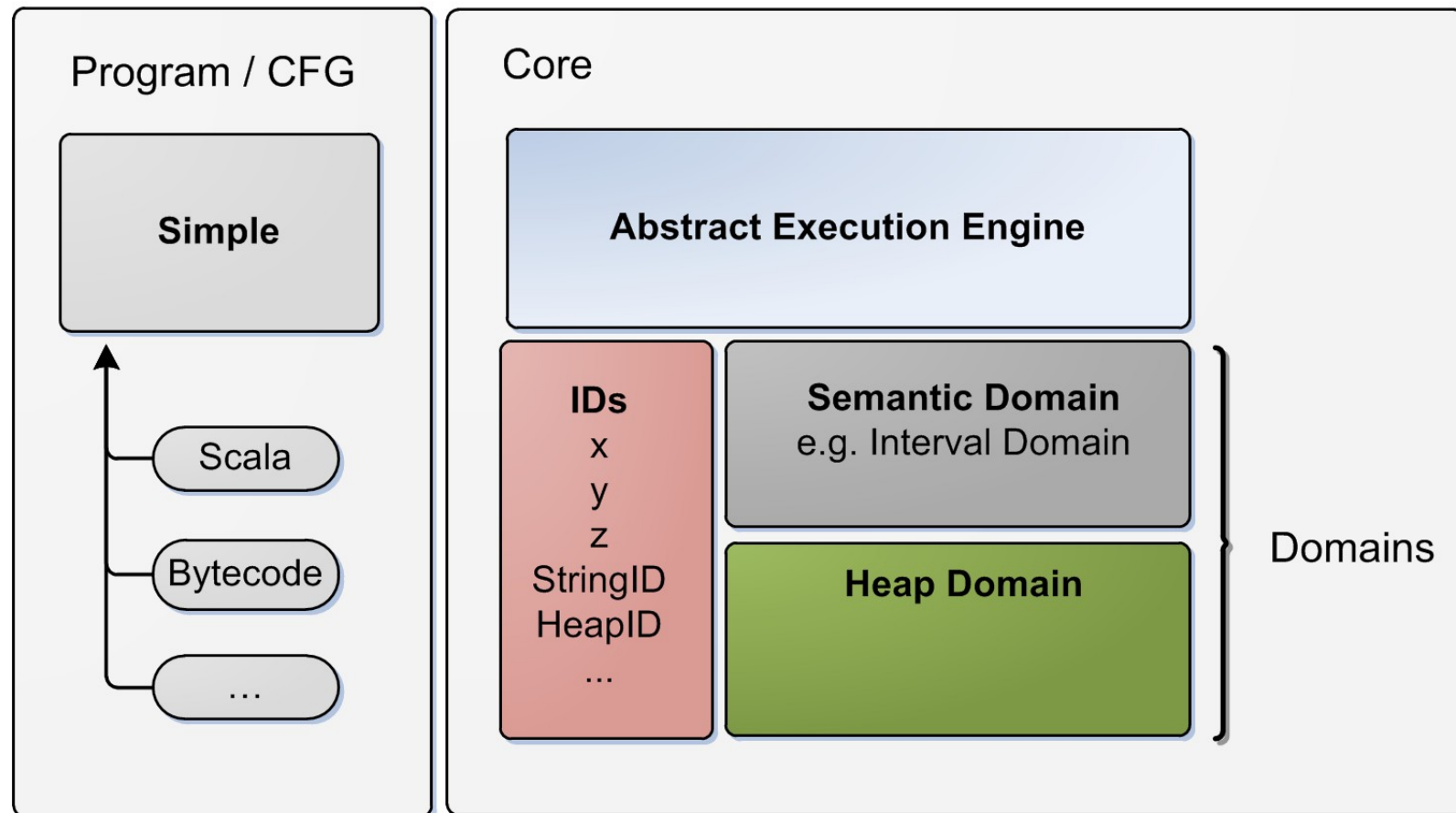


Interfacing TVLA & Sample

Raphael Fuchs

August 4, 2011

Sample: Architecture



Approach (1)

- Simple statements result in calls on heap domain
- For every heap operation, invoke TVLA
- E.g.

$x = y$

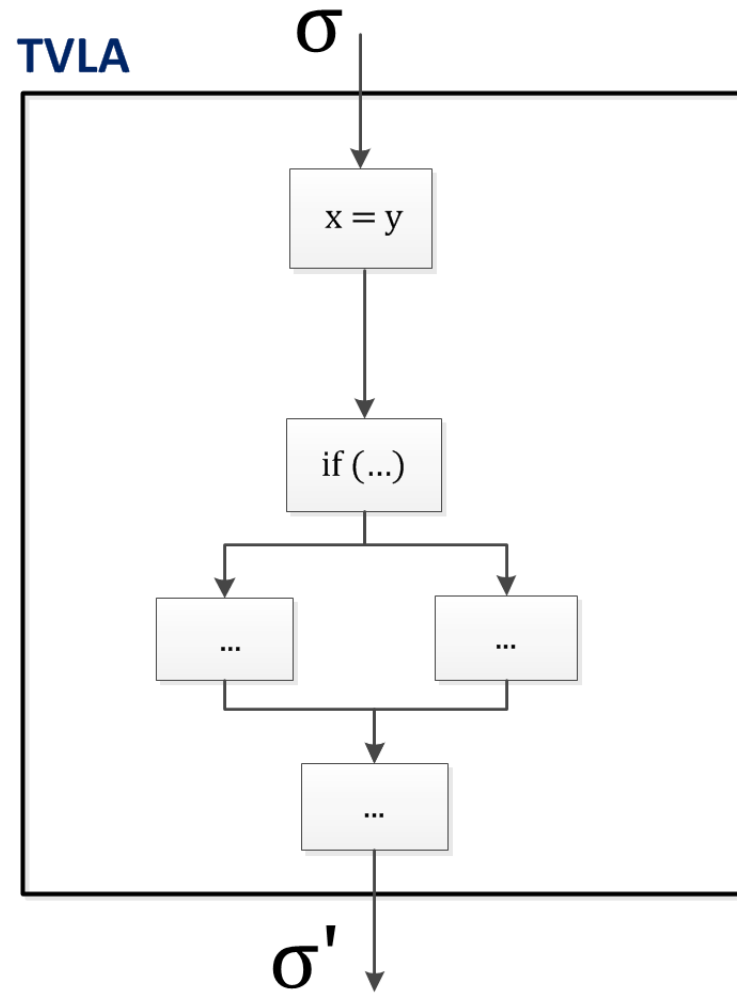


```
assignVariable(x,y)
```

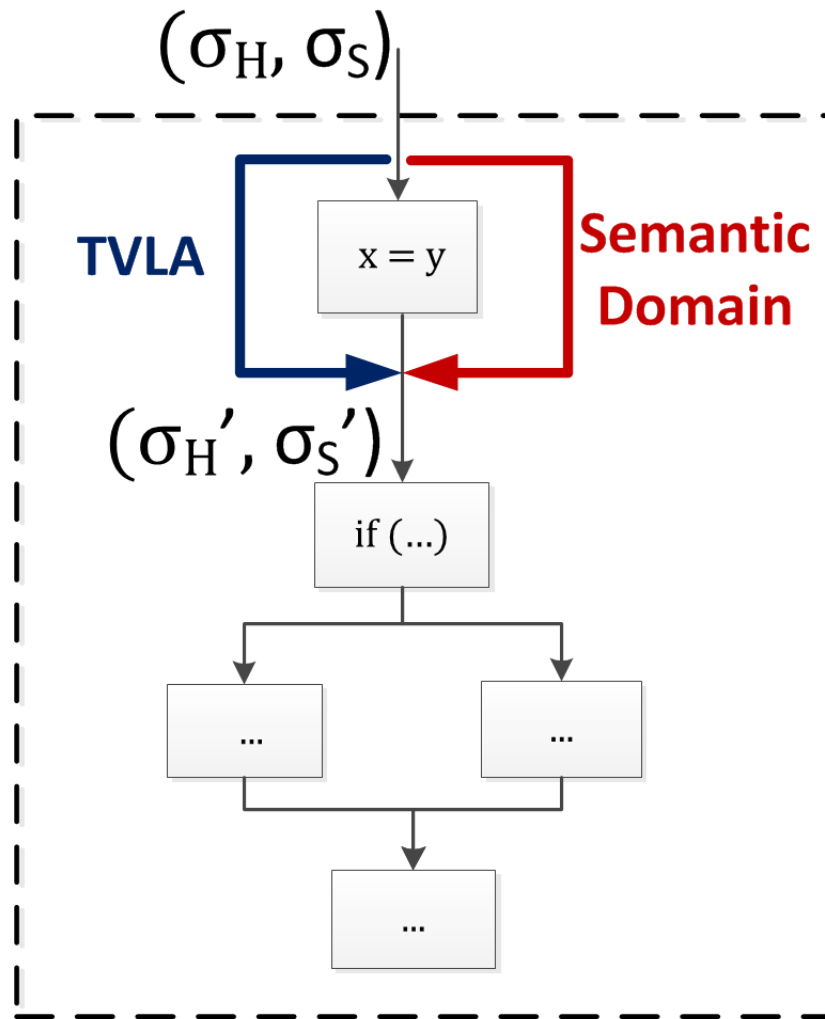
```
%action
```

```
assignVariable(target,source) {  
    %f { source(v) }  
    { target(v) = source(v) }  
}
```

Approach (2)



Approach (2)



Differences in Abstraction

Simple

- High-level language
- Object-oriented
- Lacks distinction:
expressions vs. Statements
- Types

TVP (TVLA)

- Quite low-level
- Usually C-like programs
encoded
- Semantics encoded using
FOL with TC
- Untyped

Translation not straightforward...

Heap Operations

TVP Actions for:

- Object Creation `new Foo()`
- Variable assignment (reference copy) `x = y`
- Field assignment `x.n = y`
- Field access `x.n`
- Assumptions `if (x != null) ... else ...`
- Least upper bound on heap states

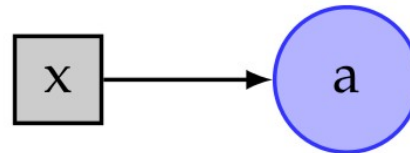
Heap State

State = Set of 3-valued logical structures

Each structure defined by:

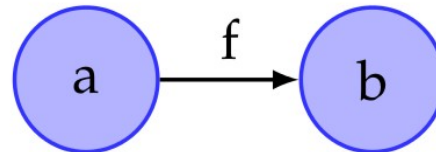
- Individuals: Blocks of allocated memory
- Unique unary predicates for program variables, e.g.

$x(a) = \text{true}$



- Binary function predicates for fields, e.g.

$f(a, b) = \text{true}$



Instrumentation Predicates

- Our analysis supports lists precisely
- Standard set of predicates
 - For every field f
 - $t[f](v1, v2)$
Transitive reflexive reachability: Can $v2$ be reached from $v1$ by following f -fields?
 - $is[f](v)$
Shared-ness: Is there more than one object whose f -field references v ?
 - For every field f and variable x
 - $r[f, x](v)$
Reachability: Can v be reached from x along f -fields?

Numerical Information

- Only heap (shape) analyzed, so far
- What about numerical information?
 - E.g. Integer fields
 - TVLA does not support this
- Idea:
 - Let every heap node be an identifier and associate some numerical info with it
 - More generally, any information from a semantic domain, for instance strings and access permissions

Heap Identifiers (1)

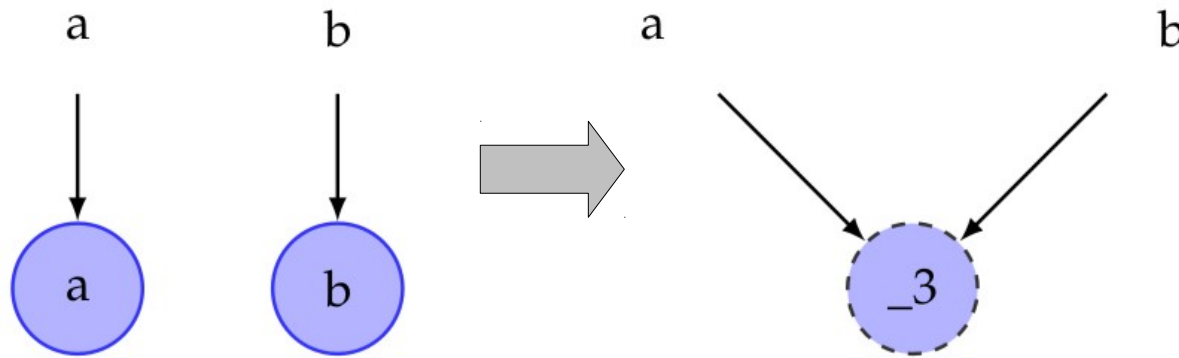
- Heap identifiers (names of nodes) are *not* static.
We have to consider
 - Summarization: Due to abstraction, nodes are merged into a summary node.
 - Materialization: Nodes can be extracted again from a summary node.
 - Merging due to heap isomorphism
- Update ID space according to these events

Heap Identifiers (2)

Problem: TVLA drops all names of individuals

- We do not obtain any information what happened
- We cannot update ID space and semantic domain

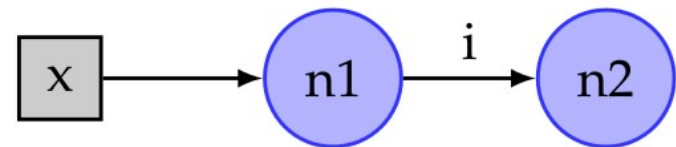
Trick: Use unary non-abstraction predicates to track identifiers (names of nodes)



Example: Intervals

```
1 def iAssign(unknown:Boolean) = {  
2   val x = new IntNode  
3  
4   if (unknown) {  
5     x.i = 1  
6   } else {  
7     x.i = 2  
8   }  
9 }
```

Interval domain end-state:



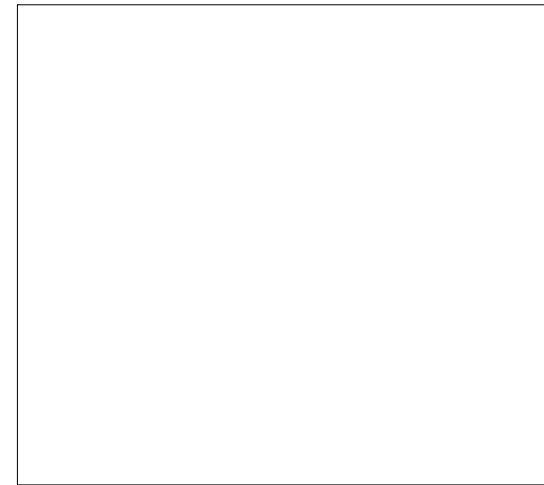
Interval domain end-state:

n1	->	\top
n2	->	[1..2]

Example: Initializing a List

```
1 def createZeroList(n: Int) = {  
2   var h: IntNode = null  
3   var cur :IntNode = null  
4   var i = 0  
5  
6   while (i < n) {  
7     cur = new IntNode  
8     cur.i = 0  
9     cur.n = h  
10    h = cur  
11    i += 1  
12  }  
13 }
```

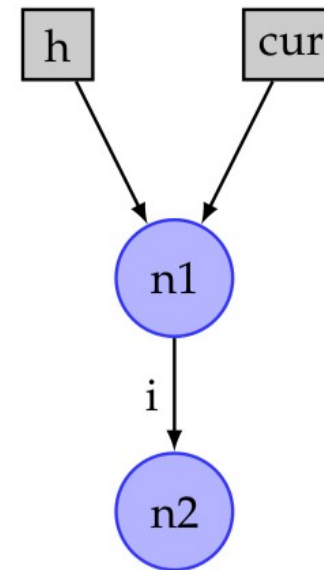
Heap Entry-State



Example: Initializing a List

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```

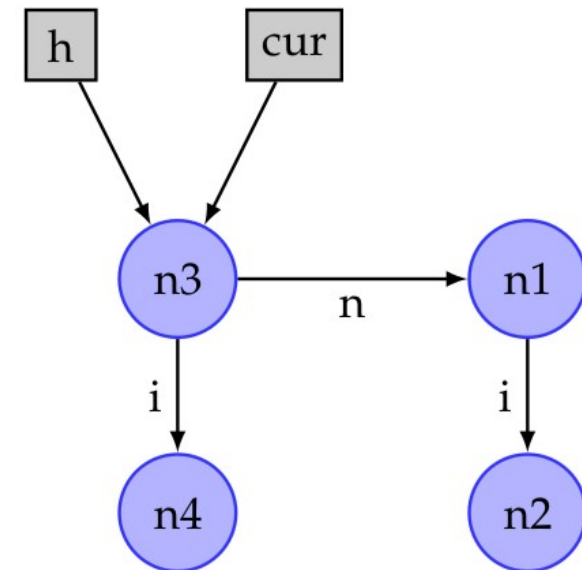
After first iteration:



Example: Initializing a List

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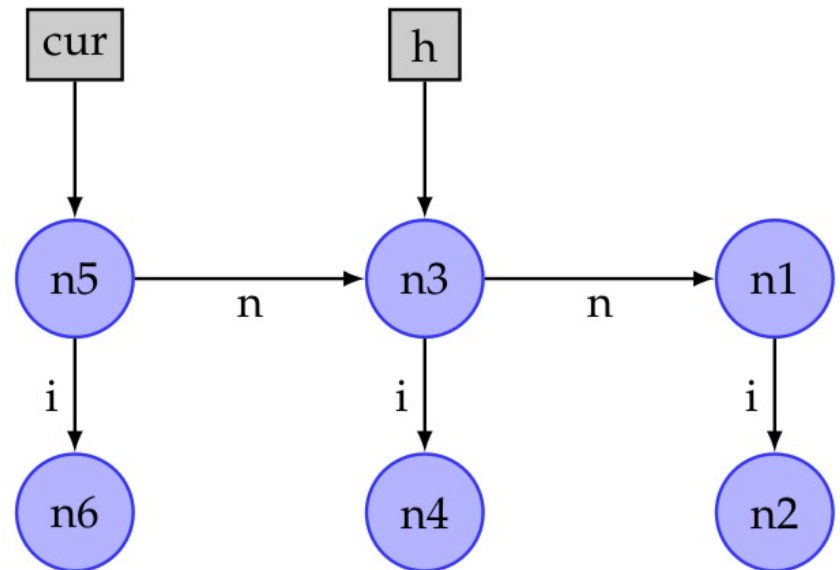
After second iteration:



Example: Initializing a List

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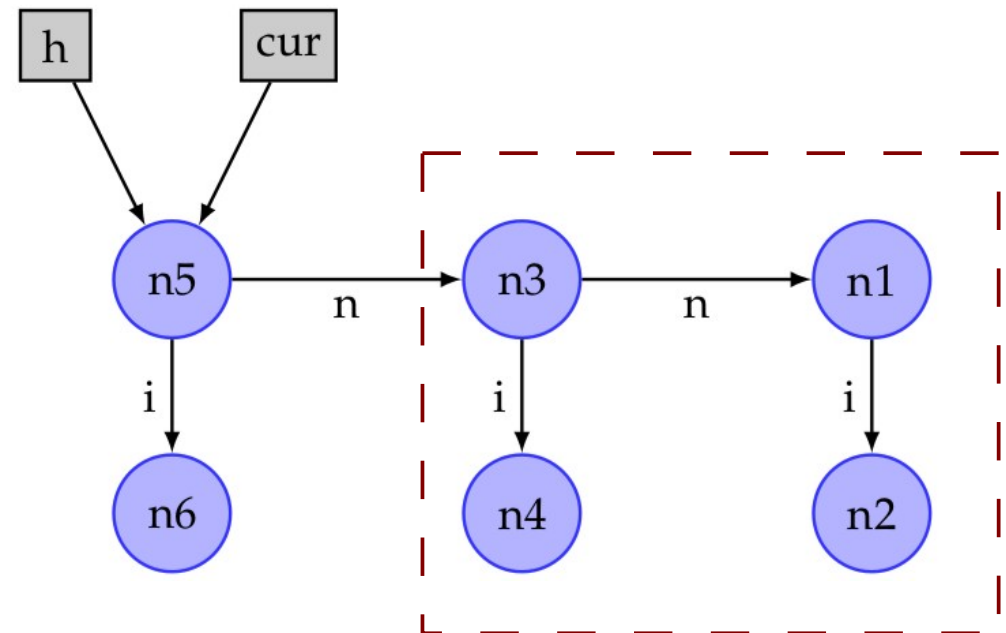
Before $h = cur$
(third iteration)



Example: Initializing a List

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11    i += 1  
12  }  
13 }
```

After $h = cur$
(third iteration)

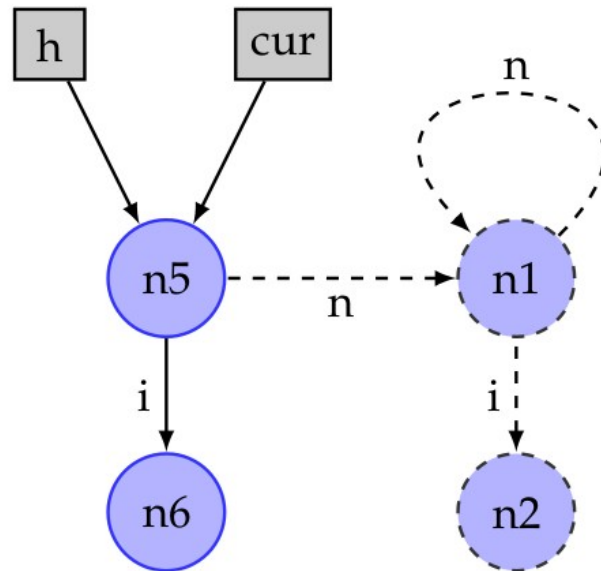


n1 and n3
n2 and n4

look the same to TVLA!

Example: Initializing a List

Summarization of n1, n3 and n2, n4

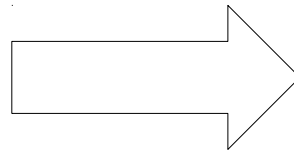


$\{n1, n3\} \rightarrow \{n1\}$

$\{n2, n4\} \rightarrow \{n2\}$

Heap ID	Semantic Domain
n2	$[0..0]$
n4	$[0..0]$
n6	$[0..0]$

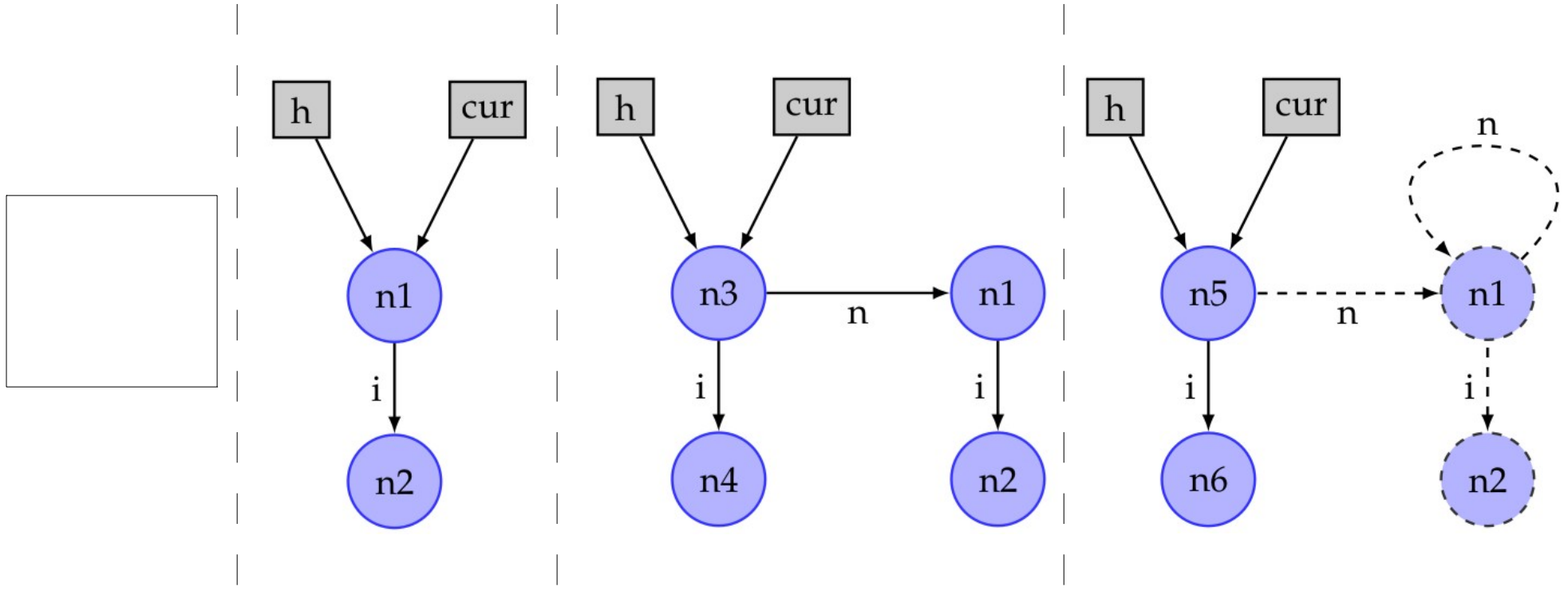
$n2 \sqcup n4$



Heap ID	Semantic Domain
n2	$[0..0]$
n6	$[0..0]$

Example: Initializing a List

Resulting Structures



Testing & Performance

- Developed test suite with ~ 40 tests to evaluate our domain.
- Total execution time: ~5 min (previously: 45)
- Calls to TVLA a severe performance bottleneck
 - JVM start-up time:
=> ~ 5x speed-up
 - Inefficiency of CFG iteration: Caching!
=> ~ 2x speed-up

Questions?