COMPUTER NETWORK ASSIGNMENT 2 REPORT

OUESTION 1

Server Code

1) The server sets up a TCP socket and listens for client connections.

Socket Creation:

- The **socket() function** creates a TCP socket (with AF_INET for IPv4 and SOCK_STREAM for TCP).
- If the socket creation fails, an error message is printed using **perror()**, and the program exits.

Binding:

- The **bind()** function associates the server socket with a specific IP address and port, allowing it to listen on that address.
- If the binding fails, the program prints an error message and exits.

Listening:

- The **listen()** function puts the server in a passive listening mode, allowing it to accept incoming client connections. The second argument specifies the maximum number of pending connections in the queue.
- If the server fails to listen, it prints an error message and exits.
- 2) The server handles multiple concurrent clients (multithreaded, concurrent server). The server accepts the client connection; hence, a new socket is created with **4 tuples** (server IP, server listening port, client IP, client port). The server creates a new thread that continues to process the client connection

Accepting Client Connections:

- The accept() function waits for incoming client connections on the server's listening socket.
 When a client connects, a new socket (new_socket) is created specifically for that client connection.
- The 4-tuple created consists of:
 - Server IP
 - Server port
 - Client IP (retrieved from the client connection)
 - Client port (assigned by the client system)

• If the accept() call fails, an error message is printed, and the server continues to listen for the next client.

Thread Creation:

- For each new client, the server creates a new thread using **pthread_create()**. This allows the server to <u>handle multiple clients concurrently.</u>
- The new_socket is passed to the thread by allocating memory for the socket pointer and passing it as an argument to the **handle client() function.**

Thread Detachment:

- **pthread_detach()** is called to ensure the thread cleans up its resources when it finishes. This allows the server to handle multiple client threads without manual thread management.
- 3) The original server socket continues to listen on the same listening port for newer incoming client connections.
- The server is placed in an infinite loop, allowing it to call accept() and handle new client connections continuously.
- Each time accept() is called, the server waits for a new client connection on the same listening port (8080). Once a client connects, the server creates a new socket (new_socket) for that specific connection, but the original socket (server fd) continues listening for additional clients.

Non-Blocking Server:

- After each client connection is accepted and handed off to a new thread via **pthread_create()**, the server immediately returns to the top of the loop and <u>waits for the next client connection</u>.
- This allows the original server socket (server_fd) to remain open on the same port and handle new incoming connections while other clients are being processed in separate threads.
- 4) The server finds out the top CPU-consuming process (user+kernel CPU time) and gathers information such as process name, pid (process id), and CPU usage of the process in user & kernel mode
 - This functionality is handled by the **get_process_info()** function, which reads the /proc directory to find running processes and gathers information such as process name, PID, and CPU usage in user and kernel modes. It calculates the total CPU time (user time + kernel time) for each process and identifies the top two CPU-consuming processes.

CLIENT

- 1) The client creates a socket and initiates the TCP connection. The client process supports initiating "n" concurrent client connection requests, where "n" is passed as a program argument.
 - The client creates a socket using **socket()** to establish a **TCP connection**. If the creation fails, it prints an error message and exits the thread.
 - The client supports initiating "n" concurrent connection requests by creating multiple threads. The number of client threads is determined by the command-line argument passed when executing the program. Each thread runs the client_thread() function, which handles the connection and communication with the server.
- 2) After the client connection is established, the client sends a request to the server to get information about the server's top TWO CPU-consuming processes.
 - After the client establishes a TCP connection, it sends a message to the server, indicating a request for information about the top two CPU-consuming processes.
- 3) The client prints this information & closes the connection.
 - The **read()** function is used to receive the server's response, which is then printed to the console.
 - After the response is printed, the client closes the connection using **close()**, indicating that the communication with the server is complete.

NOTE: taskset was used for client and server code to pin task to specific CPUs

OUESTION 2

NOTE: Readings were taken after using taskset to pin tasks to specific CPUs.

(a) Single-threaded TCP client-server

Performance counter stats for './server'

| 15.39 msec | task-clock # 0.002 CPUs utilized |
|---------------------|------------------------------------|
| 7 | context-switches # 454.911 /sec |
| 0 | cpu-migrations # 0.000/sec |
| 143 | page-faults # 9.293 K/sec |
| 26,622,471 | cpu_core/cycles/ # 1.730 GHz |
| 52,089,912 | cpu_core/instructions/ |
| 9,614,604 | cpu_core/branches/ # 624.826 M/sec |
| 68,099 | cpu_core/branch-misses/ |
| 6.487866455 seconds | time elapsed |
| 0.003354000 seconds | user |
| 0.012300000 seconds | sys |

(b) Multi-threaded TCP client-server

Performance counter stats for './server'

| 22.56 msec | task-clock # 0.004 CPUs utilized |
|-------------|------------------------------------|
| 17 | context-switches # 753.653 /sec |
| 0 | cpu-migrations # 0.000/sec |
| 1,746 | page-faults # 77.405 K/sec |
| 68,613,095 | cpu_core/cycles/ # 3.042 GHz |
| 106,328,509 | cpu_core/instructions/ |
| 19,741,646 | cpu_core/branches/ # 875.197 M/sec |
| 233,960 | cpu_core/branch-misses/ |
| | |

| 6.13455021 seconds | time elapsed |
|--------------------|--------------|
|--------------------|--------------|

| 0.003746000 seconds | user |
|---------------------|------|
| 0.008343000 seconds | sys |

(c) TCP client-server using "select"

```
whentuBubuntu-dual-mackines-/Domiloads/A2/03$ sudo taskset -c 0 perf stat ./server

Server is now listening for inconing connections.

Accepted new Client connection

Sessage received from Client: Requesting top CPU processes

Message received from Client: Requesting top CPU processes

### State of CPU processes

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### Message received from Client Requesting top CPU processes

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### Message received from Client Requesting top CPU processes

### Message received from Client Requesting top CPU processes

### Message received from Client Requesting top
```

Performance counter stats for './server'

| 28.92 msec | task-clock # 0.007 CPUs utilized |
|---------------------|------------------------------------|
| 15 | context-switches # 518.673 /sec |
| 0 | cpu-migrations # 0.000/sec |
| 1,686 | page-faults # 58.299 K/sec |
| 67,269,978 | cpu_core/cycles/ # 2.326 GHz |
| 100,160,377 | cpu_core/instructions/ |
| 18,528,735 | cpu_core/branches/ # 640.691 M/sec |
| 231,798 | cpu_core/branch-misses/ |
| 4.050153388 seconds | time elapsed |
| 0.002850000 seconds | user |
| 0.012377000 seconds | sys |

1. Task Clock and CPUs Utilized:

• Single-Threaded Server:

o Task clock: 15.39 msec

o CPUs utilized: 0.002

• Multi-threaded Server:

Task clock: 22.56 msecCPUs utilized: 0.004

• Server using Select:

Task clock: 28.92 msecCPUs utilized: 0.007

The **Server using Select** spent the longest amount of time actively executing instructions on the CPU, while the **Single-Threaded Server** spent the least time. The **Multi-threaded Server** shows moderate execution time, which aligns with its more complex structure compared to the single-threaded version.

2. Context Switches (switches between kernel and user space):

• Single Thread: 7 context switches

• **MultiThread**: 17 context switches

• **Select-based**: 15 context switches

- The **single-threaded** server has the fewest context switches, as expected. It handles requests sequentially in a single process, leading to minimal task switching.
- The **multi-threaded** server has the most context switches because each thread may need to switch contexts frequently to handle multiple requests concurrently.
- The select-based server has context switches slightly lower than the multi-threaded one. The
 select-based model, though handling multiple connections, typically does so within a single
 thread, reducing the need for frequent switching but still causing more than a single-threaded
 model.

3. Page Faults:

Single-Threaded Server: 143 (9.293 K/sec)
Multi-threaded Server: 1,746 (77.405 K/sec)
Server using Select: 1,686 (58.299 K/sec)

- The **single-threaded** server has the fewest page faults, probably because it doesn't handle many requests simultaneously and thus has fewer memory access conflicts.
- Both **multi-threaded** and **select-based** servers have significantly higher page faults. This makes sense, as both handle multiple requests simultaneously, potentially needing more memory allocation, leading to more page faults.

4. CPU Core Cycles:

• Single-Threaded Server: 26,622,471 cycles (1.730 GHz)

• Multi-threaded Server: 68,613,095 cycles (3.042 GHz)

• Server using Select: 67,269,978 cycles (2.326 GHz)

The **multi-threaded** server uses the most CPU cycles, closely followed by the select-based server. This can be explained by both servers having more tasks to manage in parallel (multiple threads or connections). The **single-threaded** server, with fewer tasks at a time, uses fewer cycles overall.

5. Instructions:

Single-Threaded Server: 52,089,912 instructions
 Multi-threaded Server: 106,328,509 instructions
 Server using Select: 100,160,377 instructions

The **Multi-threaded Server** executes more than double the number of instructions compared to the **Single-Threaded Server**, as expected due to the complexity of handling multiple threads. The **Server using Select** executes a similar number of instructions to the multi-threaded server.

6. Branch Misses:

Single-Threaded Server: 68,099 branch misses
 Multi-threaded Server: 233,960 branch misses
 Server using Select: 231,798 branch misses

Branch misses represent inefficiency in CPU operations. The multi-threaded and select-based servers both have significantly more branch misses, which is expected because they handle more complex and parallelized logic, increasing the chance of mispredictions. The single-threaded server, with simpler flow control, has far fewer branch misses

7. User and Sys Time:

• Single-Threaded Server:

User: 0.003354 secondsSys: 0.012300 seconds

• Multi-threaded Server:

User: 0.003746 secondsSys: 0.008343 seconds

• Server using Select:

User: 0.002850 secondsSys: 0.012377 seconds

8. CPU Migrations: Zero for all the three as taskset pinned the task to specific CPUs

The **Single-Threaded Server** and **Select Server** spend a slightly higher amount of time in system calls (sys time), while the **Multi-threaded Server** has more time in user space (likely due to managing multiple threads).

Summary:

- → **Single Threaded**: Best for simple, low-concurrency applications where system resources (e.g., memory and CPU cycles) must be minimized.
- → **Multi-Threaded**: Suitable for high-load applications requiring parallel processing, though it incurs a higher system overhead.
- → **Select-based**: Best for moderately loaded servers where efficient handling of multiple connections is needed without the heavy overhead of multi-threading.