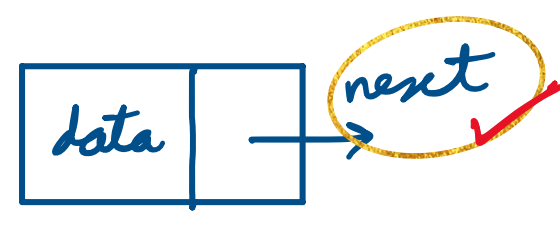


limitations on arrays  $\rightarrow$   $\left. \begin{array}{l} \Rightarrow \text{Fixed-size} \\ \Rightarrow \text{Insertion} \\ \Rightarrow \text{Deletion} \end{array} \right\}$  Dynamic Array  $\checkmark$   
linked list



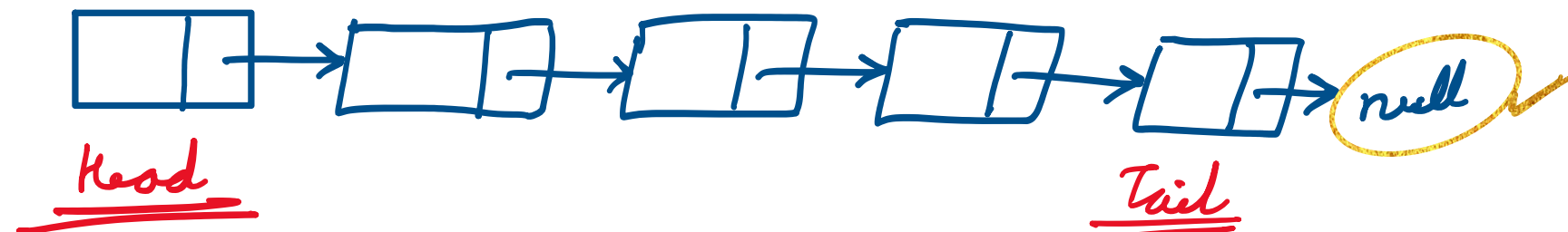
```

struct node {
    int data;
    node * next;
}
    
```

```

class node {
    int data;
    node next;
}
    
```

Non-continuous  
Memory Allocation



### Operations in LL

$\Rightarrow$  Access  $k^{\text{th}}$  element  $\rightarrow$  Array  $\rightarrow A[k-1]$   
 $TC = O(1)$

Never modify  
 head pointer  
 unless important.

```

node x = head;
for (i = 1; i < k; i++) {
    if (x == null) return null;
    x = x->next;
}
    
```

return x;  $TC = O(k) \approx O(N) \checkmark$

$\Rightarrow$  search a value v  $\rightarrow$  Array  $\rightarrow$  Linear search  $TC = O(N)$   
 Binary search (only sorted Array)  
 $TC = O(\log n)$

```

node x = head;
while (x != null) {
    if (x->data == v) return x;
    x = x->next;
}
return null;
    
```

$TC = O(N)$

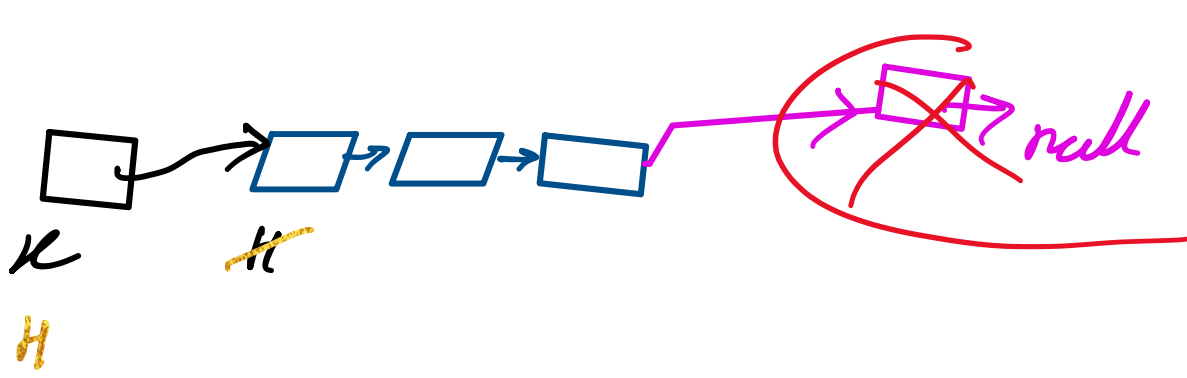
$\Rightarrow$  Insertion a value v  $\rightarrow$  Array  $\rightarrow$  Make a new array of size  $N+1$  & copy all  $N$  contents.  $TC = O(N) \checkmark$

### Insert head

```

node x = newNode(v);
x->next = head;
H = x;
return H;
    
```

$TC = O(1) \checkmark$



### Insert end

```

node x = head; // head != null
while (x->next != null) {
    x = x->next;
}
x->next = newNode(v);
    
```

$TC = O(N) \checkmark$

$\Rightarrow$  deletion of value v  $\rightarrow$  Array  $\rightarrow O(N) \checkmark$

```

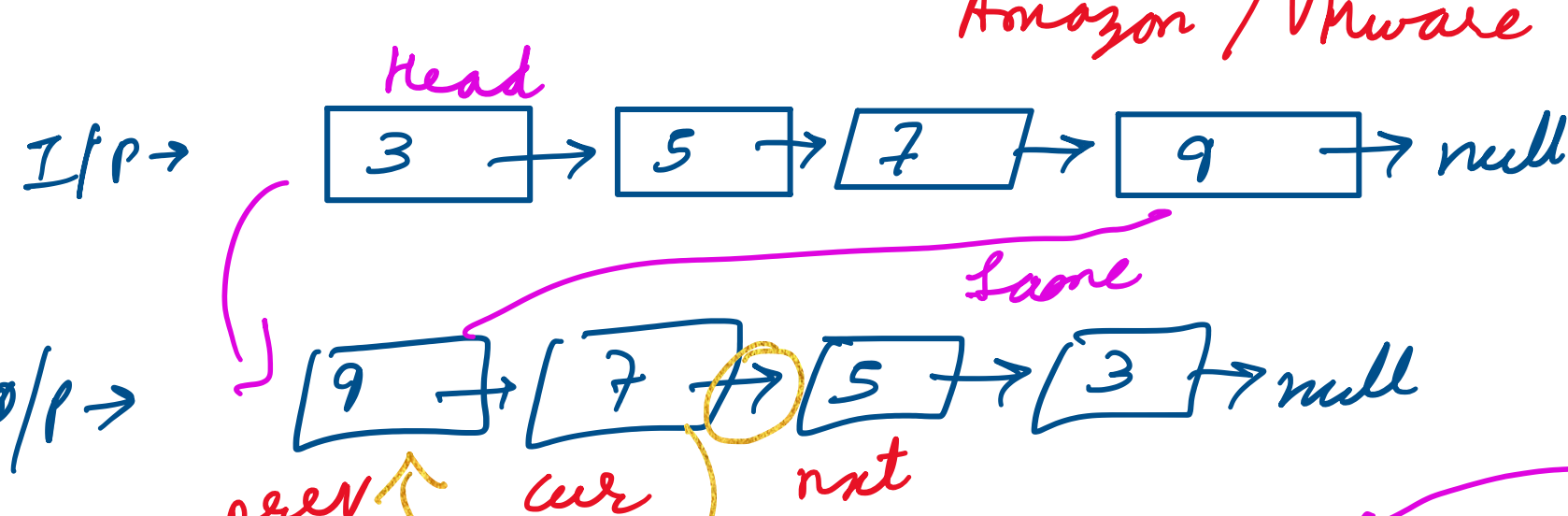
node delete (H, v) {
    if (H == null) return null;
    if (H->data == v) return H->next;
    node x = head;
    while (x->next != null) {
        if (x->next->data == v) {
            x->next = x->next->next;
            return head;
        }
        x = x->next;
    }
    return head;
}
    
```

$TC = O(N)$



### Reverse linked list

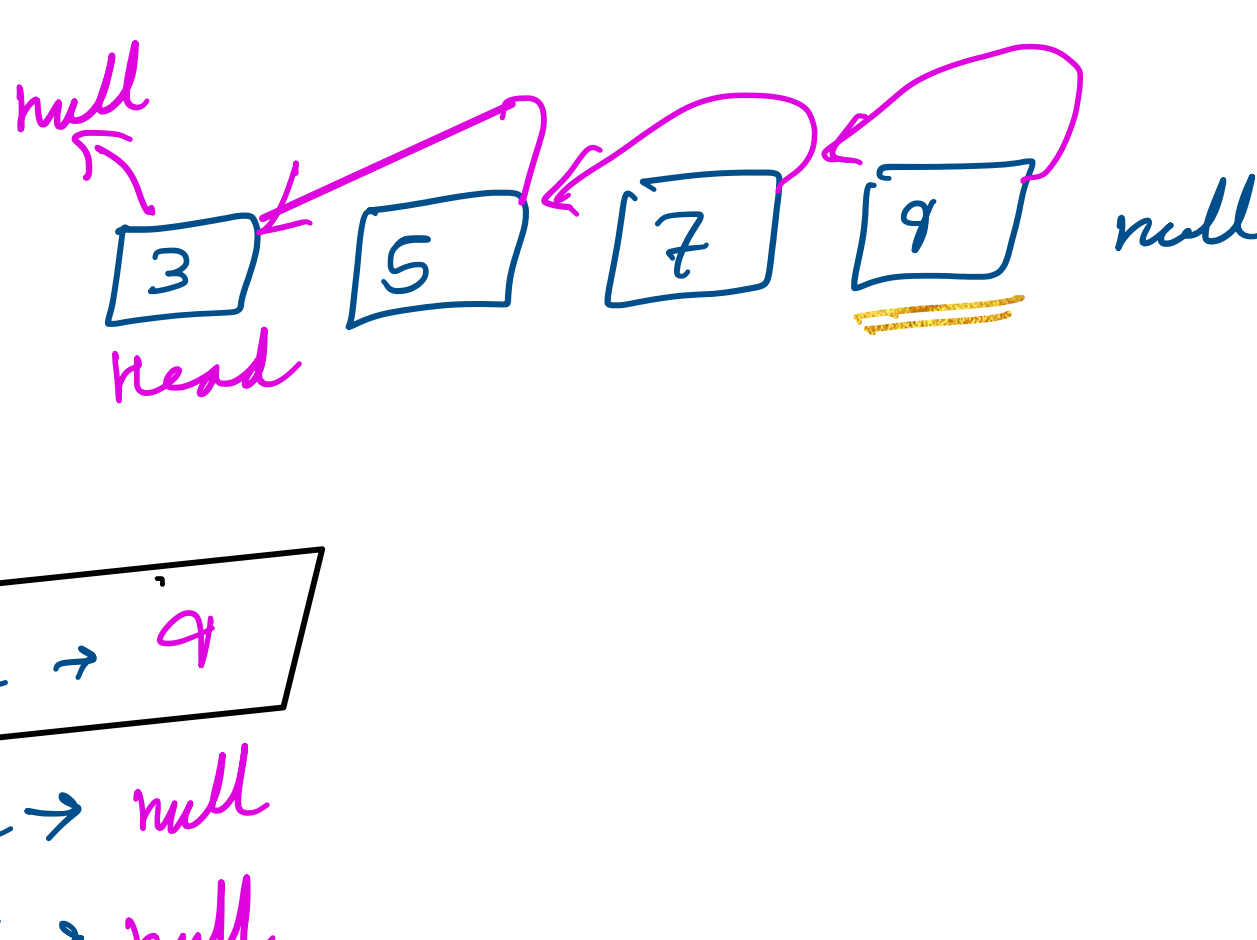
Qualcomm / Expedia / Wipro / KLL / SAP labs  
 Amazon / VMware



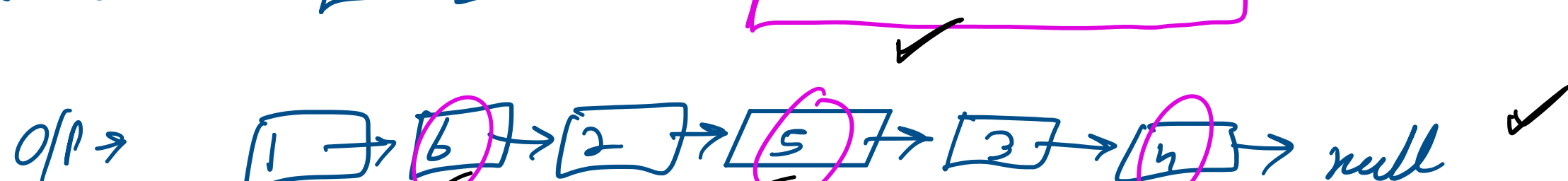
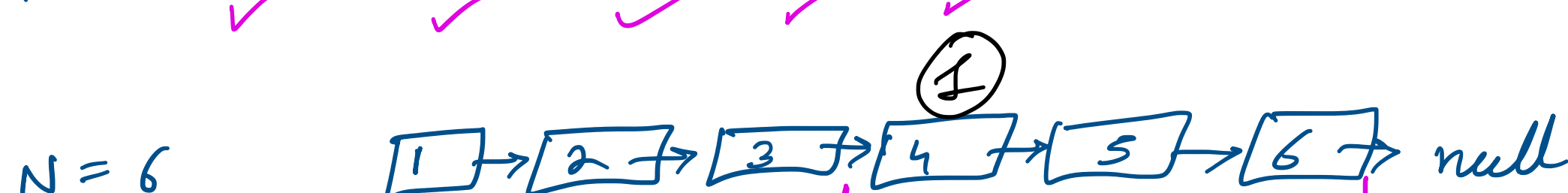
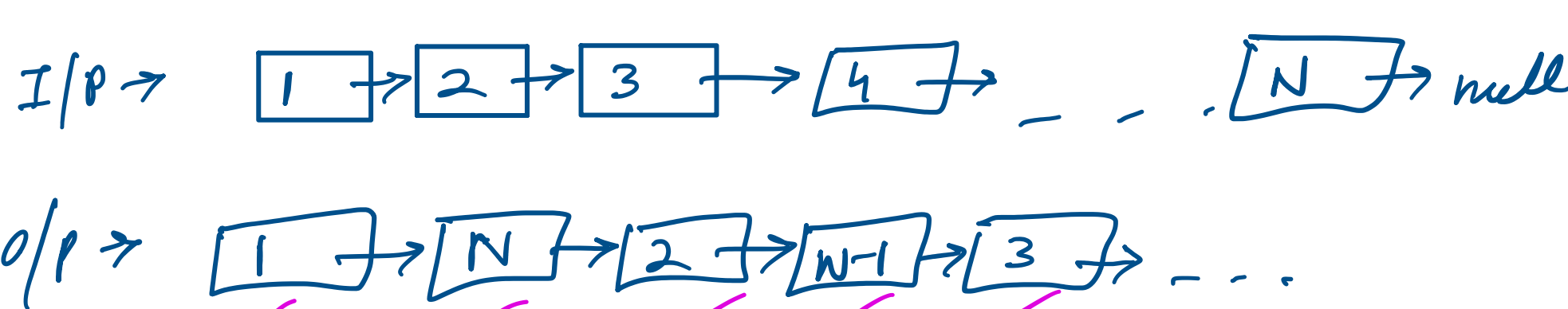
```

node pre = null, next = null;
node cur = head;
while (cur != null) {
    next = cur->next;
    cur->next = pre;
    pre = cur;
    cur = next;
}
return pre;
    
```

$TC = O(N)$



### Rearrange list

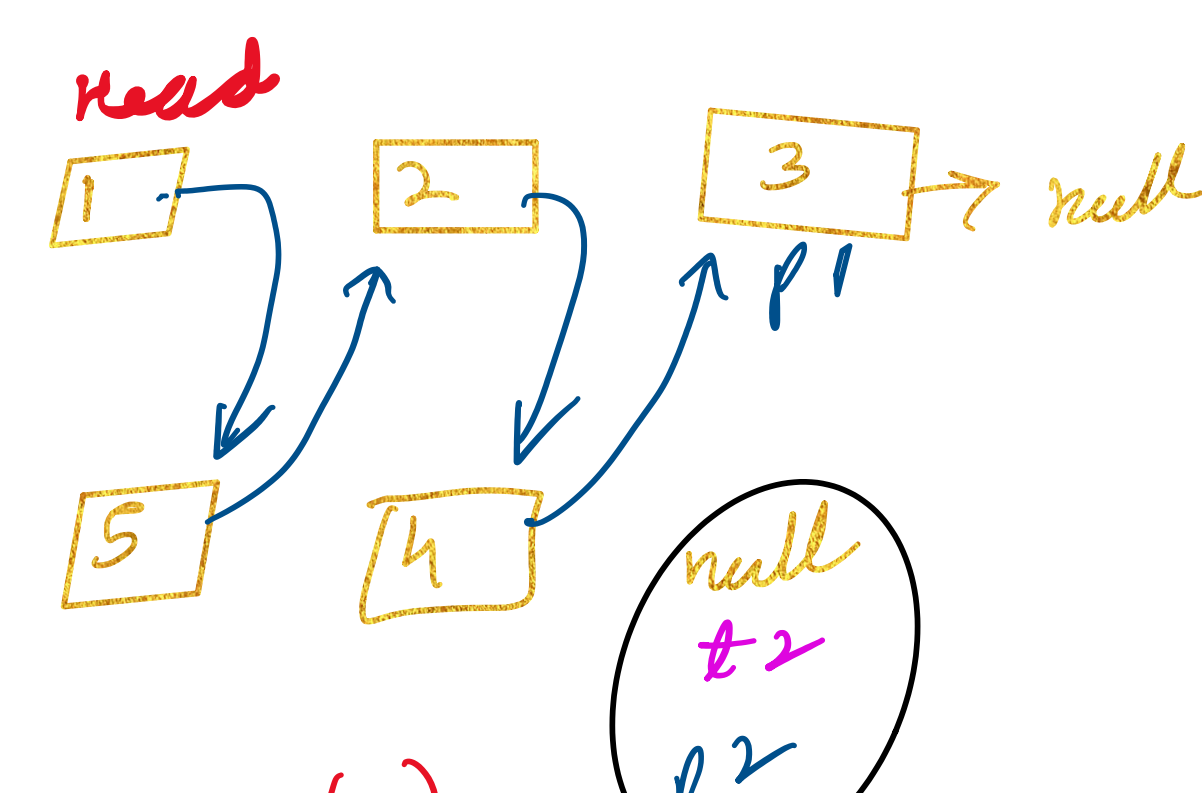
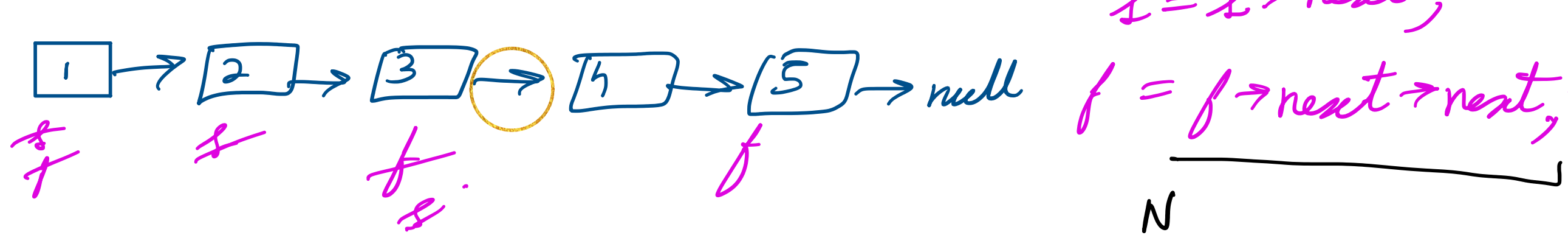


Middle element  $\rightarrow$  calculate N & middle =  $\frac{N}{2}$

$\Rightarrow$  Slow & Fast Pointers

$\downarrow$   $\downarrow$   
 $+1$   $+2$

$N/2$



$TC = O(N)$

```

x1 = p1->next;
x2 = p2->next;
p1->next = p2;
p2->next = x1;
p1 = x1;
p2 = x2;
    
```