

```

class:
    mergeSort curd:
        if len(curd) == 1:
            return curd
        mid = len(curd) // 2
        leftHalf = curd[:mid]
        rightHalf = curd[mid:]
        mergesort(leftHalf)
        mergesort(rightHalf)
        i = j = k = 0
        while i < len(leftHalf) and j < len(rightHalf):
            if leftHalf[i] < rightHalf[j]:
                arr[k] = leftHalf[i]
                i += 1
            else:
                arr[k] = rightHalf[j]
                j += 1
            k += 1
        while i < len(leftHalf):
            arr[k] = leftHalf[i]
            i += 1
            k += 1
        while j < len(rightHalf):
            arr[k] = rightHalf[j]
            j += 1
            k += 1
        arr = [27, 89, 70, 55, 62, 99, 45, 14, 10]
        print("Random list: ", arr)
        mergesort(arr)
        print("In merge sorted list: ", arr)

output:
Random list: [27, 89, 70, 55, 62, 99, 45, 14, 10]
merge sorted list [10, 14, 27, 45, 55, 62, 70, 89, 99]

```

### Practical no: 11

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Aim: To sort n list using merge sort  
 Theory: 2 type quicksort merge sort  
 algorithms. It divides input array in 2 halves and  
 then tries for merging two halves. The merging  
 is key process that assumes that arr[1...m]  
 and merges the two subarray into one. The  
 array is recursive divided in two halves till the size  
 one. The size because is the merge process comes  
 into action and starts merging back till it

Applications.

Mergesort is useful for sorting linked item of n  
 item merge sort accesses data sequentially and  
 the need of random access is low. Inversion count  
 problem  
 used in external sorting.

~~Mergesort is merge efficiently. Quicksort Sort types  
 of list of the data to be sorted can only be  
 efficiently accessed potentially~~

~~in block~~

• empirical research suggests that there is a positive correlation between quality and quantity of teaching.



Specular (Glossy)

0.5 0.5 0.5 0.5 0.5 0.5 0.5 0.5 0.5 0.5

1) Polyamide (nylon)

0.2 0.2 0.2 0.2 0.2 0.2 0.2 0.2 0.2 0.2

0.5  
0.5  
0.5  
0.5

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post order (VLR)

Step 1



Step 2



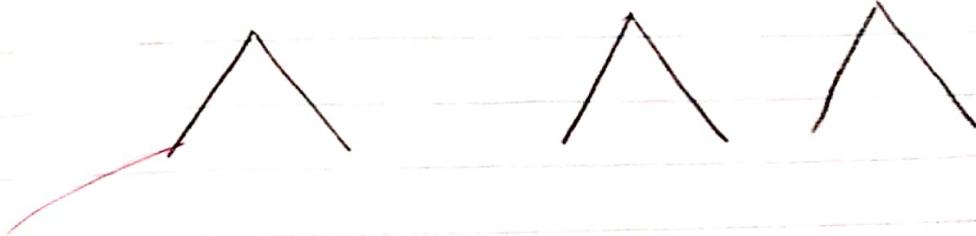
Step 3

post order : (LRV)

Step 1



Step 2



Step 3

steps define Inorder() postorder() and postorder()  
with root argument and use if statement  
that root is None and return that in all

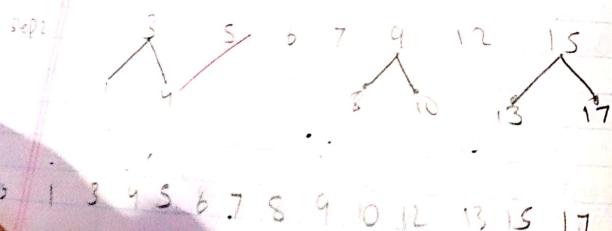
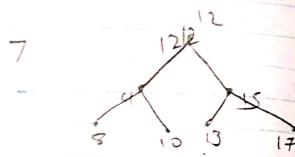
steps In order else statement used for giving that  
condition first left root and then right node

step3 for preorder we have to give condition .in else  
that first root , left and then right node;

step4 for postorder In .else part assign left then right  
and then right node

step5 display its output and input of above algorithm.

Inorder (LVR)



Steps 1 3 4 5 6 7 8 9 10 12 13 15 17

def postorder (root):

if root == None:

return

else:

- postorder (root.left)

- print (root.val)

- postorder (root.right)

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to BSTC

output :

>>> t.add(25)

Root is added successfully at 25

>>> t.add(15)

is node is added to left successfully at 25

>>> t.add(50)

so node is added to right successfully at 25

>>> t.add(10)

10 node is added to left side successfully at 15

>>> t.add(22)

22 node is added to right side successfully at 15

>>> t.add(35)

35 node is added to left side successfully at 25

>>> t.add(70)

70 node is added to right side successfully at 50

>>> t.add(4)

4 node is added to left side successfully at 10

>>> t.add(12)

4 node is added to right side successfully at 10

>>> t.add(18)

18 node is added left side successfully at 22

```

class Node:
    def __init__(self, val):
        self.val = val
        self.left = None
        self.right = None

    def insert(self, val):
        if self.val:
            if val < self.val:
                if self.left == None:
                    self.left = Node(val)
                else:
                    self.left.insert(val)
            else:
                if self.right == None:
                    self.right = Node(val)
                else:
                    self.right.insert(val)
        else:
            self.val = val

    def printTree(self):
        if self.left:
            self.left.printTree()
        print(self.val)
        if self.right:
            self.right.printTree()

def inorder(root):
    if root:
        inorder(root.left)
        print(root.val)
        inorder(root.right)

def preorder(root):
    if root:
        print(root.val)
        preorder(root.left)
        preorder(root.right)

```

### Algorithm:

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- Step 1: Define class node and define init() method with 2 arguments. Initialize value in this method.
- Step 2: Again Define a class BST that is Binary search tree with init() method with self argument and assign d.root is none.
- Step 3: Define add() method for adding a node. Define a variable p that p=node(value).
- Step 4: Use IF statement for checking the condition that if node is less than the main node then put it in leftside.
- Step 5: use while loop for checking if node is less than or greater than the main node and break the loop if it is not satisfying.
- Step 6: use if statement within that else statement for checking that node is greater than main root then put it into rightside.
- Step 7: A for this left subtree and right subtree repeat this method to arrange the node according to binary search tree.

No. Practical - a  
Aim: program based on binary search tree by implementing Inorder, Preorder & Postorder Traversals.

Theory: Binary tree is a tree which supports maximum of 2 children for any node within the tree. Thus any particular node can have either 0 or 1 or 2 children. There is another identity of binary tree that it is ordered such that one child is identified as left child and other as right child.

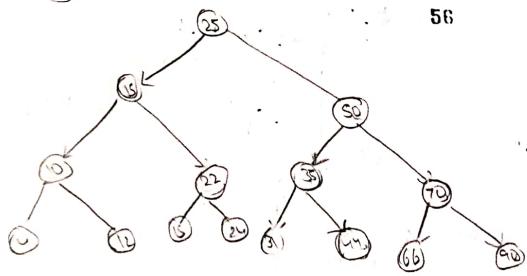
$\rightarrow$  Inorder:  
 i) Traverse its left subtree. The left subtree within might have left and right subtrees.  
 ii) Visit the root node  
 iii) Traverse its right subtree and repeat it

$\rightarrow$  Preorder:  
 i) Visit the root node  
 ii) Traverse its left subtree. The left subtree within might have left and right subtrees.  
 iii) Traverse its right subtree repeat it

$\rightarrow$  Postorder:  
~~i) Traverse its left subtree. The left subtree within might have left and right subtrees.~~  
 ii) Traverse its right subtrees  
 iii) Visit the root node

### \* Binary Search tree

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code:-

class node:

```
def __init__(self, value):
    self.left = None
    self.val = value
    self.right = None.
```

class BST:

def \_\_init\_\_(self):

self.root = None.

def add(self, value):

p: node (value)

if self.root == None:

self.root = p

print("Root is added successfully", p.val)

18  
output  
  >> start . add1 (50)  
  >> start . add2 (60)  
  >> start . add3 (70)  
  >> start . add4 (80)  
  >> start . add5 (90)  
  >>> start . add1B (30)  
  >>> start . add2B (20)  
  >> start . add3B ..  
  >> start . display ()

20  
30  
40  
50  
60  
70  
80

- 55
- Step 8: Now that current is referring to the first node, if we want to access 2nd node of list we can refer it as the next node of the 1st node.
- Step 9: BQZ du 1<sup>st</sup>.node is referred by current so we can transverse to 2<sup>nd</sup> nodes as h.h.next
- Step 10: similarly we can transverse rest of nodes in the linked list using some method by while loop
- Step 11: Our concern now is to find terminating condition for the while loop
- Step 12: The last node in the linked list is referred by link d does not have any next node. In other words the next field of the last node is None
- Step 13: So we can refer to last node of linked list self . s = None
- Step 14: We have to now see how to start traversal in linked list. I how to identify whether we reached the last node of linked list or not
- Steps: Attach the coding or input and output of above algorithm

Step 5: we should not use the head pointer to both  
the entered linked list because the head  
pointer can lead to change's which we  
cannot revert back

Step 6: we may lose the reference to the 1<sup>st</sup> node  
in our linked list. and hence most of our  
unwanted changes to the 1<sup>st</sup> node, we will  
use a temporary node to traverse the entire  
linked list.

Step 7: we will use this temporary node as a copy  
of the node we are currently traversing  
since we are making temporary node a  
copy of current node so data type of the  
temporary node should also be node

```

class linkedlist:
    def __init__(self):
        self.head = None

    def add(self, item):
        newnode = node(item)
        if self.head == None:
            self.head = newnode
        else:
            head = self.head
            while head.next != None:
                head = head.next
            head.next = newnode

    def addAt(self, item):
        newnode = node(item)
        if self.head == None:
            self.head = newnode
        else:
            newnode.next = self.head
            self.head = newnode

    def display(self):
        head = self.head
        while head.next != None:
            print(head.data)
            head = head.next
        print(head.data)

start = linkedlist()

```

### Practical-8

53

Aim: Implementation of single linked list by adding two nodes from last node position

Theory: A linked list is a linear data structure which stores the elements in a node in a linear fashion but no necessarily contiguous. The individual element of the linked list called a node. Node consists of 2 parts @data & Next. If we add / remove one element from the list, all the elements of list has to adjust itself every time we add it is very tedious task so linked list is used to solving this type of problems.

#### Algorithm:

Step 1: Traversing of a linked list means visiting all nodes in the linked list in order to perform some operation on them.

Step 2: The entire linked list can be accessed once the first node of the linked list in turn is referred by the head pointer of the linked list.

Step 3: Thus, an entire linked list can be traversed using the node which is referred by the head pointer of the linked list.

Step 4: Now that we know that we can traverse entire linked list using the head pointer, we should only use it to refer the first node of list only.

Step 7: If token is an operator \*, /, +, - or ^ it will need two operands. pop top 'p' twice. The first pop is second operand and the second pop is the first operand.

Step 8: perform the arithmetic operation, push the result back onto 'm'.

Step 9: When the input expression has been completely processed, the result is on the stack. pop 'top' of 'p' and return the value.

Step 10: print the result of string after the evaluation of prefix.

Step 11: Attach output and input of above algorithm

out put

>>> Evaluated value is: 10

52

```

No
def evaluate(s):
    s = s + " "
    n = len(s)
    stack = []
    for i in range(n):
        if s[i] - '0' >= 0 and s[i] - '9' <= 0:
            stack.append(int(s[i])))
        elif s[i] == '+':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(a) + int(b)))
        elif s[i] == '-':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) - int(a)))
        elif s[i] == '*':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) * int(a)))
        else:
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) / int(a)))
    return stack.pop()

s = "5 8 3 - +"
evaluate(s)
print("evaluated value is:", x)

```

## Practical - 7

51

### Aim post fix Implementation

#### Theory

The postfix expression is free of any parenthesis further we took care of the priorities of the operation in the program. A given postfix expression can easily be evaluated using stack. Reading the expression is always from left to right in Postfix.

#### Algorithm

- Step 1: define evaluate as function then create a empty stack in python.
- Step 2: convert the string to a list by using the string method 'split'.
- Step 3: calculated the length of string and print it.
- Step 4: use for loop to assign the range of string then give condition using If statement.
- Step 5: Scan the taken list from left to right. If token is an operand (convert it from a string to an integer) and push the value onto the 'p'.

```

>>> q.enqueue(1)
'element inserted ... , 1)
>>> q.enqueue(2)
'element inserted ... , 2)
>>> q.enqueue(3)
'element inserted ... , 3)
>>> q.enqueue(4)
queue is full
>>> q.i
[1, 2, 3]
>>> q.dequeue()
1
element deleted ...
>>> q.dequeue()
2
element deleted ...
>>> q.dequeue()
queue is empty
3
element deleted ...
>>> q.dequeue()
queue is empty
>>> q.i
[0, 0, 0]
    
```

10/10/2022

No. A

- Algorithm
- step: define a class queue and assign global variable  
then define init() method with self argument  
in init(), assignor initialize dunder value with  
the help of self argument

- step: Define an empty list and define enqueue() method  
with 1 argument assign dunder length of empty  
list

- step: use if statement that length is equal to record fun  
if else is full or else insert the element in empty  
list or display that queue element added  
successfully and increment by +1

- step: Define dequeue() with self argument under this  
use if statement that front is equal to length  
is the list then display queue is empty or else, give  
that front is at zero and using that delete an elemen  
from front side and increment it by 1

- step: Now call the above () function and give the element  
that has to be added in the empty list by using  
enqueue() and print the list after adding +1 and  
same for deleting and display the list after  
deleting the element from the list

```

class queue:
    global r
    global f
    def __init__(self):
        self.r = 0
        self.f = 0
        self.l = [0, 0, 0]
    def enqueue(self, data):
        n = len(self.l)
        if self.r < n:
            self.l[self.r] = data
            self.r += 1
            print("element inserted...", data)
        else:
            print("queue is full")
    def dequeue(self):
        n = len(self.l)
        if self.f < n:
            print(self.l[self.f])
            self.l[self.f] = 0
            print("element deleted...")
            self.f += 1
        else:
            print("queue is empty")
q = queue()

```

### Practical - 6

49

Aim: Implementing a queue using python list

Theory: Queue is a linear data structure which has 2 references front and rear. Implementing a queue using python list is the simplest as the python list provides inbuilt functions to perform the specified operations of the queue. It is based on the principle that a new element is inserted after rear and element at front is deleted which is at front. In simple terms, a queue can be described as a data structure based on first in first out FIFO principle.

- Queue(): creates a new empty queue

- Enqueue(): Insert an element at the rear of the queue and similar to that of insertion of linked list using tail.

- Dequeue(): Return the element which was at the front. the front is moved to the successive element. A dequeue operation cannot remove element if the queue is empty

output  
enter barcode

Data: S6

>>> s1

[10, 20, 30, 40, 50]

>>> s.pop()

>>> s.pop()

>>> s.l

[10, 20, 30, 40, 50]

>>> s.pop()

>>> s.pop()

>>> s.pop()

>>> s.l

[0, 0, 0, 0, 0]

>>> s.push(60)

>>> s.push(20)

>>> s.push(30)

>>> s.push(40)

>>> s.push(50)

>>> s.l

s.peek()

[10, 20, 30, 40, 50]

MR  
Output

- 1) use if statement to give the condition that range exists when print stack is full
- 2) or else print statement as insert the element into the stack and initialize the value
- 3) push method used to insert the element but pop method used to delete the element from the stack
- 4) If pop method value is less than 1 then return by stack is empty, or else delete the element from stack at topmost position
- 5) first condition checks whether the ref of element is zero while in second case when top is assigned we can be sure that stack is empty.
- a) assign the element value in push method to and print the given value as popped not
- b) return the input and output of above algorithm.

```

A.
class stack:
    global tos
    def __init__(self):
        self.i = [0, 0, 0, 0, 0]
        self.tos = -1
    def push(self, data):
        n = len(self.i)
        if self.tos == n - 1:
            print("Stack is full")
        else:
            self.i[self.tos + 1] = data
            self.tos += 1
    def pop(self):
        if self.tos < 0:
            print("Stack is empty")
        else:
            k = self.i[self.tos]
            print("data is", k)
            self.i[self.tos] = 0
            self.tos -= 1

```

~~S. 1. stack()~~

```

def peek(self):
    if self.tos < 0:
        print("Stack is empty")
    else:
        a = self.i[self.tos]
        print("data is", a)
s2 = stack()

```

### Practical - 5

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Aim: Implementation of stack using python (s)

Theory: A stack is a linear data structure that can be represented in the real world in the form of a physical stack or a pile. The elements in the stack are added or removed only from one position i.e., the topmost position. Thus the stack works on the LIFO (Last IN first OUT) principle as the element that was inserted last is popped first. The operation of adding and removing the element is known as push & pop.

#### Algorithm

- 1) Create a class stack with instance variable item
- 2) Define init method with self argument and initialize the initial value and the initialization on empty list
- 3) Define methods push and pop under the class stack

- A
- 1) As the pivot value right mark becomes less than leftmark we stop the position of rightmark is now the split point
  - 2) pivot value can be exchanged with its content
  - 3) split point and PV (pivot value) is now in place
  - 4) travelling all the items to left of split pivot are less than pivot and all the items to the right are greater than PV. The split point is greater than PV. The left and quick sort can be invoked recursively on the two halves
  - 5) The quicksort function invokes a recursive function quicksort
  - 6) Quicksort begins with some base called as merge sort
  - 7) If length of the list is less than or equal to one, it is sorted
  - 8) If it is greater than it can be partitioned and recursively sorted
  - 9) The partition function implements the process described earlier
  - 10) Display and stick the coding and output of above algorithm

for b in range(0, x):
 b = input('Enter element')
 dlist.append(b)
 h = list(dlist)
 quick(h)
 print(h)

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Output

Enter Range for list S  
 Enter element 4  
 Enter element 3  
 Enter element 2  
 Enter element 1  
 Enter element 8

{1, 2, 3, 4, 8}

P

```

def quickSort(lis):
    helpCant(lis, 0, len(lis)-1)
    def help(first, last):
        if first < last:
            split = part(lis[first], first, last)
            help(lis[first+1:last], first+1, last)
            help(lis[first], 0, first-1)
    def part(lis, first, last):
        pivot = lis[first]
        i = first + 1
        j = last
        done = False
        while not done:
            while i < j and lis[i] <= pivot:
                i += 1
            while i < j and lis[j] >= pivot and i < j:
                j -= 1
            if i < j:
                done = True
            else:
                lis[i], lis[j] = lis[j], lis[i]
                i += 1
                j -= 1
        return j
    import ("random")
    = j

```

#### Practical - 4

45

Aim Implement quick sort to sort the given list

theory: The quick sort is a recursive algorithm based on the divide and conquer technique

##### Algorithm

- ① quick sort first selects a value, which is called pivot value, first element serve as our first pivot value, since we know that first will eventually end up as last in def list.
- ② The partition process will happen next. It will either less than or greater than pivot value.
- ③ partitioning begins by selecting two position markers. Let's call them left mark and right mark at the beginning. We move side with respect to pivot value until also converging to the split point.
- ④ we begin by incrementing until we locate a value that is greater the PV we then decrement rightmark until we find value that is less than items that are out of place with respect to eventual split point.

- 
- to be sorted two to  
one element above should be replaced  
process mentioned above should  
m-1 to get the required result
- stick the output and input of above algorithm  
~~of bubble sort step wise~~

```

for a = list (input ("Enter list: ")):
    for i in range (0, len(a)-1):
        for j in range (0, len(a)-1):
            if a[i] > a [j+1]:
                a [j], a [j+1] = a [j+1], a [j]
print (a)

```

output  
 Input list : 65, 75, 30, 90, 2, 11  
 [2, 11, 30, 65, 75, 90]

### Practical -3

43

Aim : Implementation of Bubble sort program on given list .

Theory: Bubble sort is based on the idea of repeatedly comparing pairs of adjacent elements and they exist in wrong order. This is the simplest form of sorting available. In this we sort the given elements in ascending or descending order by comparing two adjacent elements at a time.

① Bubble sort algorithm start by comparing the first two elements of an array and swapping if necessary.

② If we want the elements of array in ascending order then first element is greater than second then we need to swap the element.

③ If the element is smaller than second then we do not swap the element.

④ Again second and third element are compared go on until last and second last element swapped and swapped.

~~1. To give the range in which element is found in given range~~

Step 3: use for loop in given range if element is found in element not found message is printed

Step 4: Initialize first to 0 and last to length of array. Then use else statement if element not found in range then satisfy the below condition element has to be searched

Step 5: Initialize first to 0 and last to length of array. To be searched is not found then find a middle element ( $m$ )

~~Step 7: Repeat till you found the element is still the input and output of above algorithm~~

## Practical : 2

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Aim: Implement Binary search to find an searched no. in the list.

Theory

Binary Search

Binary Search is also known as half interval search, logarithmic search or binary chop is a search algorithm that finds the position of target value within a sorted need to search an entire list in linear search which is time consuming this can be avoided by using Binary Postion Search

Algorithm

Step 1: Create empty list and assign it to variable list. use for loop, add element in list using of append () method.

Step 2: use sort () method to sort the excepted elements and assign it in increasing order list print the list after sorting

Step 8: Check if input and output of above algorithm

Step 7:- use counter if loop to print this element  
element is not found if this element is accepted  
which is accepted from user is not this

Step 6: use if loop that the element in this list is found along with the element of position  
the selected that the element is equal to the element accepted

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Searched  
3: list (input ("Enter the element"))

```
s: sort ()
print s
print ("Enter the number to be searched")
n: int(input("Enter the number"))
for i in range (len(s)):
    if (a[i] == n):
        print ("number found ! in the position", i)
        break
else:
    print ("number not found")
```

Output

Enter the element : 4, 6, 7, 12, 5, 8

[5, 6, 7, 8, 4, 12]

Enter the number to be searched = 8

number found ! in position 3

39

2) sorted

Step 1:- Create empty list and assign it to a variable  
accept total no of elements to be inserted into the  
list from user, say 'n'

Step 2:- Use for loop for using append () method to  
add n elements in the list. Use sort () method to  
sort the accepted element and assign in  
increasing order in list, then print the list

Step 3:- use If statement to give the range in which  
element not found.

Step 4:- Then use else statement if element is not  
found in range then satisfy the given  
condition

Step 5:- use for loop in range from 0 to the total  
no. of elements to be searched before doing  
this except on search no from user using  
Input statement

I solved  
Algorithm

Step: Create an empty list and assign it to a variable

Step: Accept the total no of elements to be inserted into list from user say 'n'

Step: Use for loop for adding n elements into a list

Step: Print the new list

Step: Accept an element from the user that to be searched in the list

Step: Use for loop in a range from 0 to n-1 to search the element & from the

no of elements to search the element & from the list

Step: Use for loop that if element in the list is equal

to the element accepted from user:

Step: If the element is found then print the statement & with the element position.

Step: Use another for loop to print the list if element is not found if the element given in the list.

Step: Draw the output of given algorithm

answering

s = int(input("Enter the required number"))  
a = [10, 2, 9, 14, 17, 9]

for i in range (len(a)):

if a[i] == s:

print "Required number found in position",

break

if (s) == a[i]:

print "The required number not found"

Output

Linker for required number a  
Required number found in position 2

n

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## Practical no:-1

Aim: Implement linear search to find an item in the list

Theory:

### Linear Search

Linear search is one of the simplest searching algorithm in which targeted searching algorithm in which each item in the list

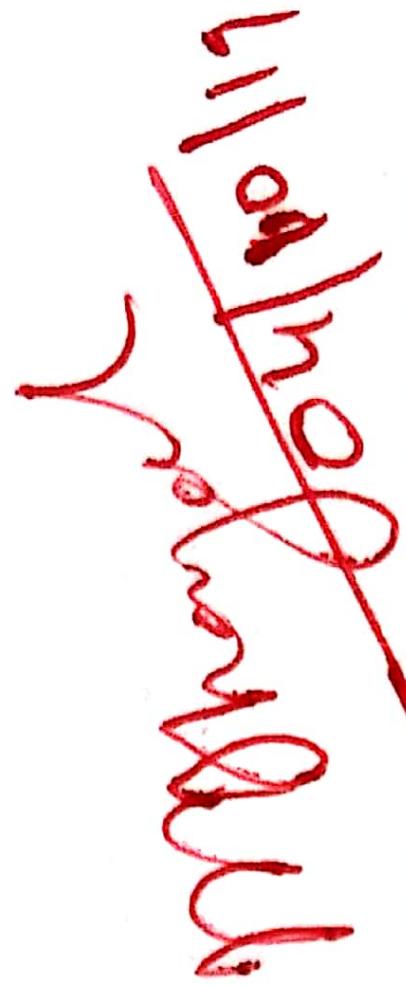
It is the worst searching algorithm with worst case time complexity. It is a brute force approach and other hard in case of list in sequence. A binary search is used which will start by examining the middle term.

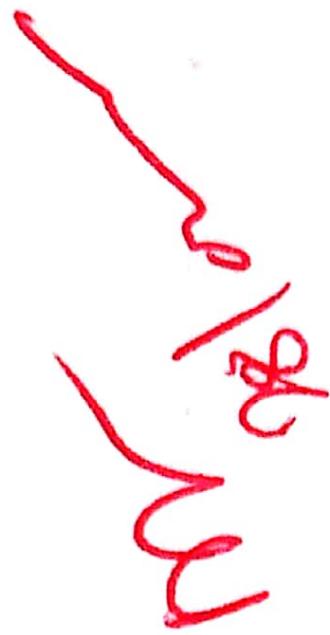
Linear search is a technique to compare each

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9.)	program based on binary Search tree by implementing Inorder, Preorder & postorder Traversal	7/2/20		

**Staff Member's  
Signature**

A red ink signature in cursive script, oriented vertically from bottom-left to top-right. The signature appears to read "Mangal Joshi" followed by "W/O".

A red ink signature in cursive script, oriented vertically from bottom-left to top-right. The signature appears to read "Ranjan Joshi".

**Remarks**

A red ink signature consisting of the letters "D.B.M.S." written in a stylized, handwritten font.

A red ink signature consisting of the letters "O.S." written in a stylized, handwritten font.