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1 Language

$e :=$	$n \mid C \mid D(e) \mid (e_1, e_2)$	Constructeurs de donnes
	x	Identificateur
	$\lambda x.e \mid \mathbf{recf}.x.e$	Fonctions (recursive ou non)
	$e_1 \ e_2$	Application
	$\mathbf{let} \ x = e_1 \ \mathbf{in} \ e_2$	Liaison
	$\mathbf{if} \ e \ \mathbf{then} \ e_1 \ \mathbf{else} \ e_2$	Expression conditionnelle
	$\mathbf{match} \ e \ \mathbf{with} \ p \rightarrow e_1 \mid x \rightarrow e_2$	Filtrage par motif
	$l : e$	Annotation d'un point d'injection

2 Operational Semantics

2.1 Values/Environments

$$\begin{aligned}
 v := & \quad n \mid b \mid C \mid D(v) \mid (v_1, v_2) && \text{Constructeurs de donnees} \\
 & < \lambda x.e, \Gamma > && \text{Fermeture} \\
 & < \mathbf{rec}f.x.e, \Gamma > && \text{Fermeture rcursive}
 \end{aligned}$$

$$\Gamma := (x_1, v_1) \oplus \dots \oplus (x_n, v_n) \quad \text{Environnement}$$

2.2 Inference Rules

$\frac{\text{OP-NUM}}{\Gamma \vdash n \mapsto n}$	$\frac{\text{OP-IDENT}}{v = \Gamma[x]} \quad \Gamma \vdash x \mapsto v$	$\frac{\text{OP-ABSTR}}{\Gamma \vdash \lambda x.e \mapsto < \lambda x.e, \Gamma >}$	$\frac{\text{OP-ABSTR-REC}}{\Gamma \vdash \mathbf{rec}f.x.e \mapsto < \mathbf{rec}f.x.e, \Gamma >}$
$\frac{\text{OP-APPLY}}{\Gamma \vdash e_1 \mapsto < \lambda x.e, \Gamma_1 > \quad (x, v_2) \oplus \Gamma_1 \vdash e \mapsto v} \quad \Gamma \vdash e_2 \mapsto v_2 \quad \Gamma \vdash e_1 e_2 \mapsto v$		$\frac{\text{OP-APPLY-REC}}{\Gamma \vdash e_1 \mapsto v_1 \quad v_1 = < \mathbf{rec}f.x.e, \Gamma_1 > \quad (f, v_1) \oplus (x, v_2) \oplus \Gamma_1 \vdash e \mapsto v} \quad \Gamma \vdash e_2 \mapsto v_2 \quad \Gamma \vdash e_1 e_2 \mapsto v$	
$\frac{\text{OP-IF-TRUE}}{\Gamma \vdash e \mapsto \mathbf{true}} \quad \Gamma \vdash e_1 \mapsto v_1 \quad \Gamma \vdash \mathbf{if} \ e \ \mathbf{then} \ e_1 \ \mathbf{else} \ e_2 \mapsto v_1$		$\frac{\text{OP-IF-FALSE}}{\Gamma \vdash e \mapsto \mathbf{false}} \quad \Gamma \vdash e_2 \mapsto v_2 \quad \Gamma \vdash \mathbf{if} \ e \ \mathbf{then} \ e_1 \ \mathbf{else} \ e_2 \mapsto v_2$	
$\frac{\text{OP-CONSTR-0}}{\Gamma \vdash C \mapsto C}$	$\frac{\text{OP-CONSTR-1}}{\Gamma \vdash e \mapsto v} \quad \Gamma \vdash D(e) \mapsto D(v)$	$\frac{\text{OP-COUPLE}}{\Gamma \vdash e_1 \mapsto v_1 \quad \Gamma \vdash e_2 \mapsto v_2} \quad \Gamma \vdash (e_1, e_2) \mapsto (v_1, v_2)$	
$\frac{\text{OP-MATCH}}{\Gamma \vdash e \mapsto v \quad v, p \vdash_p \Gamma_p \quad \Gamma_p \oplus \Gamma \vdash e_1 \mapsto v_1} \quad \Gamma \vdash \mathbf{match} \ e \ \mathbf{with} \ p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto v_1$		$\frac{\text{OP-MATCH-VAR}}{\Gamma \vdash e \mapsto v \quad v, p \vdash_p \perp \quad (x, v) \oplus \Gamma \vdash e_2 \mapsto v_2} \quad \Gamma \vdash \mathbf{match} \ e \ \mathbf{with} \ p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto v_2$	
$\frac{\text{OP-LETIN}}{\Gamma \vdash e_1 \mapsto v_1 \quad (x, v_1) \oplus \Gamma \vdash e_2 \mapsto v_2} \quad \Gamma \vdash \mathbf{let} \ x = e_1 \ \mathbf{in} \ e_2 \mapsto v_2$		$\frac{\text{OP-ANNOT}}{\Gamma \vdash e \mapsto v} \quad \Gamma \vdash l : e \mapsto v$	
$\frac{\text{OPM-CONSTR-0}}{C, C \vdash_p \{\}} \quad \frac{\text{OPM-CONSTR-1}}{D(v), D(x) \vdash_p \{(x, v)\}} \quad \frac{\text{OPM-COUPLE}}{(v_1, v_2), (x_1, x_2) \vdash_p \{(x_1, v_1); (x_2, v_2)\}}$	$\frac{\text{OPM-CONSTR-0-NOT}}{p \neq C} \quad \frac{\text{OPM-CONSTR-1-NOT}}{p \neq D'(-)} \quad \frac{\text{OPM-COUPLE-NOT}}{p \neq (-, -)} \quad \frac{}{C, p \vdash_p \perp} \quad \frac{}{D(v), p \vdash_p \perp} \quad \frac{}{(v_1, v_2), p \vdash_p \perp}$		

3 Operational Semantics with Injection

3.1 Inference Rules

OPINJ-NUM	OPINJ-IDENT $v = \Gamma[x]$	OPINJ-ABSTR	OPINJ-ABSTR-REC
$\frac{}{\Gamma \vdash_{l:v_l} n \mapsto n}$	$\frac{}{\Gamma \vdash_{l:v_l} x \mapsto v}$	$\frac{}{\Gamma \vdash_{l:v_l} \lambda x.e \mapsto \langle \lambda x.e, \Gamma \rangle}$	$\frac{}{\Gamma \vdash_{l:v_l} \mathbf{rec}f.x.e \mapsto \langle \mathbf{rec}f.x.e, \Gamma \rangle}$
OPINJ-APPLY	OPINJ-APPLY-REC		
$\frac{\Gamma \vdash_{l:v_l} e_1 \mapsto \langle \lambda x.e, \Gamma_1 \rangle \quad (x, v_2) \oplus \Gamma_1 \vdash_{l:v_l} e \mapsto v}{\Gamma \vdash_{l:v_l} e_1 e_2 \mapsto v}$	$\frac{\Gamma \vdash_{l:v_l} e_1 \mapsto v_1 \quad v_1 = \langle \mathbf{rec}f.x.e, \Gamma_1 \rangle \quad \Gamma \vdash_{l:v_l} e_2 \mapsto v_2 \quad (f, v_1) \oplus (x, v_2) \oplus \Gamma_1 \vdash_{l:v_l} e \mapsto v}{\Gamma \vdash_{l:v_l} e_1 e_2 \mapsto v}$		
OPINJ-IF-TRUE	OPINJ-IF-FALSE		
$\frac{\Gamma \vdash_{l:v_l} e \mapsto \mathbf{true} \quad \Gamma \vdash_{l:v_l} e_1 \mapsto v_1}{\Gamma \vdash_{l:v_l} \mathbf{if} e \mathbf{ then } e_1 \mathbf{ else } e_2 \mapsto v_1}$	$\frac{\Gamma \vdash_{l:v_l} e \mapsto \mathbf{false} \quad \Gamma \vdash_{l:v_l} e_2 \mapsto v_2}{\Gamma \vdash_{l:v_l} \mathbf{if} e \mathbf{ then } e_1 \mathbf{ else } e_2 \mapsto v_2}$		
OPINJ-MATCH	OPINJ-MATCH-VAR		
$\frac{\Gamma \vdash_{l:v_l} e \mapsto v \quad v, p \vdash_p \Gamma_p \quad \Gamma_p \oplus \Gamma \vdash_{l:v_l} e_1 \mapsto v_1}{\Gamma \vdash_{l:v_l} \mathbf{match} e \mathbf{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto v_1}$	$\frac{\Gamma \vdash_{l:v_l} e \mapsto v \quad v, p \vdash_p \perp \quad (x, v) \oplus \Gamma \vdash_{l:v_l} e_2 \mapsto v_2}{\Gamma \vdash_{l:v_l} \mathbf{match} e \mathbf{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto v_2}$		
OPINJ-CONSTR-0	OPINJ-CONSTR-1	OPINJ-COUPLE	
$\frac{}{\Gamma \vdash_{l:v_l} C \mapsto C}$	$\frac{}{\Gamma \vdash_{l:v_l} D(e) \mapsto D(v)}$	$\frac{\Gamma \vdash_{l:v_l} e_1 \mapsto v_1 \quad \Gamma \vdash_{l:v_l} e_2 \mapsto v_2}{\Gamma \vdash_{l:v_l} (e_1, e_2) \mapsto (v_1, v_2)}$	
OPINJ-LETIN			
$\frac{\Gamma \vdash_{l:v_l} e_1 \mapsto v_1 \quad (x, v_1) \oplus \Gamma \vdash_{l:v_l} e_2 \mapsto v_2}{\Gamma \vdash_{l:v_l} \mathbf{let} x = e_1 \mathbf{ in } e_2 \mapsto v_2}$			
OPINJ-ANNOT-SAME		OPINJ-ANNOT-OTHER	
$\frac{}{\Gamma \vdash_{l:v_l} l : e \mapsto v_l}$		$\frac{\Gamma \vdash_{l:v_l} e \mapsto v \quad l \neq l'}{\Gamma \vdash_{l:v_l} l' : e \mapsto v}$	

4 Injection Semantics

4.1 Inference Rules

INJ-NUM	INJ-LETIN
$\frac{}{t\Gamma^{oi} \vdash_{l:v_l} n \mapsto n}$	$\frac{t\Gamma^{oi} \vdash_{l:v_l} e_1 \mapsto v_1 \quad (x, \uparrow^{toi}(v_1)) \oplus t\Gamma^{oi} \vdash_{l:v_l} e_2 \mapsto v_2}{t\Gamma^{oi} \vdash_{l:v_l} \mathbf{let} x = e_1 \mathbf{ in } e_2 \mapsto v_2}$
INJ-APPLY	INJ-APPLY-REC
$\frac{t\Gamma^{oi} \vdash_{l:v_l} e_1 \mapsto \langle \lambda x.e, \Gamma_1 \rangle \quad t\Gamma^{oi} \vdash_{l:v_l} e_2 \mapsto v_2 \quad (x, \uparrow^{toi}(v_2)) \oplus \uparrow^{toi}(\Gamma_1) \vdash_{l:v_l} e \mapsto v}{t\Gamma^{oi} \vdash_{l:v_l} e_1 e_2 \mapsto v}$	$\frac{t\Gamma^{oi} \vdash_{l:v_l} e_1 \mapsto v_1 \quad v_1 = \langle \mathbf{rec}f.x.e, \Gamma_1 \rangle \quad t\Gamma^{oi} \vdash_{l:v_l} e_2 \mapsto v_2 \quad (f, \uparrow^{toi}(v_1)) \oplus (x, \uparrow^{toi}(v_2)) \oplus \uparrow^{toi}(\Gamma_1) \vdash_{l:v_l} e \mapsto v}{t\Gamma^{oi} \vdash_{l:v_l} e_1 e_2 \mapsto v}$

$$\begin{array}{c}
\text{INJ-IF-TRUE} \quad \frac{t\Gamma^{oi} \vdash_{l:v_l} e \mapsto \text{true} \quad t\Gamma^{oi} \vdash_{l:v_l} e_1 \mapsto v_1}{t\Gamma^{oi} \vdash_{l:v_l} \text{if } e \text{ then } e_1 \text{ else } e_2 \mapsto v_1} \quad \text{INJ-IF-FALSE} \quad \frac{t\Gamma^{oi} \vdash_{l:v_l} e \mapsto \text{false} \quad t\Gamma^{oi} \vdash_{l:v_l} e_2 \mapsto v_2}{t\Gamma^{oi} \vdash_{l:v_l} \text{if } e \text{ then } e_1 \text{ else } e_2 \mapsto v_2} \\
\\
\text{INJ-MATCH} \quad \frac{t\Gamma^{oi} \vdash_{l:v_l} e \mapsto v \quad v, p \vdash_p \Gamma_p \quad \uparrow^{toi}(\Gamma_p) \oplus t\Gamma^{oi} \vdash_{l:v_l} e_1 \mapsto v_1}{t\Gamma^{oi} \vdash_{l:v_l} \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto v_1} \quad \text{INJ-MATCH-VAR} \quad \frac{t\Gamma^{oi} \vdash_{l:v_l} e \mapsto v \quad v, p \vdash_p \perp \quad (x, \uparrow^{toi}(v)) \oplus t\Gamma^{oi} \vdash_{l:v_l} e_2 \mapsto v_2}{t\Gamma^{oi} \vdash_{l:v_l} \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto v_2} \\
\\
\text{INJ-CONSTR-0} \quad \frac{}{t\Gamma^{oi} \vdash_{l:v_l} C \mapsto C} \quad \text{INJ-CONSTR-1} \quad \frac{t\Gamma^{oi} \vdash_{l:v_l} e \mapsto v}{t\Gamma^{oi} \vdash_{l:v_l} D(e) \mapsto D(v)} \quad \text{INJ-COUPLE} \quad \frac{t\Gamma^{oi} \vdash_{l:v_l} e_1 \mapsto v_1 \quad t\Gamma^{oi} \vdash_{l:v_l} e_2 \mapsto v_2}{t\Gamma^{oi} \vdash_{l:v_l} (e_1, e_2) \mapsto (v_1, v_2)} \\
\\
\text{INJ-IDENT} \quad \frac{v = \downarrow^{l:v_l}(t\Gamma^{oi}[x])}{t\Gamma^{oi} \vdash_{l:v_l} x \mapsto v} \quad \text{INJ-ANNOT-SAME} \quad \frac{}{t\Gamma^{oi} \vdash_{l:v_l} l : e \mapsto v_l} \quad \text{INJ-ANNOT-OTHER} \quad \frac{t\Gamma^{oi} \vdash_{l:v_l} e \mapsto v \quad l \neq l'}{t\Gamma^{oi} \vdash_{l:v_l} l' : e \mapsto v} \\
\\
\text{INJ-ABSTR} \quad \frac{}{t\Gamma^{oi} \vdash_{l:v_l} \lambda x.e \mapsto < \lambda x.e, \downarrow^{l:v_l}(t\Gamma^{oi}) >} \quad \text{INJ-ABSTR-REC} \quad \frac{}{t\Gamma^{oi} \vdash_{l:v_l} \text{recf}.x.e \mapsto < \text{recf}.x.e, \downarrow^{l:v_l}(t\Gamma^{oi}) >}
\end{array}$$

5 Over-Instrumented Semantics

5.1 Values/Environments

$tu^{oi} := [td^{oi} \mid u^{oi}]$ Valeur avec annotation de t -dependances

$u^{oi} := [d^{oi} \mid v^{oi}]$ Valeur avec annotation de v -dependances

$v^{oi} := n \mid b \mid C \mid D(u^{oi}) \mid (u_1^{oi}, u_2^{oi})$ Constructeurs de donnees
 $< \lambda x.e, \Gamma^{oi} >$ Fermeture
 $< \text{recf}.x.e, \Gamma^{oi} >$ Fermeture rcursive

$td^{oi} := (l_1, tf_1); \dots (l_n, tf_n)$ t -dependances
 $tf : v \rightarrow \text{bool}$ t -fonction d'impact

$d^{oi} := (l_1, f_1); \dots (l_n, f_n)$ v -dependances
 $f : v \rightarrow v$ v -fonction d'impact

$t\Gamma^{oi} := (x_1, tu_1^{oi}); \dots (x_n, tu_n^{oi})$ t -environnement sur-instrument

$\Gamma^{oi} := (x_1, u_1^{oi}); \dots (x_n, u_n^{oi})$ v -environnement sur-instrument

5.2 Inference Rules

OI-NUM

$$\frac{}{t\Gamma^{oi} \vdash^{oi} n \mapsto [\emptyset \mid [\emptyset \mid n]]}$$

OI-ABSTR

$$\frac{}{t\Gamma^{oi} \vdash^{oi} \lambda x.e \mapsto [\emptyset \mid [\emptyset \mid <\lambda x.e, \uparrow_{toi}^{oi}(t\Gamma^{oi})>]]}$$

OI-IDENT

$$\frac{tu^{oi} = t\Gamma^{oi}[x]}{t\Gamma^{oi} \vdash^{oi} x \mapsto tu^{oi}}$$

OI-ABSTR-REC

$$\frac{}{t\Gamma^{oi} \vdash^{oi} \mathbf{rec}f.x.e \mapsto [\emptyset \mid [\emptyset \mid <\mathbf{rec}f.x.e, \uparrow_{toi}^{oi}(t\Gamma^{oi})>]]}$$

OI-APPLY

$$\frac{\begin{array}{l} t\Gamma^{oi} \vdash^{oi} e_1 \mapsto [td_1^{oi} \mid [d_1^{oi} \mid <\lambda x.e, \Gamma_1^{oi}>]] \\ t\Gamma^{oi} \vdash^{oi} e_2 \mapsto tu_2^{oi} \quad tu_2^{oi} = [td_2^{oi} \mid u_2^{oi}] \\ (x, tu_2^{oi}) \oplus \uparrow_{oi}^{toi}(\Gamma_1^{oi}) \vdash^{oi} e \mapsto [td^{oi} \mid [d^{oi} \mid v^{oi}]] \\ \mathit{deps_spec_apply}(tu_2^{oi}, d_1^{oi}) = (td'^{oi}, d'^{oi}) \end{array}}{t\Gamma^{oi} \vdash^{oi} e_1 e_2 \mapsto [td_1^{oi} \cup td_2^{oi} \cup td^{oi} \cup td'^{oi} \mid [d'^{oi} \cup d^{oi} \mid v^{oi}]]}$$

OI-REC-APPLY

$$\frac{\begin{array}{l} t\Gamma^{oi} \vdash^{oi} e_1 \mapsto tu_1^{oi} \quad tu_1^{oi} = [td_1^{oi} \mid [d_1^{oi} \mid <\mathbf{rec}f.x.e, \Gamma_1^{oi}>]] \\ t\Gamma^{oi} \vdash^{oi} e_2 \mapsto tu_2^{oi} \quad tu_2^{oi} = [td_2^{oi} \mid u_2^{oi}] \\ (f, tu_1^{oi}) \oplus (x, tu_2^{oi}) \oplus \uparrow_{oi}^{toi}(\Gamma_1^{oi}) \vdash^{oi} e \mapsto [td^{oi} \mid [d^{oi} \mid v^{oi}]] \\ \mathit{deps_spec_apply}(tu_2^{oi}, d_1^{oi}) = (td'^{oi}, d'^{oi}) \end{array}}{t\Gamma^{oi} \vdash^{oi} e_1 e_2 \mapsto [td_1^{oi} \cup td_2^{oi} \cup td^{oi} \cup td'^{oi} \mid [d'^{oi} \cup d^{oi} \mid v^{oi}]]}$$

OI-LETIN

$$\frac{\begin{array}{l} t\Gamma^{oi} \vdash^{oi} e_1 \mapsto tu_1^{oi} \quad tu_1^{oi} = [td_1^{oi} \mid u_1^{oi}] \\ (x, tu_1^{oi}) \oplus t\Gamma^{oi} \vdash^{oi} e_2 \mapsto [td_2^{oi} \mid u_2^{oi}] \end{array}}{t\Gamma^{oi} \vdash^{oi} \mathbf{let} x = e_1 \mathbf{in} e_2 \mapsto [td_1^{oi} \cup td_2^{oi} \mid u_2^{oi}]}$$

OI-IF-TRUE

$$\frac{\begin{array}{l} t\Gamma^{oi} \vdash^{oi} e \mapsto [td^{oi} \mid [d^{oi} \mid \mathbf{true}]] \\ t\Gamma^{oi} \vdash^{oi} e_1 \mapsto [td_1^{oi} \mid [d_1^{oi} \mid v_1^{oi}]] \\ \mathit{deps_spec_if}(t\Gamma^{oi}, e_1, e_2, d^{oi}) = (td'^{oi}, d'^{oi}) \end{array}}{t\Gamma^{oi} \vdash^{oi} \mathbf{if} e \mathbf{then} e_1 \mathbf{else} e_2 \mapsto [td'^{oi} \cup td^{oi} \cup td_1^{oi} \mid [d'^{oi} \cup d_1^{oi} \mid v_1^{oi}]]}$$

OI-IF-FALSE

$$\frac{\begin{array}{l} t\Gamma^{oi} \vdash^{oi} e \mapsto [td^{oi} \mid [d^{oi} \mid \mathbf{false}]] \\ t\Gamma^{oi} \vdash^{oi} e_2 \mapsto [td_2^{oi} \mid [d_2^{oi} \mid v_2^{oi}]] \\ \mathit{deps_spec_if}(t\Gamma^{oi}, e_1, e_2, d^{oi}) = (td'^{oi}, d'^{oi}) \end{array}}{t\Gamma^{oi} \vdash^{oi} \mathbf{if} e \mathbf{then} e_1 \mathbf{else} e_2 \mapsto [td'^{oi} \cup td^{oi} \cup td_2^{oi} \mid [d'^{oi} \cup d_2^{oi} \mid v_2^{oi}]]}$$

OI-MATCH

$$\frac{\begin{array}{l} t\Gamma^{oi} \vdash^{oi} e \mapsto tu^{oi} \quad tu^{oi} = [td^{oi} \mid [d^{oi} \mid v^{oi}]] \\ tu_p^{oi}, p \vdash_p^{oi} t\Gamma_p^{oi} \\ t\Gamma_p^{oi} \oplus t\Gamma^{oi} \vdash^{oi} e_1 \mapsto [td_1^{oi} \mid [d_1^{oi} \mid v_1^{oi}]] \\ \mathit{deps_spec_match}(t\Gamma^{oi}, p, x, e_1, e_2, d^{oi}) = (td'^{oi}, d'^{oi}) \end{array}}{t\Gamma^{oi} \vdash^{oi} \mathbf{match} e \mathbf{with} p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [td'^{oi} \cup td^{oi} \cup td_1^{oi} \mid [d'^{oi} \cup d_1^{oi} \mid v_1^{oi}]]}$$

OI-MATCH-VAR

$$\frac{\begin{array}{c} t\Gamma^{oi} \vdash^{oi} e \mapsto tu^{oi} \quad tu^{oi} = [td^{oi} \mid [d^{oi} \mid v^{oi}]] \\ tu^{oi}, p \vdash_p^{oi} \perp \\ (x, tu^{oi}) \oplus t\Gamma^{oi} \vdash^{oi} e_2 \mapsto [td_2^{oi} \mid [d_2^{oi} \mid v_2^{oi}]] \\ \text{deps_spec_match}(t\Gamma^{oi}, p, x, e_1, e_2, d^{oi}) = (td'^{oi}, d'^{oi}) \end{array}}{t\Gamma^{oi} \vdash^{oi} \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [td'^{oi} \cup td^{oi} \cup td_2^{oi} \mid [d'^{oi} \cup d_2^{oi} \mid v_2^{oi}]]}$$

OI-CONSTR-0

$$\frac{}{t\Gamma^{oi} \vdash^{oi} C \mapsto [\emptyset \mid [\emptyset \mid C]]} \quad \text{OI-CONSTR-1} \quad \frac{}{t\Gamma^{oi} \vdash^{oi} e \mapsto [td^{oi} \mid u^{oi}]} \quad \frac{}{t\Gamma^{oi} \vdash^{oi} D(e) \mapsto [td^{oi} \mid [\emptyset \mid D(u^{oi})]]}$$

OI-COUPLE

$$\frac{t\Gamma^{oi} \vdash^{oi} e_1 \mapsto [td_1^{oi} \mid u_1^{oi}] \quad t\Gamma^{oi} \vdash^{oi} e_2 \mapsto [td_2^{oi} \mid u_2^{oi}]}{t\Gamma^{oi} \vdash^{oi} (e_1, e_2) \mapsto [td_1^{oi} \cup td_2^{oi} \mid [\emptyset \mid (u_1^{oi}, u_2^{oi})]]}$$

OI-ANNOT

$$\frac{t\Gamma^{oi} \vdash^{oi} e \mapsto [td^{oi} \mid [d^{oi} \mid v^{oi}]]}{t\Gamma^{oi} \vdash^{oi} l : e \mapsto [td^{oi} \mid [(l, \text{fun } x \Rightarrow x); d^{oi} \mid v^{oi}]]}$$

OIM-CONSTR-0

$$\frac{}{[td^{oi} \mid [d^{oi} \mid C]], C \vdash_p^{oi} \{ \}}$$

OIM-CONSTR-1

$$\frac{}{[td^{oi} \mid [d^{oi} \mid D(u^{oi})]], D(x) \vdash_p^{oi} \{(x, [\emptyset \mid u^{oi}])\}}$$

OIM-COUPLE

$$\frac{}{([td_1^{oi} \mid [d_1^{oi} \mid u_1^{oi}]], [td_2^{oi} \mid [d_2^{oi} \mid u_2^{oi}]]), (x_1, x_2) \vdash_p^{oi} \{(x_1, [\emptyset \mid u_1^{oi}]); (x_2, [\emptyset \mid u_2^{oi}])\}}$$

OIM-CONSTR-0-NOT

$$\frac{p \neq C}{[td^{oi} \mid [d^{oi} \mid C]], p \vdash_p^{oi} \perp}$$

OIM-CONSTR-1-NOT

$$\frac{p \neq D'(-)}{[td^{oi} \mid [d^{oi} \mid D(u^{oi})]], p \vdash_p^{oi} \perp}$$

OIM-COUPLE-NOT

$$\frac{p \neq (-, -)}{([td_1^{oi} \mid [d_1^{oi} \mid u_1^{oi}]], [td_2^{oi} \mid [d_2^{oi} \mid u_2^{oi}]]), p \vdash_p^{oi} \perp}$$

6 Instrumented Semantics

6.1 Values/Environments

$tu^i := [td^i \mid u^i]$ Valeur avec annotation de t -dpendance

$u^i := [d^i \mid v^i]$ Valeur avec annotation de v -dpendance

$v^i := n \mid b \mid C \mid D(u^i) \mid (u_1^i, u_2^i)$ Constructeurs de donnes
 $< \lambda x. e, \Gamma^i >$ Fermeture
 $< \text{recf}. x. e, \Gamma^i >$ Fermeture rcursive

$td^i := \{l_1; \dots; l_n\}$ t -dpendances
 $d^i := \{l_1; \dots; l_m\}$ v -dpendances

$t\Gamma^i := (x_1, tu_1^i); \dots; (x_n, tu_n^i)$ t -environnement instrument

$\Gamma^i := (x_1, u_1^i); \dots; (x_n, u_n^i)$ v -environnement instrument

6.2 Inference Rules

$$\begin{array}{c}
\text{I-IDENT} \quad \frac{t\Gamma^i \vdash^i e_1 \mapsto tu_1^i \quad tu_1^i = [td_1^i \mid u_1^i]}{(x, tu_1^i) \oplus t\Gamma^i \vdash^i e_2 \mapsto [td_2^i \mid u_2^i]} \quad \text{I-LETIN} \\
\text{I-NUM} \quad \frac{}{t\Gamma^i \vdash^i n \mapsto [\emptyset \mid [\emptyset \mid n]]} \quad \frac{t\Gamma^i \vdash^i x \mapsto tu^i}{t\Gamma^i \vdash^i let \ x = e_1 \ in \ e_2 \mapsto [td_1^i \cup td_2^i \mid u_2^i]} \\
\text{I-ABSTR} \quad \frac{}{t\Gamma^i \vdash^i \lambda x. e \mapsto [\emptyset \mid [\emptyset \mid < \lambda x. e, \uparrow_{ti}^i(t\Gamma^i) >]]} \\
\text{I-ABSTR-REC} \quad \frac{}{t\Gamma^i \vdash^i \mathbf{recf}. x. e \mapsto [\emptyset \mid [\emptyset \mid < \mathbf{recf}. x. e, \uparrow_{ti}^i(t\Gamma^i) >]]} \\
\text{I-APPLY} \quad \frac{t\Gamma^i \vdash^i e_1 \mapsto [td_1^i \mid [d_1^i \mid < \lambda x. e, \Gamma_1^i >]]}{t\Gamma^i \vdash^i e_2 \mapsto tu_2^i \quad tu_2^i = [td_2^i \mid u_2^i]} \\
\frac{(x, tu_2^i) \oplus \uparrow_i^{ti}(\Gamma_1^i) \vdash^i e \mapsto [td^i \mid [d^i \mid v^i]]}{t\Gamma^i \vdash^i e_1 \ e_2 \mapsto [td_1^i \cup td_2^i \cup td^i \cup d_1^i \mid [d_1^i \cup d^i \mid v^i]]} \\
\text{I-APPLY-REC} \quad \frac{t\Gamma^i \vdash^i e_1 \mapsto tu_1^i \quad tu_1^i = [td_1^i \mid [d_1^i \mid < \mathbf{recf}. x. e, \Gamma_1^i >]]}{t\Gamma^i \vdash^i e_2 \mapsto tu_2^i \quad tu_2^i = [td_2^i \mid u_2^i]} \\
\frac{(f, tu_1^i) \oplus (x, tu_2^i) \oplus \uparrow_i^{ti}(\Gamma_1^i) \vdash^i e \mapsto [td^i \mid [d^i \mid v^i]]}{t\Gamma^i \vdash^i e_1 \ e_2 \mapsto [td_1^i \cup td_2^i \cup td^i \cup d_1^i \mid [d_1^i \cup d^i \mid v^i]]} \\
\text{I-IF-TRUE} \quad \frac{t\Gamma^i \vdash^i e \mapsto [td^i \mid [d^i \mid true]] \quad t\Gamma^i \vdash^i e_1 \mapsto [td_1^i \mid [d_1^i \mid v_1^i]]}{t\Gamma^i \vdash^i if \ e \ then \ e_1 \ else \ e_2 \mapsto [d^i \cup td^i \cup td_1^i \mid [d^i \cup d_1^i \mid v_1^i]]} \\
\text{I-IF-FALSE} \quad \frac{t\Gamma^i \vdash^i e \mapsto [td^i \mid [d^i \mid false]] \quad t\Gamma^i \vdash^i e_2 \mapsto [td_2^i \mid [d_2^i \mid v_2^i]]}{t\Gamma^i \vdash^i if \ e \ then \ e_1 \ else \ e_2 \mapsto [d^i \cup td^i \cup td_2^i \mid [d^i \cup d_2^i \mid v_2^i]]} \\
\text{I-MATCH} \quad \frac{t\Gamma^i \vdash^i e \mapsto tu^i \quad tu^i = [td^i \mid [d^i \mid v^i]] \quad tu^i, p \vdash_p^i t\Gamma_p^i}{t\Gamma_p^i \oplus t\Gamma^i \vdash^i e_1 \mapsto [td_1^i \mid [d_1^i \mid v_1^i]]} \\
\frac{}{t\Gamma^i \vdash^i match \ e \ with \ p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [d^i \cup td^i \cup td_1^i \mid [d^i \cup d_1^i \mid v_1^i]]} \\
\text{I-MATCH-VAR} \quad \frac{t\Gamma^i \vdash^i e \mapsto tu^i \quad tu^i = [td^i \mid [d^i \mid v^i]] \quad tu^i, p \vdash_p^i \perp}{(x, tu^i) \oplus t\Gamma^i \vdash^i e_2 \mapsto [td_2^i \mid [d_2^i \mid v_2^i]]} \\
\frac{}{t\Gamma^i \vdash^i match \ e \ with \ p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [d^i \cup td^i \cup td_2^i \mid [d^i \cup d_2^i \mid v_2^i]]}
\end{array}$$

$$\begin{array}{c}
\text{I-CONSTR-0} \quad \frac{}{t\Gamma^i \vdash^i C \mapsto [\emptyset \mid [\emptyset \mid C]]} \quad \text{I-CONSTR-1} \quad \frac{t\Gamma^i \vdash^i e \mapsto [td^i \mid u^i]}{t\Gamma^i \vdash^i D(e) \mapsto [td^i \mid [\emptyset \mid D(u^i)]]} \\
\text{I-COUPLE} \quad \frac{t\Gamma^i \vdash^i e_1 \mapsto [td_1^i \mid u_1^i] \quad t\Gamma^i \vdash^i e_2 \mapsto [td_2^i \mid u_2^i]}{t\Gamma^i \vdash^i (e_1, e_2) \mapsto [td_1^i \cup td_2^i \mid [\emptyset \mid (u_1^i, u_2^i)]]} \quad \text{I-ANNO} \quad \frac{t\Gamma^i \vdash^i e \mapsto [td^i \mid [d^i \mid v^i]]}{t\Gamma^i \vdash^i l : e \mapsto [td^i \mid [l; d^i \mid v^i]]} \\
\text{IM-CONSTR-0} \quad \frac{}{[td^i \mid [d^i \mid C]], C \vdash_p^i \{ \}} \quad \text{IM-CONSTR-1} \quad \frac{}{[td^i \mid [d^i \mid D(u^i)]], D(x) \vdash_p^i \{ (x, [\emptyset \mid u^i]) \}} \\
\text{IM-COUPLE} \quad \frac{}{[td^i \mid [d^i \mid (u_1^i, u_2^i)]], (x_1, x_2) \vdash_p^i \{ (x_1, [\emptyset \mid u_1^i]); (x_2, [\emptyset \mid u_2^i]) \}} \\
\text{IM-CONSTR-0-NOT} \quad \frac{p \neq C}{[td^i \mid [d^i \mid C]], p \vdash_p^i \perp} \quad \text{IM-CONSTR-1-NOT} \quad \frac{p \neq D(-)}{[td^i \mid [d^i \mid D(u^i)]], p \vdash_p^i \perp} \\
\text{IM-COUPLE-NOT} \quad \frac{p \neq (-, -)}{[td^i \mid [d^i \mid (u_1^i, u_2^i)]], p \vdash_p^i \perp}
\end{array}$$

7 Multiple Instrumented Semantics

7.1 Inference Rules

$$\text{I-MULTIPLE} \quad \frac{v^{im} = \{tu^i \mid \exists t\Gamma^i \in t\Gamma^{im}. t\Gamma^i \vdash^i e \mapsto tu^i\}}{t\Gamma^{im} \vdash^{im} e \mapsto v^{im}}$$

8 Collecting Semantics

8.1 Values/Environments

$$tu^c := [td^c \mid d^c \mid vs^i] \quad \text{Valeur collectrice}$$

$$\begin{array}{ll}
td^c := \{l_1; \dots; l_n\} & t\text{-dpendances} \\
d^c := \{l_1; \dots; l_m\} & v\text{-dpendances}
\end{array}$$

$$t\Gamma^c := (x_1, tu_1^c); \dots; (x_n, tu_n^c) \quad t\text{-environnement collecteur}$$

8.2 Inference Rules

C-NUM	$\frac{tu^c = t\Gamma^c[x]}{t\Gamma^c \vdash^c n \mapsto [\emptyset \mid \emptyset \mid \{n\}]}$	C-IDENT	$\frac{x \notin support(t\Gamma^c)}{t\Gamma^c \vdash^c x \mapsto tu^c}$	C-IDENT-EMPTY	$\frac{}{t\Gamma^c \vdash^c x \mapsto [\emptyset \mid \emptyset \mid \emptyset]}$
C-ABSTR	$\frac{vs^i = \{< \lambda x.e, \Gamma^i > \mid \Gamma^i \in \uparrow_c^{im}(t\Gamma^c)\}}{t\Gamma^c \vdash^c \lambda x.e \mapsto [\emptyset \mid \emptyset \mid vs^i]}$	C-ABSTR-REC	$\frac{vs^i = \{< \mathbf{rec}f.x.e, \Gamma^i > \mid \Gamma^i \in \uparrow_c^{im}(t\Gamma^c)\}}{t\Gamma^c \vdash^c \mathbf{rec}f.x.e \mapsto [\emptyset \mid \emptyset \mid vs^i]}$		
C-APPLY	$\frac{\begin{array}{l} t\Gamma^c \vdash^c e_1 \mapsto [td_1^c \mid d_1^c \mid vs_1^i] \quad t\Gamma^c \vdash^c e_2 \mapsto [td_2^c \mid d_2^c \mid vs_2^i] \\ v^{im} = \text{multiple_instrumented_application}(td_1^c, d_1^c, vs_1^i, td_2^c, d_2^c, vs_2^i) \\ (\forall l. l \in td^c \Leftrightarrow (\exists (td^i, d^i, v^i). l \in td^i \wedge [td^i \mid [d^i \mid v^i]] \in v^{im})) \\ (\forall l. l \in d^c \Leftrightarrow (\exists (td^i, d^i, v^i). l \in d^i \wedge [td^i \mid [d^i \mid v^i]] \in v^{im})) \\ vs^i = \{v^i \mid \exists (td^i, d^i). [td^i \mid [d^i \mid v^i]] \in v^{im}\} \end{array}}{t\Gamma^c \vdash^c e_1 e_2 \mapsto [td^c \mid d^c \mid vs^i]}$				
C-IF-TRUE	$\frac{\begin{array}{l} t\Gamma^c \vdash^c e \mapsto [td^c \mid d^c \mid vs^i] \quad true \in vs^i \quad false \notin vs^i \\ t\Gamma^c \vdash^c e_1 \mapsto [td_1^c \mid d_1^c \mid vs_1^i] \quad tu'^c = [d^c \cup td^c \cup td_1^c \mid d^c \cup d_1^c \mid vs_1^i] \end{array}}{t\Gamma^c \vdash^c \text{if } e \text{ then } e_1 \text{ else } e_2 \mapsto tu'^c}$				
C-IF-FALSE	$\frac{\begin{array}{l} t\Gamma^c \vdash^c e \mapsto [td^c \mid d^c \mid vs^i] \quad false \in vs^i \quad true \notin vs^i \\ t\Gamma^c \vdash^c e_2 \mapsto [td_2^c \mid d_2^c \mid vs_2^i] \quad tu'^c = [d^c \cup td^c \cup td_2^c \mid d^c \cup d_2^c \mid vs_2^i] \end{array}}{t\Gamma^c \vdash^c \text{if } e \text{ then } e_1 \text{ else } e_2 \mapsto tu'^c}$				
C-IF-UNKNOWN	$\frac{\begin{array}{l} t\Gamma^c \vdash^c e \mapsto [td^c \mid d^c \mid vs^i] \quad true \in vs^i \quad false \in vs^i \\ t\Gamma^c \vdash^c e_1 \mapsto [td_1^c \mid d_1^c \mid vs_1^i] \quad t\Gamma^c \vdash^c e_2 \mapsto [td_2^c \mid d_2^c \mid vs_2^i] \\ tu'^c = [d^c \cup td^c \cup td_1^c \cup td_2^c \mid d^c \cup d_1^c \cup d_2^c \mid vs_1^i \cup vs_2^i] \end{array}}{t\Gamma^c \vdash^c \text{if } e \text{ then } e_1 \text{ else } e_2 \mapsto tu'^c}$				
C-IF-EMPTY	$\frac{t\Gamma^c \vdash^c e \mapsto [td^c \mid d^c \mid vs^i] \quad true \notin vs^i \quad false \notin vs^i}{t\Gamma^c \vdash^c \text{if } e \text{ then } e_1 \text{ else } e_2 \mapsto [\emptyset \mid \emptyset \mid \emptyset]}$				
C-MATCH	$\frac{\begin{array}{l} t\Gamma^c \vdash^c e \mapsto tu^c \quad tu^c = [td^c \mid d^c \mid vs^i] \quad vs^i \cap vs_{matchable}^i \neq \emptyset \\ tu^c, p \vdash_p^c t\Gamma_p^c \quad t\Gamma_p^c \oplus t\Gamma^c \vdash^c e_1 \mapsto [td_1^c \mid d_1^c \mid vs_1^i] \end{array}}{t\Gamma^c \vdash^c \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [d^c \cup td^c \cup td_1^c \mid d^c \cup d_1^c \mid vs_1^i]}$				
C-MATCH-VAR	$\frac{\begin{array}{l} t\Gamma^c \vdash^c e \mapsto tu^c \quad tu^c = [td^c \mid d^c \mid vs^i] \quad vs^i \cap vs_{matchable}^i \neq \emptyset \\ tu^c, p \vdash_p^c \perp \quad (x, tu^c) \oplus t\Gamma^c \vdash^c e_2 \mapsto [td_2^c \mid d_2^c \mid vs_2^i] \end{array}}{t\Gamma^c \vdash^c \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [d^c \cup td^c \cup td_2^c \mid d^c \cup d_2^c \mid vs_2^i]}$				
C-MATCH-UNKNOWN	$\frac{\begin{array}{l} t\Gamma^c \vdash^c e \mapsto tu^c \quad tu^c = [td^c \mid d^c \mid vs^i] \quad vs^i \cap vs_{matchable}^i \neq \emptyset \\ tu^c, p \vdash_p^c ? \quad t\Gamma_p^c \oplus t\Gamma^c \vdash^c e_1 \mapsto [td_1^c \mid d_1^c \mid vs_1^i] \\ (x, tu^c) \oplus t\Gamma^c \vdash^c e_2 \mapsto [td_2^c \mid d_2^c \mid vs_2^i] \end{array}}{t\Gamma^c \vdash^c \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [d^c \cup td^c \cup td_1^c \cup td_2^c \mid d^c \cup d_1^c \cup d_2^c \mid vs_1^i \cup vs_2^i]}$				

$$\frac{\text{C-MATCH-EMPTY} \quad t\Gamma^c \vdash^c e \mapsto tu^c \quad tu^c = [td^c \mid d^c \mid vs^i] \quad vs^i \cap vs_{matchable}^i = \emptyset}{t\Gamma^c \vdash^c \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [\emptyset \mid \emptyset \mid \emptyset]}$$

$$\frac{\text{C-CONSTR-0}}{t\Gamma^c \vdash^c C \mapsto [\emptyset \mid \emptyset \mid \{C\}]} \quad \frac{\text{C-CONSTR-1} \quad t\Gamma^c \vdash^c e \mapsto [td^c \mid d^c \mid vs^i]}{t\Gamma^c \vdash^c D(e) \mapsto [td^c \mid \emptyset \mid \{D([d^c \mid v^i]) \mid v^i \in vs^i\}]}$$

$$\frac{\text{C-COUPLE} \quad t\Gamma^c \vdash^c e_1 \mapsto [td_1^c \mid d_1^c \mid vs_1^i] \quad t\Gamma^c \vdash^c e_2 \mapsto [td_2^c \mid d_2^c \mid vs_2^i]}{t\Gamma^c \vdash^c (e_1, e_2) \mapsto [td_1^c \cup td_2^c \mid \emptyset \mid \{([d_1^c \mid v_1^i], [d_2^c \mid v_2^i]) \mid v_1^i \in vs_1^i \wedge v_2^i \in vs_2^i\}]}$$

$$\frac{\text{C-LETIN} \quad t\Gamma^c \vdash^c e_1 \mapsto tu_1^c \quad tu_1^c = [td_1^c \mid d_1^c \mid vs_1^i] \quad vs_1^i \neq \emptyset \quad (x, tu_1^c) \oplus t\Gamma^c \vdash^c e_2 \mapsto [td_2^c \mid d_2^c \mid vs_2^i]}{t\Gamma^c \vdash^c \text{let } x = e_1 \text{ in } e_2 \mapsto [td_1^c \cup td_2^c \mid d_2^c \mid vs_2^i]}$$

$$\frac{\text{C-LETIN-EMPTY} \quad t\Gamma^c \vdash^c e_1 \mapsto [td_1^c \mid d_1^c \mid \emptyset]}{t\Gamma^c \vdash^c \text{let } x = e_1 \text{ in } e_2 \mapsto [\emptyset \mid \emptyset \mid \emptyset]} \quad \frac{\text{C-ANNOT} \quad t\Gamma^c \vdash^c e \mapsto [td^c \mid d^c \mid vs^i]}{t\Gamma^c \vdash^c l : e \mapsto [td^c \mid \{l\} \cup d^c \mid vs^i]}$$

$$vs_{matchable}^i = \{C \mid \forall C \in Constr^0\} \cup \{D(u^i) \mid \forall D \in Constr^1, \forall u^i\} \cup \{(u_1^i, u_2^i) \mid \forall u_1^i, \forall u_2^i\}$$

$$\frac{\text{CM-CONSTR-0} \quad vs^i \cap vs_{matchable}^i \subseteq \{C \mid \forall C \in Constr^0\}}{[td^c \mid d^c \mid vs^i], C \vdash_p^c \{ \}}$$

$$\frac{\text{CM-CONSTR-1} \quad vs^i \cap vs_{matchable}^i \subseteq \{D(u^i) \mid \forall D \in Constr^1, \forall u^i\} \quad vs^{i'} = \{v^i \mid \exists d^i. D([d^i \mid v^i]) \in vs^i\} \quad \forall l. l \in d'^c \Leftrightarrow \exists d^i. l \in d^i \wedge \exists v^i. D([d^i \mid v^i]) \in vs^i}{[td^c \mid d^c \mid vs^i], D(x) \vdash_p^c \{(x, [\emptyset \mid d'^c \mid vs^{i'}])\}}$$

$$\frac{\text{CM-COUPLE} \quad vs^i \cap vs_{matchable}^i \subseteq \{(u_1^i, u_2^i) \mid \forall u_1^i, \forall u_2^i\} \quad vs_1^i = \{v_1^i \mid \exists (d_1^i, d_2^i, v_2^i). ([d_1^i \mid v_1^i], [d_2^i \mid v_2^i]) \in vs^i\} \quad \forall l. l \in d_1^c \Leftrightarrow \exists d_1^i. l \in d_1^i \wedge \exists (v_1^i, d_2^i, v_2^i). ([d_1^i \mid v_1^i], [d_2^i \mid v_2^i]) \in vs^i \quad vs_2^i = \{v_2^i \mid \exists (d_1^i, v_1^i, d_2^i). ([d_1^i \mid v_1^i], [d_2^i \mid v_2^i]) \in vs^i\} \quad \forall l. l \in d_2^c \Leftrightarrow \exists d_2^i. l \in d_2^i \wedge \exists (d_1^i, v_1^i, v_2^i). ([d_1^i \mid v_1^i], [d_2^i \mid v_2^i]) \in vs^i}{[td^c \mid d^c \mid vs^i], (x_1, x_2) \vdash_p^c \{(x_1, [\emptyset \mid d_1^c \mid vs_1^i]); (x_2, [\emptyset \mid d_2^c \mid vs_2^i])\}}$$

$$\frac{\text{CM-CONSTR-0-NOT} \quad vs^i \cap \{C \mid \forall C \in Constr^0\} = \emptyset}{[td^c \mid d^c \mid vs^i], C \vdash_p^c \perp} \quad \frac{\text{CM-CONSTR-1-NOT} \quad vs^i \cap \{D(u^i) \mid \forall D \in Constr^1, \forall u^i\} = \emptyset}{[td^c \mid d^c \mid vs^i], D(x) \vdash_p^c \perp}$$

$$\frac{\text{CM-COUPLE-NOT} \quad vs^i \cap \{(u_1^i, u_2^i) \mid \forall u_1^i, \forall u_2^i\} = \emptyset}{[td^c \mid d^c \mid vs^i], (x_1, x_2) \vdash_p^c \perp}$$

CM-CONSTR-0-UNKNOWN

$$\frac{vs^i \cap vs_{matchable}^i \not\subseteq \{C \mid \forall C \in Constr^0\} \quad vs^i \cap \{C \mid \forall C \in Constr^0\} \neq \emptyset}{[td^c \mid d^c \mid vs^i], C \vdash_p^{c?} \{\}} \quad \{\}$$

CM-CONSTR-1-UNKNOWN

$$\frac{\begin{array}{l} vs^i \cap vs_{matchable}^i \not\subseteq \{D(u^i) \mid \forall D \in Constr^1, \forall u^i\} \\ vs^i \cap \{D(u^i) \mid \forall D \in Constr^1, \forall u^i\} \neq \emptyset \\ vs^i = \{v^i \mid \exists d^i. D([d^i \mid v^i]) \in vs^i \cap \{D(u^i) \mid \forall D \in Constr^1, \forall u^i\}\} \\ \forall l. l \in d^c \Leftrightarrow \exists d^i. l \in d^i \wedge \exists v^i. D([d^i \mid v^i]) \in vs^i \cap \{D(u^i) \mid \forall D \in Constr^1, \forall u^i\} \end{array}}{[td^c \mid d^c \mid vs^i], D(x) \vdash_p^{c?} \{(x, [\emptyset \mid d'^c \mid vs'^i])\}} \quad \{\}$$

CM-COUPLE-UNKNOWN

$$\frac{\begin{array}{l} vs^i \cap vs_{matchable}^i \not\subseteq \{(u_1^i, u_2^i) \mid \forall u_1^i, \forall u_2^i\} \quad vs^i \cap \{(u_1^i, u_2^i) \mid \forall u_1^i, \forall u_2^i\} \neq \emptyset \\ vs_1^i = \{v_1^i \mid \exists (d_1^i, d_2^i, v_2^i). ([d_1^i \mid v_1^i], [d_2^i \mid v_2^i]) \in vs^i \cap \{(u_1^i, u_2^i) \mid \forall u_1^i, \forall u_2^i\}\} \\ \forall l. l \in d_1^c \Leftrightarrow \exists d_1^i. l \in d_1^i \wedge \exists (v_1^i, d_2^i, v_2^i). ([d_1^i \mid v_1^i], [d_2^i \mid v_2^i]) \in vs^i \cap \{(u_1^i, u_2^i) \mid \forall u_1^i, \forall u_2^i\} \\ vs_2^i = \{v_2^i \mid \exists (d_1^i, v_1^i, d_2^i). ([d_1^i \mid v_1^i], [d_2^i \mid v_2^i]) \in vs^i \cap \{(u_1^i, u_2^i) \mid \forall u_1^i, \forall u_2^i\}\} \\ \forall l. l \in d_2^c \Leftrightarrow \exists d_2^i. l \in d_2^i \wedge \exists (d_1^i, v_1^i, v_2^i). ([d_1^i \mid v_1^i], [d_2^i \mid v_2^i]) \in vs^i \cap \{(u_1^i, u_2^i) \mid \forall u_1^i, \forall u_2^i\} \end{array}}{[td^c \mid d^c \mid vs^i], (x_1, x_2) \vdash_p^{c?} \{(x_1, [\emptyset \mid d_1^c \mid vs_1^i]); (x_2, [\emptyset \mid d_2^c \mid vs_2^i])\}} \quad \{\}$$

9 Abstract Semantics

9.1 Values/Environments

$tu^a := [td^a \mid u^a]$ Valeur avec annotation de t -dependance

$u^a := [d^a \mid v^a]$ Valeur avec annotation de v -dependance

$v^a :=$ $C \mid D(u^a) \mid (u_1^a, u_2^a)$ Constructeurs de donnes
 $< \lambda x.e, \Gamma^a >$ Fermeture
 $< \text{recf}.x.e, \Gamma^a >$ Fermeture rcursive
 \perp Aucune valeur
 \top Valeur quelconque

$td^a := \{l_1; \dots; l_n\}$ t -dependances
 $d^a := \{l_1; \dots; l_m\}$ v -dependances

$t\Gamma^a := (x_1, tu_1^a); \dots; (x_n, tu_n^a)$ t -environnement abstrait

$\Gamma^a := (x_1, u_1^a); \dots; (x_n, u_n^a)$ v -environnement abstraite

9.2 Inference Rules

$$\begin{array}{c}
\text{A-NUM} \\
\frac{}{t\Gamma^a \vdash^a n \mapsto [\emptyset \mid [\emptyset \mid \top]]} \\
\\
\text{A-IDENT} \quad \frac{tu^a = t\Gamma^a[x]}{t\Gamma^a \vdash^a x \mapsto tu^a} \quad \text{A-IDENT-EMPTY} \quad \frac{x \notin \text{support}(t\Gamma^a)}{t\Gamma^a \vdash^a x \mapsto [\emptyset \mid [\emptyset \mid \perp]]} \\
\\
\text{A-ABSTR} \\
\frac{}{t\Gamma^a \vdash^a \lambda x.e \mapsto [\emptyset \mid [\emptyset \mid <\lambda x.e, \uparrow_{ta}^a(t\Gamma^a) >]]} \\
\\
\text{A-ABSTR-REC} \\
\frac{}{t\Gamma^a \vdash^a \mathbf{rec}f.x.e \mapsto [\emptyset \mid [\emptyset \mid <\mathbf{rec}f.x.e, \uparrow_{ta}^a(t\Gamma^a) >]]} \\
\\
\text{A-LETIN} \\
\frac{\begin{array}{l} t\Gamma^a \vdash^a e_1 \mapsto tu_1^a \quad tu_1^a = [td_1^a \mid u_1^a] \\ (x, tu_1^a) \oplus t\Gamma^a \vdash^a e_2 \mapsto [td_2^a \mid u_2^a] \end{array}}{t\Gamma^a \vdash^a \text{let } x = e_1 \text{ in } e_2 \mapsto [td_1^a \cup td_2^a \mid u_2^a]} \\
\\
\text{A-APPLY} \\
\frac{\begin{array}{l} t\Gamma^a \vdash^a e_1 \mapsto [td_1^a \mid [d_1^a \mid <\lambda x.e, \Gamma_1^a >]] \\ t\Gamma^a \vdash^a e_2 \mapsto tu_2^a \quad tu_2^a = [td_2^a \mid u_2^a] \\ (x, tu_2^a) \oplus \uparrow_a^{ta}(\Gamma_1^a) \vdash^a e \mapsto [td^a \mid [d^a \mid v^a]] \end{array}}{t\Gamma^a \vdash^a e_1 e_2 \mapsto [td_1^a \cup td_2^a \cup td^a \cup d_1^a \mid [d_1^a \cup d^a \mid v^a]]} \\
\\
\text{A-APPLY-REC} \\
\frac{\begin{array}{l} t\Gamma^a \vdash^a e_1 \mapsto [td_1^a \mid [d_1^a \mid <\mathbf{rec}f.x.e, \Gamma_1^a >]] \\ tu_1^a = [td_1^a \mid [d_of_freevars(\mathbf{rec}f.x.e, t\Gamma^a) \mid \top]] \\ t\Gamma^a \vdash^a e_2 \mapsto tu_2^a \quad tu_2^a = [td_2^a \mid u_2^a] \\ (f, tu_1^a) \oplus (x, tu_2^a) \oplus \uparrow_a^{ta}(\Gamma_1^a) \vdash^a e \mapsto [td^a \mid [d^a \mid v^a]] \end{array}}{t\Gamma^a \vdash^a e_1 e_2 \mapsto [td_1^a \cup td_2^a \cup td^a \cup d_1^a \mid [d_1^a \cup d^a \mid v^a]]} \\
\\
\text{A-APPLY-UNKNOWN} \\
\frac{\begin{array}{l} t\Gamma^a \vdash^a e_1 \mapsto [td_1^a \mid [d_1^a \mid v_1^a]] \quad t\Gamma^a \vdash^a e_2 \mapsto [td_2^a \mid [d_2^a \mid v_2^a]] \\ \forall (f, x, e, \Gamma_1^a). v_1^a \neq <\lambda x.e, \Gamma_1^a > \wedge v_1^a \neq <\mathbf{rec}f.x.e, \Gamma_1^a > \end{array}}{t\Gamma^a \vdash^a e_1 e_2 \mapsto [td_1^a \cup td_2^a \cup d_1^a \cup d_2^a \mid [d_1^a \cup d_2^a \mid \top]]} \\
\\
\text{A-IF} \\
\frac{\begin{array}{l} t\Gamma^a \vdash^a e \mapsto [td^a \mid [d^a \mid v^a]] \\ t\Gamma^a \vdash^a e_1 \mapsto [td_1^a \mid [d_1^a \mid v_1^a]] \quad t\Gamma^a \vdash^a e_2 \mapsto [td_2^a \mid [d_2^a \mid v_2^a]] \end{array}}{t\Gamma^a \vdash^a \text{if } e \text{ then } e_1 \text{ else } e_2 \mapsto [d^a \cup td^a \cup td_1^a \cup td_2^a \mid [d^a \cup d_1^a \cup d_2^a \mid \top]]} \\
\\
\text{A-MATCH} \\
\frac{\begin{array}{l} t\Gamma^a \vdash^a e \mapsto tu^a \quad tu^a = [td^a \mid [d^a \mid v^a]] \quad tu^a, p \vdash_p^a t\Gamma_p^a \\ t\Gamma_p^a \oplus t\Gamma^a \vdash^a e_1 \mapsto [td_1^a \mid [d_1^a \mid v_1^a]] \end{array}}{t\Gamma^a \vdash^a \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [d^a \cup td^a \cup td_1^a \mid [d^a \cup d_1^a \mid v_1^a]]} \\
\\
\text{A-MATCH-VAR} \\
\frac{\begin{array}{l} t\Gamma^a \vdash^a e \mapsto tu^a \quad tu^a = [td^a \mid [d^a \mid v^a]] \quad tu^a, p \vdash_p^a \perp \\ (x, tu^a) \oplus t\Gamma^a \vdash^a e_2 \mapsto [td_2^a \mid [d_2^a \mid v_2^a]] \end{array}}{t\Gamma^a \vdash^a \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [d^a \cup td^a \cup td_2^a \mid [d^a \cup d_2^a \mid v_2^a]]}
\end{array}$$

A-MATCH-UNKNOWN

$$\frac{\begin{array}{c} t\Gamma^a \vdash^a e \mapsto tu^a \quad tu^a = [td^a \mid [d^a \mid v^a]] \quad tu^a, p \vdash_p^a? t\Gamma_p^a \\ t\Gamma_p^a \oplus t\Gamma^a \vdash^a e_1 \mapsto [td_1^a \mid [d_1^a \mid v_1^a]] \\ (x, tu^a) \oplus t\Gamma^a \vdash^a e_2 \mapsto [td_2^a \mid [d_2^a \mid v_2^a]] \end{array}}{t\Gamma^a \vdash^a \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [d^a \cup td^a \cup td_1^a \cup td_2^a \mid [d^a \cup d_1^a \cup d_2^a \mid \top]]}$$

A-MATCH-ERROR

$$\frac{t\Gamma^a \vdash^a e \mapsto [td^a \mid [d^a \mid v^a]] \quad tu^a, p \vdash_p^a \times}{t\Gamma^a \vdash^a \text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2 \mapsto [\emptyset \mid [\emptyset \mid \perp]]}$$

A-CONSTR-0

A-CONSTR-1

$$\frac{}{t\Gamma^a \vdash^a C \mapsto [\emptyset \mid [\emptyset \mid C]]} \quad \frac{t\Gamma^a \vdash^a e \mapsto [td^a \mid u^a]}{t\Gamma^a \vdash^a D(e) \mapsto [td^a \mid [\emptyset \mid D(u^a)]]}$$

A-COUPLE

A-ANNOT

$$\frac{t\Gamma^a \vdash^a e_1 \mapsto [td_1^a \mid u_1^a] \quad t\Gamma^a \vdash^a e_2 \mapsto [td_2^a \mid u_2^a]}{t\Gamma^a \vdash^a (e_1, e_2) \mapsto [td_1^a \cup td_2^a \mid [\emptyset \mid (u_1^a, u_2^a)]]} \quad \frac{t\Gamma^a \vdash^a e \mapsto [td^a \mid [d^a \mid v^a]]}{t\Gamma^a \vdash^a l : e \mapsto [td^a \mid [l; d^a \mid v^a]]}$$

$$d_of_freevars(e, t\Gamma^a) := d_of_aux(e, t\Gamma^a, \emptyset, \emptyset)$$

$$d_of_aux(n, t\Gamma^a, bvars, acc) := acc$$

$$d_of_aux(C, t\Gamma^a, bvars, acc) := acc$$

$$d_of_aux(D(e), t\Gamma^a, bvars, acc) := d_of_aux(e, t\Gamma^a, bvars, acc)$$

$$d_of_aux(x, t\Gamma^a, bvars, acc) := acc \text{ si } x \in bvars$$

$$d_of_aux(x, t\Gamma^a, bvars, acc) := d^a \cup acc \text{ si } x \notin bvars \wedge t\Gamma^a[x] = [td^a \mid [d^a \mid v^a]]$$

$$d_of_aux(\lambda x.e, t\Gamma^a, bvars, acc) := d_of_aux(e, t\Gamma^a, x; bvars, acc)$$

$$d_of_aux(\text{recf}.x.e, t\Gamma^a, bvars, acc) := d_of_aux(e, t\Gamma^a, f; x; bvars, acc)$$

$$d_of_aux(e_1 \ e_2, t\Gamma^a, bvars, acc) := d_of_aux(e_2, t\Gamma^a, bvars, acc_1)$$

$$d_of_aux(\text{if } e \text{ then } e_1 \text{ else } e_2, t\Gamma^a, bvars, acc) := \text{pour } acc_1 = d_of_aux(e_1, t\Gamma^a, bvars, acc)$$

$$d_of_aux(\text{match } e \text{ with } p \rightarrow e_1 \mid x \rightarrow e_2, t\Gamma^a, bvars, acc) := \text{pour } acc_1 = d_of_aux(e_1, t\Gamma^a, bvars, acc_0) \\ \text{et } acc_0 = d_of_aux(e, t\Gamma^a, bvars, acc) \\ \text{pour } acc_1 = d_of_aux(e_2, t\Gamma^a, x; bvars, acc_0) \\ \text{et } acc_0 = d_of_aux(e, t\Gamma^a, bvars, acc)$$

$$d_of_aux((e_1, e_2), t\Gamma^a, bvars, acc) := d_of_aux(e_2, t\Gamma^a, bvars, acc_1)$$

$$d_of_aux(l : e, t\Gamma^a, bvars, acc) := d_of_aux(e, t\Gamma^a, bvars, acc)$$

$$d_of_aux(\text{let } x = e_1 \text{ in } e_2, t\Gamma^a, bvars, acc) := d_of_aux(e_2, t\Gamma^a, x; bvars, acc_1)$$

$$d_of_aux(\text{let } x = e_1 \text{ in } e_2, t\Gamma^a, bvars, acc) := \text{pour } acc_1 = d_of_aux(e_1, t\Gamma^a, bvars, acc)$$

$$binders_of(C) := \{\}$$

$$binders_of(D(x)) := \{x\}$$

$$binders_of((x, y)) := \{x; y\}$$

AM-CONSTR-0

AM-CONSTR-1

$$\frac{}{[td^a \mid [d^a \mid C]], C \vdash_p^a \{\}} \quad \frac{}{[td^a \mid [d^a \mid D(u^a)]], D(x) \vdash_p^a \{(x, [\emptyset \mid u^a])\}}$$

AM-COUPLE

$$\frac{}{[td^a \mid [d^a \mid (u_1^a, u_2^a)]], (x_1, x_2) \vdash_p^a \{(x_1, [\emptyset \mid u_1^a]); (x_2, [\emptyset \mid u_2^a])\}}$$

AM-CONSTR-0-NOT

AM-CONSTR-1-NOT

$$\frac{p \neq C}{[td^a \mid [d^a \mid C]], p \vdash_p^a \perp} \quad \frac{p \neq D(-)}{[td^a \mid [d^a \mid D(u^a)]], p \vdash_p^a \perp}$$

AM-COUPLE-NOT

$$\frac{p \neq (-, -)}{[td^a \mid [d^a \mid (u_1^a, u_2^a)]], p \vdash_p^a \perp}$$

AM-CONSTR-0-UNKNOWN

AM-CONSTR-1-UNKNOWN

$$\frac{[td^a \mid [d^a \mid \top]], C \vdash_p^a? \{ \}}{[td^a \mid [d^a \mid \top]], D(x) \vdash_p^a? \{ (x, [\emptyset \mid [\emptyset \mid \top]]) \}}$$

AM-COUPLE-UNKNOWN

$$\frac{[td^a \mid [d^a \mid \top]], (x_1, x_2) \vdash_p^a? \{ (x_1, [\emptyset \mid [\emptyset \mid \top]]); (x_2, [\emptyset \mid [\emptyset \mid \top]]) \}}{[td^a \mid [d^a \mid v^a]], p \vdash_p^a \times}$$

AM-ERROR

$$\frac{\forall C. v^a \neq C \quad \forall D. v^a \neq D(-) \quad v^a \neq (-, -) \quad v^a \neq \top}{[td^a \mid [d^a \mid v^a]], p \vdash_p^a \times}$$