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VHP286
EE360C

G-S Project 1

Pre-Loop

- Create Hashmap for each internships matching
- an arraylist to store engagement of each student
- Hashmap for the slots in each internship
- inverse preference list for internships for constant access time
- Arraylist to store the place of each student in their proposals

Loop:

outer -

Sets the current student

the loop runs till students are satisfied

Inner - get the internship from the place of his next proposal

If the internship has empty spots:

- Student engaged to that internship
- decrease remaining slots
- update priority for the internship
- increase happy students by 1
- decrease slots for that internship

else

- check if the student can replace another student
- get lowest pref student of that internship
- If switch needed switch
- update all data structures

do this for all proposals

increment happy student if ^{no} remaining proposals left

return

Analysis:

The algorithm is not $O(m \cdot n)$ because when I want to switch students I check for the ^{lowest preferred} of the current paired students. Because of this In worst case I have to check all students to find the lowest preferred paired student. However this case is not common, generally the algorithm runs closer to $O(m \cdot n)$ because not ~~all~~ ^{all} students have 1 internship.