

Key Generation -> how 16 key is generated

## 4 2 3 1 Double DES :-

Since DES attack was vulnerable to brute force attack, variations of DES called multiple DES were introduced.

- Use 2 DEF keys (56+56) = 112 bit key

→ Double encryption  
occurs as follows

$$P \rightarrow E(K_1, P)$$

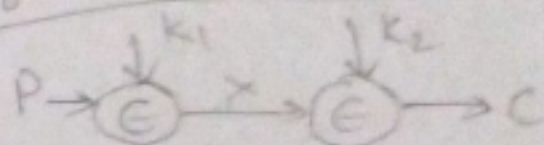
↓

$$E(K_2, E(K_1, P)) = \text{Cipher}$$

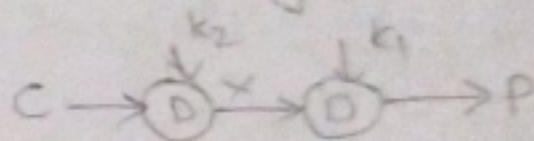
→ For Decryption

1<sup>st</sup> Decryption using Key  $K_2$   
which produce single  
encrypted cipher text

- This 64 bit middle text  
is then decrypted using  
the Key  $K_1$  to get the  
plain text.



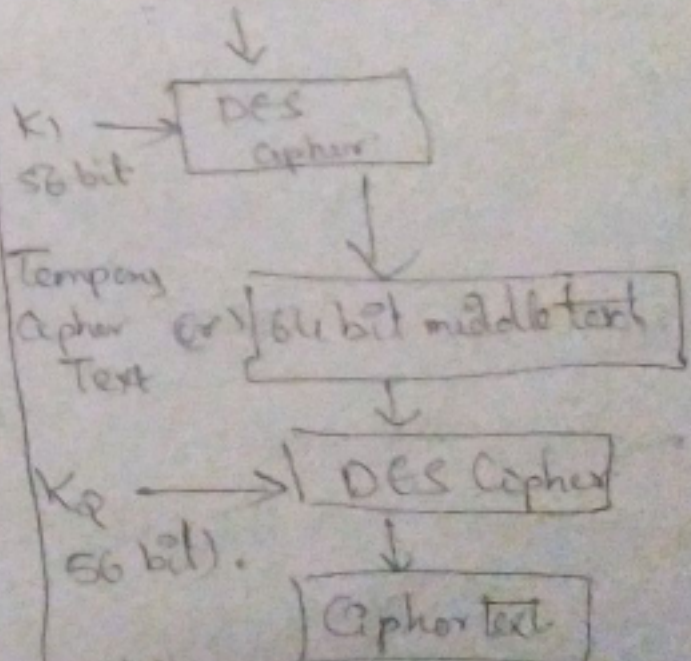
Encryption



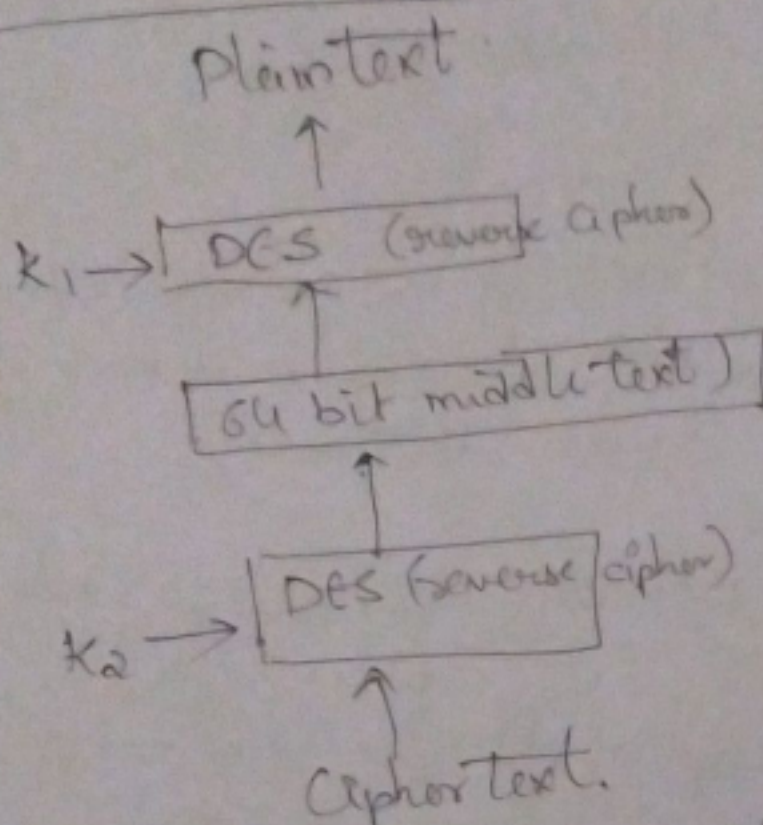
Decryption.

Data Encryption.

64 bit P.T.



## Double DES Decryption:-



The decryption, key are applied in reverse order

$$\boxed{\text{Plain} = D(K_1, D(K_2, C))}$$

First this will happen

## Drawback of Double DES

Meet in the middle attack.

This attack involves encryption from one end & decryption in other end then "making the results in the middle and ~~reverse~~ hence the name

MIM attack



Some pairs

plain text known & Cipher text known

↓  
Encrypt pairs for  
all  $2^{56}$  possible  
values of  $K_1$

↓  
decrypt pairs of  
all  $2^{56}$  possible  
values of  $K_2$

no. of  
rows in  
Table =  
no. of  
possible  
secret  
keys.

Plain	Cipher	Key
ABCD	VYMP	
####	MXTP	
XT		

sort the  
result  
in  
table

Cipher	Middle	Cipher
S		

$2^{56}$  possible combinations  
of keys

Result in  
corresponding value for  $K_2$

→  $(K_1, K_2)$

→  $(K_1, K_2)$

We will compare these values with values of  
the 1st table computed earlier.

$$\text{Decrypt}(K_2, C) = \text{Encrypt}(K_1, P)$$

∴  $(K_1, K_2)$  is the key pair used.

This attack requires some pairs of  
Plaintext / Ciphertext pairs.

Let us assume plaintext = P ; Ciphertext = C.

This attack proceeds as follows:

(i) Encrypt P for all  $2^{56}$  possible values of  $K_1$  & store the results in a table & sort it.

(ii) Now, decrypt C using all  $2^{56}$  possible values of  $K_2$ . As each <sup>undone</sup> result is produced, check against the table for a match.

(iii) When there is a match, we have located a possibly correct pair of keys.

Note, Now more than 1 pair keys may result in a match, but their no. of pairs will be small, we should.

try each possible pair of keys.

- So, it takes twice as long to break double DES using brute force attack

Beoz, DES has  $2^{56}$  bit security, Double DES

has  $2 \times 2^{56} = 2^{57}$  key security

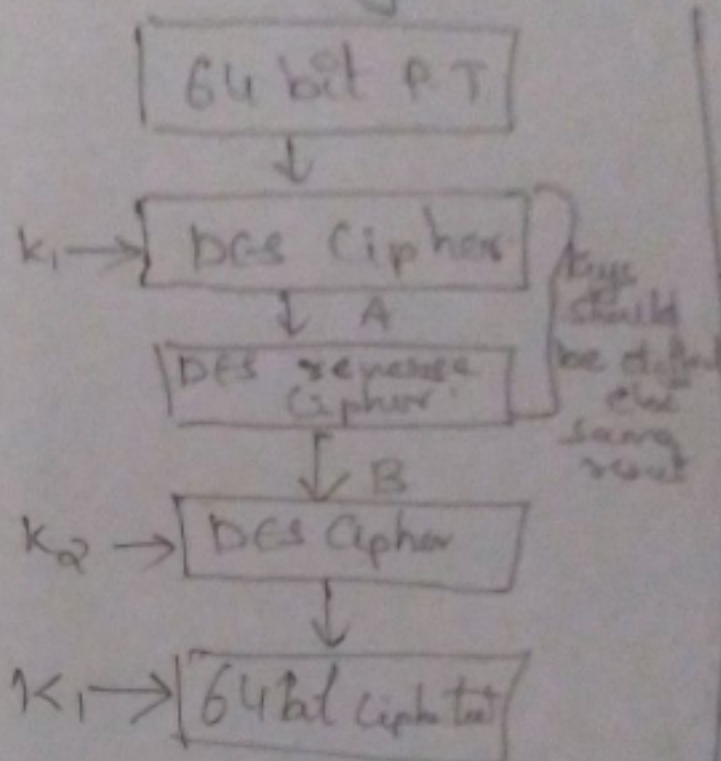


4.3.4

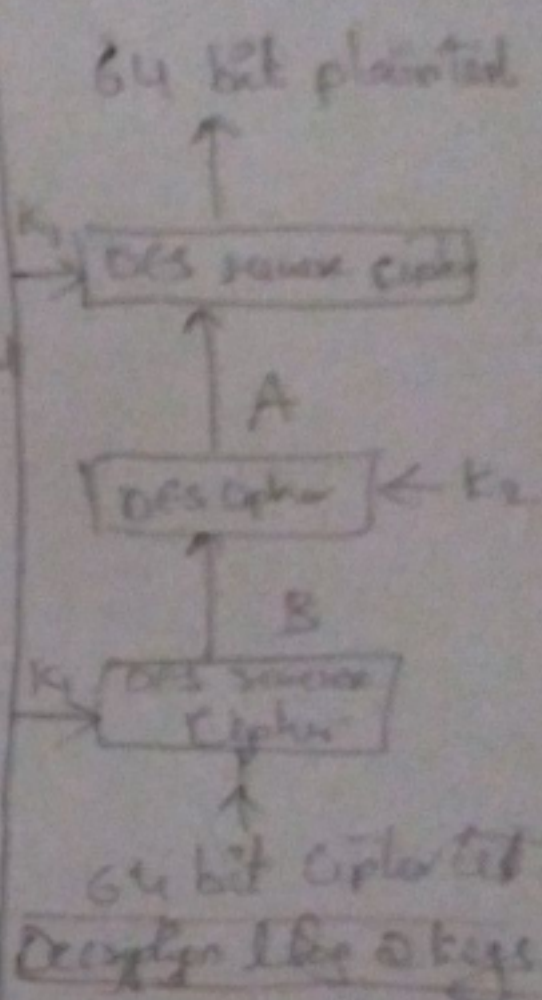
# TRIPLE DES: 3DES

2 or 3 keys are used

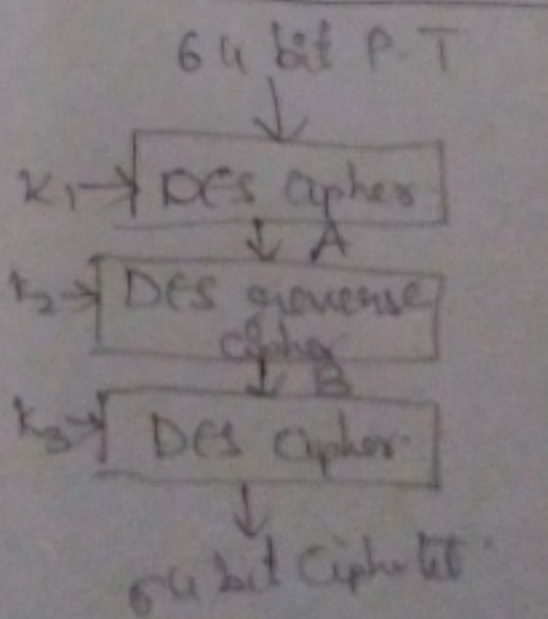
- Much stronger than 2DES



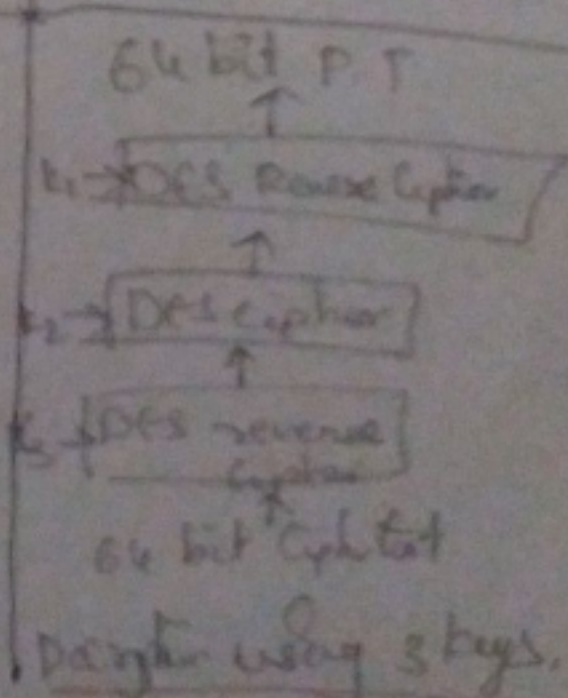
Encryption using 2 keys



Decryption using 2 keys



Encryption using 3 keys



Decryption using 3 keys