CS170 Discussion Section 11: 4/12

Definitions

- 1. P: the set of all problems that can be solved (not just verified) in polynomial time.
- 2. **NP**: the set of all problems that can be verified in polynomial time.
- 3. **NP**-complete: a problem to which all other problems in NP reduce. In other words, if we could efficiently solve any **NP**-complete problem, we could efficiently solve any problem in NP.

Reduction Basics

Assume A reduces to B in polynomial time. In each part you will be given a fact about one of the problems. Determine what, if anything, this allows you to determine about the other problem. (You can answer each part in one sentence.)

- 1. A is in **P**.
- 2. B is in **P**.
- 3. A is **NP**-hard.
- 4. B is **NP**-hard.

Stingy SAT

Stingy SAT is the following problem: given a set of clauses (each a disjunction of literals), and an integer k, find a satisfying assignment in which at most k variables are True, if such an assignment exists. Prove that Stingy SAT is NP-Complete.

More Reductions

1. Give a reduction from Vertex Cover to Set Cover. In Vertex Cover your input is a graph G and a budget b, and you want to find b vertices that touch every edge. In Set Cover, we are given a set E and several subsets of it $S_1, ..., S_m$, along with a budget b. We are asked to select b of these subsets so that their union is E.

2. Give a reduction from 3D Matching to Set Cover.

Reliable Network

Reliable Network is the following problem: We are given two $n \times n$ matrices, a distance matrix d_{ij} , and a connectivity requirement matrix r_{ij} , as well as a budget b; we must find a graph $G = (\{1, ..., n\}, E)$ such that the total cost of all edges is b or less and between any two distinct vertices i and j there are r_{ij} vertex-disjoint paths. Show that Reliable Network is NP-Complete.