

CS 221 HW 8 - Logic

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$$4a) \quad KB = \{(A \vee B) \rightarrow C, A\}$$

$$\begin{aligned} (A \vee B) \rightarrow C &= \neg(A \vee B) \vee C \\ &= (\neg A \wedge \neg B) \vee C \\ &= (\neg A \vee C) \wedge (\neg B \vee C) \\ &= (A \rightarrow C) \wedge (B \rightarrow C) \end{aligned}$$

As proved by the derivation above, we can introduce a new generalized rule as below:

$$\begin{aligned} (A_1 \vee A_2 \vee \dots \vee A_n) \rightarrow B \\ = (A_1 \rightarrow B) \wedge (A_2 \rightarrow B) \wedge \dots \wedge (A_n \rightarrow B) \end{aligned}$$

Using the new rule above and modus ponens, we can derive C from KB as follows.

$$(A \vee B) \rightarrow C = (A \rightarrow C) \wedge (B \rightarrow C)$$

The above result can be split and put back into KB .

$$KB = \{A \rightarrow C, B \rightarrow C, A\}$$

Using modus ponens we get,

$$\frac{A \rightarrow C, A}{C}$$

~~Thus~~ The KB now is $\{A \rightarrow C, B \rightarrow C, A, C\}$

Thus we derived C .

$$\begin{aligned}
 4b) \quad KB &= \{A \vee B, B \rightarrow C, (A \vee C) \rightarrow D\} \\
 &= \{A \vee B, \neg B \vee C, \neg(A \vee C) \vee D\} \\
 &= \{A \vee B, \neg B \vee C, (\neg A \wedge \neg C) \vee D\} \\
 &= \{A \vee B, \neg B \vee C, (\neg A \vee D) \wedge (\neg C \vee D)\}
 \end{aligned}$$

Resolving A and $\neg A$,

$$= \{B, \neg B \vee C, D \wedge (\neg C \vee D)\}$$

Resolving B and $\neg B$,

$$= \{C, D \wedge (\neg C \vee D)\}$$

Resolving C and $\neg C$,

$$= \{D \wedge D\}$$

$$= \{D\}$$

Thus we derived D from KB using on converting to CNF and using resolution rule.

5b)

Let us consider the finite model below for the original 6 constraints:

$\{1, 2\}$

1's successor is 2

2's successor is 1

$2 > 1 > 2$

All constraints regarding odd/even are satisfied.

Thus the finite model above is consistent for the original set of 6 constraints.

Now if we add the 7th constraint

"A number is not larger than itself", the finite model above is not consistent as the transitive property on $2 > 1 > 2$ would derive $2 > 2$ violating the new 7th constraint.

Generalizing from the finite model above, any finite model satisfying the original 6 constraints can be reduced to a standoff b/w two numbers using the transitive property as " $x > y > x$ ".

Now, enforcing the new 7th constraint on the reduced ~~form~~ form " $x > y > x$ " would violate the consistency.

However, ~~for~~ the 7 constraints can be applied ^{may not} successfully on an infinite model as it ~~cannot~~ be reduced to a form " $x > y > x$ ".