**Experiment**

**Aim:** Write a program to implement Knapsack Problem**.**

**Theory:**

**Problem Definition :**

We are given a empty knapsack of capacity „W‟ and we are given „n‟ different objects from i=1,2………n. Each object „i‟ has some positive weight „wi‟ and some profit value is associated with each object which is denoted as „pi‟. We want to fill the knapsack with total capacity „W‟ such that profit earned is maximum. When we solve this problem main goal is :

1. Choose only those objects that give maximum profit.

2. The total weight of selected objects should be <= W

Most problem have n input and require us to obtain a subset that satisfies some constraints is

called as a feasible solution. We are required to find a feasible solution that optimize (minimize or maximizes ) a given objective function. The feasible solution that does this is called an optimal solution.

**Example :**

n = 3

profit ->Pi = (25,24,15)

Weight (in kg) -> Wi = (18,15,10) where i = 1,2,3

and knapsack capacity = 20

**Feasible solutions :**

1. As kanpsack capacity is 20 , we first put W1( i.e 18 kg) into kanpsack , hence gained

profit is 25 ; but now we can put only 2 kg , in knapsack ,so we will put 2 kgs from W2 , hence profit from this 2kg of w2 is 3.2. Thus , total profit is 25+ 3.2 = 28.2

2. Now , we will put W3 (i.e 10 KG) into knapsack , hence gained profit is 15 , but we can put only 10 kg in knapsack , so we will put 10 kg form w2 , hence profit from this 10 kg of w2 is 16. Thus, total profit is 15+ 16 = 31

3. Now , we will put W2 (i.e 15 KG) into knapsack , hence gained profit is 24 , but we can put only 5 kg in knapsack , so we will put 5 kg form w3 , hence profit from this 5 kg of w3 is 7.5. Thus , total profit is 24+ 7.5 = 31.5 Hence , optimal solution for given problem is 3rd solution, which is giving total profit as 31.5.

The knapsack problems are often solved using dynamic programming, though no polynomialtime algorithm is known. Both the knapsack problem(s) and the subset sum problem are NPhard.

**Algorithm :**

1. Let „W‟ be the maximum weight of the knapsack

2. Let „wi‟ and „pi‟ be the weight and profit of individual items i.e. for i=1……n

3. Calculate pi / wi ratio and arrange that in decreasing order.

4. initially weight=0 and profit = 0

5. for i=1 to n

{

add item in knapsack till weight <= W

profit = profit + pi

}

6. Stop

**Conclusion :** Most problems have n inputs and require us to obtain a subset that satisfies some constraints is called as feasible solution. In knapsack problem, greedy strategy, we find out feasible solution that optimize (minimizes or maximizes) a given objective function. A, feasible solution that does this is called an optimal solution. Hence, in knapsack problem we found out an optimal solution. The knapsack algorithm takes θ(*nw*) times, broken up as follows: θ(*nw*) times to fill the *c*-table, which has (*n* +*1*).(*w*+1) entries, each requiring θ(1) time to compute. *O*(*n*) time to trace the solution, because the tracing process starts in row *n* of the table and moves up 1 row at each step.