

DATABASES, CATEGORIES AND FUNCTORS

INTRODUCTION TO CATEGORY THEORY

LET'S LOOK AT A DATABASE

Employee	WorksIn	Manager
1	101	2
2	101	3
3	102	3

Department	Secretary
101	2
102	3

LET'S LOOK AT A DATABASE

Employee	WorksIn	Manager
1	101	2
2	101	3
3	102	3

Department	Secretary
101	2
102	3

Schemas

EMPLOYEES

Employee: EmployeeID

WorksIn: DepartmentID

Manager: EmployeeID

DEPARTMENTS

Department: DepartmentID

Secretary: EmployeeID

DATABASE SCHEMA AS A GRAPH

EMPLOYEES

Employee: EmployeeID

WorksIn: DepartmentID

Manager: EmployeeID

DEPARTMENTS

Department: DepartmentID

Secretary: EmployeeID

DATABASE SCHEMA AS A GRAPH

EMPLOYEES

Employee: EmployeeID
WorksIn: DepartmentID
Manager: EmployeeID

DEPARTMENTS

Department: DepartmentID
Secretary: EmployeeID

Each ID column
becomes a node
(vertex)

Employee

Department

DATABASE SCHEMA AS A GRAPH

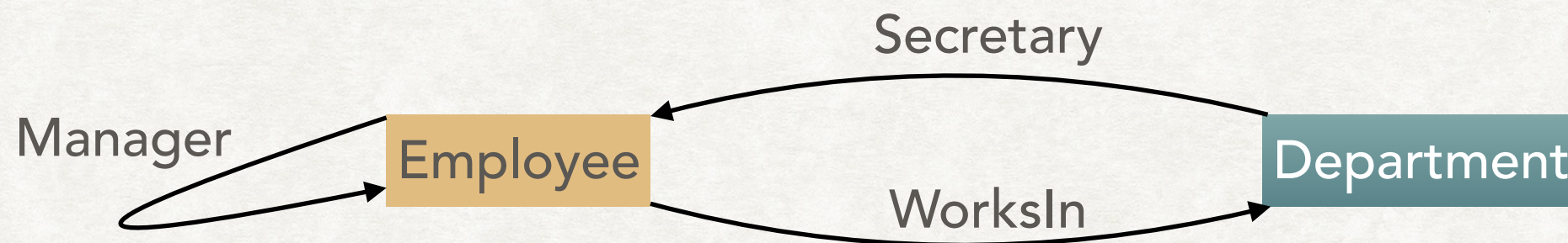
EMPLOYEES

Employee: EmployeeID
WorksIn: DepartmentID
Manager: EmployeeID

DEPARTMENTS

Department: DepartmentID
Secretary: EmployeeID

Each non-ID column
becomes an edge



DATABASE SCHEMA AS A GRAPH: IMPOSING EXTRA RULES

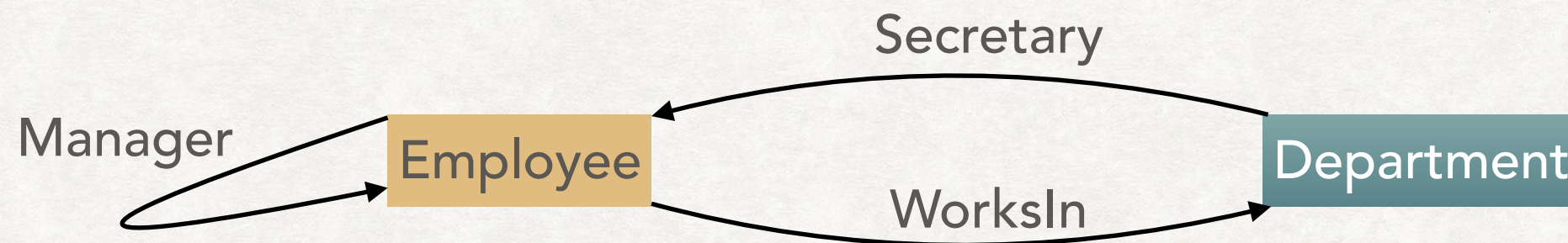
EMPLOYEES

Employee: EmployeeID
WorksIn: DepartmentID
Manager: EmployeeID

DEPARTMENTS

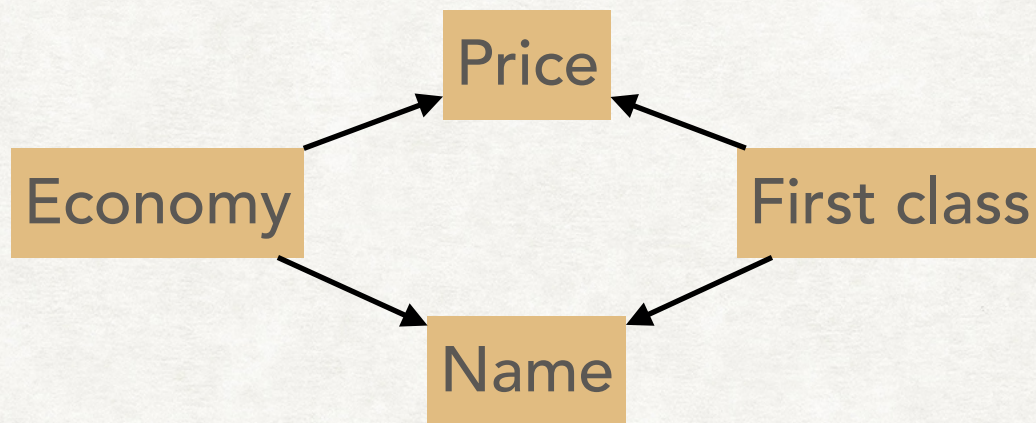
Department: DepartmentID
Secretary: EmployeeID

Each non-ID column
becomes an edge

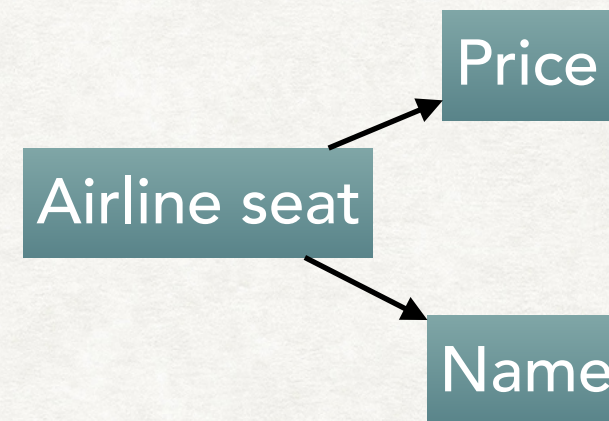


Department.Secretary.Department = Department
Employee.Manager.WorksIn = Employee.WorksIn

ANOTHER EXAMPLE



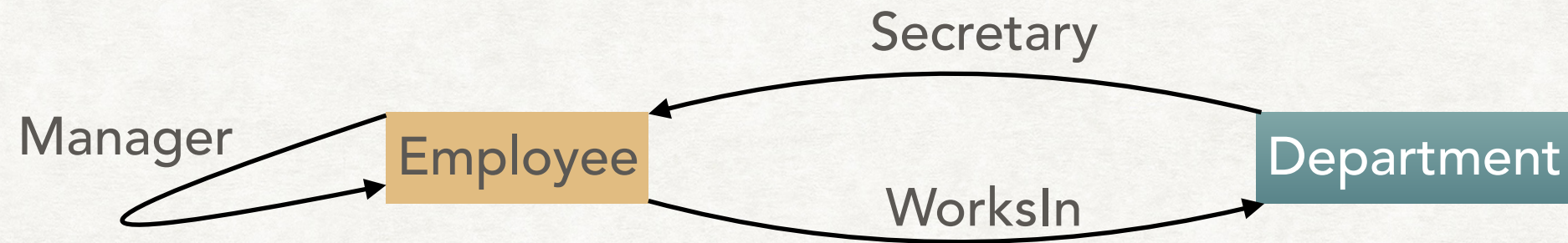
Schema 1



Schema 2

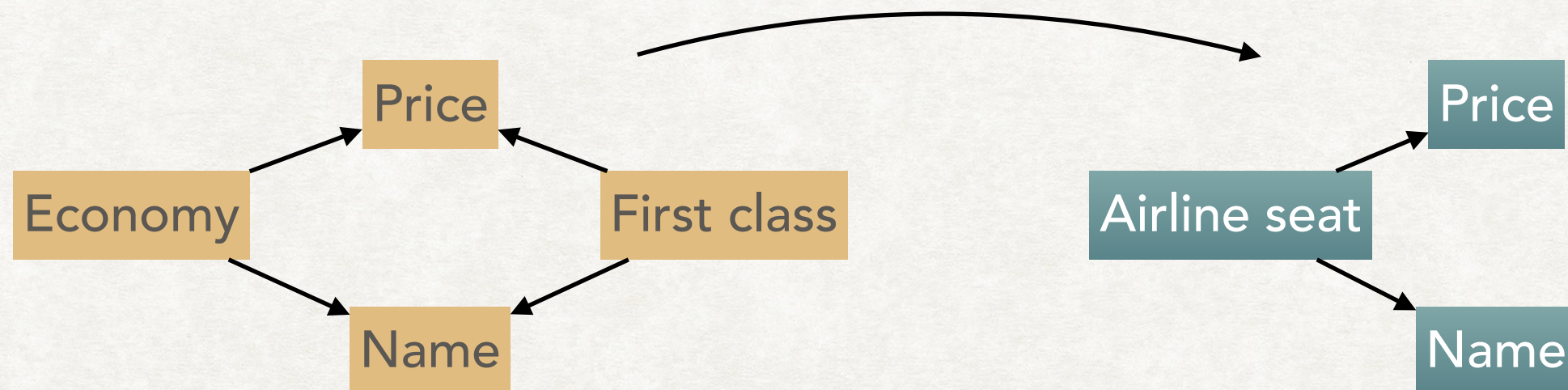
How does one migrate data between two databases?

DATABASE SCHEMA AS A GRAPH: SO WHAT?

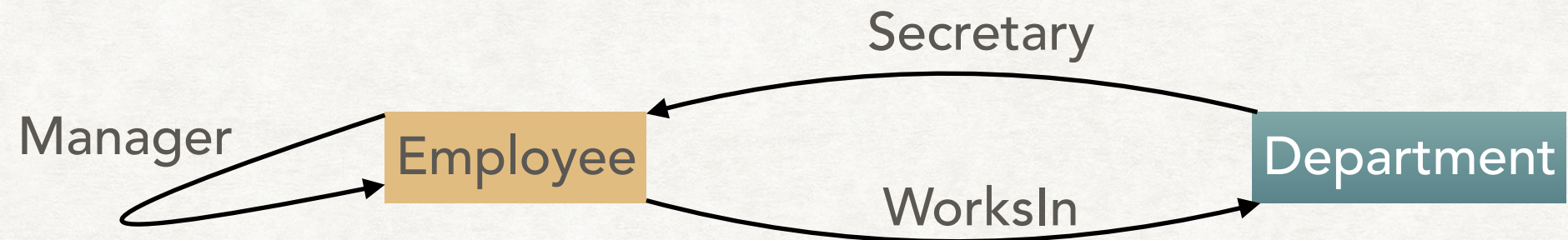


Department.Secretary.Department = Department

Employee.Manager.Employee = Employee



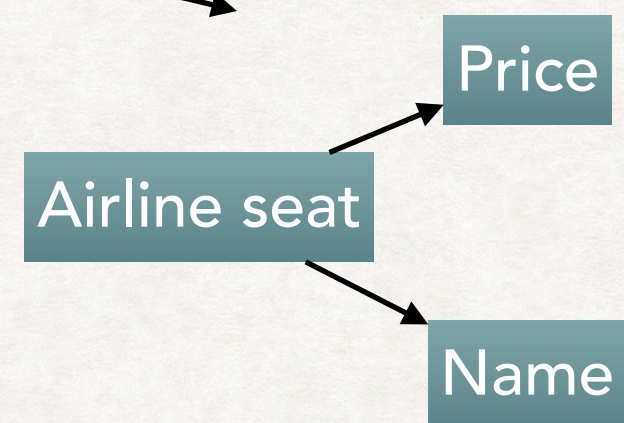
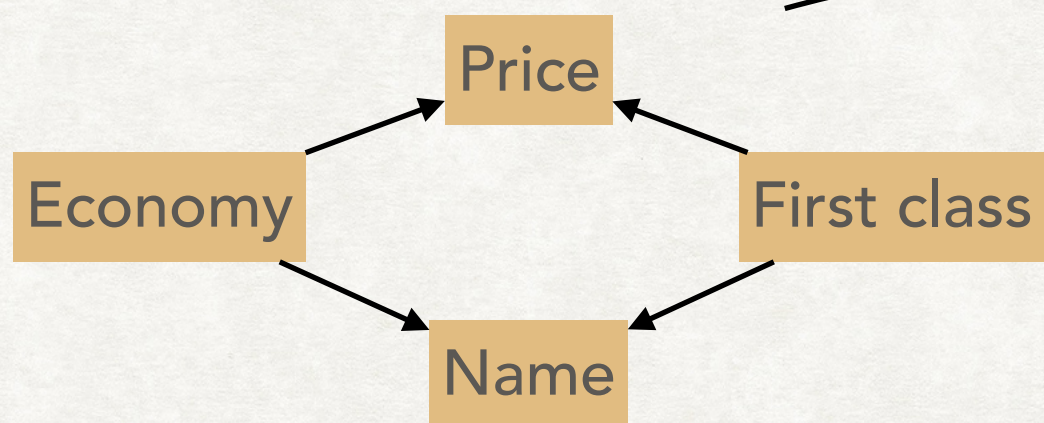
THESE PICTURES ARE CATEGORIES



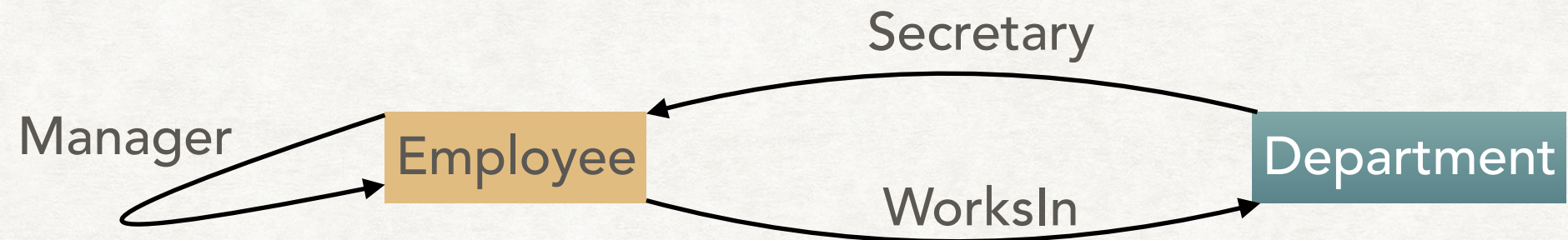
Department.Secretary.Department = Department

Employee.Manager.Employee = Employee

Employee.WorksIn.Department = Employee.WorksIn



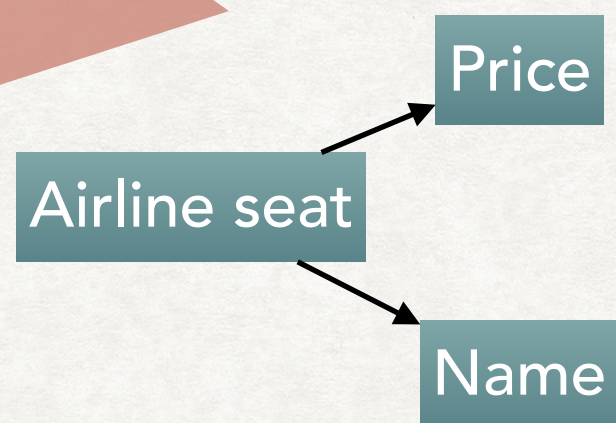
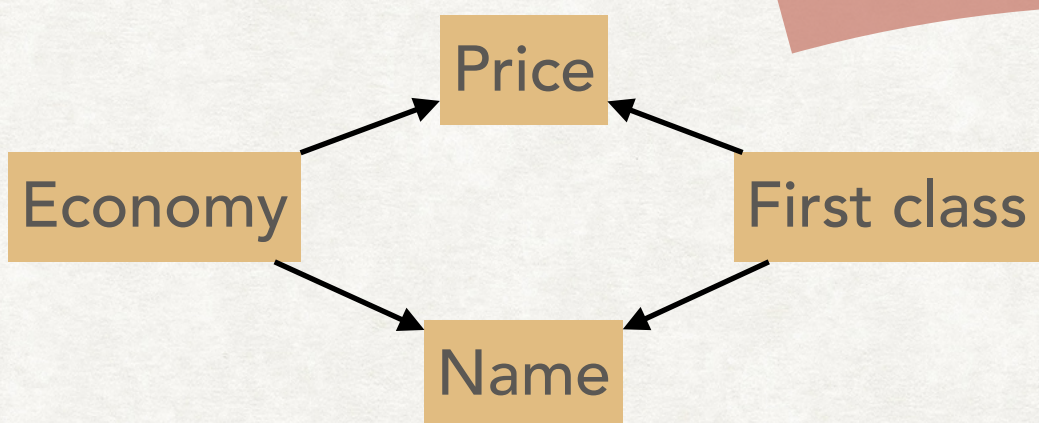
ARROW IS FUNCTOR



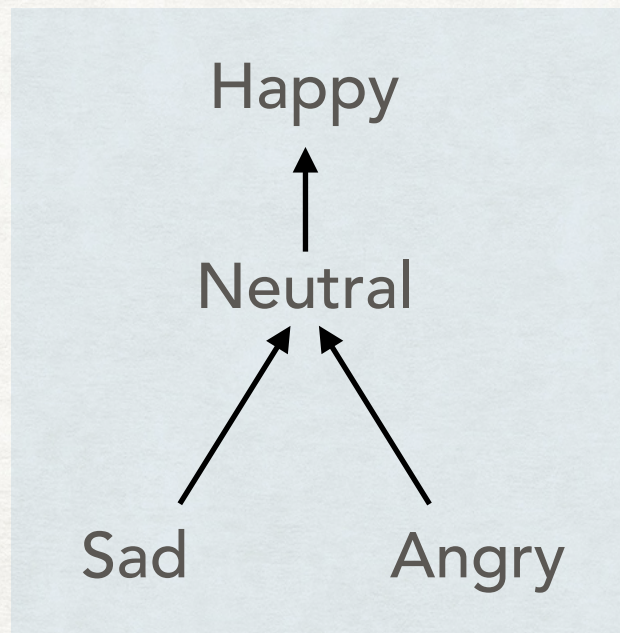
Department.Secretary.Department = Department

Employee.Manager.WorksIn = Employee.WorksIn

FUNCTOR

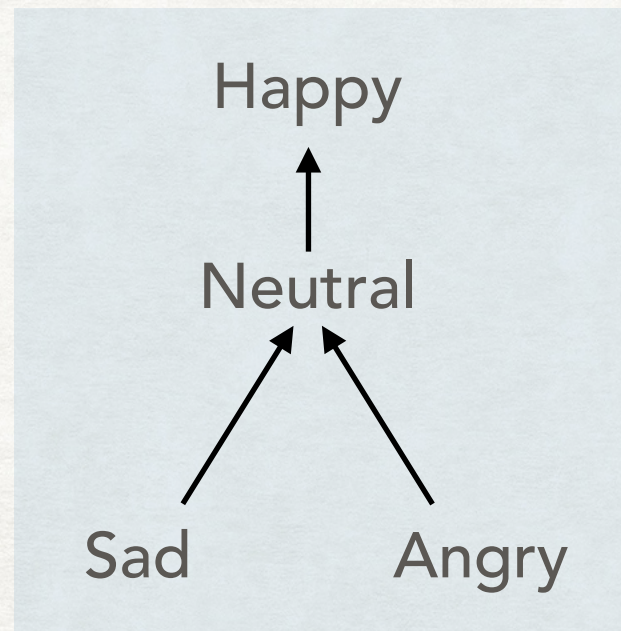


WHAT IS A CATEGORY? LET'S START WITH FAMILIAR EXAMPLES



Partial order

WHAT IS A CATEGORY? LET'S START WITH FAMILIAR EXAMPLES



Partial order

COLLECTION OF OBJECTS

Happy, Neutral, Sad, Angry

EVERY PAIR OF OBJECTS HAS A "RELATIONSHIP"

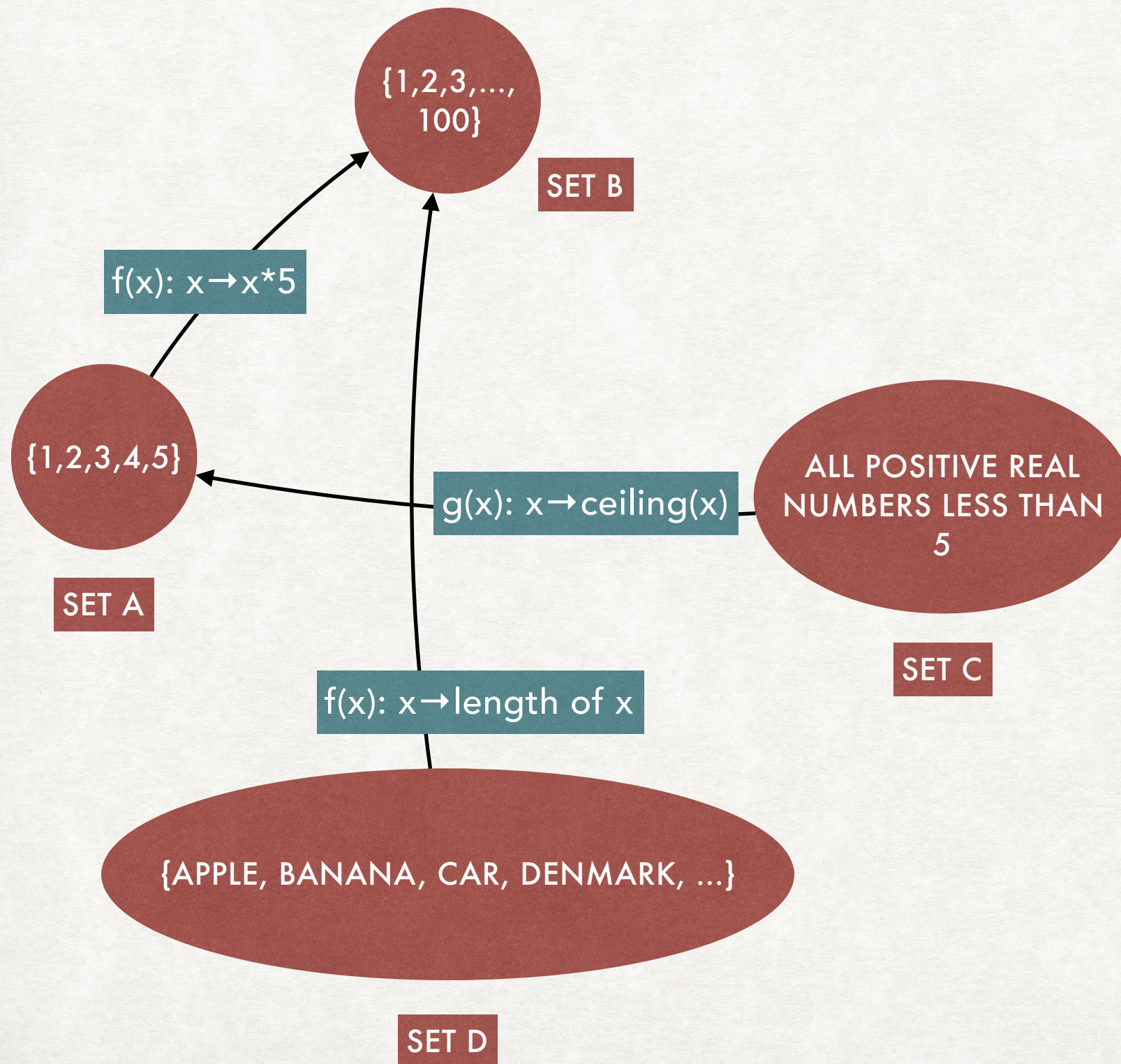
Sad, Neutral: \leq

Neutral, Happy: \leq

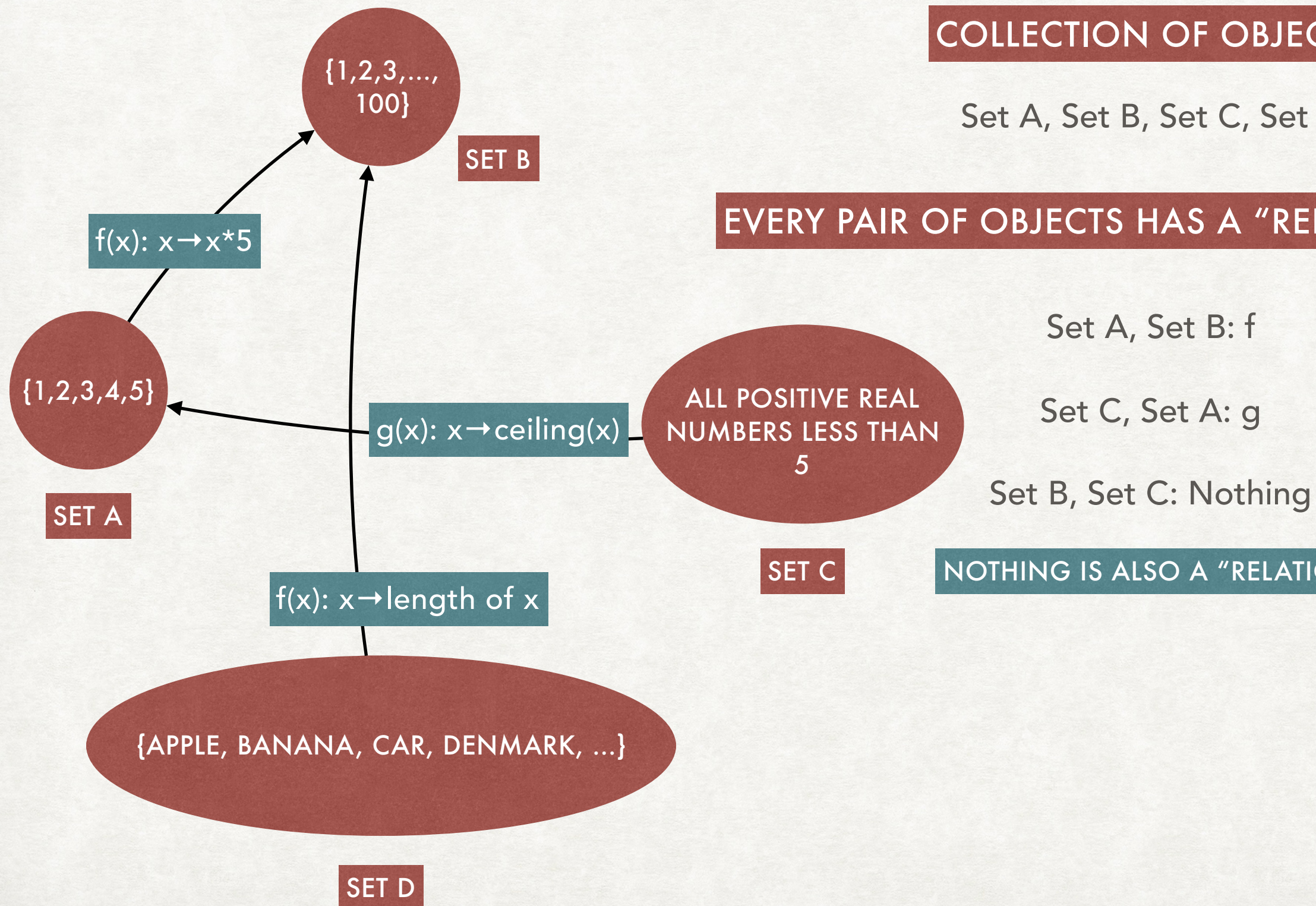
Sad, Angry: Nothing

NOTHING IS ALSO A "RELATIONSHIP"

WHAT IS A CATEGORY? LET'S START WITH FAMILIAR EXAMPLES



WHAT IS A CATEGORY? LET'S START WITH FAMILIAR EXAMPLES



COLLECTION OF OBJECTS

Set A, Set B, Set C, Set D

EVERY PAIR OF OBJECTS HAS A "RELATIONSHIP"

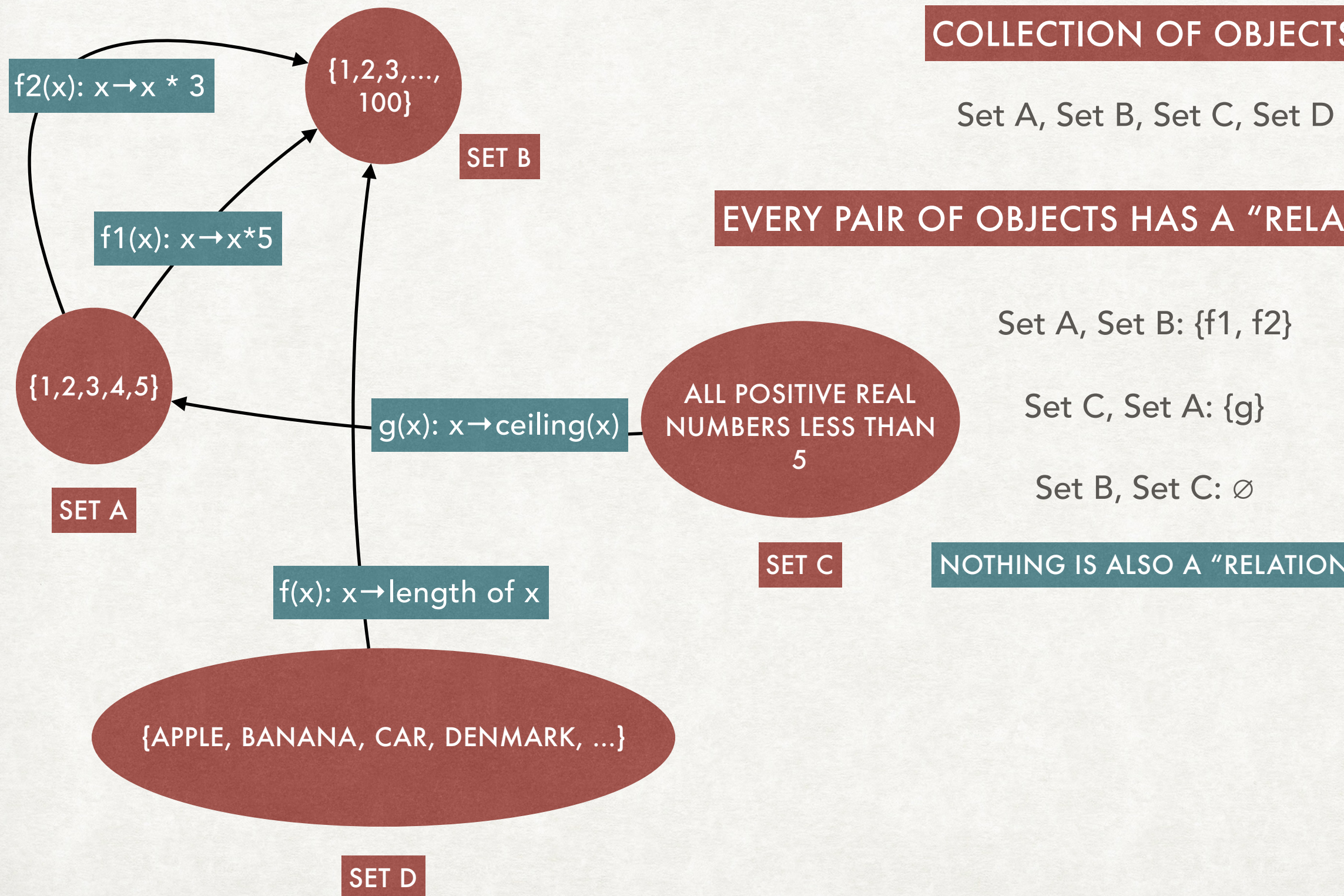
Set A, Set B: f

Set C, Set A: g

Set B, Set C: Nothing

NOTHING IS ALSO A "RELATIONSHIP"

WHAT IS A CATEGORY? LET'S START WITH FAMILIAR EXAMPLES



COLLECTION OF OBJECTS

Set A, Set B, Set C, Set D

EVERY PAIR OF OBJECTS HAS A "RELATIONSHIP"

Set A, Set B: $\{f_1, f_2\}$

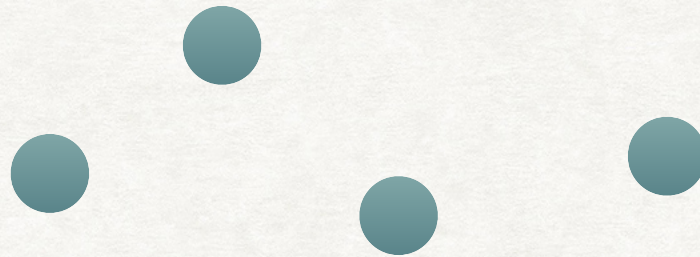
Set C, Set A: $\{g\}$

Set B, Set C: \emptyset

NOTHING IS ALSO A "RELATIONSHIP"

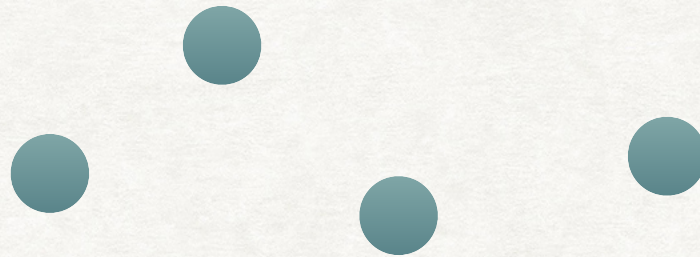
WHAT IS A CATEGORY? DEFINITION

COLLECTION OF OBJECTS

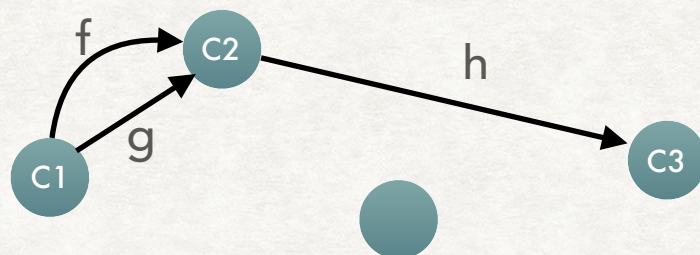


WHAT IS A CATEGORY? DEFINITION

COLLECTION OF OBJECTS



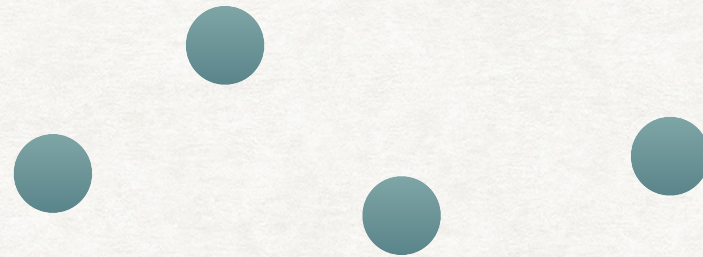
EVERY PAIR OF OBJECTS A AND B HAS A "RELATIONSHIP":
MORPHISM BETWEEN A AND B



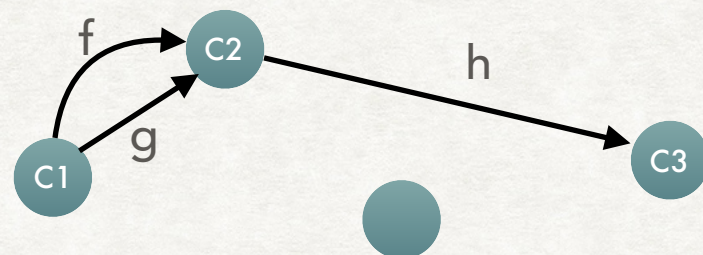
Set of morphisms
between C1 and C2
 $= \{f1, f2\}$

WHAT IS A CATEGORY? DEFINITION

COLLECTION OF OBJECTS

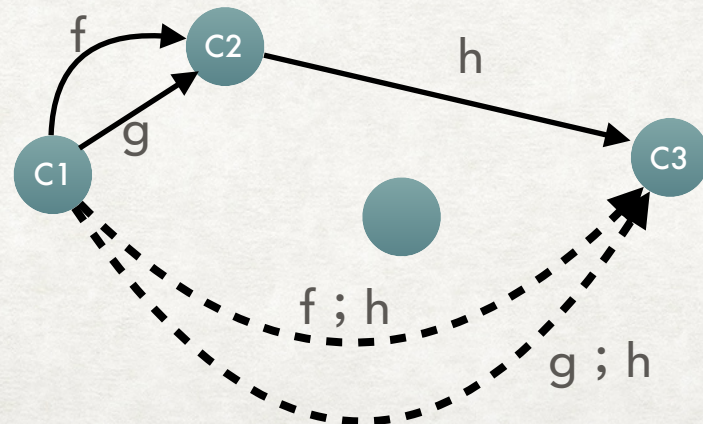


EVERY PAIR OF OBJECTS A AND B HAS A "RELATIONSHIP":
MORPHISM BETWEEN A AND B



Set of morphisms
between C1 and C2
 $= \{f, g\}$

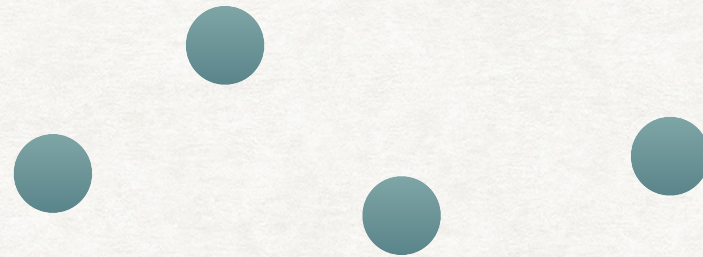
TWO CONSECUTIVE MORPHISMS CAN BE COMPOSED TOGETHER



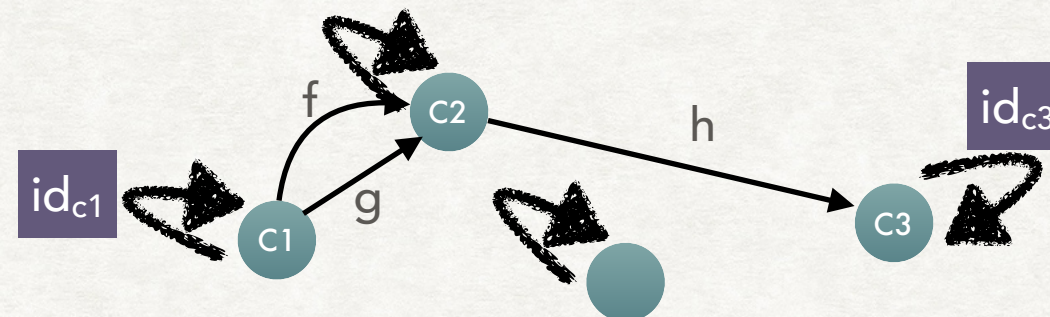
Composition of
morphisms f and h
 $=$ A morphism
between C1 and C3

WHAT IS A CATEGORY? DEFINITION

COLLECTION OF OBJECTS

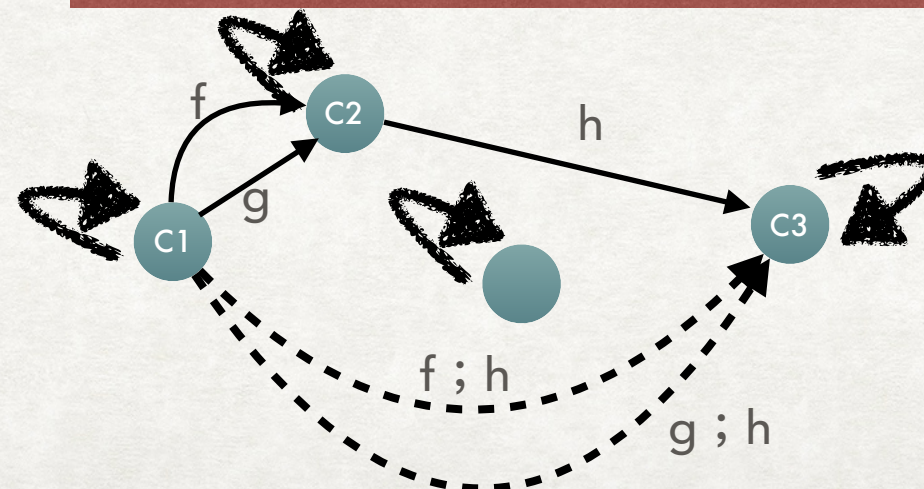


THERE IS A SPECIAL MORPHISM BETWEEN AN ELEMENT AND ITSELF:
MORPHISM $(C1, C1)$ = IDENTITY MORPHISM OF $C1$



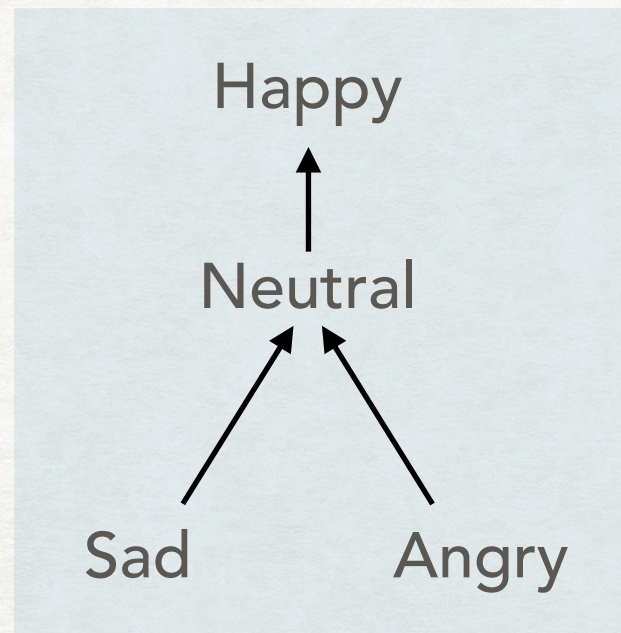
$$id_{C1} ; f = f$$
$$f ; id_{C2} = f$$

MORPHISMS ARE ASSOCIATIVE



$$(f ; g) ; h = f ; (g ; h)$$

WHAT IS A CATEGORY? LET'S START WITH FAMILIAR EXAMPLES



Partial order

COLLECTION OF OBJECTS

Happy, Neutral, Sad, Angry

EVERY PAIR OF OBJECTS HAS A MORPHISM

Sad, Neutral: $\{\leq\}$

Neutral, Happy: $\{\leq\}$

Sad, Angry: \emptyset

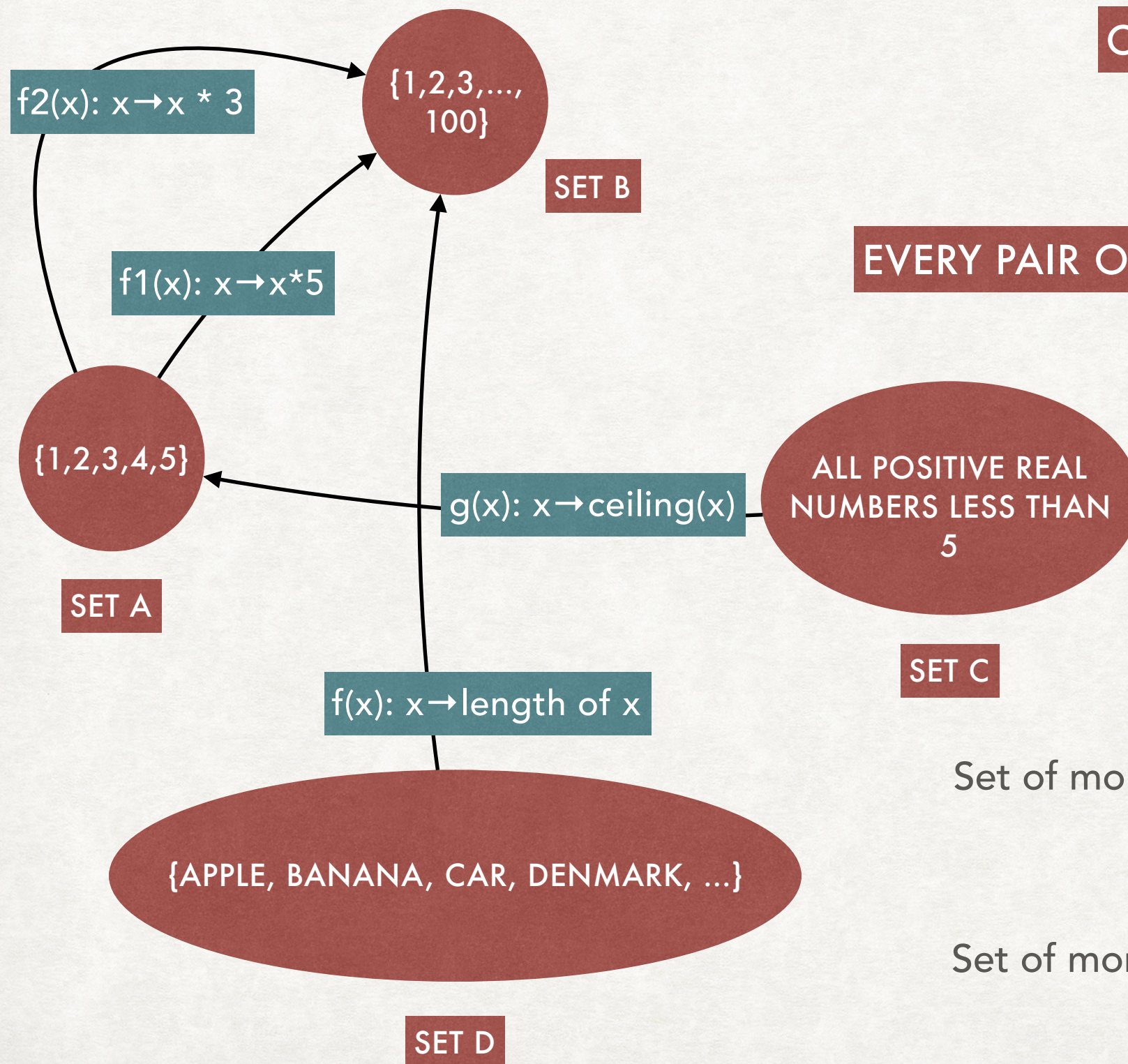
PARTIAL ORDERS HAVE ONLY ONE KIND OF MORPHISM: \leq

$$\text{Id}_{\text{sad}} = \text{Morphism}(\text{Sad}, \text{Sad}) = \{\leq\}$$

$$\text{Morphism}(\text{Sad}, \text{Neutral}) ; \text{Morphism}(\text{Neutral}, \text{Happy}) = \text{Morphism}(\text{Sad}, \text{Happy})$$

$\text{Sad} \leq \text{Neutral}$ and $\text{Neutral} \leq \text{Happy}$
implies
 $\text{Sad} \leq \text{Happy}$

WHAT IS A CATEGORY? LET'S START WITH FAMILIAR EXAMPLES



COLLECTION OF OBJECTS

Set A, Set B, Set C, Set D

EVERY PAIR OF SETS HAS A SET OF MORPHISMS

Set A, Set B: {f1, f2}

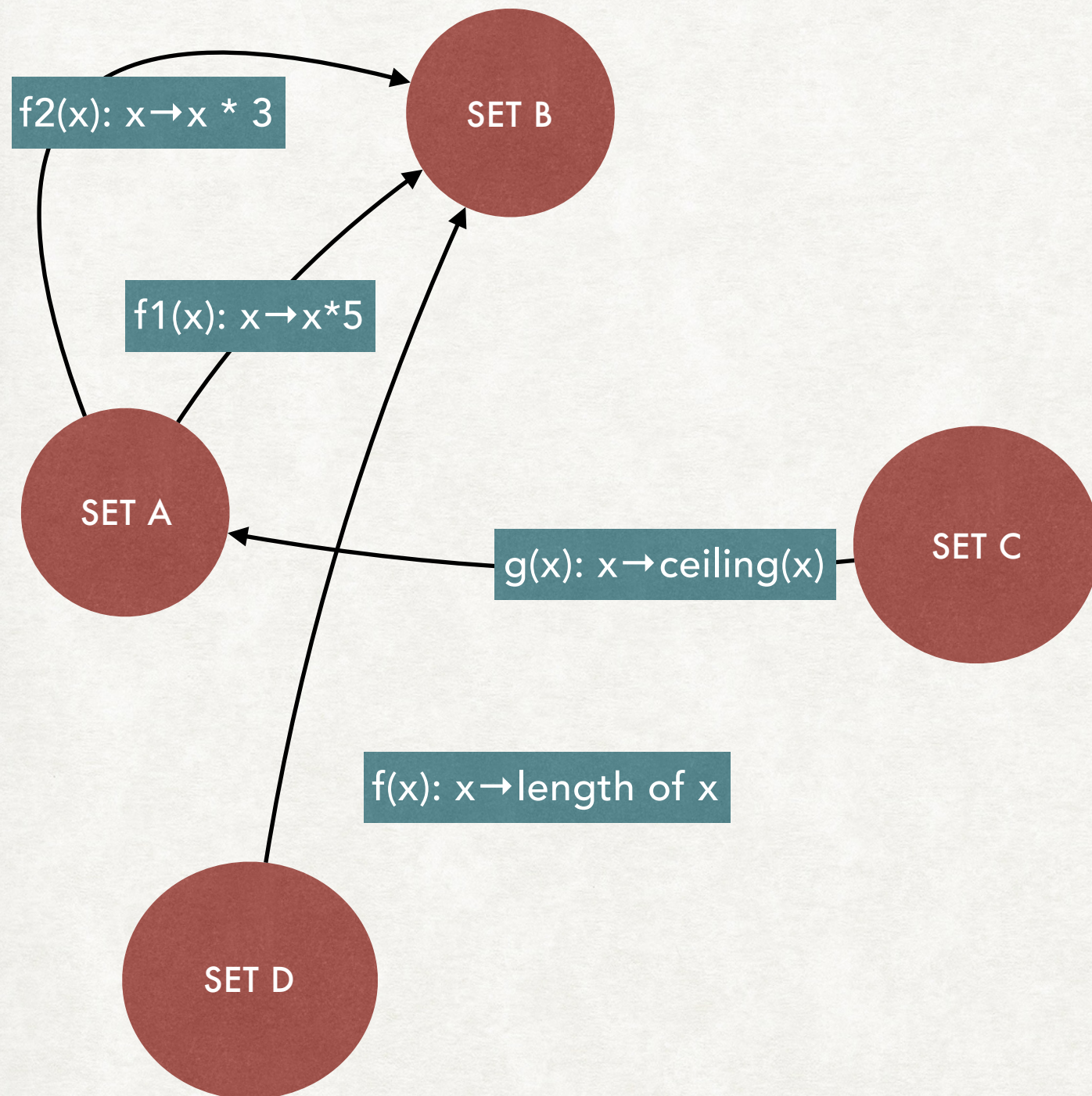
Set C, Set A: {g}

Set B, Set C: \emptyset

Set of morphisms between (Set C, Set B) =
 $\{g ; f_1, g ; f_2\}$

Set of morphisms between (Set A, Set A) =
 $\text{id}_A = x \rightarrow x$

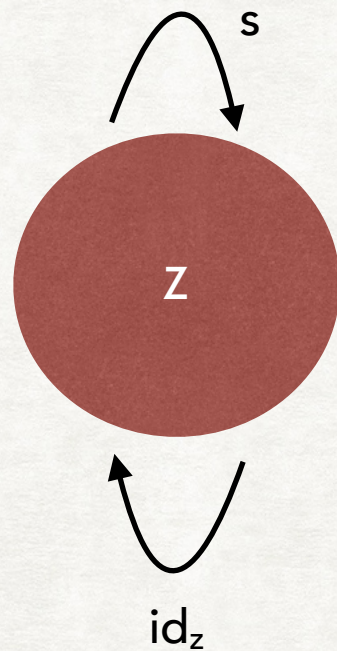
WHAT IS A CATEGORY? LET'S START WITH FAMILIAR EXAMPLES



WHEN YOU THINK OF OBJECTS
IN A CATEGORY, YOU DON'T
CARE WHAT'S INSIDE EACH
OBJECT.

YOU ONLY CARE ABOUT
RELATIONSHIPS BETWEEN
OBJECTS

CATEGORY WITH A SINGLE OBJECT



How many morphisms does it have?

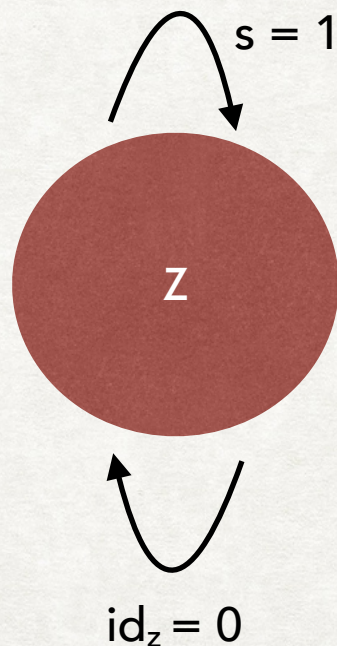
$$s = s ; \text{id}_Z = \text{id}_Z ; s$$

$$s ; s$$

$$s ; s ; s$$

$$s ; s ; s ; s$$

$$\circ \circ \circ$$



Define $a ; b = a + b$

How many morphisms does it have?

$$1 = 1 ; 0 = 0 ; 1$$

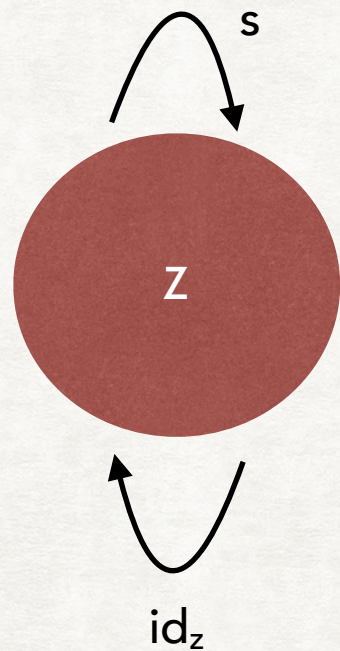
$$1 ; 1 = 2$$

$$1 ; 1 ; 1 = 2 ; 1 = 3$$

$$1 ; 1 ; 1 ; 1 = 4$$

$$\circ \circ \circ$$

CATEGORY WITH A SINGLE OBJECT: WE CAN CONSTRAIN MORPHISMS



How many morphisms does it have?

$$s ; s = s$$

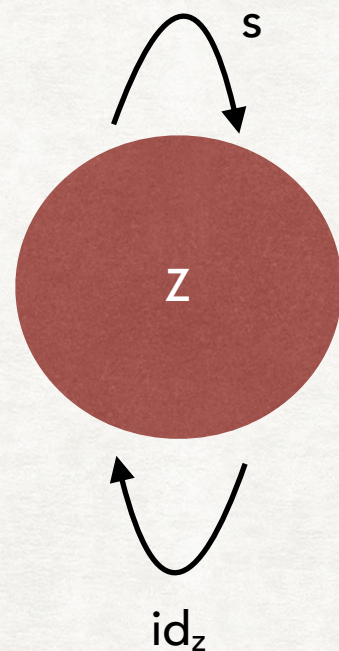
$$s = s ; \text{id}_Z = \text{id}_Z ; s$$

$$s ; s = s$$

$$s ; s ; s = s ; s = s$$

Thus it has only two morphisms: id_Z and s

CATEGORY WITH A SINGLE OBJECT: WE CAN CONSTRAIN MORPHISMS



$$s ; s = s$$

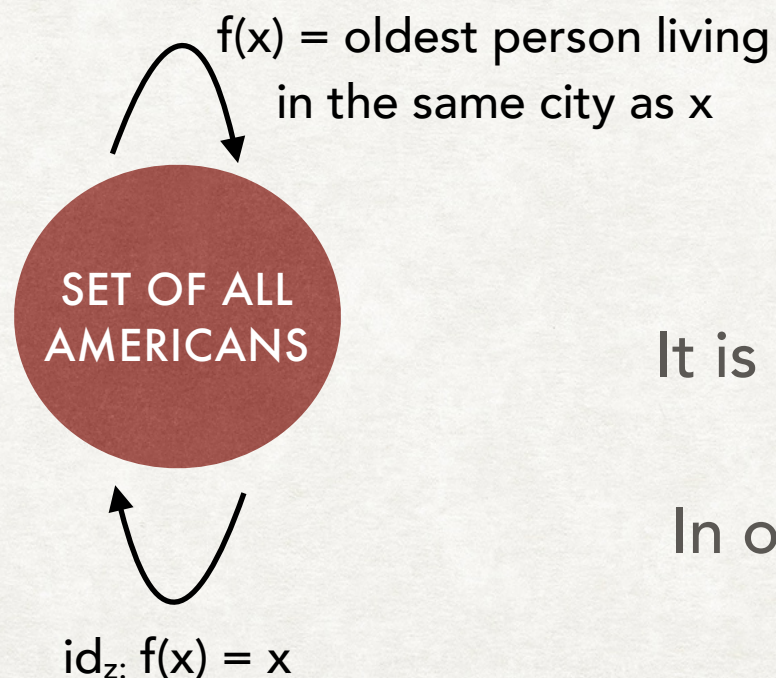
How many morphisms does it have?

$$s = s ; \text{id}_Z = \text{id}_Z ; s$$

$$s ; s = s$$

$$s ; s ; s = s ; s = s$$

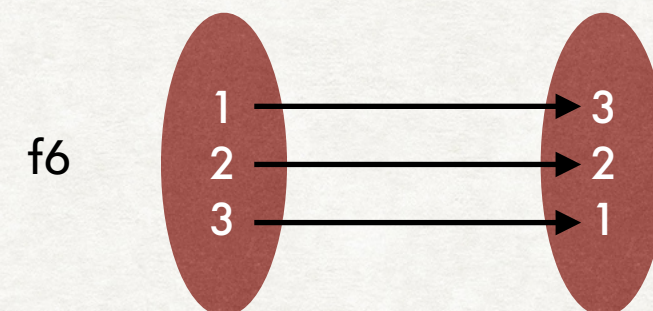
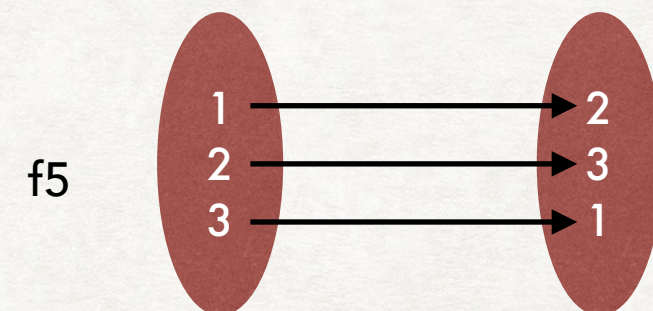
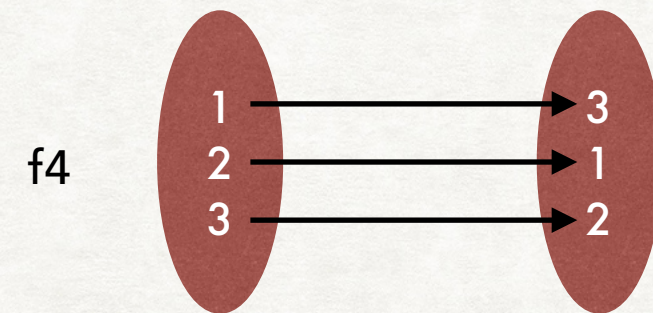
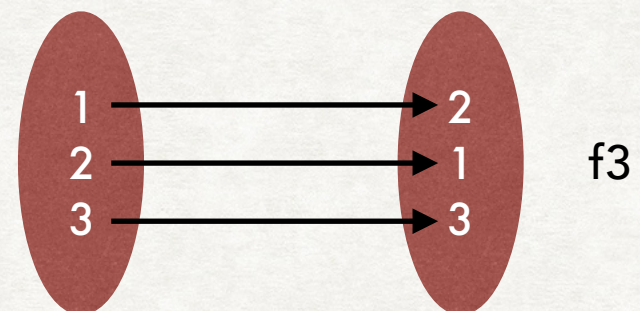
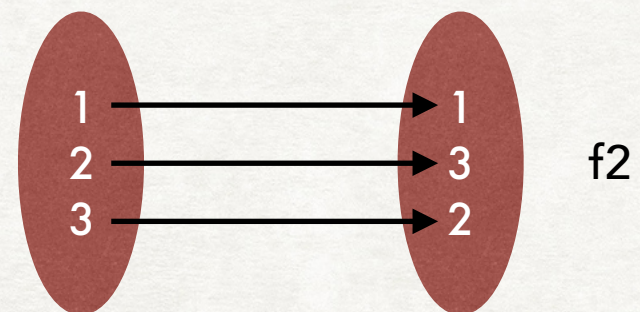
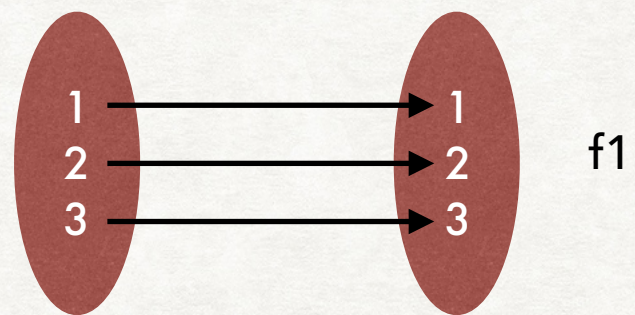
Thus it has only two morphisms: id_Z and s



It is easy to see that $f ; f = f$

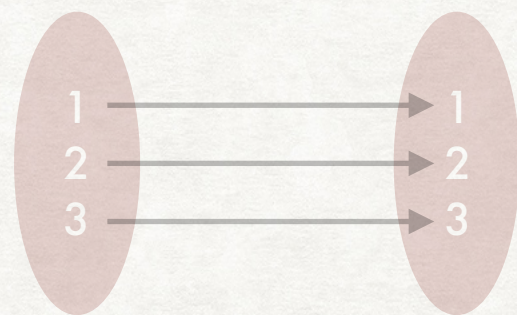
In other words, $f(f(x)) = f(x)$

WHAT IS A CATEGORY? PERMUTATION GROUP AS A CATEGORY

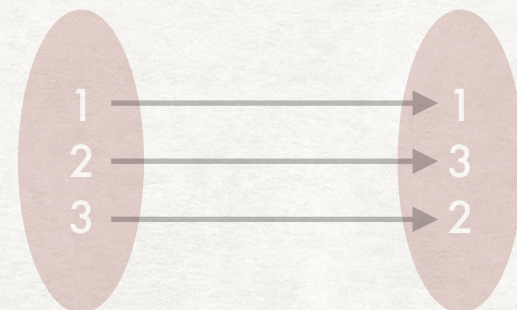


$\{f1, f2, f3, f4, f5, f6\}$
is a group of permutations

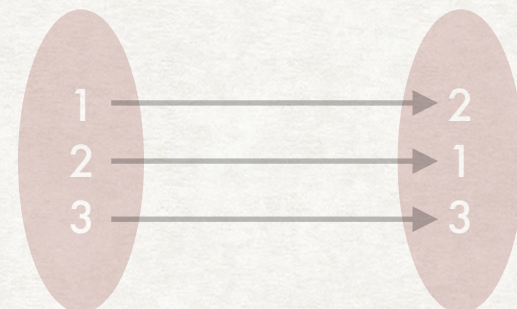
WHAT IS A CATEGORY? PERMUTATION GROUP AS A CATEGORY



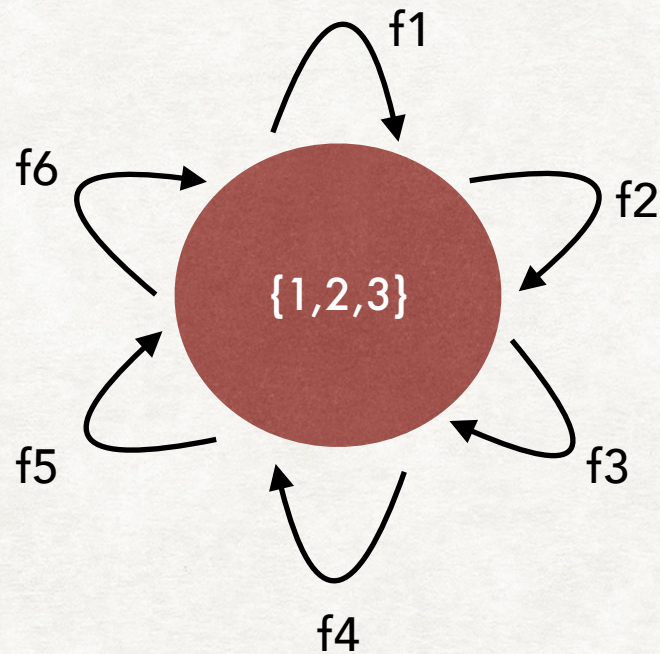
f1



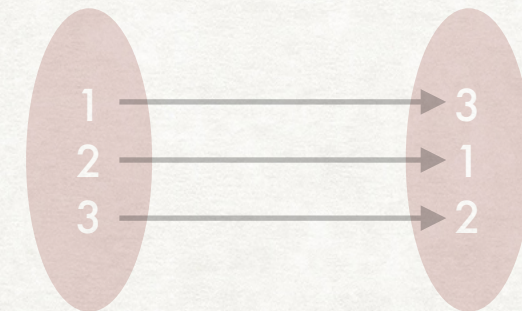
f2



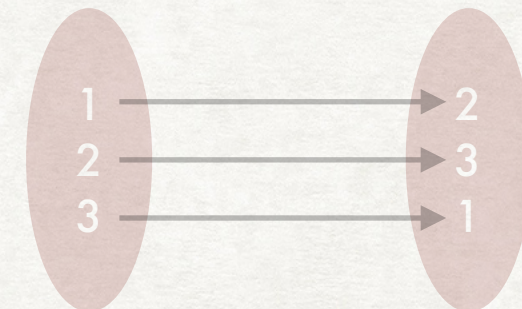
f3



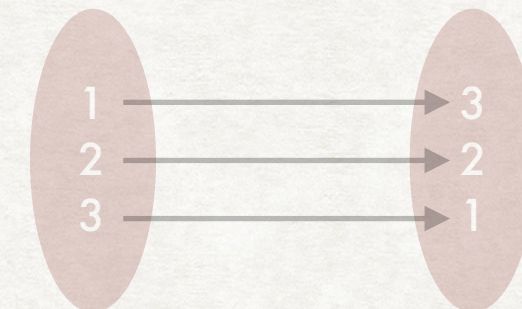
f4



f5



f6

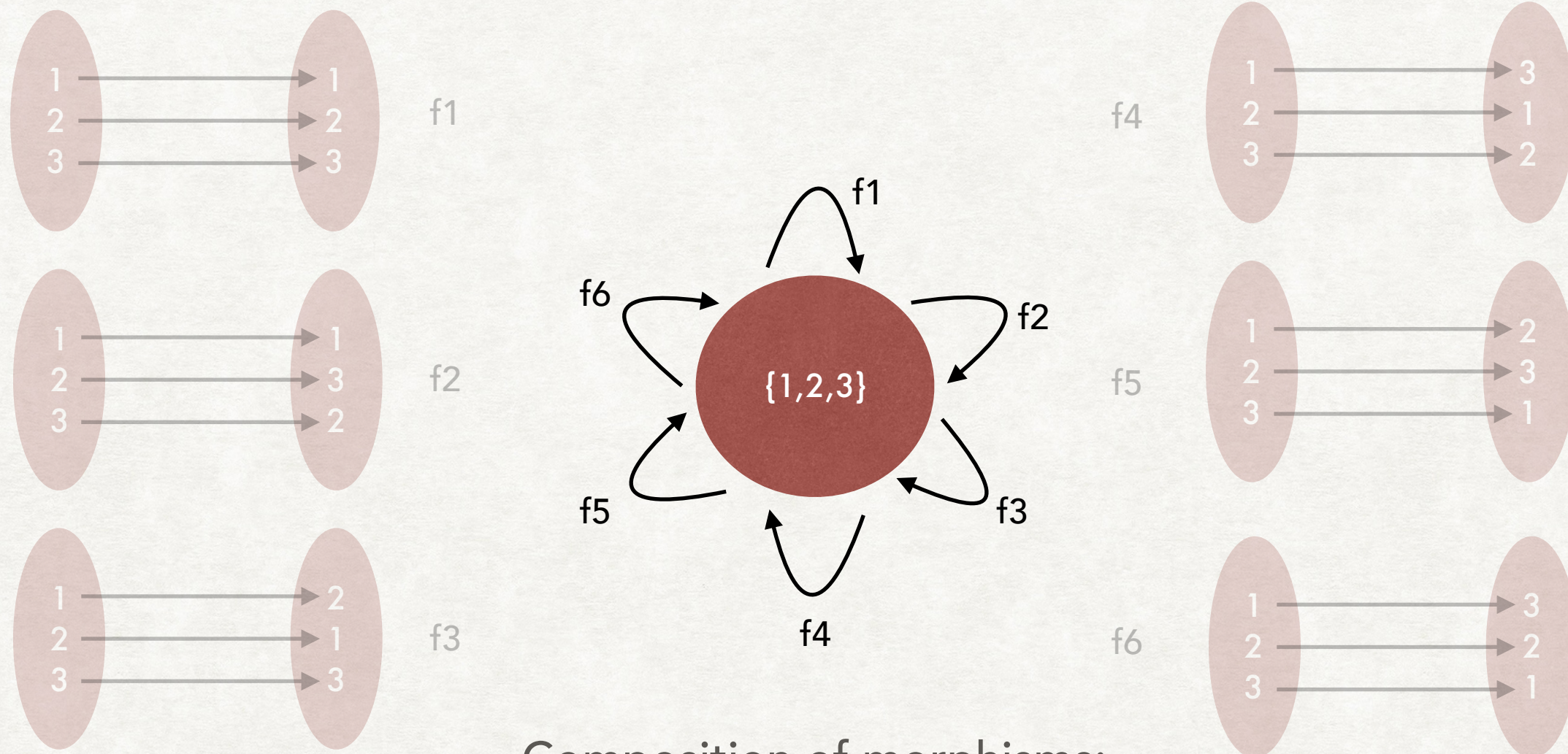


Composition of morphisms:

$$f5 ; f2 = f6$$

$\{f1, f2, f3, f4, f5, f6\}$
is a group of permutations

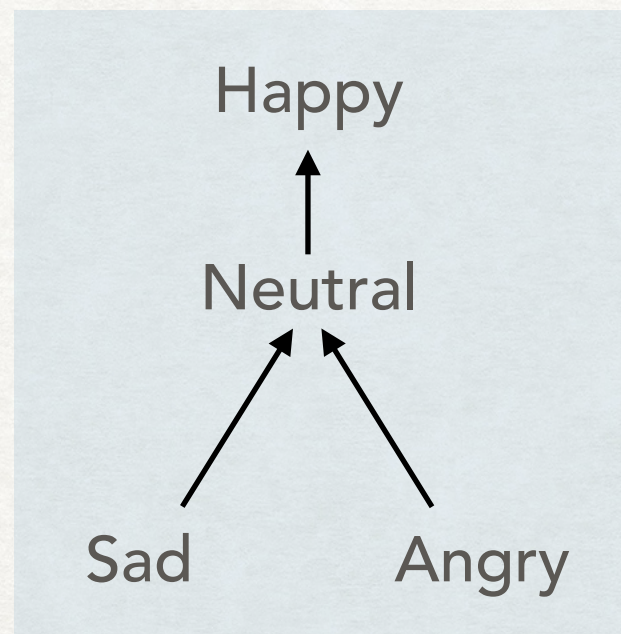
WHAT IS A CATEGORY? PERMUTATION GROUP AS A CATEGORY



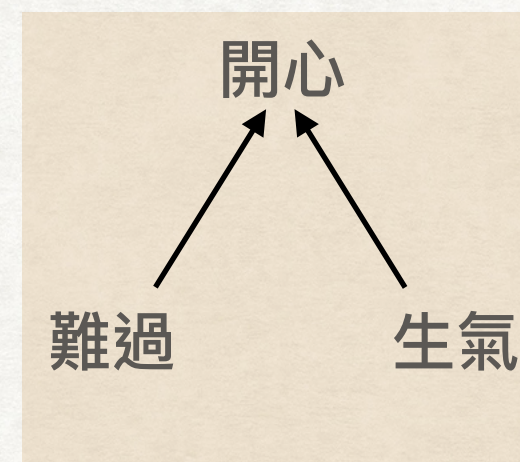
Composition of morphisms:
 $f_5 ; f_2 = f_6$

Category with one object and six morphisms

WHAT IS A FUNCTOR? LET'S START WITH A FAMILIAR EXAMPLE

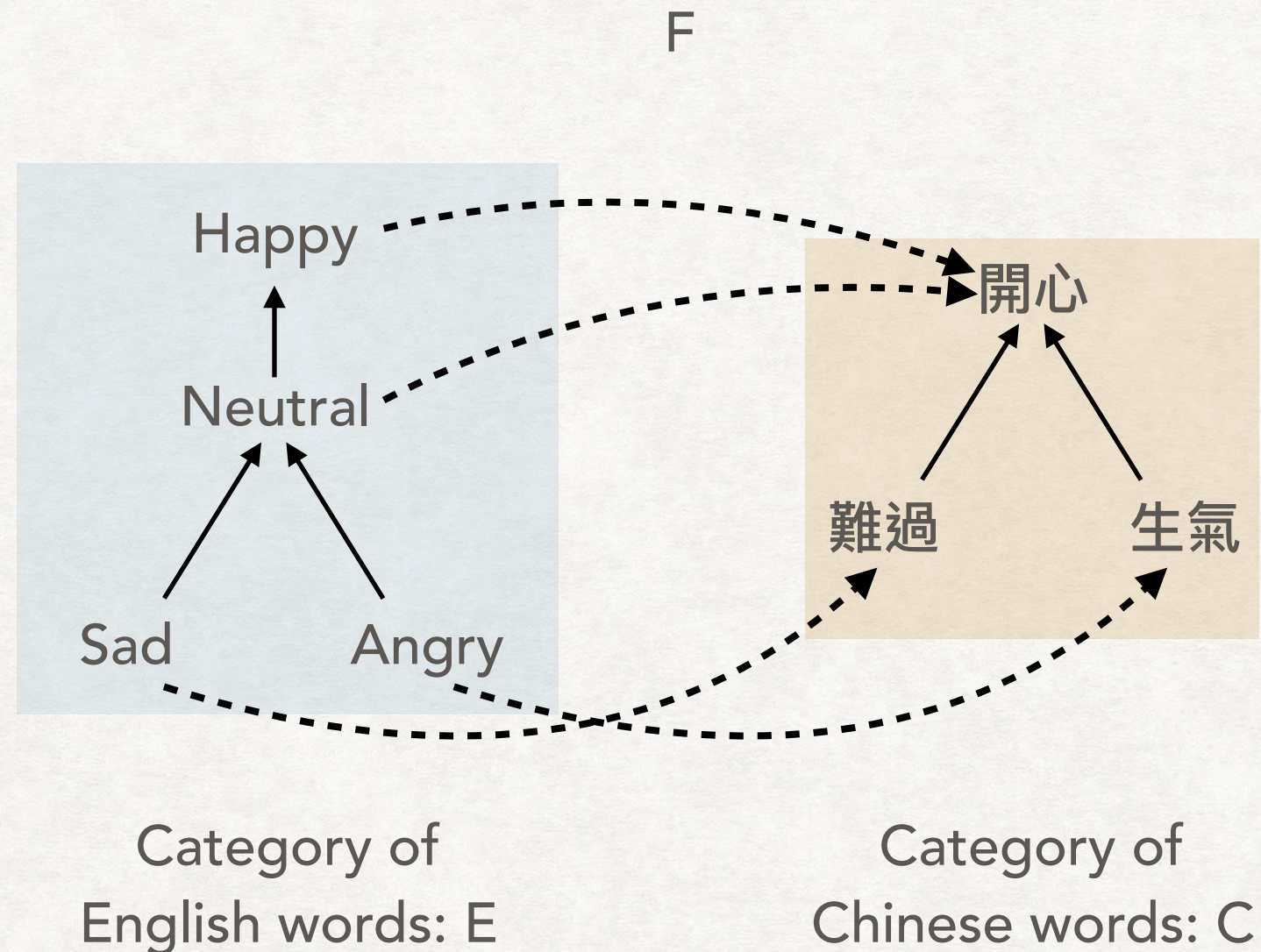


Category of
English words: E



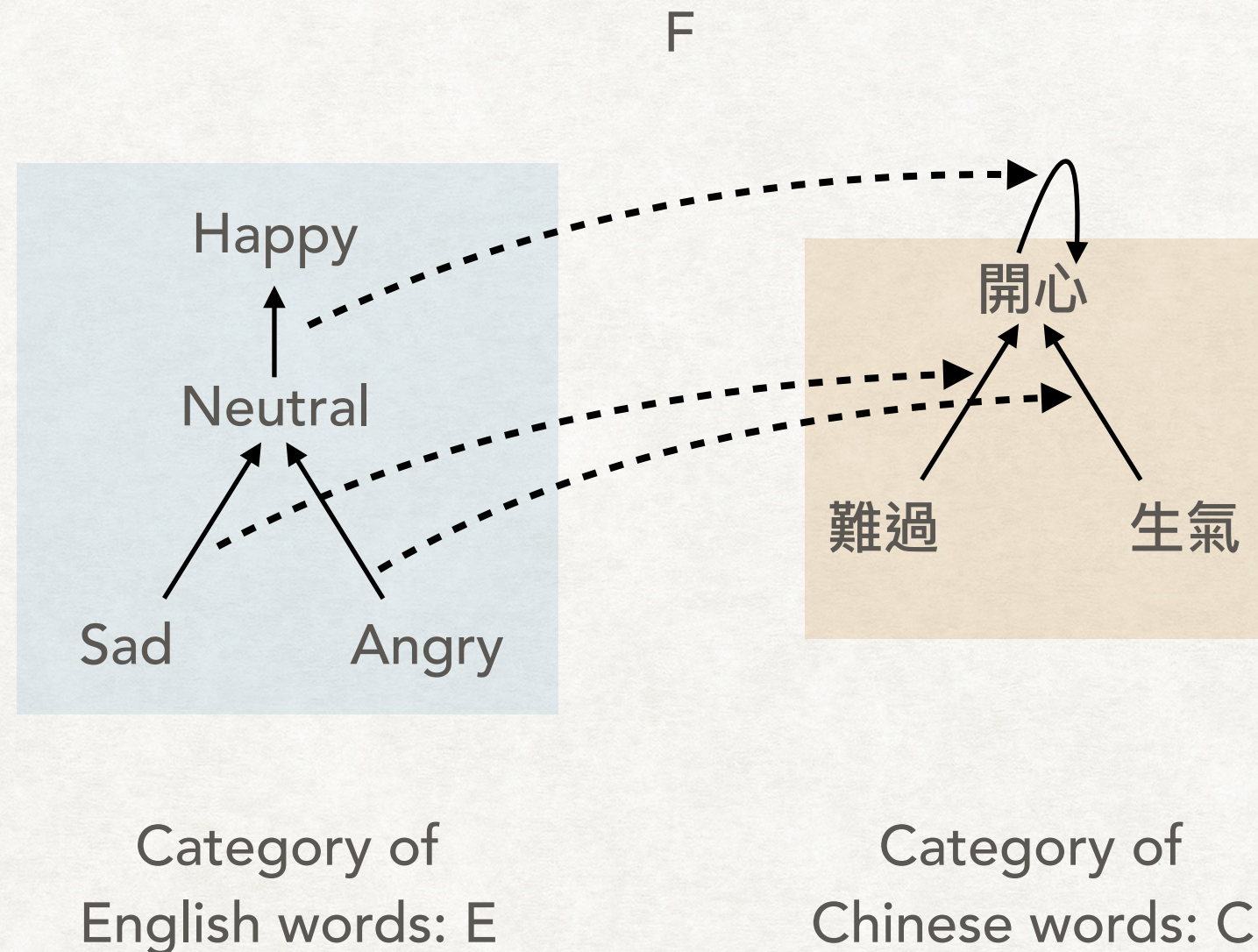
Category of
Chinese words: C

WHAT IS A FUNCTOR? LET'S START WITH A FAMILIAR EXAMPLE



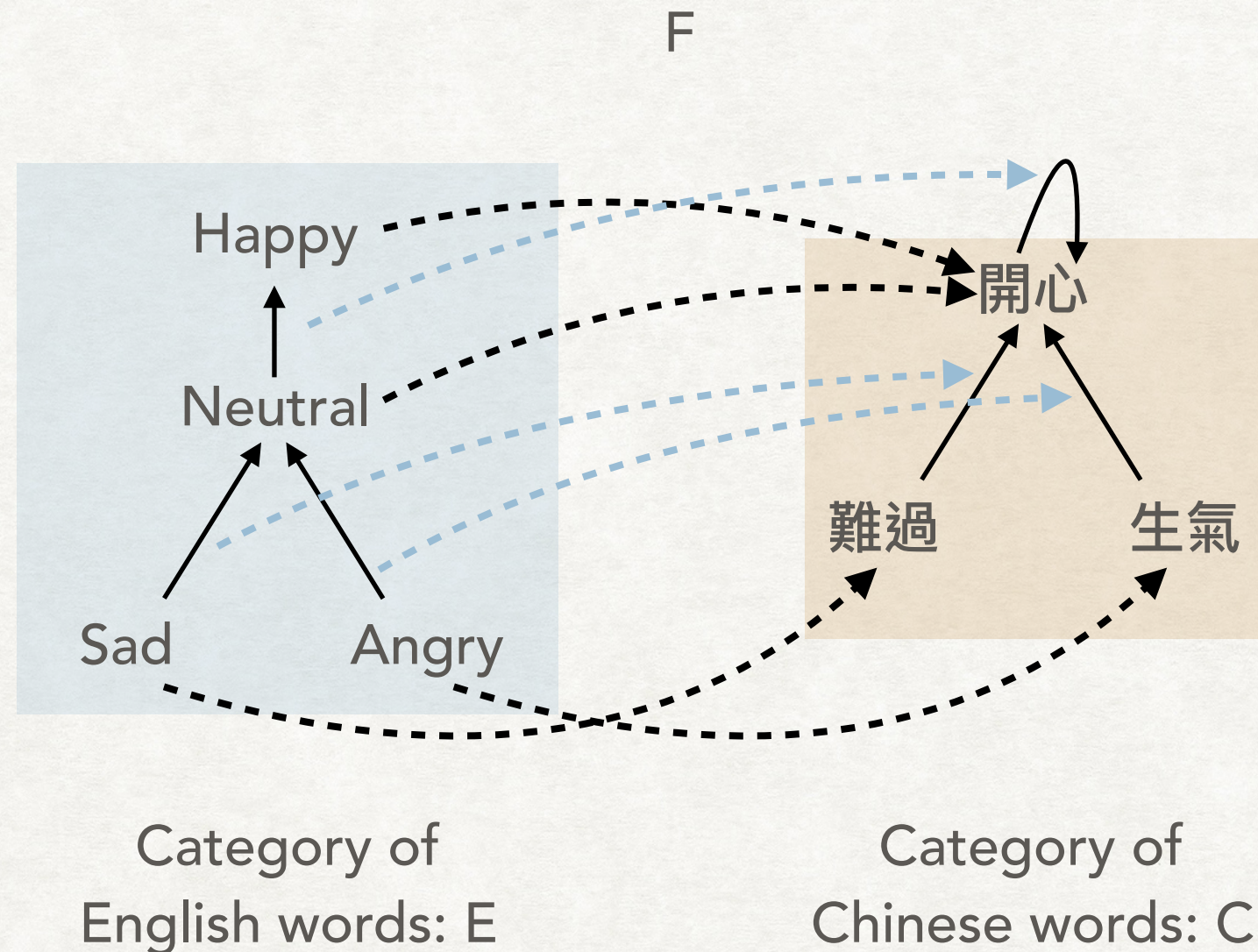
MAP EVERY OBJECT IN E TO AN OBJECT IN C

WHAT IS A FUNCTOR? LET'S START WITH A FAMILIAR EXAMPLE



MAP EVERY MORPHISM IN E TO A MORPHISM IN C

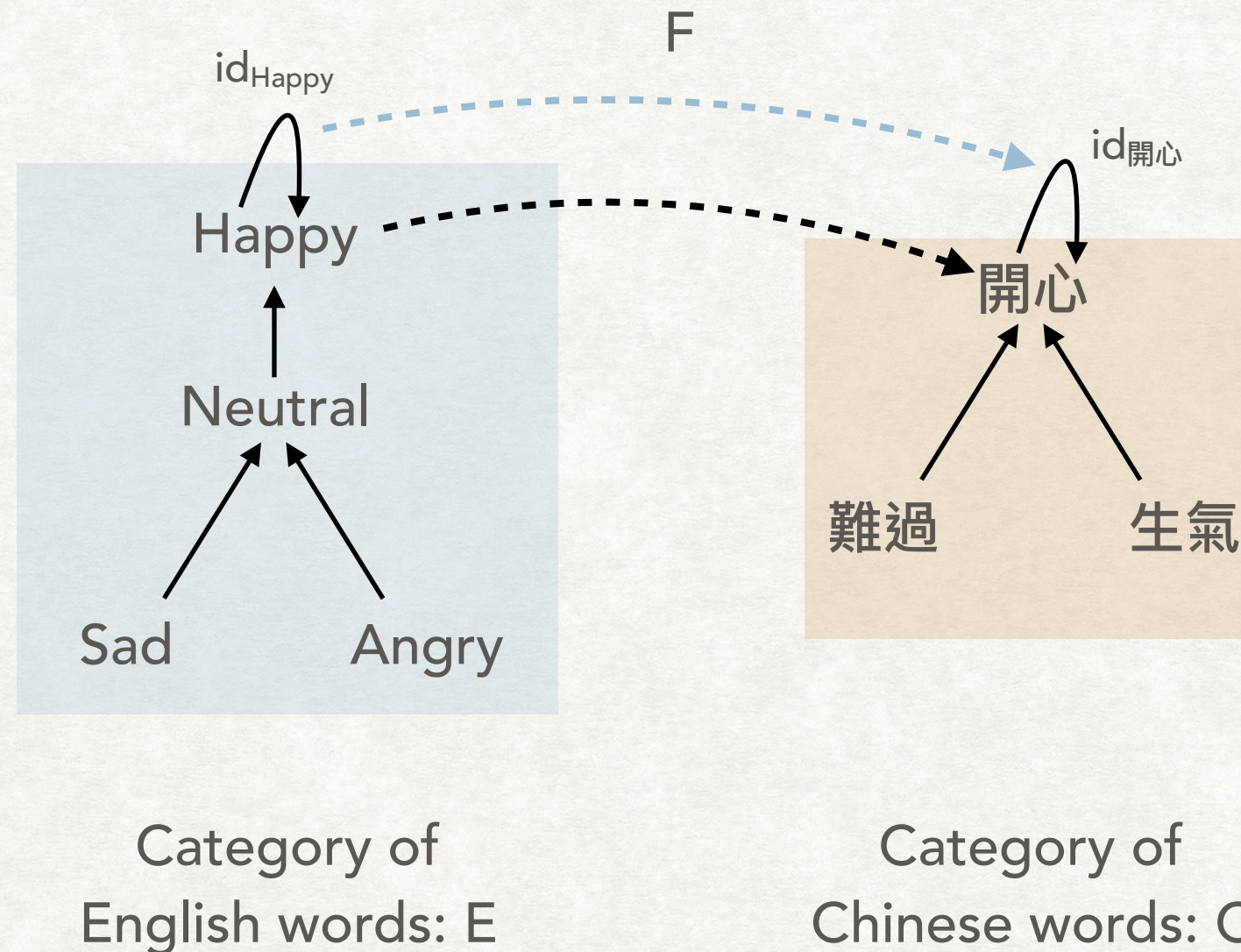
WHAT IS A FUNCTOR? LET'S START WITH A FAMILIAR EXAMPLE



A FUNCTOR BASICALLY MAPS EACH OBJECT/MORPHISM FROM ONE CATEGORY TO AN OBJECT/MORPHISM IN ANOTHER CATEGORY

BUT THERE ARE TWO RULES

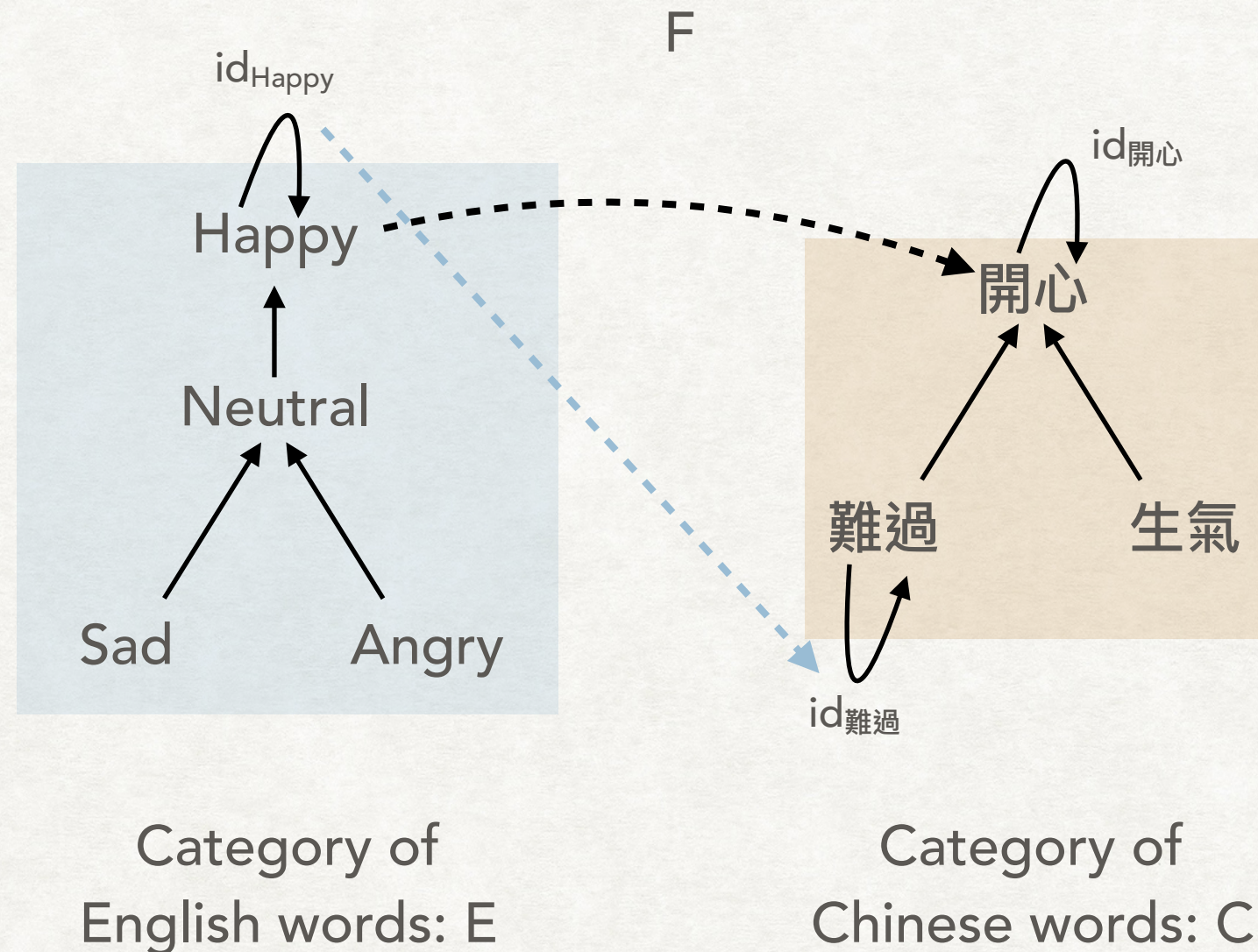
WHAT IS A FUNCTOR? LET'S START WITH A FAMILIAR EXAMPLE



$$F(id_{Happy}) = id_{F(Happy)}$$

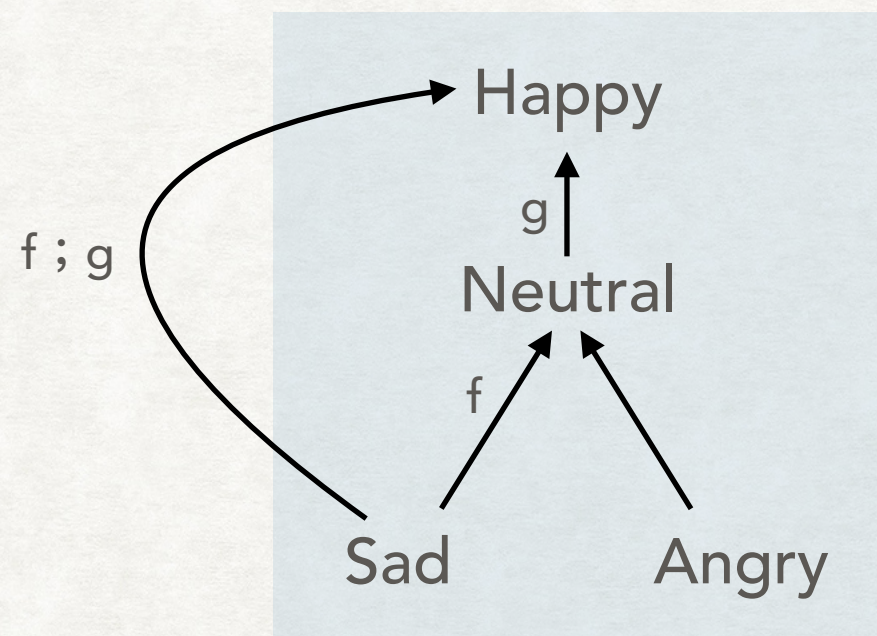
Same rule applies for each object and its identity morphism

WHAT IS A FUNCTOR? LET'S START WITH A FAMILIAR EXAMPLE

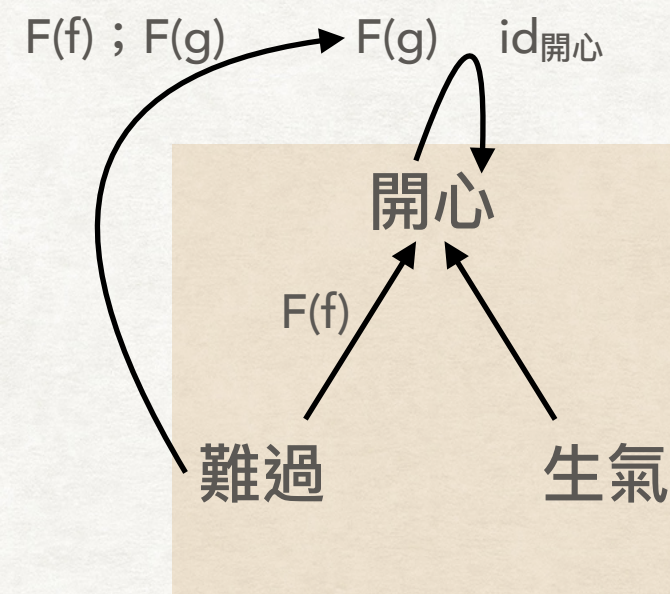


Is $F(\text{ID}_{\text{Happy}}) = \text{ID}_{F(\text{Happy})}$?
 $F(\text{ID}_{\text{Happy}}) = \text{id}_{\text{難過}}$ but
 $\text{ID}_{(F(\text{Happy}))} = \text{id}_{\text{開心}}$
Hence this mapping is not a valid functor

WHAT IS A FUNCTOR? LET'S START WITH A FAMILIAR EXAMPLE



Category of
English words: E

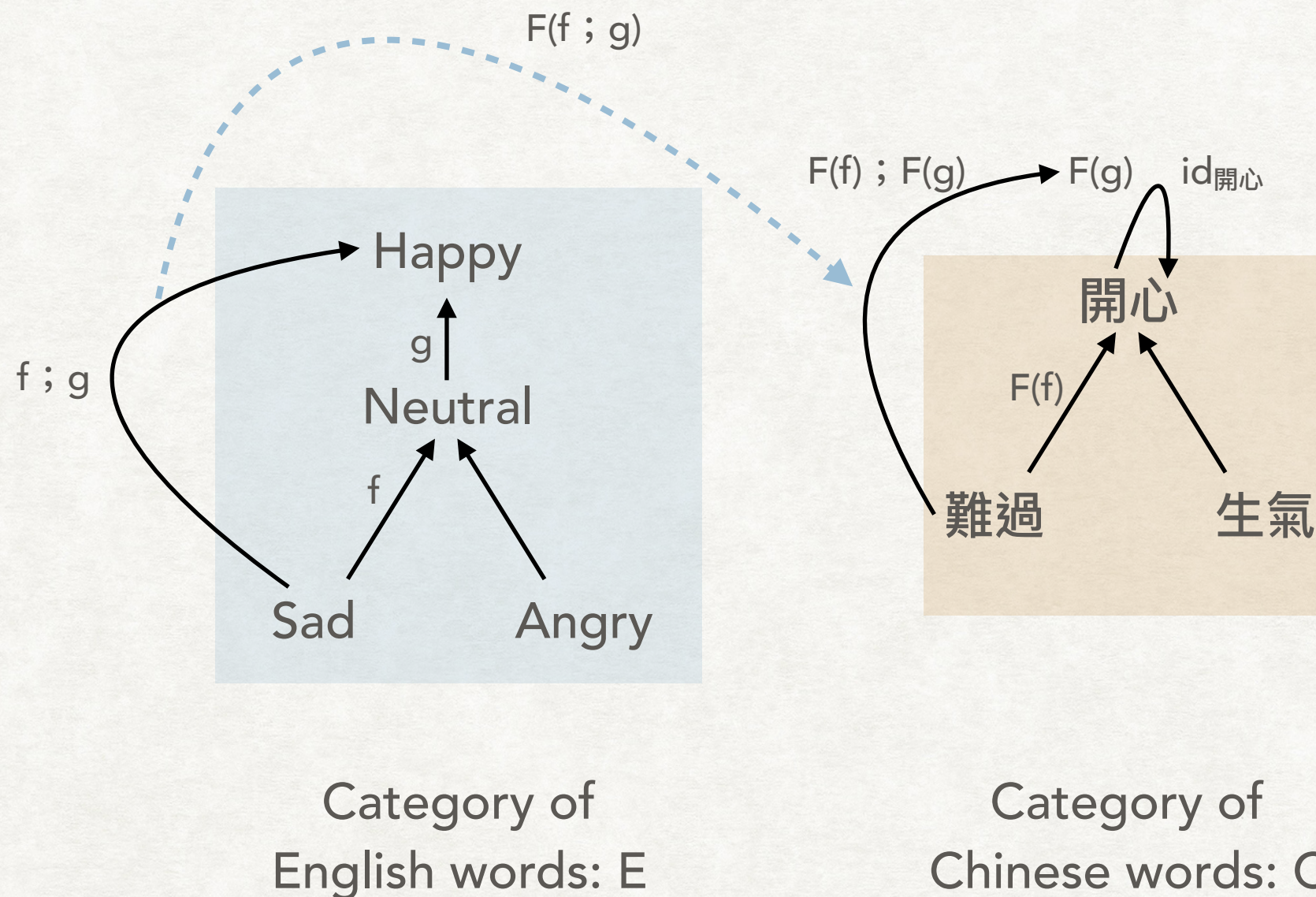


Category of
Chinese words: C

$$F(f ; g) = F(f) ; F(g)$$

Same rule applies for each object and its identity morphism

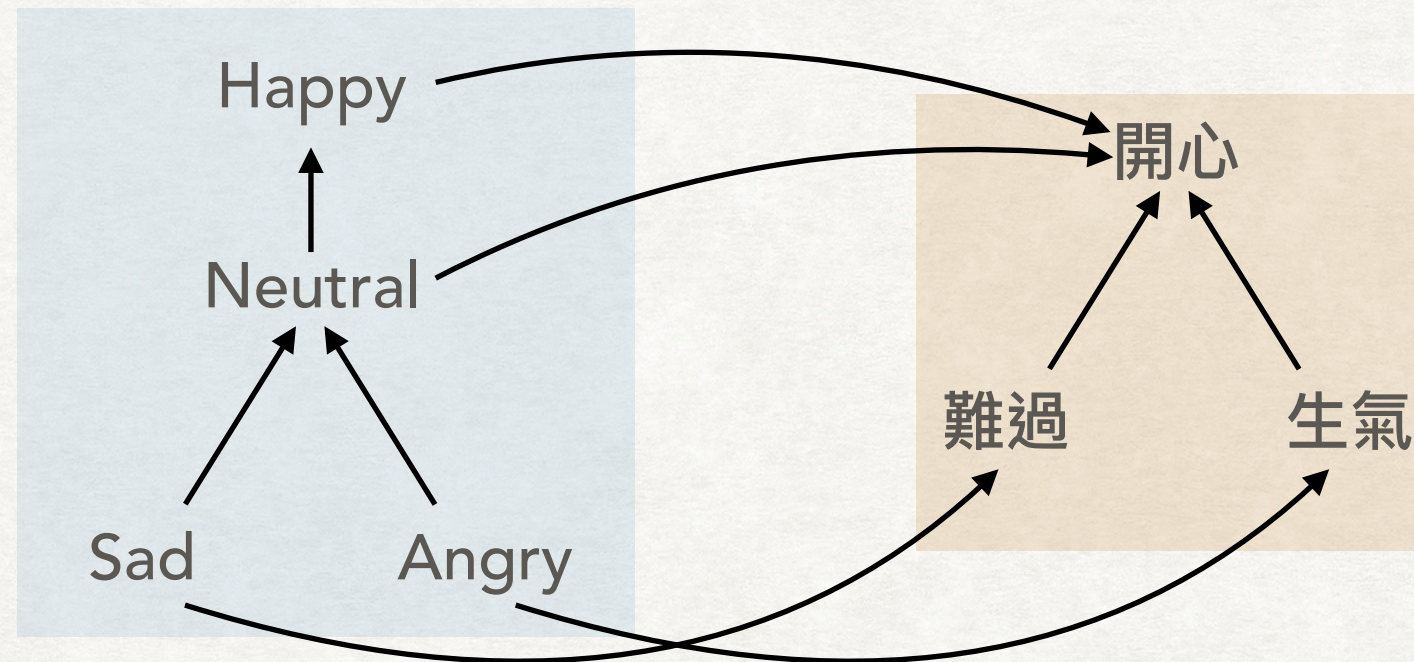
WHAT IS A FUNCTOR? LET'S START WITH A FAMILIAR EXAMPLE



$$F(f; g) = F(f); F(g)$$

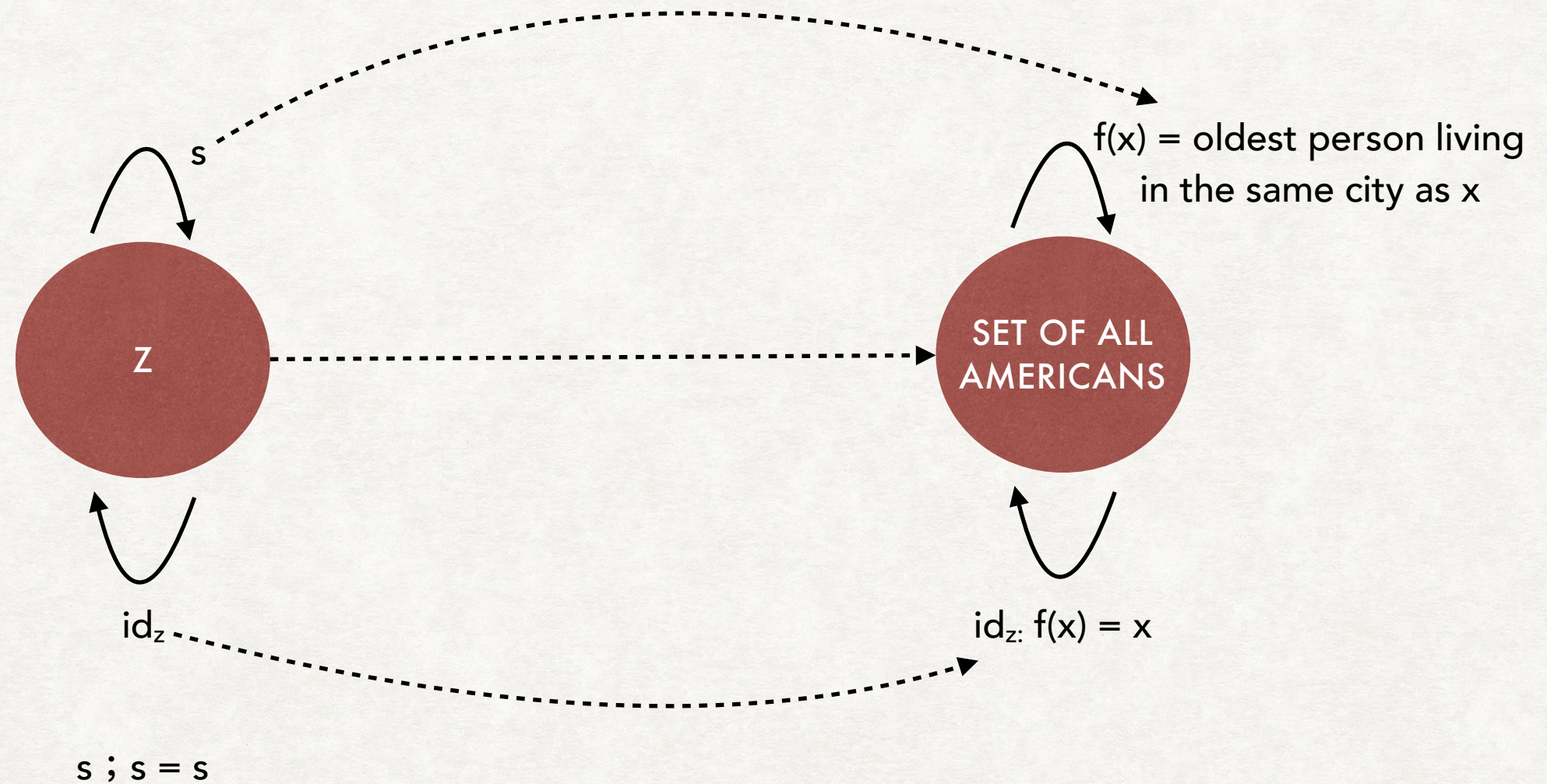
Same rule applies for each object and its identity morphism

WHAT IS A FUNCTOR? LET'S START WITH A FAMILIAR EXAMPLE

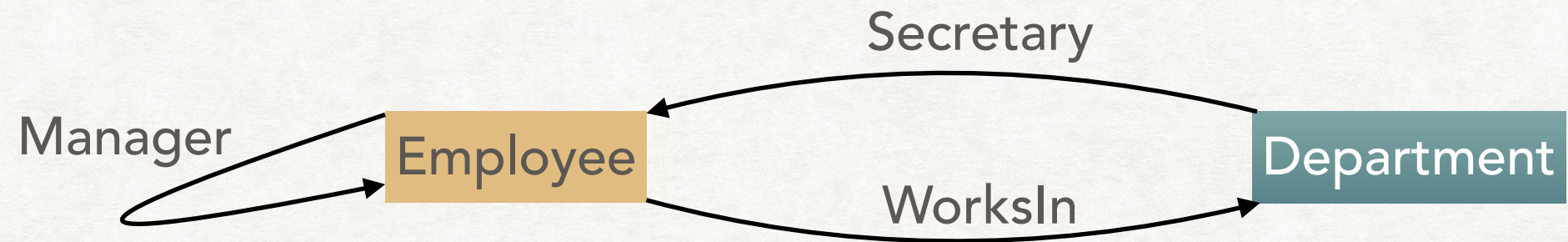


FUNCTORS BETWEEN PARTIAL ORDERS ARE MONOTONE FUNCTIONS

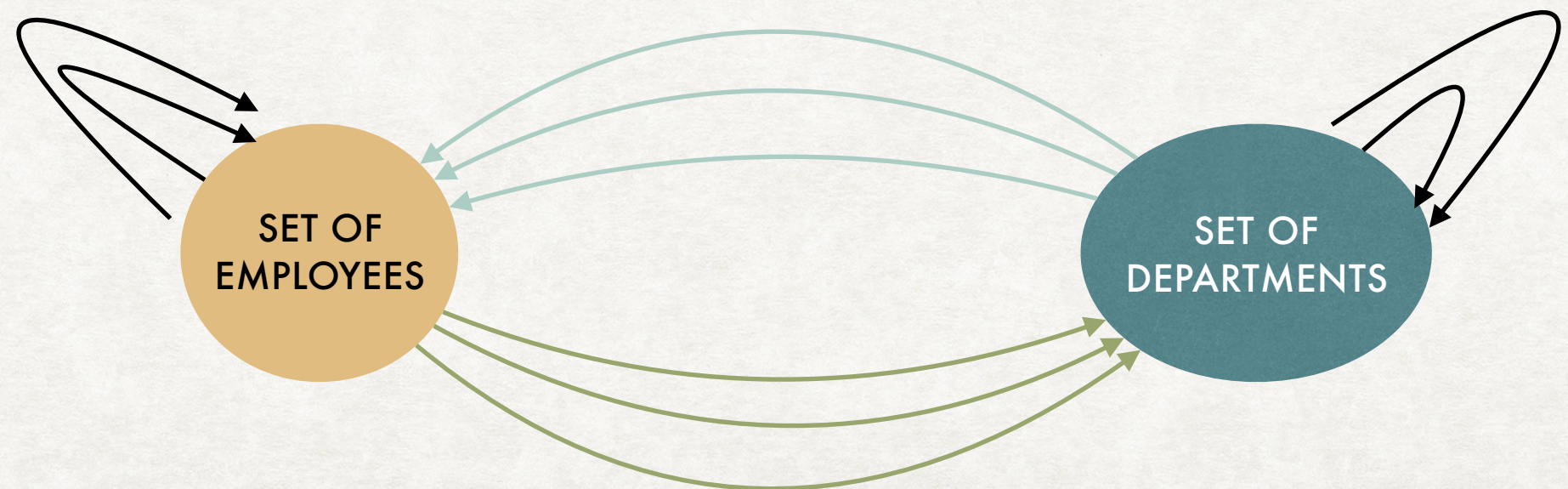
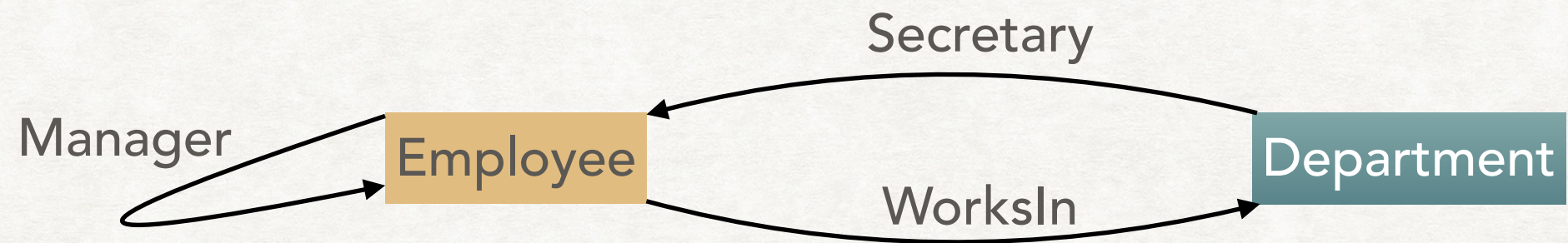
FUNCTOR: ONE MORE EXAMPLE



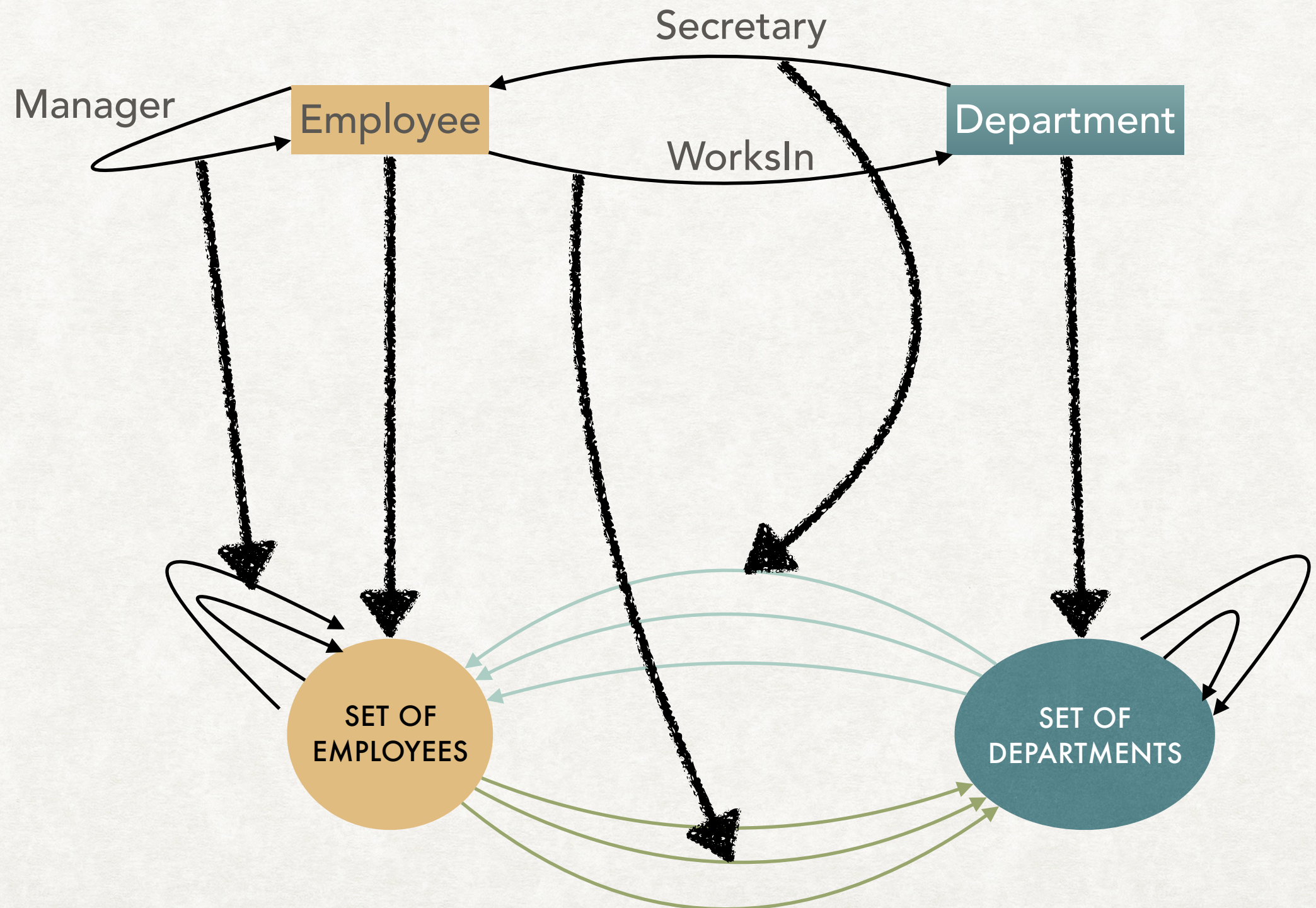
CATEGORY OF DATABASE SCHEMAS



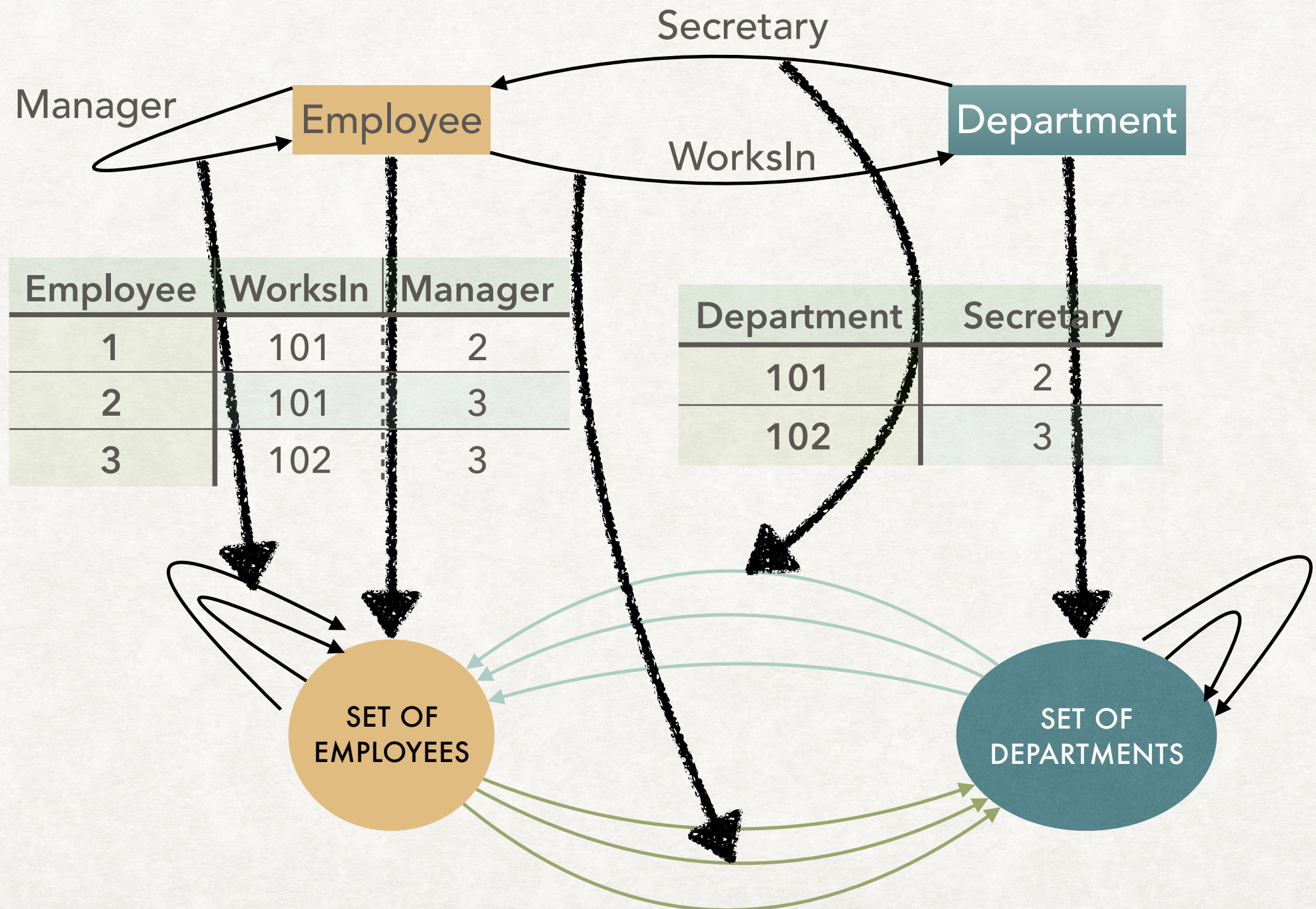
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CATEGORY OF DATABASE SCHEMAS



CATEGORY OF DATABASE SCHEMAS



THANK YOU