From MS08-067 to EternalBlue

BsidesMCR, 2017

#whoami

- Denis Isakov a.k.a vinegrep
- Exploit development and RE enthusiast
- Doing pentest as a daily job
- OSCP/OSWP/OSCE/GXPN/GREM/BBQ/WTF and other acronyms

What this talk is about?

•Follow growing exploit development difficulty based on 3 "zero-to-hero" exploits.

 Analyze how these exploits influenced exploit mitigation techniques.

 Should we expect something similar in recent time?

What this talk is not about?

Not going to blame Microsoft.

 Not going to discuss conspiracy theories why Microsoft patch MS17-010 so quickly.

Step-by-step exploit development.

SMB protocol internals



MS08-067 general info

- "Wormable" vulnerability in SMB that was triggered by malformed RPC request.
- Buffer overflow was in the way how Server service processes the PathName argument to the NetprPathCanonicalize function.
- Problem was in directory traversal sequences result in traversing past the root path.

Kudos:

https://blogs.technet.microsoft.com/johnla/2015/09/26/the-inside-story-behind-ms08-067/26/the-inside-story-behind-ms08-06/26/the-insid

https://labs.mwrinfosecurity.com/assets/BlogFiles/hello-ms08-067-my-old-friend.pdf

https://www.mysonicwall.com/sonicalert/searchresults.aspx?ev=article&id=74

https://dontstuffbeansupyournose.com/2008/10/23/looking-at-ms08-067/

MS08-067 exploitation

```
Registers (FPU)
EAX 61736964
ECX 0108F4B2
EDX 0108F508 ASCII "GGGGHHHH"
EBX 0108005C
ESP 0108F47C ASCII "disadisadisadisa
ЕВР 44444444
ESI 0108F4B6
EDI 0108F464
EIP 41414141
C 0 ES 0023 32bit 0(FFFFFFFF)
 1 CS 001B 32bit 0(FFFFFFFF)
   SS 0023 32bit 0(FFFFFFFF)
   DS 0023 32bit 0(FFFFFFFF)
    FS 003B 32bit 7FFAB000(FFF)
T 0 GS 0000 NULL
D 0
O O Lasterr ERROR_SUCCESS (0000000)
0108F45C
         000000000 ....
0108F460
         0108F4B4 'ô
0108F464 4444005C \.DD
0108F468
         4444444 DDDD
0108F46C 4444444 DDDD
0108F470 4444444 DDDD
0108F474 4444444 DDDD
0108F478 41414141 AAAA
0108F47C 61736964 disa
0108F480 61736964 disa
0108F484 61736964 disa
0108F488 61736964 disa
0108F48C 61736964 disa
0108F490 61736964 disa
0108F494 61736964 disa
0108F498 61736964 disa
0108F49C 61736964 disa
0108F4A0 61736964 disa
0108F4A4 61736964 disa
0108F4A8
         47474747 GGGG
0108F4AC
         48484848 HHHH
0108F4B0 000C0000 ....
0108F4B4
         002E005C \...
0108F4B8 005C002E ..\.
```

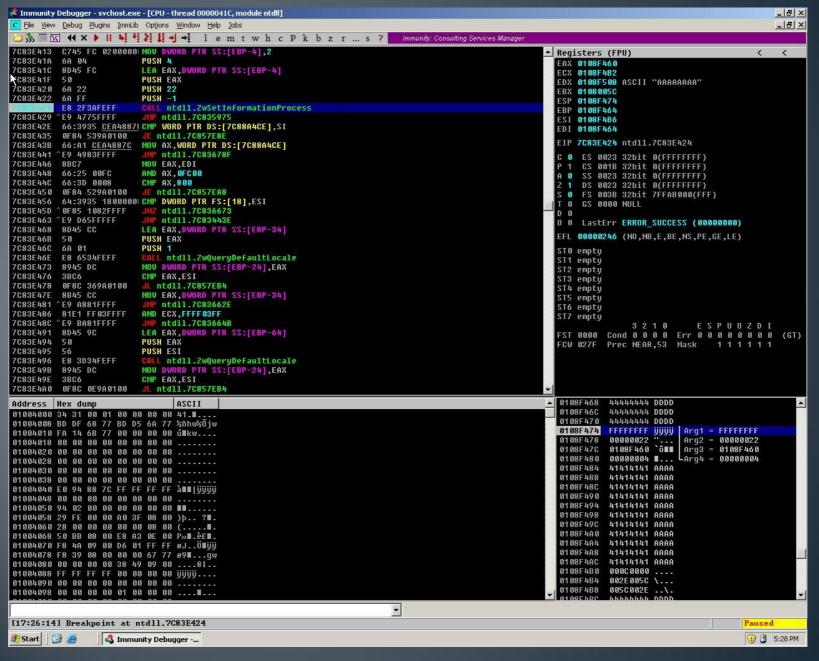
MS08-067 exploitation (2)

- Crash service by overwriting EIP at 18 byte offset.
- ESP and EDX points to our buffer, EAX and EBP was overwritten with our buffer content.
- Overwrite EIP with address of instruction "JMP EDX", as there is enough space between EIP and address where EDX points.
- Store NOP sled + egghunter after our return address.
- Jump to the end of our controller buffer.
- There is a short jump back instruction that will pass execution to egghunter.
- Egghunter will find final shellcode and PROFIT.

MS08-067 NX bypass

- From XP SP2 and 2003 SP1 we have NX bit enabled. NX bit separates between code and data.
- Excellent paper from skape and Skywing how to bypass it http://uninformed.org/?v=2&a=4&t=txt
- Idea: disable DEP for process during execution.
- Search memory for opcodes and check whether they are in executable area.
- LdprCheckNXCompatibility in ntdll.dll is used to check whether NX should be enabled and then based on this check call to ZwSetInformationProcess is used to enable/disable DEP.
- Return to controlled buffer, jump back, egg hunt and final shellcode.

MS08-067 NX bypass before call



MS08-067 NX bypass after call

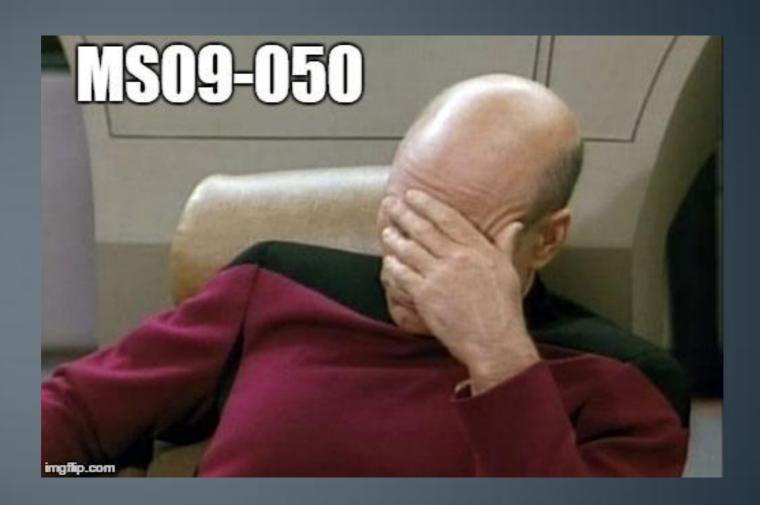
```
7C83E413 C745 FC 0200000 MOV DWORD PTR $5:[EBP-4],2
                                                                                                           ▲ Registers (FPU)
7C83E41A 6A 04
                          PUSH 4
                                                                                                              EAX 00000000
7C83E41C 8D45 FC
                         LEA EAX, DWORD PTR SS:[EBP-4]
                                                                                                              ECX 0108F46C
7C83E41F 50
                          PUSH EAX
                                                                                                              EDX 7C82ED54 ntdll.KiFastSystemCallRet
                          PUSH 22
7C83E420 6A 22
7C83E422 6A FF
                          PUSH -1
                                                                                                              ESP 0108F484 ASCII "AAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
7C83E424 E8 2F3AFEFF
                         CALL ntdll.ZwSetInformationProcess
                                                                                                              EBP 0108F464
         ^E9 4775FFFF
                          JMP ntd11.70835975
                                                                                                              ESI 0108F4B6
7C83E42E 66:3935 CEA48871 CMP WORD PTR DS:[7C88A4CE],SI
                                                                                                              EDI 0108F464
7C83E435 0F84 539A0100 JE ntdl1.7C857E8E
                                                                                                              EIP 7C83E429 ntd11.7C83E429
7C83E43B 66:A1 CEA4887C MOU AX, WORD PTR DS:[7C88A4CE]
7C83E441 ^E9 4983FFFF
                          JMP ntd11.7C83678F
                                                                                                                  ES 0023 32bit 0(FFFFFFFF)
7C83E446 8BC7
                          MOV EAX, EDI
                                                                                                                  CS 001B 32bit 0(FFFFFFFF)
7C83E448 66:25 00FC
                          AND AX, OFCOO
                                                                                                                  SS 0023 32bit 0(FFFFFFFF)
7C83E44C 66:3D 0008
                          CMP AX,800
                                                                                                              Z 1 DS 0023 32bit 0(FFFFFFFF)
7C83E450 0F84 529A0100
                        JE ntd11.7C857EA8
                                                                                                              S 0 FS 003B 32bit 7FFAB000(FFF)
7C83E456 64:3935 1800000 CMP DWORD PTR FS:[18],ESI
                                                                                                              T 0 GS 0000 NULL
                         JNZ ntd11.70836673
7C83E45D ^ 0F85 1082FFFF
7C83E463 ^E9 D65FFFFF
                          JMP ntd11.7083443E
                                                                                                              0 0 Lasterr ERROR SUCCESS (00000000)
7C835975=ntd11.7C835975
                                                                                                              EFL 00000246 (NO, NB, E, BE, NS, PE, GE, LE)
```

MS08-067 conclusions

Wrong path validation leads to overflow.

 Easy to exploit vulnerability without any protections.

 Even on NX-enabled versions also exploitable.



MS09-050 general info

- Found by Laurent Gaffié in 3 seconds of fuzzing SMBv2. Reliable exploit by Piotr Bania, first MSF version by Stephen Fewer.
- EAX register is initialized with a word from the SMB2 packet. In the next instruction our controlled value is used as an array index with one safety check whether the value is NULL. Accidentally there is no check if array index exceeds the number of elements in the array. Later, the location pointed by our controlled value is executed using the call instruction.
- Kudos:

.http://g-laurent.blogspot.be/2009/09/windows-vista7-smb20-negotiate-protocol.html

.http://g-laurent.blogspot.be/2009/10/more-explication-on-cve-2009-3103.html

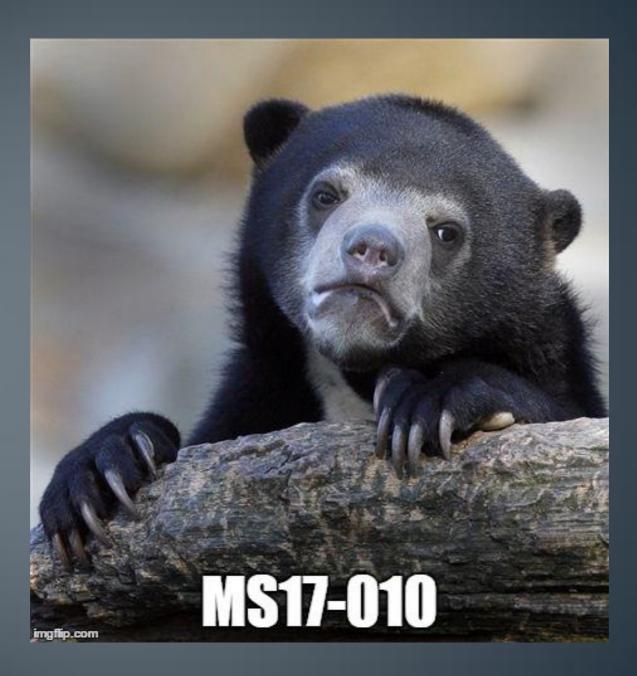
.http://piotrbania.com/all/articles/smb2_351_packets.pdf

MS09-050 exploitation

```
kd> g
Breakpoint 3 hit
srv2!SrvSnapShotScavengerTimer:
8e32eea0 6a01
                          push
kd> q
Breakpoint 2 hit
srv2!SrvProcCompleteRequest+0xd2:
8e33bb91 ffd0
                          call
                                  eax
kd> r eax
eax=ffd00d09
kd> q
Breakpoint 0 hit
ffd00d09 5e
                                  esi
                          pop
kd> u
ffd00d09 5e
                          pop
                                  esi
ffd00d0a c3
                          ret
ffd00d0b 7c58
                                  ffd00d65
                          il
ffd00d0d 0fb75102
                                  edx, word ptr [ecx+2]
                          MOVZX
ffd00d11 0fb77002
                          MOVZX
                                  esi, word ptr [eax+2]
ffd00d15 663bd6
                                  dx.si
                          CMP
                                  ffd00d07
ffd00d18 7fed
                          19
ffd00d1a 7c49
                          il
                                  ffd00d65
kd> p
ffd00d0a c3
                          ret
kd> p
84910bf8 00544d42
                          add
                                  byte ptr [ebp+ecx*2+42h],dl
kd> q
```

MS09-050 conclusions

- Quick to find, fuzzer code from Laurent is straight-forward.
- Requires efforts to make exploitation reliable: Piotr's trampoline vs hard-coded addresses.
- ASLR were bypassed as address from kernel HAL memory region was used.
- NX HAL heap and NonPagedPoolNx from Windows 8



MS17-010 general info

- Released by Shadow Brokers, developed by Equation Group.
- Complicated exploit requires several steps to trigger vulnerability.

Kudos:

http://risksense.com/_api/filesystem/466/EternalBlue_RiskSense-Exploit-Analysis-and-Port-to-Microsoft-Windows-10_v1_2.pdf

http://blog.trendmicro.com/trendlabs-security-intelligence/ms17-010-eternalblue/

http://blogs.360.cn/360safe/2017/04/17/nsa-eternalblue-smb/

https://github.com/worawit/MS17-010

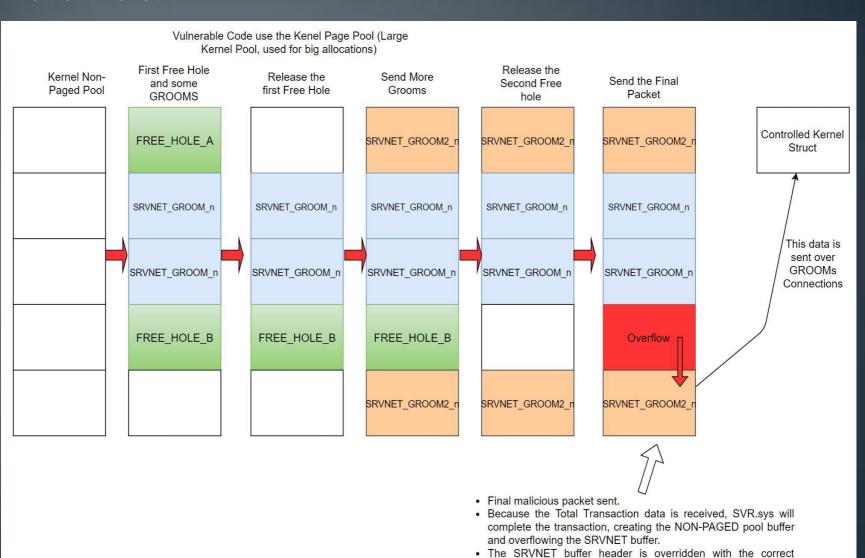
https://packetstormsecurity.com/files/142548/MS17-010-EternalBlue-SMB-Remote-Windows-Kernel-Pool-Corruption.html

MS17-010 general concept

- There is a buffer overflow before second memmove/memcpy operation in Srv!SrvOs2FeaToNt. The size is calculated in Srv!SrvOs2FeaListSizeToNt, with mathematical error where a DWORD is subtracted into a WORD.
- Size of NTFEA list requires Large Kernel Pool Allocation in SRV.sys
- Last iteration overwrite next memory area (SRVNET.sys). To achieve it both buffers should be aligned.
- Kernel Pool should be sprayed.
- Server returns 0xC00000D (status_invalid_parameter).
- Exploit may require several tries.

MS17-010 exploitation

From TrendMicro:



values

structure.

KNOWN kernel structure.

Now the SRVNET buffer is pointing to the KNOWN kernel

· Next, the data sent to the SRVNET buffer will override the

MS17-010 exploitation (2)

```
Command - Kernel 'com:pipe,resets=0,reconnect,port=\\.\pipe\kd Windows Server 2008 R2 x64' - WinDbg:6.12.0002.633 AMD64
1: kd> bp srv!SrvOs2FeaListToNt+0x108
1: kd> q
Breakpoint 4 hit
srv!SrvOs2FeaListToNt+0x108:
fffff880`03ec2a68 bb0d0000c0
                                            ebx,0C000000Dh
                                    MOV
1: kd> p
srv!SrvOs2FeaListToNt+0x10d:
fffff880°03ec2a6d ebe5
                                            srv!SrvOs2FeaListToNt+0xf4 (ffffff880`03ec2a54)
                                    jmp
1: kd> g
Break instruction exception - code 80000003 (first chance)
fffffffffffd00201 cc
                                    int
0: kd > dd
ffffffffffffd00201 41c931cc b9c301e2 c0000082 bb48320f
ffffffff ffd00211 ffd00ff8 ffffffff 89045389 058d4803
ffffffffffffd00221 0000000a 48c28948 0f20eac1 010fc330
ffffffffffffd00231 894865f8 00102524 48650000 a825248b
   fffff ffd00241 50000001 56525153 50415557 52415141
fffffffffffffd00251 54415341 56415541 2b6a5741 2534ff65
fffffffffffffd00261 00000010 336a5341 d1894c51 08ec8348
ffffffffffffd00271 ec814855 00000158 24ac8d48 00000080
```

MS17-010 conclusions

- Sophisticated exploit which requires several steps to achieve reliable execution.
- Most of kernel structures used by exploit are undocumented.
- In-depth SMB protocol knowledge required.
- Redstone 1 update randomization for page table entries prevents the DEP bypass.
- Redstone 2 update (after the MS17-010 patch), the HAL heap is randomized, defeating the ASLR bypass.

Final thoughts

- Comparing these exploits we can see how Microsoft improved
 OS security in past few years.
- New exploit mitigation techniques make reliable exploitation much harder to achieve.
- On recently updated Windows 10 I don't expect that something similar will arise in future until new class of vulnerabilities would be introduced.
- Great work by Microsoft to introduce new mitigations based on vulnerability classes from previous OS releases.

Any questions?

Thank you

- Email: vinegrep@protonmail.com
- Twitter: @vinegrep