# Rábalabaxúrias [UFMG]

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	7.3 Automato de Sufixo	68	<pre>int centroid(int k, int last, int size) {    for (int i = 0; i &lt; (int) g[k].size(); i++) {</pre>	
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	7.5 Suffix Array Rafael	69	<pre>if (subsize[u] &gt; size / 2)     return centroid(u, k, size);</pre>	
	7.6 KMP	70	} // k eh o centroid return k;	
8	Extra	<b>7</b> 1	}	
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#### 1.2 Tarjan para Pontes

```
// Computa pontos de articulação
// e pontes
//
// O(n+m)
int in[MAX];
int low[MAX];
int parent[MAX];
vector < int > g[MAX];
bool is_art[MAX];
void dfs_art(int v, int p, int &d){
    parent[v] = p;
    low[v] = in[v] = d++;
    is_art[v] = false;
    for (int j : g[v]){
        if (j == p) continue;
        if (in[j] == -1){
            dfs_art(j, v, d);
            if (low[j] >= in[v]) is_art[v] = true;
            //if (low[j] > in[v]) this edge is a bridge
```

```
low[v] = min(low[v], low[j]);
        else low[v] = min(low[v], in[j]);
    }
    if (p == -1){
        is_art[v] = false;
        int k = 0;
        for (int j : g[v])
            k += (parent[j] == v);
        if (k > 1) is_art[v] = true;
    }
}
int d = 0;
memset(in, -1, sizeof in);
dfs_art(1, -1, d);
1.3 Dominator Tree - Kawakami
// Se vira pra usar ai
//
// build - O(n)
// dominates - O(1)
int n;
namespace DTree {
    vector < int > g[MAX];
    // The dominator tree
    vector<int> tree[MAX]:
    int dfs_l[MAX], dfs_r[MAX];
    // Auxiliary data
    vector < int > rg[MAX], bucket[MAX];
    int idom[MAX], sdom[MAX], prv[MAX], pre[MAX];
    int ancestor[MAX], label[MAX];
    vector<int> preorder;
    void dfs(int v) {
```

static int t = 0;
pre[v] = ++t;

```
sdom[v] = label[v] = v:
    preorder.push_back(v);
    for (int nxt: g[v]) {
        if (sdom[nxt] == -1) {
            prv[nxt] = v;
            dfs(nxt);
        rg[nxt].push_back(v);
    }
}
int eval(int v) {
    if (ancestor[v] == -1) return v;
    if (ancestor[v]] == -1) return label[v];
    int u = eval(ancestor[v]);
    if (pre[sdom[u]] < pre[sdom[label[v]]]) label[v] = u;</pre>
    ancestor[v] = ancestor[u]:
    return label[v]:
}
void dfs2(int v) {
    static int t = 0;
    dfs_1[v] = t++;
    for (int nxt: tree[v]) dfs2(nxt);
    dfs_r[v] = t++;
}
void build(int s) {
    for (int i = 0; i < n; i++) {</pre>
        sdom[i] = pre[i] = ancestor[i] = -1;
        rg[i].clear();
        tree[i].clear();
        bucket[i].clear();
    }
    preorder.clear();
    dfs(s):
    if (preorder.size() == 1) return;
    for (int i = int(preorder.size()) - 1; i >= 1; i--) {
        int w = preorder[i];
        for (int v: rg[w]) {
            int u = eval(v);
            if (pre[sdom[u]] < pre[sdom[w]]) sdom[w] =</pre>
               sdom[u];
        }
        bucket[sdom[w]].push_back(w);
```

```
ancestor[w] = prv[w];
            for (int v: bucket[prv[w]]) {
                int u = eval(v);
                idom[v] = (u == v) ? sdom[v] : u;
            bucket[prv[w]].clear();
        for (int i = 1; i < preorder.size(); i++) {</pre>
            int w = preorder[i];
            if (idom[w] != sdom[w]) idom[w] = idom[idom[w]];
            tree[idom[w]].push_back(w);
        idom[s] = sdom[s] = -1;
        dfs2(s);
    }
    // Whether every path from s to v passes through u
    bool dominates(int u, int v) {
        if (pre[v] == -1) return 1; // vacuously true
        return dfs_1[u] <= dfs_1[v] && dfs_r[v] <= dfs_r[u];</pre>
    }
};
1.4 Kosaraju
// O(n + m)
int n;
vector < vector < int > > g(MAX);
vector < vector < int > > gi(MAX); // grafo invertido
int vis[MAX]:
stack<int> S:
int comp[MAX]; // componente conexo de cada vertice
void dfs(int k) {
    vis[k] = 1;
    for (int i = 0; i < (int) g[k].size(); i++)</pre>
        if (!vis[g[k][i]]) dfs(g[k][i]);
```

}

S.push(k);

```
void scc(int k, int c) {
    vis[k] = 1;
    comp[k] = c;
    for (int i = 0; i < (int) gi[k].size(); i++)</pre>
        if (!vis[gi[k][i]]) scc(gi[k][i], c);
}
void kosaraju() {
    for (int i = 0; i < n; i++) vis[i] = 0;</pre>
    for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
    for (int i = 0; i < n; i++) vis[i] = 0;
    while (S.size()) {
        int u = S.top();
        S.pop();
        if (!vis[u]) scc(u, u);
    }
}
    Dijkstra
1.5
// encontra menor distancia de a
// para todos os vertices
// se ao final do algoritmo d[i] = INF,
// entao a nao alcanca i
//
// O(m log(n))
int n;
vector < vector < int > > g(MAX);
vector < vector < int > > w(MAX); // peso das arestas
int d[MAX]:
void dijsktra(int a) {
    for (int i = 0; i < n; i++) d[i] = INF;</pre>
    d[a] = 0;
    priority_queue < pair < int , int > > Q;
    Q.push(make_pair(0, a));
    while (Q.size()) {
        int u = Q.top().second, dist = -Q.top().first;
```

Q.pop();

```
if (dist > d[u]) continue;
        for (int i = 0; i < (int) g[u].size(); i++) {</pre>
            int v = g[u][i];
            if (d[v] > d[u] + w[u][i]) {
                d[v] = d[u] + w[u][i];
                Q.push(make_pair(-d[v], v));
            }
        }
    }
}
    Heavy-Light Decomposition sem Update
// query de min do caminho
//
// Complexidades:
// build - O(n)
// query_path - O(log(n))
#define f first
#define s second
namespace hld {
    vector < pair < int , int > > g[MAX];
    int in[MAX], sz[MAX];
    int sobe[MAX], pai[MAX];
    int h[MAX], v[MAX], t;
    int men[MAX], seg[2*MAX];
    void build_hld(int k, int p = -1, int f = 1) {
        v[in[k] = t++] = sobe[k]; sz[k] = 1;
        for (auto& i : g[k]) if (i.f != p) {
```

sobe[i.f] = i.s; pai[i.f] = k;

i.s);

}

h[i.f] = (i == g[k][0] ? h[k] : i.f);

build\_hld(i.f, k, f); sz[k] += sz[i.f];

if (p\*f == -1) build\_hld(h[k] = k, -1, t = 0);

men[i.f] = (i == g[k][0] ? min(men[k], i.s) :

if (sz[i.f] > sz[g[k][0].f]) swap(i, g[k][0]);

```
}
    void build(int root = 0) {
        t = 0;
        build_hld(root);
        for (int i = 0; i < t; i++) seg[i+t] = v[i];</pre>
        for (int i = t-1; i; i--) seg[i] = min(seg[2*i],
           seg[2*i+1]);
    }
    int query_path(int a, int b) {
        if (a == b) return INF;
        if (in[a] < in[b]) swap(a, b);</pre>
        if (h[a] != h[b]) return min(men[a],
           query_path(pai[h[a]], b));
        int ans = INF, x = in[b]+1+t, y = in[a]+t;
        for (; x \le y; ++x/=2, --y/=2) ans = min({ans,
           seg[x], seg[y]);
        return ans;
    }
};
```

#### 1.7 Heavy-Light Decomposition - vertice

```
// SegTree de soma
// query / update de soma dos vertices
// Complexidades:
// build - O(n)
// \text{ query_path - O(log^2 (n))}
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
namespace seg {
    11 seg[4*MAX], lazy[4*MAX];
    int n, *v;
    ll build(int p=1, int l=0, int r=n-1) {
        lazy[p] = 0;
        if (1 == r) return seg[p] = v[1];
        int m = (1+r)/2;
        return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1,
```

```
r):
    }
    void build(int n2, int* v2) {
        n = n2, v = v2;
        build();
    }
    void prop(int p, int l, int r) {
        seg[p] += lazy[p]*(r-l+1);
        if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
           lazy[p];
        lazy[p] = 0;
    }
    11 query(int a, int b, int p=1, int l=0, int r=n-1) {
        prop(p, 1, r);
        if (a <= l and r <= b) return seg[p];</pre>
        if (b < 1 \text{ or } r < a) \text{ return } 0;
        int m = (1+r)/2:
        return query (a, b, 2*p, 1, m) + query (a, b, 2*p+1, m)
           m+1, r);
    }
    ll update(int a, int b, int x, int p=1, int l=0, int
       r=n-1) {
        prop(p, 1, r);
        if (a <= 1 and r <= b) {</pre>
            lazy[p] += x;
            prop(p, 1, r);
            return seg[p];
        if (b < 1 or r < a) return seg[p];</pre>
        int m = (1+r)/2;
        return seg[p] = update(a, b, x, 2*p, 1, m) +
             update(a, b, x, 2*p+1, m+1, r);
    }
};
namespace hld {
    vector < int > g[MAX];
    int in[MAX], out[MAX], sz[MAX];
    int peso[MAX], pai[MAX];
    int h[MAX], v[MAX], t;
    void build_hld(int k, int p = -1, int f = 1) {
```

```
v[in[k] = t++] = peso[k]; sz[k] = 1;
    for (auto& i : g[k]) if (i != p) {
        pai[i] = k;
       h[i] = (i == g[k][0] ? h[k] : i);
        build_hld(i, k, f); sz[k] += sz[i];
        if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
    }
    out[k] = t;
    if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
void build(int root = 0) {
    t = 0:
    build_hld(root);
    seg::build(t, v);
ll query_path(int a, int b) {
    if (a == b) return seg::query(in[a], in[a]);
    if (in[a] < in[b]) swap(a, b);</pre>
    if (h[a] == h[b]) return seg::query(in[b], in[a]);
    return seg::query(in[h[a]], in[a]) +
       query_path(pai[h[a]], b);
void update_path(int a, int b, int x) {
    if (a == b) return (void) seg::update(in[a], in[a],
       x);
    if (in[a] < in[b]) swap(a, b);</pre>
    if (h[a] == h[b]) return (void)seg::update(in[b],
       in[a], x);
    seg::update(in[h[a]], in[a], x);
       update_path(pai[h[a]], b, x);
}
11 query_subtree(int a) {
    if (in[a] == out[a]-1) return seg::query(in[a],
       in[a]):
    return seg::query(in[a], out[a]-1);
void update_subtree(int a, int x) {
    if (in[a] == out[a]-1) return
       (void)seg::update(in[a], in[a], x);
```

```
seg::update(in[a], out[a]-1, x);
    }
    int lca(int a, int b) {
        if (in[a] < in[b]) swap(a, b);</pre>
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
    }
};
1.8 LCA com HLD
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// Para buildar pasta chamar build(root)
// Complexidades:
// build - O(n)
// lca - O(log(n))
vector < vector < int > > g(MAX);
int in[MAX], h[MAX], sz[MAX];
int pai[MAX], t;
void build(int k, int p = -1, int f = 1) {
    in[k] = t++; sz[k] = 1;
    for (int& i : g[k]) if (i != p) {
        pai[i] = k;
        h[i] = (i == g[k][0] ? h[k] : i);
        build(i, k, f); sz[k] += sz[i];
        if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
    }
    if (p*f == -1) t = 0, h[k] = k, build(k, -1, 0);
}
int lca(int a, int b) {
    if (in[a] < in[b]) swap(a, b);</pre>
    return h[a] == h[b] ? b : lca(pai[h[a]], b);
}
```

#### 1.9 LCA com RMQ

```
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// Complexidades:
// build - O(n) + build_RMQ
// lca - RMQ
int n;
vector < vector < int > > g(MAX);
                // pos[i] : posicao de i em v (primeira
int pos[MAX];
   aparicao
int ord[2 * MAX]; // ord[i] : i-esimo vertice na ordem de
   visitação da dfs
int v[2 * MAX]; // vetor de alturas que eh usado na RMQ
int p;
void dfs(int k, int l) {
    ord[p] = k;
    pos[k] = p;
    v[p++] = 1;
    for (int i = 0; i < (int) g[k].size(); i++)</pre>
        if (pos[g[k][i]] == -1) {
            dfs(g[k][i], l + 1);
            ord[p] = k;
            v[p++] = 1;
        }
}
void build(int root) {
    for (int i = 0; i < n; i++) pos[i] = -1;
    p = 0;
    dfs(root, 0);
    build_RMQ();
}
int lca(int u, int v) {
    int a = pos[u], b = pos[v];
    if (a > b) swap(a, b);
    return ord[RMQ(a, b)];
}
```

#### 1.10 LCA com binary lifting

```
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
//
// Complexidades:
// build - O(n log(n))
// lca - O(log(n))
vector < vector < int > > g(MAX);
int n, p;
int pai[MAX2][MAX];
int in[MAX], out[MAX];
void dfs(int k) {
    in[k] = p++;
    for (int i = 0; i < (int) g[k].size(); i++)</pre>
        if (in[g[k][i]] == -1) {
            pai[0][g[k][i]] = k;
            dfs(g[k][i]);
    out[k] = p++;
}
void build(int raiz) {
    for (int i = 0; i < n; i++) pai[0][i] = i;</pre>
    p = 0, memset(in, -1, sizeof in);
    dfs(raiz);
    // pd dos pais
    for (int k = 1; k < MAX2; k++) for (int i = 0; i < n;
       i++)
        pai[k][i] = pai[k - 1][pai[k - 1][i]];
}
bool anc(int a, int b) { // se a eh ancestral de b
    return in[a] <= in[b] and out[a] >= out[b];
}
int lca(int a, int b) {
    if (anc(a, b)) return a;
```

```
if (anc(b, a)) return b;

// sobe a
for (int k = MAX2 - 1; k >= 0; k--)
    if (!anc(pai[k][a], b)) a = pai[k][a];

return pai[0][a];
}
```

#### 1.11 Heavy-Light Decomposition - aresta

```
// SegTree de soma
// query / update de soma das arestas
// Complexidades:
// build - O(n)
// query_path - O(log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
#define f first
#define s second
namespace seg {
    11 \text{ seg}[4*MAX], lazy[4*MAX];
    int n, *v;
    ll build(int p=1, int l=0, int r=n-1) {
        lazy[p] = 0;
        if (1 == r) return seg[p] = v[1];
        int m = (1+r)/2;
        return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1,
           r);
    }
    void build(int n2, int* v2) {
        n = n2, v = v2;
        build();
    }
    void prop(int p, int l, int r) {
        seg[p] += lazy[p]*(r-l+1);
        if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
```

```
lazy[p];
        lazy[p] = 0;
    ll query(int a, int b, int p=1, int l=0, int r=n-1) {
        prop(p, l, r);
        if (a <= 1 and r <= b) return seg[p];</pre>
        if (b < 1 or r < a) return 0;
        int m = (1+r)/2;
        return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1,
           m+1, r);
    }
    ll update(int a, int b, int x, int p=1, int l=0, int
       r=n-1) {
        prop(p, l, r);
        if (a <= 1 and r <= b) {</pre>
            lazy[p] += x;
            prop(p, l, r);
            return seg[p];
        if (b < l or r < a) return seg[p];</pre>
        int m = (1+r)/2;
        return seg[p] = update(a, b, x, 2*p, 1, m) +
            update(a, b, x, 2*p+1, m+1, r);
   }
};
namespace hld {
   vector<pair<int, int> > g[MAX];
    int in[MAX], out[MAX], sz[MAX];
    int sobe[MAX], pai[MAX];
    int h[MAX], v[MAX], t;
    void build_hld(int k, int p = -1, int f = 1) {
        v[in[k] = t++] = sobe[k]; sz[k] = 1;
        for (auto& i : g[k]) if (i.f != p) {
            sobe[i.f] = i.s; pai[i.f] = k;
            h[i.f] = (i == g[k][0] ? h[k] : i.f);
            build_hld(i.f, k, f); sz[k] += sz[i.f];
            if (sz[i.f] > sz[g[k][0].f]) swap(i, g[k][0]);
        }
        out[k] = t;
```

```
if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
    }
    void build(int root = 0) {
        t = 0;
        build_hld(root);
                                                                  //
        seg::build(t, v);
    }
    ll query_path(int a, int b) {
        if (a == b) return 0;
        if (in[a] < in[b]) swap(a, b);</pre>
        if (h[a] == h[b]) return seg::query(in[b]+1, in[a]);
        return seg::query(in[h[a]], in[a]) +
           query_path(pai[h[a]], b);
    void update_path(int a, int b, int x) {
        if (a == b) return;
        if (in[a] < in[b]) swap(a, b);</pre>
        if (h[a] == h[b]) return (void)seg::update(in[b]+1,
           in[a], x);
        seg::update(in[h[a]], in[a], x);
           update_path(pai[h[a]], b, x);
    11 query_subtree(int a) {
        if (in[a] == out[a]-1) return 0;
        return seg::query(in[a]+1, out[a]-1);
    void update_subtree(int a, int x) {
        if (in[a] == out[a]-1) return;
        seg::update(in[a]+1, out[a]-1, x);
    int lca(int a, int b) {
        if (in[a] < in[b]) swap(a, b);</pre>
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
    }
     Bellman-Ford
// Calcula a menor distancia
```

#### 1.12

};

```
// entre a e todos os vertices e
```

```
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
// O(nm)
int n, m;
int d[MAX];
vector<pair<int, int> > ar; // vetor de arestas
vector < int > w;
                            // peso das arestas
bool bellman_ford(int a) {
    for (int i = 0; i < n; i++) d[i] = INF;</pre>
    d[a] = 0:
    for (int i = 0; i <= n; i++)
        for (int j = 0; j < m; j++) {</pre>
            if (d[ar[j].second] > d[ar[j].first] + w[j]) {
                if (i == n) return 1;
                d[ar[j].second] = d[ar[j].first] + w[j];
            }
        }
    return 0;
      Tarjan para SCC
1.13
// O(n + m)
vector < vector < int > > g(MAX);
stack<int> s;
int vis[MAX], comp[MAX];
int id[MAX], p;
int dfs(int k) {
    int lo = id[k] = p++;
    s.push(k);
    vis[k] = 2; // ta na pilha
```

```
// calcula o menor cara q ele alcanca
    // que ainda nao esta em um scc
    for (int i = 0; i < g[k].size(); i++) {</pre>
        if (!vis[g[k][i]])
            lo = min(lo, dfs(g[k][i]));
        else if (vis[g[k][i]] == 2)
            lo = min(lo, id[g[k][i]]);
    }
    // nao alcanca ninguem menor -> comeca scc
    if (lo == id[k]) while (1) {
        int u = s.top();
        s.pop(); vis[u] = 1;
        comp[u] = k;
        if (u == k) break;
    }
    return lo;
}
void tarjan() {
    memset(vis, 0, sizeof(vis));
    p = 0;
    for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
}
1.14 Dinic (Dilson)
// O(n^2 m)
// Grafo bipartido -> O(sqrt(n)*m)
template <class T> struct dinic {
    struct edge {
        int v, rev;
        edge(int v_, T cap_, int rev_) : v(v_), cap(cap_),
           rev(rev ) {}
    };
    vector < vector < edge >> g;
    vector < int > level;
```

```
queue < int > q;
T flow;
int n;
dinic(int n_) : g(n_), level(n_), n(n_) {}
void add_edge(int u, int v, T cap) {
    if (u == v)
        return;
    edge e(v, cap, int(g[v].size()));
    edge r(u, 0, int(g[u].size()));
    g[u].push_back(e);
    g[v].push_back(r);
}
bool build_level_graph(int src, int sink) {
    fill(level.begin(), level.end(), -1);
    while (not q.empty())
        q.pop();
    level[src] = 0;
    q.push(src);
    while (not q.empty()) {
        int u = q.front();
        q.pop();
        for (auto e = g[u].begin(); e != g[u].end();
           ++e) {
            if (not e->cap or level[e->v] != -1)
                continue:
            level[e->v] = level[u] + 1;
            if (e->v == sink)
                return true;
            q.push(e->v);
        }
    }
    return false;
}
T blocking_flow(int u, int sink, T f) {
    if (u == sink or not f)
        return f:
    T fu = f;
    for (auto e = g[u].begin(); e != g[u].end(); ++e) {
        if (not e->cap or level[e->v] != level[u] + 1)
            continue:
```

```
T mincap = blocking_flow(e->v, sink, min(fu,
               e->cap));
            if (mincap) {
                g[e->v][e->rev].cap += mincap;
                e->cap -= mincap;
                fu -= mincap;
            }
        }
        if (f == fu)
            level[u] = -1;
        return f - fu;
    T max_flow(int src, int sink) {
        flow = 0;
        while (build_level_graph(src, sink))
            flow += blocking_flow(src, sink,
               numeric_limits <T>::max());
        return flow;
   }
};
```

#### 1.15 Centro da Arvore

```
// Centro eh o vertice que minimiza
// a maior distancia dele pra alguem
// O centro fica no meio do diametro
// A funcao center retorna um par com
// o diametro e o centro
//
// O(n+m)
vector < vector < int > > g(MAX);
int n, vis[MAX];
int d[2][MAX];
// retorna ultimo vertice visitado
int bfs(int k, int x) {
        queue < int > q; q.push(k);
    memset(vis, 0, sizeof(vis));
    vis[k] = 1;
    d[x][k] = 0;
    int last = k;
```

```
while (q.size()) {
       int u = q.front(); q.pop();
        last = u;
        for (int i : g[u]) if (!vis[i]) {
            vis[i] = 1;
            q.push(i);
            d[x][i] = d[x][u] + 1;
        }
    }
    return last;
}
pair<int, int> center() {
    int a = bfs(0, 0);
    int b = bfs(a, 1);
    bfs(b, 0);
    int c, mi = INF;
    for (int i = 0; i < n; i++) if (max(d[0][i], d[1][i]) <
        mi = max(d[0][i], d[1][i]), c = i;
    return {d[0][a], c};
1.16 Sack (DSU em arvores)
// Responde queries de todas as sub-arvores
// offline
//
// O(n log(n))
int sz[MAX], cor[MAX], cnt[MAX];
vector < vector < int > > g(MAX);
void build(int k, int d=0) {
    sz[k] = 1;
    for (auto& i : g[k]) {
        build(i, d+1); sz[k] += sz[i];
       if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
    }
}
```

```
void compute(int k, int x, bool dont=1) {
    cnt[cor[k]] += x;
    for (int i = dont; i < g[k].size(); i++)</pre>
        compute(g[k][i], x, 0);
}
void solve(int k, bool keep=0) {
    for (int i = int(g[k].size())-1; i >= 0; i--)
        solve(g[k][i], !i);
    compute(k, 1);
        // agora cnt[i] tem quantas vezes a cor
        // i aparece na sub-arvore do k
    if (!keep) compute(k, -1, 0);
}
1.17 Centroid
// Computa os 2 centroids da arvore
//
// O(n)
int n, subsize[MAX];
vector < vector < int > > g(MAX);
void dfs(int k, int p=-1) {
    subsize[k] = 1;
    for (int i : g[k]) if (i != p) {
        dfs(i, k);
        subsize[k] += subsize[i];
}
int centroid(int k, int p=-1, int size=-1) {
    if (size == -1) size = subsize[k];
    for (int i : g[k]) if (i != p) if (subsize[i] > size/2)
        return centroid(i, k, size, t);
    return k;
}
pair < int , int > centroids(int k=0) {
```

```
dfs(k):
   int i = centroid(k), i2 = i;
    for (int j : g[i]) if (2*subsize[j] == subsize[k]) i2 =
       j;
   return {i, i2};
}
1.18 Kruskal
// Gera AGM a partir do vetor de arestas
// O(m log(n))
int n;
vector < pair < int , pair < int , int > > ar; // vetor de arestas
int v[MAX];
// Union-Find em O(log(n))
void build();
int find(int k);
void une(int a, int b);
void kruskal() {
    build();
    sort(ar.begin(), ar.end());
   for (int i = 0; i < (int) ar.size(); i++) {</pre>
        int a = ar[i].s.f, b = ar[i].s.s;
        if (find(a) != find(b)) {
            une(a, b);
            // aresta faz parte da AGM
       }
   }
}
1.19 Blossom - matching maximo em grafo geral
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
```

```
vector < vector < int > > g(MAX);
int match[MAX]; // match[i] = com quem i esta matchzado ou -1
int n, pai[MAX], base[MAX], vis[MAX];
queue < int > q;
void contract(int u, int v, bool first = 1) {
    static vector < bool > bloss;
    static int 1;
    if (first) {
        bloss = vector < bool > (n, 0);
        vector < bool > teve(n, 0);
        int k = u; l = v;
        while (1) {
            teve[k = base[k]] = 1;
            if (match[k] == -1) break;
            k = pai[match[k]];
        }
        while (!teve[l = base[l]]) l = pai[match[l]];
    while (base[u] != 1) {
        bloss[base[u]] = bloss[base[match[u]]] = 1;
        pai[u] = v;
        v = match[u];
        u = pai[match[u]];
    }
    if (!first) return;
    contract(v, u, 0);
    for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
        base[i] = 1;
        if (!vis[i]) q.push(i);
        vis[i] = 1;
    }
}
int getpath(int s) {
    for (int i = 0; i < n; i++) base[i] = i, pai[i] = -1,
       vis[i] = 0;
    vis[s] = 1; q = queue < int > (); q.push(s);
    while (q.size()) {
        int u = q.front(); q.pop();
        for (int i : g[u]) {
            if (base[i] == base[u] or match[u] == i)
```

```
continue:
            if (i == s or (match[i] != -1 and pai[match[i]]
                ! = -1))
                contract(u, i);
            else if (pai[i] == -1) {
                pai[i] = u;
                if (match[i] == -1) return i;
                i = match[i];
                vis[i] = 1; q.push(i);
            }
        }
    }
    return -1;
}
int blossom() {
    int ans = 0:
    memset(match, -1, sizeof(match));
   for (int i = 0; i < n; i++) if (match[i] == -1)</pre>
        for (int j : g[i]) if (match[j] == -1) {
            match[i] = i;
            match[j] = i;
            ans++;
            break;
        }
    for (int i = 0; i < n; i++) if (match[i] == -1) {
        int j = getpath(i);
        if (j == -1) continue;
        ans++;
        while (j != -1) {
            int p = pai[j], pp = match[p];
            match[p] = j;
            match[j] = p;
            j = pp;
        }
    }
    return ans;
}
1.20 Max flow com lower bound
```

// Manda passar pelo menos 'lb' de fluxo

```
// em cada aresta
// O(dinic)
struct lb_max_flow : dinic {
    vector<int> d;
    vector<int> e;
    lb_max_flow(int n):dinic(n + 2), d(n, 0){}
    void add(int a, int b, int c, int lb = 0){
        c = 1b;
        d[a] -= lb;
        d[b] += 1b;
        dinic::add(a, b, c);
    }
    bool check_flow(int src, int snk, int F){
        int n = d.size();
        d[src] += F;
        d[snk] -= F;
        for (int i = 0; i < n; i++){
            if (d[i] > 0){
                dinic::add(n, i, d[i]);
            } else if (d[i] < 0){</pre>
                dinic::add(i, n+1, -d[i]);
            }
        }
        int f = max_flow(n, n+1);
        return (f == F);
};
1.21 Isomorfismo de Arvores
// Duas arvores T1 e T2 sao isomorfas
// sse T1.getHash() = T2.getHash()
// O(n log(n))
map < vector < int > , int > mapp;
```

struct tree {

```
int n:
vector < vector < int > > g;
vector<int> subsize;
tree(int n) {
    g.resize(n);
    subsize.resize(n);
}
void dfs(int k, int p=-1) {
    subsize[k] = 1;
    for (int i : g[k]) if (i != p) {
        dfs(i, k);
        subsize[k] += subsize[i];
    }
}
int centroid(int k, int p=-1, int size=-1) {
    if (size == -1) size = subsize[k];
    for (int i : g[k]) if (i != p)
        if (subsize[i] > size/2)
            return centroid(i, k, size);
    return k;
}
pair<int, int> centroids(int k=0) {
    dfs(k);
    int i = centroid(k), i2 = i;
    for (int j : g[i]) if (2*subsize[j] == subsize[k])
       i2 = j;
    return {i, i2};
}
int hashh(int k, int p=-1) {
    vector < int > v;
    for (int i : g[k]) if (i != p) v.push_back(hashh(i,
       k));
    sort(v.begin(), v.end());
    if (!mapp.count(v)) mapp[v] = int(mapp.size());
    return mapp[v];
}
11 getHash(int k=0) {
    pair < int , int > c = centroids(k);
    11 a = hashh(c.first), b = hashh(c.second);
    if (a > b) swap(a, b);
    return (a << 30) + b;
```

```
};
```

#### 1.22 MinCostMaxFlow Papa

```
/*
  s e t pre-definidos como MAXN - 1 e MAXN - 2.
   minCostFlow(f) computa o par (fluxo, custo) com
      max(fluxo) <= f que tenha min(custo).</pre>
   minCostFlow(INF) -> Fluxo maximo de custo minimo.
  Se tomar TLE, aleatorizar a ordem dos vertices no SPFA
*/
const int MAXN = 230;
template < typename T > struct MCMF {
    struct edge {
        int to, rev, flow, cap, residual;
        T cost;
        edge() { to = 0; rev = 0; flow = 0; cap = 0; cost =
           0; residual = false; }
        edge(int _to, int _rev, int _flow, int _cap, T
           _cost, bool _residual) {
            to = _to; rev = _rev;
            flow = _flow; cap = _cap;
            cost = _cost;
            residual = _residual;
        }
    };
    int s = MAXN - 1, t = MAXN - 2;
    vector < edge > G[MAXN];
    void addEdge(int u, int v, int w, T cost) {
        edge t = edge(v, G[v].size(), 0, w, cost, false);
        edge r = edge(u, G[u].size(), 0, 0, -cost, true);
        G[u].push_back(t);
        G[v].push_back(r);
```

```
deque < int > Q;
bool is_inside[MAXN];
int par_idx[MAXN], par[MAXN];
T dist[MAXN];
bool spfa() {
    for(int i = 0; i < MAXN; i++)</pre>
        dist[i] = INF;
    dist[t] = INF;
    Q.clear();
    dist[s] = 0;
    is_inside[s] = true;
    Q.push_back(s);
    while(!Q.empty()) {
        int u = Q.front();
        is_inside[u] = false;
        Q.pop_front();
        for (int i = 0; i < (int)G[u].size(); i++)
            if(G[u][i].cap > G[u][i].flow && dist[u] +
                G[u][i].cost < dist[G[u][i].to]) {</pre>
                 dist[G[u][i].to] = dist[u] +
                    G[u][i].cost;
                 par_idx[G[u][i].to] = i;
                 par[G[u][i].to] = u;
                 if(is_inside[G[u][i].to]) continue;
                 if(!Q.empty() && dist[G[u][i].to] >
                    dist[Q.front()])
                    Q.push_back(G[u][i].to);
                 else Q.push_front(G[u][i].to);
                 is_inside[G[u][i].to] = true;
            }
    }
    return dist[t] != INF;
}
pair < int , T > minCostFlow(int flow) {
    int f = 0:
```

```
T ret = 0:
    while(f <= flow && spfa()) {</pre>
        int mn_flow = flow - f, u = t;
        while(u != s){
            mn_flow = min(mn_flow,
               G[par[u]][par_idx[u]].cap -
               G[par[u]][par_idx[u]].flow);
            u = par[u];
        }
        u = t;
        while(u != s) {
            G[par[u]][par_idx[u]].flow += mn_flow;
            G[u][G[par[u]][par_idx[u]].rev].flow -=
            ret += G[par[u]][par_idx[u]].cost * mn_flow;
            u = par[u];
        }
        f += mn_flow;
    }
    return make_pair(f, ret);
}
/*
   Opcional.
   Retorna todas as arestas originais por onde passa
      fluxo = capacidade.
 */
vector<pair<int ,int > > recover() {
    vector<pair<int, int > > used;
    for(int i = 0; i < MAXN; i++)</pre>
        for(edge e : G[i])
            if(e.flow == e.cap && !e.residual)
                used.push_back({i, e.to});
    return used;
}
```

#### Dinic (Bruno)

};

```
// O(n^2 m)
// Grafo com capacidades 1 -> O(sqrt(n)*m)
struct dinic{
    struct edge {
        int p, c, id; // para, capacidade, id
        edge(int p_, int c_, int id_) : p(p_), c(c_),
           id(id_) {}
    };
    vector < vector < edge >> g;
    vector < int > lev;
    dinic(int n): g(n){}
    void add(int a, int b, int c) { // de a pra b com cap. c
        g[a].pb(edge(b, c, g[b].size()));
        g[b].pb(edge(a, 0, g[a].size()-1));
    }
    bool bfs(int s, int t) {
        lev = vector\langle int \rangle(g.size(), -1); lev[s] = 0;
        queue < int > q; q.push(s);
        while (q.size()) {
            int u = q.front(); q.pop();
            for (auto& i : g[u]) {
                if (lev[i.p] != -1 or !i.c) continue;
                lev[i.p] = lev[u] + 1;
                if (i.p == t) return 1;
                q.push(i.p);
            }
        return 0;
    }
    int dfs(int v, int s, int f = INF){
        if (v == s) return f;
        int tem = f;
        for (auto& i : g[v]) {
            if (lev[i.p] != lev[v] + 1 or !i.c) continue;
            int foi = dfs(i.p, s, min(tem, i.c));
            tem -= foi, i.c -= foi, g[i.p][i.id].c += foi;
        if (f == tem) lev[v] = -1;
```

```
return f - tem;
    }
    int max_flow(int s, int t) {
        int f = 0;
        while (bfs(s, t)) f += dfs(s, t);
        return f;
};
1.24 Line Tree
// Reduz min-query em arvore para RMQ
// Se o grafo nao for uma arvore, as queries
// sao sobre a arvore geradora maxima
// Queries de minimo
//
// build - O(n log(n))
// query - O(log(n))
int n;
namespace linetree {
    int id[MAX], seg[2*MAX], pos[MAX];
    vector < int > v[MAX], val[MAX];
    vector<pair<int, pair<int, int> > ar;
    void add(int a, int b, int p) { ar.pb({p, {a, b}}); }
    void build() {
        sort(ar.rbegin(), ar.rend());
        for (int i = 0; i < n; i++) id[i] = i, v[i] = {i},</pre>
           val[i].clear():
        for (auto i : ar) {
            int a = id[i.second.first], b =
               id[i.second.second];
            if (a == b) continue;
            if (v[a].size() < v[b].size()) swap(a, b);</pre>
            for (auto j : v[b]) id[j] = a, v[a].push_back(j);
            val[a].push_back(i.first);
            for (auto j : val[b]) val[a].push_back(j);
            v[b].clear(), val[b].clear();
        }
```

```
for (int i = 0; i < n; i++) pos[v[id[0]][i]] = i;</pre>
        for (int i = n; i < 2*n-1; i++) seg[i] =
           val[id[0]][i-n];
        for (int i = n-1; i > 0; i--) seg[i] = min(seg[2*i],
            seg[2*i+1]);
    }
    int query(int a, int b) {
        a = pos[a]+n, b = pos[b]+n;
        if (a > b) swap(a, b);
        int ans = INF; b--;
        for (; a <= b; ++a/=2, --b/=2) ans = min({ans,
            seg[a], seg[b]});
        return ans;
    }
};
1.25 	 2-SAT
// Retorna se eh possivel atribuir valores
// Grafo tem que caber 2n vertices
// add(x, y) adiciona implicacao x -> y
// Para adicionar uma clausula (x ou y)
// chamar add(nao(x), y)
// Se x tem que ser verdadeiro, chamar add(nao(x), x)
// O tarjan deve computar o componente conexo
// de cada vertice em comp
//
// O(|V|+|E|)
vector < vector < int > > g(MAX);
int n:
int nao(int x) { return (x + n) \% (2*n); }
// x \rightarrow y = !x ou y
void add(int x, int y){
    g[x].pb(y);
    // contraposicao
    g[nao(y)].pb(nao(x));
}
bool doisSAT(){
```

```
tarjan();
for (int i = 0; i < m; i++)
      if (comp[i] == comp[nao(i)]) return 0;
return 1;
}</pre>
```

#### 1.26 Floyd-Warshall

```
// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta
// (i, j) de peso w, INF caso contrario
//
// O(n^3)
int n;
int d[MAX][MAX];
bool floyd_warshall() {
    for (int k = 0; k < n; k++)
    for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n; j++)
        d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
    for (int i = 0; i < n; i++)</pre>
        if (d[i][i] < 0) return 1;</pre>
    return 0;
}
```

#### 2 Matematica

#### 2.1 Ordem de elemento do grupo

```
// Calcula a ordem de a em Z_n
// O grupo Zn eh ciclico sse n =
// 1, 2, 4, p^k ou 2 p^k, p primo impar
// Retorna -1 se nao achar
```

```
// O(sqrt(n) log(n))
int tot(int n); // totiente em O(sqrt(n))
int expo(int a, int b, int m); // (a^b)%m em O(log(b))
// acha todos os divisores ordenados em O(sqrt(n))
vector<int> div(int n) {
    vector<int> ret1, ret2;
    for (int i = 1; i*i <= n; i++) if (n % i == 0) {
        ret1.pb(i);
        if (i*i != n) ret2.pb(n/i);
    }
    for (int i = ret2.size()-1; i+1; i--) ret1.pb(ret2[i]);
    return ret1;
}
int ordem(int a, int n) {
    vector < int > v = div(tot(n));
    for (int i : v) if (expo(a, i, n) == 1) return i;
    return -1;
}
2.2 Pollard's Rho Alg
// Usa o algoritmo de deteccao de ciclo de Brent
// A fatoracao nao sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
// Eh recomendado chamar srand(time(NULL)) na main
//
// Complexidades (considerando mul constante):
// rho - esperado O(n^{(1/4)}) no pior caso
// fact - esperado menos que O(n^{(1/4)} \log(n)) no pior caso
11 mdc(11 a, 11 b) { return !b ? a : mdc(b, a % b); }
11 mul(ll a, ll b, ll m) {
    return (a*b-ll(a*(long double)b/m+0.5)*m+m)%m;
}
```

```
11 exp(ll a, ll b, ll m) {
    if (!b) return 1;
    ll ans = exp(mul(a, a, m), b/2, m);
    return b%2 ? mul(a, ans, m) : ans;
}
bool prime(ll n) {
    if (n < 2) return 0;
    if (n <= 3) return 1;</pre>
    if (n % 2 == 0) return 0;
    11 d = n - 1;
    int r = 0;
    while (d \% 2 == 0) \{
        r++;
        d /= 2;
    }
    int a[9] = \{2, 3, 5, 7, 11, 13, 17, 19, 23\};
    for (int i = 0; i < 9; i++) {
        if (a[i] >= n) break;
        ll x = exp(a[i], d, n);
        if (x == 1 \text{ or } x == n - 1) \text{ continue};
        bool deu = 1;
        for (int j = 0; j < r - 1; j++) {
            x = mul(x, x, n);
            if (x == n - 1) {
                 deu = 0;
                 break;
            }
        if (deu) return 0;
    }
    return 1;
}
11 \text{ rho}(11 \text{ n}) 
    if (n == 1 or prime(n)) return n;
    if (n % 2 == 0) return 2;
    while (1) {
```

```
11 x = 2, y = 2;
        11 \ ciclo = 2, i = 0;
        ll c = (rand() / (double) RAND_MAX) * (n - 1) + 1;
        11 d = 1;
        while (d == 1) {
            if (++i == ciclo) ciclo *= 2, y = x;
            x = (mul(x, x, n) + c) % n;
            if (x == y) break;
            d = mdc(abs(x - y), n);
        }
        if (x != y) return d;
    }
}
void fact(ll n, vector<ll>& v) {
    if (n == 1) return;
    if (prime(n)) v.pb(n);
    else {
        11 d = rho(n);
        fact(d, v);
        fact(n / d, v);
    }
}
    Algoritmo de Euclides extendido
// acha x e y tal que ax + by = mdc(a, b)
//
// O(log(min(a, b)))
int mdce(int a, int b, int *x, int *y){
    if(!a){
        *x = 0;
        *y = 1;
        return b;
    }
```

```
int X, Y;
int mdc = mdce(b % a, a, &X, &Y);
*x = Y - (b / a) * X;
*y = X;

return mdc;
```

#### 2.4 Algoritmo de Euclides

}

```
// O(log(min(a, b)))
int mdc(int a, int b) {
    return !b ? a : mdc(b, a % b);
}
```

#### 2.5 Divisão de Polinomios

```
// Divide p1 por p2
// Retorna um par com o quociente e o resto
// Os coeficientes devem estar em ordem
// decrescente pelo grau. Ex:
// 3x^2 + 2x - 1 \rightarrow [3, 2, -1]
// O(nm), onde n e m sao os tamanhos dos
// polinomios
typedef vector<int> vi;
pair < vi , vi > div(vi p1, vi p2) {
    vi quoc, resto;
    int a = p1.size(), b = p2.size();
    for (int i = 0; i <= a - b; i++) {</pre>
        int k = p1[i] / p2[0];
        quoc.pb(k);
        for (int j = i; j < i + b; j++)
            p1[j] = k * p2[j - i];
    }
    for (int i = a - b + 1; i < a; i++)
        resto.pb(p1[i]);
```

```
return mp(quoc, resto);
2.6 FFT
// Exemplos na main
// Soma O(n) & Multiplicacao O(nlogn)
const int MAX = 5e5;
const int MAX2 = (1 << 20); //(1 << (ceil(log2(MAX)) + 1)) + 1
const double PI = acos(-1);
struct cplx{
    double r, i;
    cplx(double r_ = 0, double i_ = 0):r(r_), i(i_){}
    const cplx operator+(const cplx &x) const{
        return cplx(r + x.r, i + x.i);
    const cplx operator - (const cplx &x) const{
        return cplx(r - x.r, i - x.i);
    const cplx operator*(double a) const {
        return cplx(r*a, i*a);
    const cplx operator/(double a)const {
        return cplx(r/a, i/a);
    const cplx operator*(const cplx x) const {
        return cplx(r*x.r - i*x.i, r*x.i + i*x.r);
    }
};
const cplx I(0, 1);
cplx X[MAX2], Y[MAX2];
int rev[MAX2];
cplx roots[MAX2];
cplx rt(int i, int n){
    double alpha = (2*i*PI)/n;
```

```
return cplx(cos(alpha), sin(alpha));
}
void fft(cplx *a, bool f, int N){
    for (int i = 0; i < N; i++)</pre>
        if (i < rev[i])</pre>
            swap(a[i], a[rev[i]]);
    int 1, r, m;
    for (int n = 2; n \le N; n *= 2) {
        cplx root = (f ? rt(1, n) : rt(-1, n));
        roots[0] = 1;
        for (int i = 1; i < n/2; i++)
            roots[i] = roots[i-1]*root;
        for (int pos = 0; pos < N; pos += n) {
            1 = pos+0, r = pos+n/2, m = 0;
            while (m < n/2) {
                 auto t = roots[m]*a[r];
                a[r] = a[1] - t;
                a[1] = a[1] + t;
                l++; r++; m++;
            }
        }
    if (f) for(int i = 0; i < N; i++) a[i] = a[i]/N;</pre>
}
//p(x) = at(i)*x^i
template < typename T > struct poly : vector < T > {
    poly(const vector<T> &coef):vector<T>(coef){}
    poly(unsigned size):vector<T>(size){}
    poly(){}
    T operator()(T x){
        T ans = 0, curr_x(1);
        for (auto c : *this) {
            ans = ans+c*curr_x;
            curr_x = curr_x*x;
        }
        return ans;
    poly <T> operator+(const poly <T> &r) {
        const poly<T> &l = *this;
```

```
int sz = max(l.size(), r.size());
    poly <T> ans(sz);
    for (unsigned i = 0; i < 1.size(); i++)</pre>
         ans[i] = ans[i]+1[i];
    for (unsigned i = 0; i < r.size(); i++)</pre>
        ans[i] = ans[i]+r[i];
    return ans;
}
poly <T> operator - (poly <T> &r) {
    for (auto &it : r) it = -it;
    return (*this)+r;
}
void fix(int k){
    if (k < this->size()) throw
       logic_error("normalizando errado");
    while (this->size() < k) this->push_back(0);
}
pair < poly < T > , poly < T >> split() {
    const poly<T> &p = *this;
    poly < T > 1, r;
    for (int i = 0; i < p.size(); i++)</pre>
        if (i&1) l.push_back(p[i]);
        else r.push_back(p[i]);
    return {1, r};
}
poly<T> operator*(const poly<T> r){
    const poly<T> &1 = *this;
    int ln = 1.size(), rn = r.size();
    int N = ln+rn+1;
    int n = 1, log_n = 0;
    while (n <= N) { n <<= 1; log_n++; }</pre>
    for (int i = 0; i < n; ++i){</pre>
        rev[i] = 0;
        for (int j = 0; j < log_n; ++j)</pre>
             if (i & (1<<j))</pre>
                 rev[i] = 1 << (log_n-1-j);
    for (int i = 0; i < ln; i++) X[i] = l[i];</pre>
    for (int i = ln; i < n; i++) X[i] = 0;</pre>
    for (int i = 0; i < rn; i++) Y[i] = r[i];</pre>
    for (int i = rn; i < n; i++) Y[i] = 0;</pre>
    fft(X, false, n);//call dft if possible
```

```
fft(Y, false, n);
        for (int i = 0; i < n; i++)</pre>
             Y[i] = X[i]*Y[i];
        fft(Y, true, n);
        poly<T> ans(N);
        for (int i = 0; i < N; i++){</pre>
             ans[i] = floor(Y[i].r + 0.25); //if T is integer
             //ans[i] = Y[i].r; //if T is floating point
        }
         while (!ans.empty() && ans.back() == 0)
             ans.pop_back();
         return ans;
    }
    pair < poly < T > , T > briot_ruffini(T r) \{//\text{for p} = Q(x - r) + \}
        R. returns (Q, R)
        const poly<T> &l = *this;
        int sz = 1.size();
        if (sz == 0) return \{poly < T > (0), 0\};
        poly < T > q(sz - 1);
        q.back() = 1.back();
        for (int i = q.size()-2; i >= 0; i--){
             cout << i << "~" << q.size() << endl;</pre>
             q[i] = q[i+1]*r + l[i+1];
        return {q, q[0]*r + 1[0]};
    }
};
template < typename T > ostream& operator << (ostream &out, const</pre>
   poly < T > &p) {
    if (p.empty()) return out;
    out << p.at(0);
    for (int i = 1; i < p.size(); i++)</pre>
        out << " + " << p.at(i) << "x^" << i;
    out << endl;</pre>
    return out;
}
int main(){ _
```

```
poly < int > p(\{-2, -1, 2, 1\});
    poly<int> q({1, 1, 1});
    poly < int > sum = p+q;
    poly<int> mult = p*q;
    cout << "p: " << p << endl;</pre>
    cout << "q: " << q << endl;</pre>
    cout << "pq: " << mult << endl;</pre>
    for (int i = 1; i <= 50; i++){</pre>
        auto P = p(i), Q = q(i), M = mult(i);
        cout << P*Q << "\t\tvs\t\t" << M << endl;</pre>
        if (abs(P*Q - M) > 1e-5) throw logic_error("bad
            implementation :(");
    }
    cout << "sucesso!" << endl;</pre>
    exit(0):
    for (int root : {1, -1, 2, -2, 3}){
        poly<int> t; int r;
        tie(t, r) = p.briot_ruffini(root);
        cout << p << "/" << poly<int>({-root, 1});
        cout << " = " << endl;
        cout << t << " + " << r << endl;
        cout << endl;</pre>
    }
    exit(0);
     Miller-Rabin
// Testa se n eh primo, n \leq 3 * 10^18
// O(log(n)), considerando multiplicacao
// e exponenciacao constantes
// multiplicacao modular
ll mul(ll x, ll y, ll m); // x*y mod m
ll exp(ll x, ll y, ll m); // x^y mod m;
bool prime(ll n) {
    if (n < 2) return 0;
    if (n <= 3) return 1;</pre>
```

if (n % 2 == 0) return 0;

```
11 d = n - 1;
    int r = 0;
    while (d \% 2 == 0) r++, d /= 2;
    // com esses primos, o teste funciona garantido para n
       <= 3*10^18
    // funciona para n <= 3*10^24 com os primos ate 41
    int a[9] = \{2, 3, 5, 7, 11, 13, 17, 19, 23\};
    // outra opcao para n <= 2^64:</pre>
    // int a[7] = {2, 325, 9375, 28178, 450775, 9780504,
       1795265022};
    for (int i = 0; i < 9; i++) {
        if (a[i] >= n) break;
        ll x = exp(a[i], d, n);
        if (x == 1 \text{ or } x == n - 1) \text{ continue};
        bool deu = 1;
        for (int j = 0; j < r - 1; j++) {
            x = mul(x, x, n);
            if (x == n - 1) {
                deu = 0;
                break;
            }
        }
        if (deu) return 0;
    return 1;
2.8 Inverso Modular
// Computa o inverso de a modulo b
// Se b eh primo, basta fazer
// a^{(b-2)}
```

long long inv(long long a, long long b){

return 1 < a ? b - inv(b%a,a)\*b/a : 1;

}

}

#### 2.9 Variações do crivo de Eratosthenes

```
// "O" crivo
//
// Encontra maior divisor primo
// Um numero eh primo sse div[x] == x
// fact fatora um numero <= lim
// A fatoracao sai ordenada
// \text{ crivo } - O(n \log(\log(n)))
// fact - O(log(n))
int divi[MAX];
void crivo(int lim) {
    for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
    for (int i = 2; i <= lim; i++) if (divi[i] == 1)</pre>
        for (int j = i; j <= lim; j += i) divi[j] = i;</pre>
}
void fact(vector<int>& v, int n) {
    if (n != divi[n]) fact(v, n/divi[n]);
    v.push_back(divi[n]);
}
// Crivo de divisores
//
// Encontra numero de divisores
// ou soma dos divisores
//
// O(n log(n))
int divi[MAX];
void crivo(int lim) {
    for (int i = 1; i <= lim; i++) divi[i] = 1;
    for (int i = 2; i <= lim; i++)</pre>
        for (int j = i; j <= lim; j += i) {</pre>
            // para numero de divisores
             divi[j]++;
```

#### 2.10 Totiente

```
// O(sqrt(n))
int tot(int n){
   int ret = n;

   for (int i = 2; i*i <= n; i++) if (n % i == 0) {
      while (n % i == 0) n /= i;
      ret -= ret / i;
   }
   if (n > 1) ret -= ret / n;

return ret;
}
```

#### 2.11 Exponenciacao rapida

```
// (x^y mod m) em O(log(y))
```

#### 2.12 Produto de dois long long mod m

```
// 0(1)
typedef long long int ll;
ll mul(ll a, ll b, ll m) { // a*b % m
    return (a*b-ll(a*(long double)b/m+0.5)*m+m)%m;
}
```

### 3 Primitivas

#### 3.1 Primitivas de Polinomios

```
#include <bits/stdc++.h>
using namespace std;
namespace algebra {
   const int inf = 1e9;
   const int magic = 500; // threshold for sizes to run the
   naive algo
```

```
namespace fft {
    const int maxn = 1 << 18;</pre>
    typedef double ftype;
    typedef complex <ftype > point;
    point w[maxn];
    const ftype pi = acos(-1);
    bool initiated = 0;
    void init() {
        if(!initiated) {
            for(int i = 1; i < maxn; i *= 2) {</pre>
                 for(int j = 0; j < i; j++) {</pre>
                     w[i + j] = polar(ftype(1), pi * j /
                        i):
                }
            }
            initiated = 1;
        }
    }
    template < typename T>
        void fft(T *in, point *out, int n, int k = 1) {
            if(n == 1) {
                 *out = *in;
            } else {
                 n /= 2;
                 fft(in, out, n, 2 * k);
                 fft(in + k, out + n, n, 2 * k);
                 for(int i = 0; i < n; i++) {</pre>
                     auto t = out[i + n] * w[i + n];
                     out[i + n] = out[i] - t;
                     out[i] += t;
                }
            }
        }
    template < typename T>
        void mul_slow(vector<T> &a, const vector<T> &b) {
            vector<T> res(a.size() + b.size() - 1);
            for(size_t i = 0; i < a.size(); i++) {</pre>
                 for(size_t j = 0; j < b.size(); j++) {</pre>
```

```
res[i + j] += a[i] * b[j];
            }
        }
        a = res;
template < typename T>
    void mul(vector<T> &a, const vector<T> &b) {
        if(min(a.size(), b.size()) < magic) {</pre>
            mul_slow(a, b);
            return;
        }
        init();
        static const int shift = 15, mask = (1 <<</pre>
           shift) - 1:
        size_t n = a.size() + b.size() - 1;
        while(__builtin_popcount(n) != 1) {
            n++:
        }
        a.resize(n);
        static point A[maxn], B[maxn];
        static point C[maxn], D[maxn];
        for(size_t i = 0; i < n; i++) {</pre>
            A[i] = point(a[i] & mask, a[i] >> shift);
            if(i < b.size()) {</pre>
                B[i] = point(b[i] & mask, b[i] >>
                    shift);
            } else {
                B[i] = 0;
            }
        fft(A, C, n); fft(B, D, n);
        for(size t i = 0: i < n: i++) {</pre>
            point c0 = C[i] + conj(C[(n - i) \% n]);
            point c1 = C[i] - conj(C[(n - i) \% n]);
            point d0 = D[i] + conj(D[(n - i) \% n]);
            point d1 = D[i] - conj(D[(n - i) \% n]);
            A[i] = c0 * d0 - point(0, 1) * c1 * d1;
            B[i] = c0 * d1 + d0 * c1;
        fft(A, C, n); fft(B, D, n);
```

```
reverse (C + 1, C + n):
            reverse (D + 1, D + n);
            int t = 4 * n;
            for(size_t i = 0; i < n; i++) {</pre>
                int64_t A0 = llround(real(C[i]) / t);
                T A1 = llround(imag(D[i]) / t);
                T A2 = llround(imag(C[i]) / t);
                a[i] = A0 + (A1 << shift) + (A2 << 2 *
                   shift):
            }
            return;
       }
template < typename T>
   T bpow(T x, size_t n) {
        return n ? n \% 2 ? x * bpow(x, n - 1) : bpow(x *
           x, n / 2) : T(1);
template < typename T>
   T bpow(T x, size_t n, T m) {
        return n ? n % 2 ? x * bpow(x, n - 1, m) % m :
           bpow(x * x % m, n / 2, m) : T(1);
template < typename T>
   T gcd(const T &a, const T &b) {
        return b == T(0)? a : gcd(b, a % b);
   }
template < typename T>
   T nCr(T n, int r) { // runs in O(r)}
       T res(1):
       for(int i = 0; i < r; i++) {
            res *= (n - T(i));
            res /= (i + 1);
        }
        return res;
   }
template < int m>
    struct modular {
        int64_t r;
        modular() : r(0) {}
        modular(int64_t rr) : r(rr) \{if(abs(r) >= m) r\}
```

```
%= m: if(r < 0) r += m:
        modular inv() const {return bpow(*this, m - 2);}
        modular operator * (const modular &t) const
           {return (r * t.r) % m;}
        modular operator / (const modular &t) const
           {return *this * t.inv();}
        modular operator += (const modular &t) {r +=
           t.r; if(r >= m) r -= m; return *this;}
        modular operator -= (const modular &t) {r -=
           t.r; if(r < 0) r += m; return *this;}
        modular operator + (const modular &t) const
           {return modular(*this) += t;}
        modular operator - (const modular &t) const
           {return modular(*this) -= t;}
        modular operator *= (const modular &t) {return
           *this = *this * t;}
        modular operator /= (const modular &t) {return
           *this = *this / t;}
        bool operator == (const modular &t) const
           {return r == t.r;}
        bool operator != (const modular &t) const
           {return r != t.r;}
        operator int64_t() const {return r;}
   };
template < int T>
   istream& operator >> (istream &in, modular <T> &x) {
        return in >> x.r;
   }
template < typename T>
    struct poly {
        vector <T> a:
        void normalize() { // get rid of leading zeroes
            while(!a.empty() && a.back() == T(0)) {
                a.pop_back();
        }
```

```
f}()vlog
polv(T a0) : a{a0}{normalize();}
poly(vector <T> t) : a(t) { normalize();}
poly operator += (const poly &t) {
    a.resize(max(a.size(), t.a.size()));
    for(size_t i = 0; i < t.a.size(); i++) {</pre>
        a[i] += t.a[i];
    }
    normalize():
    return *this;
poly operator -= (const poly &t) {
    a.resize(max(a.size(), t.a.size()));
    for(size_t i = 0; i < t.a.size(); i++) {</pre>
        a[i] -= t.a[i]:
    }
    normalize();
    return *this;
poly operator + (const poly &t) const {return
   polv(*this) += t;}
poly operator - (const poly &t) const {return
   poly(*this) -= t;}
poly mod_xk(size_t k) const { // get same
   polynomial mod x^k
   k = min(k, a.size());
   return vector < T > (begin(a), begin(a) + k);
poly mul_xk(size_t k) const { // multiply by x^k
    poly res(*this);
    res.a.insert(begin(res.a), k, 0);
    return res:
poly div_xk(size_t k) const { // divide by x^k,
   dropping coefficients
   k = min(k, a.size());
    return vector <T > (begin(a) + k, end(a));
poly substr(size_t l, size_t r) const { //
   return mod xk(r).div xk(l)
```

```
1 = min(1, a.size()):
   r = min(r, a.size());
    return vector<T>(begin(a) + 1, begin(a) + r);
poly inv(size_t n) const { // get inverse series
   mod x^n
    assert(!is_zero());
    poly ans = a[0].inv();
    size t a = 1:
    while (a < n) {
        poly C = (ans * mod_xk(2 * a)).substr(a,
        ans -= (ans * C).mod_xk(a).mul_xk(a);
        a *= 2;
    return ans.mod_xk(n);
}
poly operator *= (const poly &t) {fft::mul(a,
   t.a); normalize(); return *this;}
polv operator * (const polv &t) const {return
   poly(*this) *= t;}
poly reverse(size_t n, bool rev = 0) const { //
   reverses and leaves only n terms
    polv res(*this);
   if(rev) { // If rev = 1 then tail goes to
        res.a.resize(max(n, res.a.size()));
    std::reverse(res.a.begin(), res.a.end());
    return res.mod_xk(n);
}
pair < poly , poly > divmod_slow(const poly &b)
   const { // when divisor or quotient is small
    vector <T> A(a):
   vector <T> res:
    while(A.size() >= b.a.size()) {
        res.push_back(A.back() / b.a.back());
        if(res.back() != T(0)) {
            for(size_t i = 0; i < b.a.size();</pre>
```

```
i++) {
                A[A.size() - i - 1] -=
                    res.back() * b.a[b.a.size() -
                   i - 1];
            }
        A.pop_back();
    std::reverse(begin(res), end(res));
    return {res, A};
}
pair < poly , poly > divmod(const poly &b) const {
   // returns quotiend and remainder of a mod b
    if(deg() < b.deg()) {</pre>
        return {poly{0}, *this};
    }
    int d = deg() - b.deg();
    if(min(d, b.deg()) < magic) {</pre>
        return divmod_slow(b);
    }
    poly D = (reverse(d + 1) * b.reverse(d +
       1).inv(d + 1)).mod_xk(d + 1).reverse(d +
       1, 1);
   return {D, *this - D * b};
}
poly operator / (const poly &t) const {return
   divmod(t).first;}
poly operator % (const poly &t) const {return
   divmod(t).second;}
poly operator /= (const poly &t) {return *this =
   divmod(t).first;}
poly operator %= (const poly &t) {return *this =
   divmod(t).second:}
poly operator *= (const T &x) {
   for(auto &it: a) {
        it *= x;
    normalize();
    return *this;
}
```

```
poly operator /= (const T &x) {
    for(auto &it: a) {
        it /= x;
    normalize();
    return *this;
poly operator * (const T &x) const {return
   poly(*this) *= x;}
poly operator / (const T &x) const {return
   poly(*this) /= x;}
void print() const {
    for(auto it: a) {
        cout << it << ', ';
    cout << endl;</pre>
T eval(T x) const { // evaluates in single point
   T res(0);
    for(int i = int(a.size()) - 1; i >= 0; i--) {
        res *= x;
        res += a[i];
    }
    return res;
}
T& lead() { // leading coefficient
    return a.back();
}
int deg() const { // degree
    return a.empty() ? -inf : a.size() - 1;
bool is_zero() const { // is polynomial zero
    return a.empty();
T operator [](int idx) const {
    return idx >= (int)a.size() || idx < 0 ?</pre>
       T(0) : a[idx];
}
```

```
T& coef(size_t idx) { // mutable reference at
   coefficient
    return a[idx];
bool operator == (const poly &t) const {return a
   == t.a;}
bool operator != (const poly &t) const {return a
poly deriv() { // calculate derivative
    vector<T> res;
    for(int i = 1; i <= deg(); i++) {</pre>
        res.push_back(T(i) * a[i]);
    return res;
poly integr() { // calculate integral with C = 0
    vector < T > res = \{0\};
    for(int i = 0; i <= deg(); i++) {</pre>
        res.push_back(a[i] / T(i + 1));
    }
    return res;
size_t leading_xk() const { // Let p(x) = x^k *
   t(x), return k
    if(is_zero()) {
        return inf;
    int res = 0;
    while (a[res] == T(0)) {
        res++;
    return res;
poly log(size_t n) { // calculate log p(x) mod
    assert(a[0] == T(1));
    return (deriv().mod_xk(n) *
       inv(n)).integr().mod_xk(n);
poly exp(size_t n) { // calculate exp p(x) mod
   x^n
```

```
if(is_zero()) {
        return T(1);
    assert(a[0] == T(0));
    poly ans = T(1);
   size_t a = 1;
    while (a < n) {
        poly C = ans.log(2 * a).div_xk(a) -
           substr(a, 2 * a);
        ans -= (ans * C).mod_xk(a).mul_xk(a);
        a *= 2:
   return ans.mod_xk(n);
poly pow_slow(size_t k, size_t n) { // if k is
   small
   return k ? k % 2 ? (*this * pow_slow(k - 1,
       n)).mod_xk(n) : (*this *
       *this).mod_xk(n).pow_slow(k / 2, n):
       T(1);
poly pow(size_t k, size_t n) { // calculate
   p^k(n) mod x^n
   if(is_zero()) {
        return *this;
   if(k < magic) {</pre>
        return pow_slow(k, n);
   int i = leading_xk();
   T j = a[i];
   poly t = div_xk(i) / j;
   return bpow(j, k) * (t.log(n) *
       T(k) .exp(n).mul_xk(i * k).mod_xk(n);
poly mulx(T x) { // component-wise
   multiplication with x^k
   T cur = 1;
   polv res(*this);
   for(int i = 0; i <= deg(); i++) {</pre>
        res.coef(i) *= cur;
```

```
cur *= x:
    }
    return res;
poly mulx_sq(T x) { // component-wise
   multiplication with x^{k^2}
    T cur = x;
    T \text{ total} = 1;
    T xx = x * x;
    polv res(*this);
    for(int i = 0; i <= deg(); i++) {</pre>
        res.coef(i) *= total;
        total *= cur;
        cur *= xx;
    }
    return res;
}
vector <T> chirpz_even(T z, int n) { // P(1),
   P(z^2), P(z^4), ..., P(z^2(n-1))
    int m = deg();
    if(is_zero()) {
        return vector <T>(n, 0);
    vector <T> vv(m + n);
    T zi = z.inv();
    T zz = zi * zi;
    T cur = zi;
    T \text{ total} = 1;
    for (int i = 0; i \le max(n - 1, m); i++) {
        if(i \le m) \{vv[m - i] = total:\}
        if(i < n) {vv[m + i] = total;}</pre>
        total *= cur;
        cur *= zz;
    poly w = (mulx_sq(z) * vv).substr(m, m +
       n).mulx_sq(z);
    vector <T> res(n);
    for(int i = 0; i < n; i++) {</pre>
        res[i] = w[i];
    return res;
}
```

```
vector\langle T \rangle chirpz(T z, int n) \{ // P(1), P(z), 
   P(z^2), ..., P(z^{(n-1)})
    auto even = chirpz_even(z, (n + 1) / 2);
    auto odd = mulx(z).chirpz_even(z, n / 2);
    vector <T> ans(n);
    for(int i = 0; i < n / 2; i++) {</pre>
        ans [2 * i] = even[i];
        ans[2 * i + 1] = odd[i];
    }
    if(n % 2 == 1) {
        ans[n - 1] = even.back();
    return ans;
template < typename iter >
    vector<T> eval(vector<poly> &tree, int v,
       iter 1, iter r) { // auxiliary evaluation
       function
       if(r - 1 == 1) {
            return {eval(*1)};
        } else {
            auto m = 1 + (r - 1) / 2;
            auto A = (*this % tree[2 *
               v]).eval(tree, 2 * v, 1, m);
            auto B = (*this \% tree[2 * v +
               1]).eval(tree, 2 * v + 1, m, r);
            A.insert(end(A), begin(B), end(B));
            return A;
        }
vector<T> eval(vector<T> x) { // evaluate
   polynomial in (x1, ..., xn)
   int n = x.size();
   if(is zero()) {
        return vector <T>(n, T(0));
    vector < poly > tree (4 * n);
    build(tree, 1, begin(x), end(x));
    return eval(tree, 1, begin(x), end(x));
template < typename iter >
    poly inter(vector<poly> &tree, int v, iter
```

```
1, iter r, iter ly, iter ry) { //
               auxiliary interpolation function
                if(r - 1 == 1) {
                    return {*ly / a[0]};
                } else {
                    auto m = 1 + (r - 1) / 2;
                    auto my = ly + (ry - ly) / 2;
                    auto A = (*this % tree[2 *
                       v]).inter(tree, 2 * v, 1, m, ly,
                       my);
                    auto B = (*this \% tree[2 * v +
                       1]).inter(tree, 2 * v + 1, m, r,
                       my, ry);
                    return A * tree[2 * v + 1] + B *
                       tree[2 * v]:
                }
            }
   };
template < typename T>
    poly<T> operator * (const T& a, const poly<T>& b) {
        return b * a;
   }
template < typename T>
   poly<T> xk(int k) { // return x^k
        return poly<T>{1}.mul_xk(k);
   }
template < typename T>
   T resultant(poly<T> a, poly<T> b) { // computes
       resultant of a and b
       if(b.is zero()) {
            return 0;
        } else if(b.deg() == 0) {
            return bpow(b.lead(), a.deg());
        } else {
            int pw = a.deg();
            a \%= b;
            pw -= a.deg();
           T mul = bpow(b.lead(), pw) * T((b.deg() &
               a.deg() & 1) ? -1 : 1);
            T ans = resultant(b, a);
```

```
return ans * mul:
            }
    template < typename iter >
        poly<typename iter::value_type> kmul(iter L, iter R)
           { // \text{ computes } (x-a1)(x-a2)...(x-an) \text{ without }
           building tree
            if(R - L == 1) {
                return vector < typename
                    iter::value_type>{-*L, 1};
            } else {
                iter M = L + (R - L) / 2;
                return kmul(L, M) * kmul(M, R);
            }
    template < typename T, typename iter >
        poly<T> build(vector<poly<T>> &res, int v, iter L,
           iter R) { // builds evaluation tree for
           (x-a1)(x-a2)...(x-an)
            if(R - L == 1) {
                return res[v] = vector <T>{-*L, 1};
            } else {
                iter M = L + (R - L) / 2;
                return res[v] = build(res, 2 * v, L, M) *
                    build(res, 2 * v + 1, M, R);
            }
    template < typename T>
        poly<T> inter(vector<T> x, vector<T> y) { //
            interpolates minimum polynomial from (xi, yi)
           pairs
            int n = x.size();
            vector<poly<T>> tree(4 * n);
            return build(tree, 1, begin(x),
                end(x)).deriv().inter(tree, 1, begin(x),
                end(x), begin(y), end(y));
        }
};
using namespace algebra;
const int mod = 1e9 + 7;
```

```
typedef modular < mod > base;
typedef poly < base > polyn;
using namespace algebra;
signed main() {
    ios::sync_with_stdio(0);
    cin.tie(0);
    int n = 100000;
    polyn a;
    vector < base > x;
    for(int i = 0; i <= n; i++) {</pre>
        a.a.push_back(1 + rand() % 100);
        x.push_back(1 + rand() \% (2 * n));
    sort(begin(x), end(x));
    x.erase(unique(begin(x), end(x)), end(x));
    auto b = a.eval(x);
    cout << clock() / double(CLOCKS_PER_SEC) << endl;</pre>
    auto c = inter(x, b);
    polyn md = kmul(begin(x), end(x));
    cout << clock() / double(CLOCKS_PER_SEC) << endl;</pre>
    assert(c == a % md);
    return 0;
}
```

#### 3.2 Primitivas Geometricas

```
#include <bits/stdc++.h>
using namespace std;

#define sc(a) scanf("%d", &a)
#define sc2(a,b) scanf("%d %d", &a, &b)
#define pri(x) printf("%d\n", x)
#define prie(x) printf("%d ", x)
#define sz(x) (int)((x).size())
#define mp make_pair
#define pb push_back
#define f first
#define s second
#define sq(x) ((x)*(x))
```

```
#define BUFF ios::sync_with_stdio(false)
typedef long long int 11;
typedef double ld;
typedef pair<int, int> ii;
typedef vector<int> vi;
typedef vector < vi > vvi;
typedef vector<ii> vii;
const int INF = 0x3f3f3f3f;
const 11 LINF = 0x3f3f3f3f3f3f3f3f3f11;
const ld DINF = 1e18;
const ld pi = acos(-1.0);
const ld eps = 1e-9;
bool eq(ld a, ld b) {
    return abs(a - b) <= eps;</pre>
}
struct pt { // ponto
    ld x, v;
    pt() {}
    pt(ld x, ld y) : x(x), y(y) {}
    bool operator < (const pt p) const {</pre>
        if (!eq(x, p.x)) return x < p.x;</pre>
        return y < p.y;</pre>
    }
    bool operator == (const pt p) const {
        return eq(x, p.x) and eq(y, p.y);
    pt operator + (const pt p) const { return pt(x+p.x,
       y+p.y); }
    pt operator - (const pt p) const { return pt(x-p.x,
       y-p.y); }
    pt operator * (const ld c) const { return pt(x*c , y*c
       ); }
    pt operator / (const ld c) const { return pt(x/c , y/c
       ); }
};
struct line { // reta
    pt p, q;
    line() {}
```

```
line(pt p, pt q) : p(p), q(q) {}
};
// PONTO & VETOR
ld dist(pt p, pt q) { // distancia
    return sqrt(sq(p.x - q.x) + sq(p.y - q.y));
}
ld dist2(pt p, pt q) { // quadrado da distancia
    return sq(p.x - q.x) + sq(p.y - q.y);
}
ld norm(pt v) { // norma do vetor
    return dist(pt(0, 0), v);
}
pt normalize(pt v) { // vetor normalizado
    if (!norm(v)) return v;
    v = v / norm(v);
    return v;
}
ld dot(pt u, pt v) { // produto escalar
    return u.x * v.x + u.y * v.y;
}
ld cross(pt u, pt v) { // norma do produto vetorial
    return u.x * v.y - u.y * v.x;
}
ld sarea(pt p, pt q, pt r) { // area com sinal
    return cross(q - p, r - q) / 2;
}
bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
    return eq(sarea(p, q, r), 0);
}
int paral(pt u, pt v) { // se u e v sao paralelos
    u = normalize(u);
    v = normalize(v):
```

```
if (eq(u.x, v.x)) and eq(u.y, v.y) return 1;
    if (eq(u.x, -v.x) \text{ and } eq(u.y, -v.y)) \text{ return } -1;
    return 0;
}
bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
   return sarea(p, q, r) > 0;
pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
    return pt(p.x * cos(th) - p.y * sin(th),
            p.x * sin(th) + p.y * cos(th));
}
pt rotate90(pt p) { // rotaciona 90 graus
    return pt(-p.v, p.x);
}
// R.F.T.A
bool isvert(line r) { // se r eh vertical
    return eq(r.p.x, r.q.x);
ld getm(line r) { // coef. ang. de r
    if (isvert(r)) return DINF;
   return (r.p.y - r.q.y) / (r.p.x - r.q.x);
ld getn(line r) { // coef. lin. de r
    if (isvert(r)) return DINF;
    return r.p.y - getm(r) * r.p.x;
}
bool lineeq(line r, line s) { // r == s
    return col(r.p, r.q, s.p) and col(r.p, r.q, s.q);
}
bool paraline(line r, line s) { // se r e s sao paralelas
    if (isvert(r) and isvert(s)) return 1;
    if (isvert(r) or isvert(s)) return 0;
    return eq(getm(r), getm(s));
```

```
}
bool isinline(pt p, line r) { // se p pertence a r
    return col(p, r.p, r.q);
}
bool isinseg(pt p, line r) { // se p pertence ao seg de r
    if (p == r.p or p == r.q) return 1;
    return paral(p - r.p, p - r.q) == -1;
}
pt proj(pt p, line r) { // projecao do ponto p na reta r
    if (r.p == r.q) return r.p;
   r.q = r.q - r.p; p = p - r.p;
    pt proj = r.q * (dot(p, r.q) / dot(r.q, r.q));
    return proj + r.p;
}
pt inter(line r, line s) { // r inter s
    if (paraline(r, s)) return pt(DINF, DINF);
    if (isvert(r)) return pt(r.p.x, getm(s) * r.p.x +
       getn(s));
    if (isvert(s)) return pt(s.p.x, getm(r) * s.p.x +
       getn(r));
    1d x = (getn(s) - getn(r)) / (getm(r) - getm(s));
    return pt(x, getm(r) * x + getn(r));
}
bool interseg(line r, line s) { // se o seg de r intercepta
   o seg de s
    if (paraline(r, s)) {
        return isinseg(r.p, s) or isinseg(r.q, s)
            or isinseg(s.p, r) or isinseg(s.q, r);
    }
    pt i = inter(r, s);
    return isinseg(i, r) and isinseg(i, s);
}
ld disttoline(pt p, line r) { // distancia do ponto a reta
```

```
return dist(p, proj(p, r));
}
ld disttoseg(pt p, line r) { // distancia do ponto ao seg
    if (isinseg(proj(p, r), r))
        return disttoline(p, r);
   return min(dist(p, r.p), dist(p, r.q));
}
ld distseg(line a, line b) { // distancia entre seg
    if (interseg(a, b)) return 0;
    ld ret = DINF;
   ret = min(ret, disttoseg(a.p, b));
   ret = min(ret, disttoseg(a.q, b));
   ret = min(ret, disttoseg(b.p, a));
    ret = min(ret, disttoseg(b.q, a));
    return ret;
}
// POLIGONO
ld polper(vector<pt> v) { // perimetro do poligono
   ld ret = 0;
    for (int i = 0; i < sz(v); i++)
        ret += dist(v[i], v[(i + 1) \% sz(v)]);
    return ret;
}
ld polarea(vector<pt> v) { // area do poligono
   ld ret = 0;
   for (int i = 0; i < sz(v); i++)
        ret += sarea(pt(0, 0), v[i], v[(i + 1) \% sz(v)]);
    return abs(ret):
}
bool onpol(pt p, vector<pt> v) { // se um ponto esta na
   fronteira do poligono
   for (int i = 0; i < sz(v); i++)
        if (isinseg(p, line(v[i], v[(i + 1) % sz(v)])))
           return 1:
```

```
return 0:
}
bool inpol(pt p, vector<pt> v) { // se um ponto pertence ao
   poligono
    if (onpol(p, v)) return 1;
    int c = 0;
    line r = line(p, pt(DINF, pi * DINF));
    for (int i = 0; i < sz(v); i++) {
        line s = line(v[i], v[(i + 1) \% sz(v)]);
        if (interseg(r, s)) c++;
    return c & 1;
}
bool interpol(vector<pt> v1, vector<pt> v2) { // se dois
   poligonos se interceptam
    for (int i = 0; i < sz(v1); i++) if (inpol(v1[i], v2))</pre>
       return 1:
    for (int i = 0; i < sz(v2); i++) if (inpol(v2[i], v1))
       return 1;
    return 0;
}
ld distpol(vector<pt> v1, vector<pt> v2) { // distancia
   entre poligonos
    if (interpol(v1, v2)) return 0;
    ld ret = DINF;
    for (int i = 0; i < sz(v1); i++) for (int j = 0; j <
       sz(v2); j++)
        ret = min(ret, distseg(line(v1[i], v1[(i + 1) %
           sz(v1)]).
                    line(v2[j], v2[(j + 1) \% sz(v2)])));
    return ret:
}
vector<pt> convexhull(vector<pt> v) { // convex hull
    vector < pt > 1, u;
    sort(v.begin(), v.end());
```

```
for (int i = 0; i < sz(v); i++) {</pre>
        while (sz(1) > 1 \text{ and } !ccw(v[i], 1[sz(1) - 1],
           l[sz(1) - 2]))
            1.pop_back();
        1.pb(v[i]);
    }
    for (int i = sz(v) - 1; i \ge 0; i--) {
        while (sz(u) > 1 \text{ and } !ccw(v[i], u[sz(u) - 1],
           u[sz(u) - 2])
            u.pop_back();
        u.pb(v[i]);
    }
    1.pop_back(); u.pop_back();
    for (int i = 0; i < sz(u); i++) 1.pb(u[i]);
    return 1;
}
// CIRCULO
pt getcenter(pt a, pt b, pt c) { // centro da circunferencia
   dado 3 pontos
    b = (a + b) / 2;
    c = (a + c) / 2;
    return inter(line(b, b + rotate90(a - b)),
            line(c, c + rotate90(a - c)));
}
circle minCirc(vector < PT > v) { // minimum enclosing circle
    int n = v.size():
    random_shuffle(v.begin(), v.end());
    PT p = PT(0.0, 0.0);
    circle ret = circle(p, 0.0);
    for(int i = 0; i < n; i++) {</pre>
        if(!inside(ret, v[i])) {
            ret = circle(v[i], 0);
            for(int j = 0; j < i; j++) {
                if(!inside(ret, v[j])) {
```

```
ret = circle((v[i] + v[j]) / 2.0,
                        sqrt(dist2(v[i], v[j])) / 2.0);
                    for(int k = 0; k < j; k++) {
                         if(!inside(ret, v[k])) {
                             p = best0f3(v[i], v[j], v[k]);
                             ret = circle(p, sqrt(dist2(p,
                                v[i])));
                        }
                    }
                }
            }
        }
    }
    return ret;
}
// comparador pro set para fazer sweep angle com segmentos
double ang;
struct cmp {
    bool operator () (const line& a, const line& b) {
        line r = line(pt(0, 0), rotate(pt(1, 0), ang));
        return norm(inter(r, a)) < norm(inter(r, b));</pre>
};
3.3 Primitivas de matriz (Rafael)
11 mod(l1 v) { return (v + MOD) % MOD; }
11 sum(ll 1, ll r) { return mod(l+r); }
11 mult(11 1, 11 r) { return mod(1*r); }
ll inverse(ll 1) { return inv(1, MOD); }
bool equal(ll\ l,\ ll\ r) { return mod(l-r) == 0; }
template < typename T> struct matrix {
    vector < vector < T >> in;
    int row, col;
    void print(){//
```

for (int i = 0; i < row; i++){</pre>

cout << endl;</pre>

for (int j = 0; j < col; j++)

cout << in[i][j] << " ";</pre>

```
}
}
matrix(int row, int col, int op = 0):row(row), col(col),
   in(row, vector <T>(col, 0)){
    if (op) for (int i = 0; i < row; i++) in[i][i] = 1;</pre>
}
matrix(initializer_list<initializer_list<T>> c):
    row(c.size()), col((*c.begin()).size()){
        in = vector < vector < T >> (row, vector < T > (col, 0));
        int i, j;
        i = 0;
        for (auto &it : c){
            j = 0;
            for (auto &jt : it){
                 in[i][j] = jt;
                 j++;
            i++;
        }
    }
T &operator()(int i, int j){ return in[i][j]; }
//in case of a transposed matrix, swap i and j
matrix<T>& operator*=(T t){
    matrix<T> &l = *this;
    for (int i = 0; i < row; i++)</pre>
        for (int j = 0; j < col; j++)</pre>
            l(i, j) = mult(l(i, j), t); //% MOD) % MOD;
    return 1;
}
matrix <T> operator + (matrix <T> &r) {
    matrix<T> &l = *this;
    matrix<T> m(row, col, 0);
    for (int i = 0; i < row; i++)</pre>
        for (int j = 0; j < col; j++)
            m(i, j) = sum(l(i, j), r(i, j)); //% MOD) %
    return m;
matrix <T> operator*(matrix <T> &r){
    matrix < T > &l = *this;
    int row = 1.row;
```

```
int col = r.col;
    int K = 1.col;
    matrix<T> m(row, col, 0);
    for (int i = 0; i < row; i++)</pre>
        for (int j = 0; j < col; j++)
            for (int k = 0; k < K; k++)
                 m(i, j) = sum(m(i, j), mult(l(i, k),
                    r(k, j)));
    return m;
matrix<T> operator^(long long e){
    matrix < T > &m = (*this);
    if (e == 0) return matrix(m.row, m.row, 1);
    if (e == 1) return m;
    if (e == 2) return m*m;
    auto m_{-} = m^{(e/2)}; m_{-} = m_{*m_{-}};
    if (e\%2 == 1) m_{-} = m_{-} * m;
    return m_;
void multiply_r(int i, T k){
    matrix < T > &m = (*this);
    for (int j = 0; j < col; j++)</pre>
        m(i, j) = mult(m(i, j), k);
void multiply_c(int j, T k){
    matrix < T > &m = (*this);
    for (int i = 0; i < row; j++)
        m(i, j) = mult(m(i, j), k);
}
void sum_r(int i1, int i2, T k){
    matrix < T > \&m = (*this);
    for (int j = 0; j < col; j++)</pre>
        m(i1, j) = sum(m(i1, j), mult(k, m(i2, j)));
bool gaussian(int I, int J){
    matrix < T > \&m = (*this);
    T \text{ tmp} = m(I, J);
    if (equal(tmp, 0)) return false;
    multiply_r(I, inverse(tmp));
    for (int i = 0; i < row; i++)</pre>
        if (i != I) sum_r(i, I, mult(-1, m(i, J)));
```

```
multiply_r(I, tmp);
    return true;
}
T determinant(){
    matrix<T> m = (*this);
    for (int i = 0; i < row; i++)
        if (!m.gaussian(i, i)) return 0;

T ans = 1;
    for (int i = 0; i < row; i++)
        ans = mult(ans, m(i, i));
    return ans;
}
};</pre>
```

### 4 Estruturas

## 4.1 Sparse Table

```
// Resolve RMQ
// MAX2 = log(MAX)
//
// Complexidades:
// build - O(n log(n))
// query - 0(1)
namespace sparse {
    int m[MAX2][MAX], n;
    void build(int n2, int* v) {
        n = n2:
        for (int i = 0; i < n; i++) m[0][i] = v[i];
        for (int j = 1; (1<<j) <= n; j++) for (int i = 0;
           i+(1<<j) <= n; i++)
            m[j][i] = min(m[j-1][i], m[j-1][i+(1<<(j-1))]);
    }
    int query(int a, int b) {
        int j = __builtin_clz(0) - __builtin_clz(b-a+1) - 1;
        return min(m[j][a], m[j][b-(1<<j)+1]);</pre>
    }
```

### 4.2 Treap

```
// Usar static treap<int> t;
// Para usar, chamar o Rafael
// Complexidades:
// insert - O(log(n))
// erase - O(log(n))
// query - O(log(n))
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
template < typename T > struct treap {
    struct node {
        int p;
        int 1, r;
        T v;
        int sz;
        T sum;
        bool rev;
        node(){}
        node(T v):p(rng()), l(-1), r(-1), v(v), sz(1),
           rev(false){}
    } t[MAX];
    int it;
    //vector < node > t;
    treap(){ it = 0; }
    int size(int i){
        if (i == -1) return 0;
        return t[i].sz;
    }
    void fix(int i){
        if (i == -1) return;
        if (t[i].rev) {
            int &1 = t[i].1;
            int &r = t[i].r;
            swap(1, r);
            t[i].sz = 1 + size(1) + size(r);
            if (1 != -1)
                t[1].rev ^= true;
```

```
if (r != -1)
            t[r].rev ^= true;
        t[i].rev = false;
    }
}
void update(int i){
    if (i == -1) return;
    t[i].sum = t[i].v;
    int l = t[i].1;
    int r = t[i].r;
    t[i].sz = 1 + size(1) + size(r);
}
void split_value(int i, int k, int &1, int &r){ //values
   must be ordered
    if (i == -1){
        1 = -1; r = -1;
        return;
    }
    fix(i);
    if (t[i].v < k){
        split_value(t[i].r, k, l, r);
        t[i].r = 1;
        1 = i;
    }
    else{
        split_value(t[i].1, k, 1, r);
        t[i].l = r;
        r = i;
    }
    update(i);
}
//implicit
void split(int i, int k, int &l, int &r, int sz = 0){
   //key
    if (i == -1){
        1 = -1; r = -1;
        return;
    }
    fix(i);
    int inc = size(t[i].1); //quantidade elementos menor
```

```
aue k
    if (sz+inc < k){
        split(t[i].r, k, l, r, sz+inc+1);
        t[i].r = 1;
       1 = i;
    }
    else{
        split(t[i].1, k, 1, r, sz);
       t[i].l = r;
       r = i;
    }
    update(i);
}
int merge(int 1, int r){ //priority
    if (1 == -1) return r;
    if (r == -1) return 1;
    fix(1); fix(r);
    if (t[1].p > t[r].p){
        t[1].r = merge(t[1].r, r);
        update(1);
       return 1;
   }
    else{
       t[r].1 = merge(1, t[r].1);
        update(r);
        return r;
    }
}
void insert(int &root, T v, int pos){
    int m = it++;
    t[m] = node(v);
    if (root == -1){
        root = m;
        return;
    }
    int 1, r;
    split(root, pos, 1, r);
   l = merge(l, m);
    1 = merge(1, r);
    root = 1;
}
```

```
T query(int &root, int M){
        int 1, m, r;
        split(root, M+1, m, r);
        split(m, M, 1, m);
        T ans = t[m].v;
        l = merge(l, m);
        l = merge(l, r);
        root = 1;
        return ans;
    }
    void reverse(int &root, int L, int R){
        int 1, m, r;
        split(root, R+1, m, r);
        split(m, L, 1, m);
        t[m].rev ^= 1;
        1 = merge(1, m);
        l = merge(l, r);
        root = 1;
    }
    void print(int i, int &size){
        if (i == -1) return;
        print(t[i].1, size);
        cout << "#" << size << ": " << t[i].v << endl;</pre>
        size++;
        print(t[i].r, size);
    }
};
4.3 SQRT-decomposition
// Resolve RMQ
// O-indexed
// MAX2 = sqrt(MAX)
//
// O bloco da posicao x eh
// sempre x/q
// Complexidades:
// build - O(n)
// query - 0(sqrt(n))
```

```
int n, q;
int v[MAX];
int bl[MAX2];
void build() {
    q = (int) sqrt(n);
     // computa cada bloco
    for (int i = 0; i <= q; i++) {</pre>
        bl[i] = INF;
        for (int j = 0; j < q and q * i + j < n; j++)
            bl[i] = min(bl[i], v[q * i + j]);
    }
}
int query(int a, int b) {
    int ret = INF;
    // linear no bloco de a
    for (; a <= b and a % q; a++) ret = min(ret, v[a]);</pre>
    // bloco por bloco
    for (; a + q <= b; a += q) ret = min(ret, bl[a / q]);</pre>
    // linear no bloco de b
    for (; a <= b; a++) ret = min(ret, v[a]);</pre>
    return ret;
}
4.4 BIT 2D
// BIT de soma 1-based
// Para mudar o valor da posicao (x, y) para k,
// faca: poe(x, y, k - sum(x, y, x, y))
// Complexidades:
// poe - O(log^2(n))
// query - O(log^2(n))
int n;
```

```
int bit[MAX][MAX];
void poe(int x, int y, int k) {
   for (int y2 = y; x <= n; x += x & -x)
        for (y = y2; y \le n; y += y & -y)
            bit[x][y] += k;
}
int sum(int x, int y) {
    int ret = 0;
    for (int y2 = y; x; x -= x & -x)
        for (y = y2; y; y -= y & -y)
            ret += bit[x][y];
    return ret;
}
int query(int x, int y, int z, int w) {
    return sum(z, w) - sum(x-1, w)
        - sum(z, y-1) + sum(x-1, y-1);
}
4.5 MergeSort Tree
// query(a, b, val) retorna numero de
// elementos em [a, b] <= val</pre>
// Usa O(n log(n)) de memoria
//
// Complexidades:
// build - O(n log(n))
// \text{ query - O(log^2(n))}
#define ALL(x) x.begin(),x.end()
int v[MAX], n;
vector < vector < int > > tree(4*MAX);
void build(int p, int l, int r) {
   if (1 == r) return tree[p].push_back(v[1]);
    int m = (1+r)/2;
   build(2*p, 1, m), build(2*p+1, m+1, r);
    merge(ALL(tree[2*p]), ALL(tree[2*p+1]),
```

```
back_inserter(tree[p]));
}
int query(int a, int b, int val, int p=1, int l=0, int
   r=n-1) {
    if (b < l or r < a) return 0; // to fora</pre>
    if (a \le 1 \text{ and } r \le b) // \text{ to totalmente dentro}
        return lower_bound(ALL(tree[p]), val+1) -
            tree[p].begin();
    int m = (1+r)/2;
    return query(a, b, val, 2*p, 1, m) + query(a, b, val,
       2*p+1, m+1, r);
}
4.6 SegTree
// Recursiva com Lazy Propagation
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
//
// Complexidades:
// build - O(n)
// \text{ query - O(log(n))}
// update - O(log(n))
namespace seg {
    11 seg[4*MAX], lazy[4*MAX];
    int n, *v;
    ll build(int p=1, int l=0, int r=n-1) {
        lazy[p] = 0;
        if (1 == r) return seg[p] = v[1];
        int m = (1+r)/2;
        return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1,
           r);
    }
    void build(int n2, int* v2) {
        n = n2, v = v2;
        build();
    void prop(int p, int l, int r) {
        if (!lazy[p]) return;
```

```
int m = (1+r)/2:
        seg[2*p] += lazy[p]*(m-l+1);
        seg[2*p+1] += lazy[p]*(r-(m+1)+1);
        lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
        lazy[p] = 0;
    }
    11 query(int a, int b, int p=1, int l=0, int r=n-1) {
        if (b < 1 or r < a) return 0;
        if (a <= l and r <= b) return seg[p];</pre>
        prop(p, 1, r);
        int m = (1+r)/2;
        return query (a, b, 2*p, 1, m) + query (a, b, 2*p+1, m)
           m+1. r):
    }
    ll update(int a, int b, int x, int p=1, int l=0, int
        if (b < l or r < a) return seg[p];</pre>
        if (a \le 1 \text{ and } r \le b)
            seg[p] += (11)x*(r-1+1);
            lazy[p] += x;
            return seg[p];
        }
        prop(p, 1, r);
        int m = (1+r)/2;
        return seg[p] = update(a, b, x, 2*p, 1, m) +
            update(a, b, x, 2*p+1, m+1, r);
    }
};
4.7 SegTree Esparca
// Query: soma do range [a, b]
// Update: flipa os valores de [a, b]
//
// Complexidades:
// build - O(n)
// query - O(log^2(n))
// update - 0(log^2(n))
typedef long long 11;
namespace seg {
```

```
unordered_map<11, int> t, laz;
void build() { t.clear(), lazy.clear(); }
void prop(ll p, int l, int r) {
    if (!lazv[p]) return;
    t[p] = r-1+1-t[p];
    if (1 != r) lazy[2*p]^=lazy[p], lazy[2*p+1]^=lazy[p];
    lazv[p] = 0;
}
int query(int a, int b, ll p=1, int l=0, int r=N-1) {
    prop(p, 1, r);
    if (b < 1 or r < a) return 0;</pre>
    if (a <= l and r <= b) return t[p];</pre>
    int m = 1+r >> 1;
    return query(a, b, 2*p, 1, m)+query(a, b, 2*p+1,
       m+1, r);
}
int update(int a, int b, ll p=1, int l=0, int r=N-1) {
    prop(p, 1, r);
    if (b < 1 or r < a) return t[p];</pre>
    if (a \le 1 \text{ and } r \le b) \{
        lazv[p] ^= 1;
        prop(p, 1, r);
        return t[p];
    }
    int m = 1+r >> 1;
    return t[p] = update(a, b, 2*p, 1, m) + update(a, b, a)
        2*p+1, m+1, r);
}
```

## 4.8 SegTree Iterativa com Lazy Propagation

};

```
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Para mudar, mudar as funcoes junta, poe e query
// LOG = ceil(log2(MAX))
//
```

```
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
namespace seg {
    11 \text{ seg}[2*MAX], lazy[2*MAX];
    ll junta(ll a, ll b) {
        return a+b;
    }
    // soma x na posicao p de tamanho tam
    void poe(int p, ll x, int tam, bool prop=1) {
        seg[p] += x*tam;
        if (prop and p < n) lazy[p] += x;</pre>
    }
    // atualiza todos os pais da folha p
    void sobe(int p) {
        for (int tam = 2; p /= 2; tam *= 2) {
            seg[p] = junta(seg[2*p], seg[2*p+1]);
            poe(p, lazy[p], tam, 0);
        }
    }
    // propaga o caminho da raiz ate a folha p
    void prop(int p) {
        int tam = 1 << (LOG-1);</pre>
        for (int s = LOG; s; s--, tam /= 2) {
            int i = p \gg s;
            if (lazy[i]) {
                poe(2*i, lazy[i], tam);
                poe(2*i+1, lazy[i], tam);
                lazy[i] = 0;
            }
        }
    }
    void build(int n2, int* v) {
        n = n2;
```

```
for (int i = 0; i < n; i++) seg[n+i] = v[i];
        for (int i = n-1; i; i--) seg[i] = junta(seg[2*i],
            seg[2*i+1]);
        for (int i = 0; i < 2*n; i++) lazy[i] = 0;</pre>
    }
    ll query(int a, int b) {
        11 \text{ ret} = 0;
        for (prop(a+=n), prop(b+=n); a \le b; ++a/=2, --b/=2)
           {
            if (a%2 == 1) ret = junta(ret, seg[a]);
            if (b%2 == 0) ret = junta(ret, seg[b]);
        }
        return ret;
    }
    void update(int a, int b, int x) {
        int a2 = a += n, b2 = b += n, tam = 1;
        for (; a <= b; ++a/=2, --b/=2, tam *= 2) {
            if (a\%2 == 1) poe(a, x, tam);
            if (b\%2 == 0) poe(b, x, tam);
        sobe(a2), sobe(b2);
    }
};
     SegTree Beats
// \text{ query}(a, b) - \{\{\min(v[a..b]), \max(v[a..b])\}, \sup(v[a..b])\}
// updatemin(a, b, x) faz com que v[i] \leftarrow min(v[i], x),
// para i em [a, b]
// updatemax faz o mesmo com max, e updatesum soma x
// em todo mundo do intervalo [a, b]
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log^2 (n)) amortizado
// (se nao usar updatesum, fica log(n) amortizado)
#define f first
```

#define s second

```
typedef long long 11;
const 11 LINF = 0x3f3f3f3f3f3f3f3f11;
namespace beats {
    struct node {
        int tam;
        ll sum, lazy; // lazy pra soma
        ll mi1, mi2, mi; // mi = #mi1
        ll ma1, ma2, ma; // ma = #ma1
        node(11 x = 0) {
            sum = mi1 = ma1 = x;
            mi2 = LINF, ma2 = -LINF;
            mi = ma = tam = 1;
            lazy = 0;
        node(const node& 1, const node& r) {
            sum = 1.sum + r.sum, tam = 1.tam + r.tam;
            lazy = 0;
            if (1.mi1 > r.mi1) {
                mi1 = r.mi1, mi = r.mi;
                mi2 = min(1.mi1, r.mi2);
            } else if (l.mi1 < r.mi1) {</pre>
                mi1 = l.mi1, mi = l.mi;
                mi2 = min(r.mi1, l.mi2);
            } else {
                mi1 = 1.mi1, mi = 1.mi+r.mi;
                mi2 = min(1.mi2, r.mi2);
            if (1.ma1 < r.ma1) {
                ma1 = r.ma1, ma = r.ma;
                ma2 = max(1.ma1, r.ma2);
            } else if (1.ma1 > r.ma1) {
                ma1 = l.ma1, ma = l.ma;
                ma2 = max(r.ma1, l.ma2);
            } else {
                ma1 = 1.ma1, ma = 1.ma+r.ma;
                ma2 = max(1.ma2, r.ma2);
            }
        void setmin(ll x) {
            if (x >= ma1) return;
```

```
sum += (x - ma1)*ma:
        if (mi1 == ma1) mi1 = x;
        if (mi2 == ma1) mi2 = x;
        ma1 = x;
    }
    void setmax(ll x) {
        if (x <= mi1) return;</pre>
        sum += (x - mi1)*mi;
        if (ma1 == mi1) ma1 = x;
       if (ma2 == mi1) ma2 = x;
        mi1 = x:
    }
    void setsum(ll x) {
        mi1 += x, mi2 += x, ma1 += x, ma2 += x;
        sum += x*tam:
        lazy += x;
    }
};
node seg[4*MAX];
int n, *v;
node build(int p=1, int l=0, int r=n-1) {
    if (1 == r) return seg[p] = {v[1]};
    int m = (1+r)/2;
    return seg[p] = \{build(2*p, 1, m), build(2*p+1, m+1,
       r)};
void build(int n2, int* v2) {
    n = n2, v = v2;
    build();
void prop(int p, int l, int r) {
    if (1 == r) return:
    for (int k = 0; k < 2; k++) {
        if (seg[p].lazy) seg[2*p+k].setsum(seg[p].lazy);
        seg[2*p+k].setmin(seg[p].ma1);
        seg[2*p+k].setmax(seg[p].mi1);
    seg[p].lazv = 0;
}
pair < pair < 11, 11>, 11> query (int a, int b, int p=1, int
```

```
l=0. int r=n-1) {
    if (b < l or r < a) return \{\{INF, -INF\}, 0\};
    if (a <= l and r <= b) return {{seg[p].mi1,</pre>
        seg[p].ma1}, seg[p].sum};
    prop(p, 1, r);
    int m = (1+r)/2;
    auto L = query(a, b, 2*p, 1, m), R = query(a, b,
        2*p+1, m+1, r);
    return {{min(L.f.f, R.f.f), max(L.f.s, R.f.s)},
       L.s+R.s};
}
node updatemin(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
    if (b < l or r < a or seg[p].ma1 <= x) return seg[p];</pre>
    if (a \le 1 \text{ and } r \le b \text{ and } seg[p].ma2 < x) {
        seg[p].setmin(x);
        return seg[p];
    prop(p, 1, r);
    int m = (1+r)/2;
    return seg[p] = \{updatemin(a, b, x, 2*p, 1, m),
                     updatemin(a, b, x, 2*p+1, m+1, r)};
}
node updatemax(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
    if (b < l or r < a or seg[p].mi1 >= x) return seg[p];
    if (a \le 1 \text{ and } r \le b \text{ and } seg[p].mi2 > x) {
        seg[p].setmax(x);
        return seg[p];
    }
    prop(p, 1, r);
    int m = (1+r)/2;
    return seg[p] = \{updatemax(a, b, x, 2*p, 1, m),
                     updatemax(a, b, x, 2*p+1, m+1, r)};
}
node updatesum(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
    if (b < 1 or r < a) return seg[p];</pre>
    if (a <= 1 and r <= b) {</pre>
        seg[p].setsum(x);
        return seg[p];
    }
```

## 4.10 SegTree Iterativa

```
// Consultas 0-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b]
// Update: muda o valor da posicao p para x
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
int seg[2 * MAX];
int n;
void build() {
    for (int i = n - 1; i; i--) seg[i] = seg[2*i] +
       seg[2*i+1];
}
int query(int a, int b) {
    int ret = 0:
    for (a += n, b += n; a <= b; ++a /= 2, --b /= 2) {
        if (a % 2 == 1) ret += seg[a];
        if (b \% 2 == 0) ret += seg[b];
    return ret;
}
void update(int p, int x) {
    seg[p += n] = x;
    while (p /= 2) seg[p] = seg[2*p] + seg[2*p+1];
}
```

## 4.11 SegTree Persistente

```
// SegTree de soma, update de somar numa posicao
//
// query(a, b, t) retorna a query de [a, b] depois de
// t updates
// update(a, x, t) faz um update v[a]+=x a partir de
// como era depois de t updates
// build - O(n)
// query - O(log(n))
// update - O(log(n))
const int MAX = 3e4+10, UPD = 2e5+10, LOG = 20;
const int MAXS = 4*MAX+UPD*LOG;
namespace perseg {
    11 seg[MAXS], lazy[MAXS];
    int T[UPD], L[MAXS], R[MAXS], cnt, t;
    int n, *v;
    11 build(int p, int l, int r) {
        if (1 == r) return seg[p] = v[1];
        L[p] = cnt++, R[p] = cnt++;
        int m = (1+r)/2;
        return seg[p] = build(L[p], 1, m) + build(R[p], m+1,
           r);
    }
    void build(int n2, int* v2) {
        n = n2, v = v2;
        T[0] = cnt++;
        build(0, 0, n-1);
    }
    ll query(int a, int b, int p, int l, int r) {
        if (b < 1 or r < a) return 0;
        if (a <= l and r <= b) return seg[p];</pre>
        int m = (1+r)/2;
        return query(a, b, L[p], 1, m) + query(a, b, R[p],
           m+1, r);
    }
    11 query(int a, int b, int tt) {
        return query(a, b, T[tt], 0, n-1);
```

### 4.12 SegTree 2D Iterativa

```
// Consultas 0-based
// Um valor inicial em (x, y) deve ser colocado em
   seg[x+n][y+n]
// Query: soma do retangulo ((x1, y1), (x2, y2))
// Update: muda o valor da posicao (x, y) para val
// Nao pergunte como que essa coisa funciona
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
//
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
//
// Complexidades:
// build - O(n^2)
// \text{ query - } O(\log^2(n))
// update - O(log^2(n))
int seg[2*MAX][2*MAX], n;
void build() {
    for (int x = 2*n; x; x--) for (int y = 2*n; y; y--) {
        if (x < n) seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
```

```
if (y < n) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
     }
 }
 int query(int x1, int y1, int x2, int y2) {
     int ret = 0, y3 = y1 + n, y4 = y2 + n;
     for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2)
          for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /=
             2) {
              if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x1][y1];
              if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x1][y2];
              if (x2\%2 == 0 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x2][y1];
              if (x2\%2 == 0 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x2][y2];
         }
     return ret;
 }
 void update(int x, int y, int val) {
     int y2 = y += n;
     for (x += n; x; x /= 2, y = y2) {
          if (x \ge n) seg[x][y] = val;
          else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
          while (y /= 2) seg[x][y] = seg[x][2*y] +
             seg[x][2*y+1];
     }
 }
 4.13 BIT
 // BIT de soma 1-based, v 0-based
// Para mudar o valor da posicao p para x,
 // faca: poe(x - query(p, p), p)
 // l_bound(x) retorna o menor p tal que
 // query(1, p+1) > x (0 based!)
 //
 // Complexidades:
 // build - O(n)
 // poe - O(log(n))
 // query - O(log(n))
// l_bound - O(log(n))
```

```
int n;
int bit[MAX];
int v[MAX];
void build() {
    bit[0] = 0;
    for (int i = 1; i <= n; i++) bit[i] = v[i - 1];</pre>
    for (int i = 1; i <= n; i++) {</pre>
        int j = i + (i & -i);
        if (j <= n) bit[j] += bit[i];</pre>
    }
}
// soma x na posicao p
void poe(int x, int p) {
    for (; p <= n; p += p & -p) bit[p] += x;</pre>
}
// soma [1, p]
int pref(int p) {
    int ret = 0;
    for (; p; p -= p & -p) ret += bit[p];
    return ret;
}
// soma [a, b]
int query(int a, int b) {
    return pref(b) - pref(a - 1);
}
int l_bound(ll x) {
    int p = 0;
    for (int i = MAX2; i+1; i--) if (p + (1 << i) <= n
        and bit [p + (1 << i)] <= x) x -= bit <math>[p += (1 << i)];
    return p;
}
4.14 Order Statistic Set
```

// Funciona do C++11 pra cima

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <class T>
    using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
// para declarar:
ord_set < int > s;
// coisas do set normal funcionam:
for (auto i : s) cout << i << endl;</pre>
cout << s.size() << endl;</pre>
// k-esimo maior elemento O(\log|s|):
// k=0: menor elemento
cout << *s.find_by_order(k) << endl;</pre>
// quantos sao menores do que k O(log|s|):
cout << s.order_of_key(k) << endl;</pre>
// Para fazer um multiset, tem que
// usar ord_set<pair<int, int> > com o
// segundo parametro sendo algo para diferenciar
// os ementos iguais.
// s.order_of_key({k, -INF}) vai retornar o
// numero de elementos < k
4.15 Min queue
// Tudo O(1) amortizado
template < class T> struct minstack {
    stack<pair<T, T> > s;
    void push(T x) {
        if (!s.size()) s.push({x, x});
        else s.push({x, std::min(s.top().second, x)});
    T top() { return s.top().first; }
    T pop() {
        T ans = s.top().first;
        s.pop();
        return ans;
```

```
T size() { return s.size(); }
    T min() { return s.top().second; }
};
template < class T> struct minqueue {
    minstack <T> s1, s2;
    void push(T x) { s1.push(x); }
    void move() {
        if (s2.size()) return;
        while (s1.size()) {
            T x = s1.pop();
            s2.push(x);
        }
    T front() { return move(), s2.top(); }
    T pop() { return move(), s2.pop(); }
    T size() { return s1.size()+s2.size(); }
    T min() {
        if (!s1.size()) return s2.min();
        else if (!s2.size()) return s1.min();
        return std::min(s1.min(), s2.min());
    }
};
```

#### 4.16 DSU Persistente

```
// Complexidades:
// build - O(n)
// find - O(log(n))
// une - O(log(n))

int n, p[MAX], sz[MAX], ti[MAX];

void build() {
   for (int i = 0; i < n; i++) {
      p[i] = i;
      sz[i] = 1;
      ti[i] = -INF;
   }
}</pre>
```

```
int find(int k, int t) {
    if (p[k] == k or ti[k] > t) return k;
   return find(p[k], t);
}
void une(int a, int b, int t) {
    a = find(a); b = find(b);
   if (a == b) return;
    if (sz[a] > sz[b]) swap(a, b);
    sz[b] += sz[a];
    p[a] = b;
    ti[a] = t;
}
4.17 SQRT Tree
// RMQ em O(log log n) com O(n log log n) pra buildar
// Funciona com qualquer operacao associativa
// Tao rapido quanto a sparse table, mas usa menos memoria
// (log log (1e9) < 5, entao a query eh praticamente O(1))
//
// build - O(n log log n)
// query - O(log log n)
namespace sqrtTree {
    int n, *v;
    int pref[4][MAX], sulf[4][MAX], getl[4][MAX],
       entre [4] [MAX], sz [4];
    int op(int a, int b) { return min(a, b); }
   inline int getblk(int p, int i) { return
       (i-getl[p][i])/sz[p]; }
    void build(int p, int l, int r) {
        if (1+1 >= r) return;
        for (int i = 1; i <= r; i++) getl[p][i] = 1;</pre>
        for (int L = 1; L <= r; L += sz[p]) {</pre>
            int R = min(L+sz[p]-1, r);
            pref[p][L] = v[L], sulf[p][R] = v[R];
            for (int i = L+1; i <= R; i++) pref[p][i] =</pre>
               op(pref[p][i-1], v[i]);
```

```
for (int i = R-1; i >= L; i--) sulf[p][i] =
               op(v[i], sulf[p][i+1]);
            build(p+1, L, R);
        for (int i = 0; i <= sz[p]; i++) {</pre>
            int at = entre[p][l+i*sz[p]+i] =
               sulf[p][l+i*sz[p]];
            for (int j = i+1; j <= sz[p]; j++)</pre>
               entre[p][1+i*sz[p]+j] = at =
                    op(at, sulf[p][1+j*sz[p]]);
        }
    void build(int n2, int* v2) {
        n = n2, v = v2;
        for (int p = 0; p < 4; p++) sz[p] = n2 = sqrt(n2);
        build(0, 0, n-1);
    }
    int query(int 1, int r) {
        if (1+1 >= r) return 1 == r ? v[1] : op(v[1], v[r]);
        int p = 0;
        while (getblk(p, 1) == getblk(p, r)) p++;
        int ans = sulf[p][l], a = getblk(p, l)+1, b =
           getblk(p, r)-1;
        if (a \le b) ans = op(ans,
           entre[p][getl[p][1]+a*sz[p]+b]);
        return op(ans, pref[p][r]);
}
      Wavelet Tree
```

#### 4.18

```
// Usa O(sigma + n log(sigma)) de memoria,
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
// count(i, j, x, y) retorna o numero de elementos de
// v[i, j) que pertencem a [x, y]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v[i, j), se ele fosse ordenado
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
```

```
//
// Complexidades:
// build - O(n log(sigma))
// count - O(log(sigma))
// kth - O(log(sigma))
// sum - O(log(sigma))
// sumk - O(log(sigma))
int n, v[MAX];
vector < vector < int > > esq(4*(MAXN-MINN)), pref(4*(MAXN-MINN));
void build(int b = 0, int e = n, int p = 1, int l = MINN,
   int r = MAXN) {
    int m = (1+r)/2; esq[p].push_back(0);
       pref[p].push_back(0);
    for (int i = b; i < e; i++) {
        esq[p].push_back(esq[p].back()+(v[i]<=m));</pre>
        pref[p].push_back(pref[p].back()+v[i]);
    }
    if (1 == r) return;
    int m2 = stable_partition(v+b, v+e, [=](int i){return i
       <= m; \}) - v;
    build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
}
int count(int i, int j, int x, int y, int p = 1, int l =
   MINN, int r = MAXN) {
    if (y < 1 or r < x) return 0;
    if (x <= l and r <= y) return j-i;</pre>
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    return count(ei, ej, x, y, 2*p, 1, m)+count(i-ei, j-ej,
       x, y, 2*p+1, m+1, r);
}
int kth(int i, int j, int k, int p=1, int l = MINN, int r =
   MAXN) {
    if (1 == r) return 1;
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
    return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
}
```

```
int sum(int i, int j, int x, int y, int p = 1, int l = MINN,
   int r = MAXN) {
   if (y < 1 or r < x) return 0;</pre>
    if (x <= 1 and r <= y) return pref[p][j]-pref[p][i];</pre>
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei, j-ej, x,
       y, 2*p+1, m+1, r);
}
int sumk(int i, int j, int k, int p = 1, int l = MINN, int r
   = MAXN)
    if (1 == r) return 1*k;
    int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
    if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);</pre>
    return pref[2*p][ej]-pref[2*p][ei]+sumk(i-ei, j-ej,
       k-(ej-ei), 2*p+1, m+1, r);
}
4.19 Trie
// N deve ser maior ou igual ao numero de nos da trie
// fim indica se alguma palavra acaba nesse no
//
// Complexidade:
// Inserir e conferir string S -> O(|S|)
// usar static trie T
// T.insert(s) para inserir
// T.find(s) para ver se ta
// T.prefix(s) printa as strings
// que tem s como prefixo
struct trie{
    map < char , int > t[MAX+5];
    int p;
    trie(){
        p = 1;
    void insert(string s){
        s += '$';
        int i = 0;
        for (char c : s){
```

```
auto it = t[i].find(c);
        if (it == t[i].end())
            i = t[i][c] = p++;
        else
            i = it->second;
   }
}
bool find(string s){
    s += '$':
    int i = 0;
    for (char c : s){
        auto it = t[i].find(c);
        if (it == t[i].end()) return false;
        i = it->second;
    }
    return true;
}
void prefix(string &l, int i){
    if (t[i].find('$') != t[i].end())
        cout << " " << 1 << endl;
    for (auto p : t[i]){
        l += p.first;
        prefix(l, p.second, k);
        1.pop_back();
    }
}
void prefix(string s){
    int i = 0;
    for (char c : s){
        auto it = t[i].find(c);
        if (it == t[i].end()) return;
        i = it->second;
    }
    int k = 0;
    prefix(s, i, k);
}
```

};

# 5 Papa

#### 5.1 BIT Persistente

```
#include <bits/stdc++.h>
using namespace std;
typedef long long int 11;
const 11 LINF = 0x3f3f3f3f3f3f3f3f11;
#define MAXN 100010
vector < pair < int , ll > > FT [MAXN];
int n;
void clear()
    for(int i=1;i<=n;i++)</pre>
        FT[i].clear();
        FT[i].push_back({-1,0});
}
void add(int i,int v,int time)
    for(;i<=n;i+=i&(-i))</pre>
        11 last=FT[i].back().second;
        FT[i].push_back({time,last+v});
    }
}
ll get(int i,int time)
{
    ll ret=0;
    for(;i>0;i-=i&(-i))
        int pos = upper_bound(FT[i].begin(),FT[i].end(),
                     make_pair(time,LINF))-FT[i].begin()-1;
        ret+=FT[i][pos].second;
    return ret;
}
```

```
11 getRange(int a,int b,int time)
    return get(b,time)-get(a-1,time);
}
     Baby step Giant step
// Resolve Logaritmo Discreto a^x = b mod m, m primo em
   0(sqrt(n)*hash(n))
// Meet In The Middle, decompondo x = i * ceil(sqrt(n)) - j,
   i,j<=ceil(sqrt(n))
int babyStep(int a,int b,int m)
    unordered_map < int , int > mapp;
    int sq=sqrt(m)+1;
    11 asq=1;
    for(int i=0; i<sq; i++)</pre>
        asq=(asq*a)%m;
    11 curr=asq;
    for(int i=1; i<=sq; i++)</pre>
        if(!mapp.count(curr))
             mapp[curr]=i;
        curr=(curr*asq)%m;
    }
    int ret=INF;
    curr=b;
    for(int j=0; j<=sq; j++)</pre>
        if (mapp.count(curr))
             ret=min(ret,(int)(mapp[curr]*sq-j));
        curr = (curr * a) %m;
    if(ret<INF) return ret;</pre>
    return -1;
}
int main()
    int a,b,m;
    while(cin>>a>>b>>m,a or b or m)
```

```
int x=babyStep(a,m,b);
         if (x!=-1)
             cout << x << endl;</pre>
         else
             cout << "No Solution" << endl;</pre>
    return 0;
}
5.3 LIS Rec. Resp.
#include < bits / stdc++.h>
using namespace std;
#define sc(a) scanf("%d", &a)
typedef long long int 11;
const int INF = 0x3f3f3f3f;
#define MAXN 100100
int aux[MAXN], endLis[MAXN];
//usar upper_bound se puder >=
vector < int > LisRec(vector < int > v) {
    int n=v.size();
    int lis=0;
    for (int i = 0; i < n; i++){
        int it = lower_bound(aux, aux+lis, v[i]) - aux;
        endLis[i] = it+1;
        lis = max(lis, it+1);
        aux[it] = v[i];
    vector<int> resp;
    int prev=INF;
    for(int i=n-1;i>=0;i--){
        if (endLis[i] == lis && v[i] <= prev) {</pre>
             lis--;
             prev=v[i];
             resp.push_back(i);
        }
    reverse(resp.begin(),resp.end());
    return resp;
```

```
int main()
    int n;
    sc(n);
    vector < int > v(n);
    for(int i=0;i<n;i++)</pre>
        sc(v[i]);
    cout << LisRec(v).size() << endl;</pre>
    return 0;
}
5.4 Aho Corasick
const int N=100010;
const int M=26;
//N= tamanho da trie, M tamanho do alfabeto
int to[N][M], Link[N], fim[N];
int idx = 1;
void add_str(string &s)
    int v = 0;
    for (int i = 0; i < s.size(); i++) {</pre>
        if (!to[v][s[i]]) to[v][s[i]] = idx++;
        v = to[v][s[i]];
    }
    fim[v] = 1;
}
void process()
    queue < int > fila;
    fila.push(0);
    while (!fila.empty()) {
        int cur = fila.front();
        fila.pop();
        int l = Link[cur];
        fim[cur] |= fim[1];
        for (int i = 0; i < M; i++) {</pre>
             if (to[cur][i]) {
                 if (cur != 0) {
                     Link[to[cur][i]] = to[1][i];
```

```
}
                 else
                     Link[to[cur][i]] = 0;
                 fila.push(to[cur][i]);
            }
             else {
                 to[cur][i] = to[1][i];
        }
}
int resolve(string &s)
{
    int v = 0, r = 0;
    for (int i = 0; i < s.size(); i++) {</pre>
        v = to[v][s[i]];
        if (fim[v]) r++, v = 0;
    return r;
}
```

## 6 Problemas

## 6.1 Inversion Count

```
// O(n log(n))
int n;
int v[MAX];

// bit de soma
void poe(int p);
int query(int p);

// converte valores do array pra
// numeros de 1 a n
void conv() {
   vector<int> a;
   for (int i = 0; i < n; i++) a.push_back(v[i]);</pre>
```

```
sort(a.begin(), a.end());
    for (int i = 0; i < n; i++)</pre>
        v[i] = 1 + (lower_bound(a.begin(), a.end(), v[i]) -
           a.begin());
}
long long inv() {
    conv();
    build();
    long long ret = 0;
    for (int i = n - 1; i \ge 0; i - -) {
        ret += query(v[i] - 1);
        poe(v[i]);
    }
    return ret;
}
    Merge Sort Rafael
// Melhor do Brasil, segundo o autor
//
// O(n log(n))
long long merge_sort(int 1, int r, vector<int> &t){
    if (1 >= r) return 0;
    int m = (1+r)/2;
   auto ans = merge_sort(1, m, t) + merge_sort(m+1, r, t);
    static vector < int > aux; if (aux.size() != t.size())
       aux.resize(t.size());
    for (int i = 1; i <= r; i++) aux[i] = t[i];
   int i_l = 1, i_r = m+1, i = 1;
    auto move_1 = [&](){
        t[i++] = aux[i_1++];
    auto move_r = [\&](){
        t[i++] = aux[i_r++];
    };
```

```
while (i \le r){
        if (i_1 > m) move_r();
        else if (i_r > r) move_l();
        elsef
            if (aux[i_1] <= aux[i_r]) move_1();</pre>
            else{
                 move_r();
                 ans += m - i_1 + 1;
            }
        }
    return ans;
}
//inversions to turn r into l
template < typename T > 11 inv_count(vector < T > &1, vector < T >
   &r){
    int n = 1.size();
    map < T , int > occ;
    map<pair<T, int>, int> rk;
    for (int i = 0; i < n; i++)</pre>
        rk[make_pair(l[i], occ[l[i]]++)] = i;
    occ.clear();
    vector < int > v(n);
    for (int i = 0; i < n; i++)</pre>
        v[i] = rk[make_pair(r[i], occ[r[i]]++)];
    return merge_sort(0, n-1, v);
}
6.3 RMQ com Divide and Conquer
// Responde todas as queries em
// O(n log(n))
typedef pair<pair<int, int>, int> iii;
#define f first
#define s second
int n, q, v[MAX];
iii qu[MAX];
```

int ans[MAX], pref[MAX], sulf[MAX];

```
void solve(int l=0, int r=n-1, int ql=0, int qr=q-1) {
    if (1 > r or q1 > qr) return;
    int m = (1+r)/2;
    int qL = partition(qu+ql, qu+qr+1, [=](iii x){return
       x.f.s < m; \}) - qu;
    int qR = partition(qu+qL, qu+qr+1, [=](iii x){return
       x.f.f <=m;}) - qu;
    pref[m] = sulf[m] = v[m];
    for (int i = m-1; i >= 1; i--) pref[i] = min(v[i],
       pref[i+1]);
    for (int i = m+1; i <= r; i++) sulf[i] = min(v[i],</pre>
       sulf[i-1]);
    for (int i = qL; i < qR; i++)</pre>
        ans[qu[i].s] = min(pref[qu[i].f.f], sulf[qu[i].f.s]);
    solve(1, m-1, ql, qL-1), solve(m+1, r, qR, qr);
}
6.4 SOS DP
// O(n 2^n)
//iterative version
for(int mask = 0; mask < (1 << N); ++mask){
    dp[mask][-1] = A[mask]; //handle base case separately
       (leaf states)
    for(int i = 0;i < N; ++i){</pre>
        if (mask & (1<<i))</pre>
            dp[mask][i] = dp[mask][i-1] +
                dp[mask^(1<<i)][i-1];</pre>
        else dp[mask][i] = dp[mask][i-1];
    }
    F[mask] = dp[mask][N-1];
}
//memory optimized, super easy to code.
for(int i = 0; i<(1<<N); ++i) F[i] = A[i];</pre>
for(int i = 0; i < N; ++i) for(int mask = 0; mask < (1<<N);
   ++mask){
```

```
if (mask & (1<<i))</pre>
        F[mask] += F[mask^(1<<i)];
}
6.5 LIS2
// O(n log(n))
template < typename T > int lis(vector < T > &v){
    vector < T > ans;
    for (T t : v){
        auto it = upper_bound(ans.begin(), ans.end(), t);
        if (it == ans.end()) ans.push_back(t);
        else *it = t;
    return ans.size()
}
6.6 Convex Hull Trick (Rafael)
// linear
struct CHT {
    int it;
    vector<ll> a, b;
    CHT(): it(0){}
    ll eval(int i, ll x){
        return a[i]*x + b[i];
    bool useless(){
        int sz = a.size();
        int r = sz-1, m = sz-2, 1 = sz-3;
        return (b[1] - b[r])*(a[m] - a[1]) <
            (b[1] - b[m])*(a[r] - a[1]);
    }
    void add(ll A, ll B){
        a.push_back(A); b.push_back(B);
        while (!a.empty()){
```

if ((a.size() < 3) || !useless()) break;</pre>

a.erase(a.end() - 2);

b.erase(b.end() - 2);

```
}
    }
    ll get(ll x){
        it = min(it, int(a.size()) - 1);
        while (it+1 < a.size()){</pre>
            if (eval(it+1, x) > eval(it, x)) it++;
            else break;
        return eval(it, x);
   }
};
     Mininum Enclosing Circle Vasek
// O(n) com alta probabilidade
const long double EPS = 1e-12;
struct pt {
    long double x, y;
    pt() {}
    pt(long double x, long double y) : x(x), y(y) {}
    pt(const pt \& p) : x(p.x), y(p.y) \{ \}
    pt operator + (const pt& p) const { return pt(x+p.x,
       y+p.y); }
   pt operator - (const pt& p) const { return pt(x-p.x,
       y-p.y); }
    pt operator * (long double c) const { return pt(x*c, y*c
       ); }
    pt operator / (long double c) const { return pt(x/c, y/c
       ); }
};
long double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
long double dist2(pt p, pt q) { return dot(p-q, p-q); }
long double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
pt rotate90(pt p) { return pt(p.y, -p.x); }
pt interline(pt a, pt b, pt c, pt d) {
    b = b-a; d = c-d; c = c-a;
    return a+b*cross(c, d)/cross(b, d);
```

```
}
pt center(pt a, pt b, pt c) {
    b = (a+b)/2;
    c = (a+c)/2;
    return interline(b, b+rotate90(a-b), c, c+rotate90(a-c));
}
struct circle {
    pt cen;
    long double r;
    circle() {}
    circle(pt cen, long double r) : cen(cen), r(r) {}
};
bool inside(circle& c, pt& p) {
    return c.r*c.r+1e-9 > dist2(p, c.cen);
}
pt bestof3(pt a, pt b, pt c) {
    if (dot(b-a, c-a) < 1e-9) return (b+c)/2;
    if (dot(a-b, c-b) < 1e-9) return (a+c)/2;
    if (dot(a-c, b-c) < 1e-9) return (a+b)/2;
    return center(a, b, c);
}
circle minCirc(vector<pt> v) {
    int n = v.size();
    random_shuffle(v.begin(), v.end());
    pt p = pt(0, 0);
    circle ret = circle(p, 0);
    for (int i = 0; i < n; i++) if (!inside(ret, v[i])) {</pre>
        ret = circle(v[i], 0);
        for (int j = 0; j < i; j++) if (!inside(ret, v[j])) {</pre>
            ret = circle((v[i]+v[j])/2, sqrt(dist2(v[i],
               v[i]))/2);
            for (int k = 0; k < j; k++) if (!inside(ret,
               v[k])) {
                p = bestof3(v[i], v[j], v[k]);
                ret = circle(p, sqrt(dist2(p, v[i])));
            }
        }
```

```
return ret;
6.8 Nim
// Calcula movimento otimo do jogo classico de Nim
// Assume que o estado atual eh perdedor
// Funcao move retorna um par com a pilha (0 indexed)
// e quanto deve ser tirado dela
// XOR deve estar armazenado em x
// Para mudar um valor, faca insere(novo_valor),
// atualize o XOR e mude o valor em v
// MAX2 = teto do log do maior elemento
// possivel nas pilhas
//
// O(log(n)) amortizado
int v[MAX], n, x;
stack<int> pi[MAX2];
void insere(int p) {
    for (int i = 0; i < MAX2; i++) if (v[p] & (1 << i))
       pi[i].push(p);
}
pair<int, int> move() {
    int bit = 0; while (x >> bit) bit++; bit--;
    // tira os caras invalidos
    while ((v[pi[bit].top()] & (1 << bit)) == 0)</pre>
       pi[bit].pop();
    int cara = pi[bit].top();
    int tirei = v[cara] - (x^v[cara]);
    v[cara] -= tirei;
    insere(cara);
```

return make\_pair(cara, tirei);

```
// Acha o movimento otimo baseado
// em v apenas
//
// O(n)

pair<int, int> move() {
   int x = 0;
   for (int i = 0; i < n; i++) x ^= v[i];

   for (int i = 0; i < n; i++) if ((v[i]^x) < v[i])
        return make_pair(i, v[i] - (v[i]^x));
}</pre>
```

### 6.9 Distinct Range Query com Update

```
// build - O(n log^2(n))
// \text{ query - O(log^2(n))}
// update - 0(log^2(n))
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <class T>
    using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
int v[MAX], n, nxt[MAX], prv[MAX];
map<int, set<int> > ocor;
namespace seg {
    ord_set < ii > seg [4*MAX];
    void build(int p=1, int l=0, int r=n-1) {
        if (l == r) return (void)seg[p].insert({nxt[1], 1});
        int m = (1+r)/2;
        build(2*p, 1, m), build(2*p+1, m+1, r);
        for (ii i : seg[2*p]) seg[p].insert(i);
        for (ii i : seg[2*p+1]) seg[p].insert(i);
    int query(int a, int b, int x, int p=1, int l=0, int
       r=n-1) {
```

```
if (b < 1 \text{ or } r < a) \text{ return } 0;
        if (a <= 1 and r <= b) return
            seg[p].order_of_key({x, -INF});
        int m = (1+r)/2;
        return query(a, b, x, 2*p, 1, m)+query(a, b, x,
            2*p+1, m+1, r);
    void update(int a, int x, int p=1, int l=0, int r=n-1) {
        if (a < 1 or r < a) return;
        seg[p].erase({nxt[a], a});
        seg[p].insert({x, a});
        if (1 == r) return;
        int m = (1+r)/2;
        update(a, x, 2*p, 1, m), update(a, x, 2*p+1, m+1, r);
    }
}
void build() {
    for (int i = 0; i < n; i++) nxt[i] = INF;</pre>
    for (int i = 0; i < n; i++) prv[i] = -INF;</pre>
    vector<ii> t;
    for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
    sort(t.begin(), t.end());
    for (int i = 0; i < n; i++) {</pre>
        if (i and t[i].f == t[i-1].f) prv[t[i].s] = t[i-1].s;
        if (i+1 < n \text{ and } t[i].f == t[i+1].f) nxt[t[i].s] =
            t[i+1].s;
    }
    for (int i = 0; i < n; i++) ocor[v[i]].insert(i);</pre>
    seg::build();
}
void muda(int p, int x) {
    seg::update(p, x);
    nxt[p] = x;
}
int query(int a, int b) {
    return b-a+1 - seg::query(a, b, b+1);
}
```

```
void update(int p, int x) { // mudar valor na pos. p para x
    if (prv[p] > -INF) muda(prv[p], nxt[p]);
    if (nxt[p] < INF) prv[nxt[p]] = prv[p];</pre>
    ocor[v[p]].erase(p);
    if (!ocor[x].size()) {
        muda(p, INF);
        prv[p] = -INF;
    } else if (*ocor[x].rbegin() < p) {</pre>
        int i = *ocor[x].rbegin();
        prv[p] = i;
        muda(p, INF);
        muda(i, p);
    } else {
        int i = *ocor[x].lower_bound(p);
        if (prv[i] > -INF) {
            muda(prv[i], p);
            prv[p] = prv[i];
        } else prv[p] = -INF;
        prv[i] = p;
        muda(p, i);
    v[p] = x; ocor[x].insert(p);
}
6.10 LIS1
// Calcula uma LIS
// Para ter o tamanho basta fazer lis().size()
// Implementacao do algotitmo descrito em:
// https://goo.gl/HiFkn2
// O(n log(n))
const int INF = 0x3f3f3f3f;
int n, v[MAX];
vector < int > lis() {
    int I[n + 1], L[n];
```

```
// pra BB funfar bacana
    I[0] = -INF;
    for (int i = 1; i <= n; i++) I[i] = INF;</pre>
    for (int i = 0; i < n; i++) {</pre>
        // BB
        int 1 = 0, r = n;
        while (1 < r) {
            int m = (1 + r) / 2;
            if (I[m] >= v[i]) r = m;
            else 1 = m + 1;
        }
        // ultimo elemento com tamanho l eh v[i]
        I[1] = v[i]:
        // tamanho da LIS terminando com o
        // elemento v[i] eh l
        L[i] = 1;
    }
    // reconstroi LIS
    vector<int> ret;
    int m = -INF, p;
    for (int i = 0; i < n; i++) if (L[i] > m) {
        m = L[i];
        p = i;
    ret.push_back(v[p]);
    int last = m;
    while (p--) if (L[p] == m - 1) {
        ret.push_back(v[p]);
        m = L[p];
    }
    reverse(ret.begin(), ret.end());
    return ret:
6.11 Mo algorithm - distinct values
// O(s*n*f + q*(n/s)*f) optimize over s, insert/erase = O(f)
```

```
// for s = sqrt(n), O((n+q)*sqrt(n)*f)
```

}

```
const int MAX = 3e4+10;
const int SQ = sqrt(MAX);
int v[MAX];
int ans, freq[MAX];
void insert(int p){
    int o = v[p];
    freq[o]++;
    ans += (freq[o] == 1);
}
void erase(int p){
    int o = v[p];
    ans -= (freq[o] == 1);
    freq[o]--;
}
vector<int> MO(vector<ii> &q){
    ans = 0;
    memset(freq, 0, sizeof freq);
    int m = q.size();
    vector < int > ord(m), ret(m);
    iota(ord.begin(), ord.end(), 0);
    sort(ord.begin(), ord.end(), [&](int 1, int r){
        int sl = q[1].first/SQ;
        int sr = q[r].first/SQ;
        if (sl != sr) return sl < sr;</pre>
        return q[1].second < q[r].second;</pre>
    }):
    int 1 = 0, r = 0;
    insert(0);
    for (int i : ord){
        int ql, qr;
        tie(ql, qr) = q[i];
        while (r < qr) insert(++r);</pre>
        while (1 > q1) insert(--1);
        while (1 < q1) erase(1++);</pre>
        while (r > qr) erase(r--);
        ret[i] = ans;
    }
```

```
return ret:
}
6.12 Distinct Range Query
// build - O(n (log n + log(sigma)))
// query - O(log(sigma))
int v[MAX], n, nxt[MAX];
namespace wav {
    vector < vector < int > > esq(4*(1+MAXN-MINN));
    void build(int b = 0, int e = n, int p = 1, int l =
       MINN, int r = MAXN) {
        if (1 == r) return;
        int m = (1+r)/2; esq[p].push_back(0);
        for (int i = b; i < e; i++)</pre>
            esq[p].push_back(esq[p].back()+(nxt[i]<=m));</pre>
        int m2 = stable_partition(nxt+b, nxt+e, [=](int
           i) {return i <= m;}) - nxt;
        build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
    }
    int count(int i, int j, int x, int y, int p = 1, int l =
       MINN, int r = MAXN) {
        if (y < 1 \text{ or } r < x) \text{ return } 0;
        if (x <= l and r <= y) return j-i;</pre>
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        return count(ei, ej, x, y, 2*p, 1, m)+count(i-ei,
           j-ej, x, y, 2*p+1, m+1, r);
    }
}
void build() {
    for (int i = 0; i < n; i++) nxt[i] = MAXN+1;
    vector<ii> t;
    for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
    sort(t.begin(), t.end());
    for (int i = 0; i < n-1; i++) if (t[i].f == t[i+1].f)
       nxt[t[i].s] = t[i+1].s;
```

```
wav::build();
}
int query(int a, int b) {
    return wav::count(a, b+1, b+1, MAXN+1);
}
```

## 6.13 Area da Uniao de Retangulos

```
// O(n log(n))
const int MAX = 1e5+10;
namespace seg {
    pair < int , 11 > seg [4*MAX];
    ll lazy[4*MAX], *v;
    int n;
    pair<int, ll> merge(pair<int, ll> l, pair<int, ll> r){
        if (1.second == r.second) return {1.first+r.first,
           1.second};
        else if (l.second < r.second) return l;</pre>
        else return r;
    }
    pair < int, ll > build(int p=1, int l=0, int r=n-1) {
        lazy[p] = 0;
        if (1 == r) return seg[p] = {1, v[1]};
        int m = (1+r)/2;
        return seg[p] = merge(build(2*p, 1, m), build(2*p+1,
           m+1, r));
    }
    void build(int n2, l1* v2) {
        n = n2, v = v2;
        build();
    }
    void prop(int p, int l, int r) {
        seg[p].second += lazy[p];
        if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] +=
           lazy[p];
        lazy[p] = 0;
    }
    pair < int, ll > query (int a, int b, int p=1, int l=0, int
```

```
r=n-1) {
        prop(p, 1, r);
        if (a <= l and r <= b) return seg[p];</pre>
        if (b < 1 or r < a) return {0, LINF};</pre>
        int m = (1+r)/2;
        return merge(query(a, b, 2*p, 1, m), query(a, b,
            2*p+1, m+1, r));
    }
    pair < int , ll > update(int a, int b, int x, int p=1, int
       1=0, int r=n-1) {
        prop(p, 1, r);
        if (a <= 1 and r <= b) {
            lazy[p] += x;
            prop(p, 1, r);
            return seg[p];
        if (b < l or r < a) return seg[p];</pre>
        int m = (1+r)/2;
        return seg[p] = merge(update(a, b, x, 2*p, 1, m),
                 update(a, b, x, 2*p+1, m+1, r));
    }
};
11 seg_vec[MAX];
ll area_sq(vector<pair<ii, ii>> &sq){
    vector < pair < ii, ii >> up;
    for (auto it : sq){
        int x1, y1, x2, y2;
        tie(x1, y1) = it.first;
        tie(x2, y2) = it.second;
        up.push_back({{x1+1, 1}, {y1, y2}});
        up.push_back(\{\{x2+1, -1\}, \{y1, y2\}\}\});
    }
    sort(up.begin(), up.end());
    memset(seg_vec, 0, sizeof seg_vec);
    11 H_MAX = MAX;
    seg::build(H_MAX-1, seg_vec);
    auto it = up.begin();
    11 \text{ ans} = 0;
    while (it != up.end()){
        11 L = (*it).first.first;
```

```
while (it != up.end() && (*it).first.first == L){
            int x, inc, y1, y2;
            tie(x, inc) = it->first;
            tie(y1, y2) = it->second;
            seg::update(y1+1, y2, inc);
            it++;
        }
        if (it == up.end()) break;
        11 R = (*it).first.first;
        11 W = R-L;
        auto jt = seg::query(0, H_MAX-1);
        11 H = H_MAX - 1;
        if (jt.second == 0) H -= jt.first;
        ans += W*H:
    return ans;
}
6.14 Area Maxima de Histograma
// Assume que todas as barras tem largura 1,
// e altura dada no vetor v
//
// O(n)
typedef long long 11;
11 area(vector<int> v) {
    11 \text{ ret} = 0;
    stack<int> s;
    // valores iniciais pra dar tudo certo
    v.insert(v.begin(), -1);
    v.insert(v.end(), -1);
    s.push(0);
    for(int i = 0; i < (int) v.size(); i++) {</pre>
        while (v[s.top()] > v[i]) {
            11 h = v[s.top()]; s.pop();
            ret = max(ret, h * (i - s.top() - 1));
        }
```

s.push(i);

```
}
    return ret;
}
6.15 Mo algorithm - DQUERY path on trees
// https://codeforces.com/blog/entry/43230
// https://www.spoj.com/problems/COT2/
// (s*2*n*f + q*(2*n/s)*f) optimize over s, insert/erase =
   O(f)
// for s = sqrt(n), O((n+q)*sqrt(n)*f)
vector < int > g[MAX];
namespace LCA { ... }
const int MAX = 40010;
const int SQ = 316;
int w[MAX];
int st[MAX], en[MAX], hst[2*MAX];
int v[2*MAX];
int ans, freq[MAX], freqv[MAX];
void dfs(int i, int p, int &t){
    v[t] = i;
    st[i] = t++;
    for (int j : g[i]){
       if (j == p) continue;
        dfs(j, i, t);
    }
    v[t] = i;
    en[i] = t++;
}
void update(int o){
   if (freqv[o] == 1){//insert w[o]
        ans += (freq[w[o]] == 0);
```

```
freq[w[o]]++;
    if (freqv[o] != 1){//erase w[o]
        ans -= (freq[w[o]] == 1);
        freq[w[o]]--;
}
void insert(int p){
    int o = v[p];
    freqv[o]++;
    update(o);
}
void erase(int p){
    int o = v[p];
    freqv[o]--;
    update(o);
}
vector<tuple<int, int, int>> make_queries(vector<ii>> &q_){
    LCA::build(0);//any LCA alg works
    vector<tuple<int, int, int>> q;
    for (auto &it : q_){
        int 1, r;
        tie(1, r) = it;
        if (st[r] < st[l]) swap(l, r);</pre>
        int p = LCA::lca(l, r);
        int init = (p == 1) ? st[1] : en[1];
        q.push_back({init, st[r], st[p]});
    }
    return q;
}
vector<int> MO(vector<ii> &q_){
    int t = 0;
    dfs(0, -1, t);
    auto q = make_queries(q_);
    ans = 0;
    memset(freq, 0, sizeof freq);
    int m = q.size();
    vector < int > ord(m), ret(m);
```

```
iota(ord.begin(), ord.end(), 0);
    sort(ord.begin(), ord.end(), [&](int 1, int r){
        int sl = get < 0 > (q[1])/SQ;
        int sr = get < 0 > (q[r])/SQ;
        if (sl != sr) return sl < sr;</pre>
        return get<1>(q[1]) < get<1>(q[r]);
    });
    int 1 = 0, r = 0;
    insert(0);
    for (int i : ord){
        int ql, qr, qp;
        tie(ql, qr, qp) = q[i];
        while (r < qr) insert(++r);</pre>
        while (1 > q1) insert(--1);
        while (1 < q1) erase(1++);
        while (r > qr) erase(r--);
        if (qp < 1 || qp > r){
            //lca out of range
            insert(qp);
            ret[i] = ans;
            erase(qp);
        }
        else ret[i] = ans;
    }
    return ret;
}
6.16
      Mininum Enclosing Circle
// O(n) com alta probabilidade
const double EPS = 1e-12;
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
struct pt {
    double x, y;
    pt(double x_{-} = 0, double y_{-} = 0) : x(x_{-}), y(y_{-}) {}
    pt operator + (const pt& p) const { return pt(x+p.x,
```

```
y+p.y); }
    pt operator - (const pt& p) const { return pt(x-p.x,
       y-p.y); }
    pt operator * (double c) const { return pt(x*c, y*c); }
    pt operator / (double c) const { return pt(x/c, y/c); }
};
double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
pt center(pt p, pt q, pt r) {
    pt a = p-r, b = q-r;
    pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
    return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x, b.x),
       c)) / cross(a, b);
}
struct circle {
    pt cen;
    double r;
    circle(pt cen_, double r_) : cen(cen_), r(r_) {}
    circle(pt a, pt b, pt c) {
        cen = center(a, b, c);
        r = dist(cen, a);
    bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
};
circle minCirc(vector<pt> v) {
    shuffle(v.begin(), v.end(), rng);
    circle ret = circle(pt(0, 0), 0);
    for (int i = 0; i < v.size(); i++) if</pre>
       (!ret.inside(v[i])) {
        ret = circle(v[i], 0);
        for (int j = 0; j < i; j++) if (!ret.inside(v[j])) {</pre>
            ret = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
            for (int k = 0; k < j; k++) if
               (!ret.inside(v[k]))
                ret = circle(v[i], v[j], v[k]);
    }
```

```
return ret;
}
```

## 6.17 Min fixed range

```
// https://codeforces.com/contest/1195/problem/E
// O(n)
// ans[i] = min_{0} <= j < k v[i+j]
vector < int > min_k (vector < int > &v, int k) {
    int n = v.size();
    deque < int > d;
    auto put = [&](int i){
        while (!d.empty() && v[d.back()] > v[i])
             d.pop_back();
        d.push_back(i);
    };
    for (int i = 0; i < k-1; i++)</pre>
        put(i);
    vector < int > ans (n-k+1);
    for (int i = 0; i < n-k+1; i++){</pre>
        put(i+k-1);
        while (i > d.front()) d.pop_front();
        ans[i] = v[d.front()];
    }
    return ans;
}
```

#### 6.18 Conectividade Dinamica

```
// Offline com Divide and Conquer e
// DSU com rollback
// O(n log^2(n))

typedef pair<int, int> T;

namespace data {
  int n, ans;
  int p[MAX], sz[MAX];
  stack<int> S;
```

```
void build(int n2) {
        n = n2;
        for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;</pre>
        ans = n:
    int find(int k) {
        while (p[k] != k) k = p[k];
        return k;
    }
    void add(T x) {
        int a = x.first, b = x.second;
        a = find(a), b = find(b);
        if (a == b) return S.push(-1);
        ans --:
        if (sz[a] > sz[b]) swap(a, b);
        S.push(a);
        sz[b] += sz[a];
        p[a] = b;
    int query() {
        return ans;
    void rollback() {
        int u = S.top(); S.pop();
        if (u == -1) return;
        sz[p[u]] -= sz[u];
        p[u] = u;
        ans++;
    }
};
int ponta[MAX]; // outra ponta do intervalo ou -1 se for
   auerv
int ans[MAX], n, q;
T qu[MAX];
void solve(int l = 0, int r = q-1) {
    if (1 >= r) {
        ans[1] = data::query(); // agora a estrutura ta certa
    }
```

```
int m = (1+r)/2, qnt = 1;
    for (int i = m+1; i <= r; i++) if (ponta[i]+1 and
       ponta[i] < 1)</pre>
        data::add(qu[i]), qnt++;
    solve(1, m);
    while (--qnt) data::rollback();
    for (int i = 1; i <= m; i++) if (ponta[i]+1 and ponta[i]
        data::add(qu[i]), qnt++;
    solve(m+1, r);
    while (qnt--) data::rollback();
}
6.19 Points Inside Polygon
// Encontra quais pontos estao
// dentro de um poligono simples nao convexo
// o poligono tem lados paralelos aos eixos
// Pontos na borda estao dentro
// Pontos podem estar em ordem horaria ou anti-horaria
//
// O(n log(n))
#define f first
#define s second
#define pb push_back
typedef long long 11;
typedef pair<int, int> ii;
const ll N = 1e9+10;
const int MAX = 1e5+10;
int ta[MAX];
namespace seg {
    unordered_map<11, int> seg;
    int query(int a, int b, ll p, ll l, ll r) {
        if (b < 1 \text{ or } r < a) \text{ return } 0;
        if (a <= l and r <= b) return seg[p];</pre>
        11 m = (1+r)/2;
        return query(a, b, 2*p, 1, m)+query(a, b, 2*p+1,
           m+1, r);
```

```
}
    int query(ll p) {
        return query(0, p+N, 1, 0, 2*N);
    int update(ll i, int x, ll p, ll l, ll r) {
        if (i < l or r < i) return seg[p];</pre>
        if (1 == r) return seg[p] += x;
        11 m = (1+r)/2;
        return seg[p] = update(i, x, 2*p, 1, m)+update(i, x,
           2*p+1, m+1, r);
    void update(ll a, ll b, int x) {
        if (a > b) return;
        update(a+N, x, 1, 0, 2*N);
        update(b+N+1, -x, 1, 0, 2*N);
    }
};
void pointsInsidePol(vector<ii>>& pol, vector<ii>>& v) {
    vector<pair<int, pair<int, ii> > ev; // {x, {tipo, {a,
       b}}}
    // -1: poe ; id: query ; 1e9: tira
    for (int i = 0; i < v.size(); i++)</pre>
        ev.pb({v[i].f, {i, {v[i].s, v[i].s}}});
    for (int i = 0; i < pol.size(); i++) {</pre>
        ii u = pol[i], v = pol[(i+1)\%pol.size()];
        if (u.s == v.s) {
            ev.pb(\{min(u.f, v.f), \{-1, \{u.s, u.s\}\}\});
            ev.pb({max(u.f, v.f), {N, {u.s, u.s}}});
            continue;
        }
        int t = N;
        if (u.s > v.s) t = -1;
        ev.pb({u.f, {t, {min(u.s, v.s)+1, max(u.s, v.s)}}});
    }
    sort(ev.begin(), ev.end());
    for (int i = 0; i < v.size(); i++) ta[i] = 0;</pre>
    for (auto i : ev) {
        pair < int , ii > j = i.s;
        if (j.f == -1) seg::update(j.s.f, j.s.s, 1);
        else if (j.f == N) seg::update(j.s.f, j.s.s, -1);
```

```
else if (seg::query(j.s.f)) ta[j.f] = 1; // ta dentro
}
```

# 7 Strings

}

## 7.1 Algoritmo Z

```
// Complexidades:
// z - O(|s|)
// \text{ match - } O(|s| + |p|)
vector<int> get_z(string s) {
    int n = s.size();
    vector < int > z(n, 0);
    // intervalo da ultima substring valida
    int 1 = 0, r = 0;
    for (int i = 1; i < n; i++) {</pre>
        // estimativa pra z[i]
        if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
        // calcula valor correto
        while (i + z[i] < n \text{ and } s[z[i]] == s[i + z[i]])
            z[i]++:
        // atualiza [l, r]
        if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
    }
    return z;
}
// quantas vezes p aparece em s
int match(string s, string p) {
    int n = s.size(), m = p.size();
    vector < int > z = get_z(p + s);
    int ret = 0:
    for (int i = m; i < n + m; i++)</pre>
        if (z[i] >= m) ret++;
```

```
return ret;
}
```

## 7.2 String hashing

```
// String deve ter valores [1, x]
// p deve ser o menor primo maior que x
// Para evitar colisao: testar mais de um
// mod; so comparar strings do mesmo tamanho
// ex : str_hash<31, 1e9+7> h(s);
      11 \text{ val} = h(10, 20);
// Complexidades:
// build - O(|s|)
// get_hash - 0(1)
typedef long long 11;
template < int P, int MOD> struct str_hash {
    int n;
    string s;
    vector<ll> h, power;
    str_hash(string s_): n(s_.size()), s(s_), h(n), power(n){
        power[0] = 1;
        for (int i = 1; i < n; i++) power[i] = power[i-1]*P</pre>
           % MOD;
        h[0] = s[0];
        for (int i = 1; i < n; i++) h[i] = (h[i-1]*P + s[i])
            % MOD;
    11 operator()(int i, int j){
        if (!i) return h[j];
        return (h[j] - h[i-1]*power[j-i+1] % MOD + MOD) %
            MOD;
};
```

#### 7.3 Automato de Sufixo

```
// Automato que aceita os sufixos de uma string // Todas as funcoes sao lineares
```

```
namespace sam {
    int cur, sz, len[2*MAX], link[2*MAX], acc[2*MAX];
    int nxt[2*MAX][26];
    void add(int c) {
        int at = cur;
        len[sz] = len[cur]+1, cur = sz++;
        while (at != -1 and !nxt[at][c]) nxt[at][c] = cur,
           at = link[at];
        if (at == -1) { link[cur] = 0; return; }
        int q = nxt[at][c];
        if (len[q] == len[at]+1) { link[cur] = q; return; }
        int qq = sz++;
        len[qq] = len[at]+1, link[qq] = link[q];
        for (int i = 0; i < 26; i++) nxt[qq][i] = nxt[q][i];</pre>
        while (at !=-1 and nxt[at][c] == q) nxt[at][c] =
           qq, at = link[at];
        link[cur] = link[q] = qq;
    }
    void build(string& s) {
        len[0] = 0, link[0] = -1, sz++;
        for (auto i : s) add(i-'a');
        int at = cur;
        while (at) acc[at] = 1, at = link[at];
    }
    // coisas que da pra fazer:
    11 distinct_substrings() {
        11 \text{ ans} = 0;
        for (int i = 1; i < sz; i++) ans += len[i] -</pre>
           len[link[i]];
        return ans;
    }
    string longest_common_substring(string& S, string& T) {
        build(S):
        int at = 0, 1 = 0, ans = 0, pos = -1;
        for (int i = 0; i < T.size(); i++) {</pre>
            while (at and !nxt[at][T[i]-'a']) at = link[at],
               1 = len[at];
            if (nxt[at][T[i]-'a']) at = nxt[at][T[i]-'a'],
               1++;
```

```
else at = 0, 1 = 0;
            if (1 > ans) ans = 1, pos = i;
        return T.substr(pos-ans+1, ans);
    }
    11 dp[2*MAX];
    11 paths(int i) {
        auto& x = dp[i];
        if (x) return x;
        x = 1;
        for (int j = 0; j < 26; j++) if (nxt[i][j]) x +=
           paths(nxt[i][j]);
        return x;
    void kth_substring(int k, int at=0) { // k=1 : menor
       substring lexicog.
        for (int i = 0; i < 26; i++) if (k and nxt[at][i]) {
            if (paths(nxt[at][i]) >= k) {
                cout << char('a'+i);</pre>
                kth_substring(k-1, nxt[at][i]);
                return;
            }
            k -= paths(nxt[at][i]);
        }
    }
};
7.4 Suffix Array
// kasai recebe o suffix array e calcula lcp[i],
// o lcp entre s[sa[i],...,n-1] e s[sa[i+1],..,n-1]
//
// Complexidades:
// suffix_array - O(n log(n))
// kasai - O(n)
vector<int> suffix_array(string s) {
```

s += "\$";

int n = s.size(), N = max(n, 260);

for(int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];</pre>

vector < int > sa(n), ra(n);

```
for (int k = 0; k < n; k ? k *= 2 : k++) {
        vector < int > nsa(sa), nra(n), cnt(N);
        for(int i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n)%n,
            cnt[ra[i]]++;
        for(int i = 1; i < N; i++) cnt[i] += cnt[i-1];</pre>
        for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]]] =
            nsa[i];
        for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r +=</pre>
           ra[sa[i]] !=
            ra[sa[i-1]] or ra[(sa[i]+k)%n] !=
                ra[(sa[i-1]+k)%n];
        ra = nra;
    }
    return vector < int > (sa.begin()+1, sa.end());
}
vector<int> kasai(string s, vector<int> sa) {
    int n = s.size(), k = 0;
    vector<int> ra(n), lcp(n);
    for (int i = 0; i < n; i++) ra[sa[i]] = i;
    for (int i = 0; i < n; i++, k -= !!k) {
        if (ra[i] == n-1) { k = 0; continue; }
        int j = sa[ra[i]+1];
        while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
        lcp[ra[i]] = k;
    }
    return lcp;
}
     Suffix Array Rafael
// O(n log^2(n))
struct suffix_array{
    string &s;
    int n;
    vector<int> p, r, aux, lcp;
```

seg\_tree < int , min\_el > st;
suffix\_array(string &s):

```
s(s), n(s.size()), p(n), r(n), aux(n), lcp(n){
        for (int i = 0; i < n; i++){</pre>
            p[i] = i;
            r[i] = s[i];
        }
        auto rank = [&](int i){
            if (i >= n) return -i;
            return r[i];
        };
        for (int d = 1; d < n; d *= 2){</pre>
            auto t = [&](int i){
                return make_pair(rank(i), rank(i+d));
            };
            sort(p.begin(), p.end(),
                     [&](int &i, int &j){
                     return t(i) < t(j);</pre>
                );
            aux[p[0]] = 0;
            for (int i = 1; i < n; i++)
                 aux[p[i]] = aux[p[i-1]] + (t(p[i]) >
                    t(p[i-1]));
            for (int j = 0; j < n; j++) r[j] = aux[j];
            if (aux[p[n-1]] == n-1) break;
        }
        int h = 0;
        for (int i = 0; i < n; i++){</pre>
            if (r[i] == n-1){
                lcp[r[i]] = 0;
                 continue;
            int j = p[r[i] + 1];
            while (i + h < n \&\& j + h < n \&\& s[i+h] ==
               s[j+h]) h++;
            lcp[r[i]] = h;
            h = \max(0, h-1);
        st = seg_tree < int, min_el > (&lcp);
int query(int 1, int r){
    return st.query(1, r);
```

```
}
    11 distinct_substrings(){
        11 \text{ ans} = p[0] + 1;
        for (int i = 1; i < n; i++)</pre>
             ans += p[i] - lcp[i-1] + 1;
        return ans;
    }
};
7.6 KMP
// Primeiro chama a funcao process com o padrao
// Depois chama match com (texto, padrao)
// Vai retornar o numero de ocorrencias do padrao
// p eh 1-based
//
// Complexidades:
// process - O(m)
// match - 0(n + m)
// n = |texto| e m = |padrao|
int p[MAX];
void process(string& s) {
    int i = 0, j = -1;
    p[0] = -1;
    while (i < s.size()) {</pre>
        while (j \ge 0 \text{ and } s[i] != s[j]) j = p[j];
        i++, j++;
        p[i] = j;
    }
}
int match(string& s, string& t) {
    process(t);
    int i = 0, j = 0, ans = 0;
    while (i < s.size()) {</pre>
        while (j \ge 0 \text{ and } s[i] != t[j]) j = p[j];
        i++, j++;
        if (j == t.size()) j = p[j], ans++;
    }
    return ans;
```

}

## 8 Extra

#### 8.1 makefile

```
CXX = g++

CXXFLAGS = -fsanitize=address -01 -fno-omit-frame-pointer -g

-Wall -Wshadow -std=c++14 -Wno-unused-result

-Wno-sign-compare

CXXFLAGS = -fsanitize=address, undefined

-fno-sanitize-recover=all -D_GLIBCXX_DEBUG -01

-fno-omit-frame-pointer -g -Wall -Wshadow -Wconversion

-std=c++14 -Wno-unused-result -Wno-sign-compare
```

## 8.2 template.cpp

```
set ts=4 si ai sw=4 number mouse=a
syntax on
```

## 8.4 stress.sh

```
make a a2 gen
for ((i = 1; ; i++)) do
    ./gen $i > in
    ./a < in > out
    ./a2 < in > out2
    if (! cmp -s out out2) then
        echo "--> entrada:"
        cat in
        echo "--> saida1:"
        cat out
        echo "--> saida2:"
        cat out2
        break;
    fi
done
```