## Yate's DP with Bitmasks

# ${\rm md.mottakin.chowdhury}$

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## 1 Theory

Given an array A of  $2^N$  integers, we need to calculate for all mask,

$$F[mask] = \sum_{i \subset mask} A[i].$$

S(mask,i) contains all subsets of mask which differ from it only in the first i bits. Then

$$S(mask,i) = \left\{ \begin{array}{l} S(mask,i-1) \ if \ i^{th} \ bit \ is \ off \\ S(mask,i-1) \cup S(mask \oplus 2^i,i-1) \ otherwise \end{array} \right.$$

### 2 Problems

#### 2.1 Problem 1

You have been given an integer array a on size n. You must report the number of ordered pairs (i,j) such that a[i]&a[j] is 0.

#### 2.2 Problem 2

You are given an array a of n integers. How many subsets of these integers have 0 when applied bitwise and on the subset?

```
const int N=(1<<20);
int n, cnt[N], a;
int dp[N+5];
11 twos[N];
void solve()
{
        for(int mask=0; mask<N; mask++)</pre>
                 dp[mask] = cnt[mask];
        for(int i=0; i<20; i++)</pre>
        {
                 for(int mask=0; mask<N; mask++)</pre>
                 {
                          if(!(mask&(1<<i)))
                                   dp[mask] +=dp[mask^(1<<i)];</pre>
        // dp[mask] - number of a[i] where a[i] @mask=mask
        twos[0]=1;
        for(int i=1; i<N; i++) twos[i]=(twos[i-1]*2LL)%mod;</pre>
        // apply inclusion-exclusion
        ll ans=0;
        for(int i=0; i<N; i++)</pre>
        {
                 int bits=__builtin_popcount(i);
                 ans+=((bits&1) ? (-1) : 1)*twos[dp[i]];
                 ans%=mod;
                 if(ans<0) ans+=mod;</pre>
```

```
}
prnt(ans);
}
```

### 2.3 Problem 3

You are given an array a. For each element of the array, output another element such that their bitwise and equals 0, otherwise output -1.

```
// same as problem 1, except we keep out[mask] which denotes an integer such that
// out[mask]&mask=0
const int MAX=(1<<22)+5;</pre>
int n, a[MAX], out[MAX], cnt[MAX];
11 dp[MAX];
void solve()
        for(int i=0; i<MAX; i++)</pre>
                 if(cnt[i])
                 {
                          dp[i]=cnt[i];
                          out[i]=i;
                 }
        }
        for(int i=0; i<=22; i++)</pre>
                 for(int mask=0; mask<(1<<22); mask++)</pre>
                 {
                          if(mask&(1<<i))
                                   dp[mask] +=dp[mask^(1<<i)];</pre>
                                   if(!out[mask])
                                   {
                                            if(cnt[mask]) out[mask]=mask;
                                            else out[mask] = out[mask^(1<<i)];</pre>
                                   }
                          }
                 }
        }
        FOR(i,0,n)
                 int idx=((1<<22)-1)^a[i];</pre>
                 if(out[idx]==0) printf("-1\n");
                 else printf("%d\n", out[idx]);
```

```
}
```

#### 2.4 Problem 4

You are given an array a of n integers and a vlue k. How many subsets of these integers have a value of k when applied bitwise or on the subset?

```
const int N=(1<<20)+5;
int n, m, k, a[MAX], cnt[N];
11 dp[N], twos[N];
void solve()
{
        for(int i=0; i<(1<<20); i++)
                 dp[i]=cnt[i]; // count of each number in input
        for(int i=0; i<20; i++)</pre>
                 for(int mask=0; mask<(1<<20); mask++)</pre>
                 {
                         if(mask&(1<<i))
                          {
                                  dp[mask] +=dp[mask^(1<<i)];</pre>
                                  dp[mask]%=mod;
                         }
                 }
        }
        ll ans=0;
        for(int i=k; i>=0; i--)
                 if((i|k)!=k) continue;
                 int curr=__builtin_popcount(k^i);
                 ans+=(curr\%2==0 ? 1 : -1)*(twos[dp[i]]-1);
                 ans%=mod;
                 if(ans<0) ans+=mod;</pre>
                 // prnt((1<<curr));
        prnt(ans);
}
```

### 2.5 Problem 5

You are given an array of n numbers and q queries. In each query, you are asked to add a number, delete a number or answer value of f(x) where f(x) denotes the value of the maximum bitwise-and of all the subsequences of length x of the array. The idea is to use Yate's dp along with SQRT Decomposition on queries.

```
const int N=1<<18;
const int bsz=512;</pre>
```

```
int cnt[N], a[MAX], n, m, k, dp[N], idx=0;
pii bucket[bsz+5];
void calc()
{
        FOR(i,0,(1<<k))
                 dp[i]=cnt[i];
        // dp[mask] = total numbers x such that x<math>mask=mask
        for(int i=0; i<k; i++)</pre>
        {
                 for(int mask=0; mask<(1<<k); mask++)</pre>
                         if(!(mask&(1<<i)))
                                  dp[mask] += dp[mask^(1 << i)];
                 }
        }
}
void add(int x)
{
        cnt[x]++;
        bucket[idx]=\{x,1\};
        idx++;
        if(idx==bsz) // bucket full, recalc dp
                 calc();
                 idx=0;
        }
}
void rem(int x)
{
        cnt[x]--;
        bucket[idx]=\{x,-1\};
        if(idx==bsz) // bucket full, recalc dp
                calc();
                 idx=0;
        }
}
// how many values t are there such that t&curr=curr?
int canget(int curr)
{
        int ret=dp[curr];
        FOR(i,0,idx)
        {
                 if((bucket[i].first & curr)==curr)
```

```
ret+=bucket[i].second;
        return ret;
}
int getans(int x)
        int ret=0;
        // try to ensure the higher bits first
        for(int i=k-1; i>=0; i--)
                 int curr=(ret^(1<<i));</pre>
                 // There are total greater equal \boldsymbol{x} numbers t such
                 // that t&curr=curr;
                 if(canget(curr)>=x) // We can achieve curr
                         ret | = (1<<i);
        }
        return ret;
}
void handleQueries()
        int t, x;
        FOR(i,0,m)
                 scanf("%d%d", &t, &x);
                 if(t==1) add(x);
                 else if(t==2) rem(x);
                 else printf("d\n", getans(x));
        }
}
int main()
{
    scanf("%d%d%d", &n, &m, &k);
    FOR(i,0,n)
        scanf("%d", &a[i]);
        cnt[a[i]]++;
    }
    calc();
    handleQueries();
    return 0;
}
```