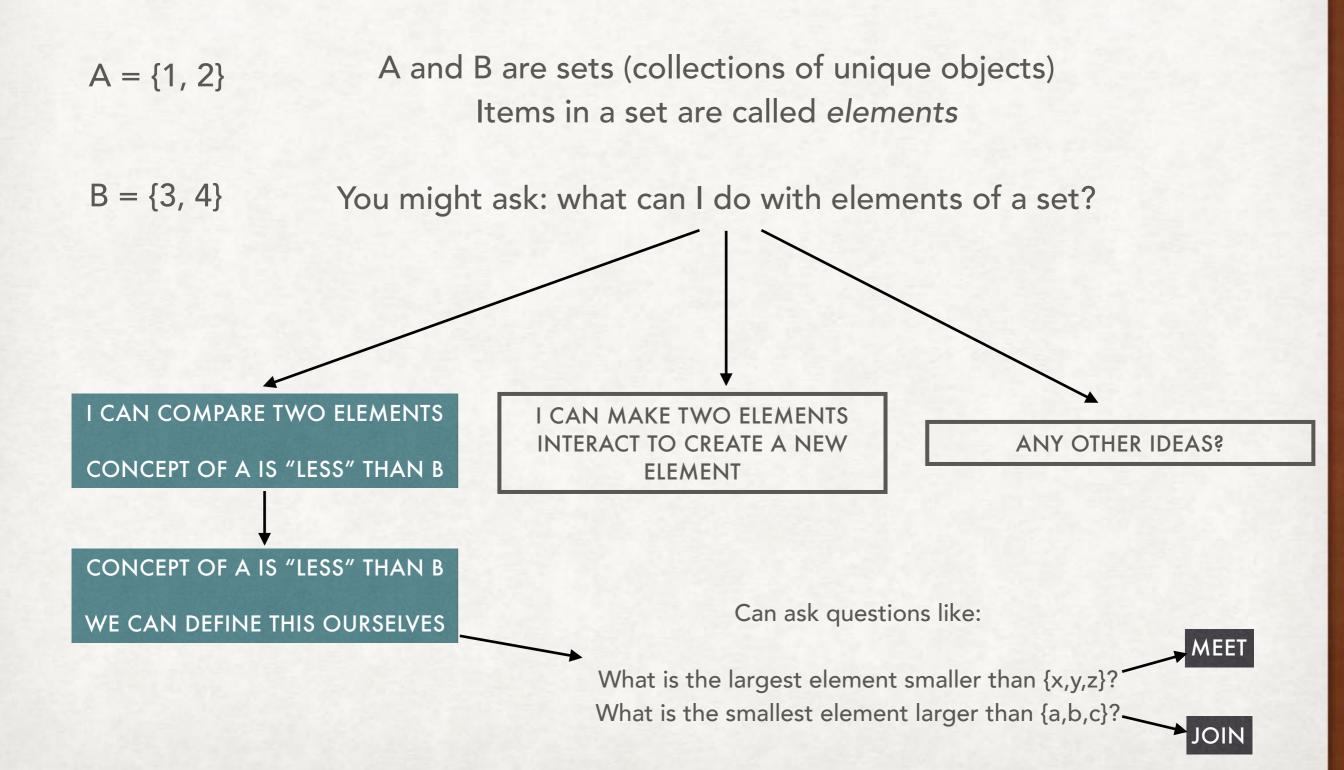
INTRODUCTION TO CATEGORY THEORY

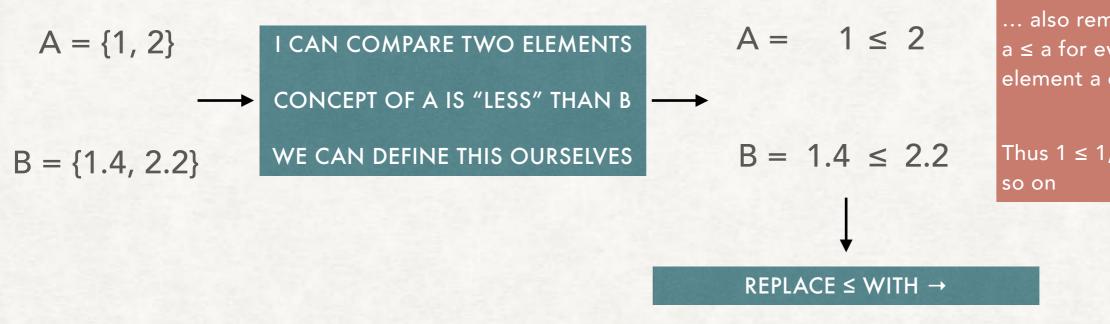
 $A = \{1, 2\}$

A and B are sets (collections of unique objects)
Items in a set are called *elements*

 $B = \{3, 4\}$

You might ask: what can I do with elements of a set?





... also remember that a ≤ a for every element a of a set

Thus $1 \le 1$, $2 \le 2$ and

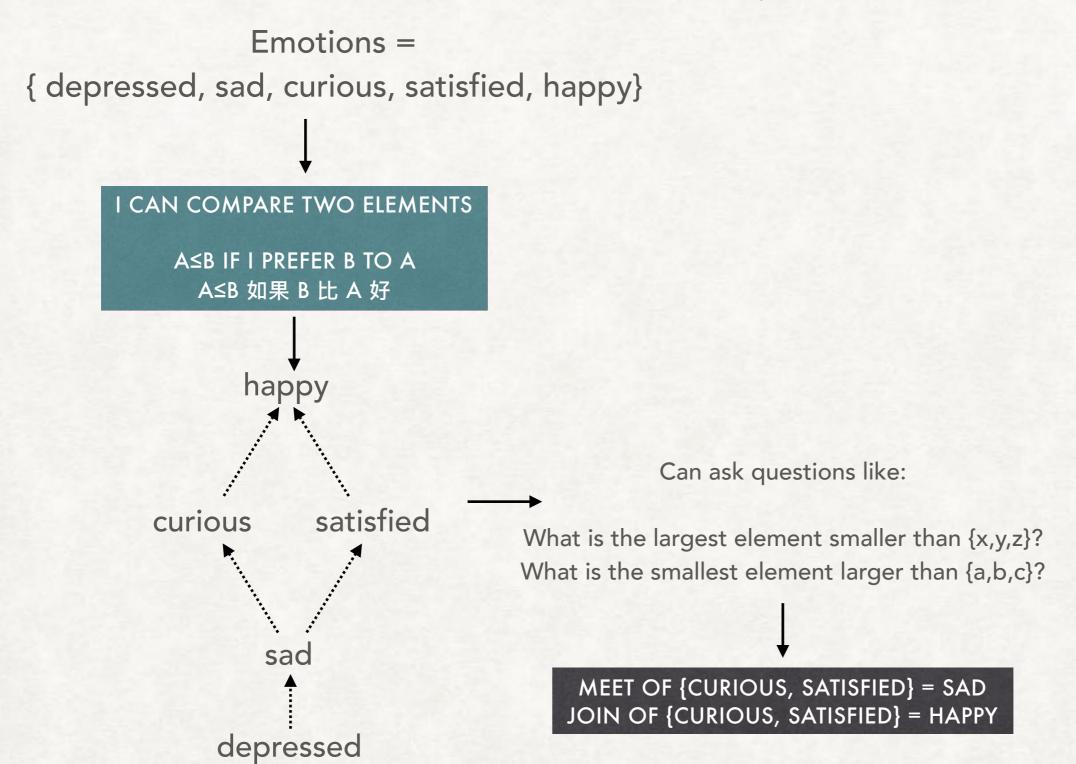
Interpret this as:

WE HAVE TWO SETS A AND B.

WE KNOW WHICH ELEMENTS ARE SMALLER THAN OR EQUAL TO WHICH OTHER ELEMENTS

$$A = 1 \longrightarrow 2$$

$$B = 1.4 \longrightarrow 2.2$$



Emotions = { depressed, sad, curious, satisfied, happy} Remember: I CAN COMPARE TWO ELEMENTS A≤B IF I PREFER B TO A sad ≤ sad, curious ≤ curious A≤B 如果 B 比 A 好 and so on happy Can ask questions like: satisfied curious What is the largest element smaller than $\{x,y,z\}$? What is the smallest element larger than {a,b,c}? sad MEET OF {CURIOUS, SATISFIED} = SAD JOIN OF {CURIOUS, SATISFIED} = HAPPY depressed

EXAMPLE #1

Preorder A $1 \longrightarrow 2$

Find nearest neighbor in preorder B

Preorder B $1.4 \longrightarrow 2.2$

Find nearest neighbor in preorder A

Preorder A $1 \longrightarrow 2$

Preorder B $1.4 \longrightarrow 2.2$

Find nearest neighbor in preorder B

Find nearest neighbor in preorder A

$$l(1) = 1.4$$

$$r(1.4) = 1$$

$$l(2) = 2.2$$

$$l(1) = 1.4$$
 $r(1.4) = 1$
 $l(2) = 2.2$ $r(2.2) = 2$

Preorder A

Find nearest neighbor in preorder B

Preorder B

Find nearest neighbor in preorder A

$$l(1) = 1.4$$

$$l(2) = 2.2$$
 $r(2.2) = 2$

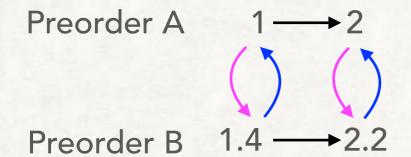
$$l(1) = 1.4$$
 $r(1.4) = 1$

$$r(2.2)=2$$

l and r are inverses of each other

$$r(l(x)) = x$$

$$l(r(y)) = y$$

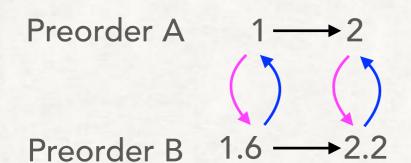


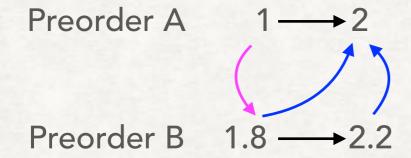
Preorder A
$$1 \longrightarrow 2$$

Find nearest neighbor in preorder B

Preorder B
$$1.8 \longrightarrow 2.2$$

Find nearest neighbor in preorder A





Find nearest neighbor in preorder B

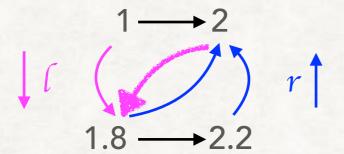
Find nearest neighbor in preorder A

$$l(1) = 1.8$$

$$r(1.8) = 2$$

$$l(2) = ?$$
 $r(2.2) = 2$

Note that the function r is not invertible here



$$l(1) = 1.8$$
 $r(1.8) = 2$
 $l(2) = 1.8$ $r(2.2) = 2$

$$r(1.8) = 2$$
$$r(2.2) = 2$$

$$\begin{array}{ccc}
1 & \longrightarrow 2 \\
\downarrow & & \uparrow \\
1.8 & \longrightarrow 2.2
\end{array}$$

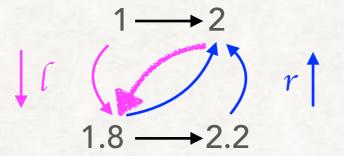
$$l(1) = 1.8$$
 $r(1.8) = 2$

$$r(1.8)=2$$

$$l(2) = 2.2$$

$$f(2) = 2.2$$
 $r(2.2) = 2$

WHICH ONE IS A BETTER CHOICE\$



$$l(1) = 1.8$$
 $r(1.8) = 2$ $l(2) = 1.8$ $r(2.2) = 2$

$$r(1.8) = 2$$
$$r(2.2) = 2$$

$$\begin{array}{c} 1 \longrightarrow 2 \\ \downarrow \\ 1.8 \longrightarrow 2.2 \end{array}$$

$$l(1) = 1.8$$

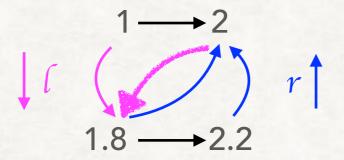
$$f(2) = 2.2$$

$$r(1.8)=2$$

$$r(2.2)=2$$

WHICH ONE IS A BETTER CHOICE\$

FUNCTIONS L AND R SHOULD PRESERVE SOMETHING FROM ISOMORPHIC MAPS



$$l(1) = 1.8$$
 $r(1.8) = 2$ $l(2) = 1.8$ $r(2.2) = 2$

$$1 \longrightarrow 2$$

$$\downarrow () r \uparrow$$

$$1.8 \longrightarrow 2.2$$

$$l(1) = 1.8 r(1.8) = 2$$

$$l(2) = 2.2 r(2.2) = 2$$

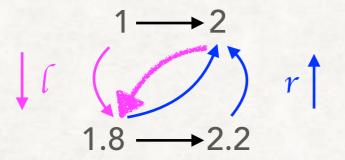
WHICH ONE IS

A BETTER

CHOICE?

FUNCTIONS L AND R SHOULD PRESERVE SOMETHING FROM ISOMORPHIC MAPS

 $L(X) \leq Y$ \Leftrightarrow $X \leq R(Y)$



$$l(1) = 1.8$$

$$\underline{l(2) = 1.8}$$

$$r(1.8) = 2$$
$$r(2.2) = 2$$

$$2 \le g(1.8) \Leftrightarrow f(2) \le 1.8$$

$$\begin{array}{ccc}
1 \longrightarrow 2 \\
\downarrow & & \uparrow \\
1.8 \longrightarrow 2.2
\end{array}$$

$$l(1) = 1.8$$

$$l(2) = 2.2$$

$$r(1.8)=2$$

$$r(2.2) = 2$$

 $2 \le g(1.8)$ but $f(2) \le 1.8$?

WHICH ONE IS

A BETTER

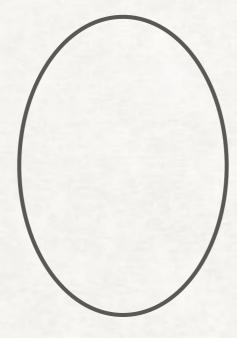
CHOICE?

FUNCTIONS L AND R SHOULD PRESERVE SOMETHING FROM ISOMORPHIC MAPS

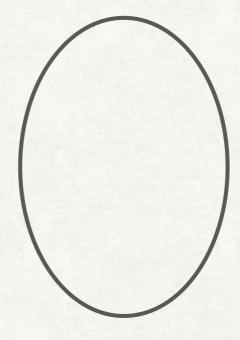
Aims to minimize disagreement

DEFINITION

Given preorders P and Q



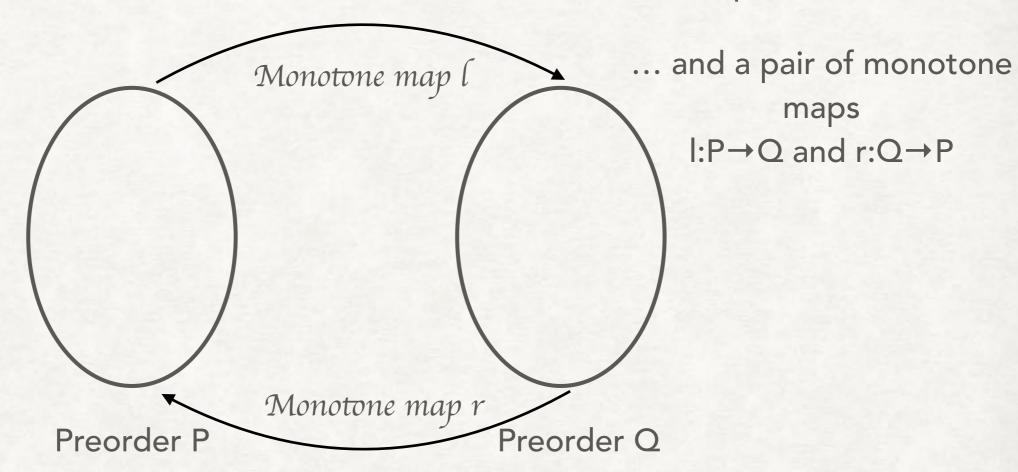
Preorder P



Preorder Q

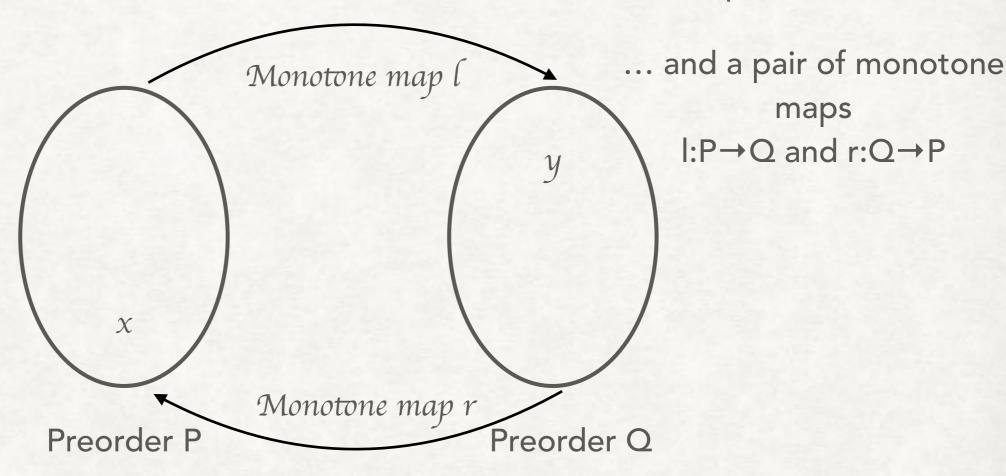
DEFINITION

Given preorders P and Q



DEFINITION

Given preorders P and Q

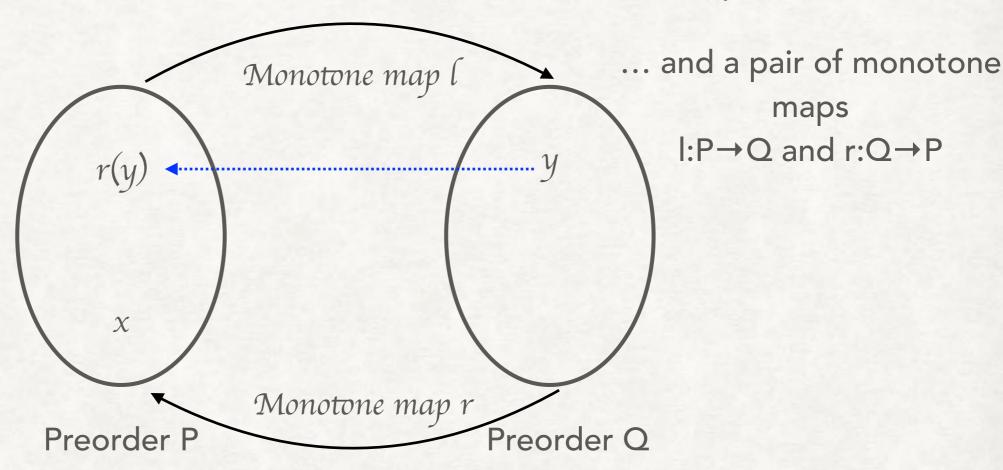


$$x \le r(y) \Leftrightarrow \ell(x) \le y$$

for all $x \in \mathcal{P}$, $y \in \mathcal{Q}$

DEFINITION

Given preorders P and Q

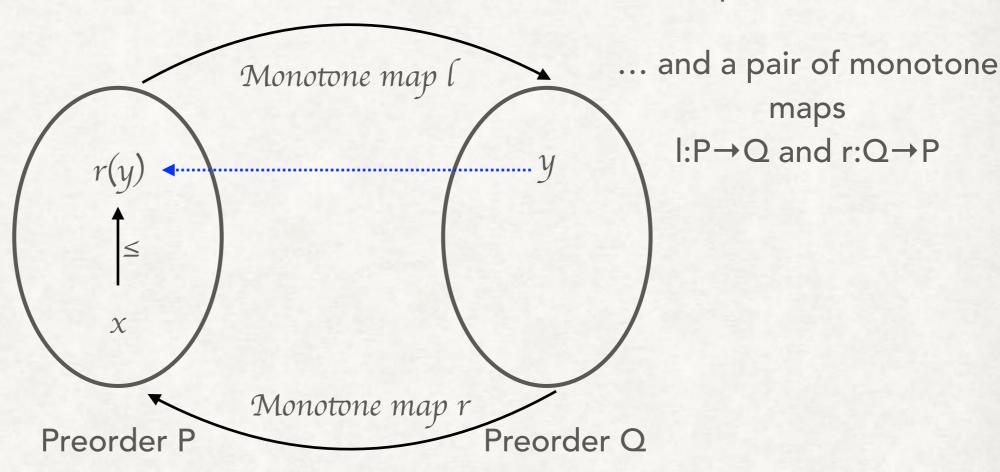


$$x \le r(y) \Leftrightarrow \ell(x) \le y$$

for all $x \in \mathcal{P}$, $y \in \mathcal{Q}$

DEFINITION

Given preorders P and Q

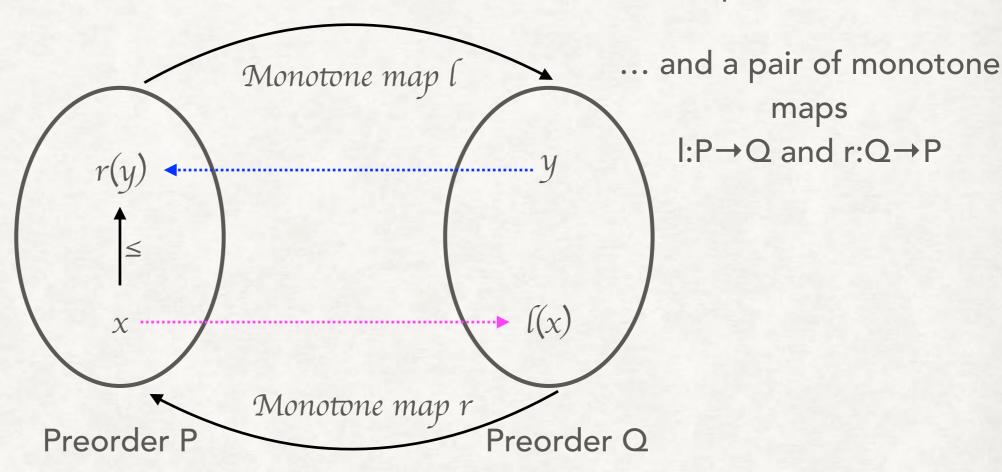


$$x \le r(y) \Leftrightarrow \ell(x) \le y$$

for all $x \in \mathcal{P}$, $y \in \mathcal{Q}$

DEFINITION

Given preorders P and Q

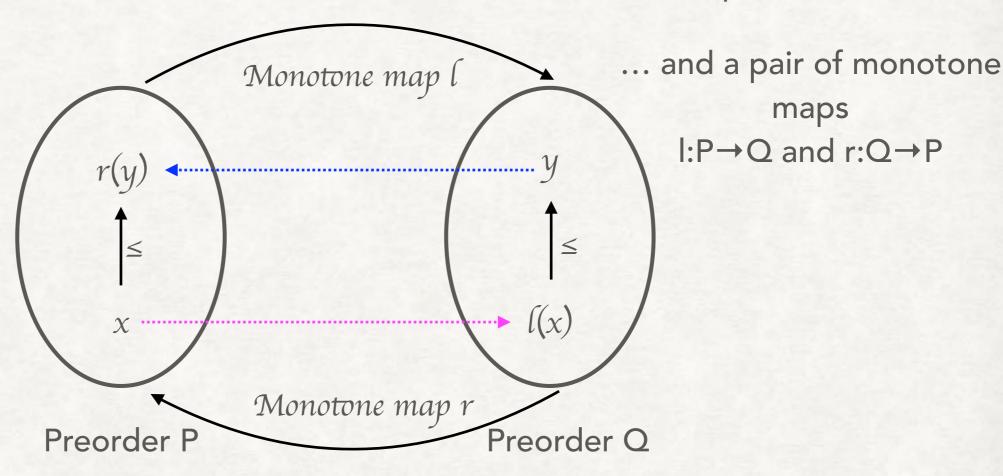


$$x \le r(y) \Leftrightarrow \ell(x) \le y$$

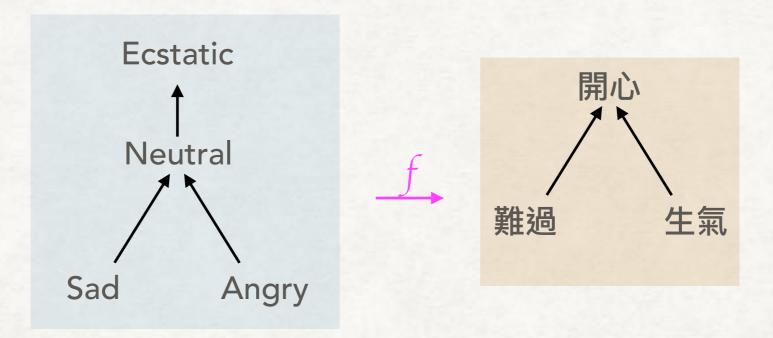
for all $x \in P$, $y \in Q$

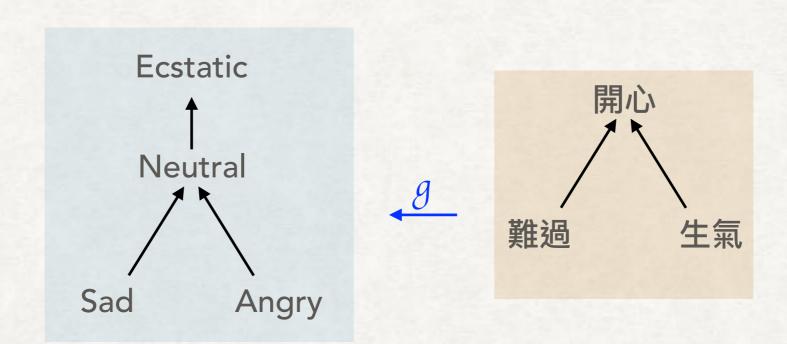
DEFINITION

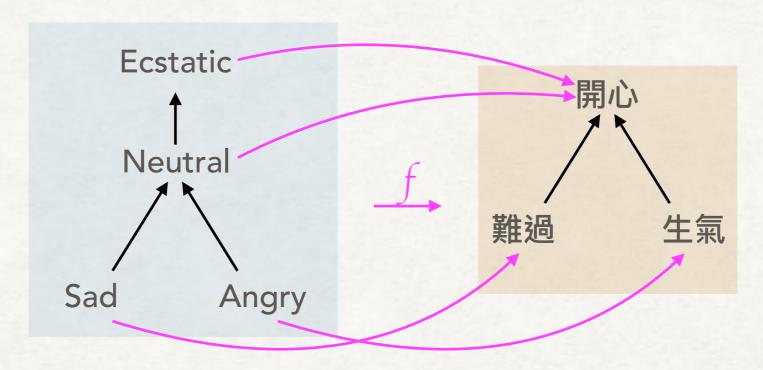
Given preorders P and Q

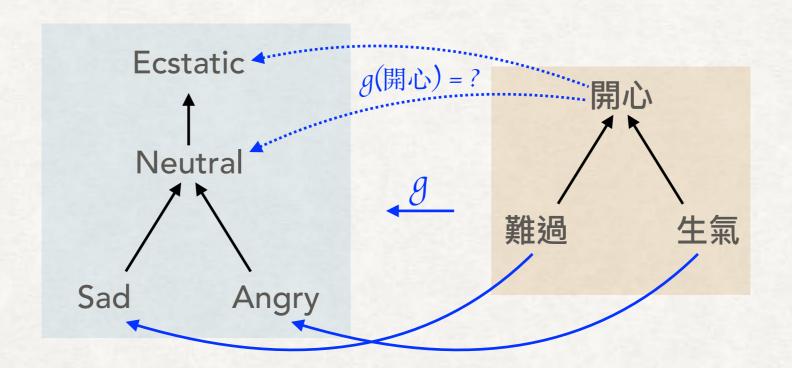


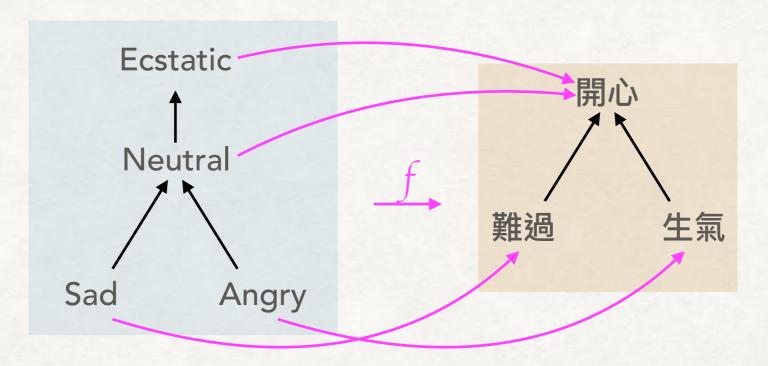
$$x \le r(y) \Leftrightarrow \ell(x) \le y$$
for all $x \in \mathcal{P}, y \in \mathcal{Q}$

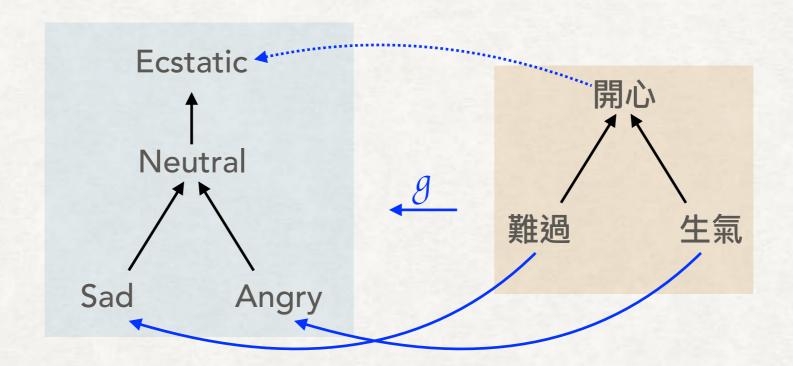


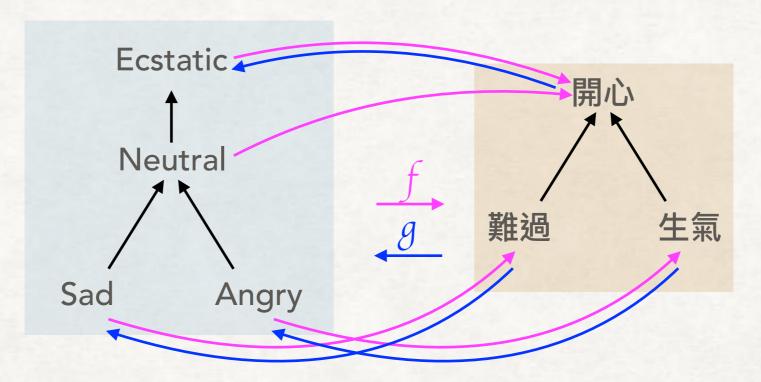












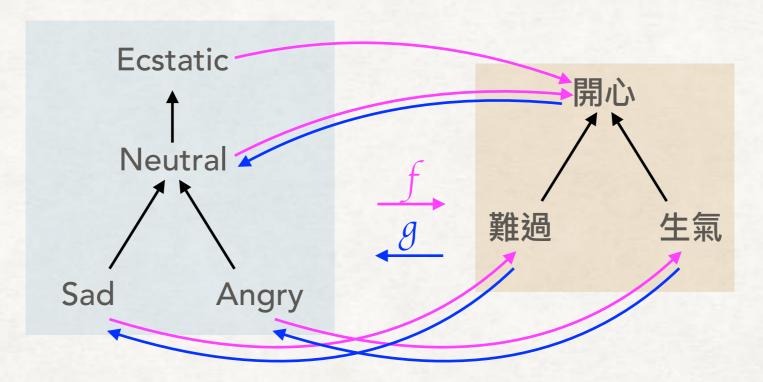
 $neutral \leq g(開心) \Leftrightarrow f(neutral) \leq 開心$

Obvious for all other terms:

 $sad \leq g(難過) \Leftrightarrow f(sad) \leq 難過$

... and so on

f: left adjoint g: right adjoint



 $g(開心) \leq ecstatic \Leftrightarrow 開心 \leq f(ecstatic)$

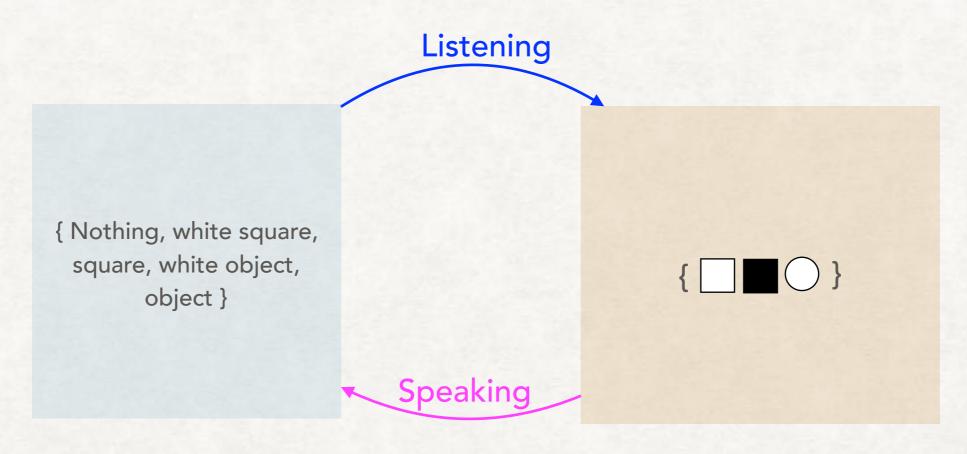
Obvious for all other terms:

 $g(難過) \leq sad \Leftrightarrow 難過 \leq f(sad)$

... and so on

f: right adjoint

g: left adjoint

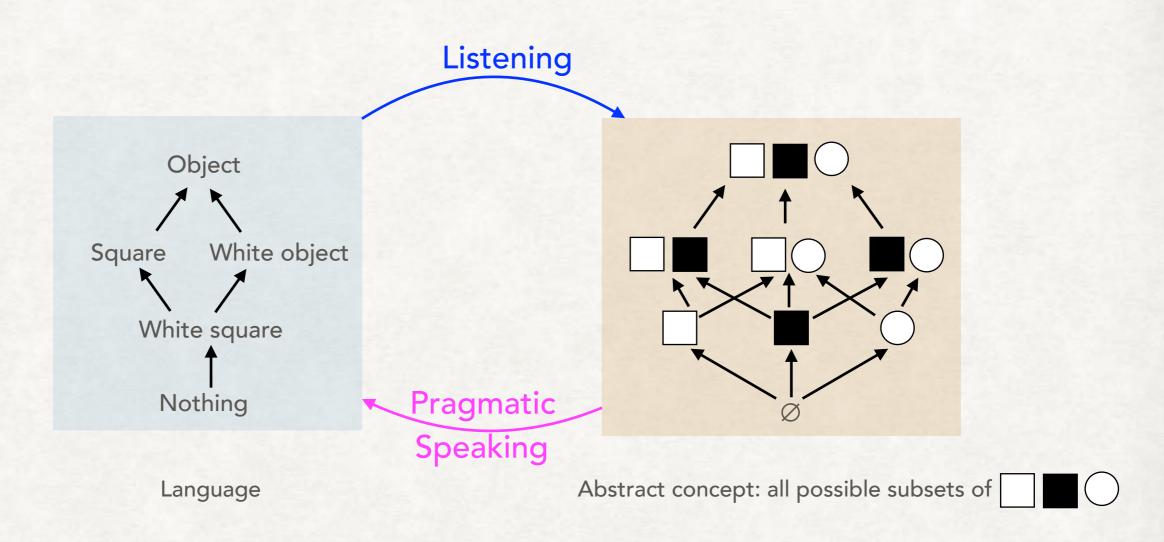


Language

Abstract concept: all possible subsets of

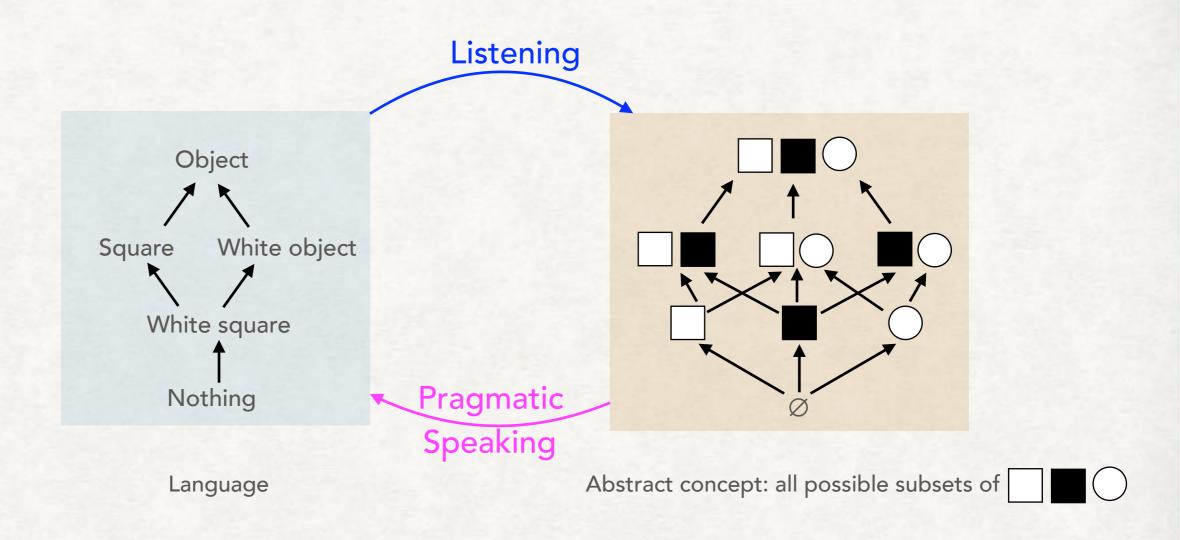
Define \leq : $a \leq b$ if a is subset of b

Define \leq : $a \leq b$ if a implies b



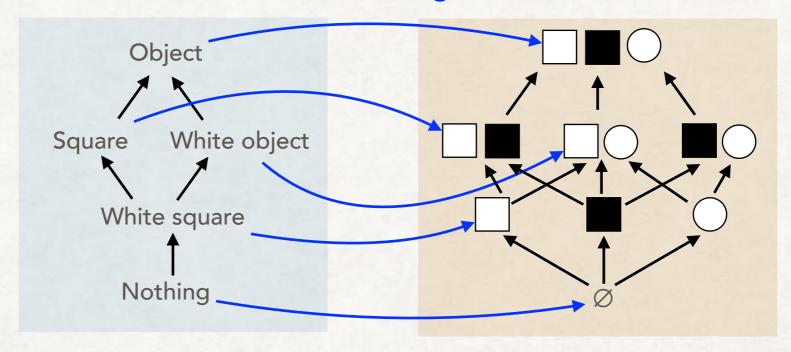
Define \leq : $a \leq b$ if a implies b

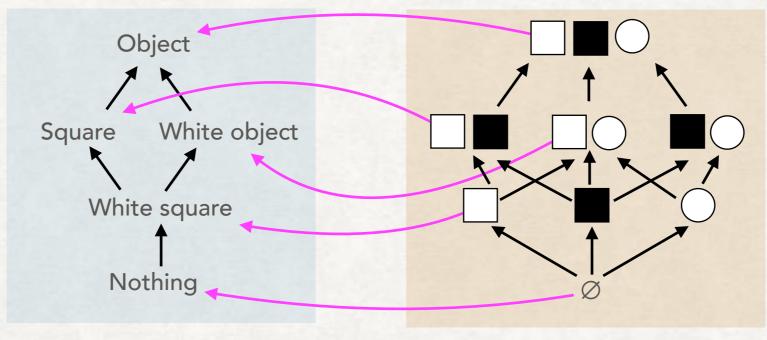
Define \leq : $a \leq b$ if a is subset of b



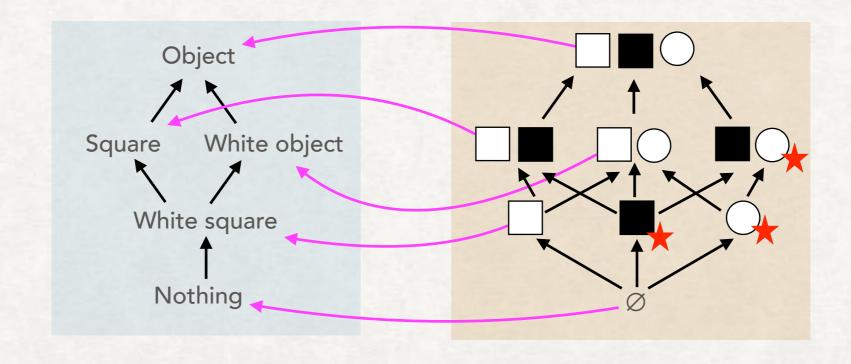
Pragmatic means the speaker wants to be correct and as possible
Quantity
Quality

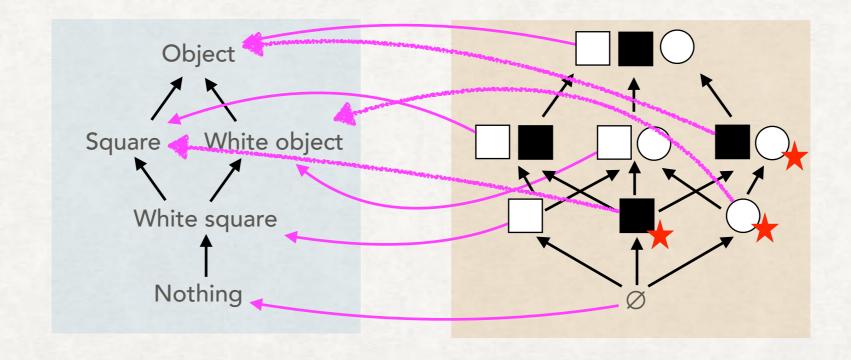
Listening



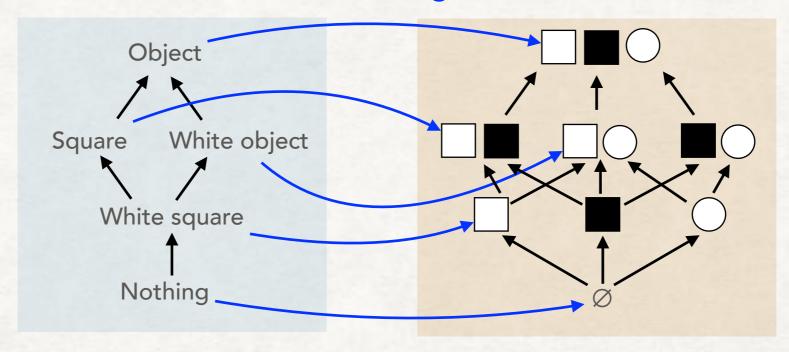


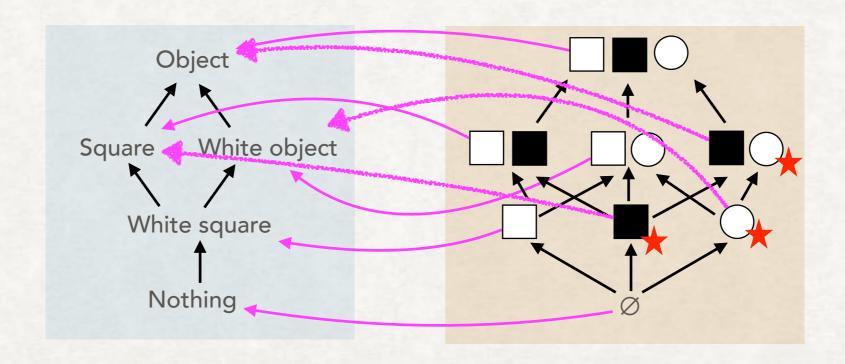
Speaking





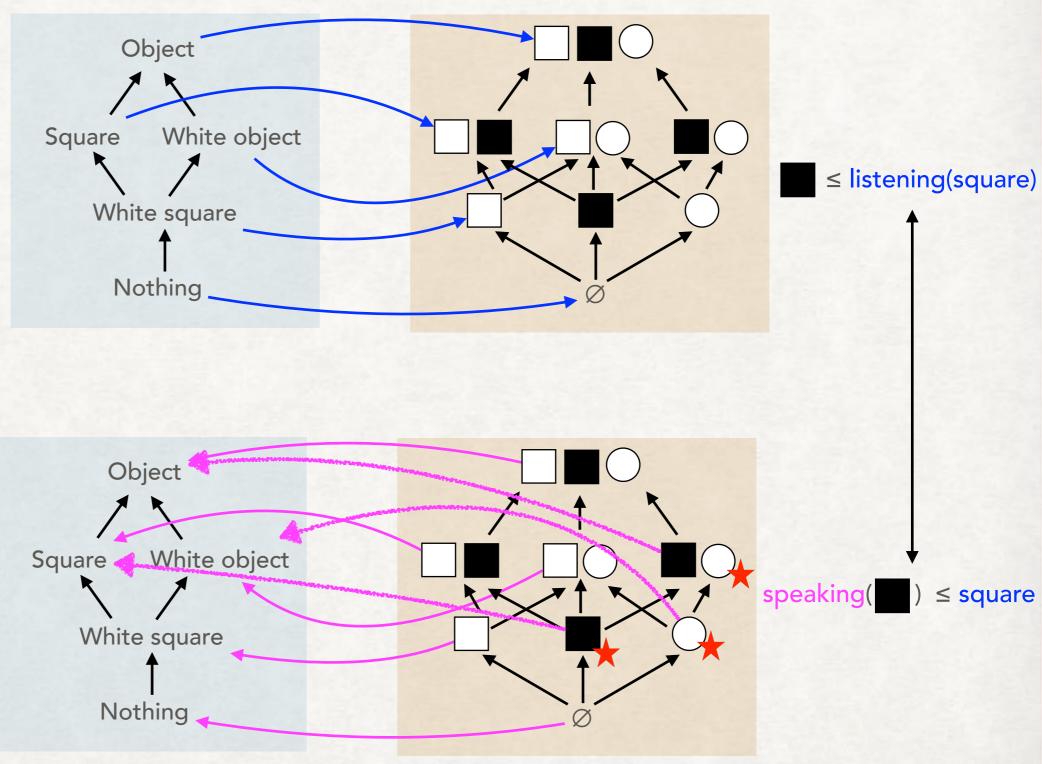
Listening





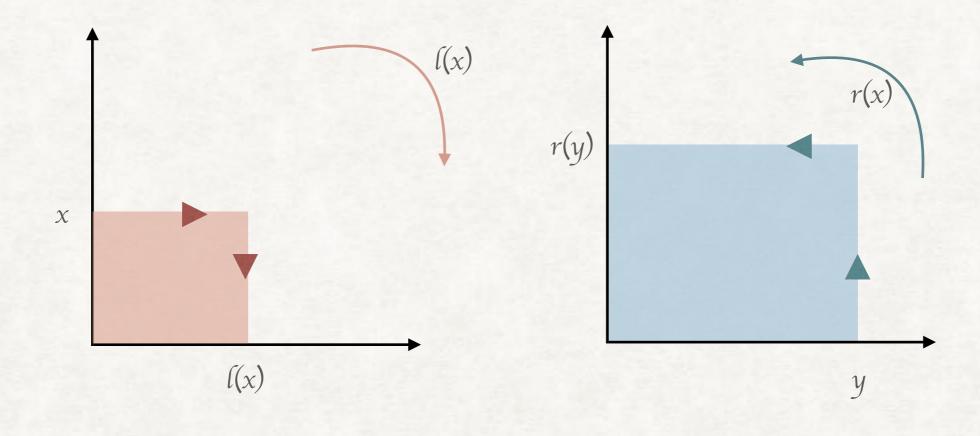
Speaking

Listening → Right adjoint



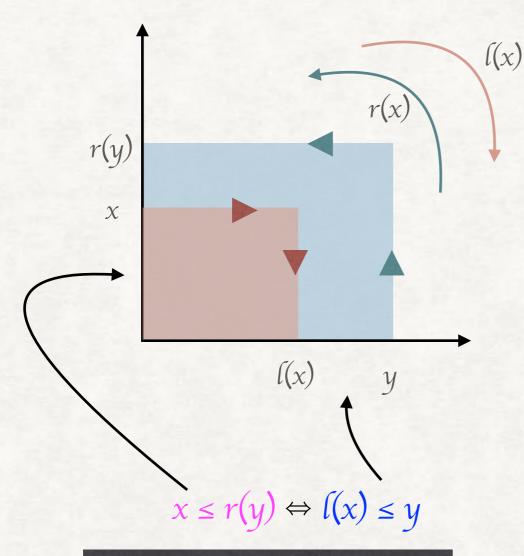
Speaking → Left adjoint

VISUAL EXPLANATION OF GALOIS CONNECTION



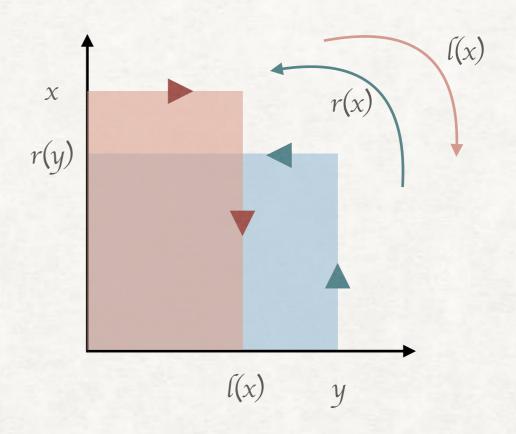
Right adjoint

Left adjoint



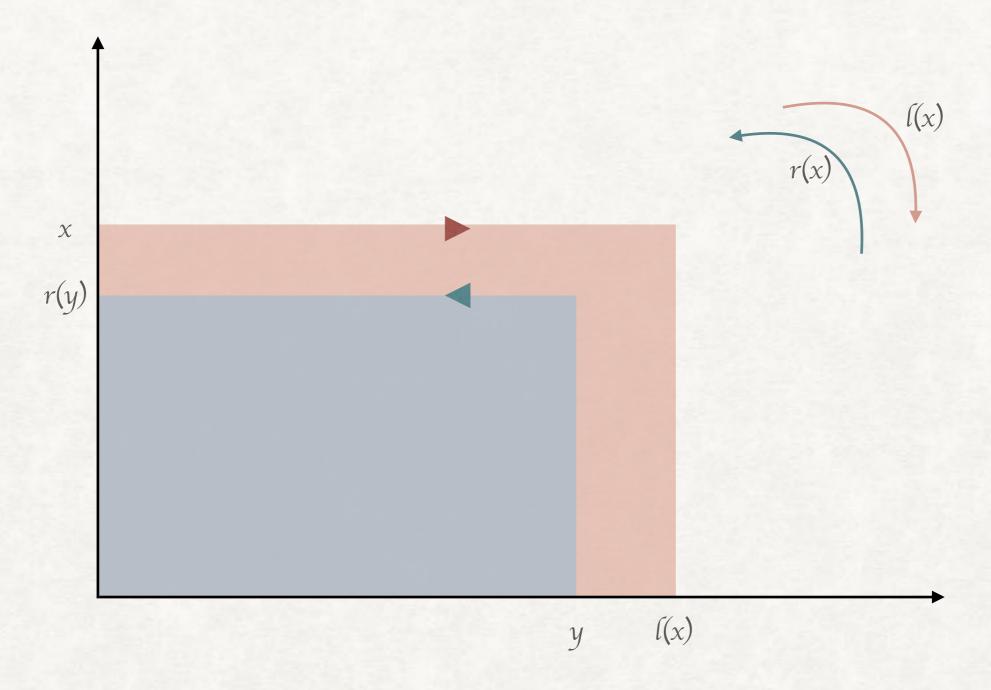
RECTANGLES DO NOT INTERSECT

RECTANGLE_{LEFT-ADJOINT}
CONTAINED BY
RECTANGLE_{RIGHT-ADJOINT}

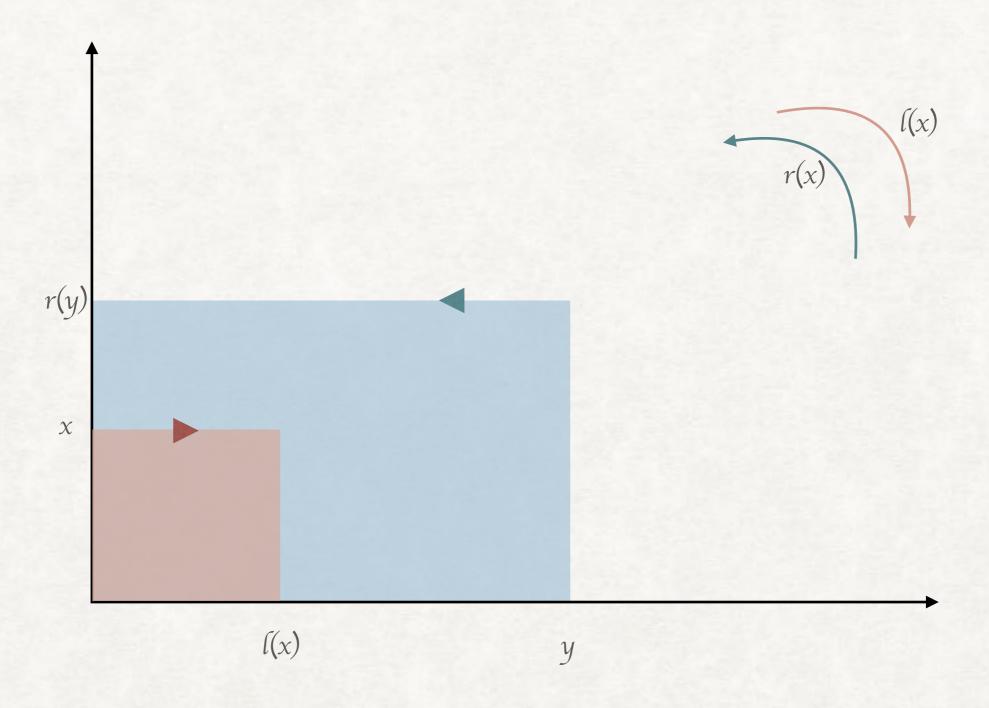


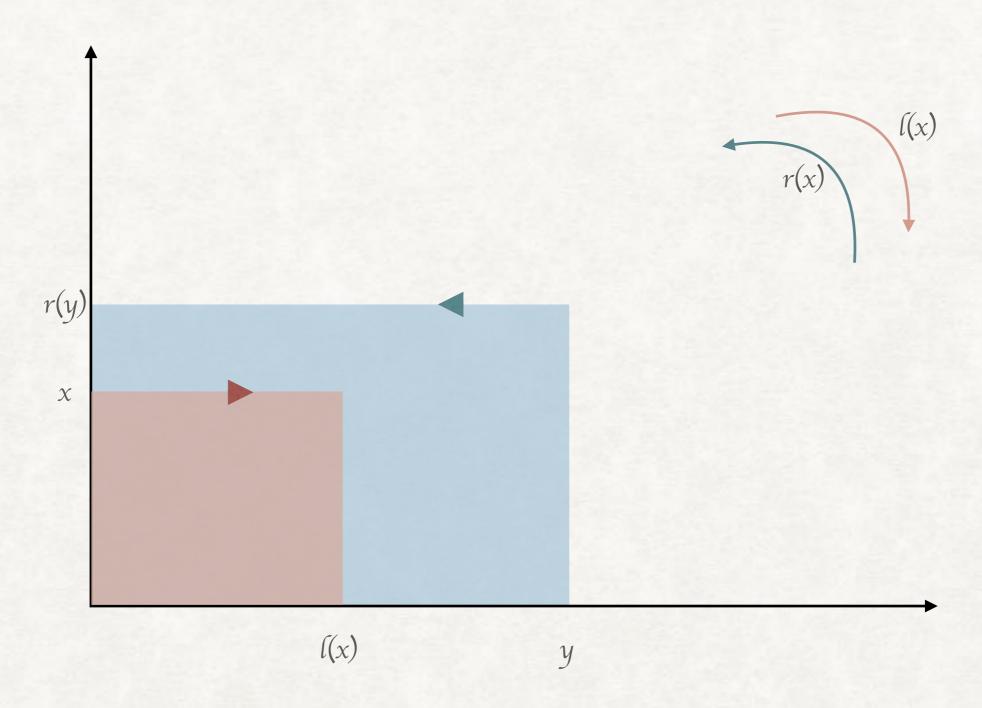
No longer true: $x \le r(y) \Leftrightarrow l(x) \le y$

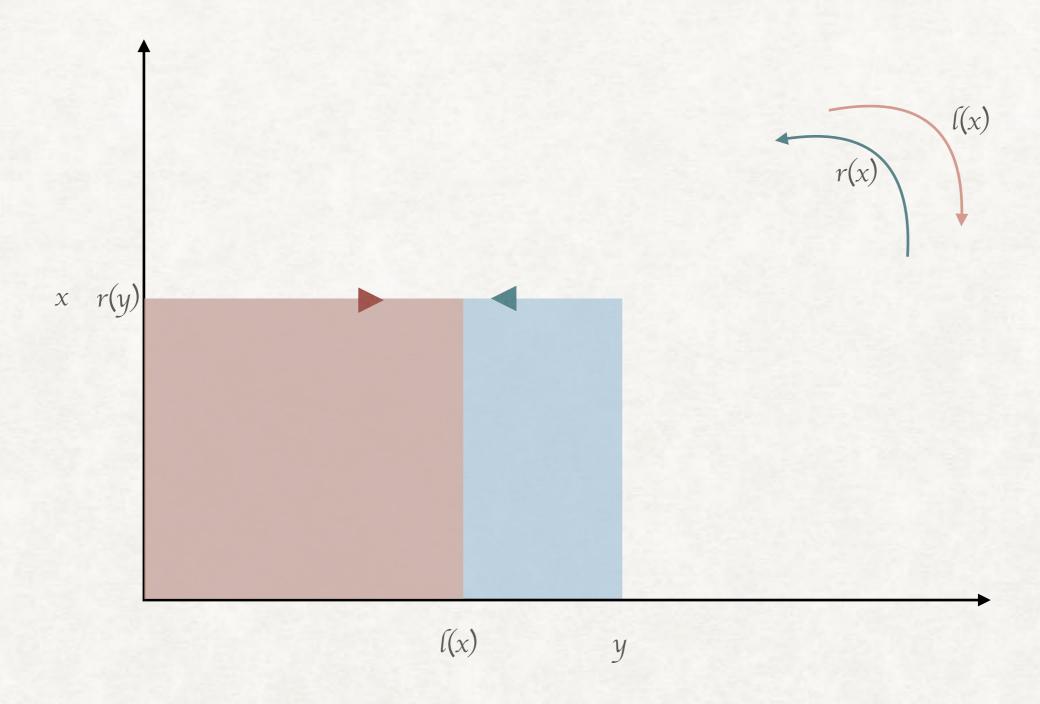
RECTANGLES INTERSECT



Is this a Galois connection?

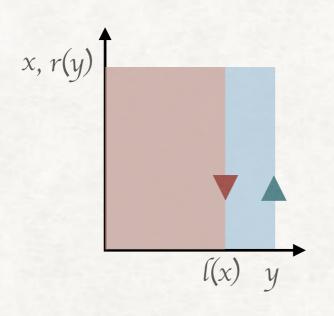


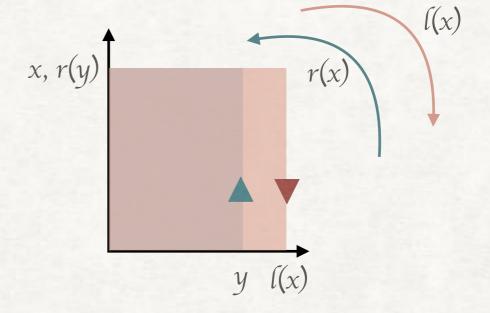


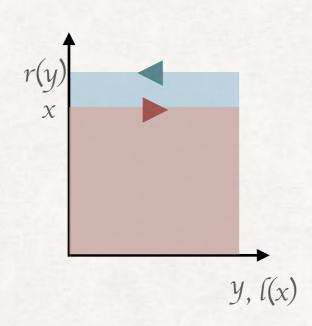


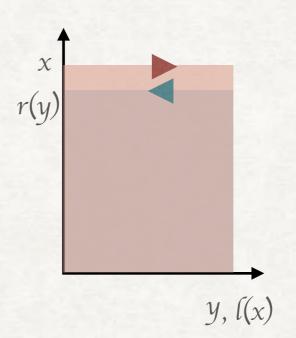
WHICH ONES ARE GALOIS CONNECTIONS?

$$x \le r(y) \Leftrightarrow \ell(x) \le y$$



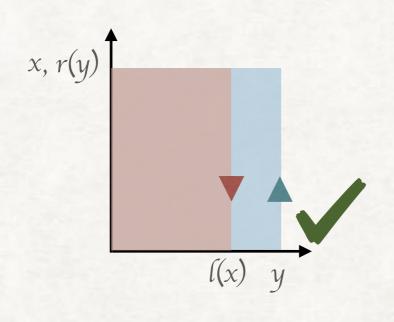


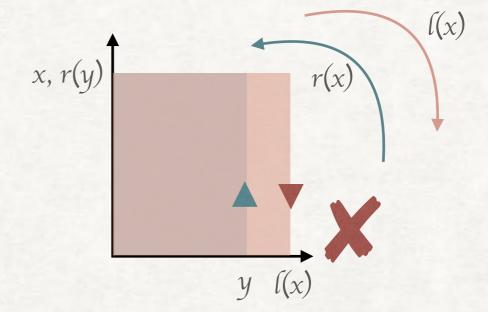


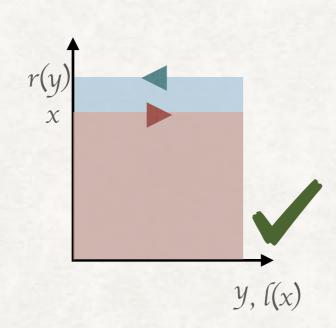


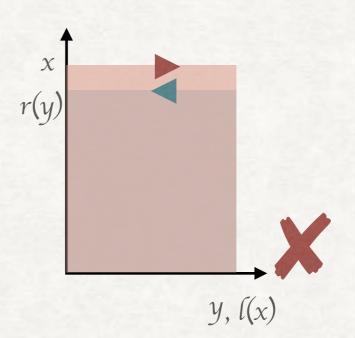
WHICH ONES ARE GALOIS CONNECTIONS?

$$x \le r(y) \Leftrightarrow \ell(x) \le y$$



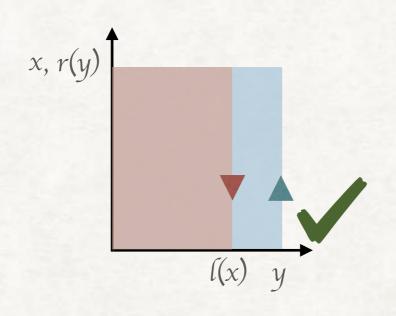


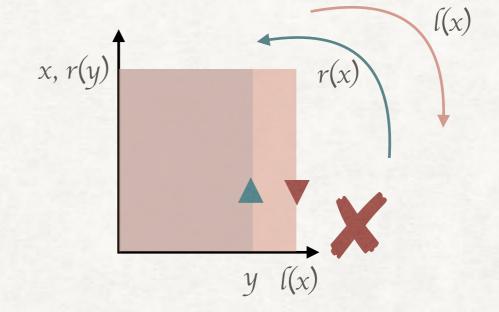


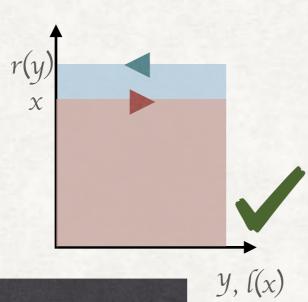


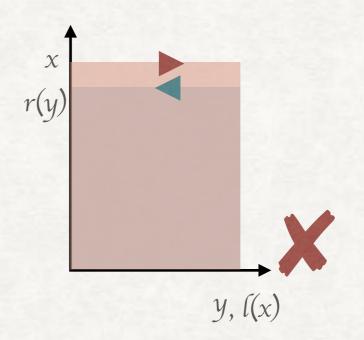
WHICH ONES ARE GALOIS CONNECTIONS?

$$x \le r(y) \Leftrightarrow \ell(x) \le y$$



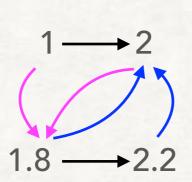


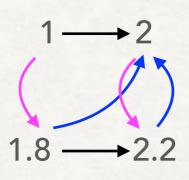


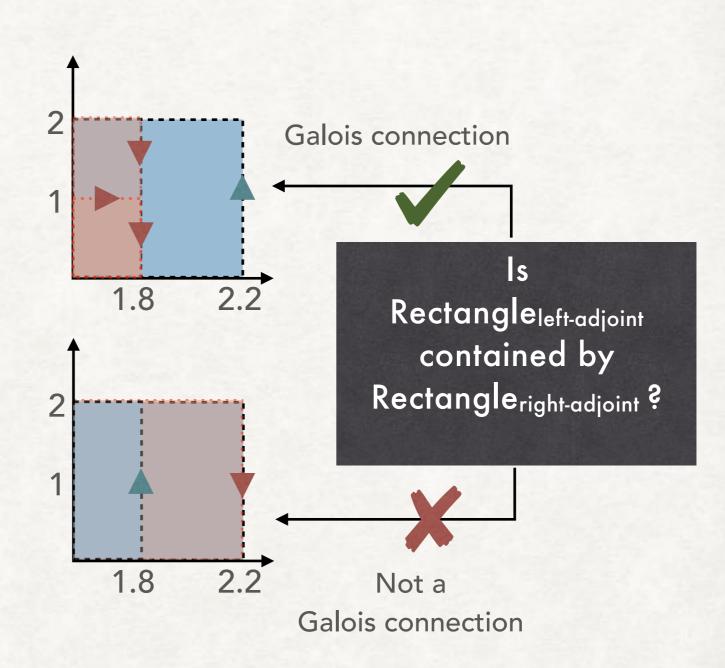


Rectangle_{left-adjoint} should share an edge with Rectangle_{right-adjoint} from the inside

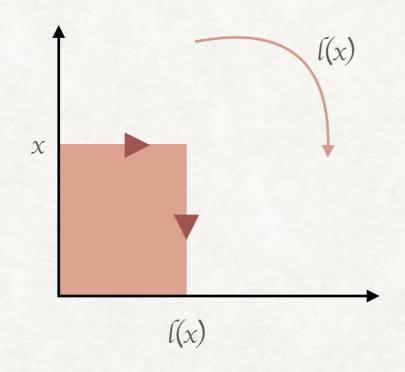
LOOKING BACK AT EXAMPLE #1

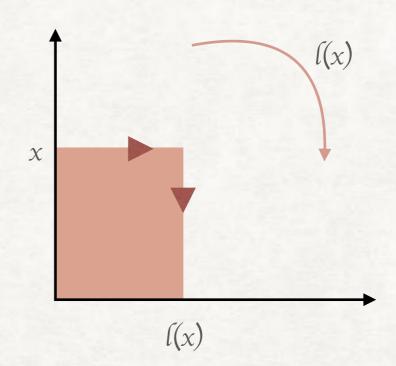




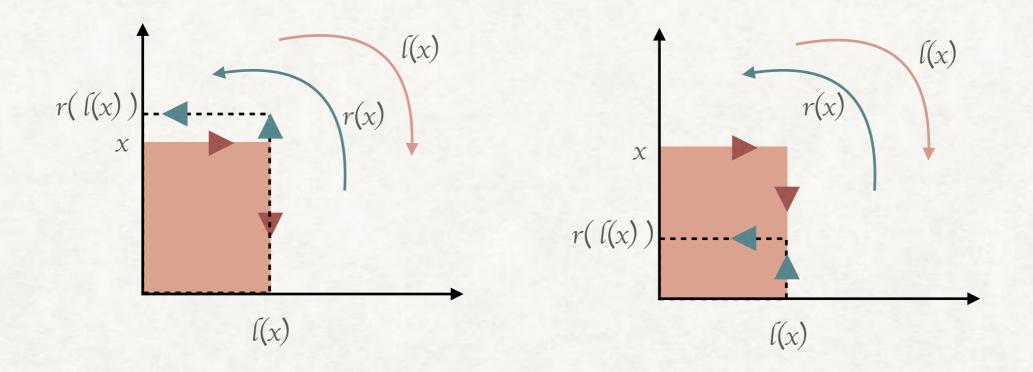


LEFT ADJOINT THEN RIGHT ADJOINT OF AN ELEMENT



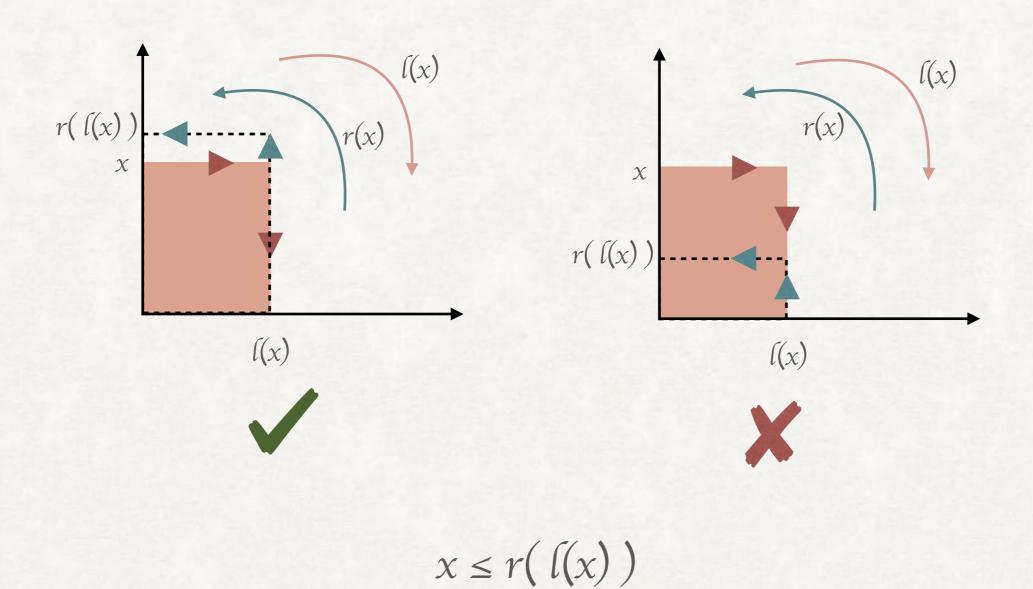


LEFT ADJOINT THEN RIGHT ADJOINT OF AN ELEMENT

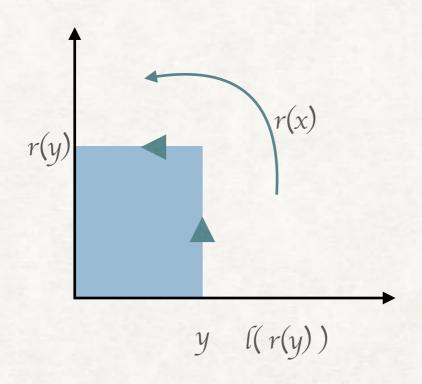


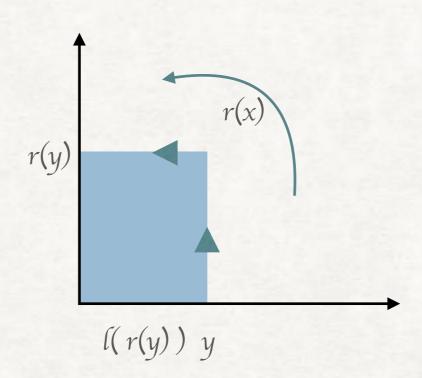
Which one is the right adjoint of a Galois connection?

LEFT ADJOINT THEN RIGHT ADJOINT OF AN ELEMENT

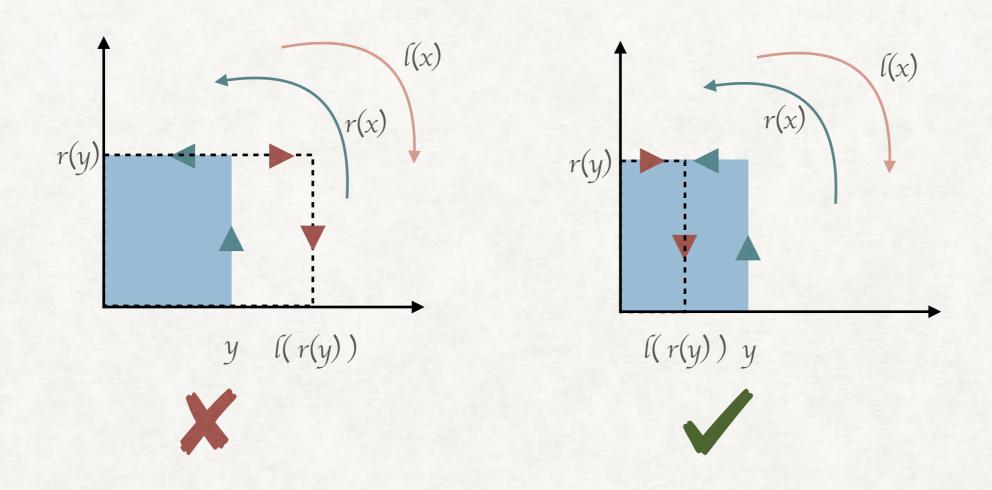


RIGHT ADJOINT THEN LEFT ADJOINT OF AN ELEMENT



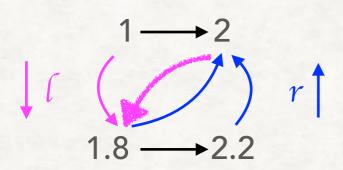


RIGHT ADJOINT THEN LEFT ADJOINT OF AN ELEMENT

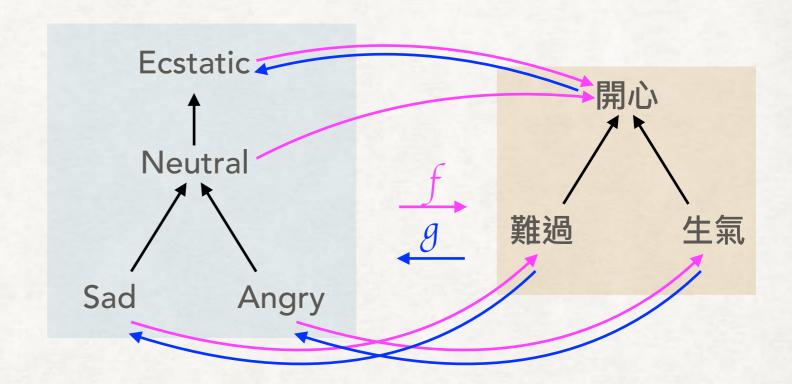


$$l(r(y)) \leq y$$

LOOKING BACK AT EXAMPLES #1 AND #2



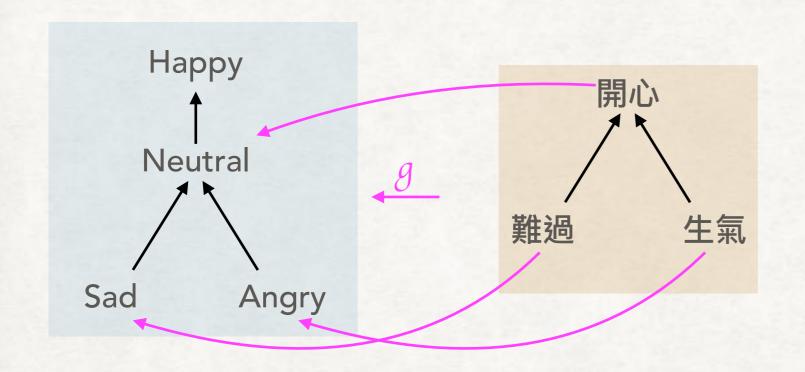
$$1 \le r(l(1))$$
$$l(r(2.2)) \le 2.2$$



$$neutral \leq g(f(neutral))$$

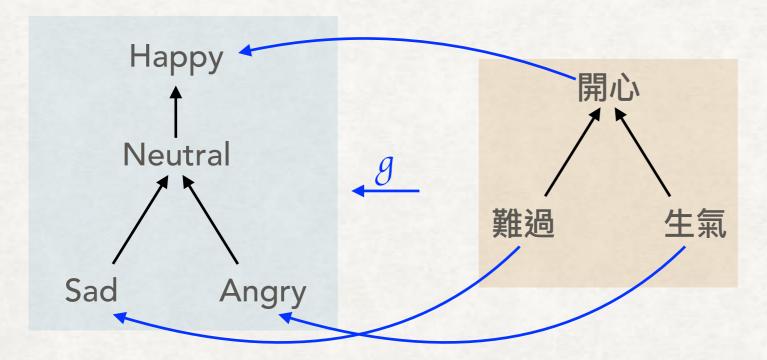
$$f(g(難過)) \leq 難過$$

$$x \le r(l(x))$$
 $l(r(y)) \le y$



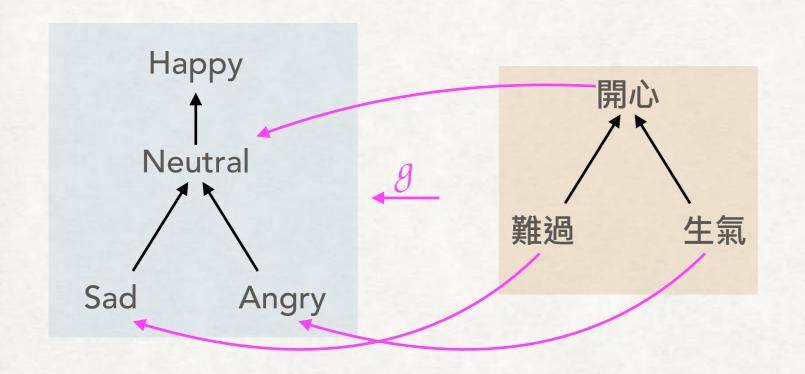
g (join(難過, 生氣))

join(g(難過),g(生氣))

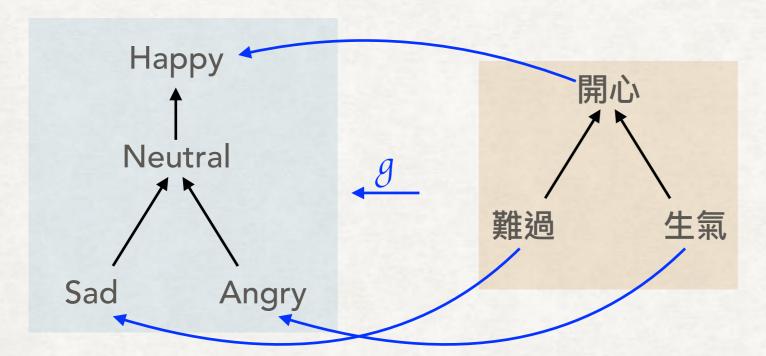


g (join(難過, 生氣))

join(g(難過),g(生氣))



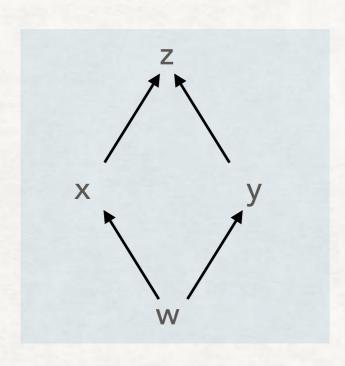
g (join(難過, 生氣)) = join(g(難過), g(生氣))



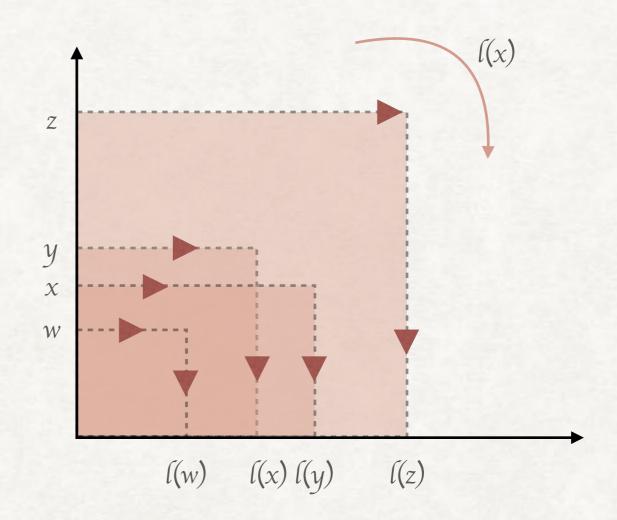
g (join(難過, 生氣))

≠
join(g(難過), g(生氣))

SIMILARLY RIGHT ADJOINT PRESERVES MEETS

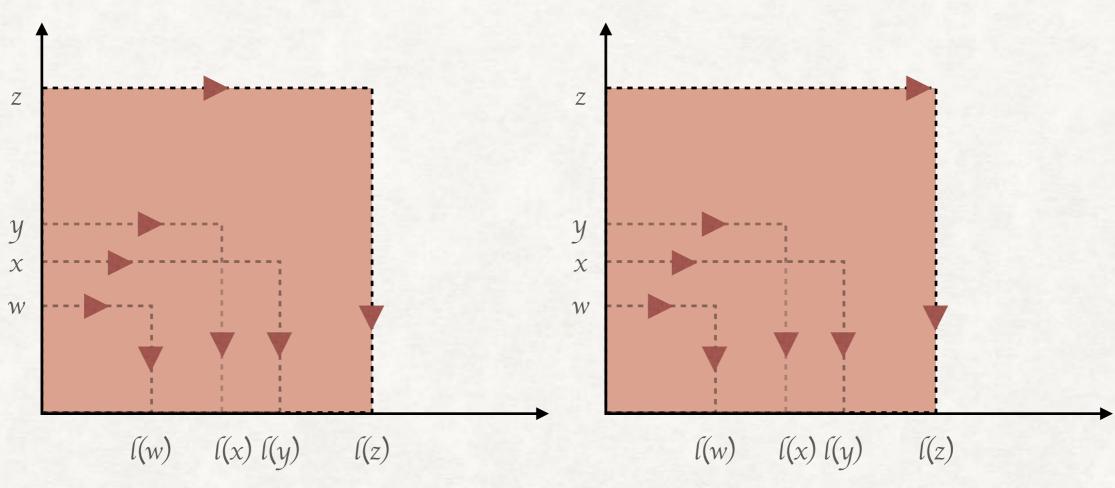


Meet(x, y) = wJoin(x, y) = z

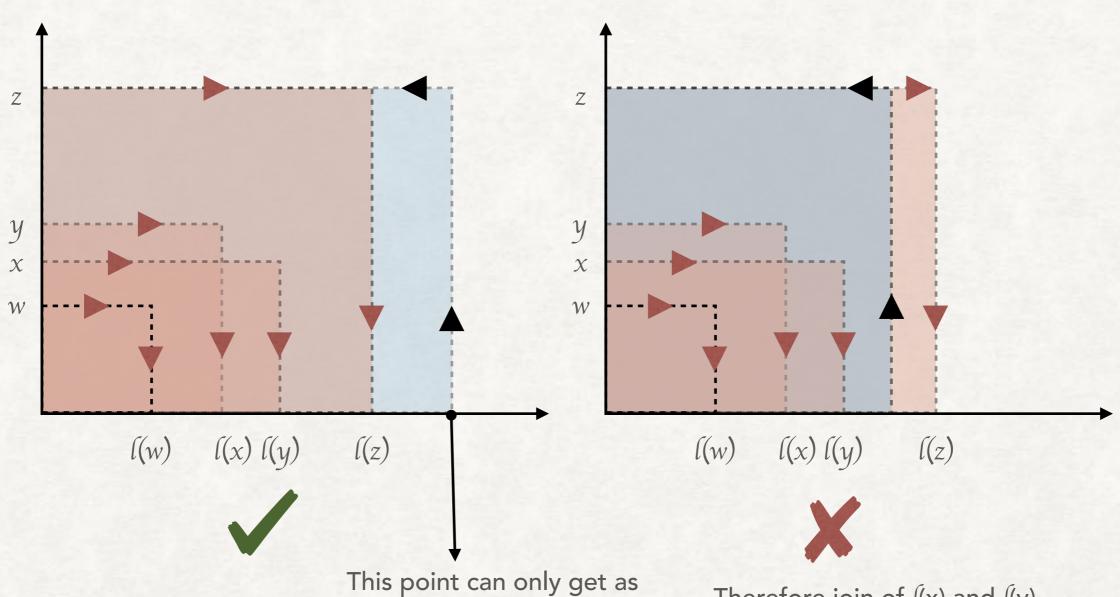


Is this a monotone map?

Simply means that: If join of x and y is z then join of $\ell(x)$ and $\ell(y)$ is $\ell(z)$



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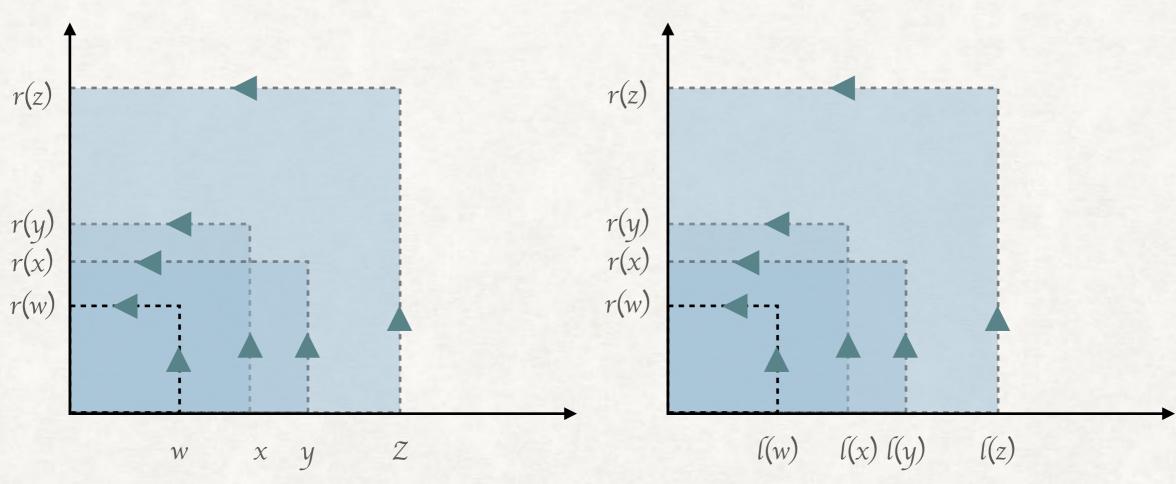


This point can only get as small as I(z) while maintaining Galois connection

Therefore join of $\ell(x)$ and $\ell(y)$ is $\ell(z)$

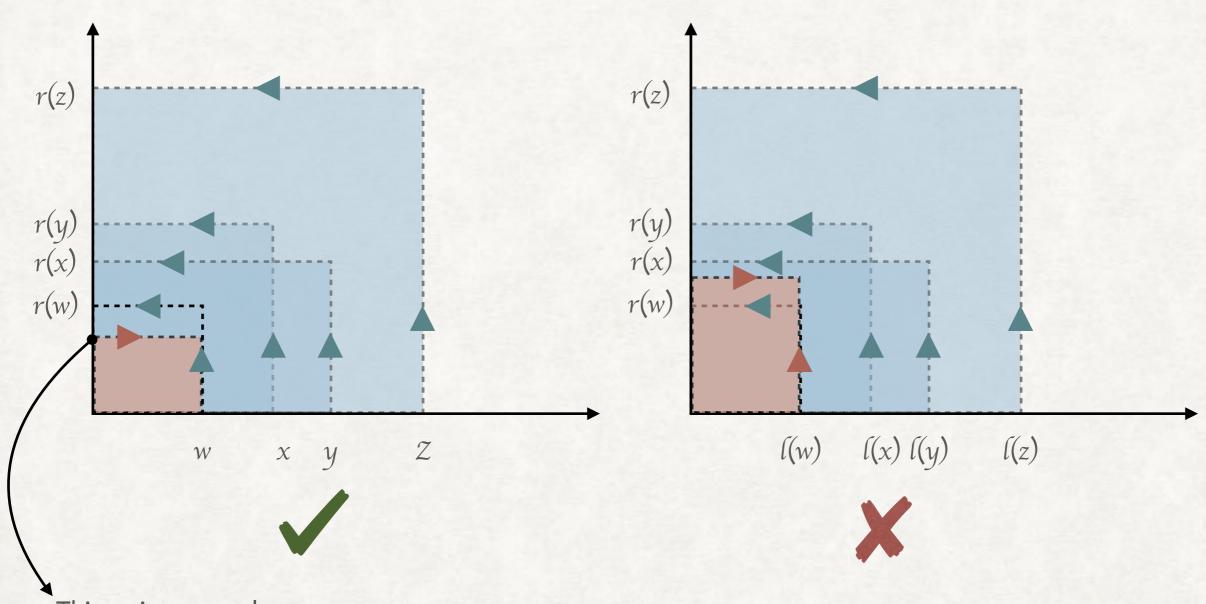
RIGHT ADJOINT PRESERVES MEETS

Simply means that: If meet of x and y is z then meet of r(x) and r(y) is r(z)



RIGHT ADJOINT PRESERVES MEETS

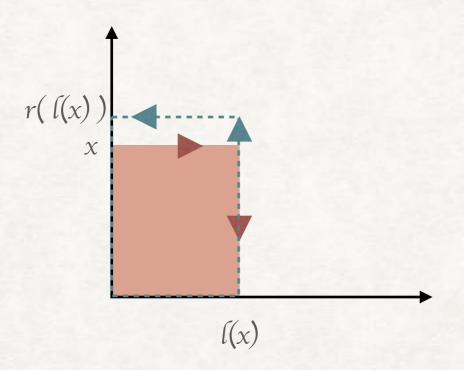
Simply means that: If meet of x and y is z then meet of r(x) and r(y) is r(z)



This point can only get as large as r(w) while maintaining Galois connection

Therefore join of r(x) and r(y) is r(z)

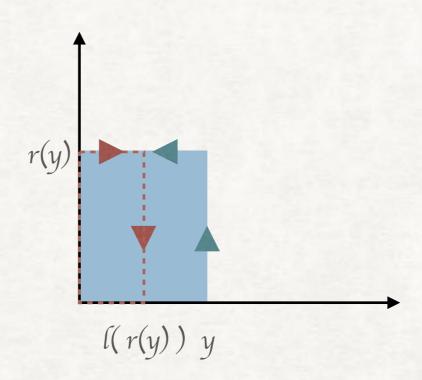
CLOSURE



$$x \le r(l(x))$$

$$x \le l \text{ then } r(x)$$

$$x \le l; r(x)$$



$$l(r(y)) \le y$$

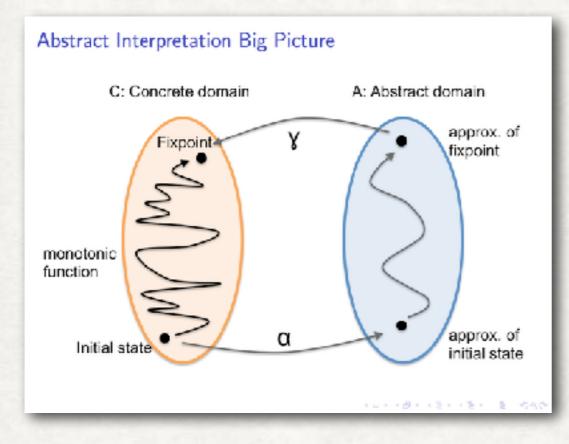
$$r \text{ then } l(y) \le y$$

$$r; l(y) \le y$$

$$l;r;l;r(x)=l;r(x)$$

APPLICATIONS

Abstract interpretation



Source: https://lara.epfl.ch/w/_media/sav17:lecturecise10.pdf

Syntax and semantics

Grammatical aspect of language

Semantic meaning of language

A statement can be syntactically correct but semantically meaningless. Eg.

Cow eats supremely

Semantics is left adjoint

Syntax is right adjoint

APPLICATIONS

Probability distribution

Cumulative distribution function

Left adjoint

 $F_X(x)$: $\mathbb{P}(X \leq x)$

Quantile function

Right adjoint

 $Q_X(p)$: inf{ $x \in \mathbb{R}$: p < F(x)}

Programming

Programming from Galois Connections

Shin-Cheng Mu^a, José Nuno Oliveira^b

^aInstitute of Information Science, Academia Sinica, Taiwan
^bHigh Assurance Software Lab / INESC TEC and Univ. Minho, Portugal

lead to specifications made of two parts: one defining a broad class of solutions (the *easy* part) and the other requesting one particular such solution, optimal in some sense (the *hard* part).

... analogous to ...

Pragmatic means the speaker wants to be correct and as specific as possible

Quantity Quality

THANK YOU