Introduction to Big Data with Apache Spark







BerkeleyX

This Lecture

Structured Data and Relational Databases

The Structured Query Language (SQL)

SQL and pySpark Joins

Review: Key Data Management Concepts

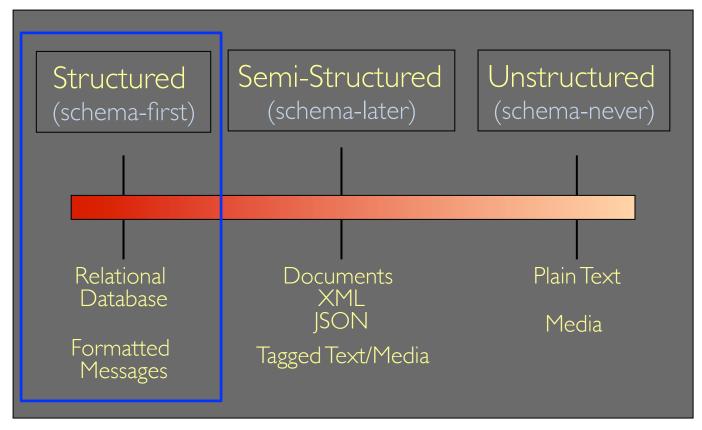
- A data model is a collection of concepts for describing data
- A schema is a description of a particular collection of data, using a given data model

Whither Structured Data?

- Conventional Wisdom:
 - » Only 20% of data is structured.
- The state of the s

- Decreasing due to:
 - » Consumer applications
 - » Enterprise search
 - » Media applications

The Structure Spectrum



This lecture

Relational Database: Definitions

- Relational database: a set of relations
- Two parts to a Relation:

Schema: specifies name of relation, plus each column's name and type

```
Students(sid: string, name: string, email: string, age: integer, gpa: real)
```

Instance: the actual data at a given time

- #rows = cardinality
- #fields = degree

Review: Key Data Management Concepts

- A data model is a collection of concepts for describing data
- A schema is a description of a particular collection of data, using a given data model
- A relational data model is the most used data model
 - » Relation, a table with rows and columns
 - » Every relation has a schema defining fields in columns

What is a Database?

- A large organized collection of data
 - » Transactions used to modify data

- Models real world, e.g., enterprise
 - » Entities (e.g., teams, games)
 - » Relationships, e.g.,
 - » A plays against B in The World Cup

Large Databases

- US Internal Revenue Service: <u>150 Terabytes</u>
- Australian Bureau of Stats: <u>250 Terabytes</u>
- AT&T call records: 312 Terabytes
- eBay database: <u>I.4 Petabytes</u>
- Yahoo click data: <u>2 Petabytes</u>
- What matters for these databases?

Large Databases

- US Internal Revenue Service: 150 Terabytes Consistency, Durability, Rich queries
- Australian Bureau of Stats: <u>250 Terabytes</u> ← Fast, Rich queries
- AT&T call records: 312 Terabytes

 Accuracy, Consistency, Durability
- eBay database: I.4 Petabytes
 Availability
 Timeliness
- Yahoo click data: <u>2 Petabytes</u>
- What matters for these databases?

Example: Instance of Students Relation

Students(sid:string, name:string, login:string, age:integer, gpa:real)

sid	name	login	age	gpa
53666	Jones	jones@eecs	18	3.4
53688	Smith	smith@statistics	18	3.2
53650	Smith	smith@math	19	3.8

- Cardinality = 3 (rows)
- Degree = 5 (columns)
- All rows (tuples) are distinct

Relational Databases

- Advantages
 - » Well-defined structure
 - » Maintains indices for high performance
 - » Consistency maintained by transactions
- Disadvantages
 - » Limited, rigid structure
 - » Most of disk space is taken up by large indices
 - » Transactions are slow
 - » Poor support for sparse data

Sparse Data

- Very sparse data is common today
 - » Want to store data with thousands of columns
 - » But, not all rows have values for all columns
- Typical database tables might have dozens of columns
- Tables are very wasteful for sparse data

SQL - A language for Relational DBs

- <u>SQL</u> = Structured Query Language
- Supported by pySpark DataFrames (<u>SparkSQL</u>)
- Some of the functionality SQL provides:
 - » Create, modify, delete relations
 - » Add, modify, remove tuples
 - » Specify queries to find tuples matching criteria

Queries in SQL

- Single-table queries are straightforward
- To find all 18 year old students, we can write:

```
SELECT *
  FROM Students S
WHERE S.age=18
```

To find just names and logins:

```
SELECT S.name, S.login
  FROM Students S
WHERE S.age=18
```

Querying Multiple Relations

Can specify a join over two tables as follows:

```
SELECT S.name, E.cid
FROM Students S, Enrolled E
WHERE S.sid=E.sid
```

Students

Enrolled

F	E.sid	E.cid	E.grade
L	53831	Physics203	Α
	53650	Topology112	А
	53341	History105	В

	S.sid	S.name	S.login	S.age	S.gpa
,	53341	Jones	jones@cs	18	3.4
	53831	Smith	smith@ee	18	3.2

First, combine the two tables, S and E

Cross Join

• Cartesian product of two tables $(E \times S)$:

Enrolled Students

F	E.sid	E.cid	E.grade
L	53831	Physics203	А
	53650	Topology112	Α
	53341	History105	В

ς	S.sid S.name		S.login	S.age	S.gpa
	53341	Jones	jones@cs	18	3.4
	53831	Smith	smith@ee	18	3.2

Cross Join

• Cartesian product of two tables $(E \times S)$:

Enrolled Students

 E.sid
 E.cid
 E.grade

 53831
 Physics203
 A

 53650
 Topology112
 A

 53341
 History105
 B

, ,	S.sid	S.name	S.login	S.age	S.gpa
•	53341	Jones	jones@cs	18	3.4
	53831	Smith	smith@ee	18	3.2

E.sid	E.cid	E.grade	S.sid	S.name	S.login	S.age	S.gpa
53831	Physics203	Α	53341	Jones	jones@cs	18	3.4
53650	Topology112	А	53341	Jones	jones@cs	18	3.4
53341	History105	В	53341	Jones	jones@cs	18	3.4
53831	Physics203	Α	53831	Smith	smith@ee	18	3.2
53650	Topology112	А	53831	Smith	smith@ee	18	3.2
53341	History105	В	53831	Smith	smith@ee	18	3.2

Where Clause

Choose matching rows using Where clause:

```
SELECT S.name, E.cid
FROM Students S, Enrolled E
WHERE S.sid=E.sid
```

E.sid	E.cid	E.grade	S.sid	S.name	S.login	S.age	S.gpa
53831	Physics203	А	53341	Jones	jones@cs	18	3.4
53650	Topology112	А	53341	Jones	jones@cs	18	3.4
53341	History105	В	53341	ones	jones@cs	18	3.4
53831	Physics203	A	53831	mith	smith@ee	18	3.2
53650	Topology112	А	53831	Smith	smith@ee	18	3.2
53341	History105	В	53831	Smith	smith@ee	18	3.2

Select Clause

• Filter columns using Select clause:

```
SELECT S.name, E.cid
FROM Students S, Enrolled E
WHERE S.sid=E.sid
```

E.sid	E.cid	E.grade	S.sid	S.name	S.login	S.age	S.gpa
53831	Physics203	А	53341	Jones	jones@cs	18	3.4
53650	Topology112	А	53341	Jones	jones@cs	18	3.4
53341	History105	В	53341	ones	jones@cs	18	3.4
53831	Physics203	A	53831	mith	smith@ee	18	3.2
53650	Topology112	А	53831	Smith	smith@ee	18	3.2
53341	History105	В	53831	Smith	smith@ee	18	3.2

Result

Can specify a join over two tables as follows:

```
SELECT S.name, E.cid
FROM Students S, Enrolled E
WHERE S.sid=E.sid
```

Students

Enrolled

F	E.sid	E.cid	E.grade
(53831	Physics203	А
	53650	Topology112	А
	53341	History105	В

ς	S.sid	S.name	S.login	S.age	S.gpa
	53341	Jones	jones@cs	18	3.4
	53831	Smith	smith@ee	18	3.2

$$Result = \begin{bmatrix} S.name & E.cid \\ Jones & History105 \\ Smith & Physics203 \end{bmatrix}$$

Explicit SQL Joins

SELECT S.name, E.classid FROM Students S INNER JOIN Enrolled E ON S.sid=E.sid



E.sid E.classid

11111 History105

11111 DataScience194

22222 French150

44444 English10

Result

S.name	E.classid
Jones	History105
Jones	DataScience194
Smith	French150

Equivalent SQL Join Notations

• Explicit Join notation (preferred):

```
SELECT S.name, E.classid
FROM Students S INNER JOIN Enrolled E ON S.sid=E.sid
```

```
SELECT S.name, E.classid FROM Students S JOIN Enrolled E ON S.sid=E.sid
```

Implicit join notation (deprecated):

```
SELECT S.name, E.cid
FROM Students S, Enrolled E
WHERE S.sid=E.sid
```

SQL Types of Joins

SELECT S.name, E.classid FROM Students S INNER JOIN Enrolled E ON S.sid=E.sid



•	E.sid	E.classid
	11111	History105
	11111	DataScience194
	22222	French150
	44444	English10

R	ക്യ	ı	l+
\Box	にろし	Л	ΙL

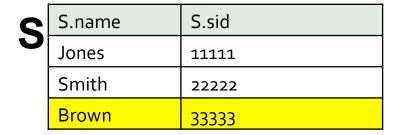
S.name	E.classid
Jones	History105
Jones	DataScience194
Smith	French150

Unmatched keys

The type of join controls how unmatched keys are handled

SQL Joins: Left Outer Join

SELECT S.name, E.classid
FROM Students S LEFT OUTER JOIN Enrolled E ON S.sid=E.sid



	E.sid	E.classid
E	11111	History105
	11111	DataScience194
	22222	French150
	44444	English10

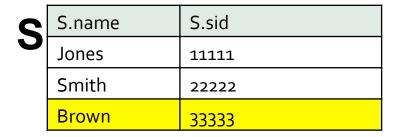
Result

-	S.name	E.classid
	Jones	History105
	Jones	DataScience194
	Smith	French150
	Brown	<null></null>

Unmatched keys

SQL Joins: Right Outer Join

SELECT S.name, E.classid
FROM Students S RIGHT OUTER JOIN Enrolled E ON S.sid=E.sid



Ε	E.sid	E.classid
	11111	History105
	11111	DataScience194
	22222	French150
	44444	English10

Result

-	S.name	E.classid
	Jones	History105
	Jones	DataScience194
	Smith	French150
	<null></null>	English10

Unmatched keys

Spark Joins

- SparkSQL and Spark DataFrames join() supports:
 - » inner, outer, left outer, right outer, semijoin
- For Pair RDDs, pySpark supports:
 - » inner join(), leftOuterJoin(), rightOuterJoin(), fullOuterJoin()

X.join(Y)

- » Return RDD of all pairs of elements with matching keys in X and Y
- » Each pair is (k, (vI, v2)) tuple, where (k, vI) is in X and (k, v2) is in Y

```
>>> x = sc.parallelize([("a", 1), ("b", 4)])
>>> y = sc.parallelize([("a", 2), ("a", 3)])
>>> sorted(x.join(y).collect())

Value: [('a', (1, 2)), ('a', (1, 3))]
```

X.leftOuterJoin(Y)

- » For each element (k, v) in X, resulting RDD will either contain
 - All pairs (k, (v, w)) for w in Y,
 - Or the pair (k, (v, None)) if no elements in Y have key k

```
>>> x = sc.parallelize([("a", 1), ("b", 4)])
>>> y = sc.parallelize([("a", 2)])
>>> sorted(x.leftOuterJoin(y).collect())

Value: [('a', (1, 2)), ('b', (4, None))]
```

- Y.rightOuterJoin(X)
 - » For each element (k, w) in X, resulting RDD will either contain
 - All pairs (k, (v, w)) for v in Y,
 - Or the pair (k, (None, w)) if no elements in Y have key k

```
>>> x = sc.parallelize([("a", 1), ("b", 4)])
>>> y = sc.parallelize([("a", 2)])
>>> sorted(y.rightOuterJoin(x).collect())

Value: [('a', (2, 1)), ('b', (None, 4))]
```

- X.fullOuterJoin(Y)
 - » For each element (k, v) in X, resulting RDD will either contain
 - All pairs (k, (v, w)) for w in Y, or (k, (v, None)) if no elements in Y have k
 - » For each element (k, v) in Y, resulting RDD will either contain
 - All pairs (k, (v, w)) for v in X, or (k, (None, w)) if no elements in X have k

```
>>> x = sc.parallelize([("a", 1), ("b", 4)])
>>> y = sc.parallelize([("a", 2), ("c", 8)])
>>> sorted(x.fullOuterJoin(y).collect())

Value: [('a', (1, 2)), ('b', (4, None)), ('c', (None, 8))]
```