then a MITM can intercept, using SSLStrip.

# SSRF

- USS services usually check the links submitted, to avoid linking to other competing USS.
- They run the check from their own server.
- If a certain website is vulnerable through a URL request, an attacker can get the USS service to issue the request on his behalf, thereby obfuscating his identity.

# Mass Surveillance

Observation by @dakotasmith :

“Twitter's URL shortener is a Colombian domain.  
Facebook's domain is in Montenegro. Bit.ly is Lybian. NSA  
only targets international traffic.”

# Brute Force Shorteners

- What can be gained by trying out random short URLs?
- This paper is from April, 2016.
- Exposed OneDrive and Google Maps short links.

## Gone in Six Characters: Short URLs Considered Harmful for Cloud Services

Martin Georgiev  
*independent* \*

Vitaly Shmatikov  
*Cornell Tech*

# Scan against official API

- Bit.ly API:
  - 2.6 queries / second, indefinitely.
  - 227 queries / second, results in temporary blocking of IP address.
- Google Maps API ([goo.gl/maps](http://goo.gl/maps)):
  - 1M queries per day, for free. Higher quote if paid.

# Brute Force Algorithm

- Generate 100 million random shortened URLs.
- Hit the API, find out which of them are valid
- Also, estimate density
- Math:  $\text{density} = \# \text{ found} / 100 \text{ M.}$
- Given the density, estimate total URLs encoded.

# OneDrive Results

- Found: 42 million working URLs on bit.ly 6 char space.
  - 2000 of these are to OneDrive, and are live
- Found: 30 million working URLs on bit.ly 7 char space
  - 25,000 of these are to OneDrive, and are live
- Other user's files on OneDrive can be traversed, given one good URL.
- 300 OneDrive documents allow *write* access

# Traverse OneDrive

- <http://1drv.ms/1xNOWV7> (operated by [bit.ly](#))
- [https://onedrive.live.com/?cid=485bef1a80539148!115&ihtint=folder,xlsx&authkey=!AOOp2TqTTSM5](https://onedrive.live.com/?cid=485bef1a80539148&id=485BEF1A80539148!115&ihtint=folder,xlsx&authkey=!AOOp2TqTTSM5)
- Use this parameters and construct the root URL.
- Then, parse the HTML of the root page, to find links to other documents.

# OneDrive Traverse Results

- bit.ly 6 char space:
  - Traversing from 2000 live OneDrive urls led to:
  - 227,000 more OneDrive readable documents.
- bit.ly 7 char space:
  - traversing 22k documents, led to:
  - 1M more documents.

# Exploit Delivery

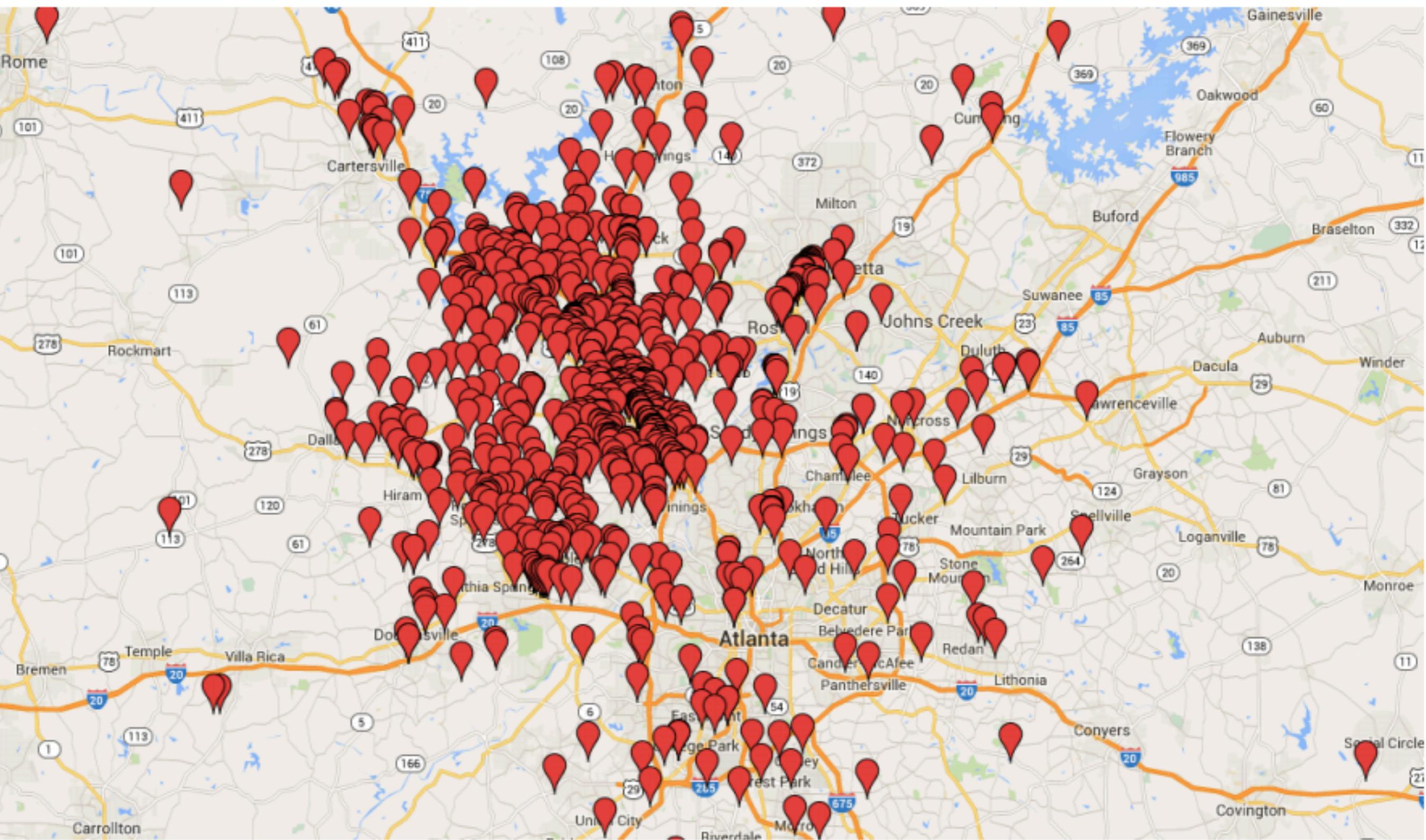
- OneDrive editable folders, allow to deliver exploits
- Exploits land on user's hard drives.
- 2000 OneDrive documents, led through traversal to 150 editable folders.
- Microsoft's Response: fixed the problem by removing shortened urls altogether, but stated that the report is not a vulnerability.

# Google Maps Results

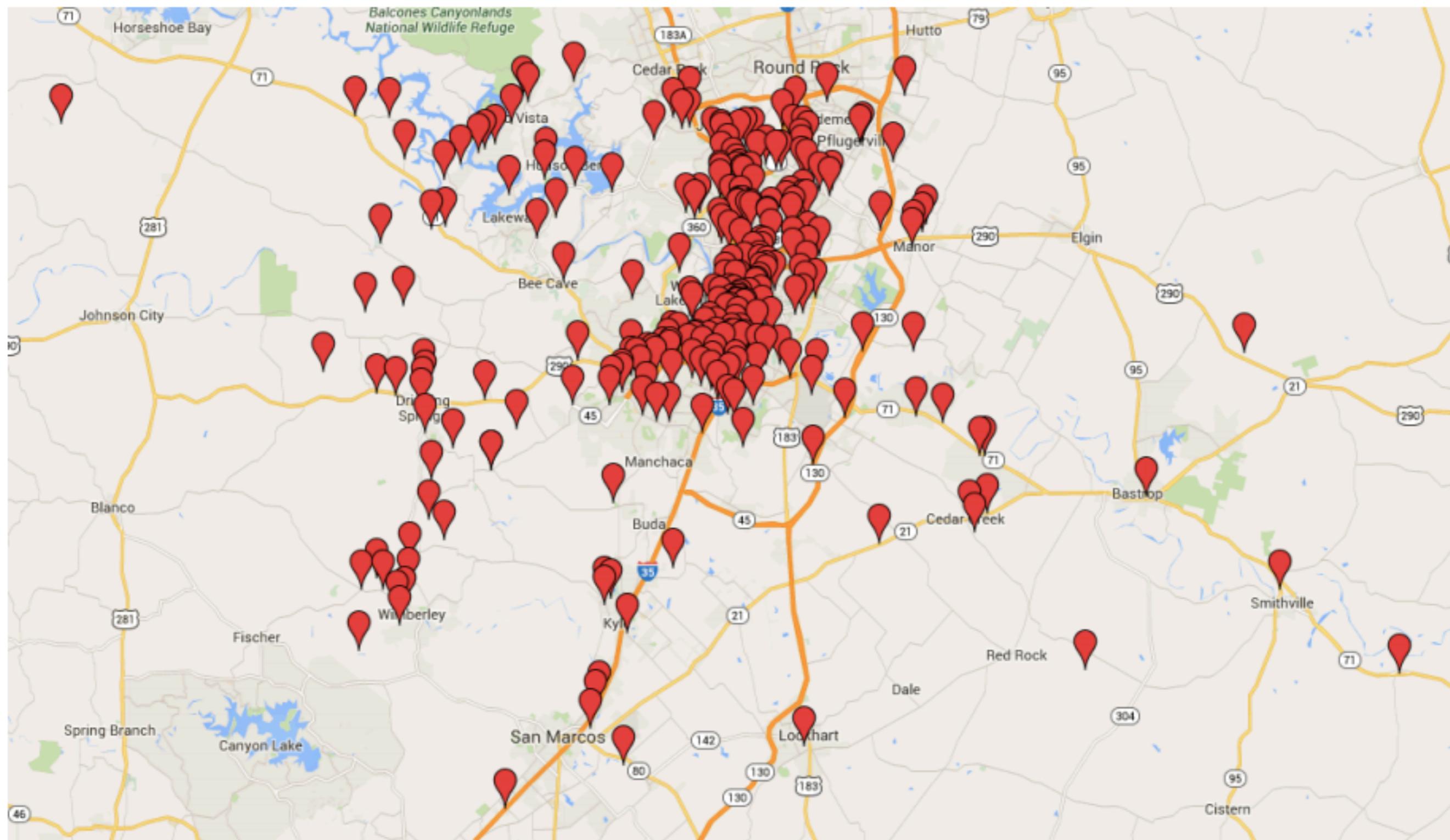
- Found: 24 million working URLs on google maps.
- 10% of them are for directions.
  - Sensitive locations: special clinics (cancer, mental illness), addiction, abortion, gentlemen's clubs, correctional facilities, payday lenders.
  - Identify the individual using one of the “home” endpoints.
- 90% are individual’s location.

# Google Maps Results

- Found: 24 million working URLs on google maps.
- 90% are locations.
- 10% of them are for directions.
  - Sensitive locations: special clinics (cancer, mental illness), addiction, abortion, gentlemen's clubs, correctional facilities, payday lenders.
  - Identify the individual using one of the “home” end-points.
  - Create a map of who has visited whom.
  - Find all clients of a particular business. Example: car towing.
  - Google API for short URLs reveals *when* the URL was created.



*Figure 5: Locations associated with D & D Autows Inc.*



*Figure 4: Locations associated with a single user in Austin, TX.*

# Responses to CVE reports

- Microsoft:
  - Fixed the problem by removing shortened urls altogether,
  - But: stated that the report is not a vulnerability.
- Google:
  - fixed, by increasing URL length to 12 characters.

# One-Time-Password

- Hash Chain
- Original idea by Leslie Lamport
- Hash OTP, HMAC-OTP, Time-OTP

# Hash Chain

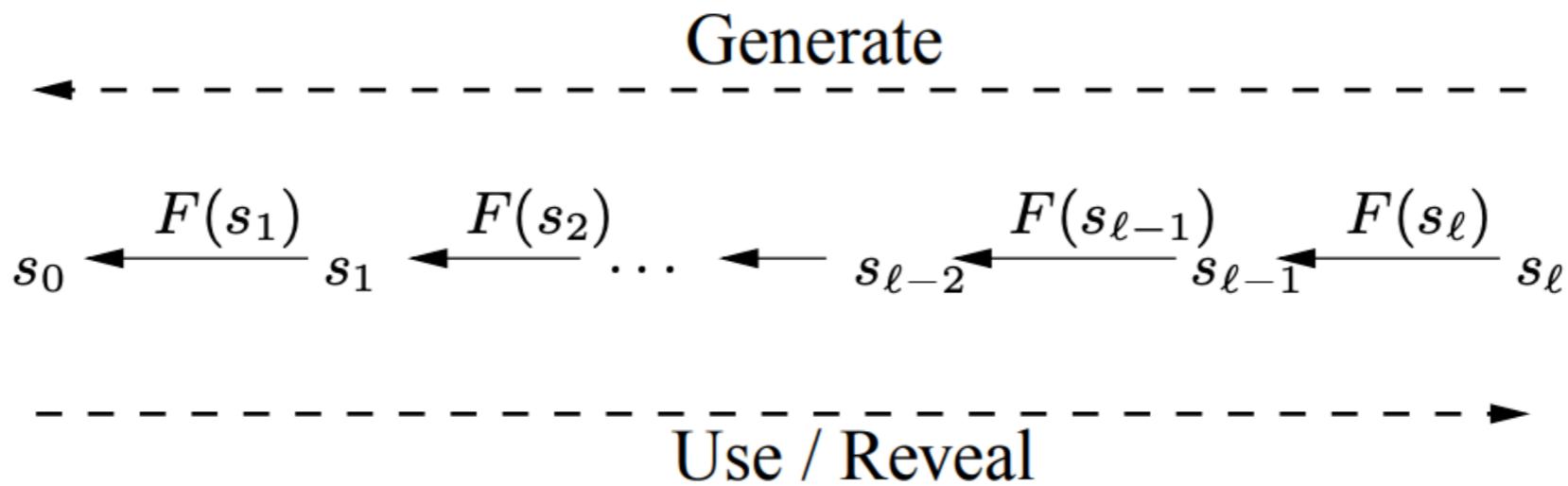


Figure 1: One-way chain example. The sender generates this chain by randomly selecting  $s_\ell$  and repeatedly applying the one-way function  $F$ . The sender then reveals the values in the opposite order.

# Key Derivation

- Authentication = identification
- Key Derivation = Authentication + end-to-end encryption

# Key Stretching

- scrypt
- bcrypt
- PBKDF
- input: short password, output: long enough key
- must be memory, computation and space expensive.

# Biometric Authentication

- Touch ID
- Face ID
- Typing speed pattern
- Voice
- Imprecise answers to a list of questions

# Biometric Authentication

- Calculate delta between expected input and received input
- If delta is small, user is authenticated.

# Biometric Key Derivation

- Idea: Quantize input and add a correction
- Registration:
  - $F_1$  is the input, generate key
  - Also, calculate  $F_2 = \text{quantize}(F_1)$ , and  $D = F_1 - F_2$
  - Store  $D$  on the server.
- Login:
  - $F_3$  is new input. Calculate:  $\text{quantize}(F_3) + D$ .

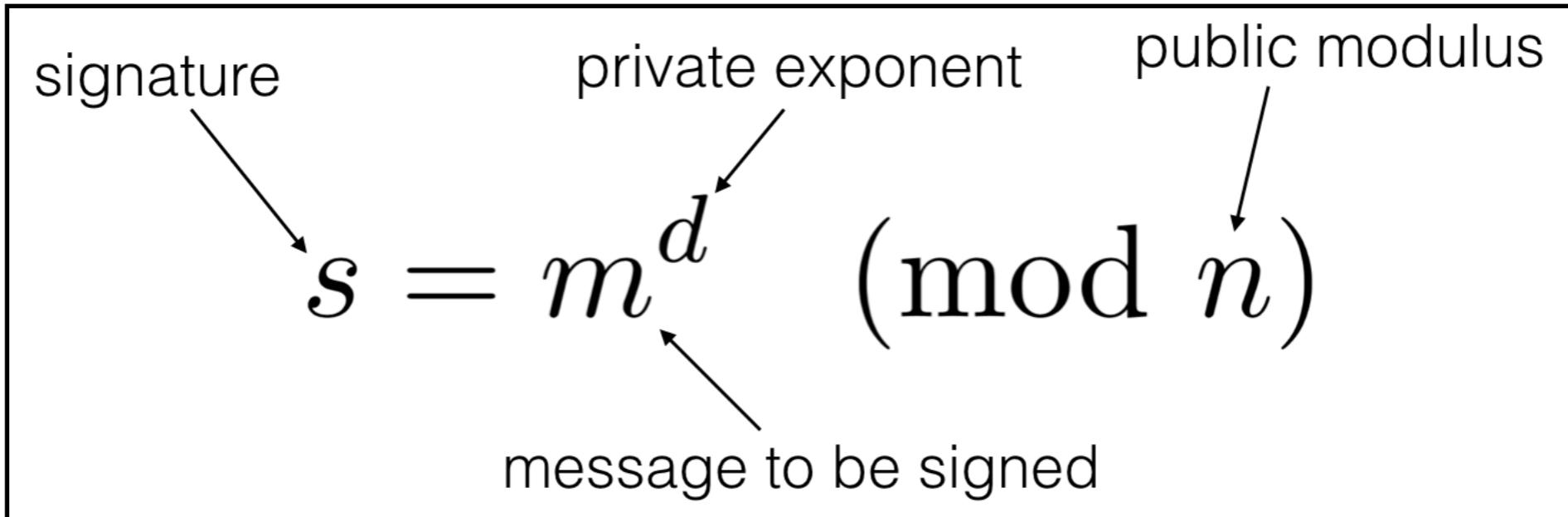
# Digital Signatures

- RSA method
- Elliptic Curve method

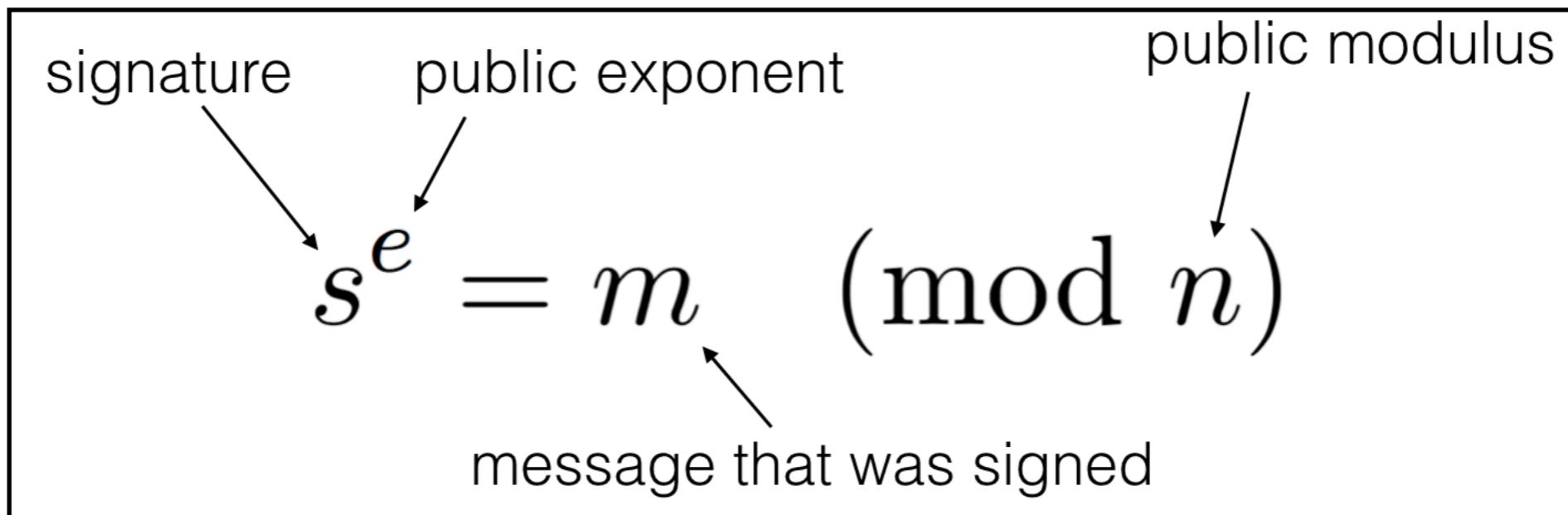
# RSA signatures

- Sender encrypts “foo” with sender’s *private key*
- Recipient decrypts signature with sender’s public key.
- Signature is correct, if recipient sees “foo”.

## Sign



## Verify



# RSA problem

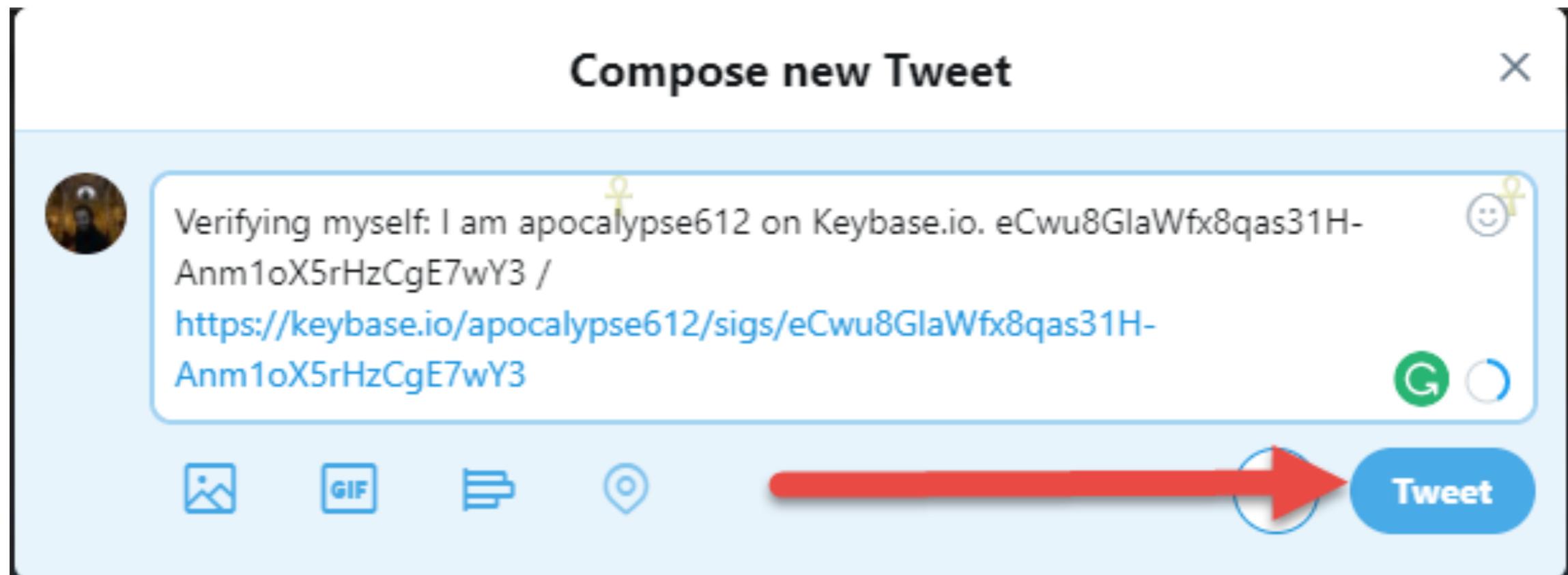
- But: Recipient doesn't know senders public key.

# Solutions

- keybase.io
- PGP servers
- GPG web of trust

# keybase.io

- Verify user's identity based on his membership on social networking websites.
- Twitter, Facebook, etc.



# Bitcoin Transaction Verification

- Uses an Elliptic Curve group secp256k1
- Group of points  $(x_1, y_1), (x_2, y_2), \dots$
- Points satisfy curve:  $y^2 = x^3 + 7 \pmod{p}$ .
- Points P and Q can be combined to produce new point P(+) $Q$  in the set, that also satisfies the curve equation.
- $P + P + P$  is written as  $3 \cdot P$ .
- All points are of the form  $n \cdot P$ .

# Zero Knowledge Proof

- Equation:  $h \cdot Q + R = s \cdot B$
- $h = \text{hash}(m, R)$ ,  $Q$  is the signature,  $B$  is base point
- Signature  $(m, R, s)$  proves that:
  - you know private scalar  $t$  that:
    - $Q = t \cdot B$
    - set  $R = t \cdot B$
    - factor LHS into  $ht \cdot (Q + B)$

# RFID Authentication

- The ID on the card
- Attack: replicate the same ID on another card
- Card cloning
- Magic card: allows to configure any ID
- Master card: opens all doors in a hotel.

# \$2 per card

CHENGKA MF1 S50 Clone Copy Back door Rewritable Blank RFID Card Chinese Magic Card 13.56MHZ Changeable MF S50 1K NFC Card

★★★★★ 5.0 (2 votes) | 8 orders



Anniversary Sale **US \$ 1.66 - 2.07**

For big savings, collect coupons now.

Starts in  
5 days

Price: ~~US \$2.40 - 3 / Set~~

Discount Price: **US \$1.68 - 2.10** / Set **-30%**

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Blank RFID card

# Proxmark3 card



NETWORKING / WIRELESS NETWORKING

## Proxmark3 RDV4 Kit



From: \$300

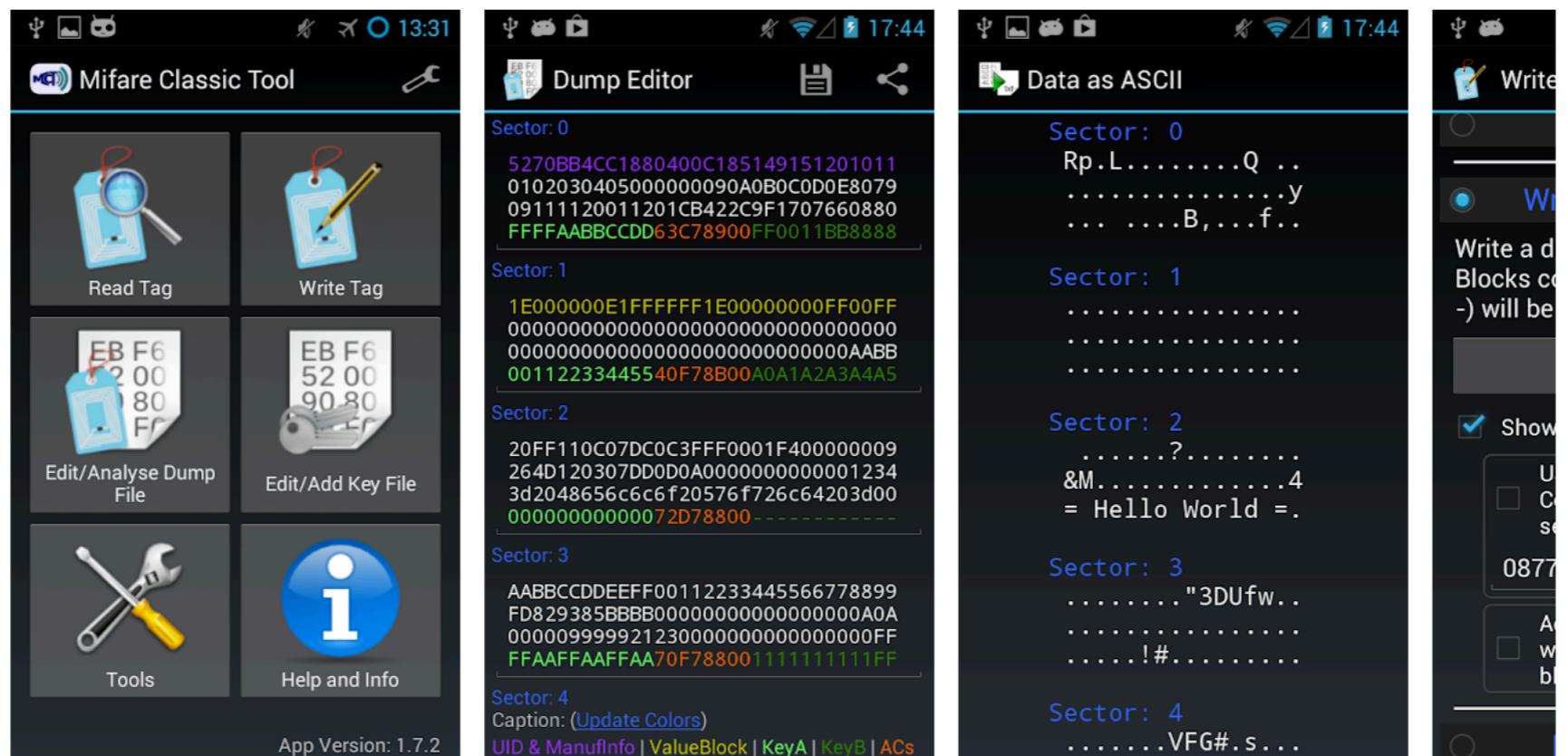
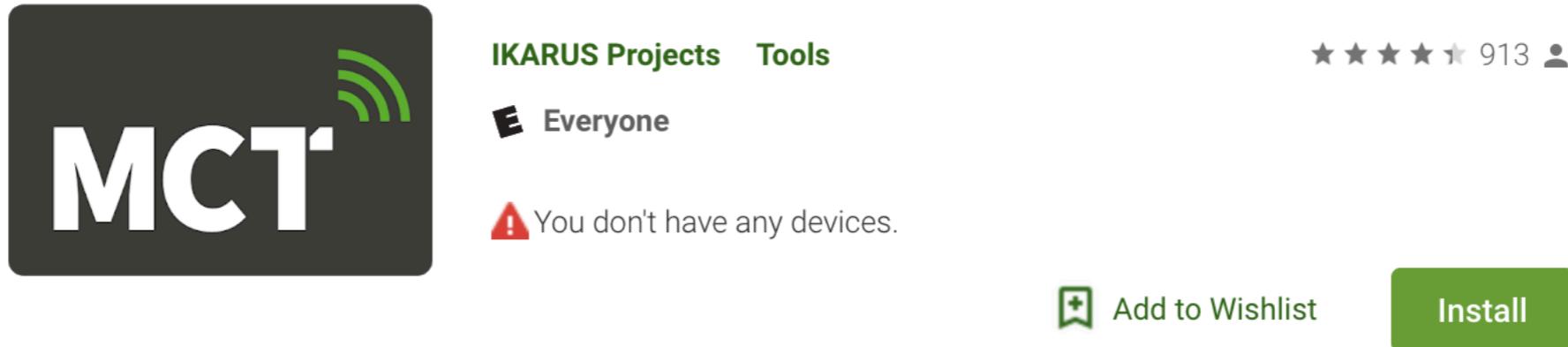


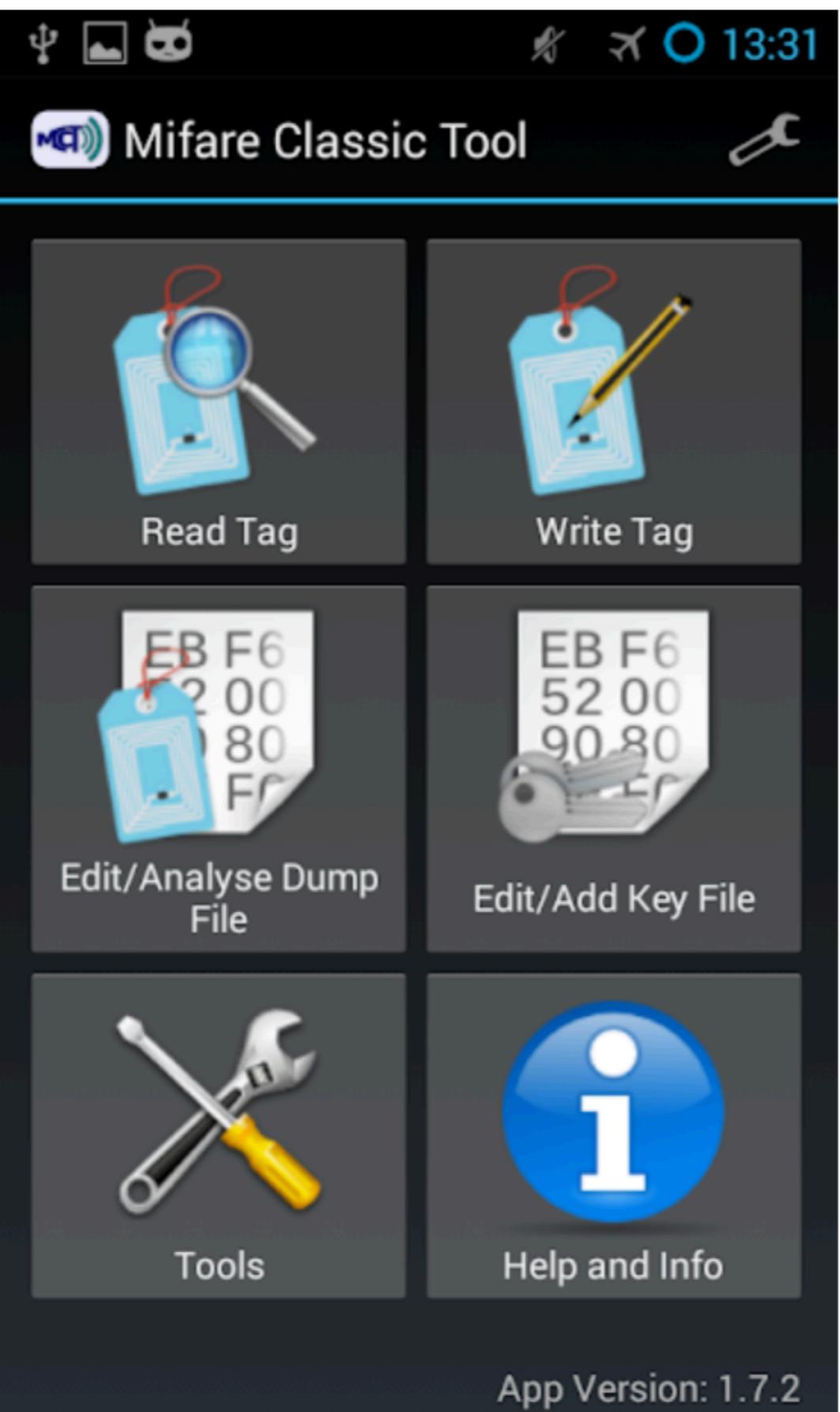
Proxmark3 RDV4 Accessory Pack 

Add for \$100

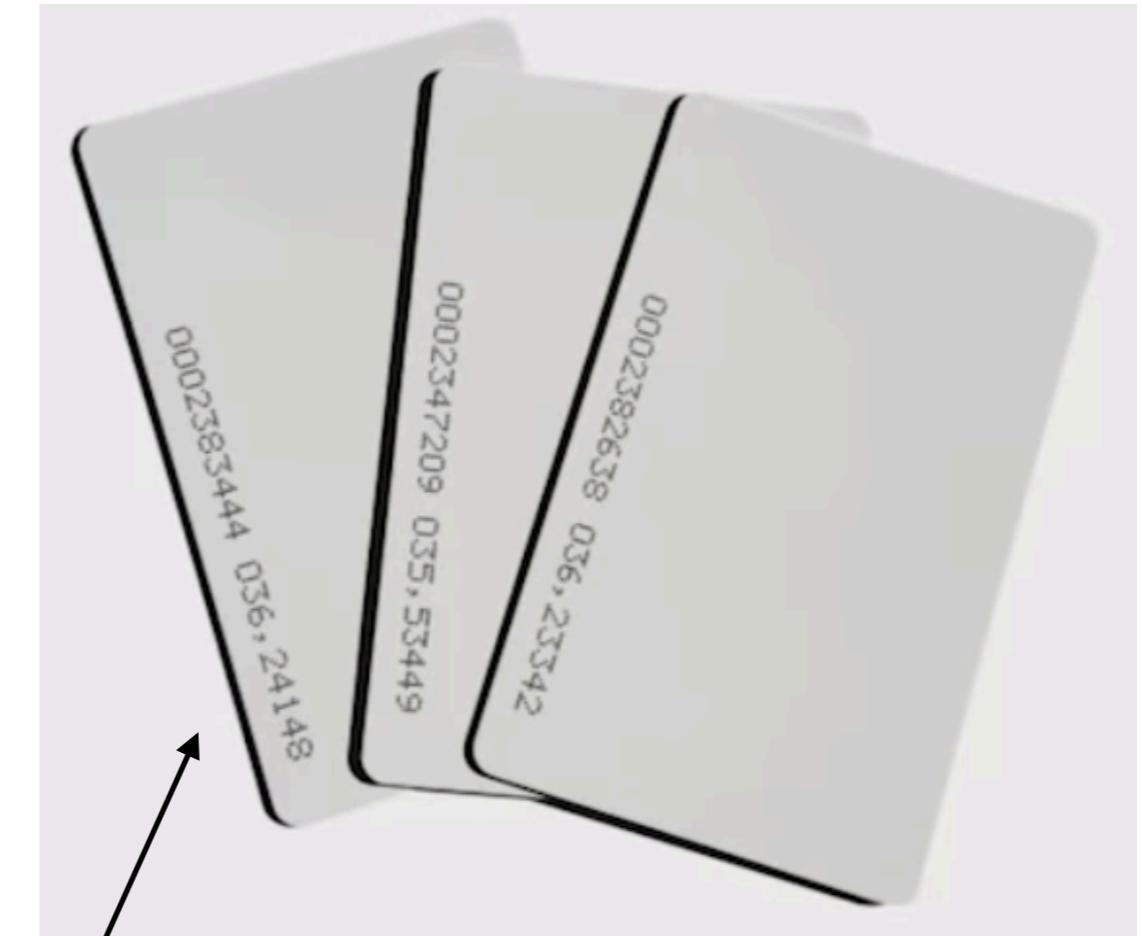
# Mifare Android App

## MIFARE Classic Tool - MCT



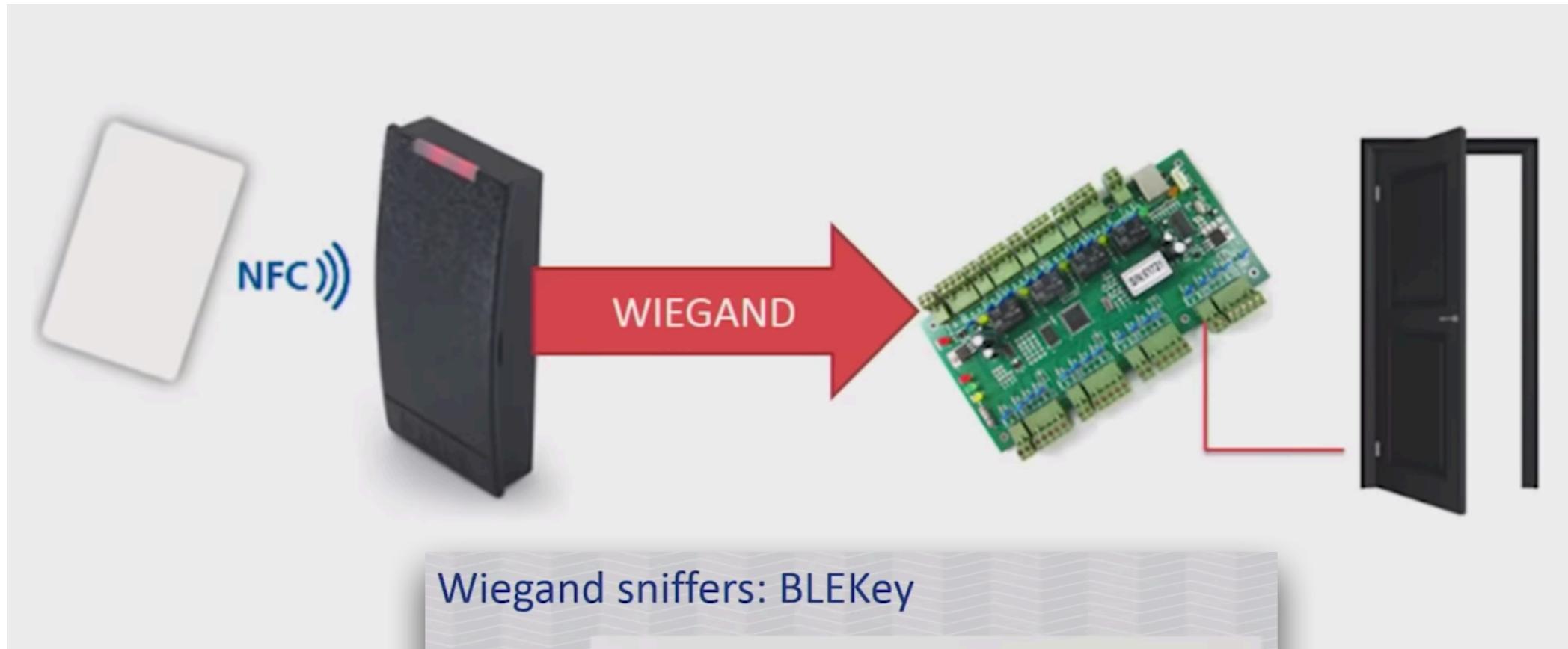


# Photo of a card



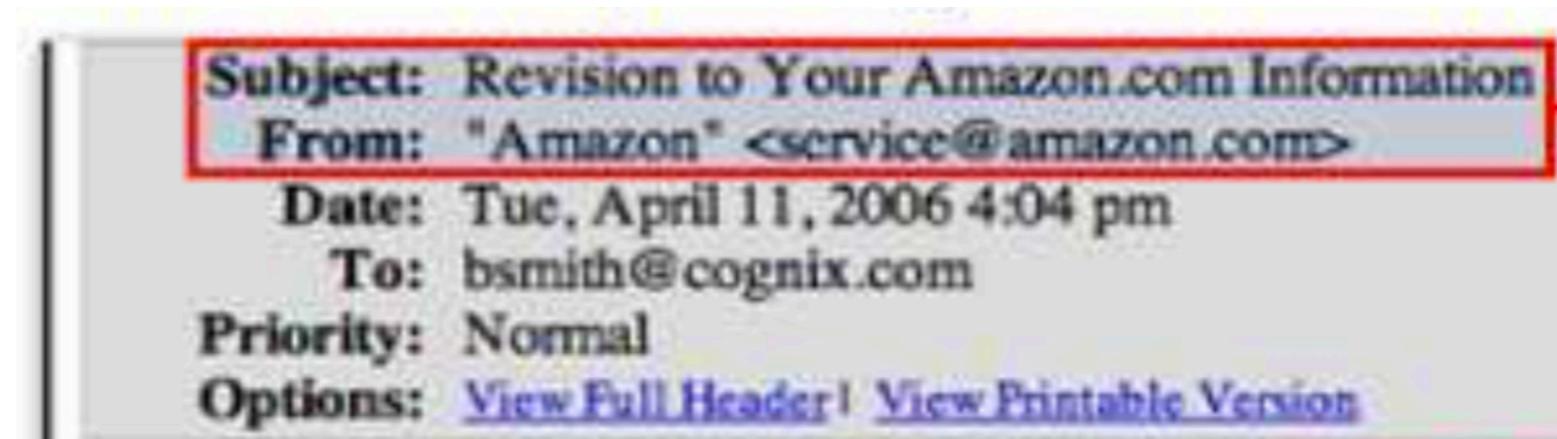
**ID is decimal**

# Capture Reader's output



# Email Fishing

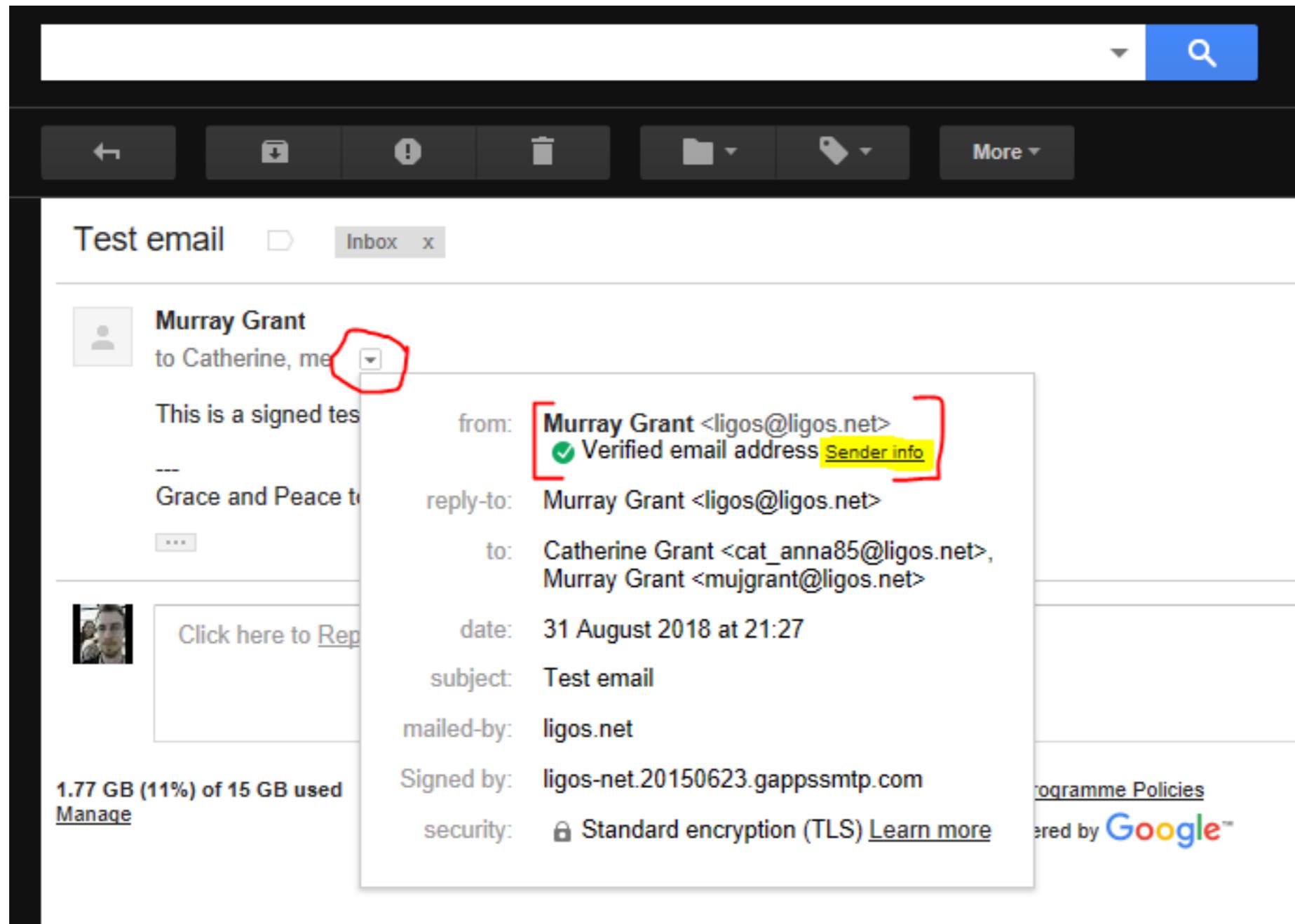
- Impersonate name of a sender
- Impersonate photo of a sender



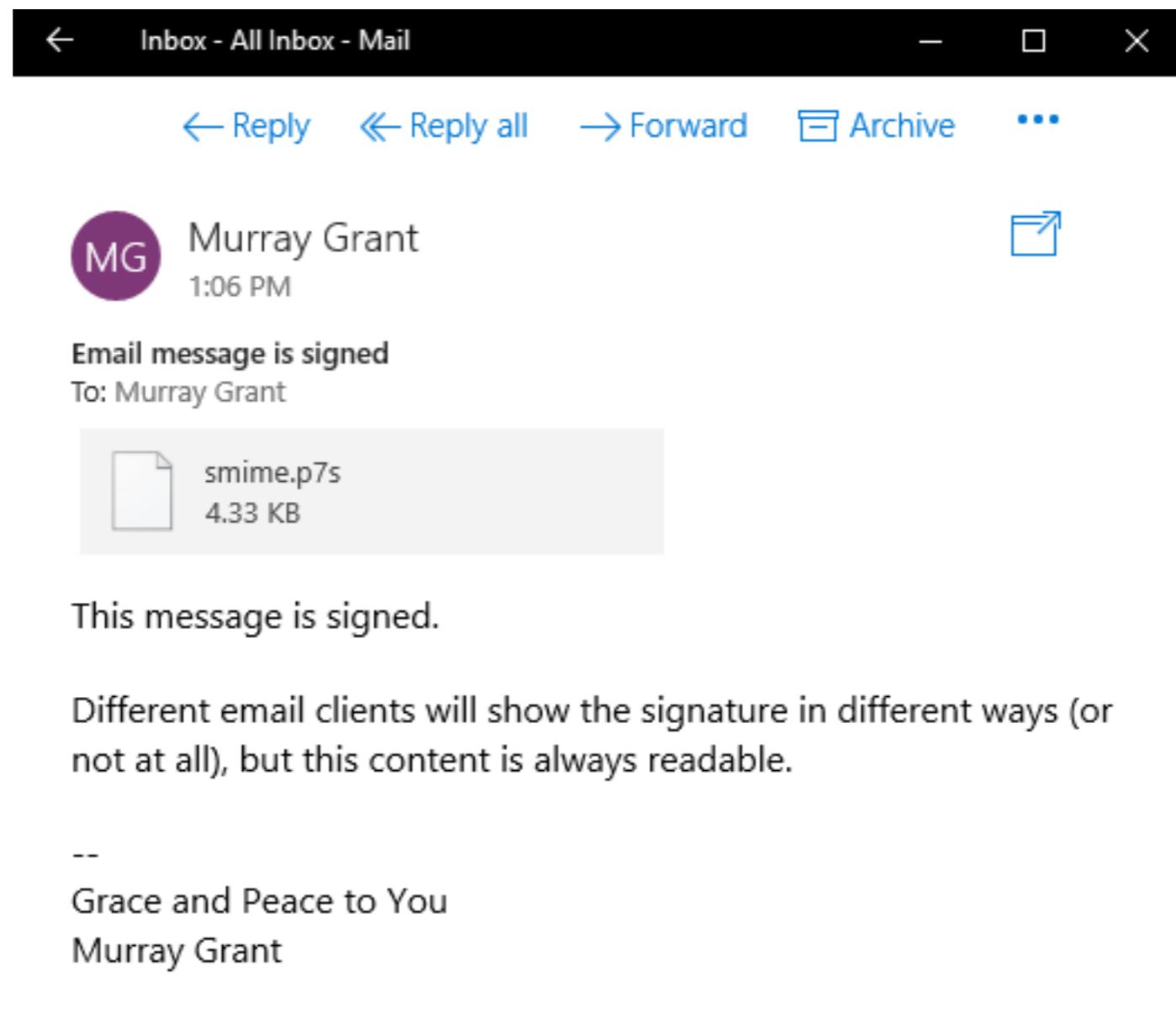
# S/Mime

- Certificate is associated with the sender.
- Recipient must trust Certificate Authority (CA) that signed the certificate.
- To send, you must have the certificate of recipient.
- However, this works well for signing.
- But, it attaches a weird attachment file.

# Gmail S/MIME signature



# Gmail S/MIME signature



# The Problem

- Who is “Murray Grant” ? Any one who can prove this name, can get a certificate like that.
- Chinese Names: there are many people with the same last name such as “Cho”.

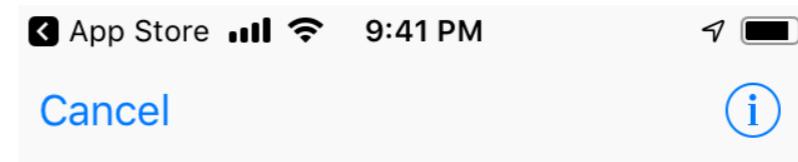
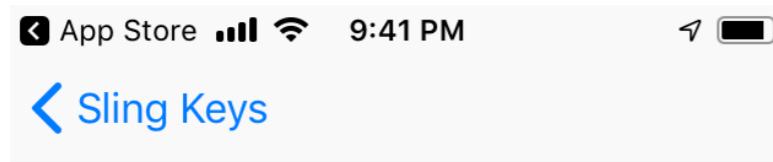
# PGP

- PGP is RSA + standard to embed things in the body of email.
- Requires distribution of public keys. How?
- GPG, and PGP servers
- PGP signing parties.
- Problem: public keys get lost, email becomes unreadable.

# SafeSlinger

- A mobile app to share contacts, in-person. A modern version of a PGP signing party.
- By Carnegie-Mellon university.

# SafeSlinger



How many users in the exchange?

Compare screens  
on 2 devices.

60

This number is used to create a unique group of users. Compare, and then enter the lowest number among

2  
3  
4

OK

Lowest

OK

1	2 ABC	3 DEF
4 GHI	5 JKL	6 MNO
7 PQRS	8 TUV	9 WXYZ
0		✖

# DomainKeys

- A mail server signs outgoing mail.
- Email header “DKIM”.

# DomainKeys

```
DKIM-Signature: v=1; a=rsa-sha1; c=relaxed/relaxed; d=meetup.com;
h=from:sender:to:subject:mime-version:content-type:content-transfer-encoding;
s=s1; bh=7rVlVxx51rnBRZGnOKcysrxyY0A=; b=PCaPuc/qf5XLgGb1T8ZuYRJ
hWChagq6JRiz7EfUrlpejC0yvcenqlKUp17H1I3XA2RzcNSDZRVn5Ig890ZC0tYZ
kkCliVFVGUr7h60sSGR+mEQ6pHIU8BTBSfdIbgqt1mPcL1Pgnk9+5QS7ZFkgo29V
hTuxhdyaqTHR+06M5vJY=
```

```
$ dig TXT s1._domainkey.meetup.com
```

```
s1._domainkey.meetup.com. 300 IN CNAME s1.domainkey.u3863915.wl166.sendgrid.net.
```

```
s1.domainkey.u3863915.wl166.sendgrid.net. 1800 IN TXT "k=rsa; t=s; p=MIGfMA0GCSqGSIb3QEBAQUAA4GNADCBiQKBgQChK0RIGEj4ahkPXIhENbmXC6Cu2Q1eVNDM6nZdrJGR2p4jWYNVGQ/EYQRC35Qu+ncvNvayv8igvCou1A9Y6xso1ls6MCMpT3LjatFo+U+qfMI9Uh6P0sQ+NNS7NAGc0GGl8bAx+mbG0AHgbgrB6DJwAz7uGd0IzjPtPdn5EuQIDAQAB"
```

# Gmail DKIM

body hash

```
DKIM-Signature: v=1; a=rsa-sha256; c=relaxed/relaxed;
d=gmail.com; s=20161025;
h=mime-version:references:in-reply-to:from:date:message-id:subject:to;
bh=0SfhqDZkUTtM2Aqq0EJJAngQdx1dHhjX4xG5mGEhy/Y=;
b=qINF14m9coo6mVB18gukDDcycH3kKD8nEwQ+xfFYU1amXDjV922gPDnIISz/rQsT86
gbjyC1jRuwsffz+POEnMcntz7hqDUA/6BddfInantTi9uZj3gEna4TBjEgE6XHuChqkt
2ME59WIVduUATG39M6OP+PVdoLZL8McLADSJZbeDT1DHOQXsLx/XY+XU42cEXOmtGbYE
lcvoqHGYnd8CAWL1y/9MwfEZ/un8CmozF9mDaUmwZvxxHWqPgzzmKBL/Wez27jdD4YSH
F9HMU3aaFDodWv6tGNVJfsjqM0lod7ihC682WD5xgc17NR7XGv1mq2E/NlvuoztxfIqc
Mibg==
```

signature

selector

```
$ dig TXT 20161025._domainkey.gmail.com
```

```
20161025._domainkey.gmail.com. 300 IN TXT "k=rsa; p=MIIBIjANBgkqhkiG9w0BAQEFAAO
AQ8AMIIIBCgKCAQEAviPGBk4ZB64UfSqWyAicdR7lodhytae+EYRQVtKDhM+1mXjEqRtP/pDT3sBhazkmA48n2k
5NJUyMEo08nc2r6sUA+/Dom5jRBZp6qDKJ0wjJ5R/0pHamlRG+YRJQqR" "tqEgSiJWG7h7efGYWmh4URhFM9k
9+rmG/CwCgwx7Et+c80MlngaLl04/bPmfjpjdEyLWyNimk761CX6KymzYiRD Nz1M0J0J70zFaS4PFbVLn0m5mf0
HVNtBpPwWuCNvaFVflUYxEyb1bB6h/oW0PGbzoSgtRA47SHV53SwZjIsVpbq4LxUW9IxAEwYzGcSgZ4n5Q8X8T
ndowsDUzoccPFGhdwIDAQAB"
```

domain key

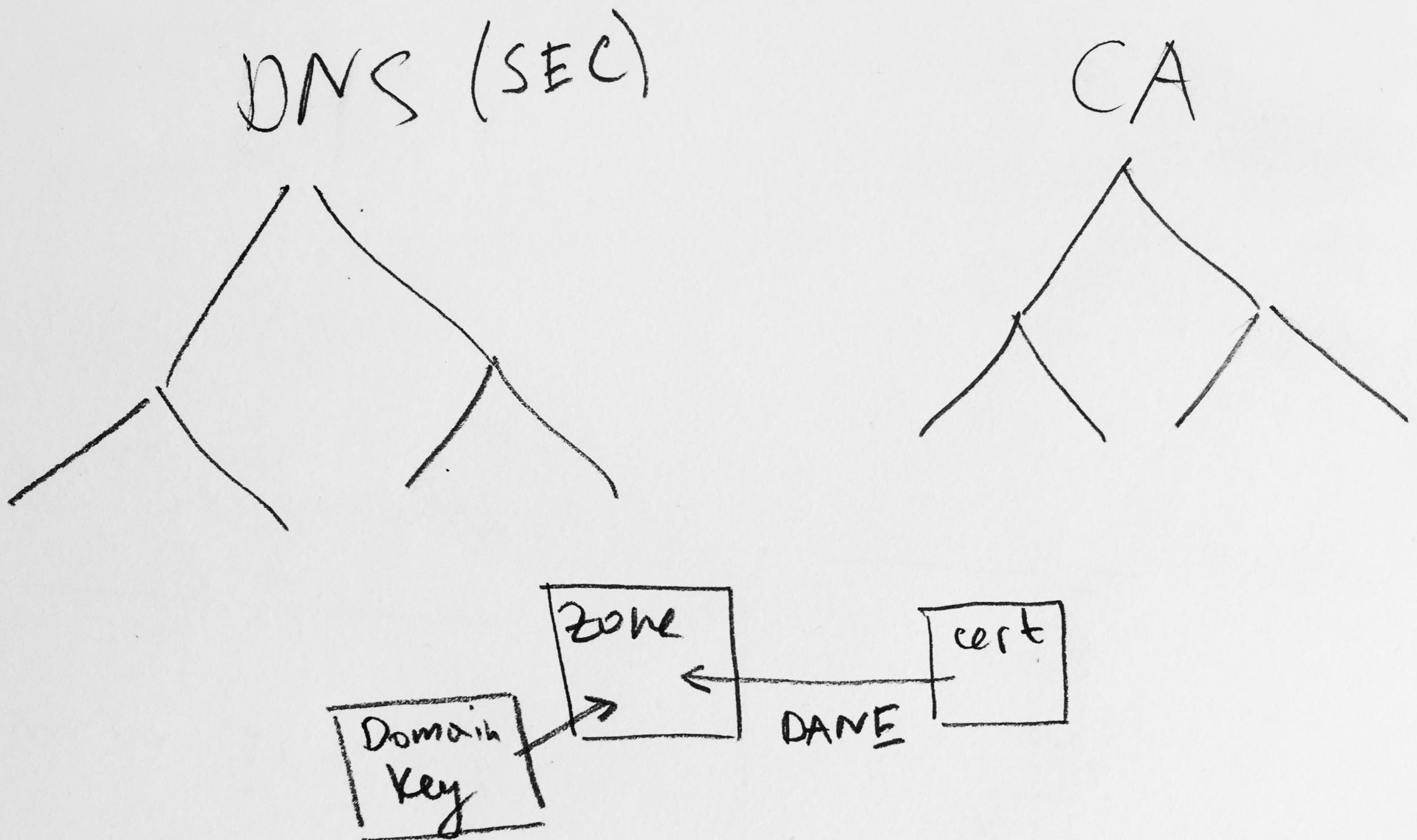
# **But:**

- Relies on the domain's DNS record to provide the public key.

# Server Authentication

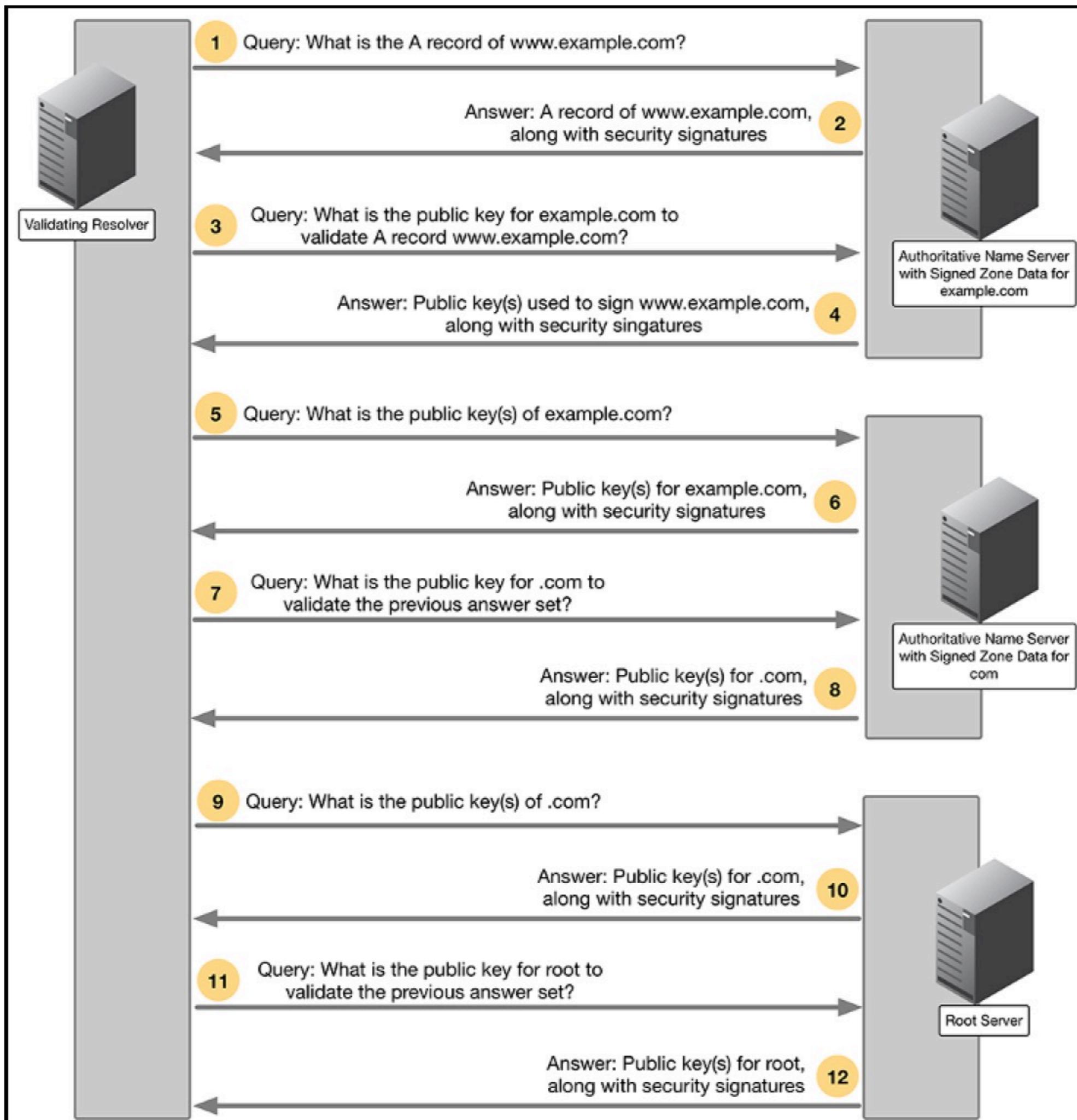
- User authenticates to server?
- Or: server authenticates to user.

# DNS, DNSSEC, Certificate Authorities



# DNSSEC

- Each zone content is signed by parent DNS server.
- Lookup IP for example.co.il
- Verify against signature in DNS for .co.il



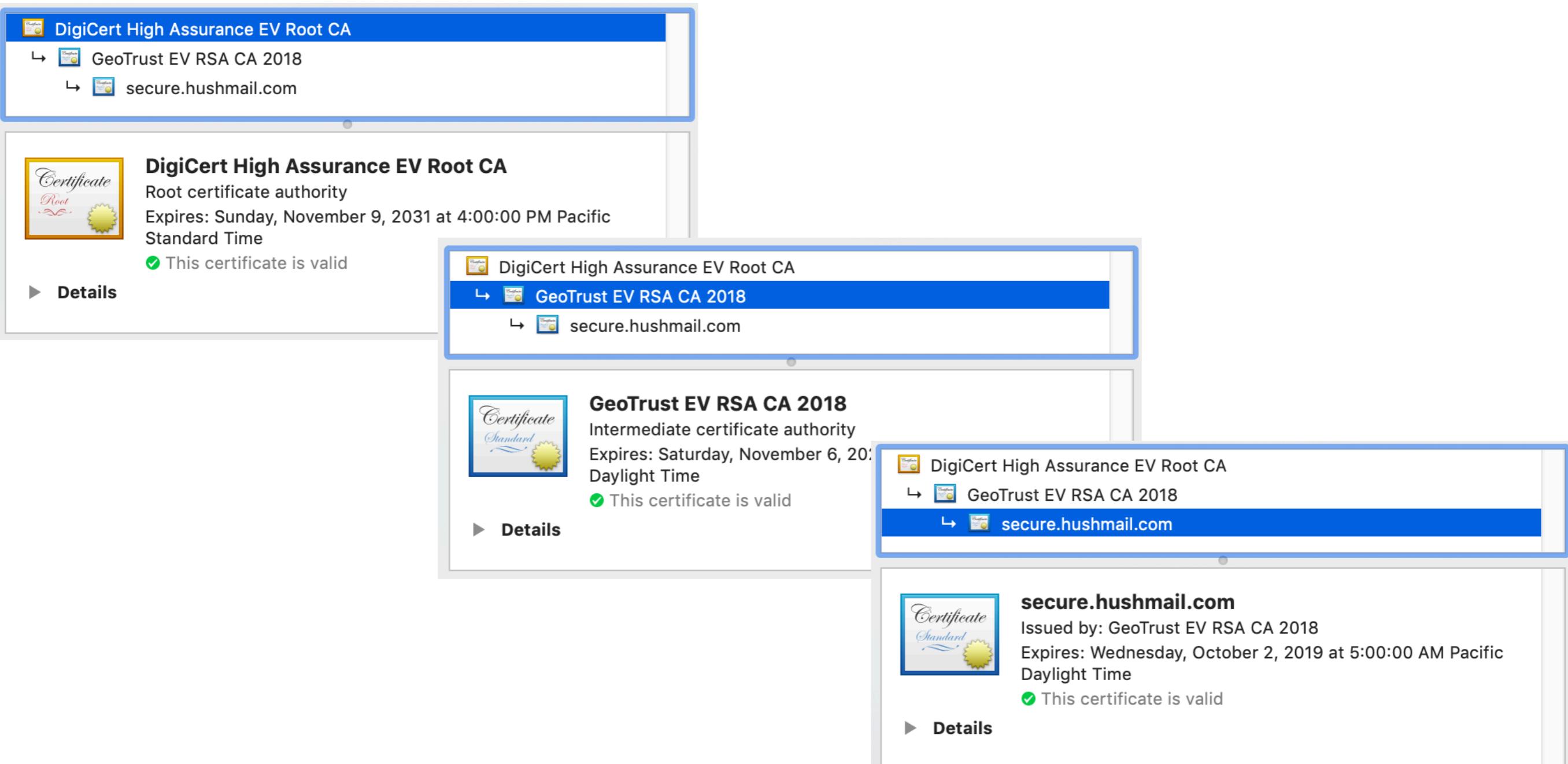
# Certificate Authorities

- Certificates signed by a tree of certificate authorities.
- Example: VeriSign, Comodo

# Chain of certificates



# Trust is placed at root signer





Mac

iPad

iPhone

Watch

TV

Music

Support



# List of available trusted root certificates in iOS 12, macOS 10.14, watchOS 5, and tvOS 12

## Trusted certificates

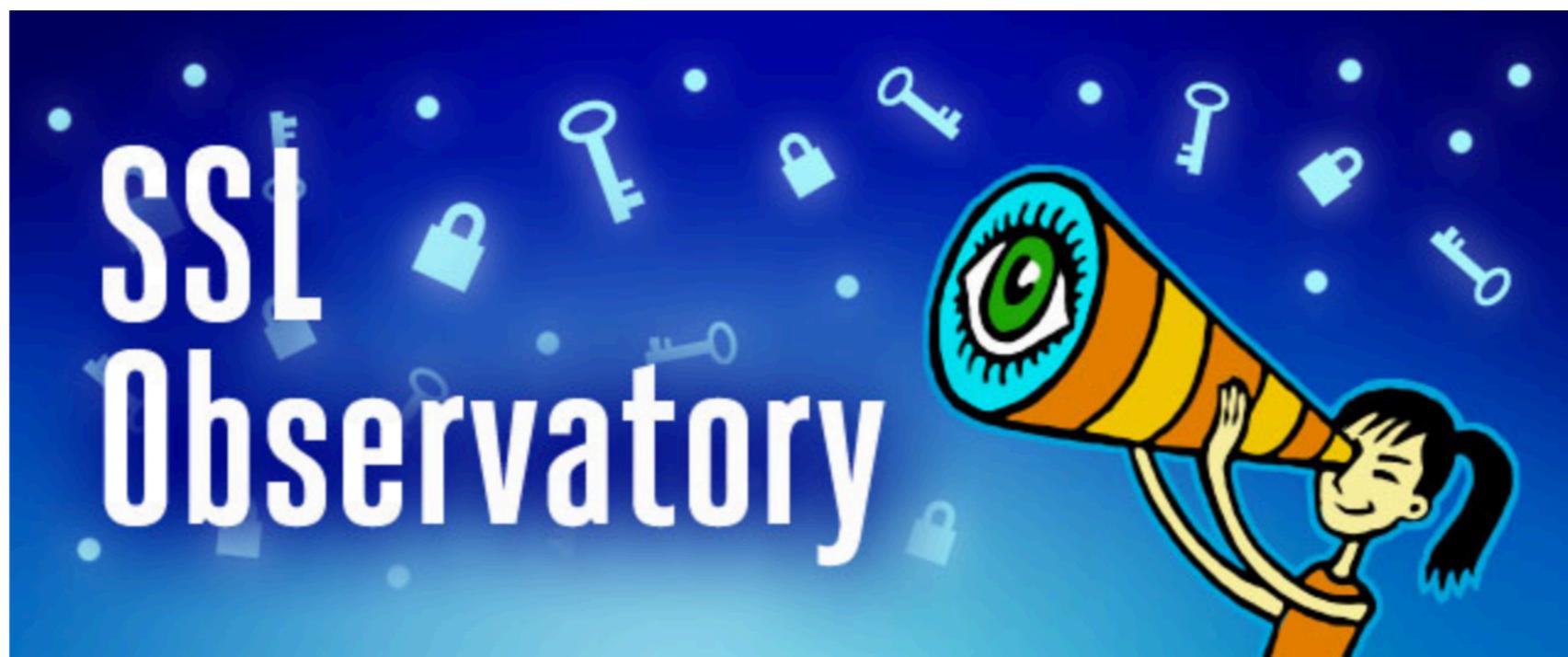
Certificate name	Issued by	Type	Key size	Sig alg	Serial number	Expires	EV policy	Fingerprint (SHA-256)
AAA Certificate Services	AAA Certificate Services	RSA	2048 bits	SHA-1	01	23:59:59 Dec 31, 2028	Not EV	D7 A7 A0 FB 5D 7E 27 31 D7 71
StartCom Certification Authority	StartCom Certification Authority	RSA	4096 bits	SHA-1	01	19:46:36 Sep 17, 2036	1.3.6.1.4.1.23223.2 1.3.6.1.4.1.23223.1.1.1	C7 66 A9 BE F2 D4 07 1C 86 3A 31 AA
Izenpe.com	Izenpe.com	RSA	4096 bits	SHA-1	06 E8 46 27 2F 1F	08:27:25 Dec 13, 2037	1.3.6.1.4.1.14777.6.1.1 1.3.6.1.4.1.14777.6.1.2	23 80 42 03 CA 45 D8 CD F7
Government Root Certification Authority	Government Root Certification Authority	RSA	4096 bits	SHA-256	00 B6 4B 88 07 E2 23 EE C8 5C 12 AD A6 0E 00 A1	15:59:59 Dec 31, 2037	Not EV	70 B9 22 BF DA 01 3F 4A 34 2E 4E E2 2D 57 9A E5 98 D0 71 CC 5B 00 00 00

# Trust govt of other countries

- CA, AT, AU, BE, BG, BM, BR, CA, CH, CL, CN, CO, CZ, DE, DK, EE, ES, EU, FI, FR, GB, HK, HUU, IE, IL, IN, IS, IT, JP, KR, LT, LV, MK, MO, MX, MY, NL, NO, PL PT, RO, RU, SE, SG, SI, SK, TR, TW, UK, US, UY, WW, ZA
- US. Department of Homeland Security
- US Defence Contractors

# Scan CAs

**The EFF SSL Observatory**



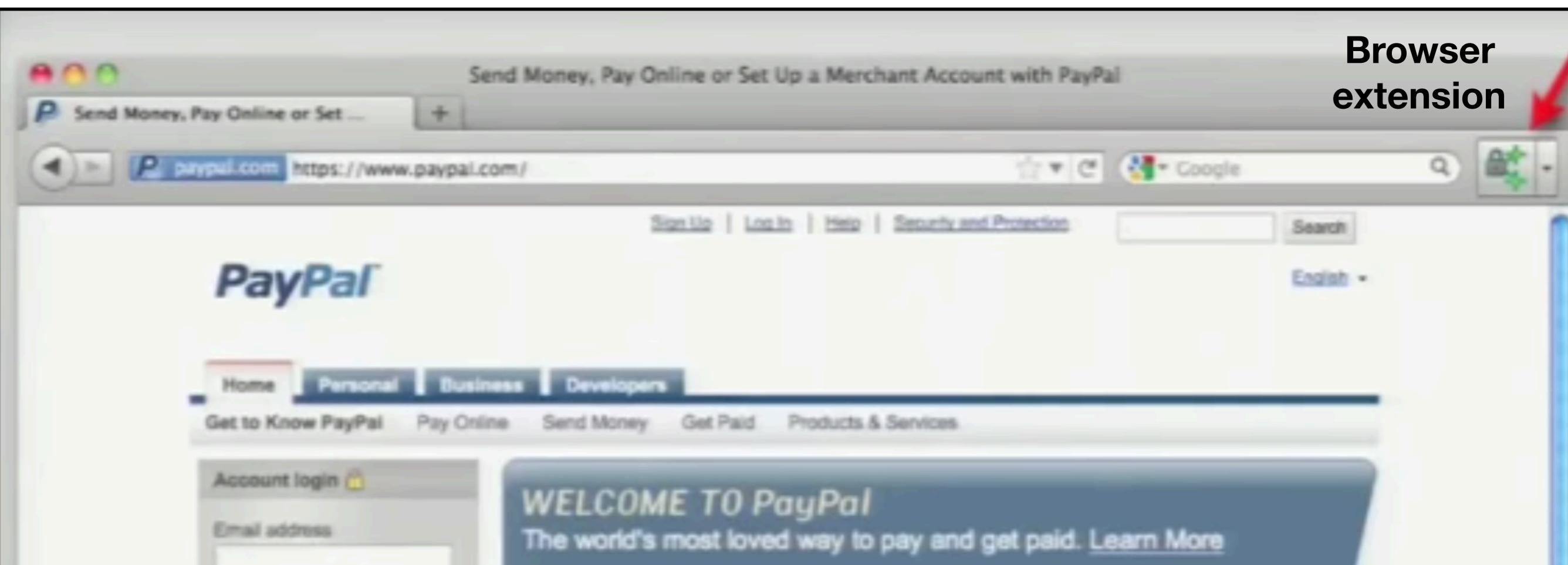
**Download all publicly served SSL certificates from websites.**

# CA backdoors

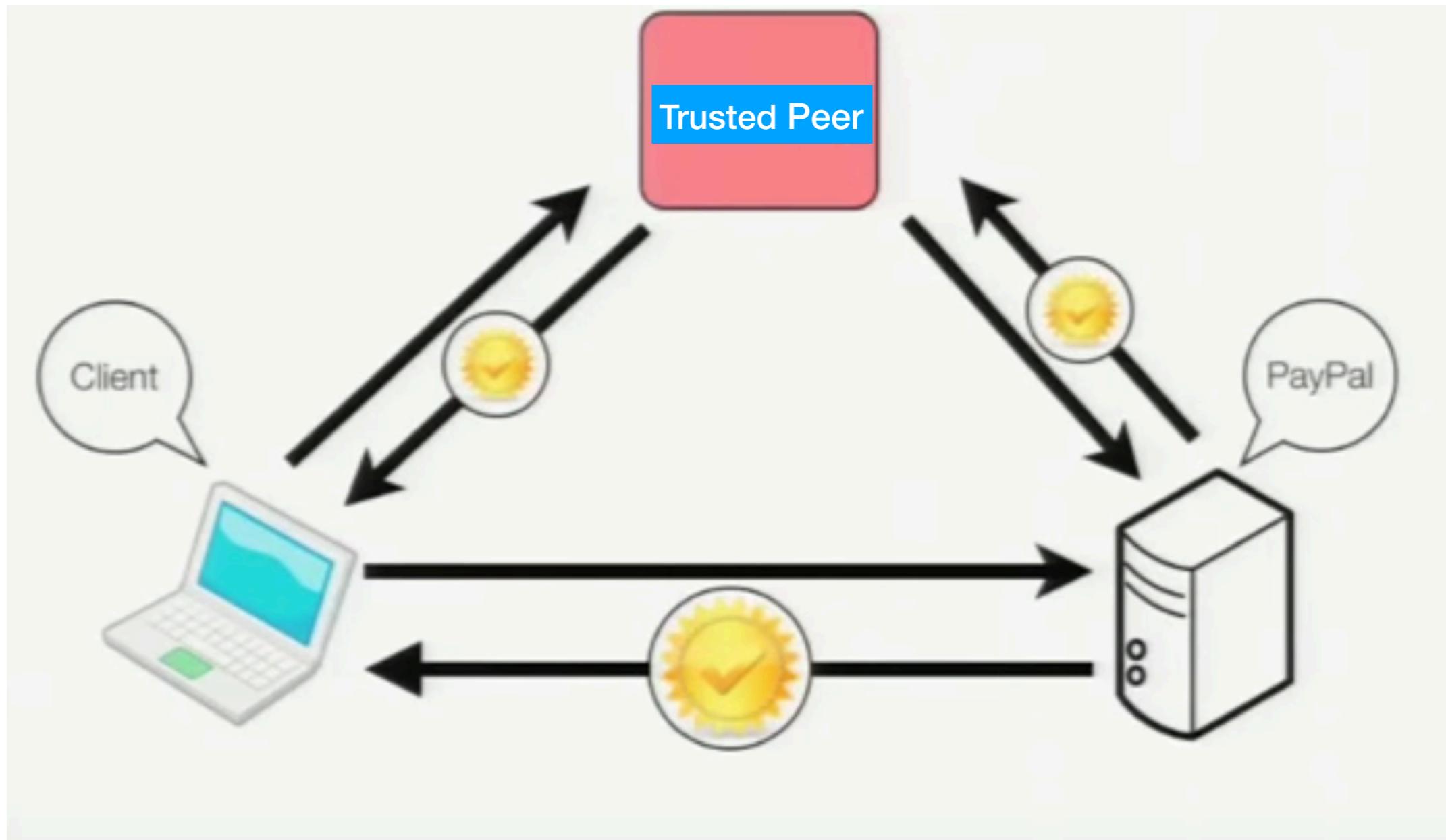
- 2009: 3 vulnerabilities due to CA mistakes
- 2010: evidence of govt compelling CAs
- 2012: Trustwave selling backdoor to private company.

# Moxie's Solution

From 2013



# Peer verifications



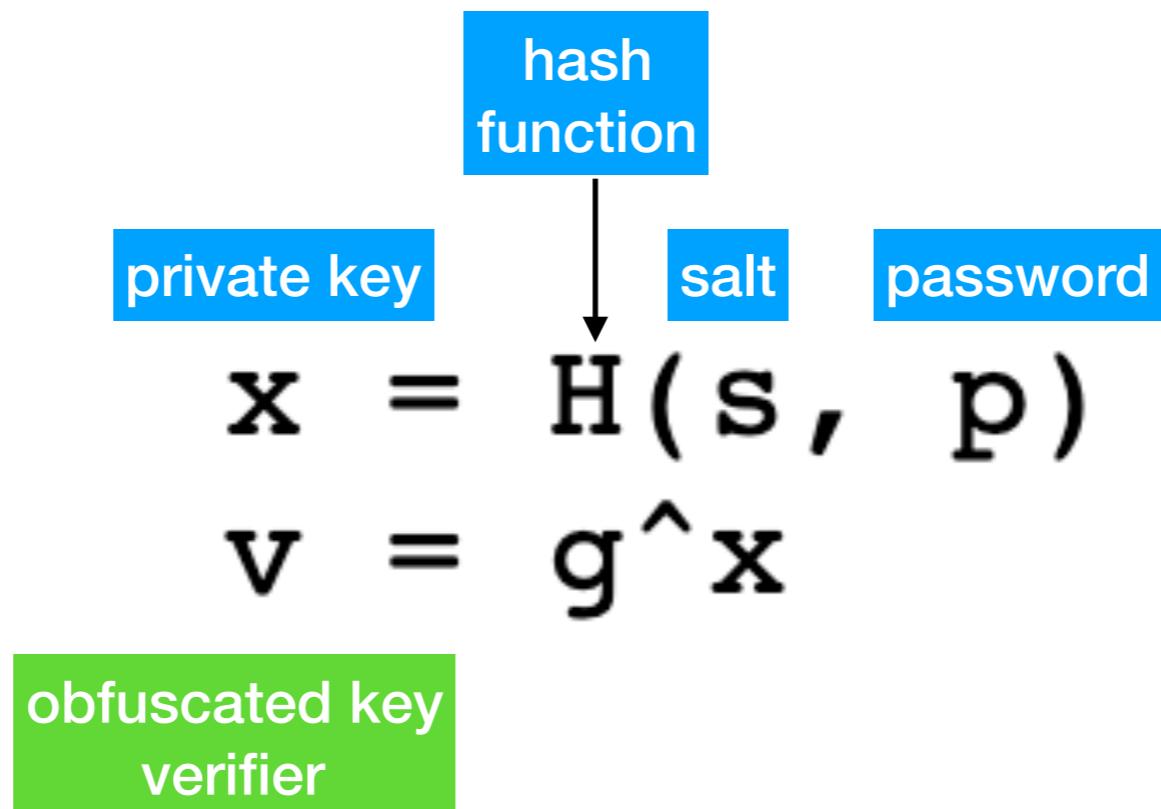
# Authenticated Key Exchange

- PAKE: password authenticated key exchange
- SRP: Secure Remote Password protocol
- OPAQUE
- For symmetric keys.
- Can use for login as special case.

# SRP

- Implemented in OpenSSL library
- Used by Apple for iCloud Key Vault.

# SRP signup



**Server stores {s, v} per username**

# SRP Login

User generates secret “a”

Host generates secret “b”

User -> Host:  $I, A = g^a$   
Host -> User:  $s, B = kv + g^b$

obfuscated  
secrets

public constant

verifier

# SRP Login math

Both:  $u = H(A, B)$

User:  $x = H(s, p)$

User:  $S = (B - kg^x) \hat{^} (a + ux)$

User:  $K = H(S)$

Host:  $S = (Av^u) \hat{^} b$

derived key

Host:  $K = H(S)$

# Problem with SRP

- No security proof
- Is based on a ring, not a group.
- RSA ok
- not ECC
- already in version 4 after some vulnerabilities found.

# OPAQUE

- Newer protocol, a replacement for SRP
- Can work in Curve25519 group.
- Hash to curve function.
- Elligator

# OPAQUE in JavaScript

 borisreitman / [opaque\\_dh\\_oprf\\_gist.js](#)

Last active 5 days ago

[Code](#)   [Revisions 3](#)   [Stars 1](#)

OPAQUE DH-OPRF in Javascript

## Protecting Passwords from Eavesdroppers



Boris Reitman

Jan 22 · 14 min read ★

*This article gets into technical security discussion of a new key-exchange protocol called OPAQUE.*

# OPAQUE protocol

RFC Draft

Crypto Forum Research Group  
Internet-Draft  
Intended status: Informational  
Expires: April 4, 2019

H. Krawczyk  
IBM Research  
October 01, 2018

The OPAQUE Asymmetric PAKE Protocol  
[draft-krawczyk-cfrg-opaque-00](https://datatracker.ietf.org/doc/draft-krawczyk-cfrg-opaque-00)

# OPAQUE Registration

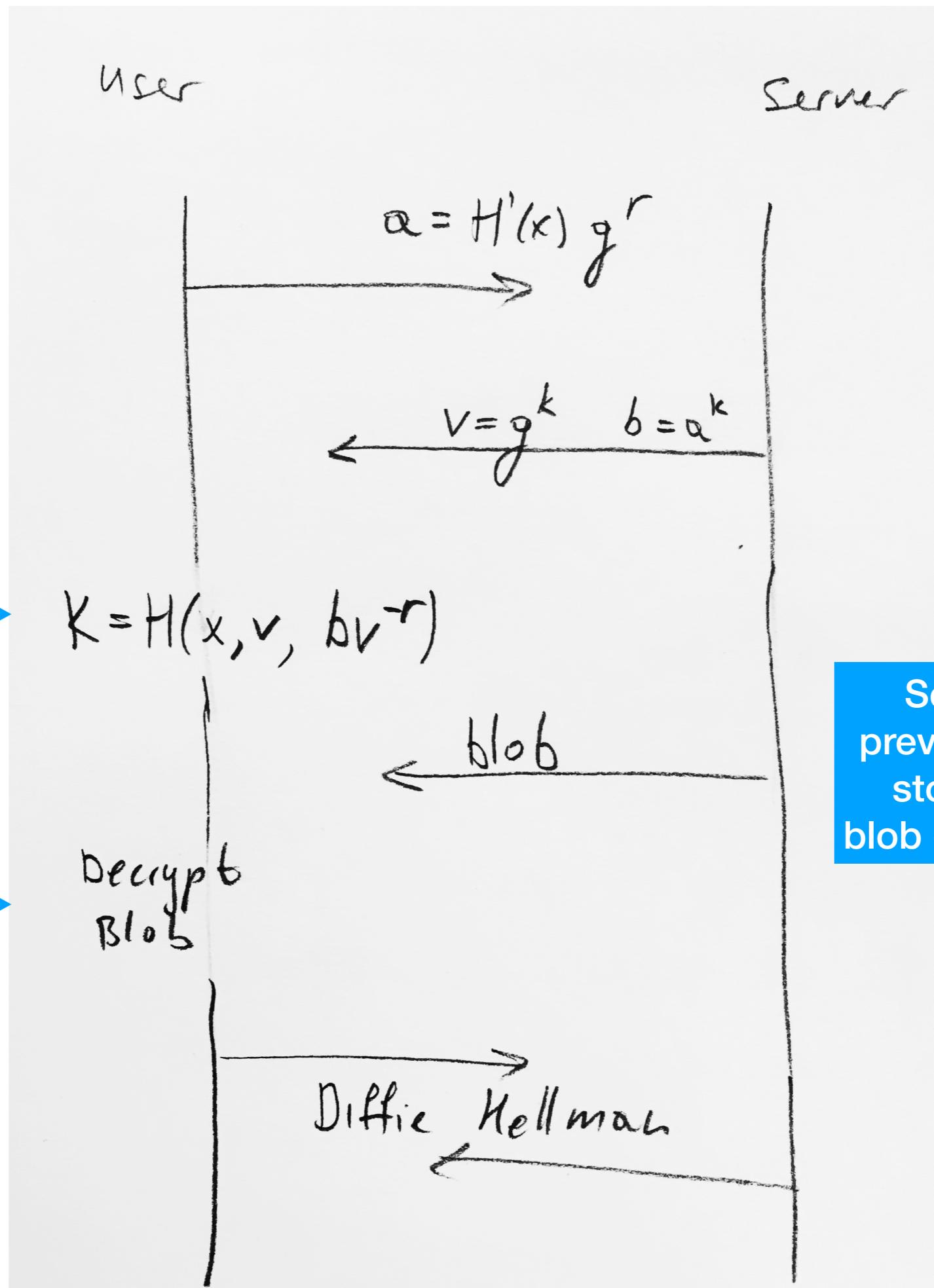
- User picks weak password “x”
- User generates salt “k”, stores it on the server.
- Generates symmetric key  $K = f(x, k)$
- Blob = Encrypt( $K$ ; { server pub key, user pub key })
- Stores the blob on the server.

# OPAQUE Login

- User generates a random value “r”
- Sends  $x^r$  to the server. (x is the password)
- Server responds with obfuscated “k” value.
- User constructs from this key K.
- Server also sends the encrypted blob.
- User decrypts it with K, and can authenticate the server.
- Also can authenticate himself to the server.

# OPAQUE Key Exchange

Stuff user stored  
about legitimate  
server



# Conclusion

- [borisreitman.com](http://borisreitman.com)
- [medium.com/@borisreitman](https://medium.com/@borisreitman)
- Other philosophical meetups