

Cryptography and Network Security

NUMBER THEORY



Session Meta Data

Author	Dr T Sree Sharmila
Reviewer	
Version Number	1.0
Release Date	30 June 2018

Revision History

Revision Date	Details	Version no.
		1.0

Agenda

- prime numbers
- Fermat's and Euler's Theorems
- Primality Testing
- Chinese Remainder Theorem
- Discrete Logarithms
- Summary
- Test your understanding
- References


Prime Numbers

- prime numbers only have divisors of 1 and self
 - they cannot be written as a product of other numbers
 - note: 1 is prime, but is generally not of interest
- eg. 2,3,5,7 are prime, 4,6,8,9,10 are not
- prime numbers are central to number theory
- list of prime number less than 200 is:

2 3 5 7 11 13 17 19 23 29 31 37 41 43 47 53 59
61 67 71 73 79 83 89 97 101 103 107 109 113 127
131 137 139 149 151 157 163 167 173 179 181 191
193 197 199

Prime Factorisation

- to **factor** a number n is to write it as a product of other numbers: $n = a \times b \times c$
- note that factoring a number is relatively hard compared to multiplying the factors together to generate the number
- the **prime factorisation** of a number n is when its written as a product of primes
 - eg. $91 = 7 \times 13$; $3600 = 2^4 \times 3^2 \times 5^2$

$$a = \prod_{p \in P} p^{a_p}$$


Relatively Prime Numbers & GCD

- two numbers a , b are **relatively prime** if have **no common divisors** apart from 1
 - eg. 8 & 15 are relatively prime since factors of 8 are 1,2,4,8 and of 15 are 1,3,5,15 and 1 is the only common factor
- conversely can determine the greatest common divisor by comparing their prime factorizations and using least powers
 - eg. $300 = 2^1 \times 3^1 \times 5^2$ $18 = 2^1 \times 3^2$ hence
 $\text{GCD}(18, 300) = 2^1 \times 3^1 \times 5^0 = 6$

Agenda

- prime numbers
- Fermat's and Euler's Theorems
- Primality Testing
- Chinese Remainder Theorem
- Discrete Logarithms
- Summary
- Test your understanding
- References

Fermat's Theorem

- $a^{p-1} \bmod p = 1$
 - where p is prime and $\gcd(a, p) = 1$
- also known as Fermat's Little Theorem
- useful in public key and primality testing

Euler Totient Function $\phi(n)$

- when doing arithmetic modulo n
- **complete set of residues** is: $0 \dots n-1$
- **reduced set of residues** is those numbers (residues) which are relatively prime to n
 - eg for $n=10$,
 - complete set of residues is $\{0,1,2,3,4,5,6,7,8,9\}$
 - reduced set of residues is $\{1,3,7,9\}$
- number of elements in reduced set of residues is called the **Euler Totient Function $\phi(n)$**

Euler Totient Function $\phi(n)$

- to compute $\phi(n)$ need to count number of elements to be excluded
- in general need prime factorization, but
 - for p (p prime) $\phi(p) = p-1$
 - for $p.q$ (p, q prime) $\phi(p.q) = (p-1)(q-1)$
- eg.
 - $\phi(37) = 36$
 - $\phi(21) = (3-1) \times (7-1) = 2 \times 6 = 12$

Euler's Theorem

- a generalisation of Fermat's Theorem
- $a^{\phi(n)} \bmod N = 1$
 - where $\gcd(a, N) = 1$
- eg.
 - $a=3; n=10; \phi(10)=4;$
 - hence $3^4 = 81 = 1 \bmod 10$
 - $a=2; n=11; \phi(11)=10;$
 - hence $2^{10} = 1024 = 1 \bmod 11$

Example

Example 1

What is the value of $\phi(13)$?

Solution

Because 13 is a prime, $\phi(13) = (13 - 1) = 12$.

Example 2

What is the value of $\phi(10)$?

Solution

We can use the third rule: $\phi(10) = \phi(2) \times \phi(5) = 1 \times 4 = 4$, because 2 and 5 are primes.

Example

Example 3

What is the value of $\phi(240)$?

Solution

We can write $240 = 2^4 \times 3^1 \times 5^1$. Then

$$\phi(240) = (2^4 - 2^3) \times (3^1 - 3^0) \times (5^1 - 5^0) = 64$$

Example 4

Can we say that $\phi(49) = \phi(7) \times \phi(7) = 6 \times 6 = 36$?

Solution

No. The third rule applies when m and n are relatively prime. Here $49 = 7^2$. We need to use the fourth rule: $\phi(49) = 7^2 - 7^1 = 42$.

Example

Example 5

What is the value of $\phi(240)$?

Solution

We can write $240 = 2^4 \times 3^1 \times 5^1$. Then

$$\phi(240) = (2^4 - 2^3) \times (3^1 - 3^0) \times (5^1 - 5^0) = 64$$

Example 6

Can we say that $\phi(49) = \phi(7) \times \phi(7) = 6 \times 6 = 36$?

Solution

No. The third rule applies when m and n are relatively prime. Here $49 = 7^2$. We need to use the fourth rule: $\phi(49) = 7^2 - 7^1 = 42$.

Agenda

- prime numbers
- Fermat's and Euler's Theorems
- Primality Testing
- Chinese Remainder Theorem
- Discrete Logarithms
- Summary
- Test your understanding
- References

Primality Testing

- often need to find large prime numbers
- traditionally **sieve** using **trial division**
 - ie. divide by all numbers (primes) in turn less than the square root of the number
 - only works for small numbers
- alternatively can use statistical primality tests based on properties of primes
 - for which all primes numbers satisfy property
 - but some composite numbers, called pseudo-primes, also satisfy the property

Miller Rabin Algorithm

- a test based on Fermat's Theorem
- algorithm is:

TEST (n) is:

1. Find integers $k, q, k > 0, q$ odd, so that $(n-1) = 2^k q$
2. Select a random integer $a, 1 < a < n-1$
3. **if** $a^q \bmod n = 1$ **then** return ("maybe prime");
4. **for** $j = 0$ **to** $k-1$ **do**
 5. **if** $(a^{2^j q} \bmod n = n-1)$
then return(" maybe prime ")
6. return ("composite")

Probabilistic Considerations

- if Miller-Rabin returns “composite” the number is definitely not prime
- otherwise is a prime or a pseudo-prime
- chance it detects a pseudo-prime is $< \frac{1}{4}$
- hence if repeat test with different random a then chance n is prime after t tests is:
 - $\Pr(n \text{ prime after } t \text{ tests}) = 1 - 4^{-t}$
 - eg. for $t=10$ this probability is > 0.99999

Prime Distribution

- prime number theorem states that primes occur roughly every $(\ln n)$ integers
- since can immediately ignore evens and multiples of 5, in practice only need test $0.4 \ln(n)$ numbers of size n before locate a prime
 - note this is only the “average” sometimes primes are close together, at other times are quite far apart

Agenda

- prime numbers
- Fermat's and Euler's Theorems
- Primality Testing
- Chinese Remainder Theorem
- Discrete Logarithms
- Summary
- Test your understanding
- References

CHINESE REMAINDER THEOREM

The Chinese remainder theorem (CRT) is used to solve a set of congruent equations with one variable but different moduli, which are relatively prime, as shown below:

$$\begin{aligned}x &\equiv a_1 \pmod{m_1} \\x &\equiv a_2 \pmod{m_2} \\&\dots \\x &\equiv a_k \pmod{m_k}\end{aligned}$$

CHINESE REMAINDER THEOREM

Example 1

The following is an example of a set of equations with different moduli:

$$\begin{aligned}x &\equiv 2 \pmod{3} \\x &\equiv 3 \pmod{5} \\x &\equiv 2 \pmod{7}\end{aligned}$$

The solution to this set of equations is given in the next section; for the moment, note that the answer to this set of equations is $x = 23$. This value satisfies all equations: $23 \equiv 2 \pmod{3}$, $23 \equiv 3 \pmod{5}$, and $23 \equiv 2 \pmod{7}$.

CHINESE REMAINDER THEOREM

Solution To Chinese Remainder Theorem

1. Find $M = m_1 \times m_2 \times \dots \times m_k$. This is the common modulus.
2. Find $M_1 = M/m_1, M_2 = M/m_2, \dots, M_k = M/m_k$.
3. Find the multiplicative inverse of M_1, M_2, \dots, M_k using the corresponding moduli (m_1, m_2, \dots, m_k) . Call the inverses $M_1^{-1}, M_2^{-1}, \dots, M_k^{-1}$.
4. The solution to the simultaneous equations is

$$x = (a_1 \times M_1 \times M_1^{-1} + a_2 \times M_2 \times M_2^{-1} + \dots + a_k \times M_k \times M_k^{-1}) \bmod M$$

CHINESE REMAINDER THEOREM

Example 2

Find the solution to the simultaneous equations:

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

Solution

We follow the four steps.

1. $M = 3 \times 5 \times 7 = 105$

2. $M_1 = 105 / 3 = 35$, $M_2 = 105 / 5 = 21$, $M_3 = 105 / 7 = 15$

3. The inverses are $M_1^{-1} = 2$, $M_2^{-1} = 1$, $M_3^{-1} = 1$

4. $x = (2 \times 35 \times 2 + 3 \times 21 \times 1 + 2 \times 15 \times 1) \bmod 105 = 23 \bmod 105$

CHINESE REMAINDER THEOREM

Example 3

Find an integer that has a remainder of 3 when divided by 7 and 13, but is divisible by 12.

Solution

This is a CRT problem. We can form three equations and solve them to find the value of x .

$$\begin{aligned}x &= 3 \pmod{7} \\x &= 3 \pmod{13} \\x &= 0 \pmod{12}\end{aligned}$$

If we follow the four steps, we find $x = 276$. We can check that $276 = 3 \pmod{7}$, $276 = 3 \pmod{13}$ and 276 is divisible by 12 (the quotient is 23 and the remainder is zero).

CHINESE REMAINDER THEOREM

Example4

Assume we need to calculate $z = x + y$ where $x = 123$ and $y = 334$, but our system accepts only numbers less than 100.

$$\begin{array}{ll} x \equiv 24 \pmod{99} & y \equiv 37 \pmod{99} \\ x \equiv 25 \pmod{98} & y \equiv 40 \pmod{98} \\ x \equiv 26 \pmod{97} & y \equiv 43 \pmod{97} \end{array}$$

Adding each congruence in x with the corresponding congruence in y gives

$$\begin{array}{ll} x + y \equiv 61 \pmod{99} & \rightarrow z \equiv 61 \pmod{99} \\ x + y \equiv 65 \pmod{98} & \rightarrow z \equiv 65 \pmod{98} \\ x + y \equiv 69 \pmod{97} & \rightarrow z \equiv 69 \pmod{97} \end{array}$$

Now three equations can be solved using the Chinese remainder theorem to find z .
One of the acceptable answers is $z = 457$.

Primitive Roots

- from Euler's theorem have $a^{\phi(n)} \bmod n = 1$
- consider $a^m \bmod n = 1$, $\text{GCD}(a, n) = 1$
 - must exist for $m = \phi(n)$ but may be smaller
 - once powers reach m , cycle will repeat
- if smallest is $m = \phi(n)$ then a is called a **primitive root**
- if p is prime, then successive powers of a "generate" the group $\bmod p$
- these are useful but relatively hard to find

Agenda

- prime numbers
- Fermat's and Euler's Theorems
- Primality Testing
- Chinese Remainder Theorem
- Discrete Logarithms
- Summary
- Test your understanding
- References

Discrete Logarithms or Indices

- the inverse problem to exponentiation is to find the **discrete logarithm** of a number modulo p
- that is to find x where $a^x = b \pmod{p}$
- written as $x = \log_a b \pmod{p}$ or $x = \text{ind}_{a,p}(b)$
- if a is a primitive root then always exists, otherwise may not
 - $x = \log_3 4 \pmod{13}$ (x st $3^x = 4 \pmod{13}$) has no answer
 - $x = \log_2 3 \pmod{13} = 4$ by trying successive powers
- whilst exponentiation is relatively easy, finding discrete logarithms is generally a **hard** problem

Agenda

- prime numbers
- Fermat's and Euler's Theorems
- Primality Testing
- Chinese Remainder Theorem
- Discrete Logarithms
- Summary
- Test your understanding
- References

Summary

- have considered:
 - prime numbers
 - Fermat's and Euler's Theorems
 - Primality Testing
 - Chinese Remainder Theorem
 - Discrete Logarithms

Agenda

- prime numbers
- Fermat's and Euler's Theorems
- Primality Testing
- Chinese Remainder Theorem
- Discrete Logarithms
- Summary
- Test your understanding
- References

Test your understanding

- Define Fermat's theorem.
- Define Euler's theorem.
- Explain CRT with example.
- List out the primality testing techniques.

Agenda

- prime numbers
- Fermat's and Euler's Theorems
- Primality Testing
- Chinese Remainder Theorem
- Discrete Logarithms
- Summary
- Test your understanding
- References

References

1. William Stallings, Cryptography and Network Security, 6th Edition, Pearson Education, March 2013.
2. Charlie Kaufman, Radia Perlman and Mike Speciner, "Network Security", Prentice Hall of India, 2002.