

Chapter 1

Introduction

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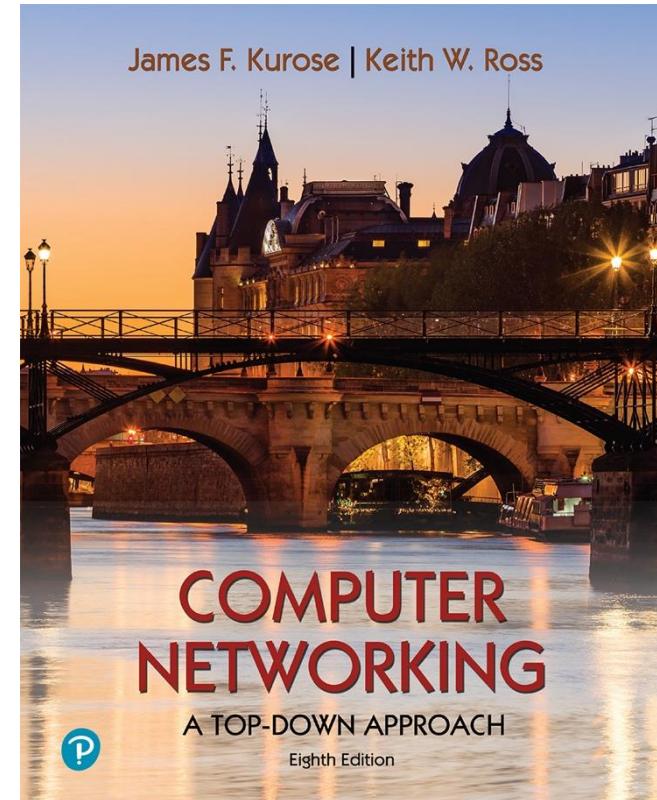
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*Computer Networking: A
Top-Down Approach*
8th edition
Jim Kurose, Keith Ross
Pearson, 2020

Chapter 1: introduction

Chapter goal:

- Get “feel,” “big picture,” introduction to terminology
 - more depth, detail *later* in course
- Approach:
 - use Internet as example



Overview/roadmap:

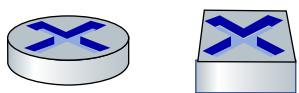
- What *is* the Internet?
- What *is* a protocol?
- **Network edge:** hosts, access network, physical media
- **Network core:** packet/circuit switching, internet structure
- **Performance:** loss, delay, throughput
- Security
- Protocol layers, service models
- History

The Internet: a “nuts and bolts” view



Billions of connected computing *devices*:

- *hosts* = end systems
- running *network apps* at Internet's “edge”



Packet switches: forward packets (chunks of data)

- routers, switches



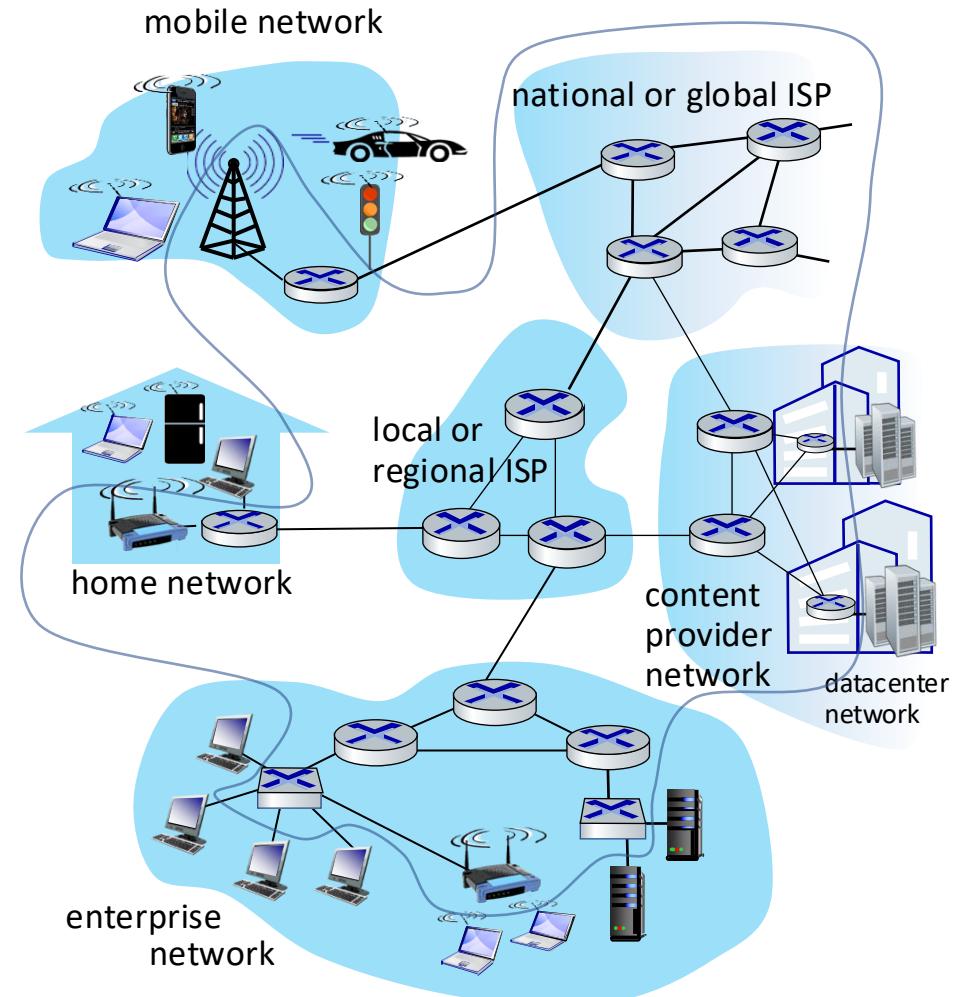
Communication links

- fiber, copper, radio, satellite
- transmission rate: *bandwidth*



Networks

- collection of devices, routers, links: managed by an organization



“Fun” Internet-connected devices



Amazon Echo



Internet refrigerator



Security Camera



IP picture frame



Slingbox: remote
control cable TV



Pacemaker & Monitor



Web-enabled toaster +
weather forecaster



Tweet-a-watt:
monitor energy use



AR devices

Internet phones



sensorized,
bed
mattress

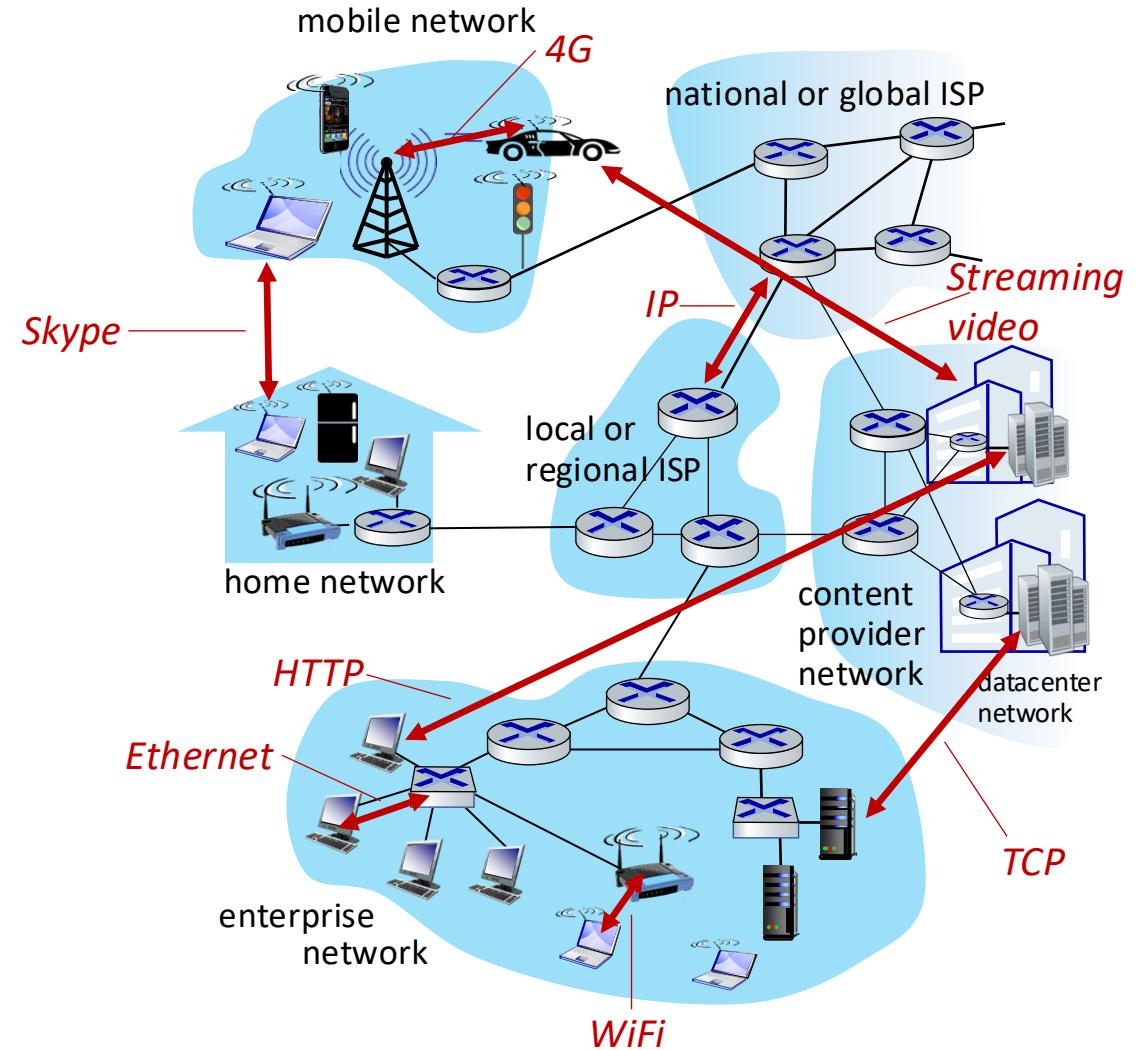


Fitbit

Others?

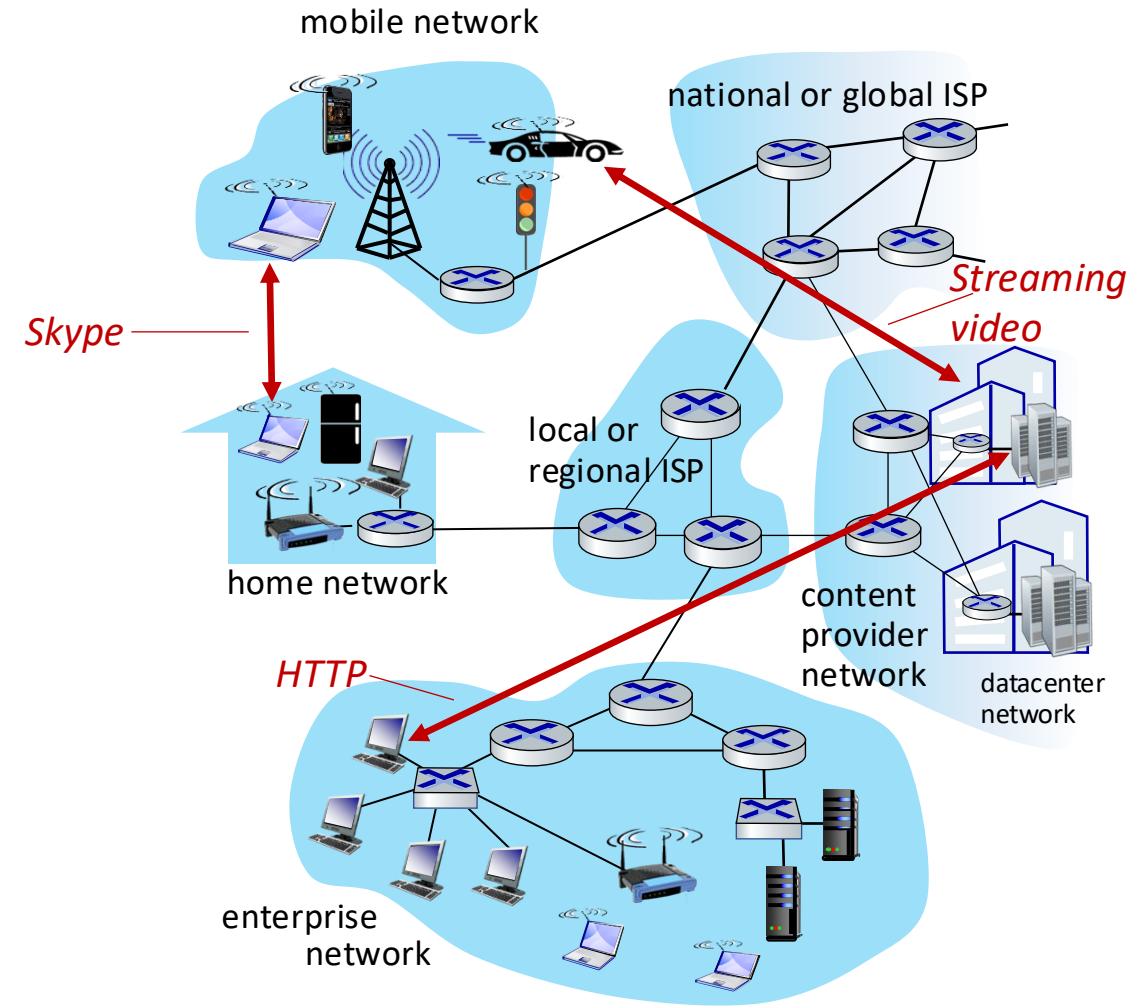
The Internet: a “nuts and bolts” view

- *Internet: “network of networks”*
 - Interconnected ISPs
- *protocols are everywhere*
 - control sending, receiving of messages
 - e.g., HTTP (Web), protocols for streaming video, Skype, TCP, IP, WiFi, 4G, Ethernet
- *Internet standards*
 - RFC: Request for Comments
 - IETF: Internet Engineering Task Force



The Internet: a “service” view

- *Infrastructure* that provides services to applications:
 - Web, streaming video, multimedia teleconferencing, email, games, e-commerce, social media, interconnected appliances, ...
- provides *programming interface* to distributed applications:
 - “hooks” allowing sending/receiving apps to “connect” to, use Internet transport service
 - provides service options, analogous to postal service



What's a protocol?

Human protocols:

- “what’s the time?”
- “I have a question”
- introductions

... specific messages sent
... specific actions taken
when message received,
or other events

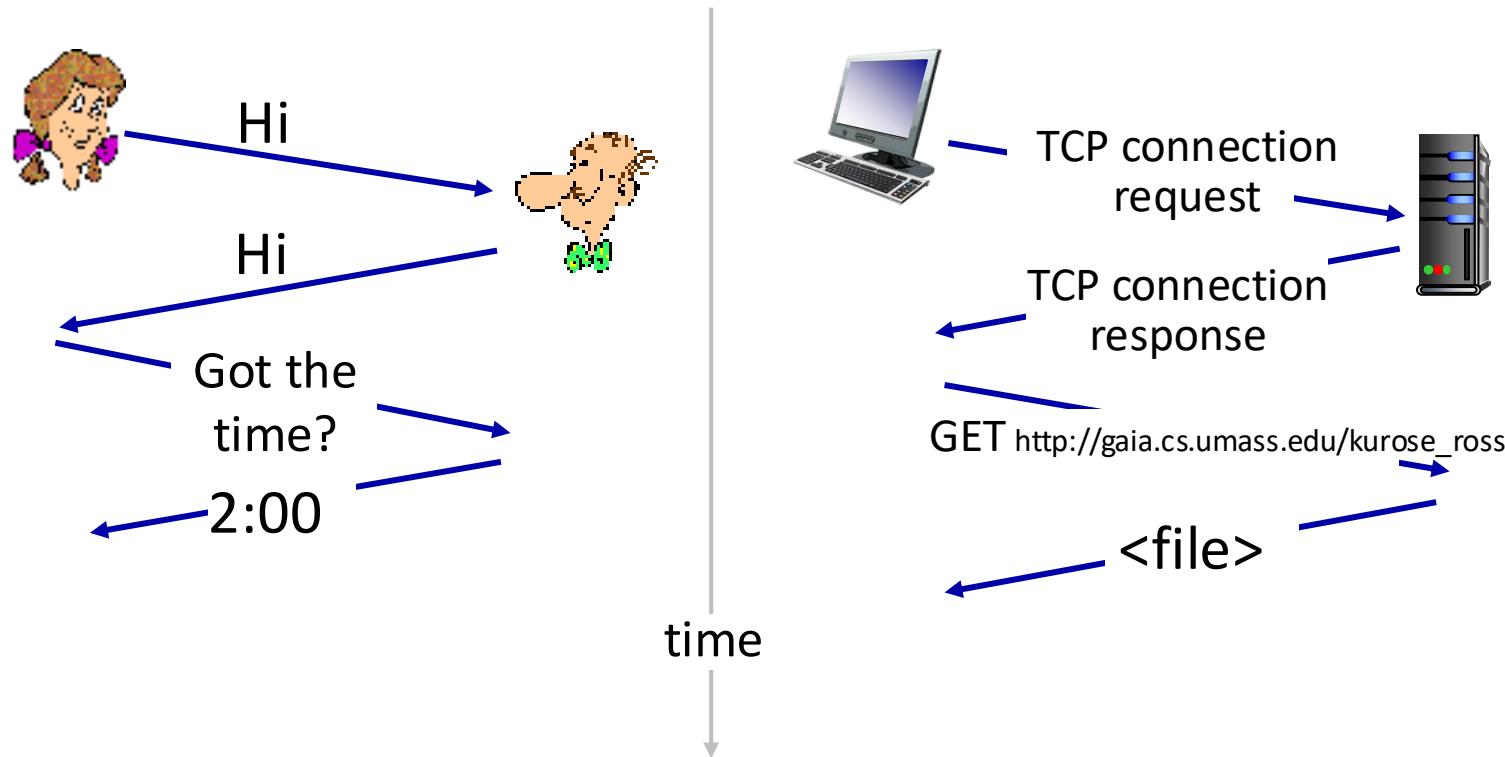
Network protocols:

- computers (devices) rather than humans
- all communication activity in Internet governed by protocols

*Protocols define the **format, order** of messages sent and received among network entities, and **actions taken** on msg transmission, receipt*

What's a protocol?

A human protocol and a computer network protocol:



Q: other human protocols?

Knowledge Check

Which of the following descriptions below correspond to a "nuts-and-bolts" view of the Internet?

1. A collection of billions of computing devices, and packet switches interconnected by links.
2. A platform for building network applications.
3. A collection of hardware and software components executing protocols that define the format and the order of messages exchanged between two or more communicating entities, as well as the actions taken on the transmission and/or receipt of a message or other event.
4. A place I go for information, entertainment, and to communicate with people.
5. A "network of networks".

Knowledge Check

Which of the following human scenarios involve a protocol?

1. A person sleeping.
2. A person reading a book.
3. A student raising her/his hand to ask a really insightful question, followed by the teacher acknowledging the student, listening carefully to the question, and responding with a clear, insightful answer. And then thanking the student for the question, since teachers *love* to get questions.
4. One person asking, and getting, the time to/from another person.
5. Two people introducing themselves to each other.

Chapter 1: roadmap

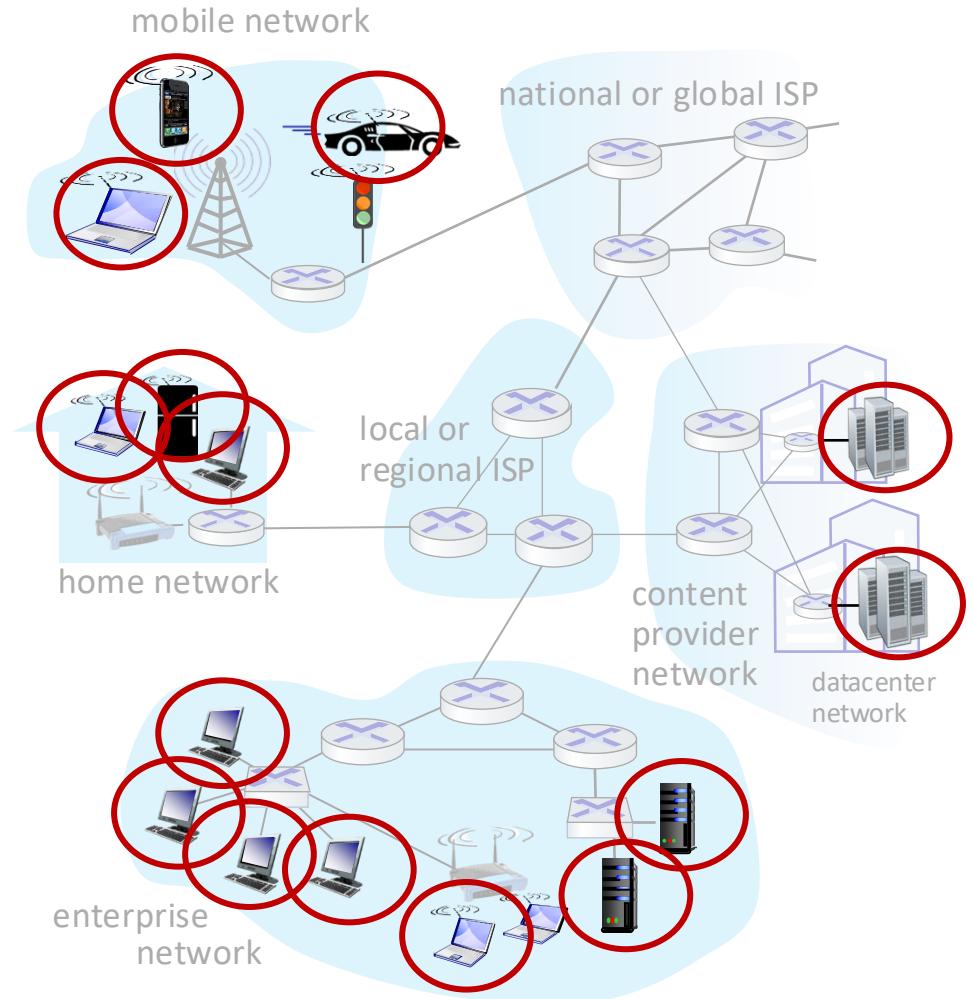
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- Network core: packet/circuit switching, internet structure
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A closer look at Internet structure

Network edge:

- hosts: clients and servers
- servers often in data centers



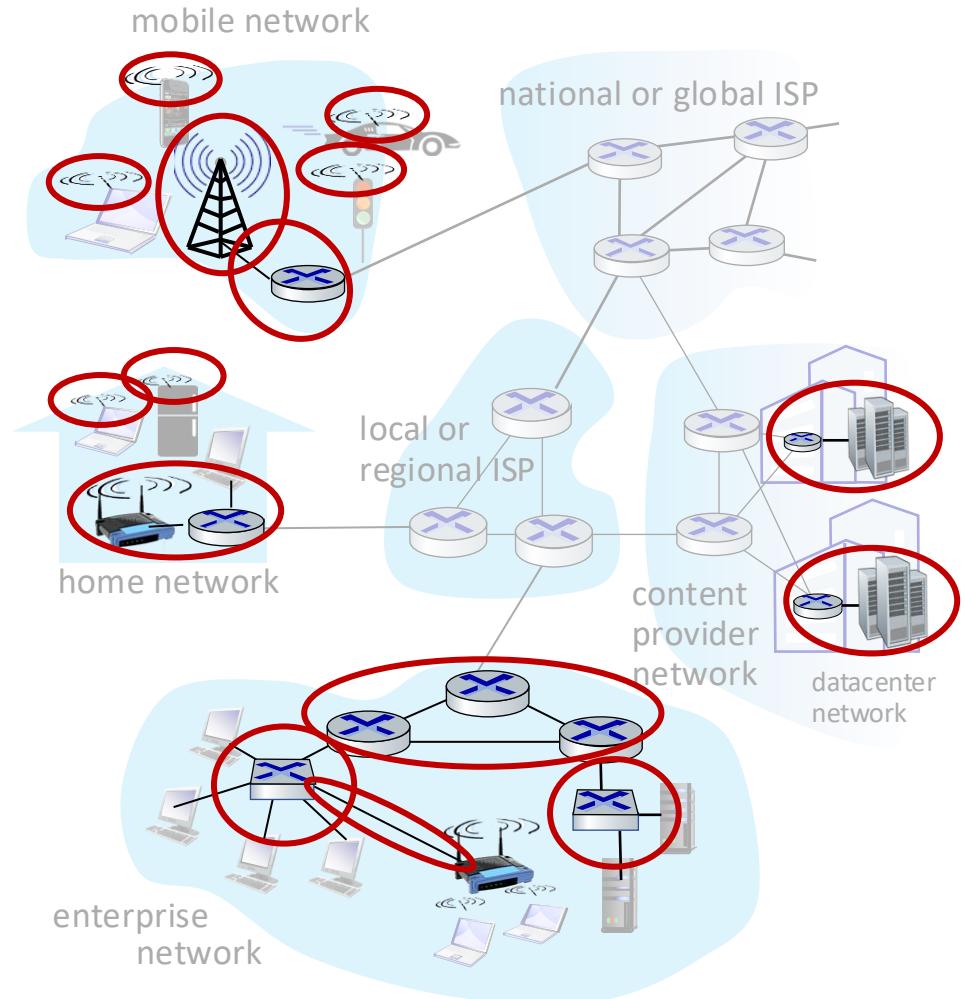
A closer look at Internet structure

Network edge:

- hosts: clients and servers
- servers often in data centers

Access networks, physical media:

- wired, wireless communication links



A closer look at Internet structure

Network edge:

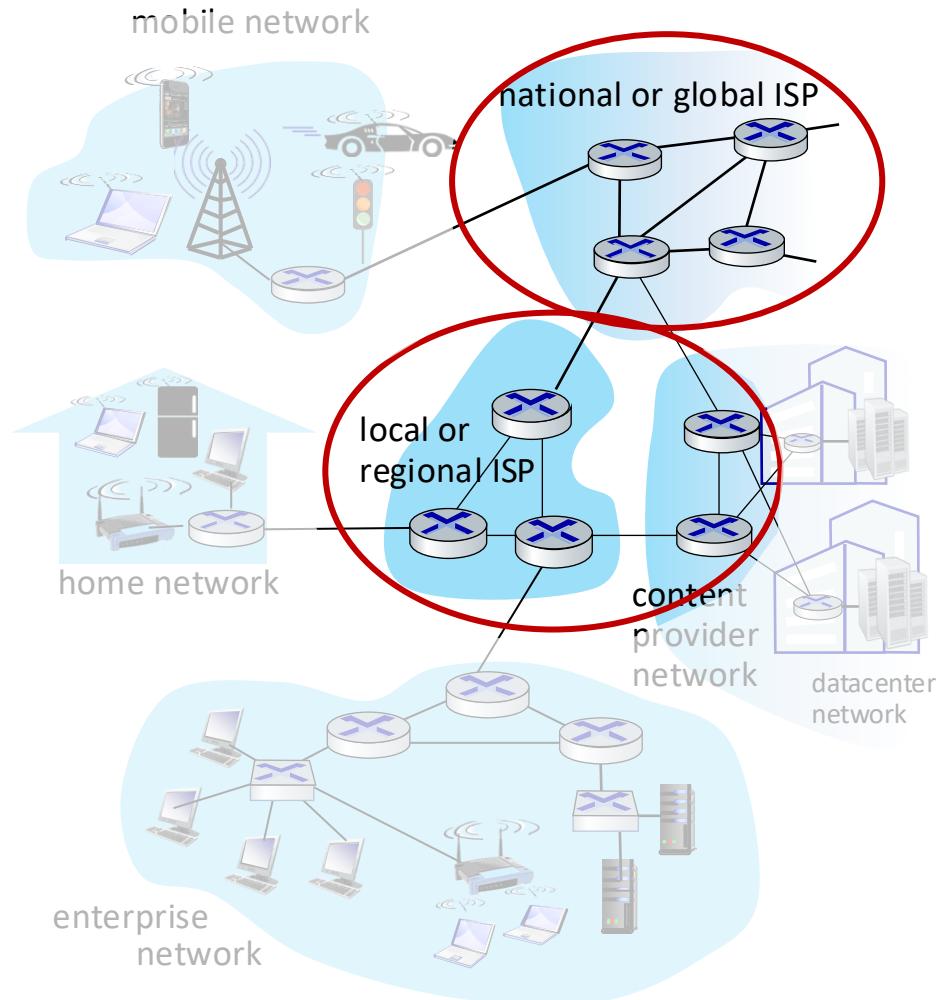
- hosts: clients and servers
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Access networks, physical media:

- wired, wireless communication links

Network core:

- interconnected routers
- network of networks



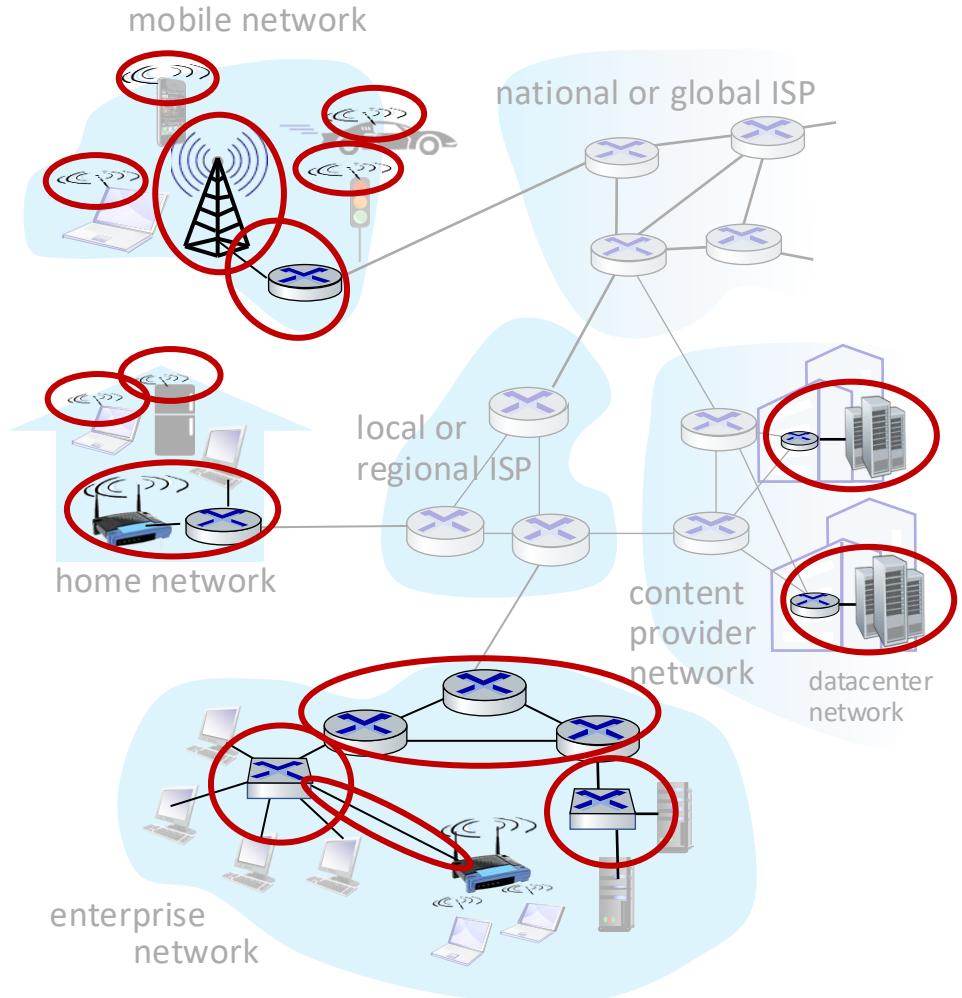
Access networks and physical media

*Q: How to connect end systems
to edge router?*

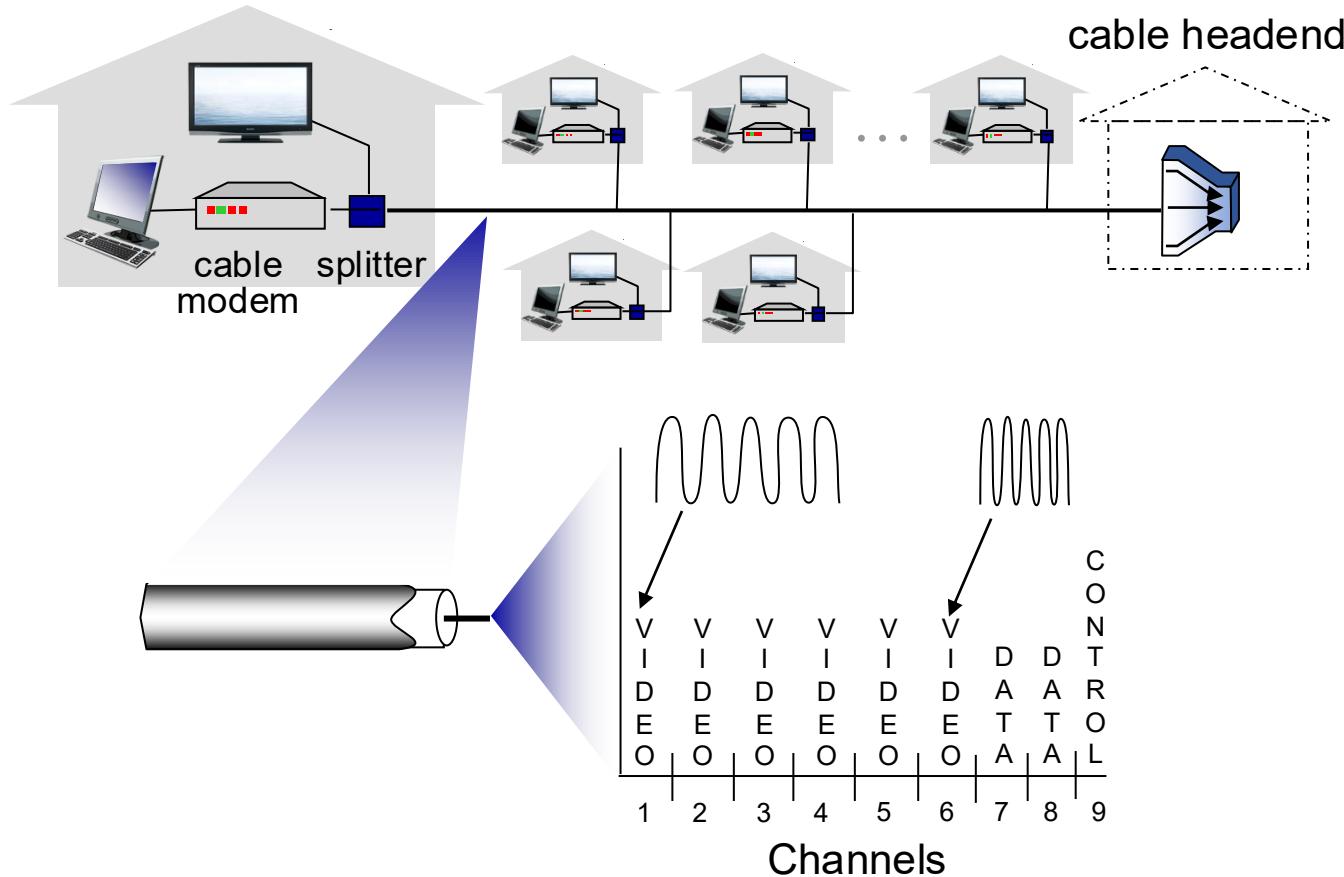
- residential access nets
- institutional access networks (school, company)
- mobile access networks (WiFi, 4G/5G)

What to look for:

- transmission rate (bits per second) of access network?
- shared or dedicated access among users?

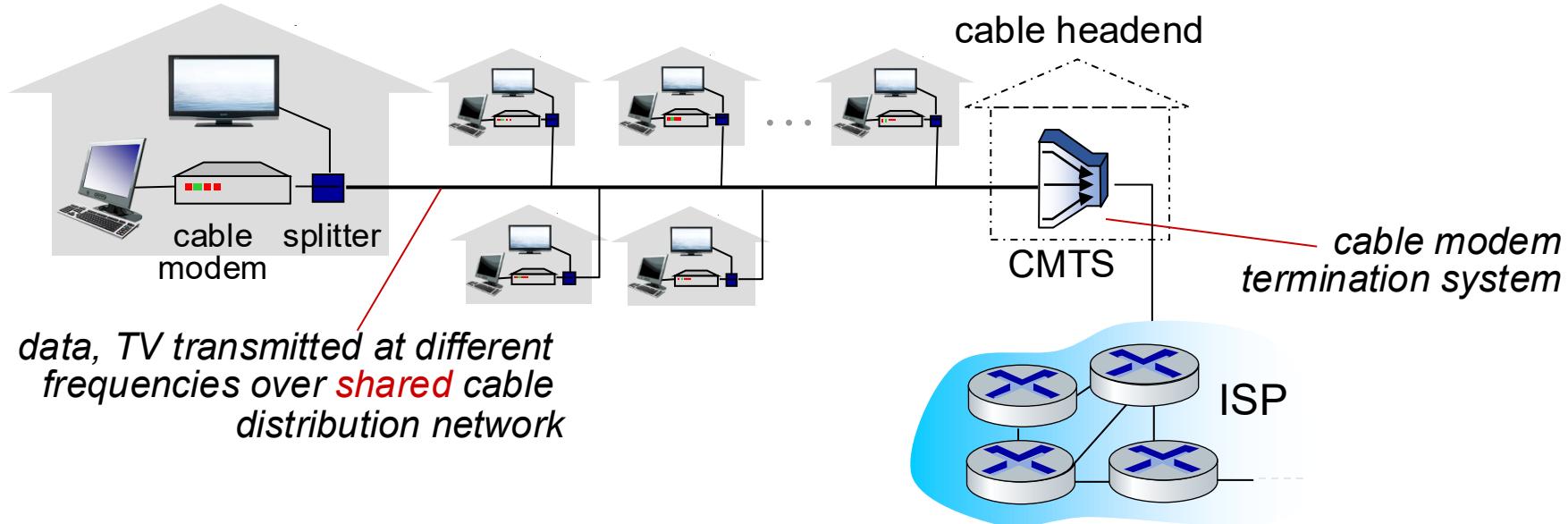


Access networks: cable-based access



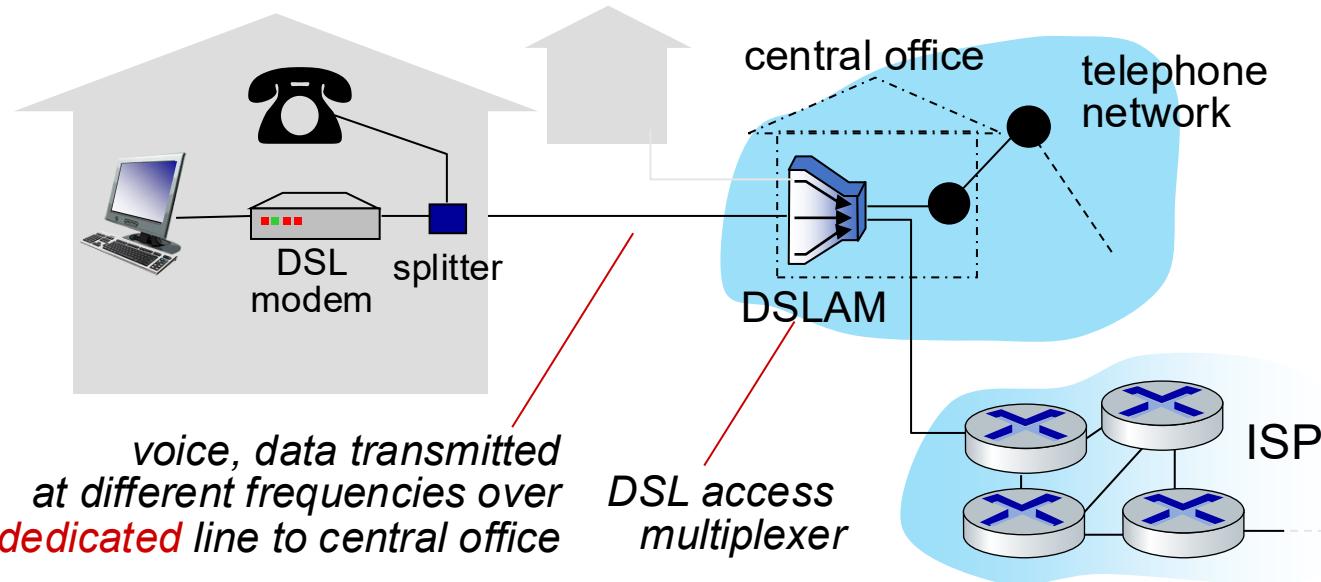
frequency division multiplexing (FDM): different channels transmitted in different frequency bands

Access networks: cable-based access



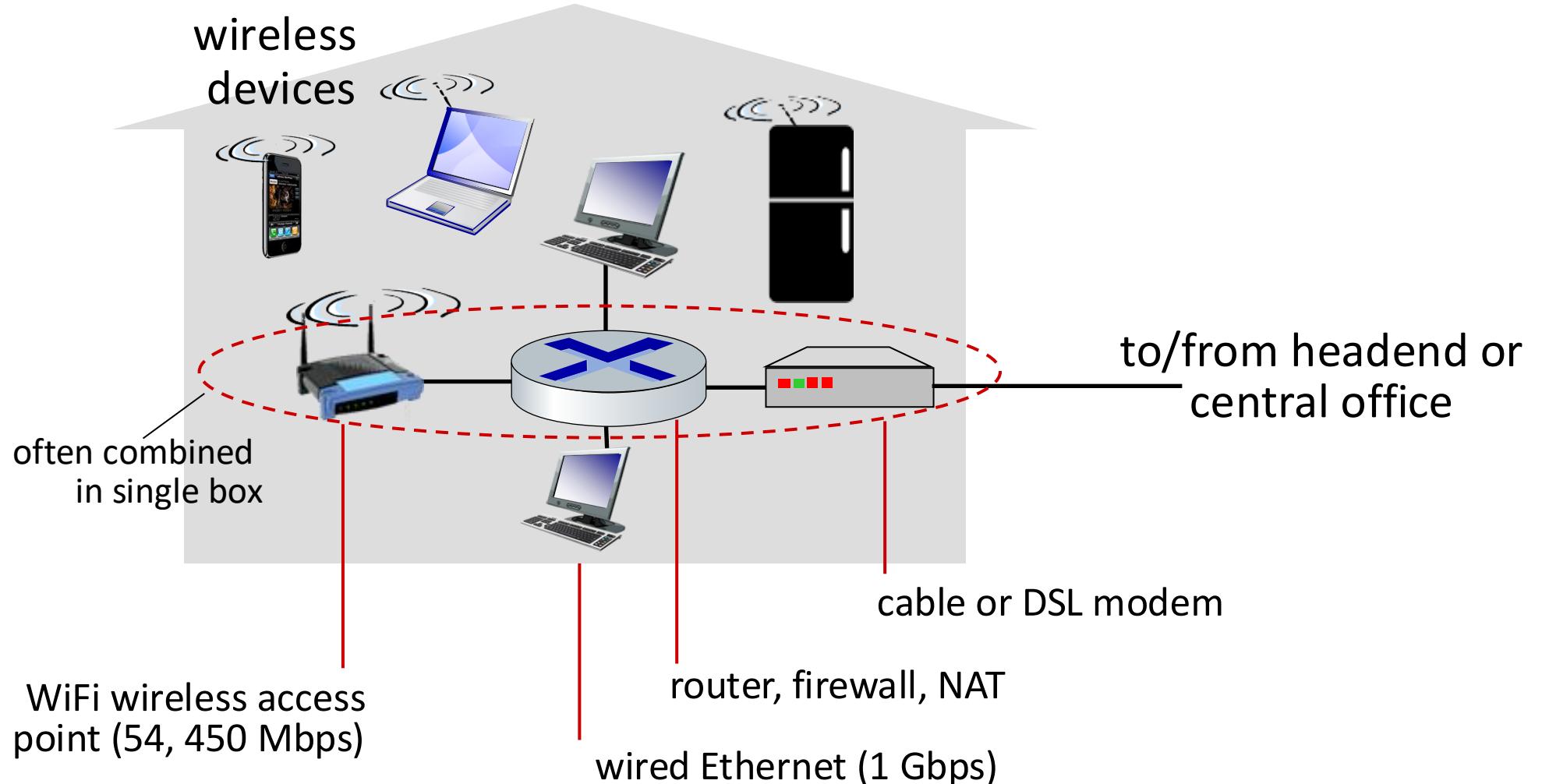
- HFC: hybrid fiber coax
 - asymmetric: up to 40 Mbps – 1.2 Gbs downstream transmission rate, 30-100 Mbps upstream transmission rate
- network of cable, fiber attaches homes to ISP router
 - homes **share access network** to cable headend

Access networks: digital subscriber line (DSL)



- use *existing* telephone line to central office DSLAM
 - data over DSL phone line goes to Internet
 - voice over DSL phone line goes to telephone net
- 24-52 Mbps (or more) dedicated downstream transmission rate
- 3.5-16 Mbps dedicated upstream transmission rate

Access networks: home networks



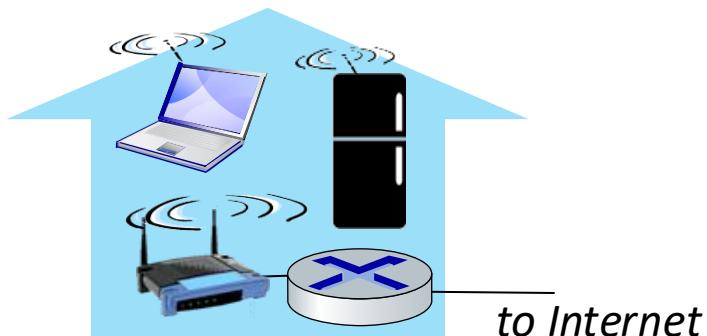
Wireless access networks

Shared *wireless* access network connects end system to router

- via base station aka “access point”

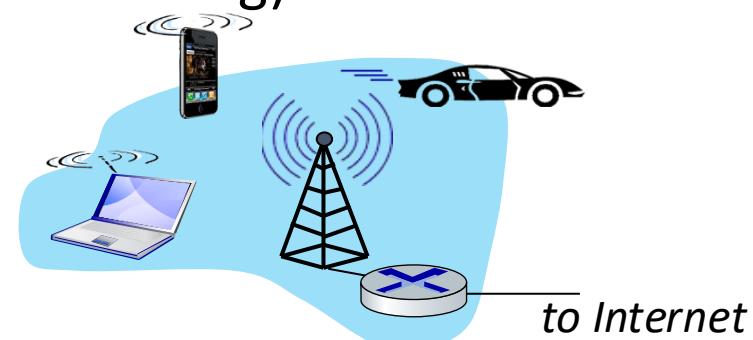
Wireless local area networks (WLANs)

- typically within or around building (~100 ft)
- 802.11b/g/n (WiFi): 11, 54, 450 Mbps transmission rate



Wide-area cellular access networks

- provided by mobile, cellular network operator (10's km)
- 10's Mbps
- 4G and 5G cellular networks (6G coming)



Actually, as of now, Wifi

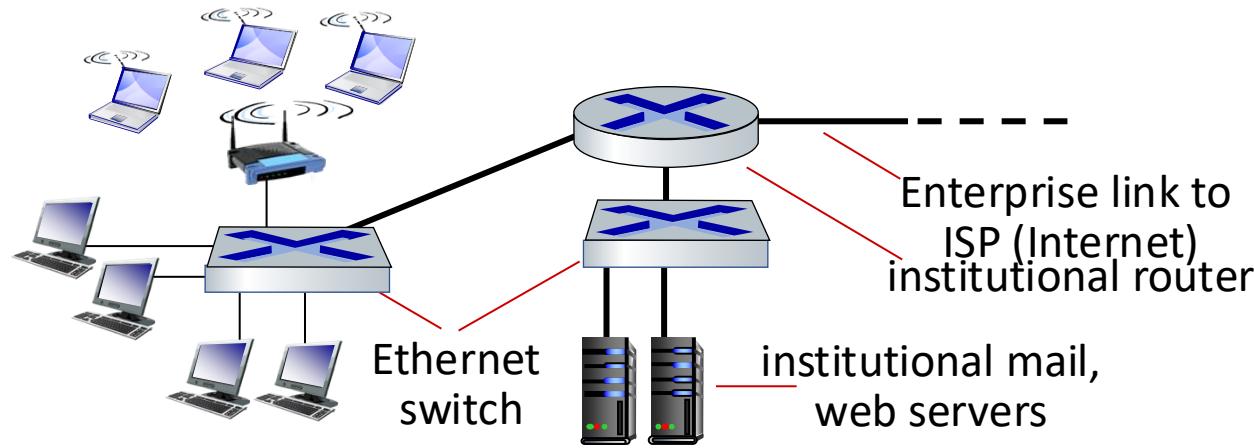
Wi-Fi Generations

Generation	IEEE Standard	Maximum Linkrate (Mbit/s)	Adopted	Radio Frequency (GHz) ^[1]
Wi-Fi 7	802.11be	40000	TBA	2.4/5/6
Wi-Fi 6E	802.11ax	600 to 9608	2020	2.4/5/6
Wi-Fi 6			2019	2.4/5
Wi-Fi 5	802.11ac	433 to 6933	2014	5
Wi-Fi 4	802.11n	72 to 600	2008	2.4/5
(Wi-Fi 3*)	802.11g	6 to 54	2003	2.4
(Wi-Fi 2*)	802.11a	6 to 54	1999	5
(Wi-Fi 1*)	802.11b	1 to 11	1999	2.4
(Wi-Fi 0*)	802.11	1 to 2	1997	2.4

Actually, as of now, Cellular Data

xG	Typical Spectrum Available (MHz) ¹	Technology	Access Scheme	Highest Modulation	Channel Bandwidth (MHz)	MIMO ²	Voice Capacity (Users/sector)	Peak Terminal DL Data Rate (Mbps)	Average DL Sector Throughput (Mbps)
2G	170	GSM	TDMA + FDMA	GMSK	5 ³	1X 1 (DL) 1X 2 (UL)	30 (AMR 5.9 codec)	0.171	0.28
		EDGE	TDMA + FDMA	8-PSK	5	1X 1 (DL) 1X 2 (UL)	N/A	0.237	0.62
		IS-95	CDMA	QPSK	1.25	1X 1 (DL) 1X 2 (UL)	20 (EVRC codec)	0.064	0.1
3G	260	CDMA 2000 1X	CDMA	QPSK	1.25	1X1 (DL) 1X 2 (UL)	35 (EVRC codec) ⁴	0.154	0.19
		EV-DO REV. 0	CDMA + TDMA	16 QAM	1.25	1X 2 (DL) 1X 2 (UL)	N/A	2.4	1
		EV-DO REV. B	CDMA + TDMA + CA	64 QAM	1.25 X 3 CA	1X 2 (DL) 1X 2 (UL)	N/A	12.9	4.5
		UMTS	CDMA	QPSK	5	1X 2 (DL) 1X 2 (UL)	82 (AMR 7.95)	2.69	0.55
		HSPA	CDMA + TDMA	16 QAM	5	1X 2 (DL) 1X 2 (UL)	N/A	14.4	3.75
		HSPA++ (Rel 7 and above)	CDMA + TDMA + CA	64 QAM	5 X 3 CA	4X 2 (DL) 1X 4 (UL)	N/A	128	14.175
4G	480	LTE (FDD)	OFDMA/SC-FDMA + TDMA	64 QAM	10 (1.4, 3, 5, 10, 15, 20 MHz supported)	4X 2 (DL) 1X 4 (UL), Multi-user MIMO	416 ⁵ (AMR-WB 12.65 codec)	60.9	13
		LTE Advanced	OFDMA/SC-FDMA + CA	64 QAM	10 X 2 CA (typical)	8X 4 (DL), 4X N (UL), Multi-user MIMO		248 ⁶	35
5G	~1-2 GHz	Sub 6 GHz	OFDMA/FILTERED + CA	≥ 64 QAM	~40 MHz	~8 BTS, ~4 Mobile		~500 ⁷	~80
		mm Wave	OFDMA / Single Carrier	QPSK / 16 QAM	~1 GHz	~ 64 Tx, Rx?	N/A	~2000	~700

Access networks: enterprise networks

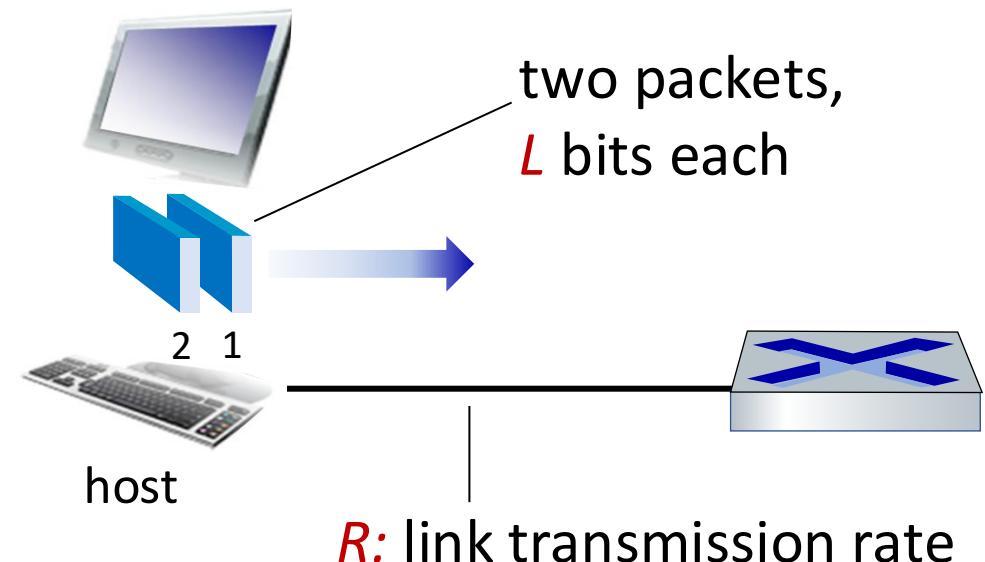


- companies, universities, etc.
- mix of wired, wireless link technologies, connecting a mix of switches and routers (we'll cover differences shortly)
 - Ethernet: wired access at 100Mbps, 1Gbps, 10Gbps
 - WiFi: wireless access points at 11, 54, 450 Mbps

Host: sends *packets* of data

host sending function:

- takes application message
- breaks into smaller chunks, known as *packets*, of length L bits
- transmits packet into access network at *transmission rate R*
 - link transmission rate, aka link *capacity, aka link bandwidth*



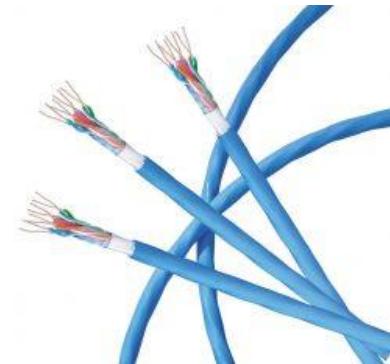
$$\text{packet transmission delay} = \frac{\text{time needed to transmit } L\text{-bit packet into link}}{R \text{ (bits/sec)}}$$

Links: physical media

- **bit**: propagates between transmitter/receiver pairs
- **physical link**: what lies between transmitter & receiver
- **guided media**:
 - signals propagate in solid media: copper, fiber, coax
- **unguided media**:
 - signals propagate freely, e.g., radio

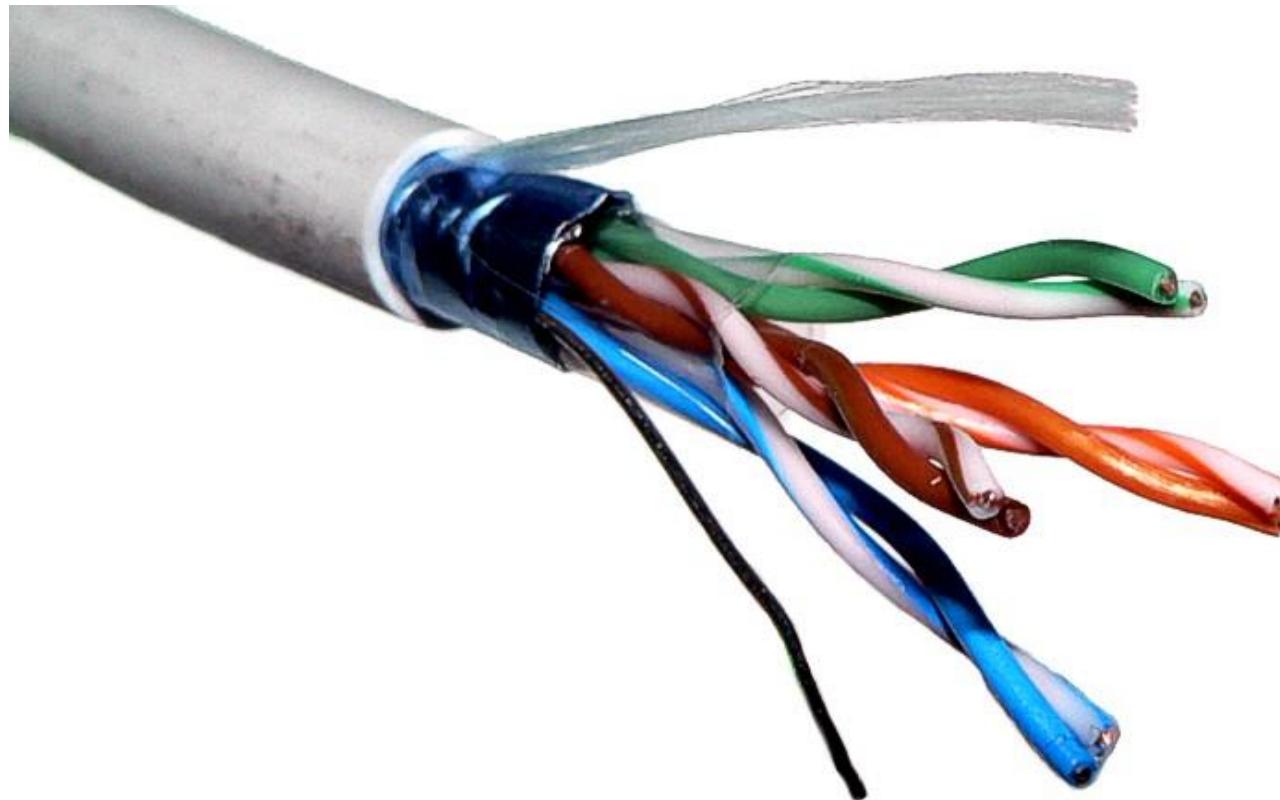
Twisted pair (TP)

- two insulated copper wires
 - Category 5: 100 Mbps, 1 Gbps Ethernet
 - Category 6: 10Gbps Ethernet



Cat 5 or Cat 6: Zoomed In

- Each cable has four twisted pairs.



Links: physical media

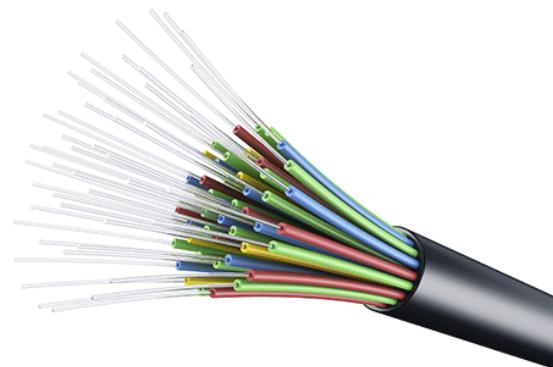
Coaxial cable:

- two concentric copper conductors
- bidirectional
- broadband:
 - multiple frequency channels on cable
 - 100's Mbps per channel

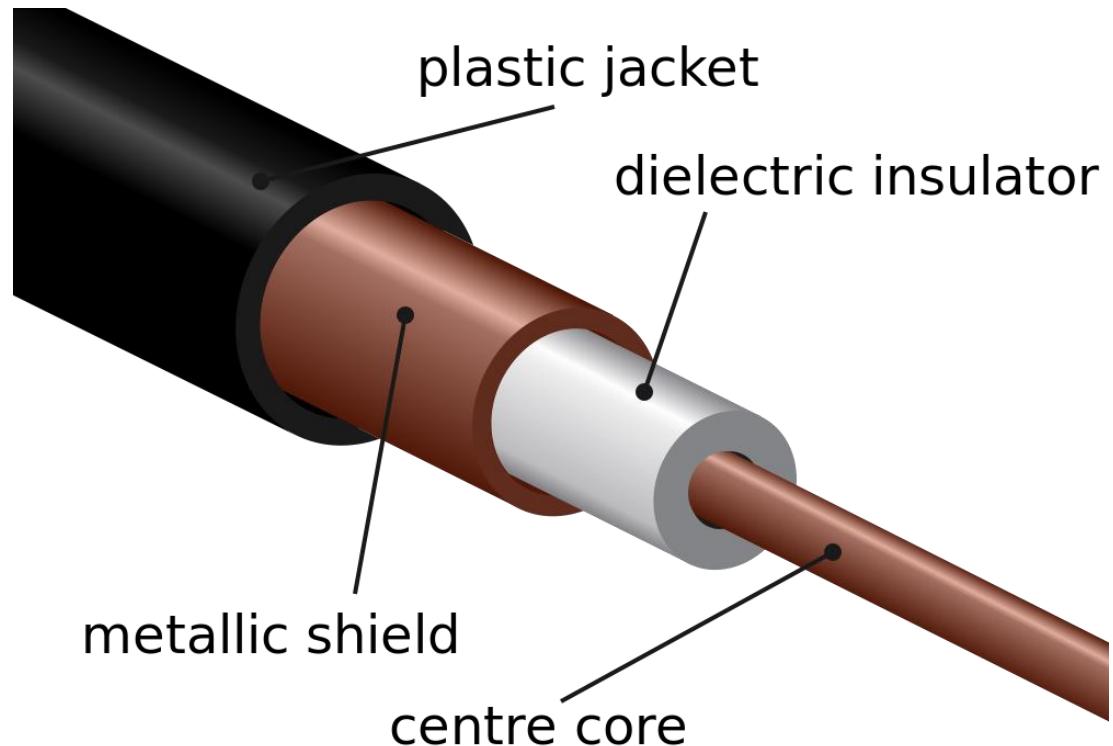


Fiber optic cable:

- glass fiber carrying light pulses, each pulse a bit
- high-speed operation:
 - high-speed point-to-point transmission (10's-100's Gbps)
- low error rate:
 - repeaters spaced far apart
 - immune to electromagnetic noise



Coaxial Cable



Links: physical media

Wireless radio

- signal carried in electromagnetic spectrum
- no physical “wire”
- **broadcast** and “half-duplex”
(sender to receiver)
- propagation environment effects:
 - reflection
 - obstruction by objects
 - interference

Radio link types:

- terrestrial microwave
 - up to 45 Mbps channels
- Wireless LAN (WiFi)
 - Up to 100's Mbps
- wide-area (e.g., cellular)
 - 4G cellular: ~ 10's Mbps
- satellite
 - up to 45 Mbps per channel
 - 270 msec end-end delay
 - geosynchronous versus low-earth-orbit

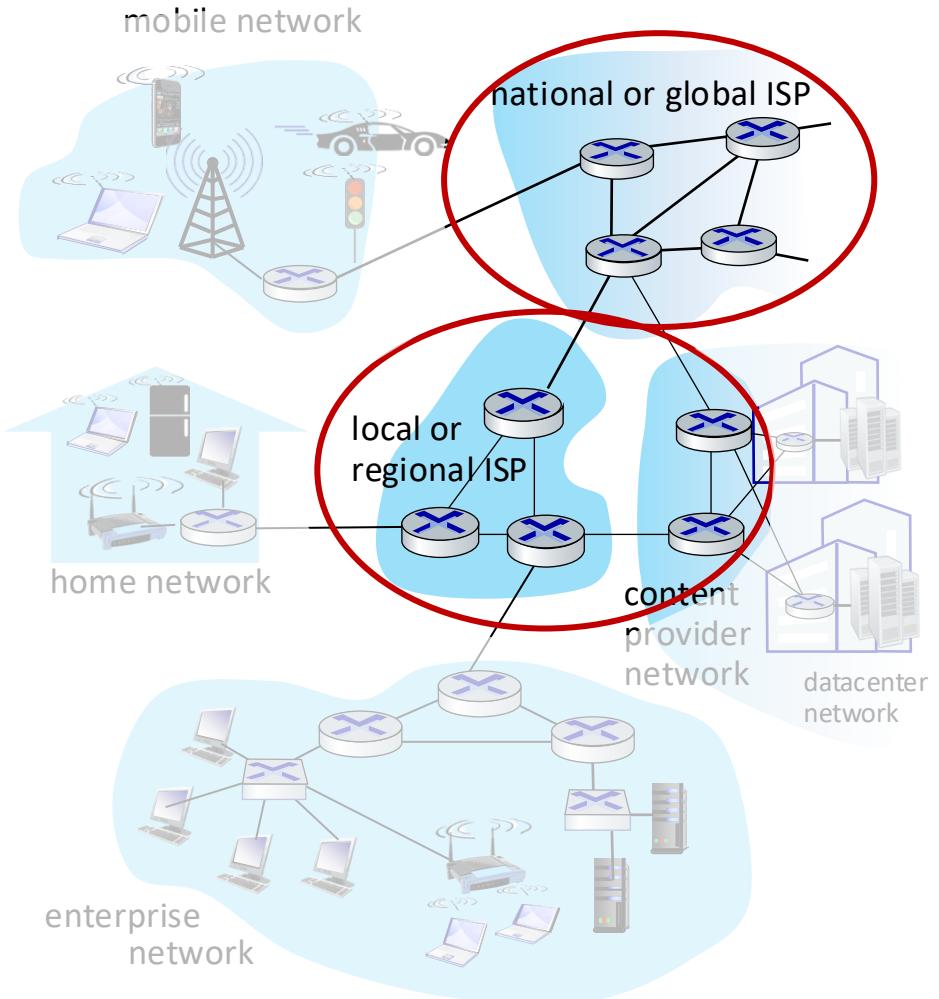
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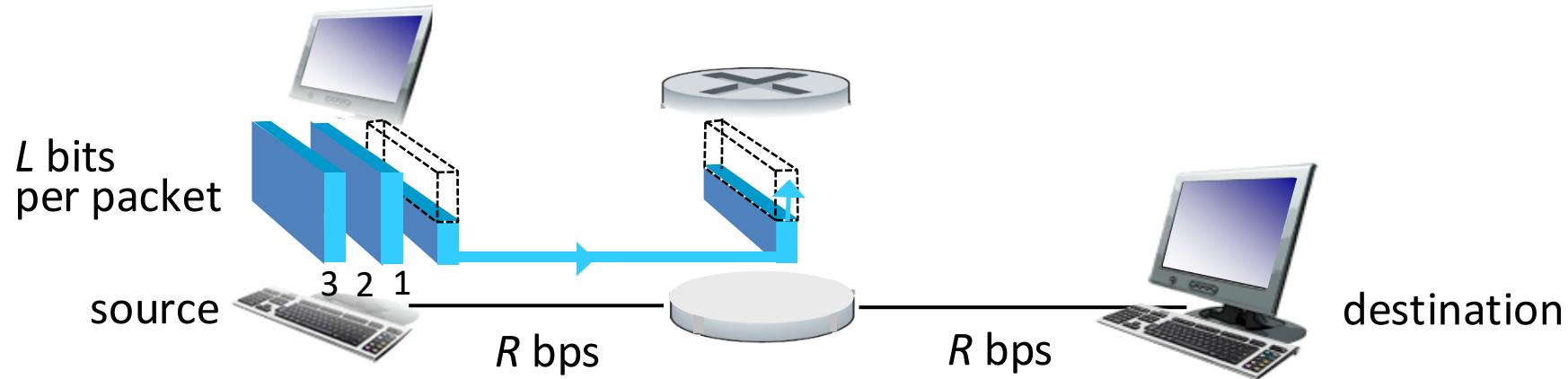


The network core

- mesh of interconnected routers
- **packet-switching:** hosts break application-layer messages into *packets*
 - forward packets from one router to the next, across links on path from source to destination
 - each packet transmitted at full link capacity



Packet-switching: store-and-forward

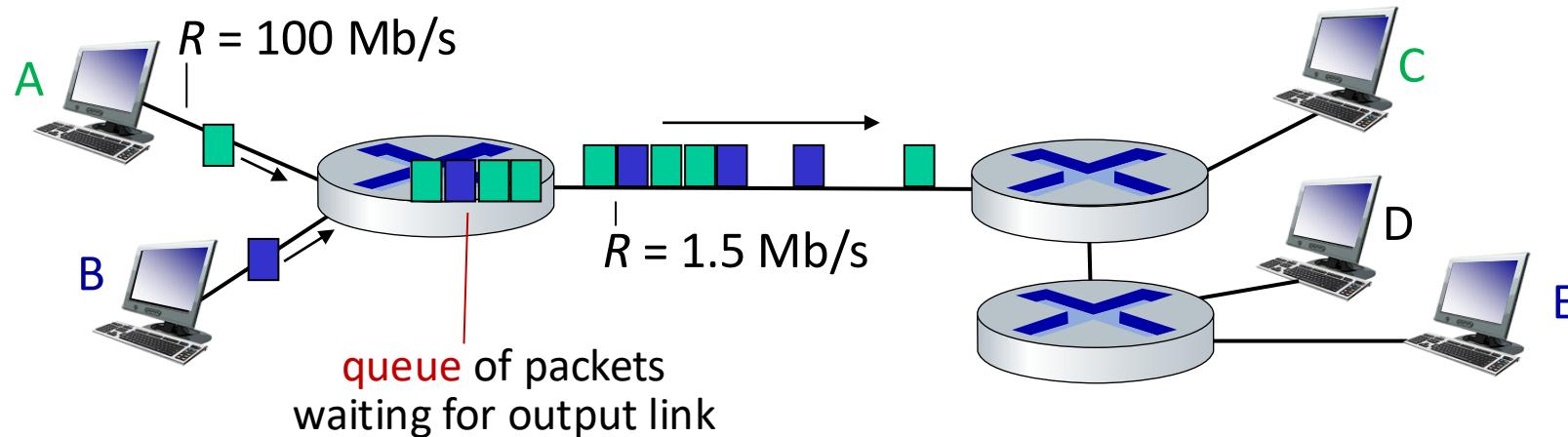


- **Transmission delay:** takes L/R seconds to transmit (push out) L -bit packet into link at R bps
- **Store and forward:** entire packet must arrive at router before it can be transmitted on next link
- **End-end delay:** $2L/R$ (above), assuming zero propagation delay (more on delay shortly)

One-hop numerical example:

- $L = 10$ Kbits
- $R = 100$ Mbps
- one-hop transmission delay = 0.1 msec

Packet-switching: queueing delay, loss



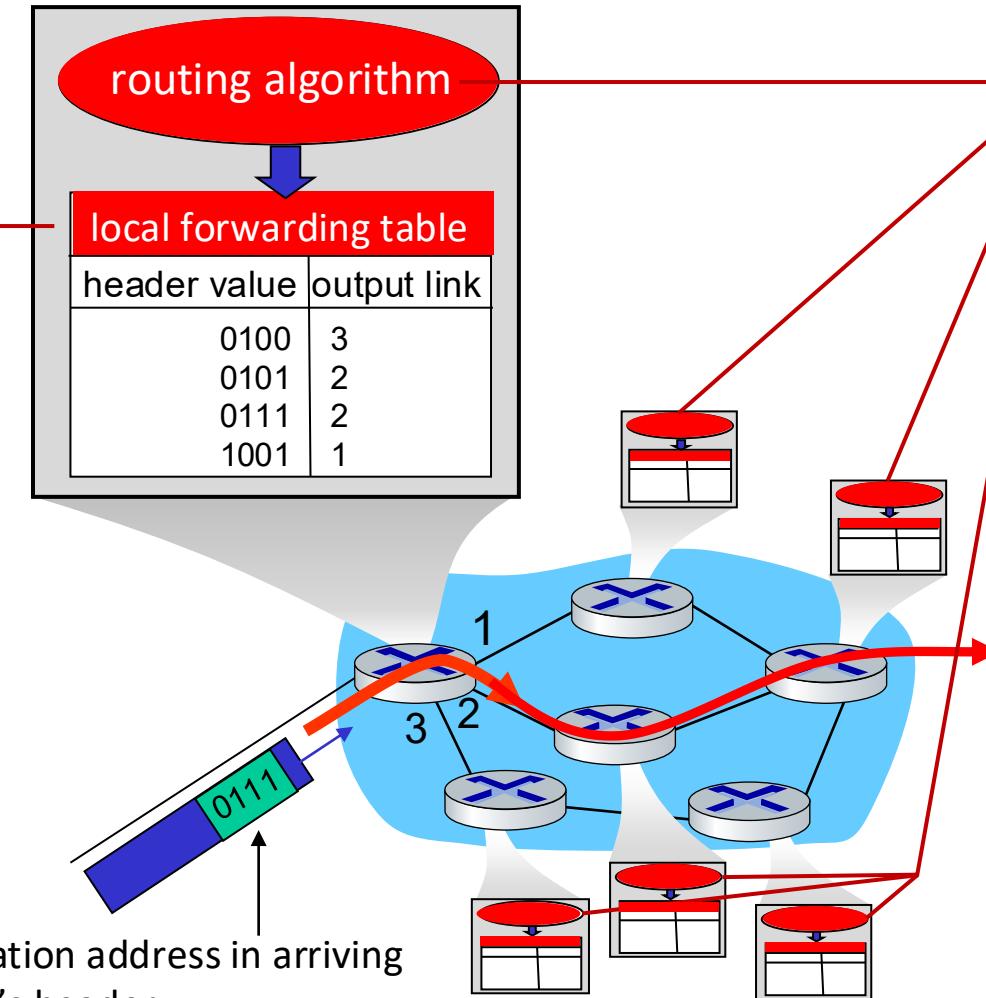
Packet queuing and loss: if arrival rate (in bps) to link exceeds transmission rate (bps) of link for a period of time:

- packets will queue, waiting to be transmitted on output link
- packets can be dropped (lost) if memory (buffer) in router fills up

Two key network-core functions

Forwarding: _____

- *local* action:
move arriving
packets from
router's input link
to appropriate
router output link



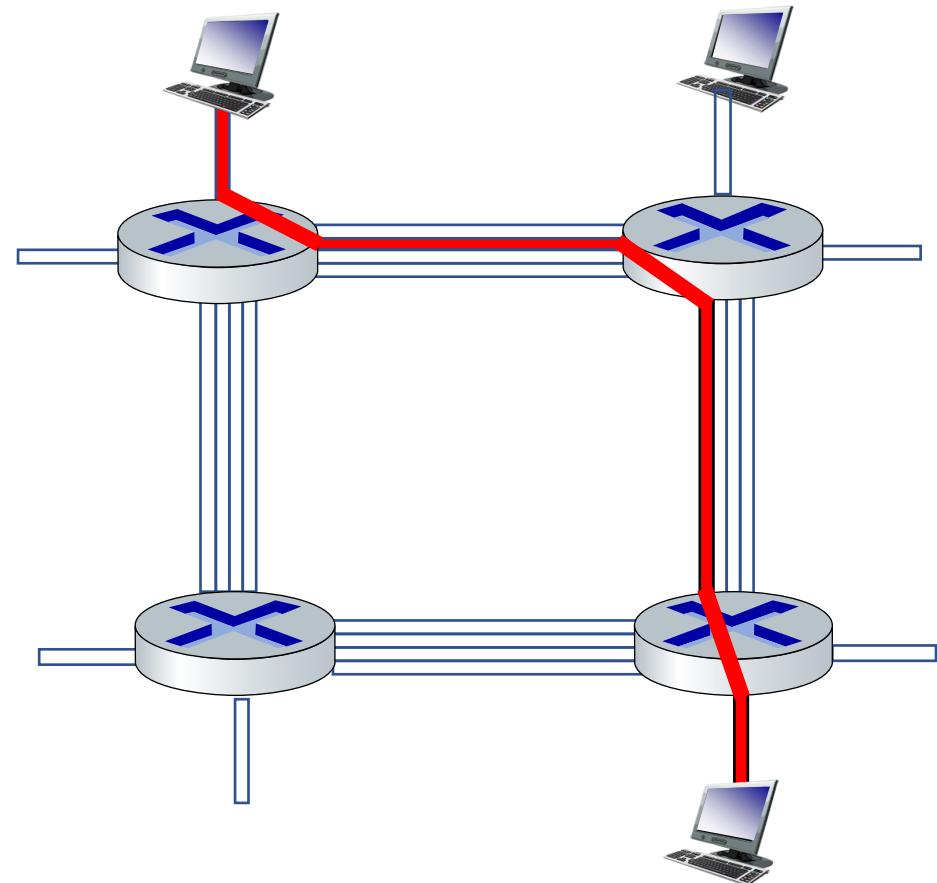
Routing:

- *global* action:
determine source-destination paths taken by packets
- routing algorithms

Alternative to packet switching: circuit switching

end-end resources allocated to,
reserved for “call” between source
and destination

- in diagram, each link has four circuits.
 - call gets 2nd circuit in top link and 1st circuit in right link.
- dedicated resources: no sharing
 - circuit-like (guaranteed) performance
- circuit segment idle if not used by call (**no sharing**)
- commonly used in traditional telephone networks
- A call is rejected/blocked when no circuit is available.



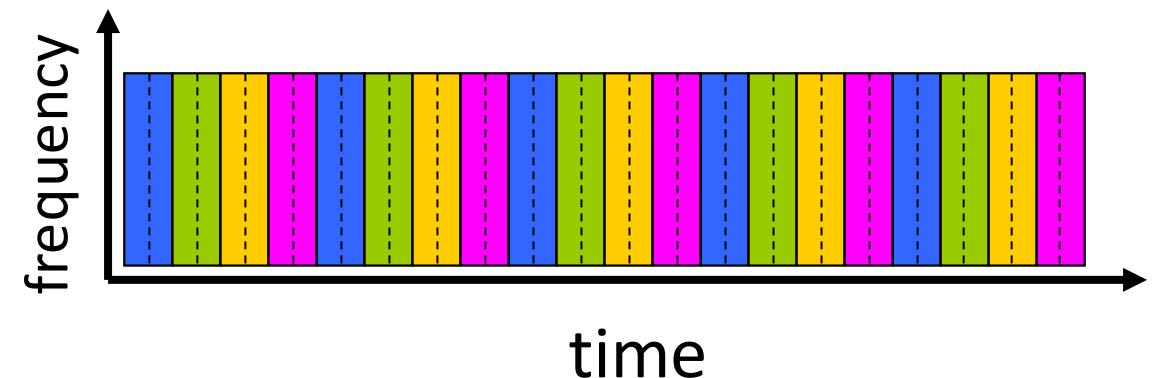
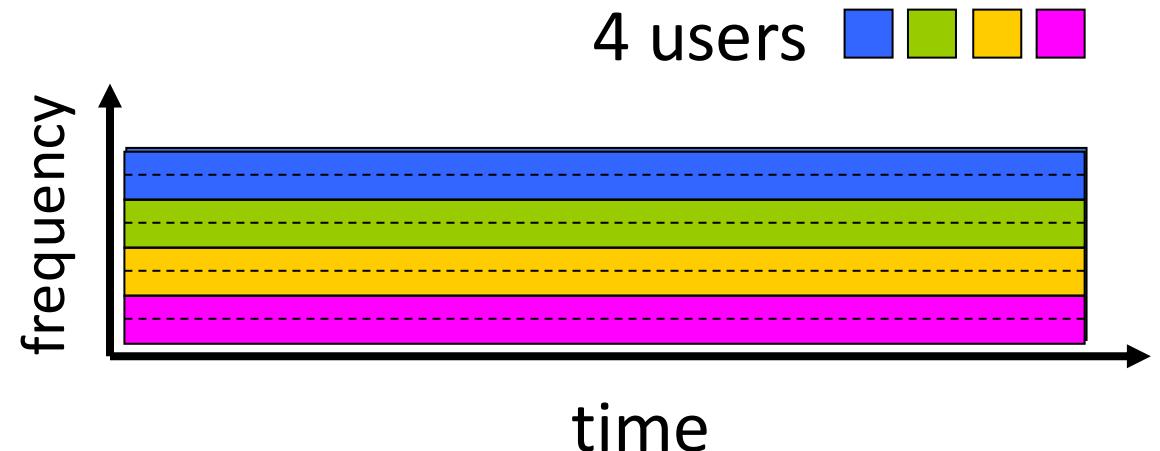
Circuit switching: FDM and TDM

Frequency Division Multiplexing (FDM)

- optical, electromagnetic frequencies divided into (narrow) frequency bands
- each call allocated its own band, can transmit at max rate of that narrow band

Time Division Multiplexing (TDM)

- time divided into slots
- each call allocated periodic slot(s), can transmit at maximum rate of (wider) frequency band, but only during its time slot(s)

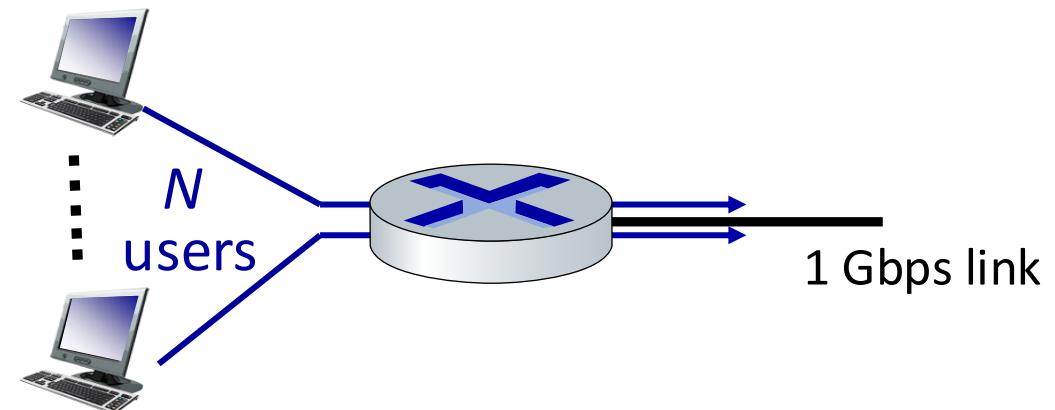


Packet switching versus circuit switching

packet switching allows more users to use network!

Example:

- 1 Gb/s link
- each user:
 - 100 Mb/s when “active”
 - active 10% of time
- *circuit-switching*: 10 users
- *packet switching*: with 35 users, probability > 10 active at same time is less than .0004 *



Q: how did we get value 0.0004?

Q: what happens if > 35 users ?

* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive

Packet switching versus circuit switching

Is packet switching a “slam dunk winner”?

- great for “bursty” data – sometimes has data to send, but at other times not
 - resource sharing
 - simpler, no call setup
- **excessive congestion possible:** packet delay and loss due to buffer overflow
 - protocols needed for reliable data transfer, congestion control
- **Q: How to provide circuit-like behavior?**
 - bandwidth guarantees traditionally used for audio/video applications

Q: human analogies of reserved resources (circuit switching) versus on-demand allocation (packet switching)?

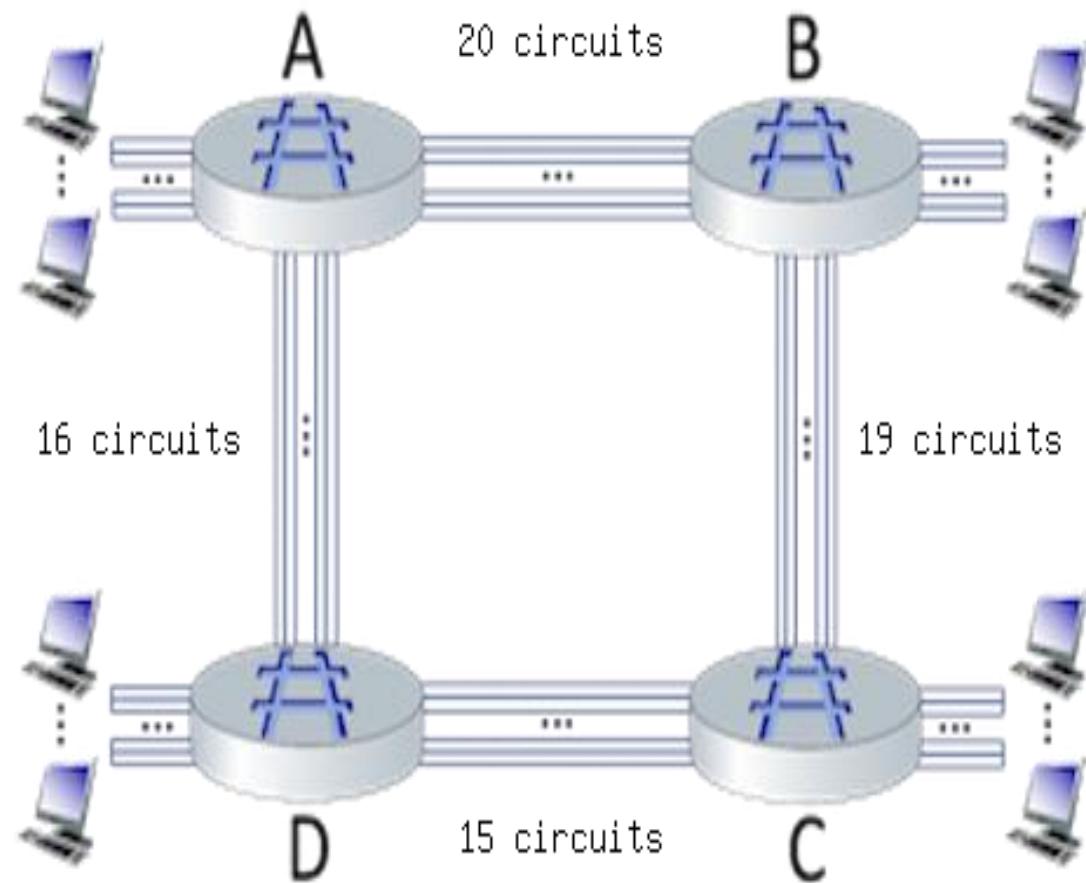
Knowledge Check

Which of the characteristics below are associated with the technique of *packet switching*?

1. Congestion loss and variable end-end delays are possible with this technique.
2. Reserves resources needed for a call from source to destination.
3. Resources are used on demand, not reserved in advance.
4. Frequency Division Multiplexing (FDM) and Time Division Multiplexing (TDM) are two approaches for implementing this technique.
5. This technique was the basis for the telephone call switching during the 20th century and into the beginning of this current century.
6. Data may be queued before being transmitted due to other user's data that's also queueing for transmission.
7. This technique is used in the Internet.

Knowledge Check

Circuit Switching: What is the maximum number of connections that can be ongoing in the network at any one time?

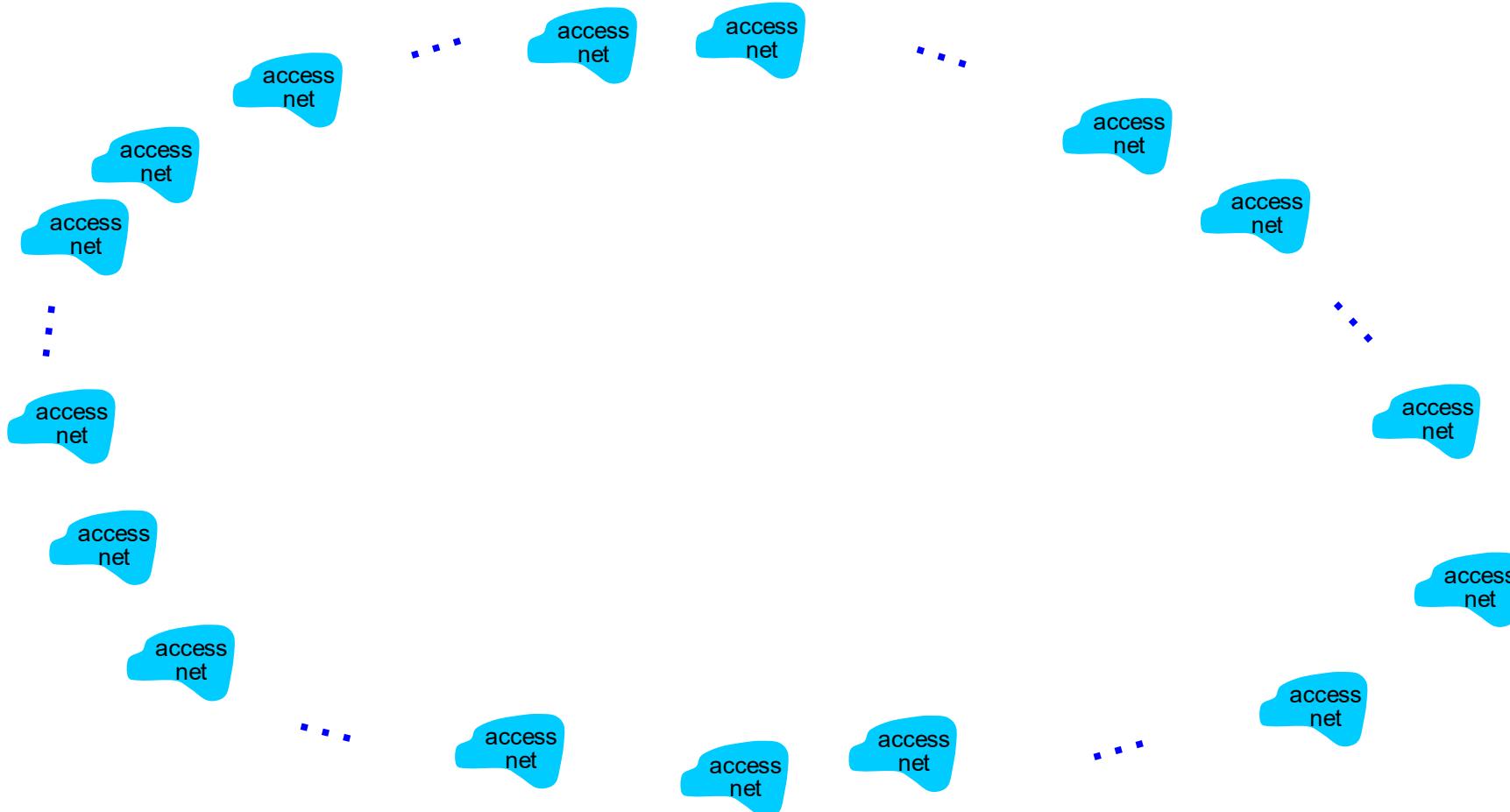


Internet structure: a “network of networks”

- Hosts connect to Internet via **access** Internet Service Providers (ISPs)
 - residential, enterprise (company, university, commercial) ISPs
- Access ISPs in turn must be interconnected
 - so that any two hosts can send packets to each other
- Resulting network of networks is very complex
 - evolution was driven by **economics** and **national policies**
- Let's take a stepwise approach to describe current Internet structure

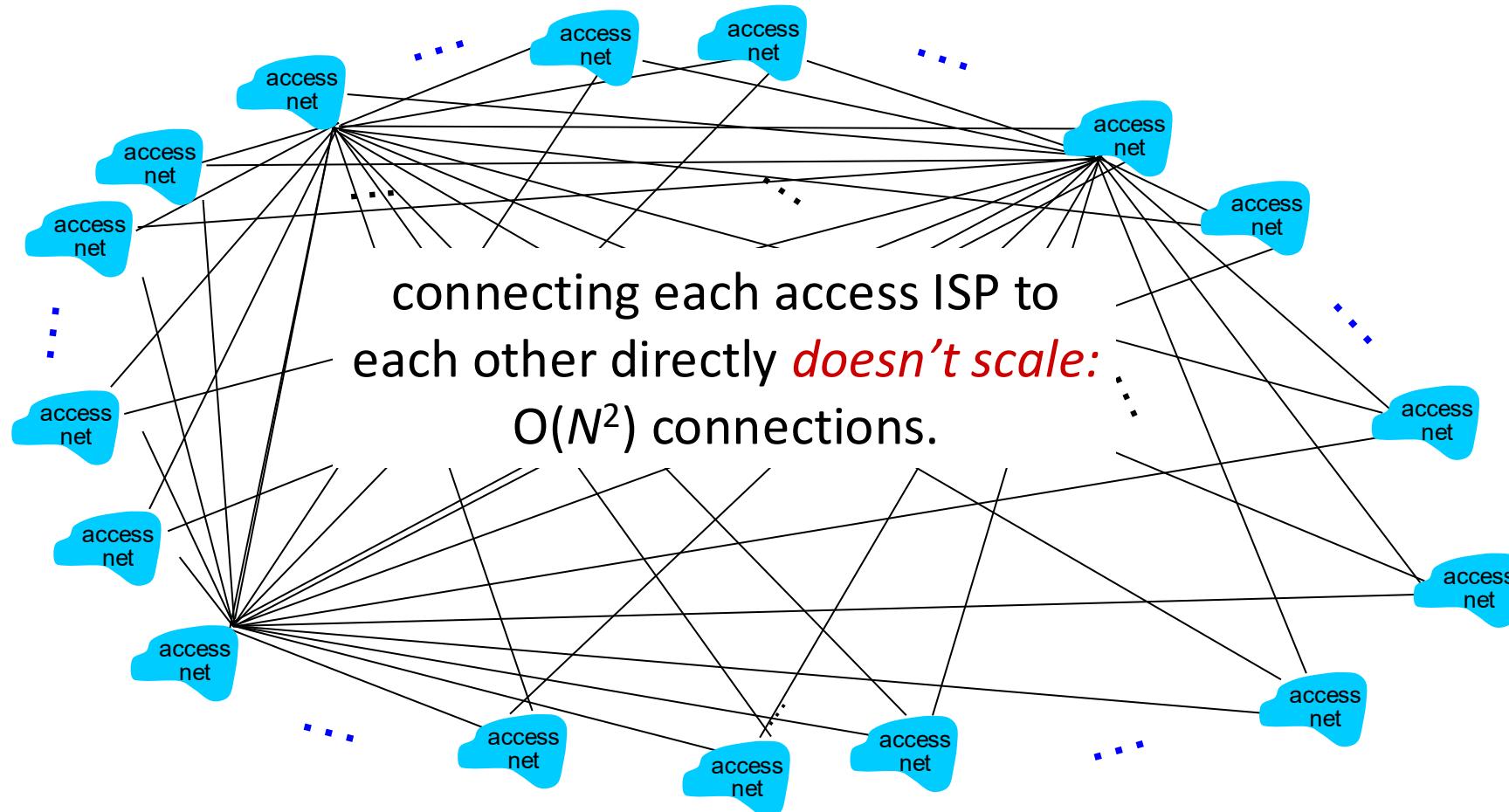
Internet structure: a “network of networks”

Question: given *millions* of access ISPs, how to connect them together?



Internet structure: a “network of networks”

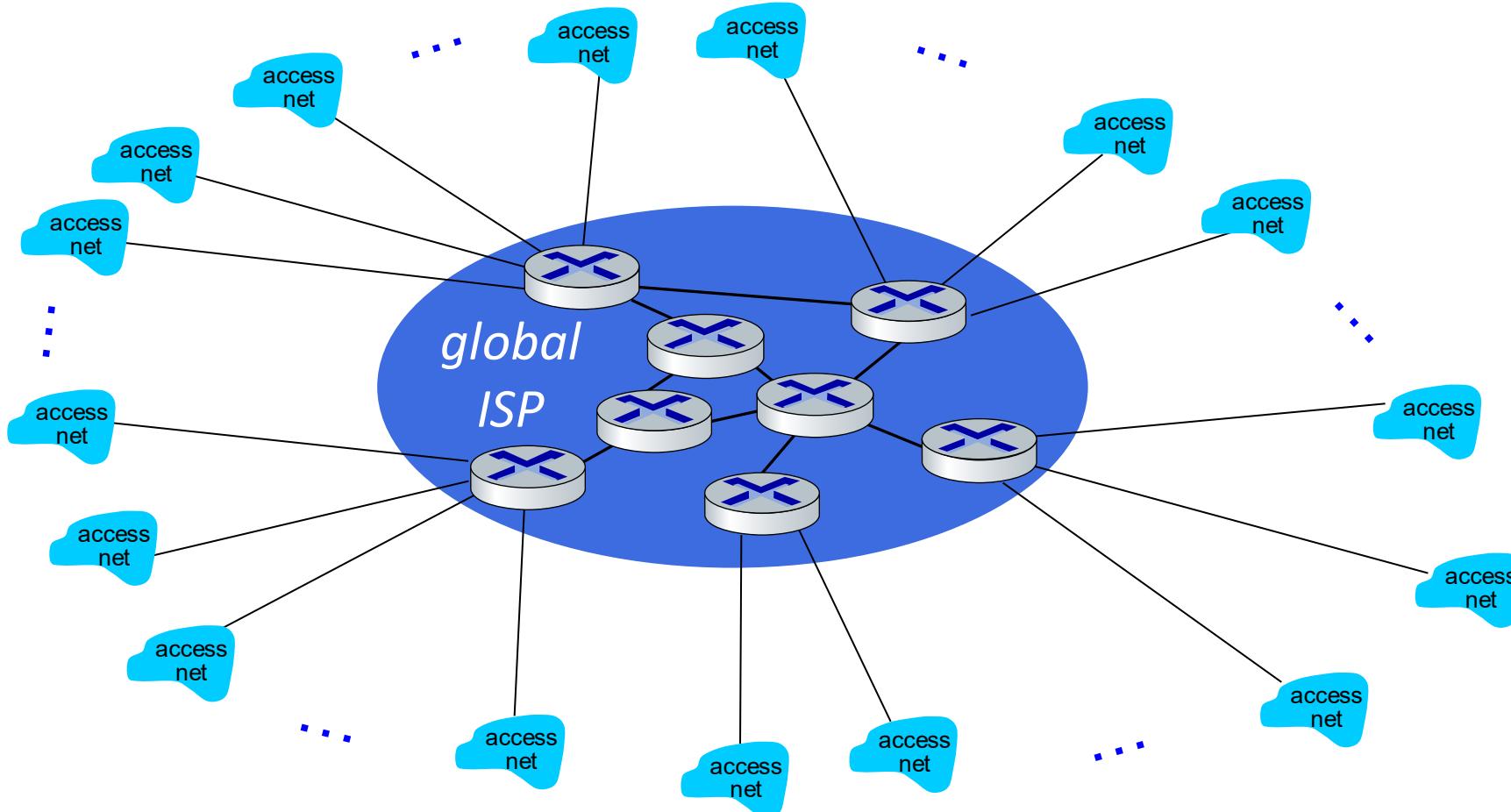
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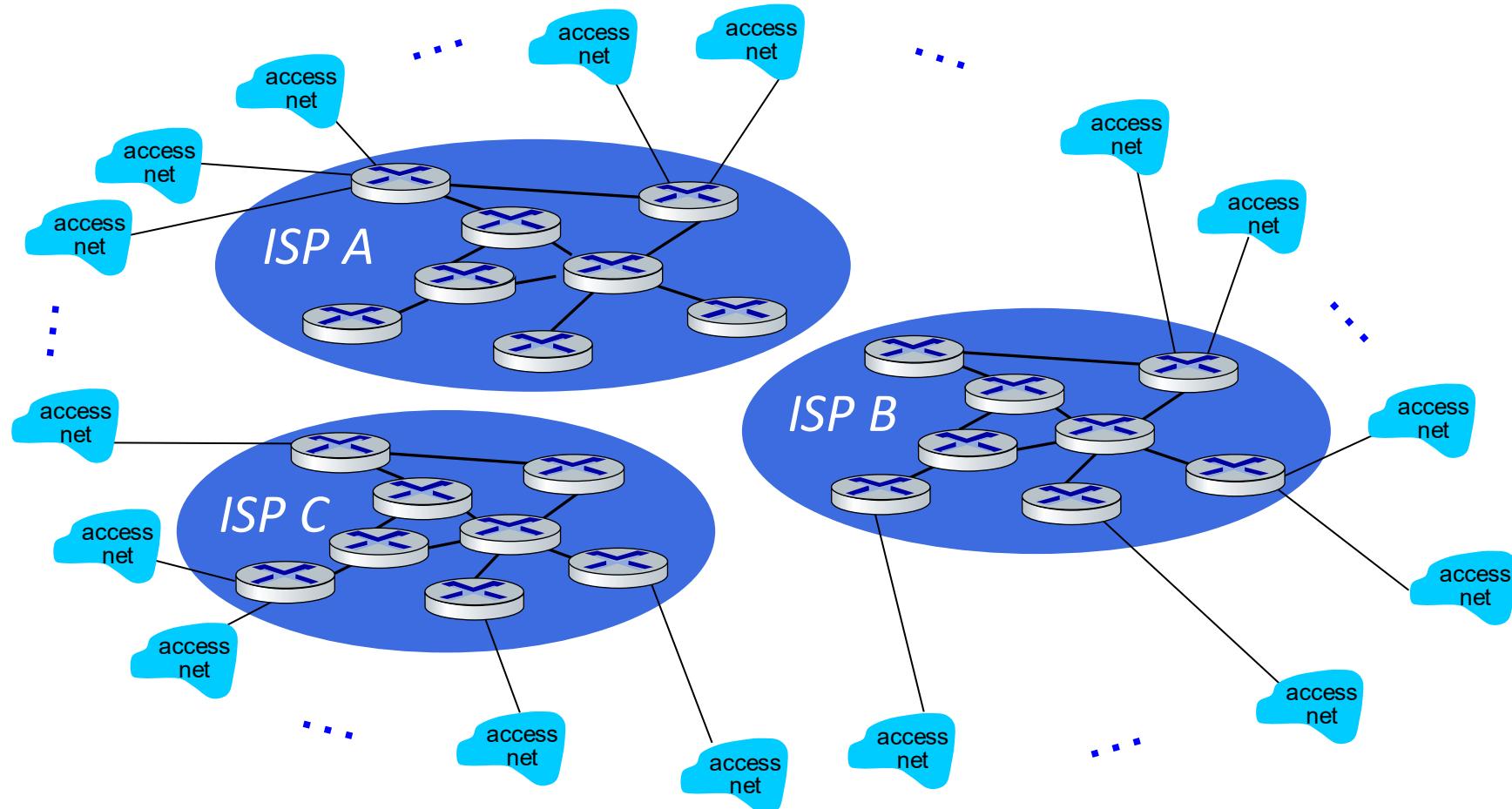
Option: connect each access ISP to one global transit ISP?

Customer and provider ISPs have economic agreement.



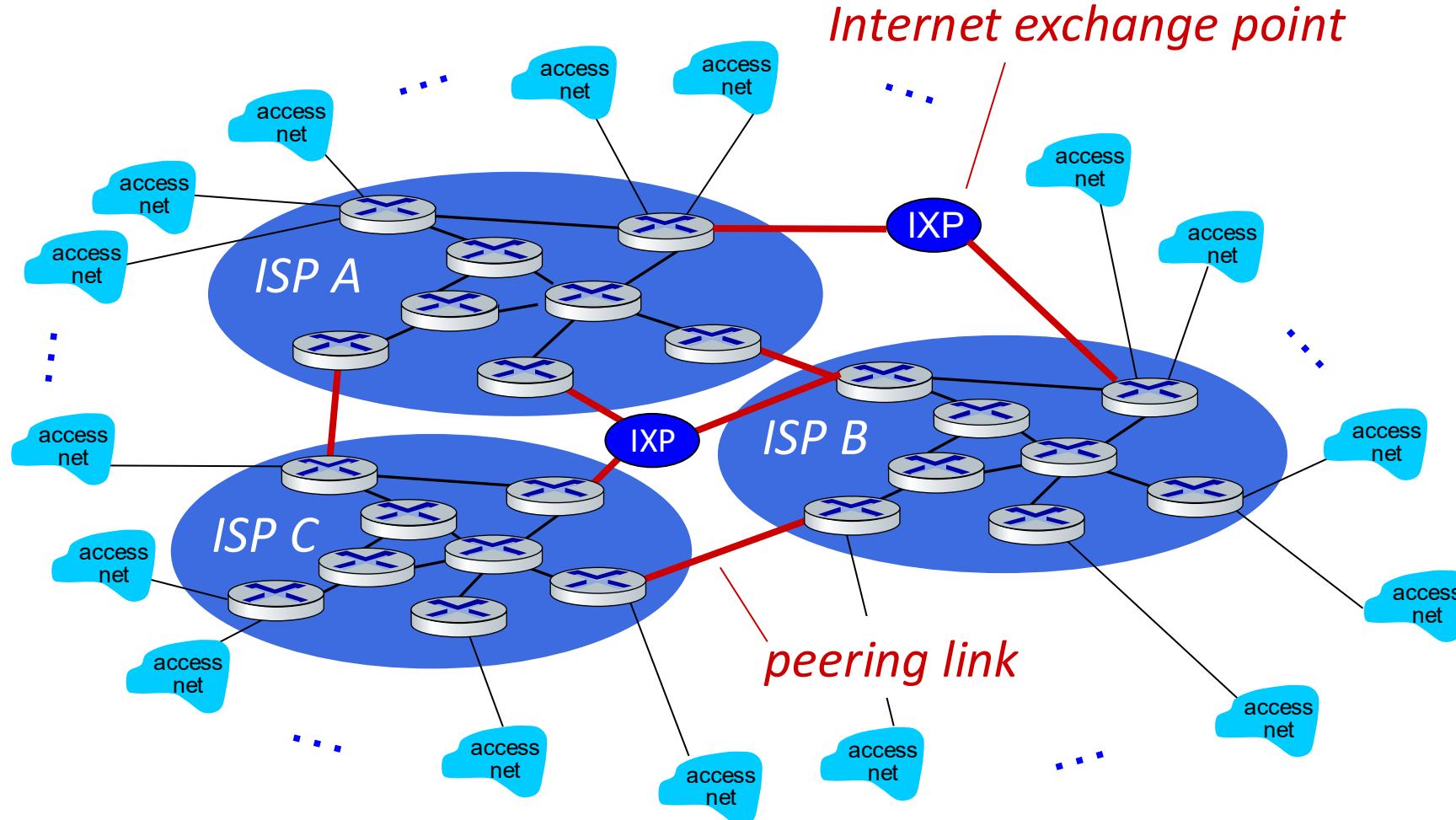
Internet structure: a “network of networks”

But if one global ISP is viable business, there will be competitors



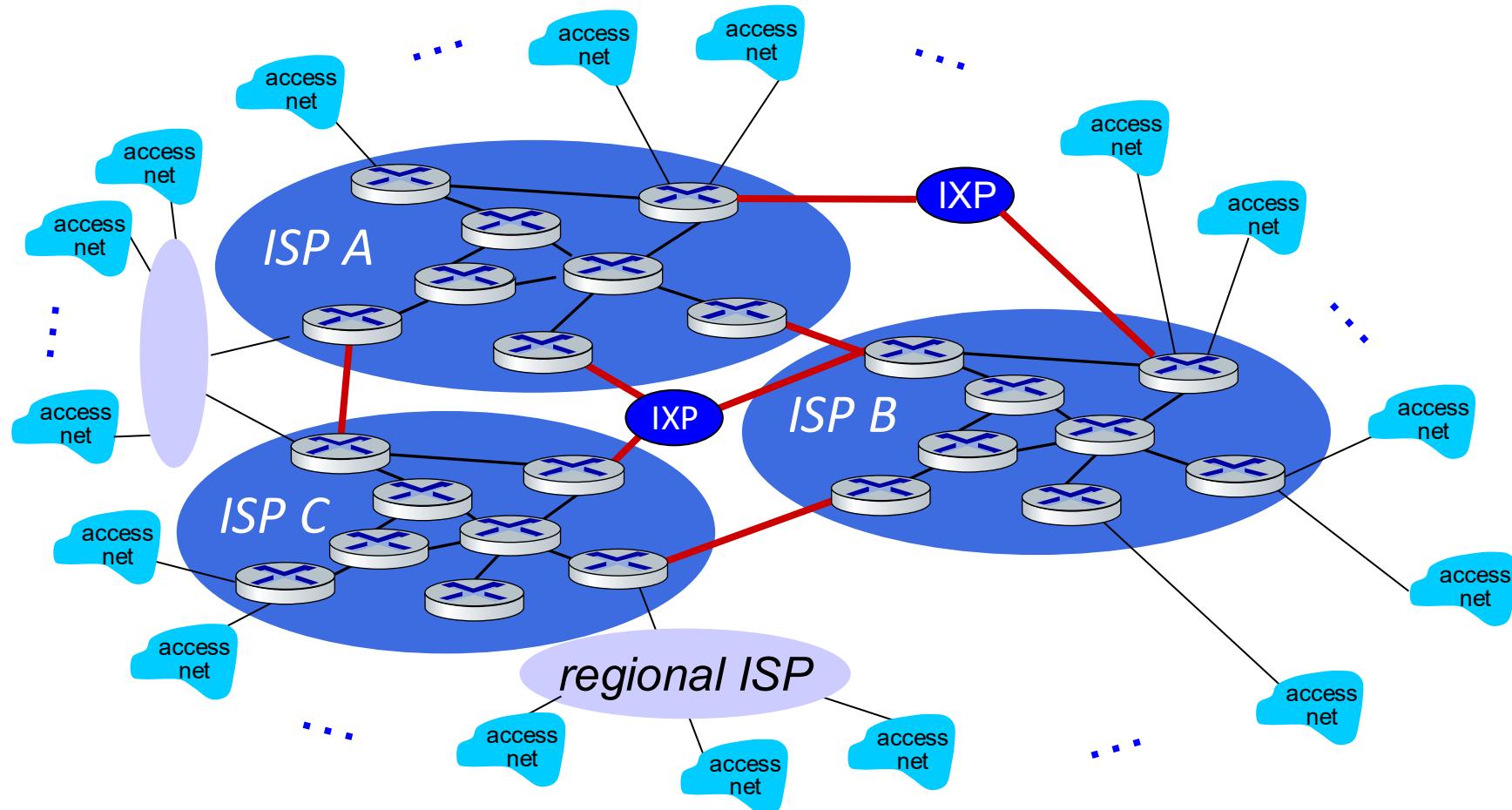
Internet structure: a “network of networks”

But if one global ISP is viable business, there will be competitors ... who will want to be connected



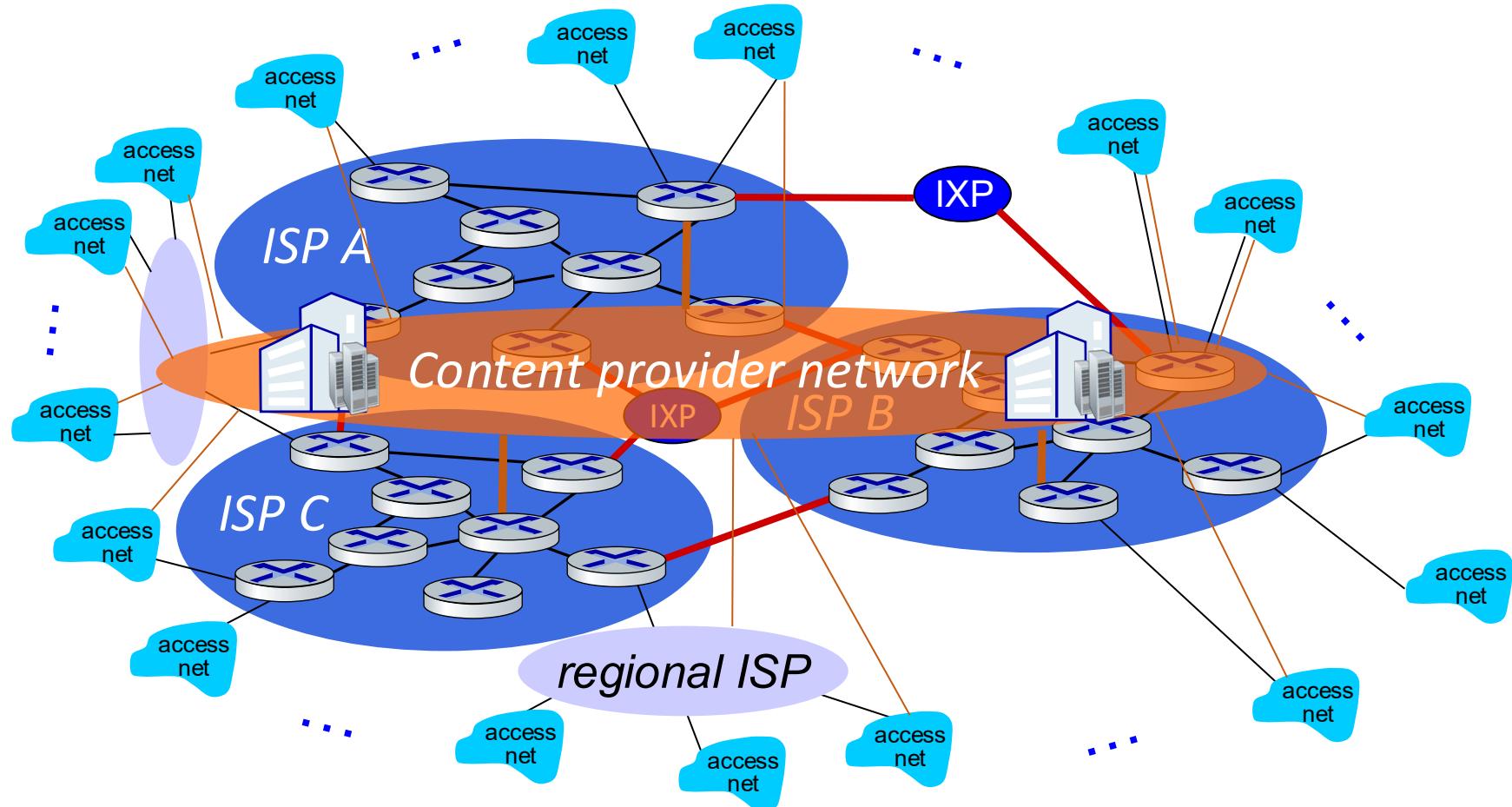
Internet structure: a “network of networks”

... and regional networks may arise to connect access nets to ISPs

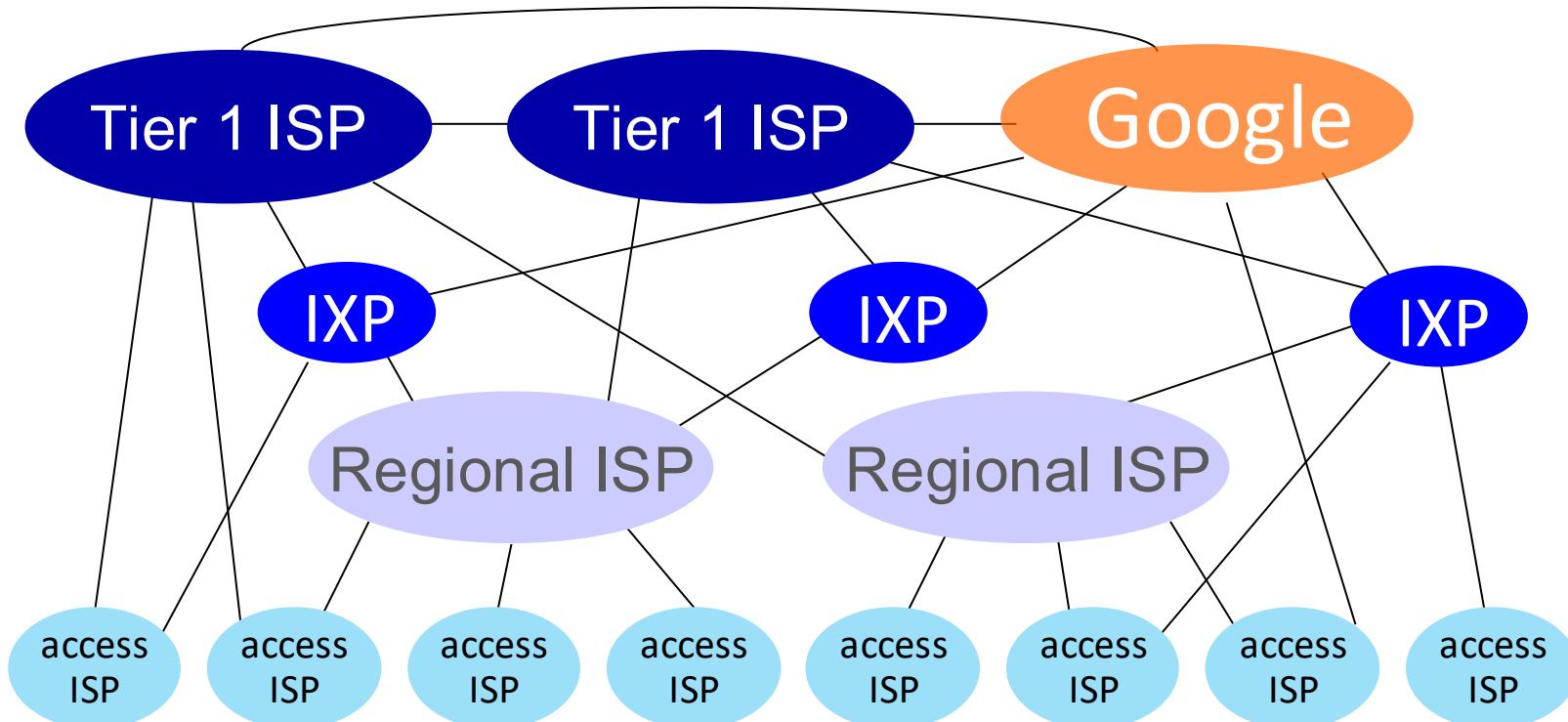


Internet structure: a “network of networks”

... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users



Internet structure: a “network of networks”



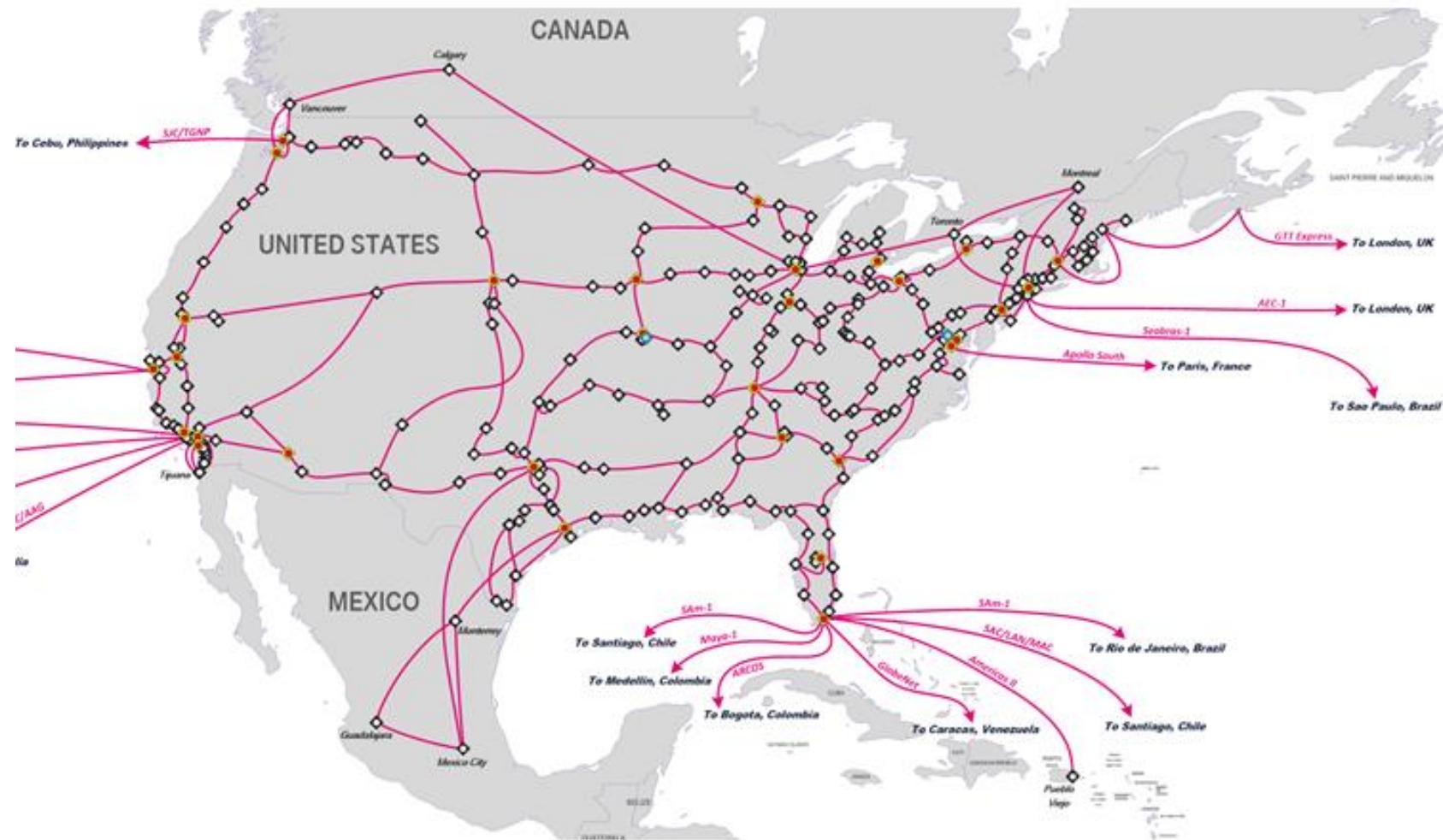
At “center”: small # of well-connected large networks

- **“tier-1” commercial ISPs** (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
- **content provider networks** (e.g., Google, Facebook): private network that connects its data centers to Internet, often bypassing tier-1, regional ISPs

Tier-1 ISP Network map: Sprint (2019)



Sprint Network (now T-Mobile) Links



Chapter 1: roadmap

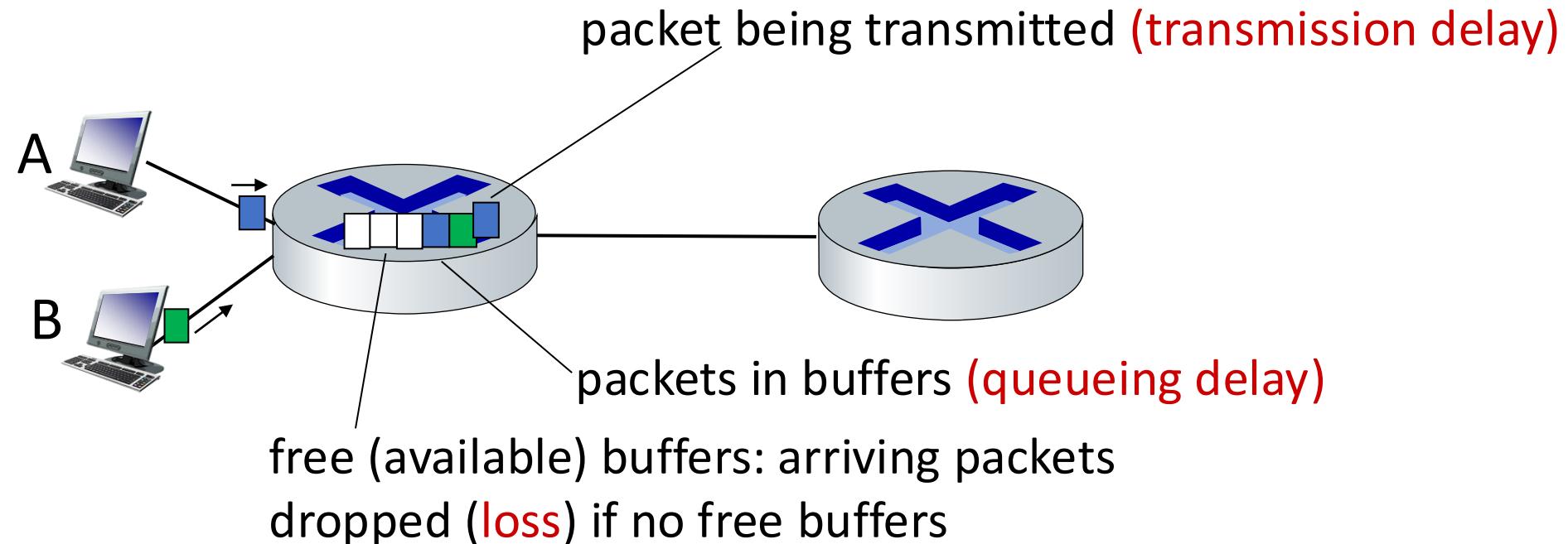
- What *is* the Internet?
- What *is* a protocol?
- Network edge: hosts, access network, physical media
- Network core: packet/circuit switching, internet structure
- **Performance: loss, delay, throughput**
- Security
- Protocol layers, service models
- History



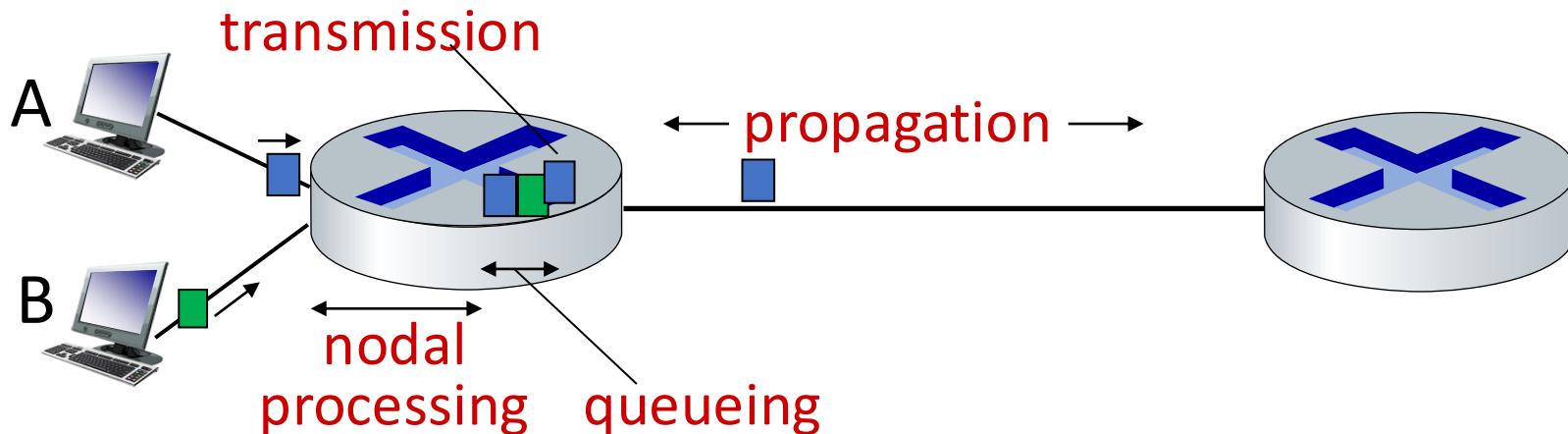
How do packet loss and delay occur?

packets *queue* in router buffers

- packets queue, wait for turn
- arrival rate to link (temporarily) exceeds output link capacity: packet loss



Packet delay: four sources



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

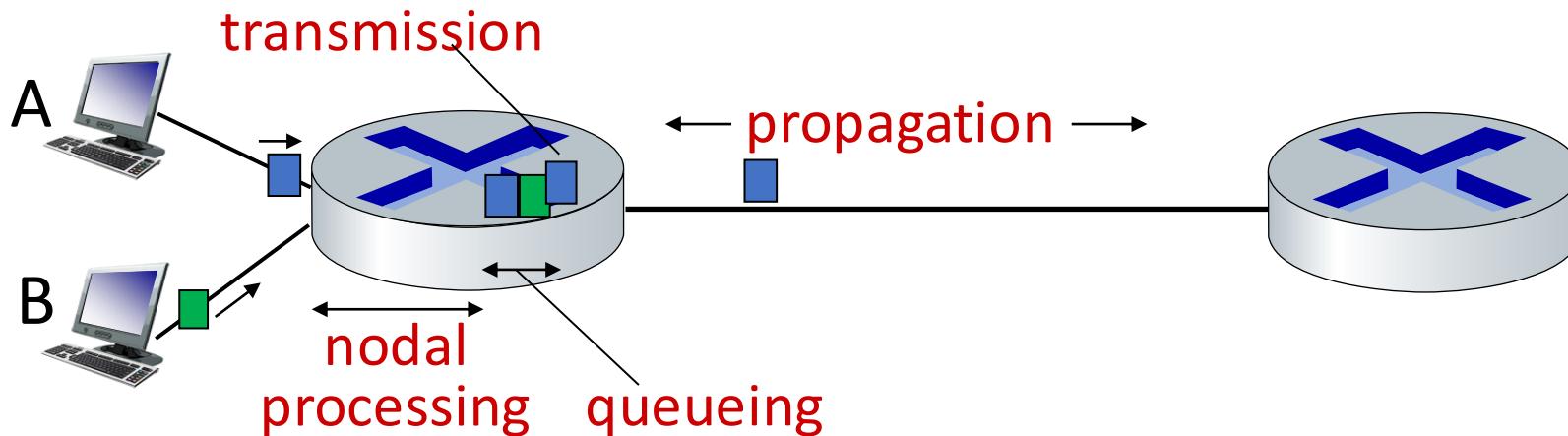
d_{proc} : nodal processing

- check bit errors
- determine output link
- typically < msec

d_{queue} : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

Packet delay: four sources



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{trans} : transmission delay:

- L : packet length (bits)
- R : link *transmission rate (bps)*
- $d_{\text{trans}} = L/R$

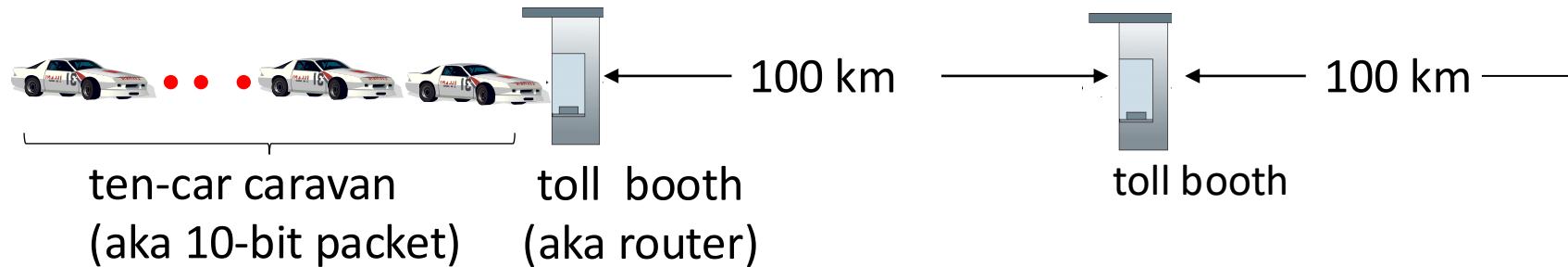
d_{trans} and d_{prop}
very different

d_{prop} : propagation delay:

- d : length of physical link
- s : propagation speed ($\sim 2 \times 10^8$ m/sec)
- $d_{\text{prop}} = d/s$

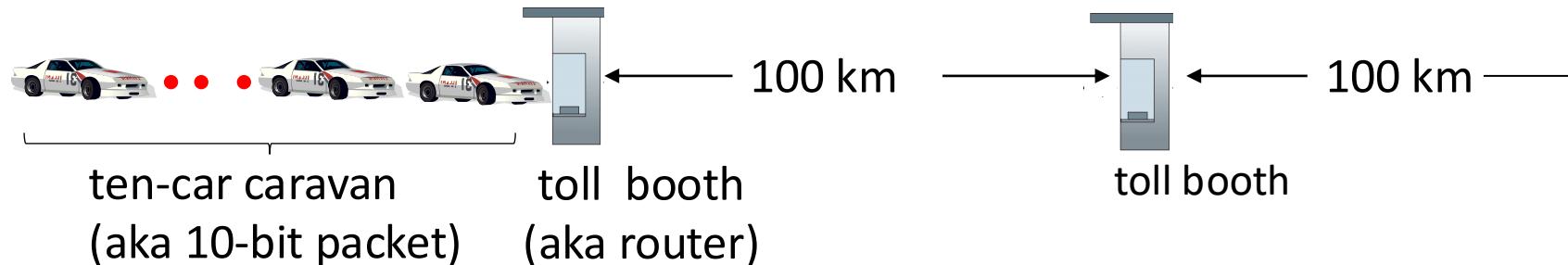
* Check out the online interactive exercises:
http://gaia.cs.umass.edu/kurose_ross

Caravan analogy



- cars “propagate” at 100 km/hr
- toll booth takes 12 sec to service car (bit transmission time)
- car ~ bit; caravan ~ packet
- **Q: How long until caravan is lined up before 2nd toll booth?**
- time to “push” entire caravan through toll booth onto highway = $12 * 10 = 120$ sec
- time for last car to propagate from 1st to 2nd toll both: $100\text{km}/(100\text{km/hr}) = 1$ hr
- **A: 62 minutes**

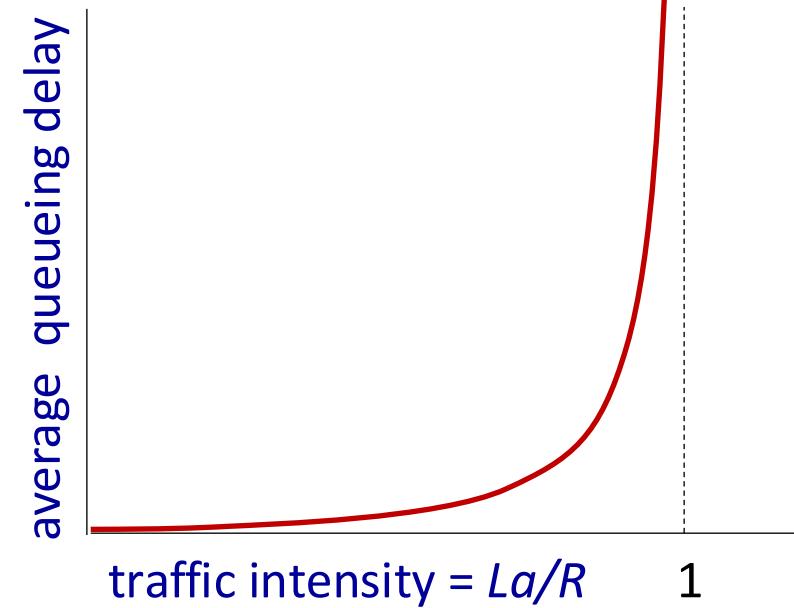
Caravan analogy



- suppose cars now “propagate” at 1000 km/hr
 - and suppose toll booth now takes one min to service a car
 - ***Q: Will cars arrive to 2nd booth before all cars serviced at first booth?***
- A: Yes!** after 7 min, first car arrives at second booth; three cars still at first booth

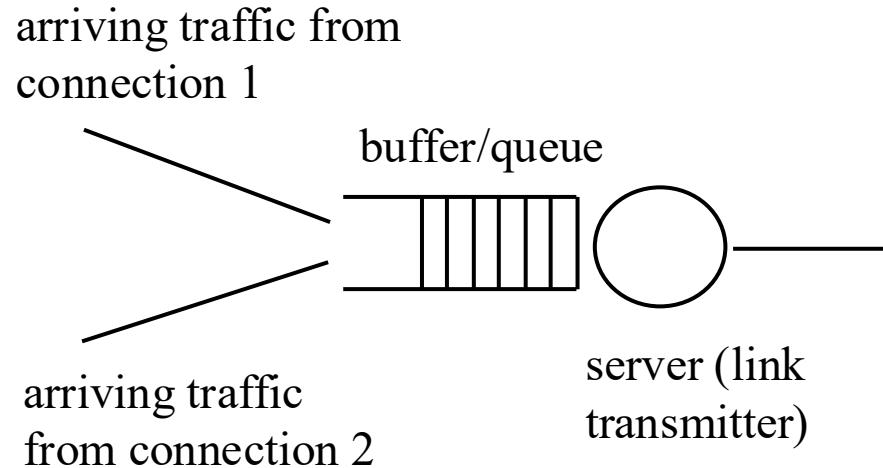
Packet queueing delay (revisited)

- R : link bandwidth (bps)
- L : packet length (bits)
- a : average packet arrival rate
- $La/R \sim 0$: avg. queueing delay small
- $La/R \rightarrow 1$: avg. queueing delay large
- $La/R > 1$: more “work” arriving is more than can be serviced - average delay infinite!



$La/R \rightarrow 1$

Homework Hints

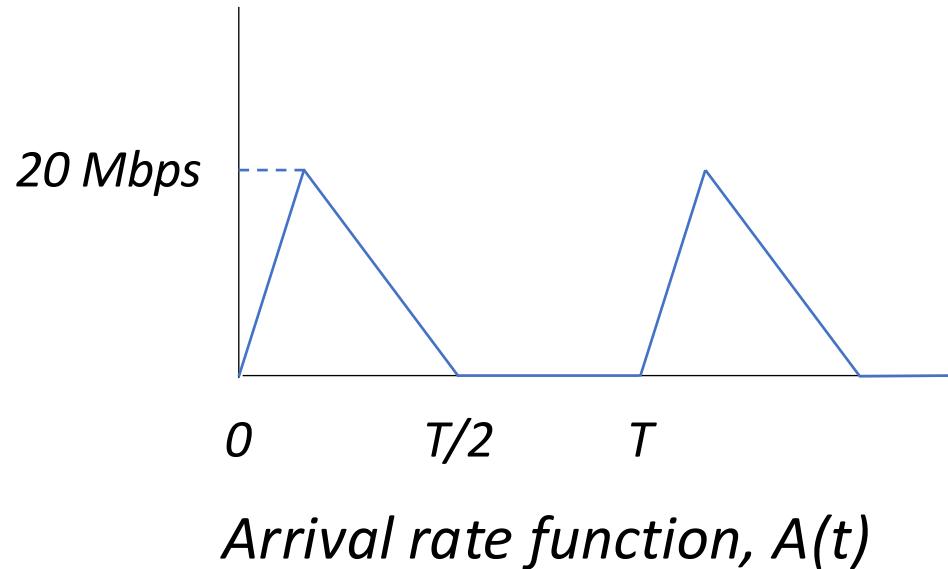


Multiplexer model of a router output link

- *Two types of models for the arrival process*
 - *Point process of packet arrivals, e.g., Poisson arrivals*
 - *Fluid model*
- *For deterministic, periodic traffic, what is the average arrival rate?*
 - *In fluid model, it is the integral of the arrival rate function over one period, divided by period length.*

Homework Hints

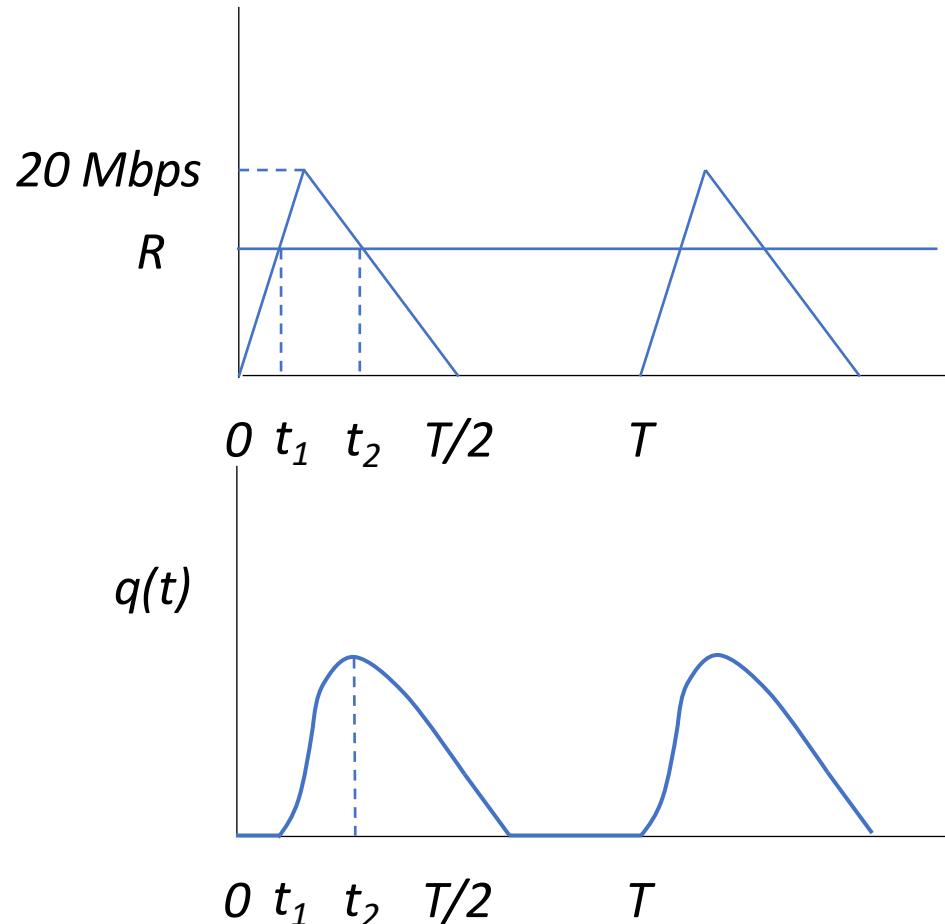
- Example: single source/connection



- Total arrival on $[0, T]$:
$$\int_0^T A(t)dt$$
- Average arrival rate:
$$\frac{1}{T} \int_0^T A(t)dt$$
, which is equal to
$$\frac{\frac{1}{2} \times \frac{T}{2} \times 20}{T} = 5 \text{ Mbps}$$
- If $R < 5 \text{ Mbps}$, the queue is unstable (goes to infinity).
- If $R \geq 5 \text{ Mbps}$, the queue is stable (finite all the time).

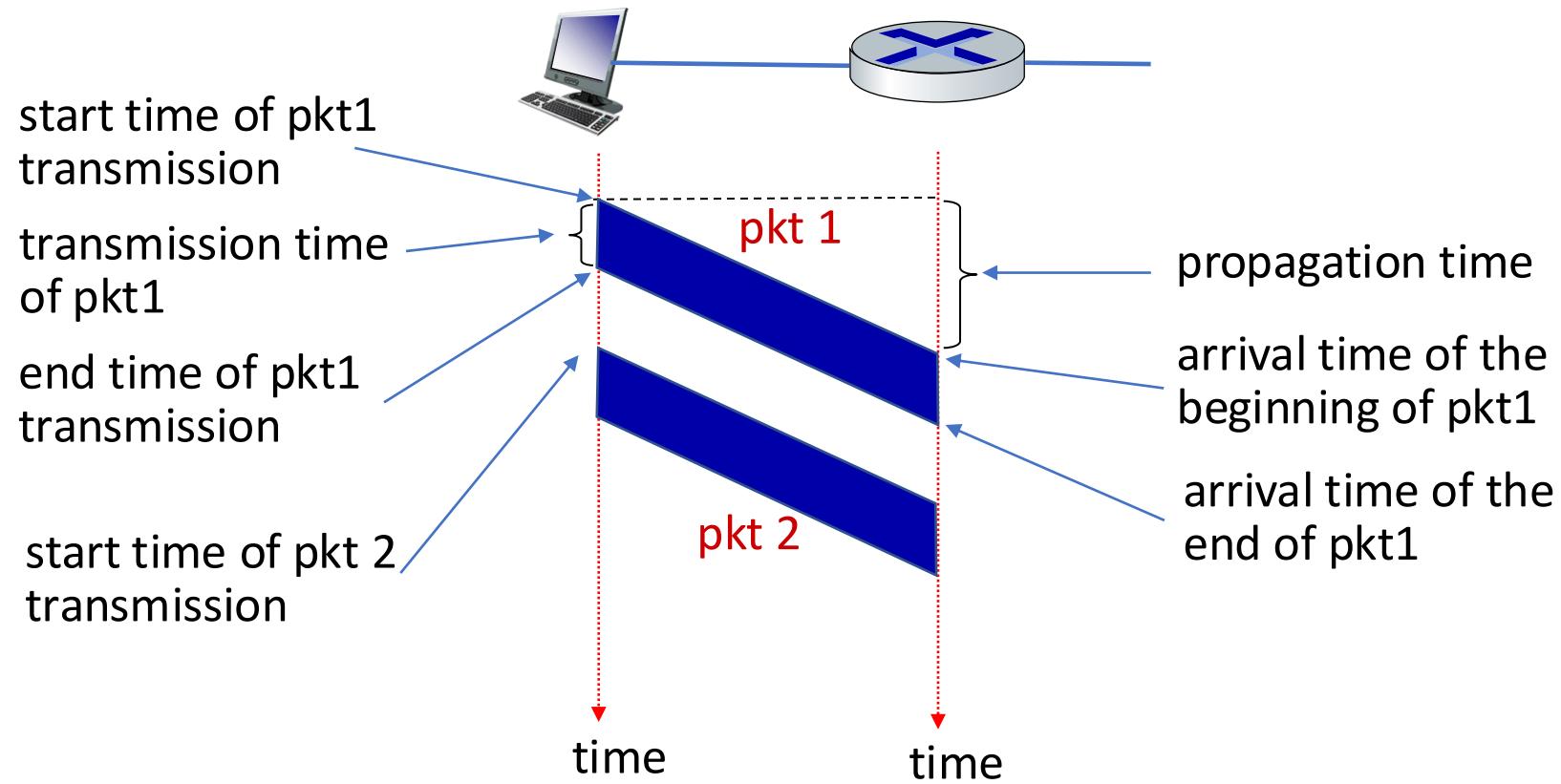
Homework Hints – Queue Process

- A.k.a. buffer occupancy process, $q(t)$



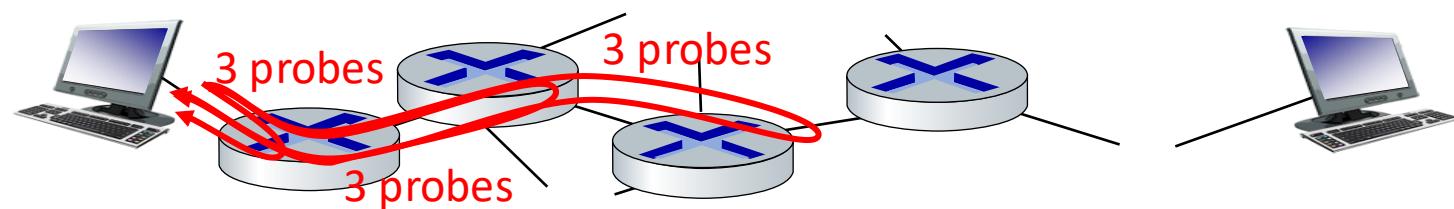
- Assume $q(0)=0$.
- On $[0, t_1]$, no queue buildup
- If $A(t) > R$, $q(t)$ increases at the rate $A(t)-R$; e.g., on (t_1, t_2) .
- How much increase on an interval $[a,b]$ where $A(t)>R$: $\int_a^b (A(t) - R) dt$
- If $A(t) < R$, then $q(t)$ decreases at the rate $R-A(t)$, as long as queue is not empty.

Homework Hints – Timing Diagram



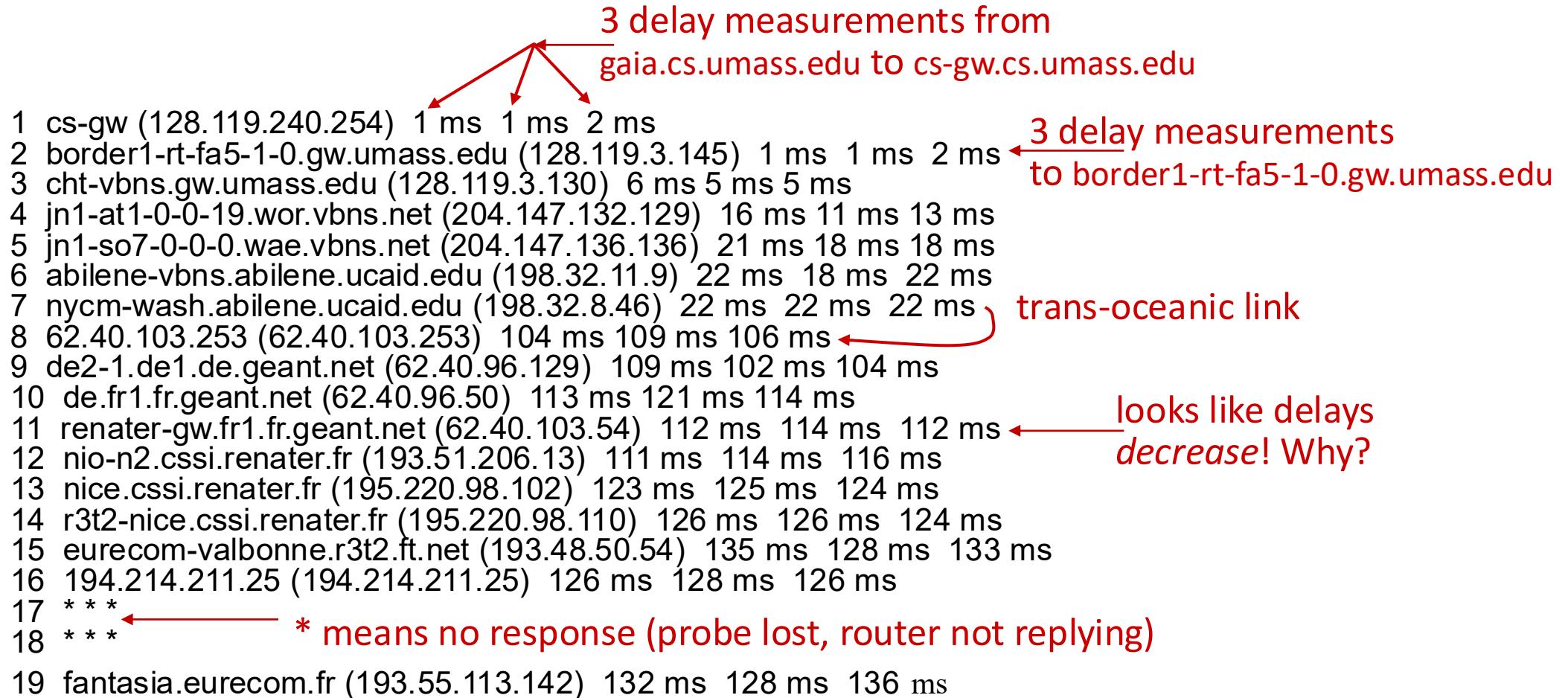
“Real” Internet delays and routes

- what do “real” Internet delay & loss look like?
- **traceroute** program: provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination (with time-to-live field value of i)
 - router i will return packets to sender
 - sender measures time interval between transmission and reply



Real Internet delays and routes

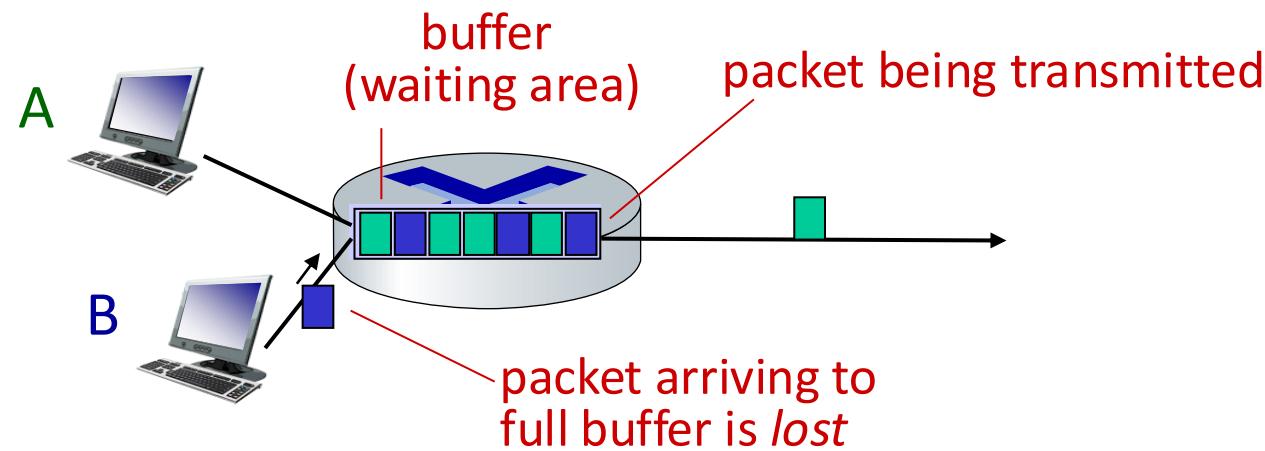
traceroute: gaia.cs.umass.edu to www.eurecom.fr



* Do some traceroutes from exotic countries at www.traceroute.org

Packet loss

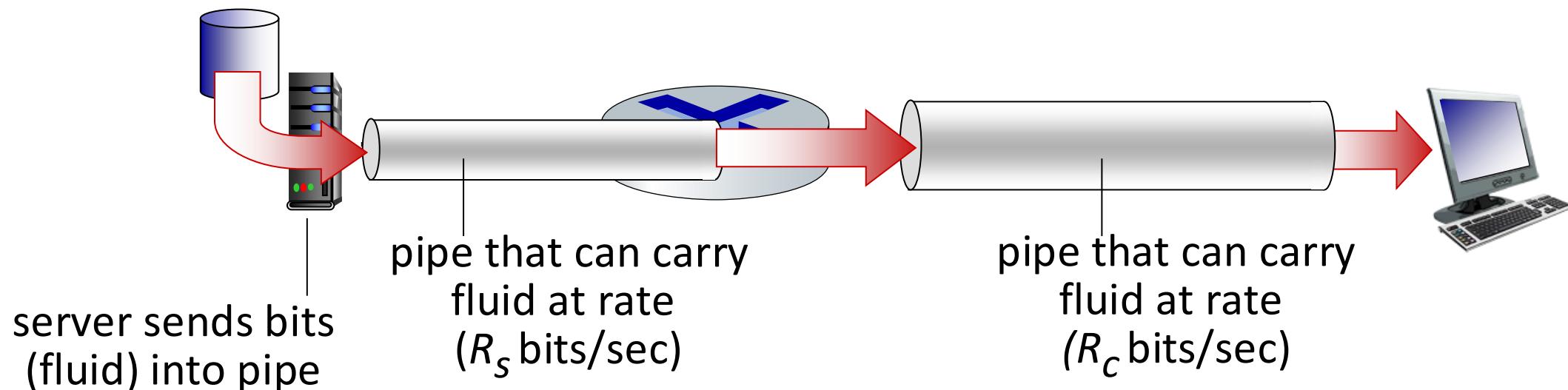
- queue (aka buffer) preceding link in buffer has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not at all



* Check out the Java applet for an interactive animation on queuing and loss

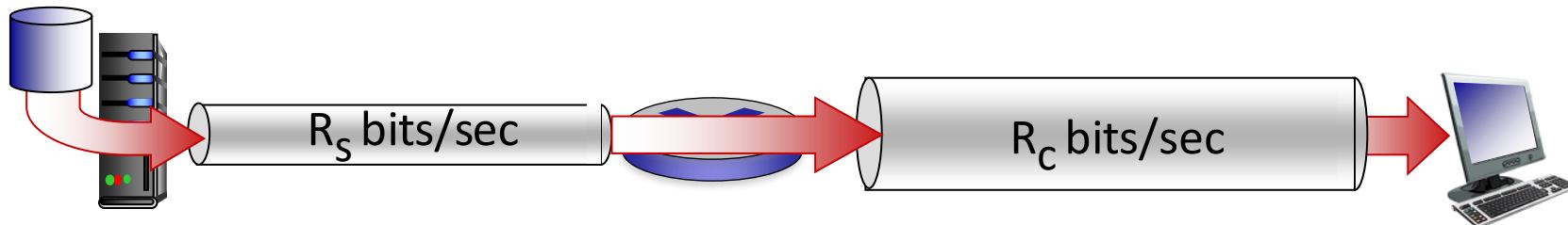
Throughput

- *throughput*: rate (bits/time unit) at which bits are being sent from sender to receiver
 - *instantaneous*: rate at given point in time
 - *average*: rate over longer period of time

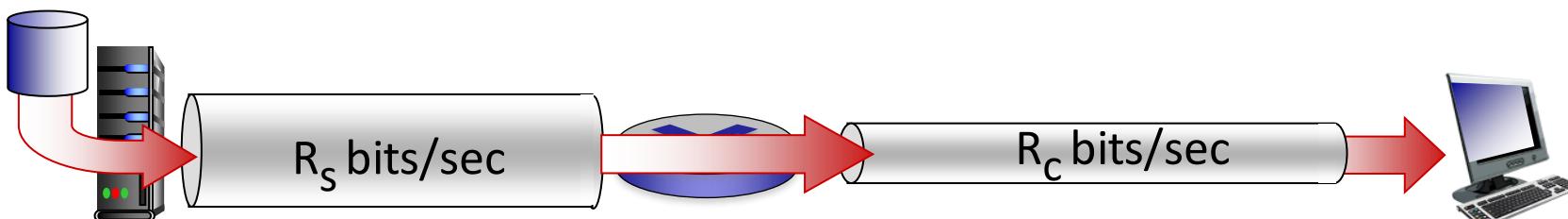


Throughput

$R_s < R_c$ What is average end-end throughput?



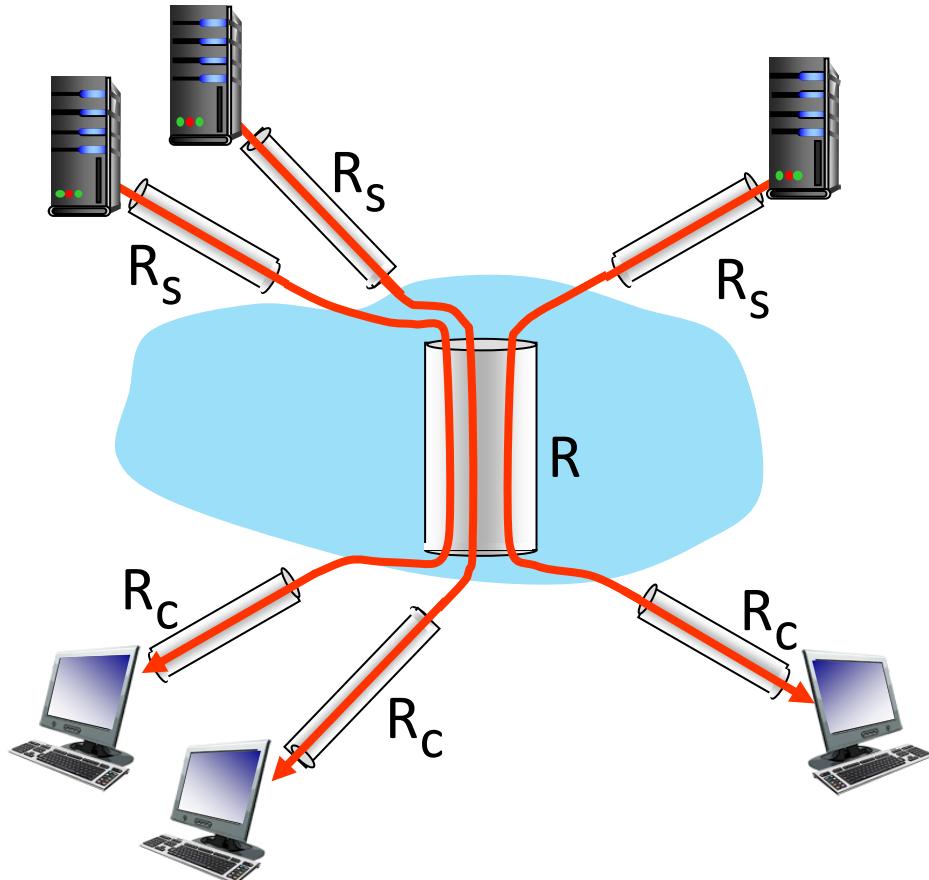
$R_s > R_c$ What is average end-end throughput?



bottleneck link

link on end-end path that constrains end-end throughput

Throughput: network scenario



10 connections (fairly) share
backbone bottleneck link R bits/sec

- per-connection end-end throughput: $\min(R_c, R_s, R/10)$
- in practice: R_c or R_s is often bottleneck

* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/

Chapter 1: roadmap

- What *is* the Internet?
- What *is* a protocol?
- Network edge: hosts, access network, physical media
- Network core: packet/circuit switching, internet structure
- Performance: loss, delay, throughput
- **Security**
- Protocol layers, service models
- History



Network security

- field of network security:
 - how bad guys can attack computer networks
 - how we can defend networks against attacks
 - how to design architectures that are immune to attacks
- Internet not originally designed with (much) security in mind
 - *original vision:* “a group of mutually trusting users attached to a transparent network” ☺
 - Internet protocol designers playing “catch-up”
 - security considerations in all layers!

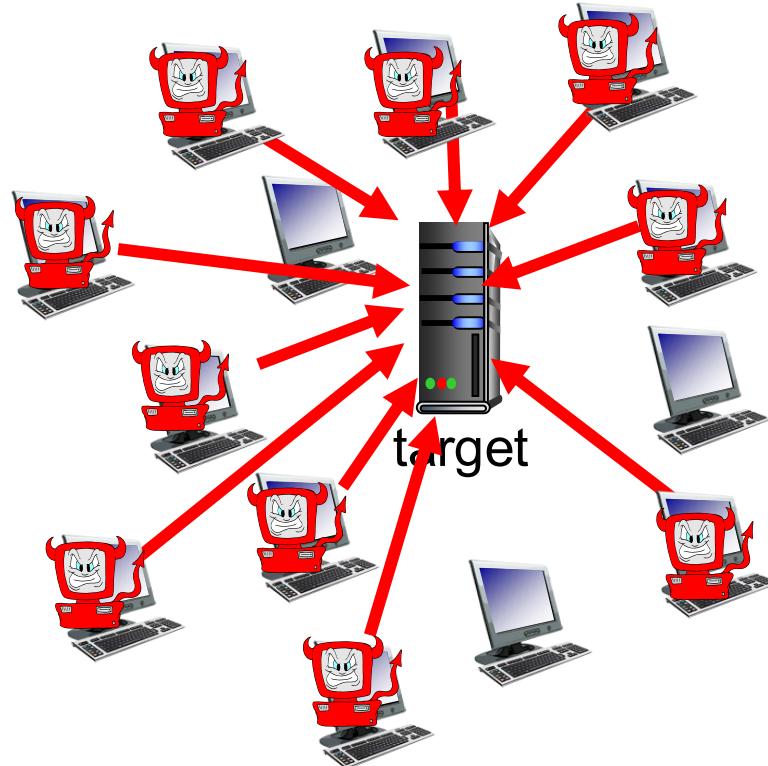
Bad guys: malware

- malware can get in host from:
 - *virus*: self-replicating infection by receiving/executing object (e.g., e-mail attachment)
 - *worm*: self-replicating infection by passively receiving object that gets itself executed
- **spyware malware** can record keystrokes, web sites visited, upload info to collection site
- infected host can be enrolled in **botnet**, used for spam or distributed denial of service (DDoS) attacks

Bad guys: denial of service

Denial of Service (DoS): attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

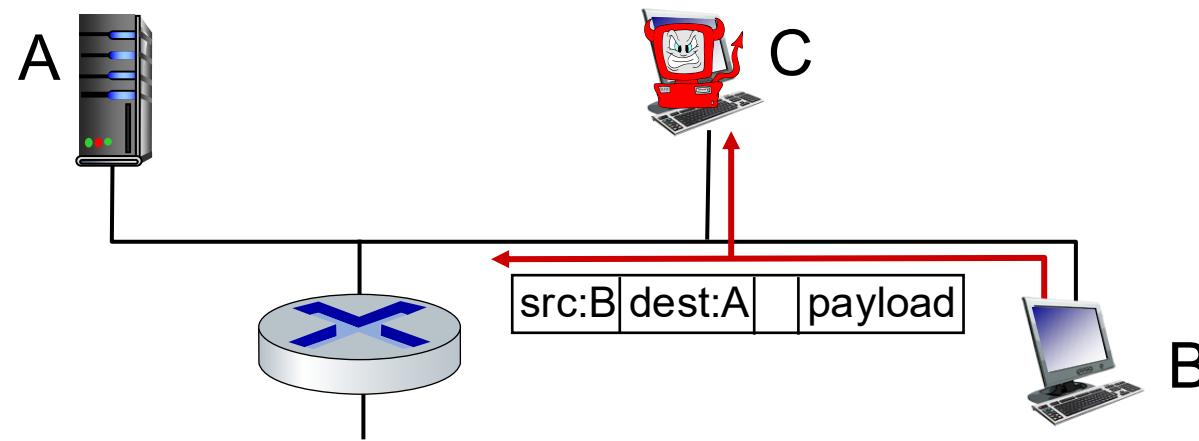
1. select target
2. break into hosts
around the network
(see botnet)
3. send packets to target
from compromised
hosts



Bad guys: packet interception

packet “sniffing”:

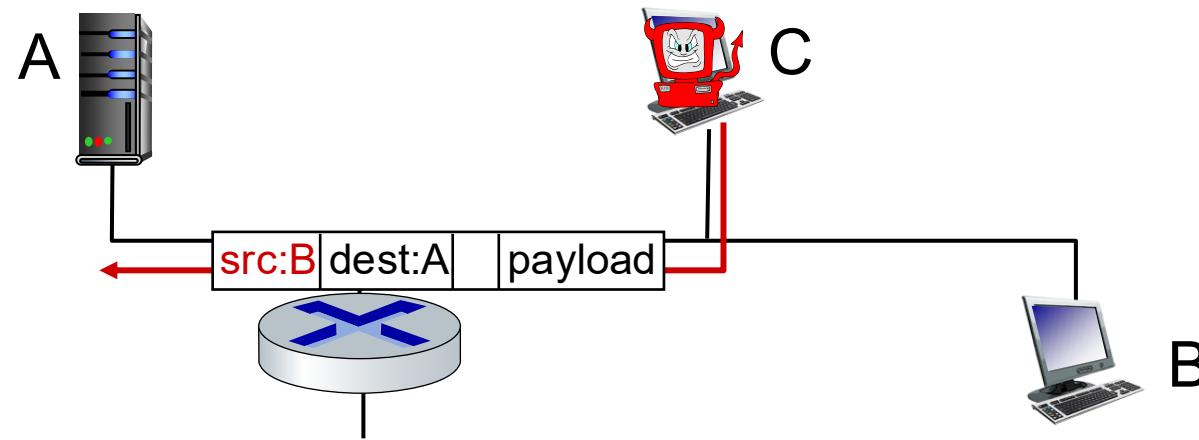
- broadcast media (shared Ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords!) passing by



Wireshark software used for our end-of-chapter labs is a (free) packet-sniffer

Bad guys: fake identity

IP spoofing: send packet with false source address



... lots more on security (throughout, Chapter 8)

Chapter 1: roadmap

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Protocol “layers” and reference models

*Networks are complex,
with many “pieces”:*

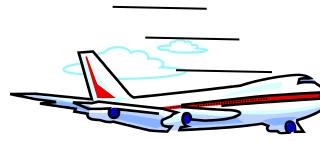
- hosts
- routers
- links of various media
- applications
- protocols
- hardware, software

Question:

is there any hope of
organizing structure of
network?

.... or at least our
discussion of networks?

Example: organization of air travel



ticket (purchase)
baggage (check)
gates (load)
runway takeoff
airplane routing

ticket (complain)
baggage (claim)
gates (unload)
runway landing
airplane routing

airplane routing

airline travel: a series of steps, involving many services

Example: organization of air travel



layers: each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

Q: describe in words
the service provided
in each layer above

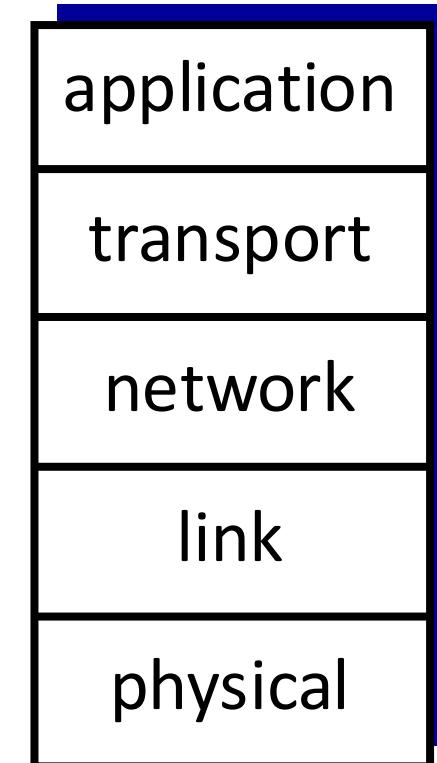
Why layering?

dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
 - layered *reference model* for discussion
- modularization eases maintenance, updating of system
 - change in layer's service *implementation*: transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system
- layering considered harmful?
- layering in other complex systems?

Internet protocol stack

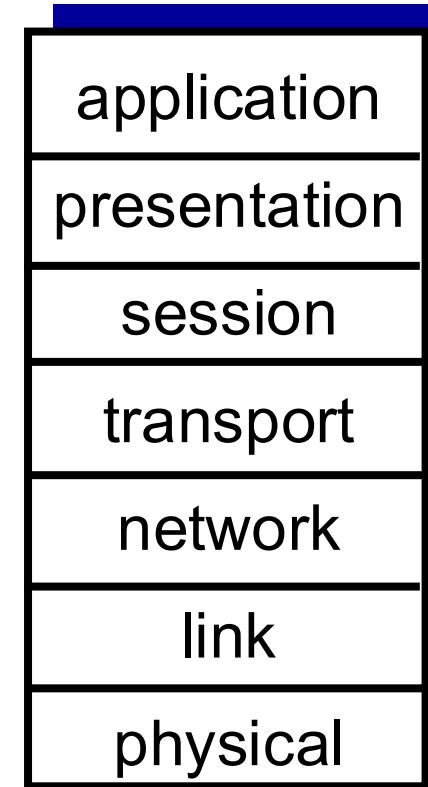
- *application*: supporting network applications
 - IMAP, SMTP, HTTP
- *transport*: process-process data transfer
 - TCP, UDP
- *network*: routing of datagrams from source to destination
 - IP, routing protocols
- *link*: data transfer between neighboring network elements
 - Ethernet, 802.11 (WiFi), PPP
- *physical*: bits “on the wire”



ISO/OSI reference model

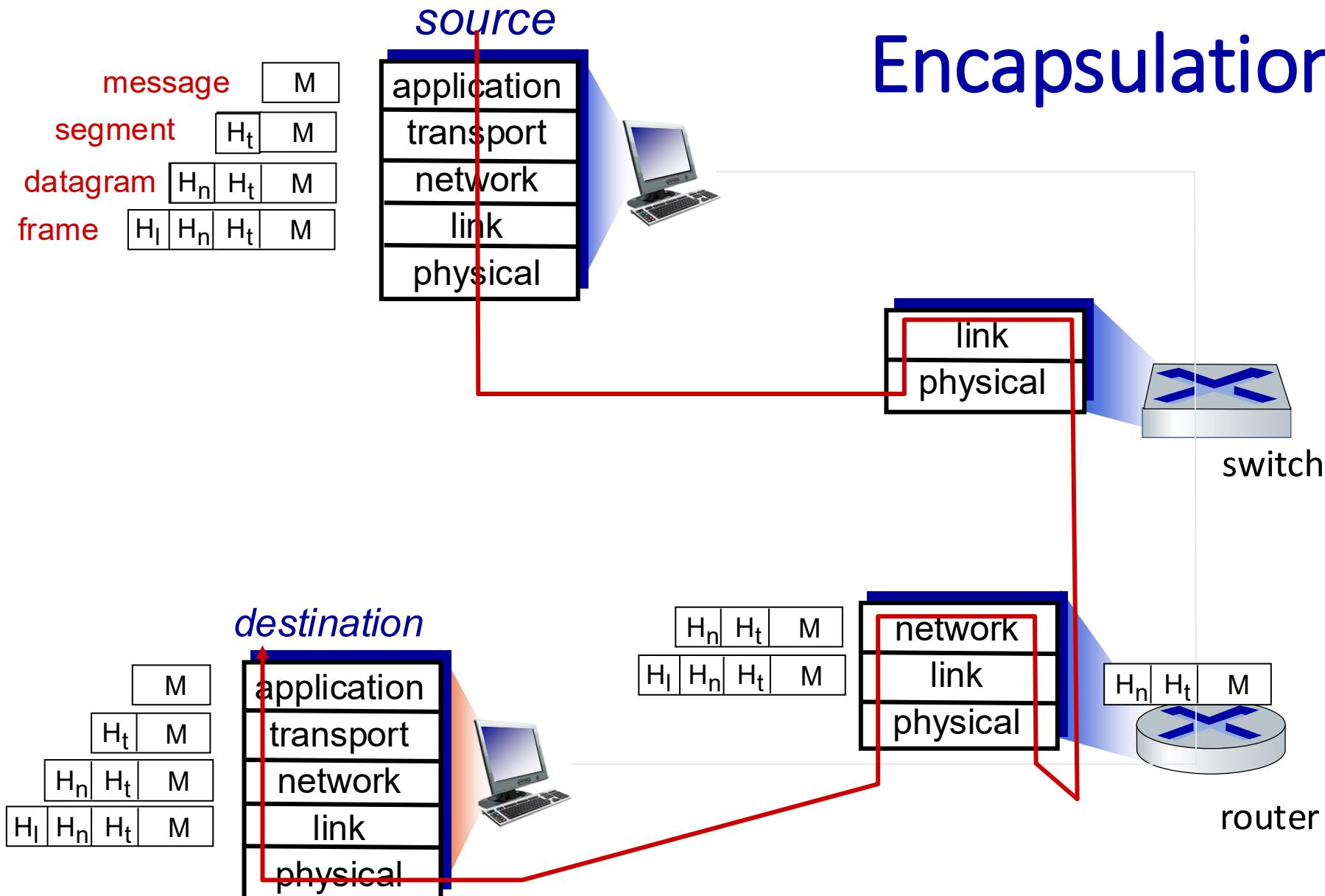
Two layers not found in Internet protocol stack!

- ***presentation***: allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- ***session***: synchronization, checkpointing, recovery of data exchange
- Internet stack “missing” these layers!
 - these services, *if needed*, must be implemented in application
 - needed?



The seven layer OSI/ISO reference model

Encapsulation



Chapter 1: roadmap

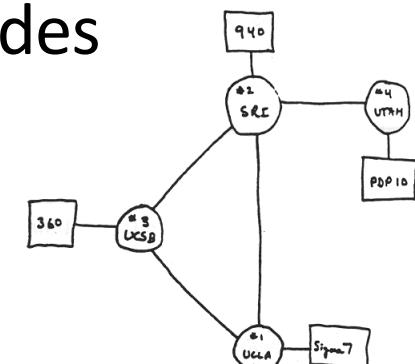
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Internet history

1961-1972: Early packet-switching principles

- 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- 1964: Baran - packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational
- 1972:
 - ARPAnet public demo
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program
 - ARPAnet has 15 nodes



THE ARPA NETWORK

Internet history

1972-1980: Internetworking, new and proprietary nets

- 1970: ALOHAnet satellite network in Hawaii
- 1974: Cerf and Kahn - architecture for interconnecting networks
- 1976: Ethernet at Xerox PARC
- late 70's: proprietary architectures: DECnet, SNA, XNA
- late 70's: switching fixed length packets (ATM precursor)
- 1979: ARPAnet has 200 nodes

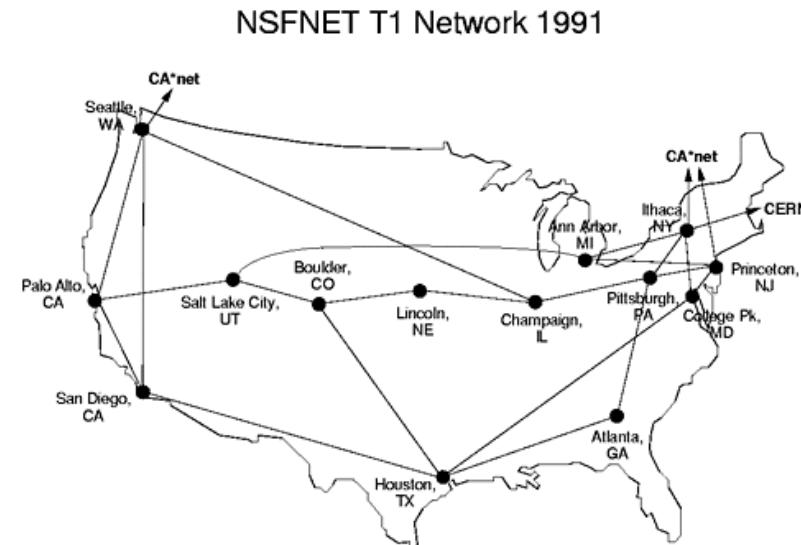
Cerf and Kahn's internetworking principles:

- minimalism, autonomy - no internal changes required to interconnect networks
 - best-effort service model
 - stateless routing
 - decentralized control
- define today's Internet architecture

Internet history

1980-1990: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: smtp e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: ftp protocol defined
- 1988: TCP congestion control
- new national networks: CSnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks



Internet history

1990, 2000s: commercialization, the Web, new applications

- early 1990s: ARPAnet decommissioned
 - 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
 - early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990s: commercialization of the Web
- late 1990s – 2000s:
- more killer apps: instant messaging, P2P file sharing
 - network security to forefront
 - est. 50 million host, 100 million+ users
 - backbone links running at Gbps

Internet history

2005-present: more new applications, Internet is “everywhere”

- ~18B devices attached to Internet (2017)
 - rise of smartphones (iPhone: 2007)
- aggressive deployment of broadband access
- increasing ubiquity of high-speed wireless access: 4G/5G, WiFi
- emergence of online social networks:
 - Facebook: ~ 2.5 billion users
- service providers (Google, FB, Microsoft) create their own networks
 - bypass commercial Internet to connect “close” to end user, providing “instantaneous” access to search, video content, ...
- enterprises run their services in “cloud” (e.g., Amazon Web Services, Microsoft Azure)

Chapter 1: summary

We've covered a "ton" of material!

- Internet overview
- what's a protocol?
- network edge, access network, core
 - packet-switching versus circuit-switching
 - Internet structure
- performance: loss, delay, throughput
- layering, service models
- security
- history

You now have:

- context, overview, vocabulary, "feel" of networking
- more depth, detail, *and fun* to follow!

Additional Chapter 1 slides

Wireshark

