

Diameter of tree \rightarrow longest distance ^{#edges} b/w any two nodes in the tree.

\Rightarrow Find diameter of binary tree.

diameter = 0;

int height(root) {

if (root == null)

return -1;

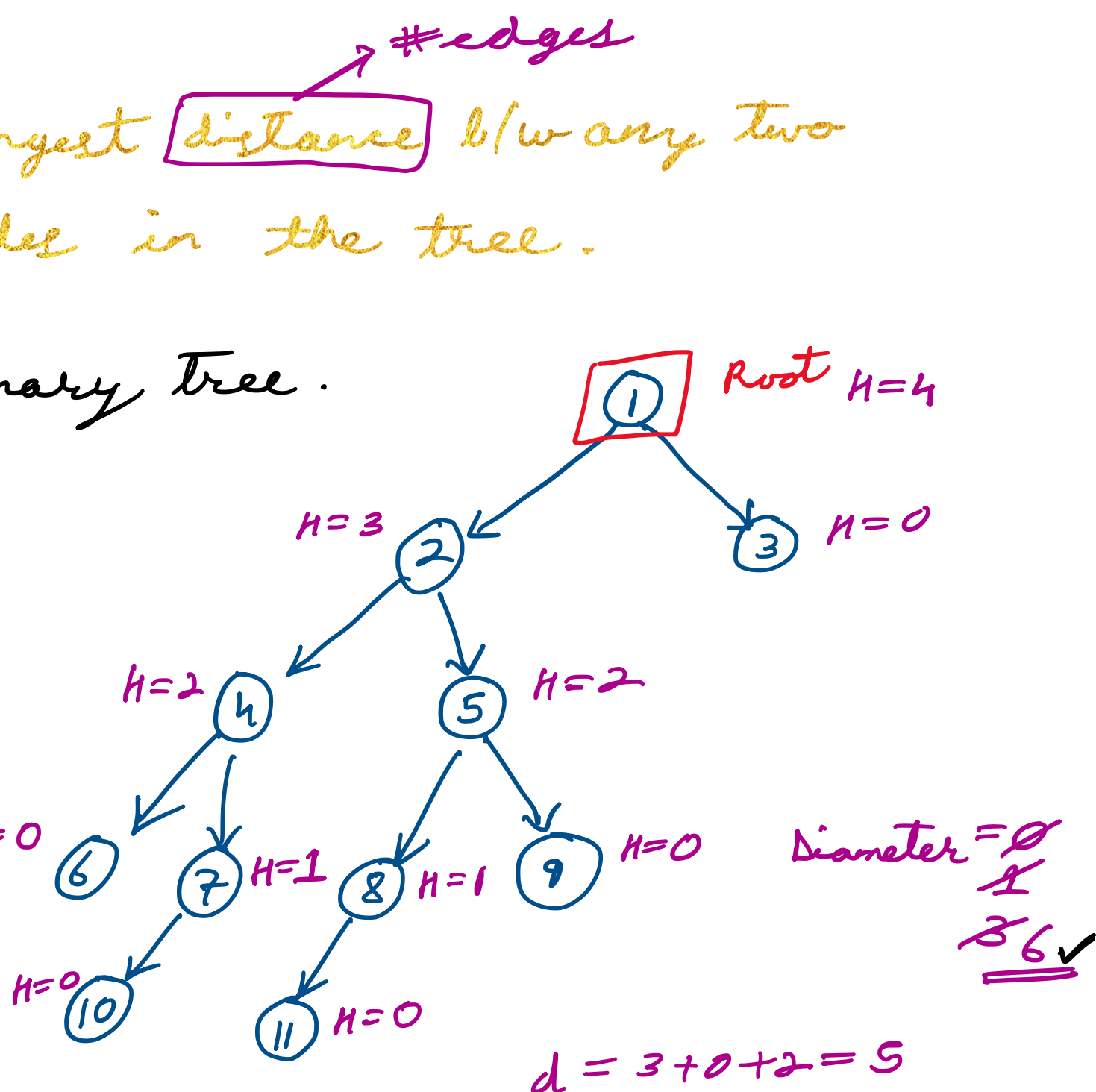
L = height(root.left);

R = height(root.right);

diameter = max(diameter, L + R + 2);

return max(L, R) + 1;

}



$$d = 3 + 0 + 2 = 5$$

$$TC = O(N)$$

$$SC = O(H)$$

Morris Inorder Traversal

SC = O(1) ✓

\rightarrow [3, 5, 8, 9, 10, 20, 50]

Left Node Right

✓ Goal \rightarrow Find a way to come back to current node once we go left.

✓ We want to come back to current node from inorder predecessor of current node.

```
node predecessor(cur) {
    temp = cur.left;
    while (temp.right != null && temp.right != cur)
        temp = temp.right;
    return temp;
}
```

cur = root;

while (cur != null) {

if (cur.left == null) {

print(cur.data);

cur = cur.right;

} else {

pre = predecessor(cur); // cur \rightarrow left \rightarrow right

if (pre.right == null) {

pre.right = cur; // DO

cur = cur.left;

} else {

pre.right = null; // UNDO

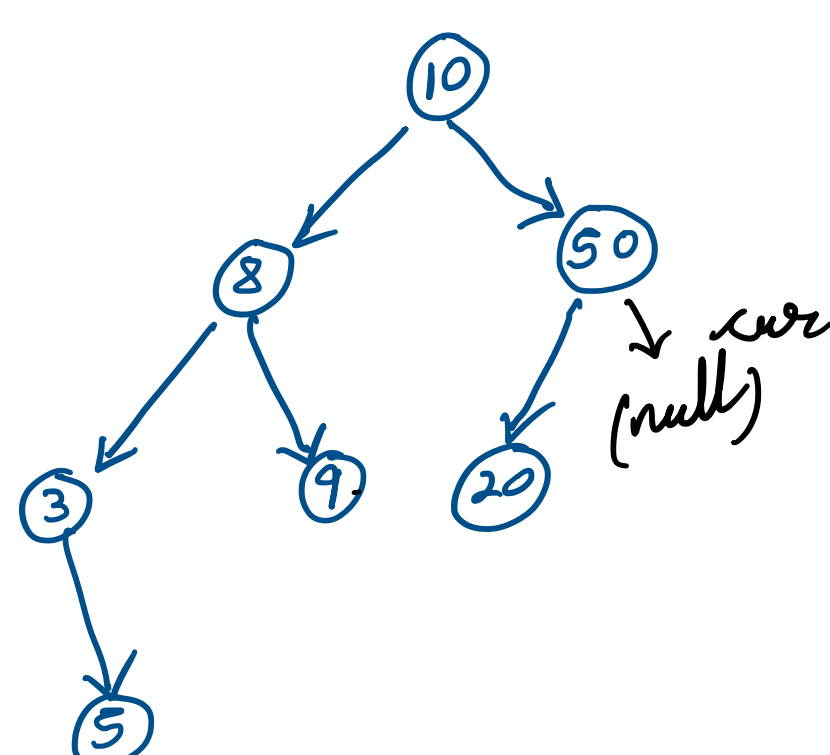
print(cur.data);

cur = cur.right;

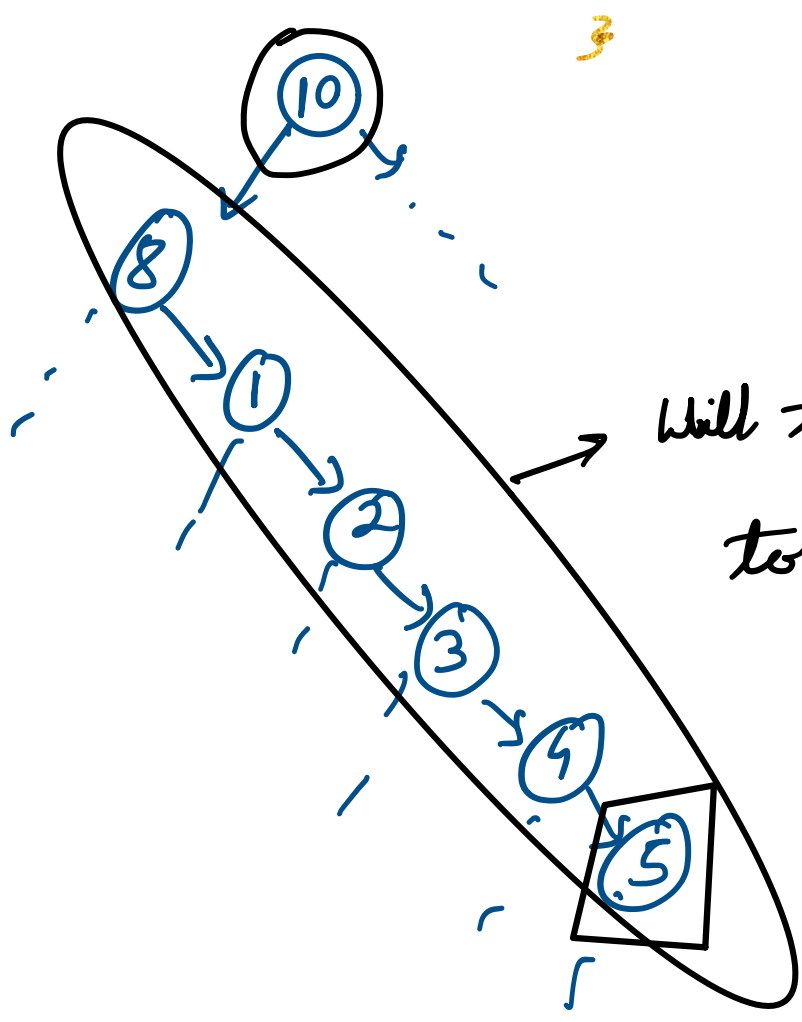
}

$$TC = O(H)$$

Inorder \rightarrow [3, 5, 8, 9, 10, 20, 50]



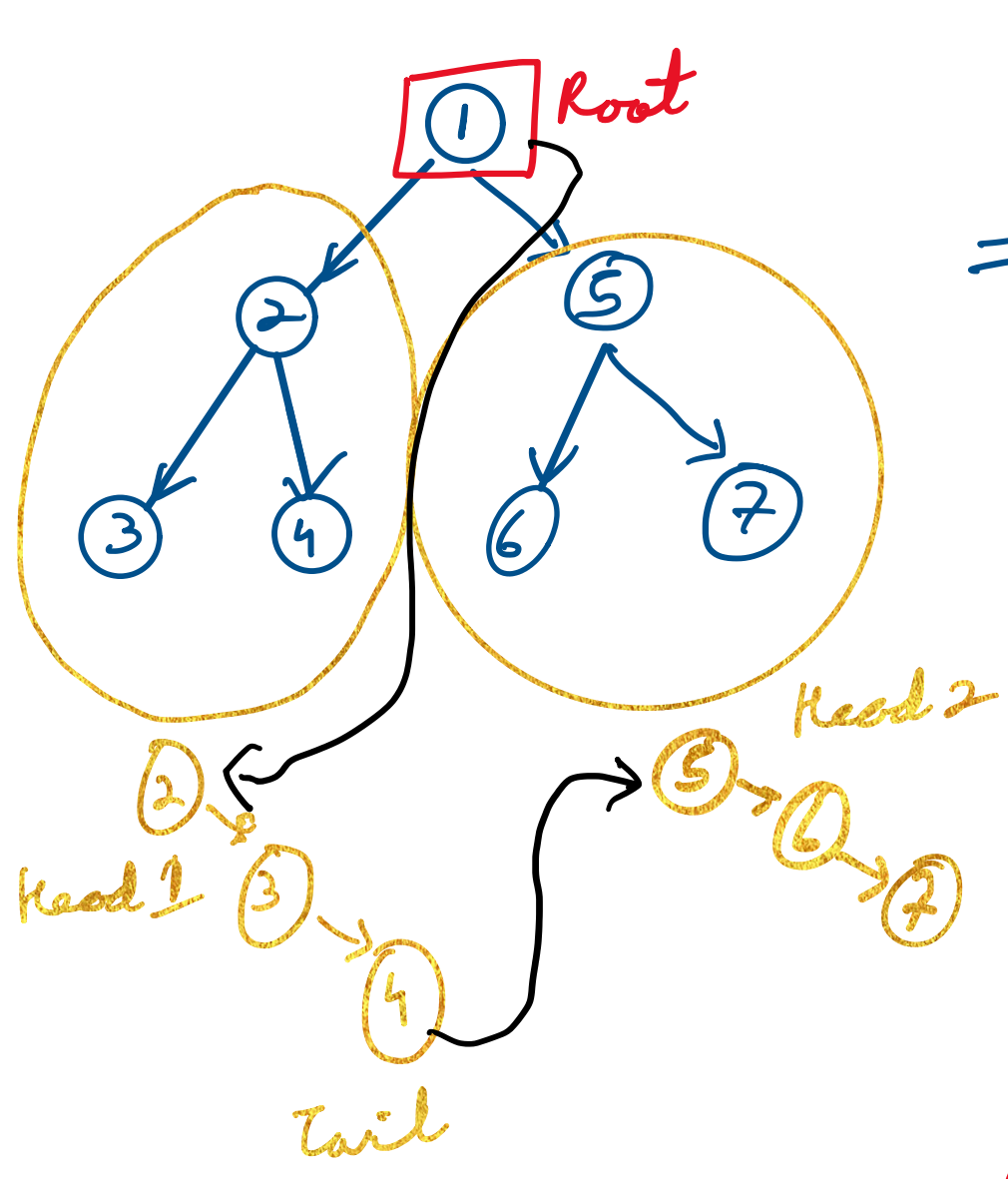
$$TC = O(N) \checkmark$$



Will these nodes be travelled over to find predecessor for any other node? No

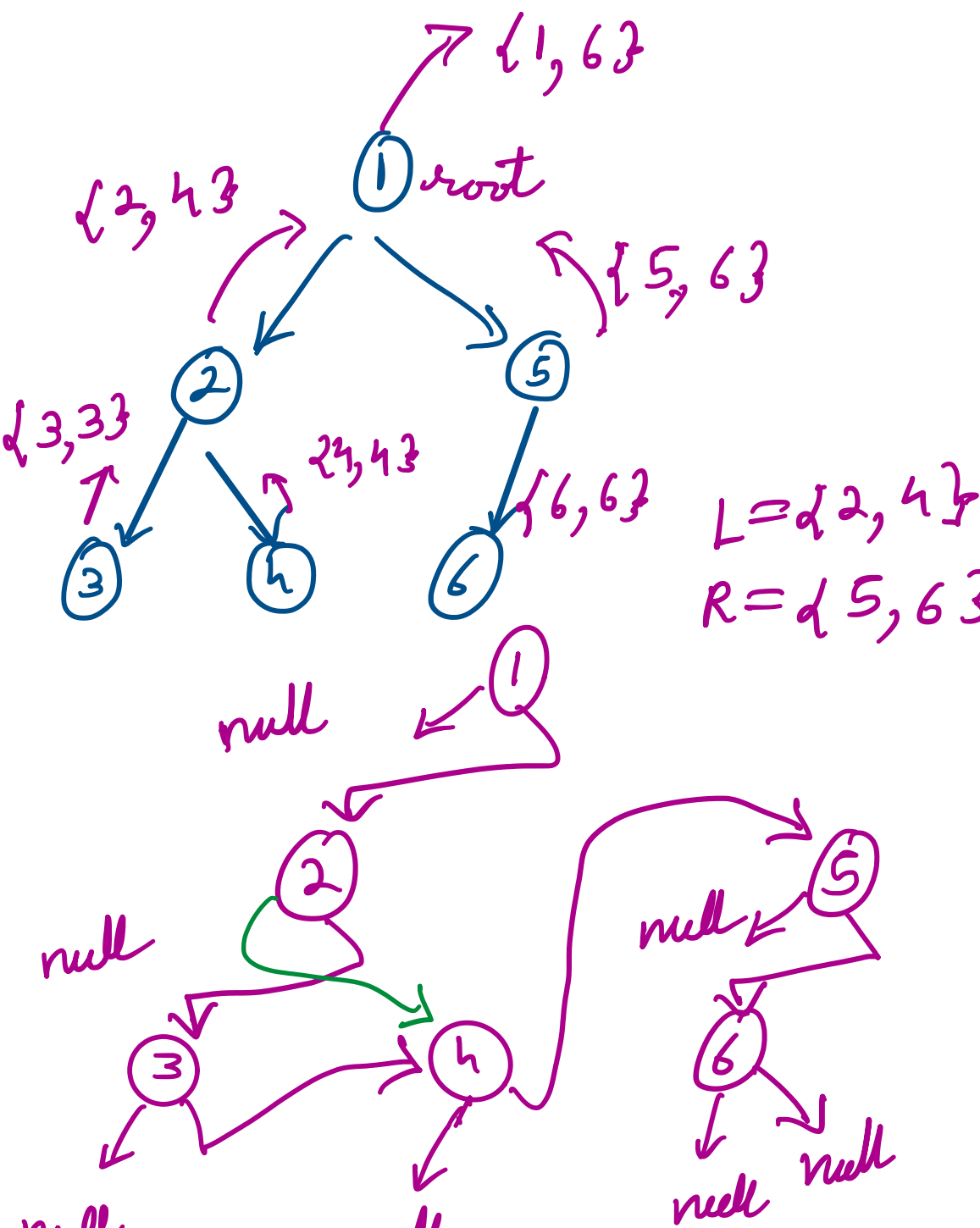
2 times for 10.

\Rightarrow Flatten Binary Tree to linked list.



$$TC = O(N)$$

first = head
second = tail



pair < node > flatten(root) {

if (root == null) return {null, null};

L = flatten(root.left);

R = flatten(root.right);

root.left = null;

if (L.head != null) {

root.right = L.head;

L.tail.right = R.head;

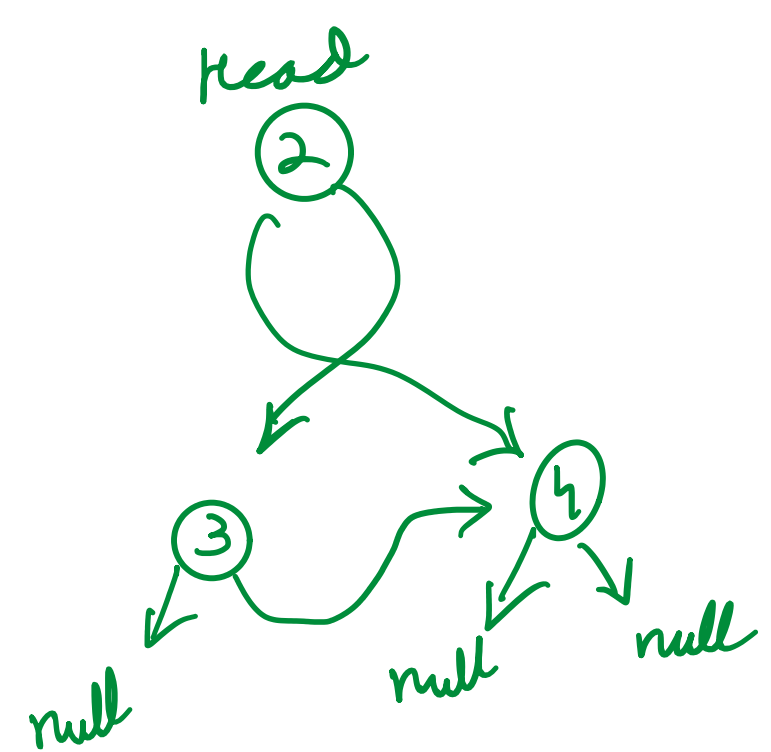
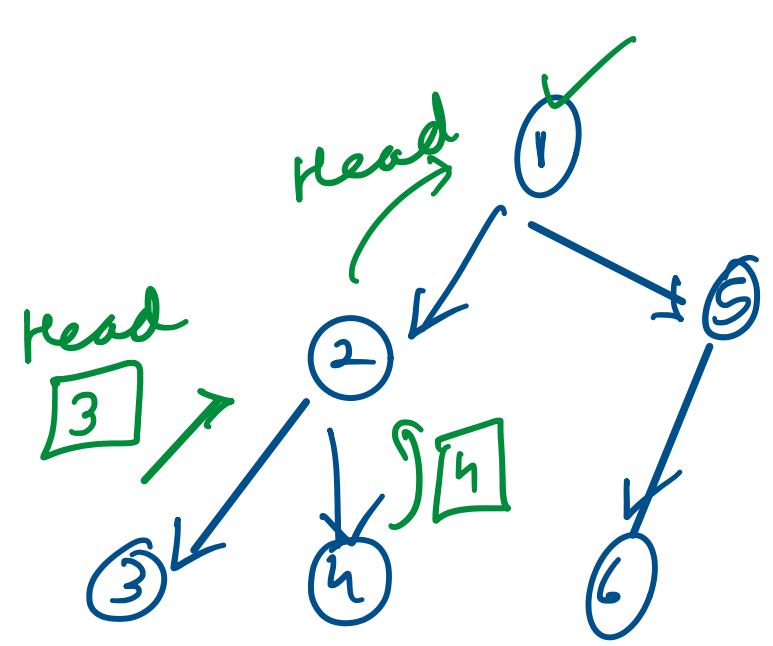
} else {

root.right = R.head;

}

return {root, Head}

(R.tail != null ? R.tail :
(L.tail != null ? L.tail :
root)) }



$$L = 3$$

$$R = 4$$