

Its work a principal LIFO (Last In first out)
 FILO

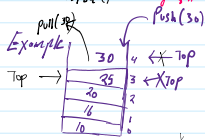
ADT of Stack → Abstract Data Type.

ADT contain data representation or operation on the stack.

- Data
1. Space for storing elements
 2. Top pointer which pointing the topmost element.

Operations

1. push()
2. pop()
3. peek()
4. Stack top() → for knowing what is topmost
5. isEmpty() → knowing empty. Value in Stack.
6. isFull() → knowing full.



peek(i) → 16
 Looking the value at index i

Imp/Obtaining

It is done by using two storing data structure.

- 1) Array
- 2) Linked List.

→ Array

First Understand Structure/Class.

<name structure or class first>

#include <iostream>
 using namespace std; // Implementation by Dynamically.

```
class Stack {
public: // for accessing is another class.
    int size; int top; int *s; // initialize
    // make constructor to initialize.
    Stack(int sz) { this->top = -1; this->size = sz;
        s = new int[sz]; // this new allocate memory
        // in heap so its dynamically.
    }
    // And Code of isEmpty, push, pop, peek.
    bool isEmpty() {
        if (top == -1)
            return true;
        return false;
    }
    void push(int val) {
        if (top == size)
            cout << "overloaded" << endl;
        else {
            // top ++;
            s[++top] = val;
        }
    }
    int peek() {
        return s[top];
    }
    int pop() {
        if (isEmpty()) {
            cout << "underflow" << endl;
            return -1;
        }
        return s[top--];
    }
};
```

// we use call stack variable outside the class the value of top ref.

```
void display(Stack st) {
    for (int i = 0; i < st.top; i++) {
        cout << st.s[i] << " ";
    }
}
```

```
int main() {
    Stack st(inputSize); // here st is object
    st.push(5);
    st.push(1000);
    st.push(100);
    cout << "deleted value -> " << st.pop() << endl; // 100
    cout << st.pop() << endl; // 1000
    display(st); // 1000 5
    return 0;
}
```

→ Linked List

Stack using LL, in this we make top as a head.
 In this first push element goes into last node.
 The top/last push/pop/peek is in the head pointer.
 Note:- peek() we can say about a node at first position.
 pop() we can say delete a first element in LL.


```

Empty
if (top == NULL)

full
Node * t = new Node;
if (t == NULL)

```

Recursion

```

struct Node {
    int data;
    Node *next;
}

Node * push(int val) {
    Node * n = new Node(val);
    if (n == NULL) return NULL;
    if (top == NULL)
        top = n;
    else
        n->next = top;
    top = n;
}

```

Infix to Postfix Conversion

- 1) What is Postfix.
- 2) Why Postfix.
- 3) Precedence.
- 4) Infix to Postfix Conversion.

There are notation:-

1) Infix: operand operator operand.
eg. $a + b$

2) Prefix: opt operand operand.
eg. $+ab$

3) Postfix: operand operand opt.
eg. abt

Note: Postfix is mostly used.

Infix to Postfix Conversion

Symbol	Precedence
$+, -$	1
$\times, /$	2
$()$	3

i) $a + b \times c$
 Prefix: $(a + [b \times c])$
 Postfix: $a + [b \times c]$
 $\Rightarrow (+abc)$ $(abc+)$

ii) $a + b + c \times d$
 Here both are same operators but give left op. more precedence.

Prefix: $a + b + [c \times d]$
 Postfix: $a + b + [c \times d]$
 $[a + b] + [c \times d]$ $[a + b] + [c \times d]$
 $+ + ab \times cd$ $ab + cd +$

iii) $(a + b) \times (c \times d)$

Prefix: $[a + b] \times [c \times d]$
 Postfix: $[a + b] \times [c \times d]$
 $* + ab - cd$ $ab + cd - *$

iv) $a \times b + c$ v) $a \times b + c$
 $a(b+c)$ $(ab)c +$
 $(a + (b \times c))$ $((ab) + c)$

Associativity

If precedence is same then go with associativity.

Op	Pre	Ass
$+$	1	L-R
$\times, /$	2	L-R
$^$	3	R-L
$-$	4	R-L
$()$	5	L-R

Ex: $a + b + c$
 L-R
 $(a + b) + c$
 $(a + b) + c$

ii) $a \times b \times c$
 R-L
 $a \times (b \times c)$

Unary Operator

i) $-a \rightarrow$ precedence 4
 \rightarrow Asso. R-L

ii) $--a$ R-L
 $(-(-a))$
 $\rightarrow --a$
 $\rightarrow -(-a) \rightarrow a$

$--a$ R-L
 $(-(-a))$
 $\xrightarrow{Pre} --a$
 $\xrightarrow{Post} (-[-a]) \rightarrow a--$

(ii) $*P$ Infix form
 Pre: $*P$ Post: $P*$
 $*+P \rightarrow (*(*P))$
 Asso: (R-L)

(iii) $n!$ Pre: $n!$ Post: $n!$
 (iv) $\log x$ Pre: $\log x$ Post: $x \log$

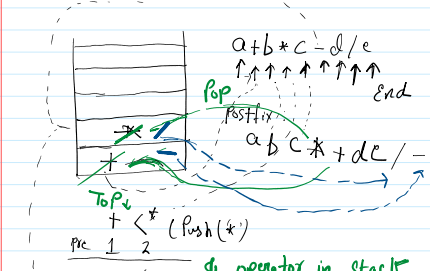
Unary $\rightarrow - + ! \log$

Example Convert into Postfix
 $--a + b * \log n!$
 Using precedence and associativity:
 R-L

$-a + b * \log[n!]$
 $-a + b * [n! \log]$
 $[a-] + b * [n! \log]$
 $[a-] + [b n! \log *]$
 $(a - b n! \log * +)$ Postfix

Infix To Postfix conversion using Stack.

$a + b * c - d / e$
 $\Rightarrow abc * + de / -$
 Sym Pre Asso
 $+ - 1$ L-R
 $* / 2$ L-R
 $a b c 3$ L-R

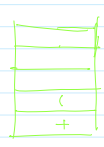


If operator in stack is in lower precedence, then push. otherwise pop.
 Pop() $+ < *$ (Push(*))
 Pre 1 2
 Pop() $+ -$
 then also pop() $-$

Now stack is empty then push() that is ' $-$ '.
 At end of expression.
 Empty the stack and put into postfix.
 if stack is not empty.
 End
 pop() (Null)

Post $abc++ \rightarrow ++abc$
 $a \pm b \pm c$
 Pre.

Lower Precedence Equal Precedence
 $a - b * c \div d / e$
 Doing postfix:
 postfix = $abc * + de / -$



④
 $a + b + (c * d + e)$
 $(cd * +) + e$
 $a + b + : cd * c -$
 $ab + + : :$
 $ab + cd * e +$
 $a b + c$

char c
 int prec () {
 if (c == '+' || c == '-')
 return 1;
 if (c == '*' || c == '/')
 return 2;
 if (c == '(')
 return -1;
 }

$+ -$
 $* /$
 $()$

Symbol	Stack	Postfix
a	a	a
+	+	a
b	+	ab
*	* +	ab
c	* +	abc
-	-	abc * +
d	-	abc * + d
/	/ -	abc * + d

	in	uol
-	-	abc*+
d	-	abc*d
/	/-	abc*+d
c	/-	abc*+dc
-		abc*+dc/-

Conclusion: All the symbols in stack which is greater or equal than $\text{pop}()$.
 * If the operator in stack is lower precedence or empty then $\text{flush}()$.

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