

Translating Chalice into SIL

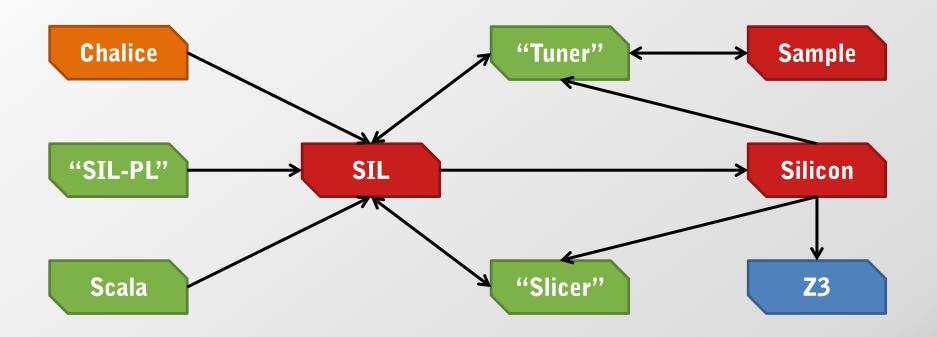
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The Semper Project

- Long term project
- Automatic program verifier for Scala
 - verify concurrent programs
 - reduce annotation overhead
 - deal with functional features (e.g., closures)



Semper Architecture Design



Semper Intermediate Language (SIL)

- Not a programming language
- A program representation for verification
- Not all constructs are executable
- High-level
- Aimed at 00

```
method C::m(this : ref) : (y:int)
 requires this ≠ null
 requires acc(this.C::f,write)
 ensures acc(this.C::f,write)
 ensures y == this.f
implementation C::m {
 entry:{
   y := this.f;
```

- Method Signatures
- Method Bodies
- Domains (data types)
 - predicates
 - functions
- Fields
- Functions
- Predicates

```
method C::cmpexc(
    this : ref,v : int,c : int)
    : (o : int)
    requires this ≠ null
    requires acc(this.C::f,write)
    ensures acc(this.C::f,write)
    ensures
    old(this.C::f == c) ⇒
        this.C::f == v
    ensures o = old(this.C::f)
```

- Method Signatures
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```
implementation C::cmpexc
 entry:{
   o := this.C::f;
 } if(this.C::f = c) goto exc
   if(this.C::f ≠ c) goto end
 exc:{
   this.C::f := v;
 } goto end
                           entry
 end:{
                       exc
                           end
```

- Method Signatures
- Method Bodies
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```
domain Pair[A,B]{
 function create(A,B)
                   : Pair[A,B];
 function getFirst(Pair[A,B])
                    : A;
 axiom getFirst = ∀ a:A,b:B ::
   getFirst(create(a,b)) = a;
domain Permission{
 function
   +(Permission, Permission)
                   : Permission
 predicate
   <(Permission, Permission);
```

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```
field C::f : int;
field L::value : int;
field L::next : ref;
function C::fGreater(a : int)
                         : bool
 requires acc(this.C::f,write)
   = C::f>a
predicate L::valid =
 acc(this.L::value,write) &&
 acc(this.L::next,write) &&
 this.L::next≠null ⇒
   acc((this.L::next).L::inv,
                          write)
```

Permissions

- Tracking
 (Thread×Field×Object) → Permission
- Permission can be
 - None ⇒ cannot access at all "0"
 - Some ⇒ can only read "]0,1["
 - Full ⇒ can read and write "1"



Permission Transfer

Transfer from caller to callee Modular verification

Fractional Read Permissions in Chalice

Explain what fractional read permissions do and that there is no direct SIL-equivalent

Implementing Fractional Permissions

Explain the basic idea behind the implementation.

- → select new fraction for every call-site
- -> collect constraints from current amount of permissions held

Implementing Fractional Permissions #2

Show naïve translation without using map $Explain\ how\ rd(f)\ \&\&\ rd(f)\ makes\ translation\ more\ complicated$

Implementing Fractional Permissions #3

Explain translation using Map[(ref,int), Permission]

Implement Fork-Join

Another Chalice feature not present in SIL Show how token object is assembled

- → Trick: how shadow field permissions are always linked to token.joinable
- → Current limitation: join mostly useless when not in same method as fork

QUICK DEMO

Conclusion



Thank you

QUESTIONS?

BACKUP/SCRAP

Annotated Methods

```
class Cell {
   var v: int;

method inc(d: int)
   requires 0 < d;
   requires acc(v);
   ensures v == old(v) + d;
   { v := v + d; }
}</pre>
```

- Annotated Methods
- Monitors

```
class Cell {
 var v: int;
 invariant acc(v) && 0 <= c;</pre>
class Program {
 method main() {
   var c:Cell := new Cell;
   c.v := 3;
   share c;
   acquire c; call c.inc(2); release c;
```

- Annotated Methods
- Monitors
- Predicates/Functions

```
class Cell {
 var v: int;
 predicate valid
 { acc(this.v) && 0 <= this.v }
 function add(d:int) requires valid;
 { unfolding valid in this.v + d; }
```

- Annotated Methods
- Monitors
- Predicates/Functions
- Fork-Join

```
class Cell { ... }
class Program {
 method main() {
   var c1:Cell := new Cell;
   var c2:Cell := new Cell;
   c1.v := 0; c2.v := 5;
   fork tk1 = c1.inc(3);
   fork tk2 = c2.inc(1);
   join f1 := tk1;
   join f2 := tk2;
```

Chalice2SIL

- First front-end for SIL
- Help establish and test the tool chain
- Ideally no changes to Chalice
- If enough time is left
 - Predicates and functions
 - Deadlock avoidance
 - Channels (Actor model)